

DNS

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DNS - Internet's Directory Service

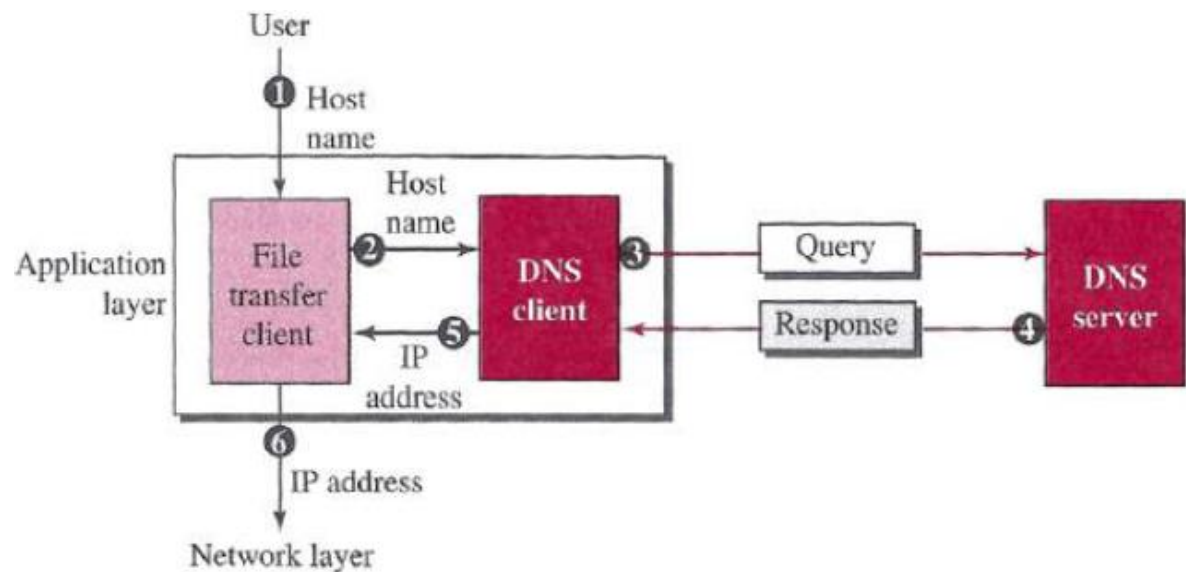


- Just as humans can be identified in many ways, so too can **Internet hosts**. Two ways:
 - **Hostname** (e.g., gmail.co.in, iitg.ac.in)
 - these are **mnemonic**,
 - user friendly for **Humans**
 - **IP Address** (e.g., 121.7.106.83, 172.17.0.10)
 - these are structured **numeric** digits,
 - user friendly for **Routers**
- The **Internet** needs to have a **directory system** that can **map a name to an address**.
- The Internet is so huge today, a **central directory system** cannot hold all the mapping.
- A **better solution** - distribute the directory information among many computers in the world.
- This method is used by the **Domain Name System (DNS)**.

Cont...

- The DNS is a combination of :
 - a **distributed database** implemented in a hierarchy of **DNS servers**, and
 - an **application-layer protocol** that **allows hosts to query** the distributed database
- Let the purpose of accessing the Internet is to make a connection between the **file transfer client** and **server**,
- but before this can happen, **another connection** needs to be made between the **DNS client** and **DNS server**

- DNS protocol runs over **UDP** and uses **port 53**.
- The **DNS servers** are often UNIX machines running the Berkeley Internet Name Domain (BIND) software



Design for DNS



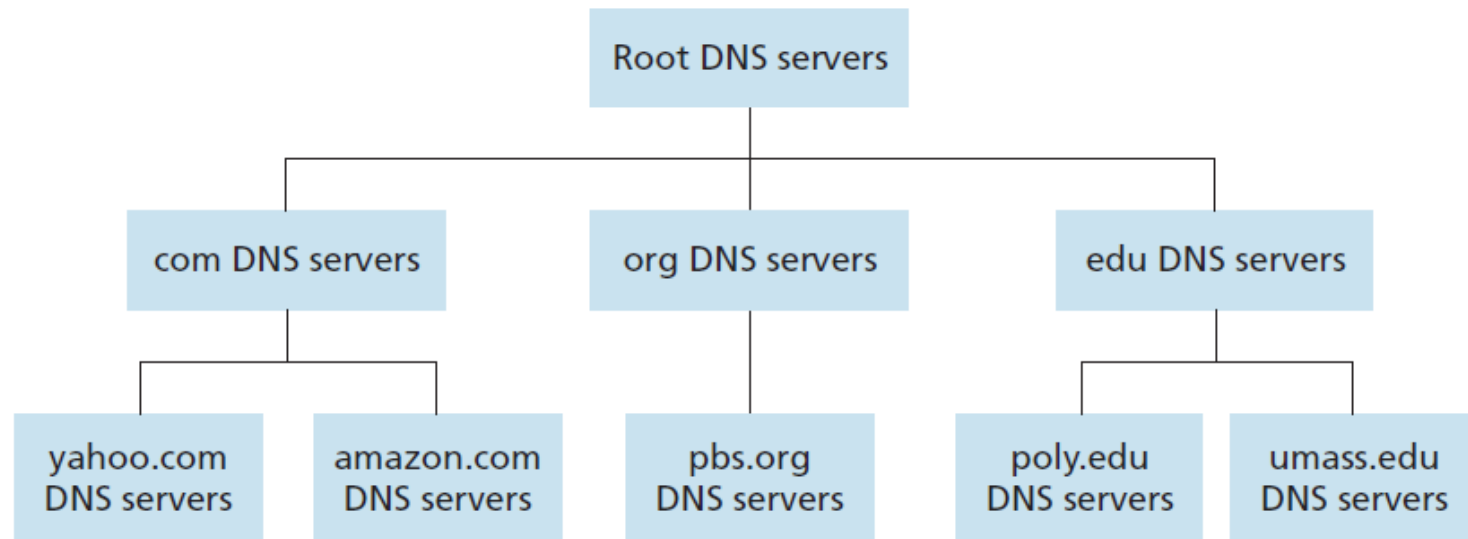
- Design for DNS:
 - Centralized / Distributed
- The problems with a **centralized design** include:
 - **A single point of failure**: DNS server crashes, so does the entire Internet!
 - **Traffic volume**: A single DNS server would have to handle all DNS queries generated from hundreds of millions of hosts
 - **Distant database**: A single DNS server cannot be “close to” all the querying clients.
 - **Maintenance**: The single DNS server would have to keep records for all Internet hosts. Management of it becomes very difficult!

