NPTEL MOOC, JAN-FEB 2015 Week 2, Module 4

DESIGN AND ANALYSIS OF ALGORITHMS

Insertion Sort

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Sorting

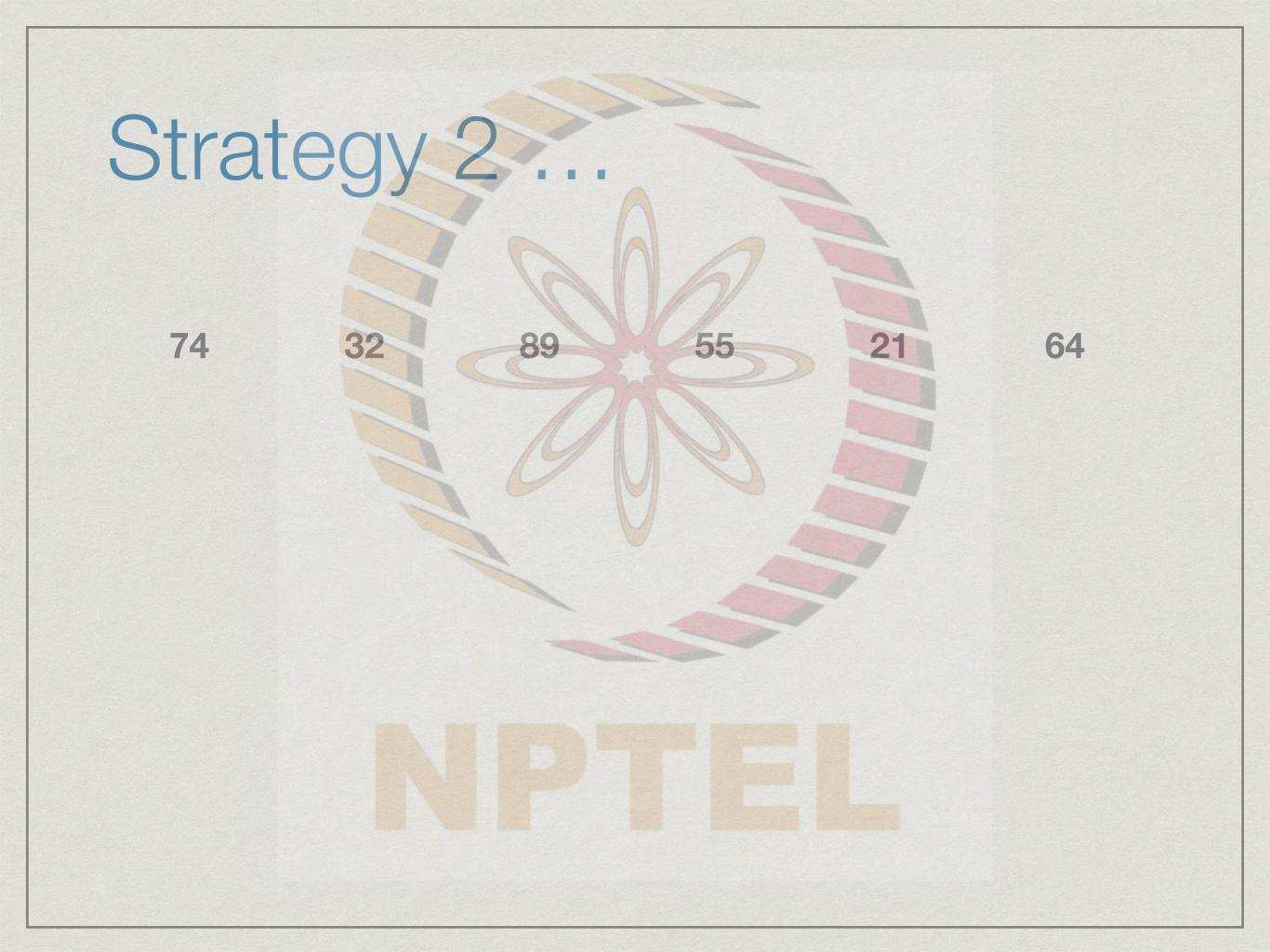
- * Searching for a value
 - * Unsorted array linear scan, O(n)
