

NPTEL MOOC, JAN-FEB 2015  
Week 1, Module 5

# DESIGN AND ANALYSIS OF ALGORITHMS

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# Analysis of algorithms

- \* Measuring efficiency of an algorithm
  - \* Time: How long the algorithm takes (running time)
  - \* Space: Memory requirement

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# Time and space

- \* Time depends on processing speed
  - \* Impossible to change for given hardware
- \* Space is a function of available memory
  - \* Easier to reconfigure, augment
- \* Typically, we will focus on time, not space



# Measuring running time

- \* Analysis independent of underlying hardware
  - \* Don't use actual time
  - \* Measure in terms of “basic operations”
- \* Typical basic operations
  - \* Compare two values
  - \* Assign a value to a variable
- \* Other operations may be basic, depending on context
  - \* Exchange values of a pair of variables



# Input size

- \* Running time depends on input size
  - \* Larger arrays will take longer to sort
- \* Measure time efficiency as function of input size
  - \* Input size  $n$
  - \* Running time  $t(n)$
- \* Different inputs of size  $n$  may each take a different amount of time
- \* Typically  $t(n)$  is **worst case estimate**



# Example 1: Sorting

- \* Sorting an array with  $n$  elements
  - \* Naïve algorithms : time proportional to  $n^2$
  - \* Best algorithms : time proportional to  $n \log n$
- \* How important is this distinction?
- \* Typical CPUs process up to  $10^8$  operations per second
  - \* Useful for approximate calculations



# Example 1: Sorting ...

- \* Telephone directory for mobile phone users in India
  - \* India has about 1 billion =  $10^9$  phones
- \* Naïve  $n^2$  algorithm requires  $10^{18}$  operations
  - \*  $10^8$  operations per second  $\Rightarrow 10^{10}$  seconds
  - \* 2778000 hours
  - \* 115700 days
  - \* 300 years!
- \* Smart  $n \log n$  algorithm takes only about  $3 \times 10^{10}$  operations
  - \* About 300 seconds, or 5 minutes



# Example 2: Video game

- \* Several objects on screen
- \* Basic step: find closest pair of objects
- \* Given  $n$  objects, naïve algorithm is again  $n^2$ 
  - \* For each pair of objects, compute their distance
  - \* Report minimum distance over all such pairs
- \* There is a clever algorithm that takes time  $n \log n$



# Example 2: Video game ...

- \* High resolution monitor has  $2500 \times 1500$  pixels
  - \* 3.75 million points
- \* Suppose we have  $500,000 = 5 \times 10^5$  objects
- \* Naïve algorithm takes  $25 \times 10^{10}$  steps = 2500 seconds
  - \* 2500 seconds = 42 minutes response time is unacceptable!
- \* Smart  $n \log n$  algorithm takes a fraction of a second



# Orders of magnitude

- \* When comparing  $t(n)$  across problems, focus on orders of magnitude
- \* Ignore constants
- \*  $f(n) = n^3$  eventually grows faster than  $g(n) = 5000 n^2$
- \* For small values of  $n$ ,  $f(n)$  is smaller than  $g(n)$
- \* At  $n = 5000$ ,  $f(n)$  overtakes  $g(n)$
- \* What happens in the limit, as  $n$  increases :  
**asymptotic complexity**



# Typical functions

- \* We are interested in orders of magnitude
- \* Is  $t(n)$  proportional to  $\log n$ , ...,  $n^2$ ,  $n^3$ , ...,  $2^n$ ?
- \* Logarithmic, polynomial, exponential ...

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# Typical functions $t(n)$ ...

Input	$\log n$	$n$	$n \log n$	$n^2$	$n^3$	$2^n$	$n!$
10	3.3	10	33	100	1000	1000	$10^6$
100	6.6	100	66	$10^4$	$10^6$	$10^{30}$	$10^{157}$
1000	10	1000	$10^4$	$10^6$	$10^9$		
$10^4$	13	$10^4$	$10^5$	$10^8$	$10^{12}$		
$10^5$	17	$10^5$	$10^6$	$10^{10}$			
$10^6$	20	$10^6$	$10^7$				
$10^7$	23	$10^7$	$10^8$				
$10^8$	27	$10^8$	$10^9$				
$10^9$	30	$10^9$	$10^{10}$				
$10^{10}$	33	$10^{10}$					