DNS Services



- **Fundamental service** : a directory service that translates hostnames to IP addresses.
- provides a few **other important services** :
 - **Host aliasing**: [relay1.west-coast.enterprise.com](#) could have, say, two aliases such as [enterprise.com](#) and [www.enterprise.com](#)
 - **Mail server aliasing**: the canonical hostname of the Hotmail server might be something like [relay1.west-coast.hotmail.com](#) but the mail server is simply [hotmail.com](#)
 - **Load distribution**: used to perform **load distribution** among replicated servers. For replicated servers, a set of IP addresses is thus associated with one canonical hostname.

Hierarchy of DNS servers



- the mappings for all of the hosts in the Internet are distributed across the DNS servers
- **three classes** of DNS servers
 - root DNS servers (until 2012, Internet has 13 root DNS servers)
 - top-level domain (TLD) DNS servers
 - authoritative DNS servers (large university or organization may have it)
 - **local DNS server** : Each ISP has one or more local DNS

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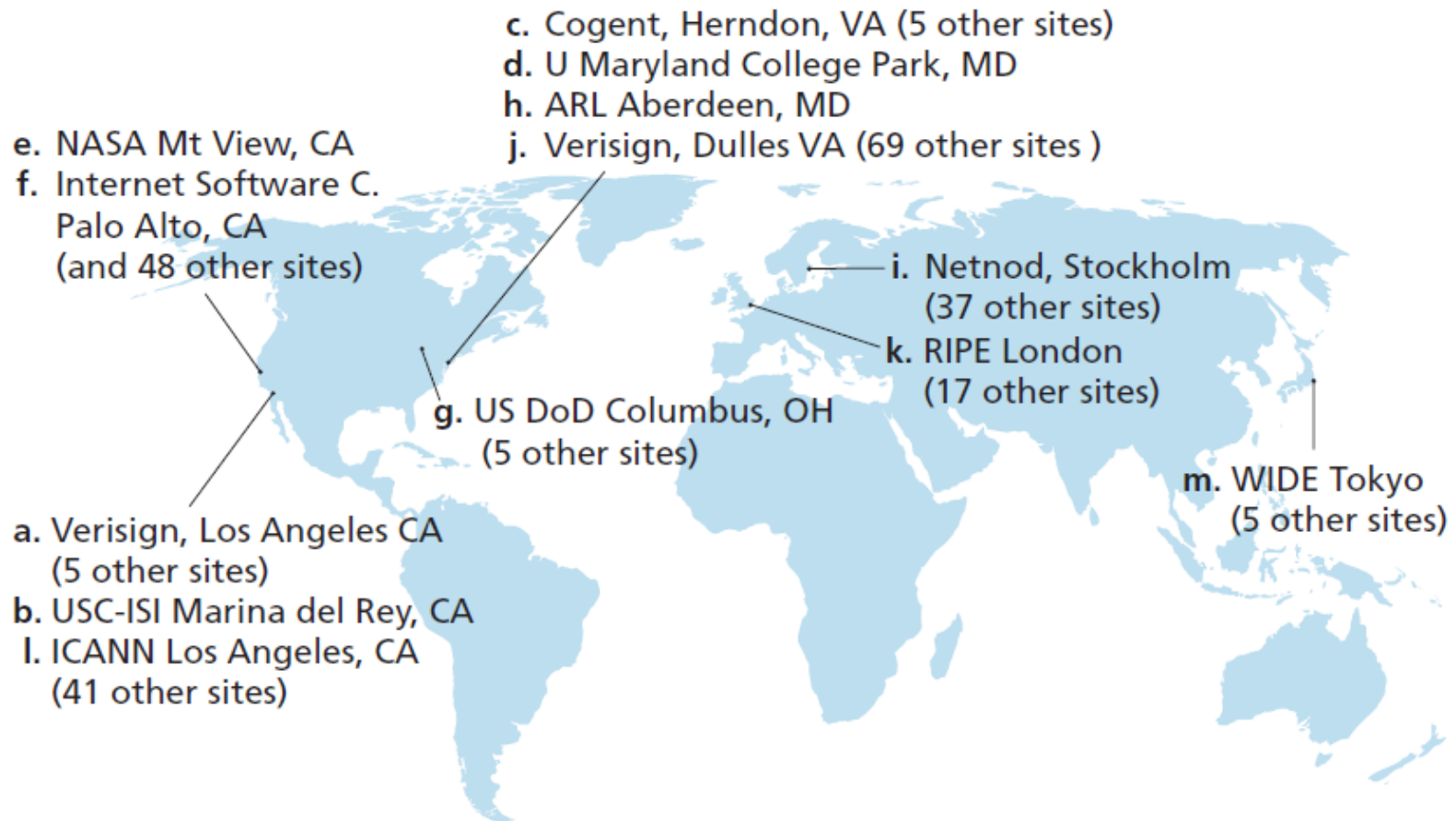
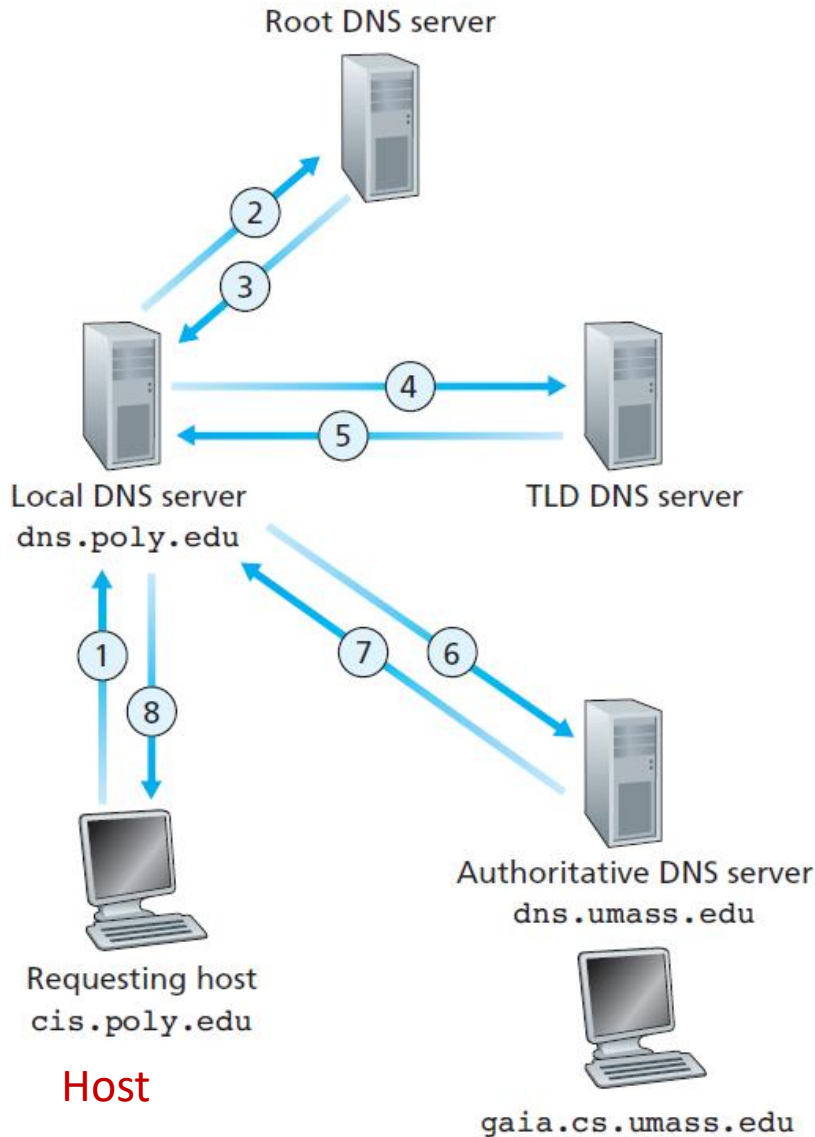


Figure 2.20 ♦ DNS root servers in 2012 (name, organization, location)

Interaction among DNS servers



- Let the **host** cis.poly.edu desires the IP address of gaia.cs.umass.edu.
- Let the Polytechnic's local DNS server is called dns.poly.edu
- Let an authoritative DNS server for gaia.cs.umass.edu is called dns.umass.edu.