 - * Sorted array binary search, O(log n)
- * Other advantages of sorting
 - * Finding median value: midpoint of sorted list
 - * Checking for duplicates
 - * Building a frequency table of values

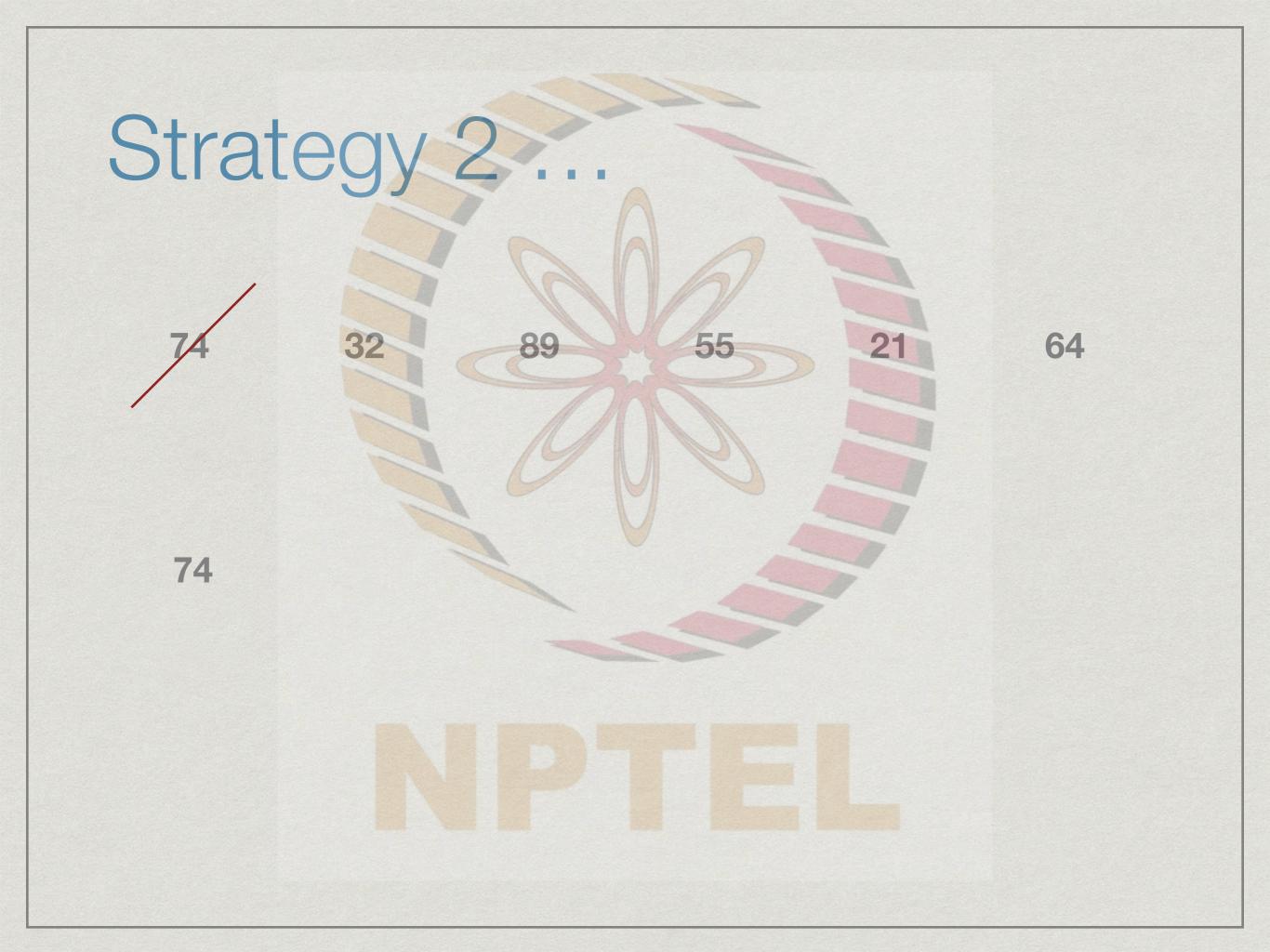
How to sort?