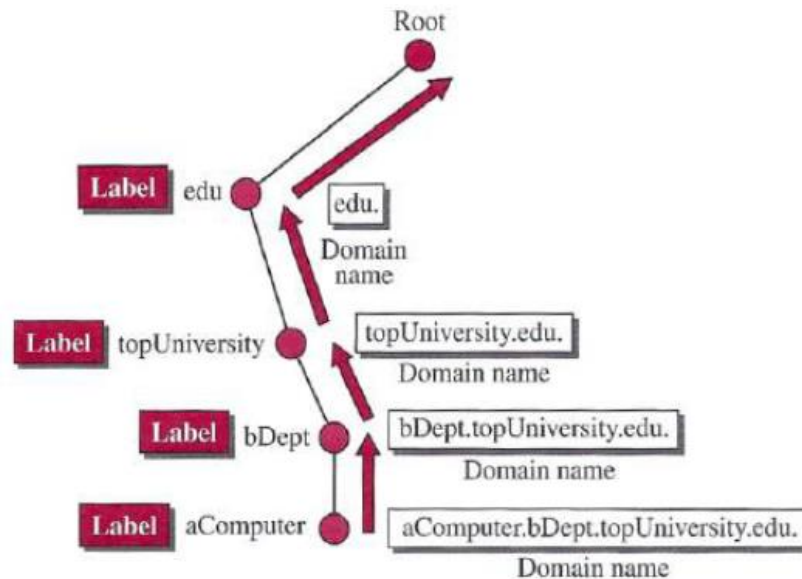
- 1) The **host** first sends a DNS query message to its **local** DNS server.
- 2) The **local** DNS server forwards the query message to a **root** DNS server.
- 3) The **root** DNS server takes note of the edu suffix and **returns** a list of IP addresses for **TLD** servers responsible for edu.
- 4) The **local** DNS server then resends the query to one of these **TLD** servers.
- 5) The **TLD** server responds with the IP address of the **authoritative** DNS server
- 6) Finally, the **local** DNS server resends the query message directly to the **authoritative** DNS server

Name Space

- the **names must be unique** because the addresses are unique.
- A **name space** that maps each address to a unique name can be organized in two ways:
 - flat
 - hierarchical
- ***flat name space***
 - a name is assigned to an address
 - a name is a sequence of characters without structure
 - The names may or may not have a common section
 - **Disadvantage:** it cannot be used in a large system such as the Internet because it must be **centrally controlled** to avoid ambiguity and duplication

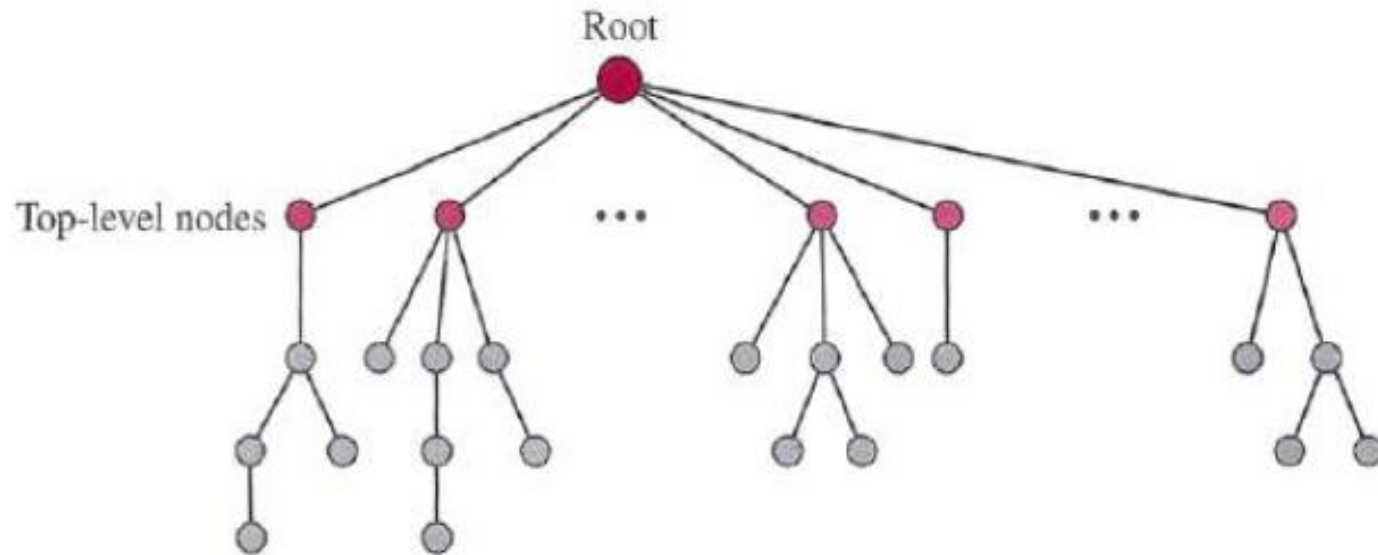
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- *hierarchical name space*: each name is made of **several parts**
 - the **first part** can define the nature of the organization
 - the **second part** can define the name of an organization
 - the **third part** can define departments in the organization
- *Advantages*
 - the authority to assign and control the name spaces can be decentralized.
 - A central authority can assign the part of the name. E.g, name & nature of the organization
Rest of the name can be assigned by the organization itself



Domain Name Space

- the names are defined in an **inverted-tree structure** with the root at the top.

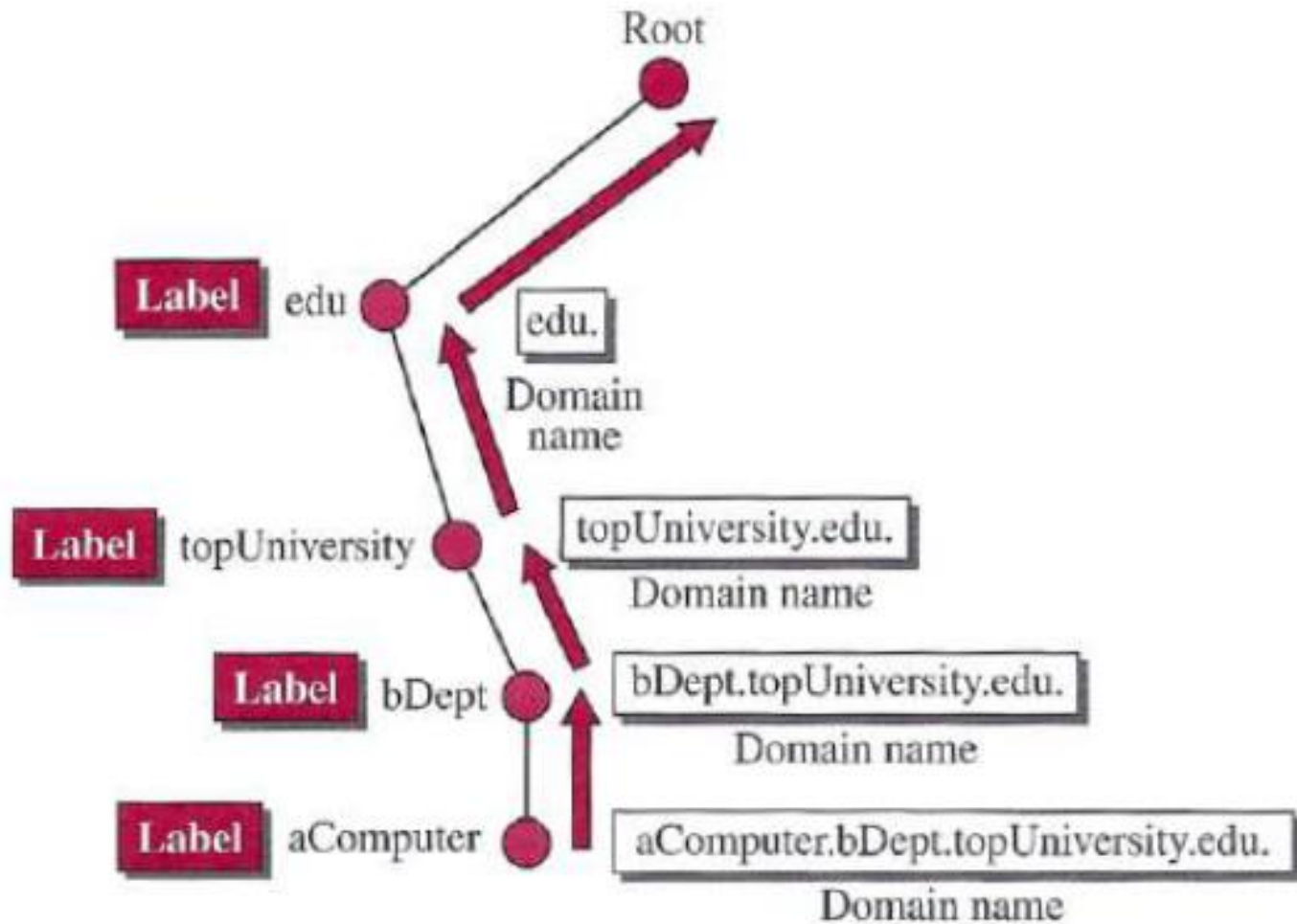


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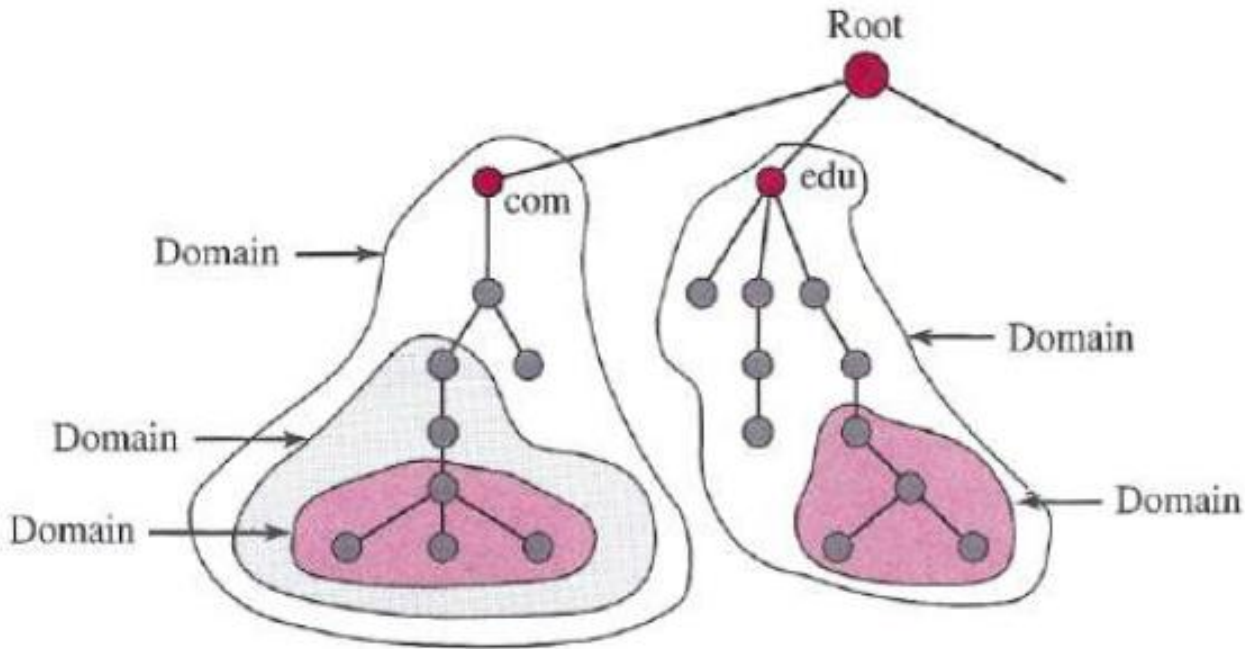
- *Label:*
 - Each node in the tree has a label, which is a string with a maximum of 63 characters.
 - The root label is a null string (empty string).
- *Domain Name:*
 - Each **node** in the tree has a **domain name**.
 - A full domain name is a sequence of labels **separated by dots** (.)