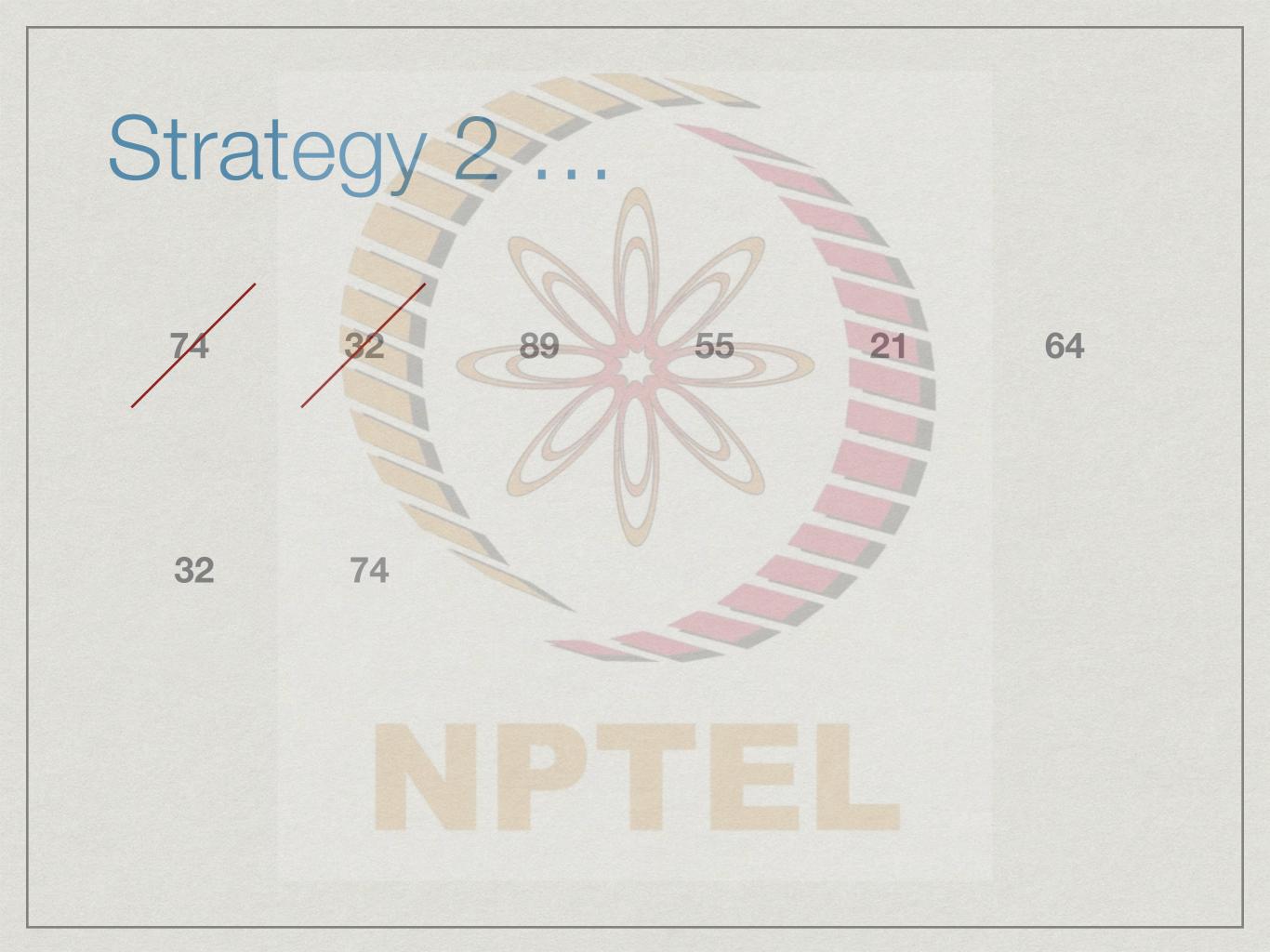
- * You are a Teaching Assistant for a course
- * The instructor gives you a stack of exam