 - The domain names are always **read from** the node up to the **root**.
 - The last label is the label of the root (null).
- **Fully qualified domain name (FQDN):**
 - If a label is terminated by a null string.
 - Else, it is Partially qualified domain name (PQDN)

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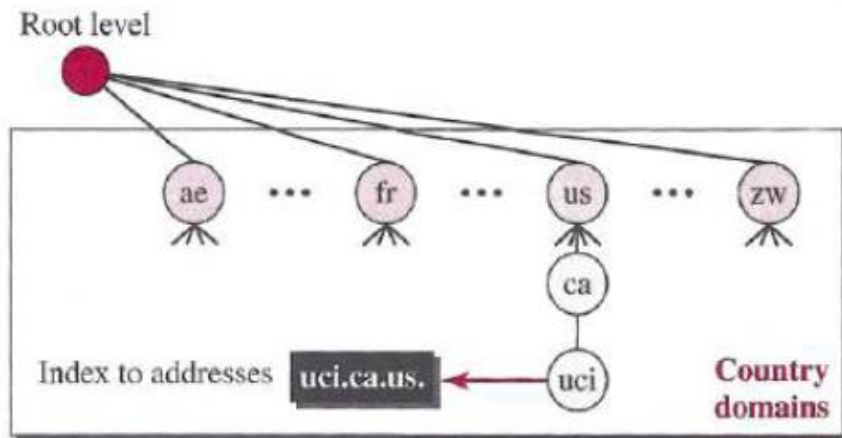
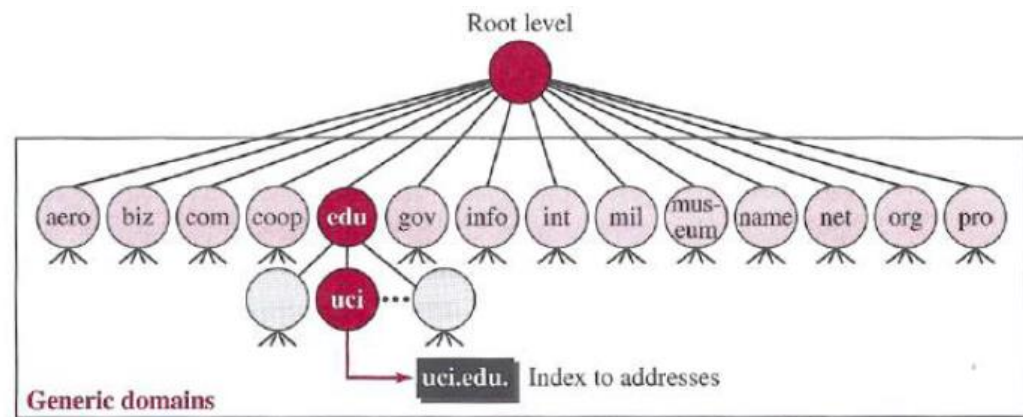
Domain