answer papers with marks, ordered randomly
- * Your task is to arrange them in descending order

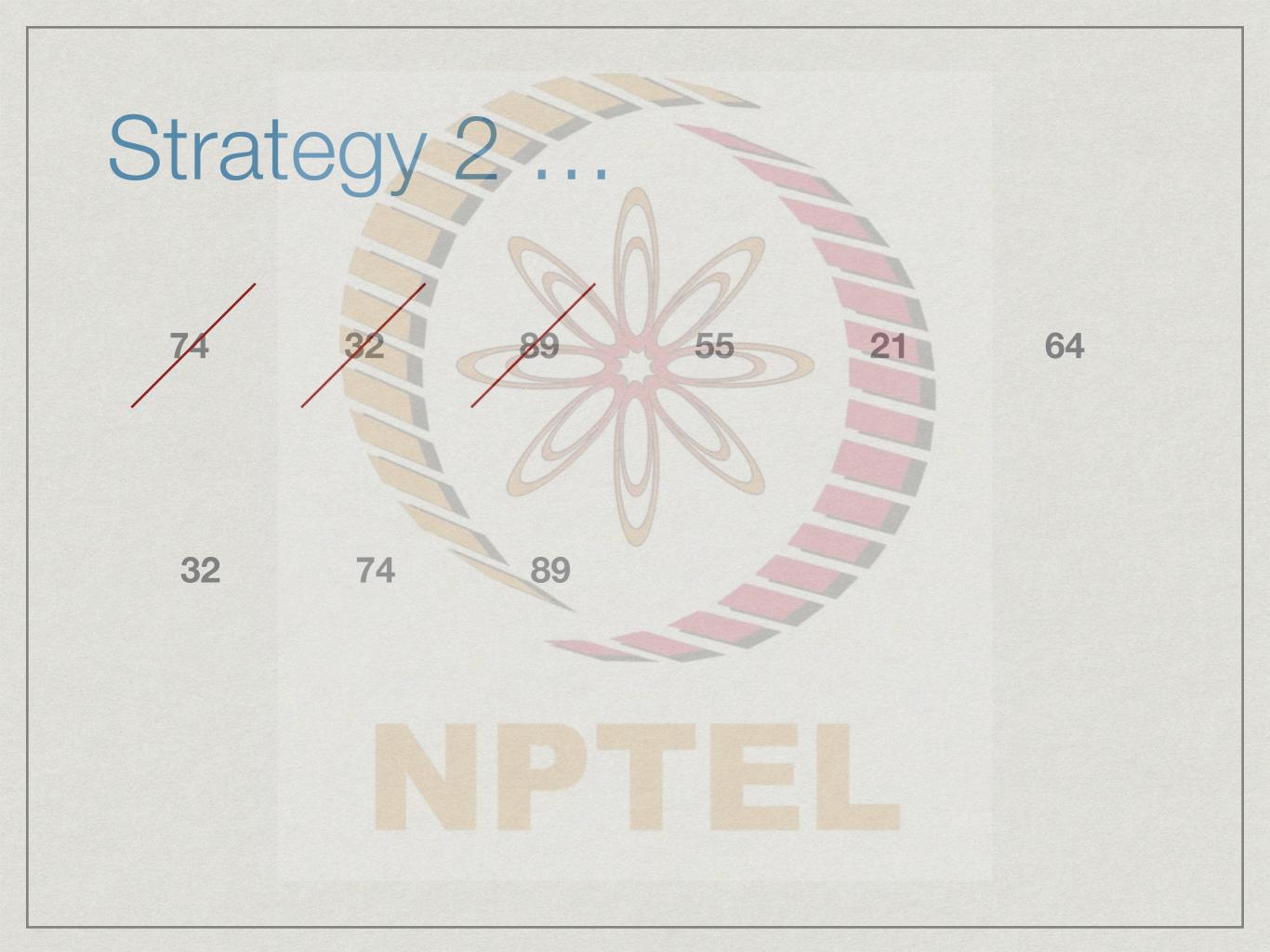