- A domain is a subtree of the domain name space



DNS in the Internet

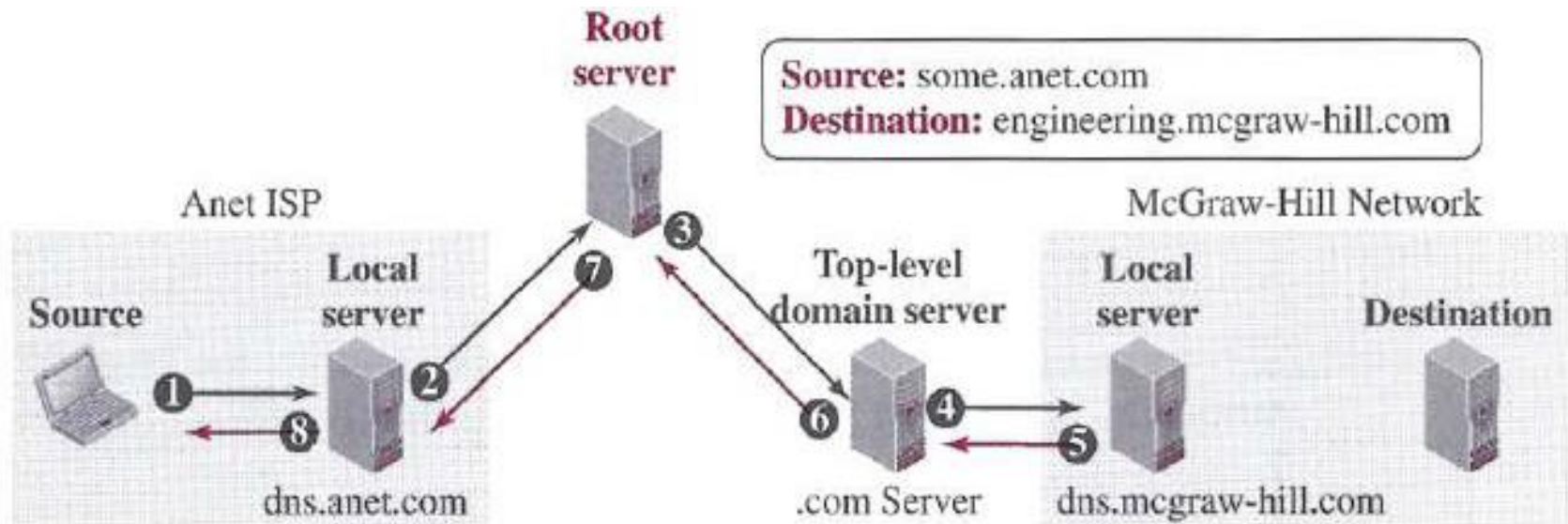
- DNS is a protocol that can be used in different platforms.
- the **domain name space (tree)** is designed by many different ways:
 - **generic** domains
 - **country** domains



- E.g, The address uci.ca.us can be translated to **University of California, Irvine**, in the state of **California** in the **United States**.

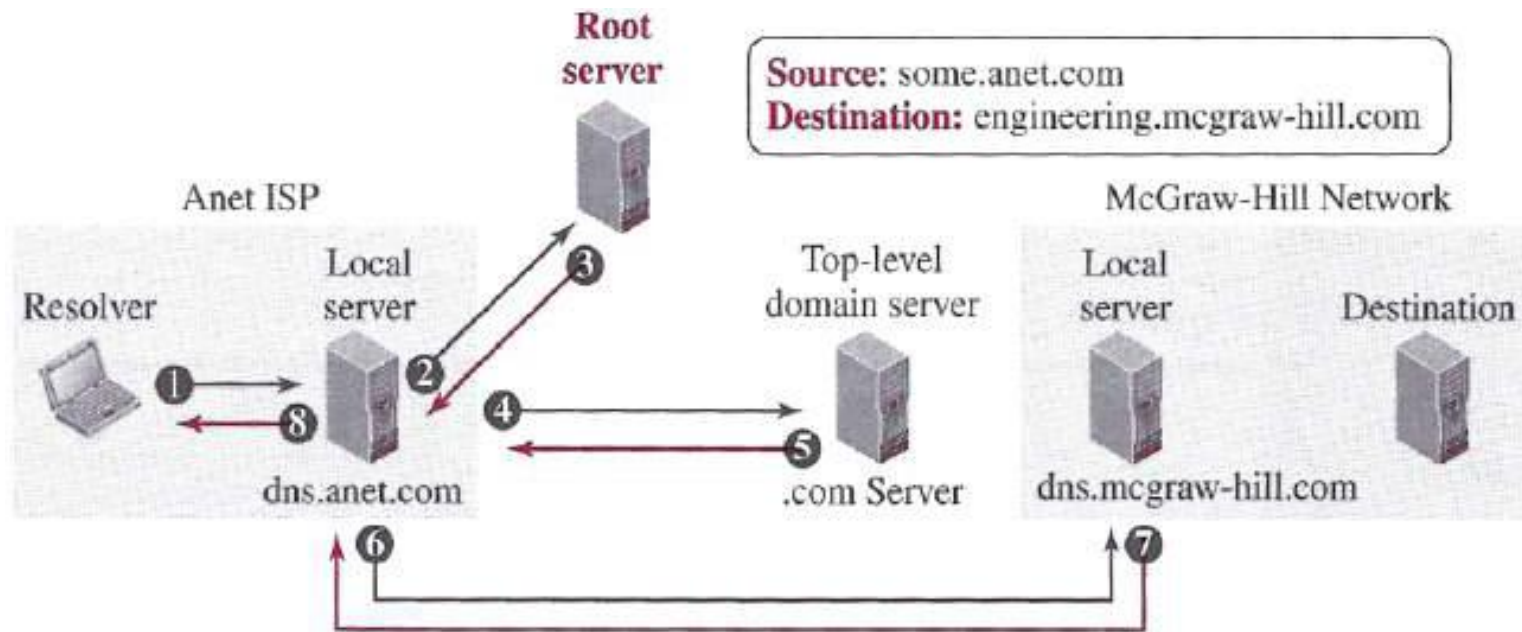
Name-Address Resolution

- Mapping a name to an address is called *name-address resolution*
- DNS is designed as a **client-server application**.
- A **resolution process** can be
 - Recursive
 - Iterative



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- Iterative Resolution:



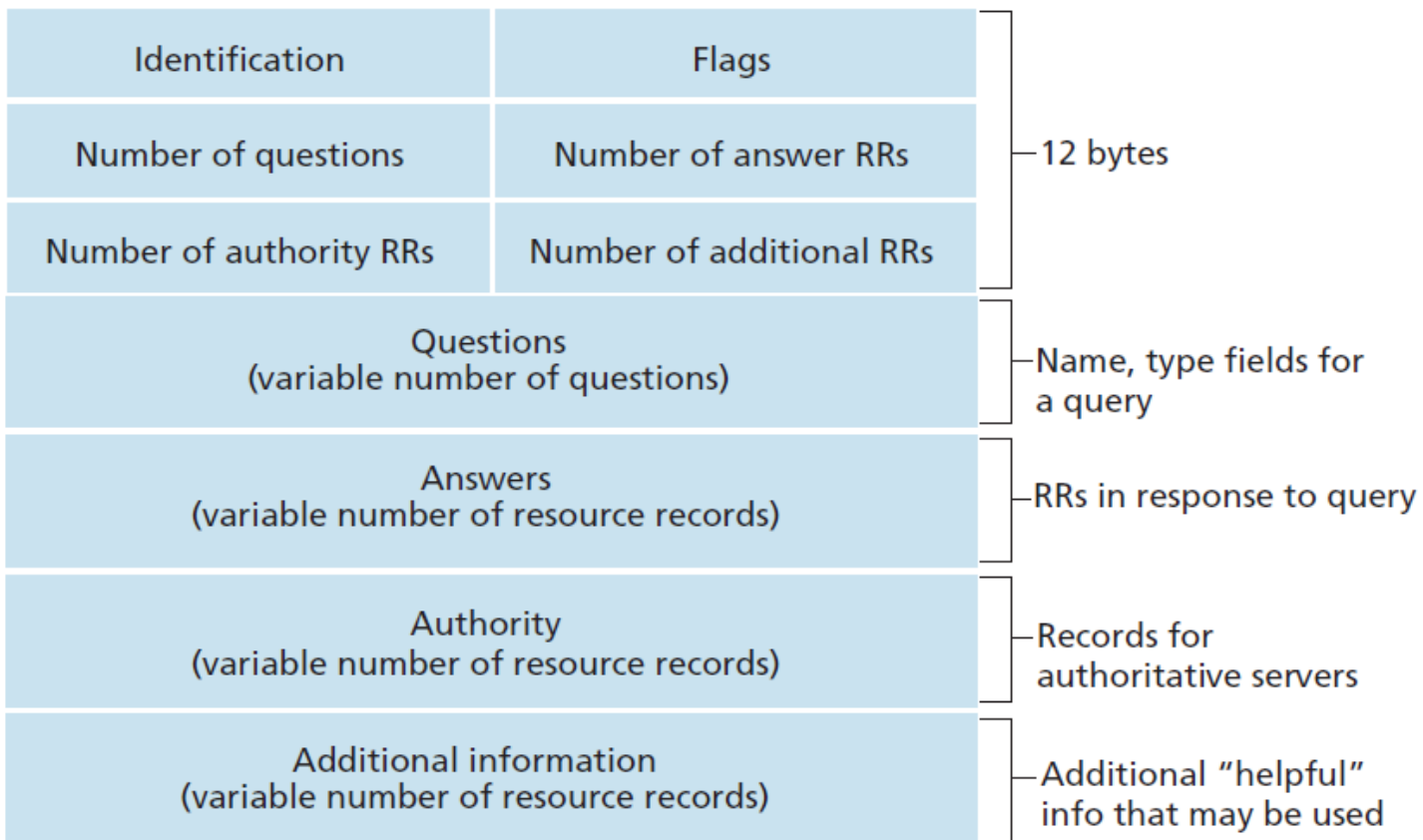
DNS Caching



- Each time a server receives a query for a name that is not in its domain, it needs to search its database for a server IP address.
- Reduction of this search time would increase efficiency.
- DNS server handles this with a mechanism called *caching*
- Caching speeds up resolution, but it can also be problematic by sending outdated mapping.
- To counter this, TTL (time-to-live) based technique is used.

DNS Messages

- The **identification field** is used by the client to match the response with the query.
- The **flag field** defines whether the message is a query or response.



Cont...



- DNS can use either **UDP** or **TCP**.
- In both cases the well-known port used by the server is port 53.
- Example:
 - In UNIX and Windows, the *nslookup* utility can be used to retrieve address/name mapping.

```
$nslookup www.forouzan.biz  
Name: www.forouzan.biz  
Address: 198.170.240.179
```

Thanks!