Strategy 2

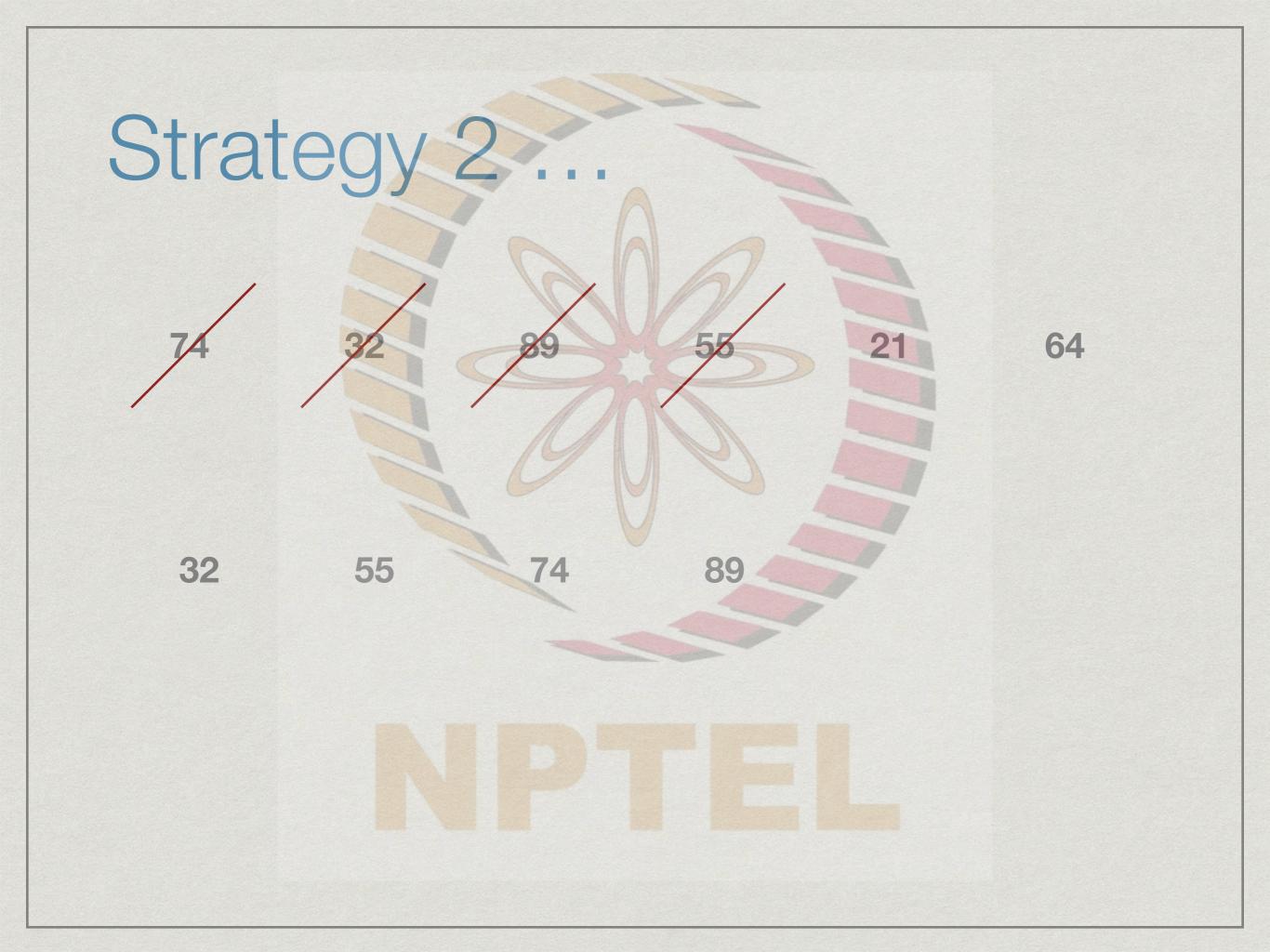
- * First paper: put in a new stack
- * Second paper:
 - * Lower marks than first? Place below first paper Higher marks than first? Place above first paper
- * Third paper
 - * Insert into the correct position with respect to first two papers
- * Do this for each subsequent paper: insert into correct position in new sorted stack

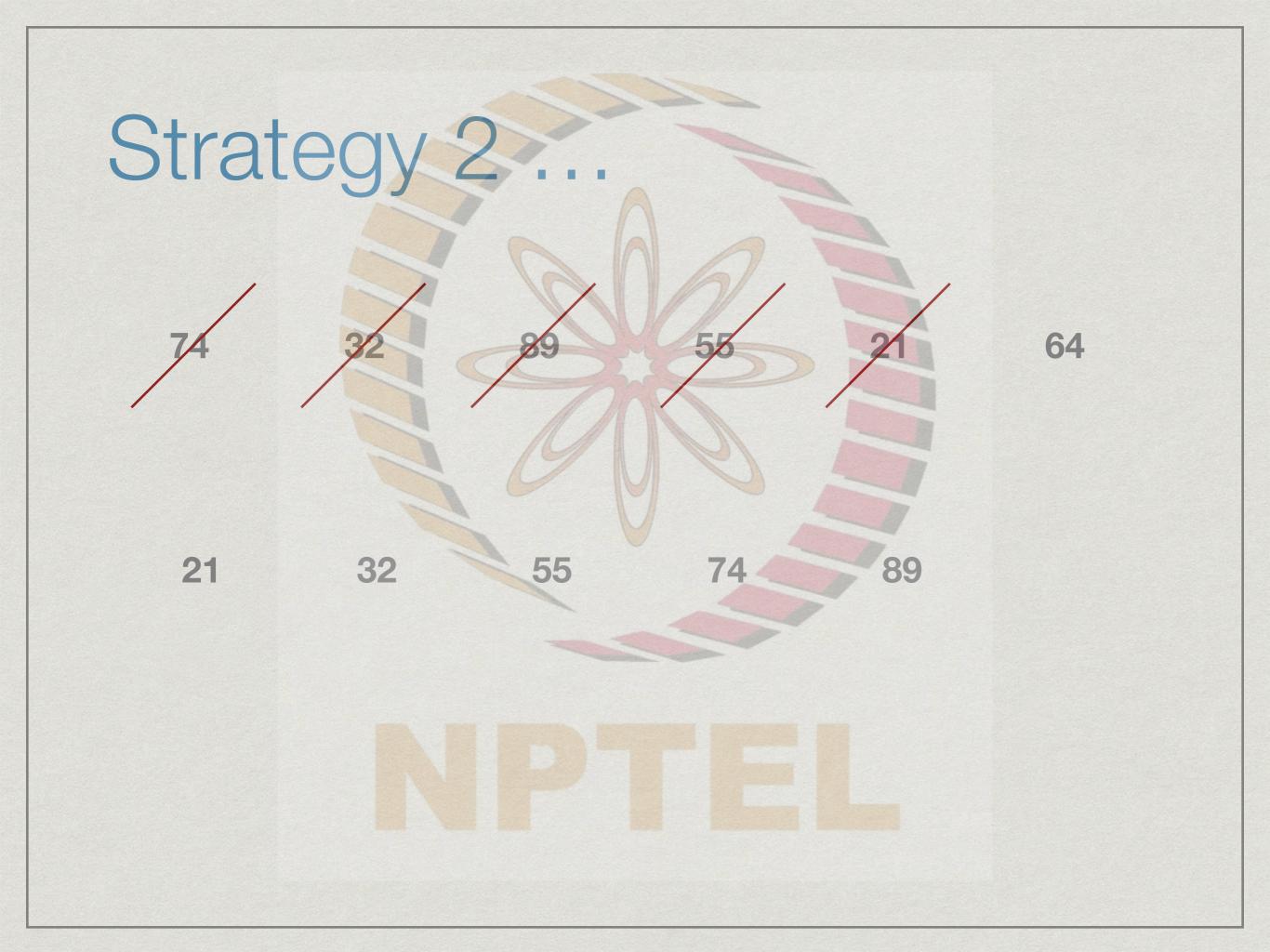


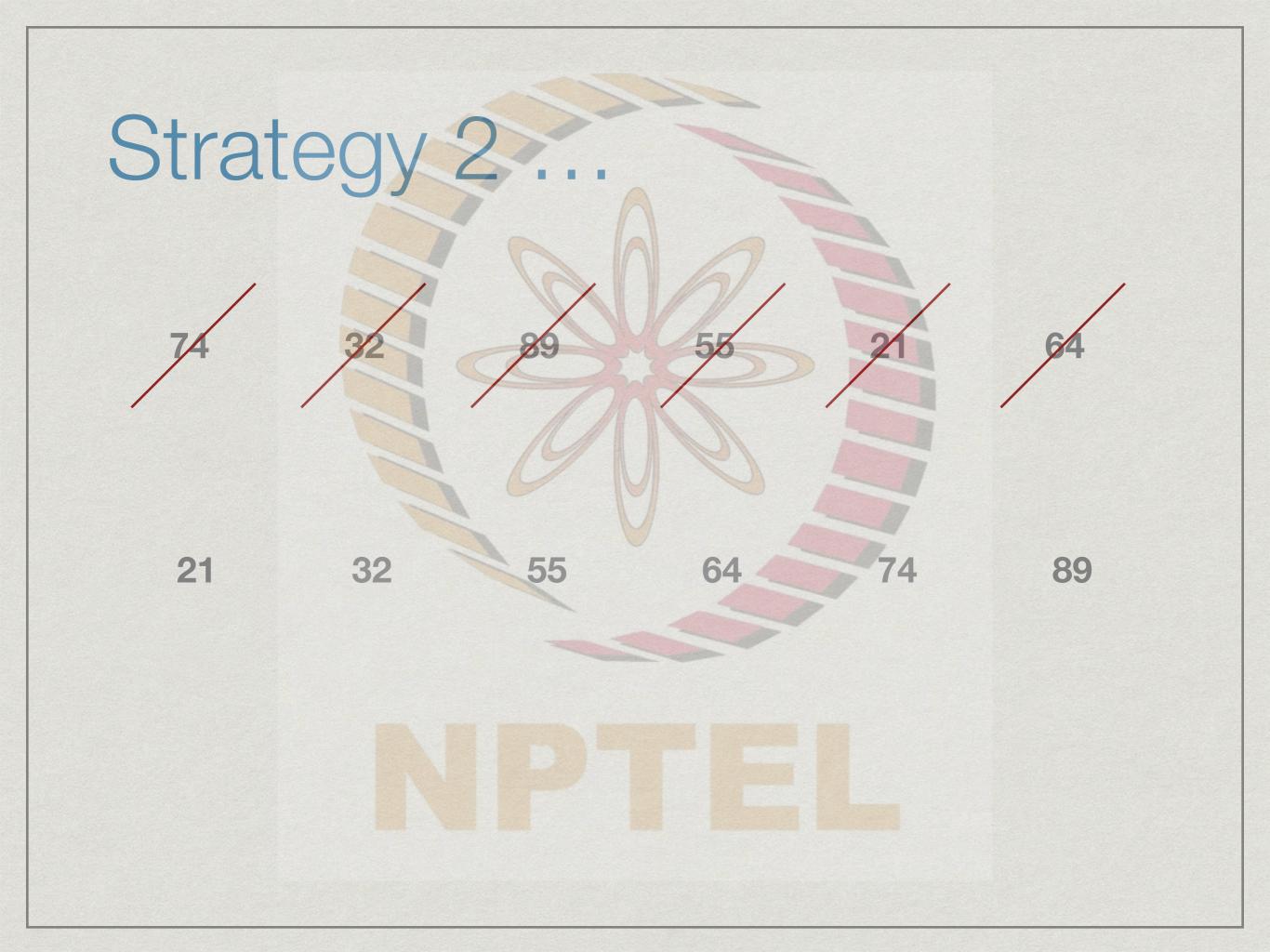












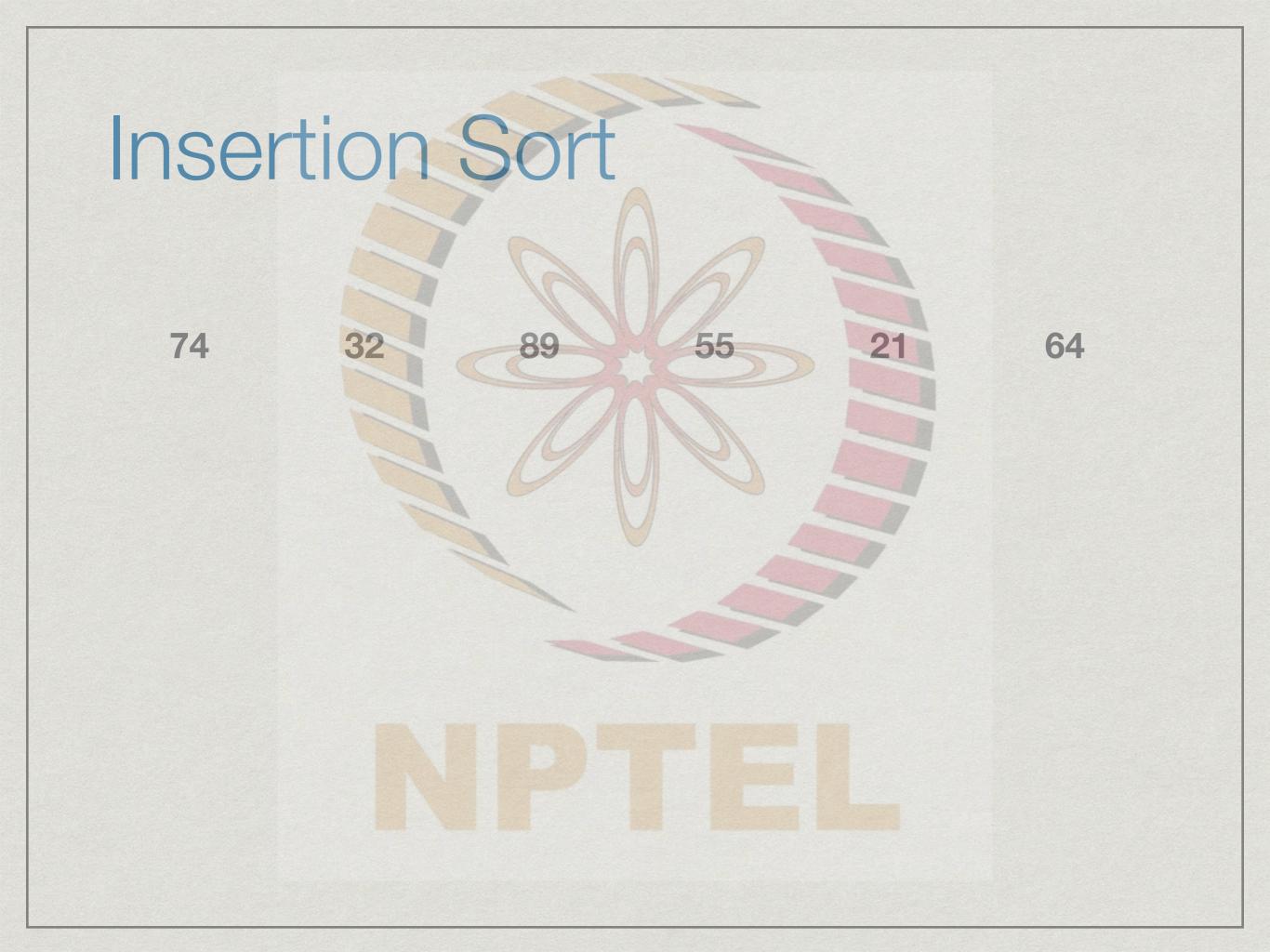
Strategy 2...

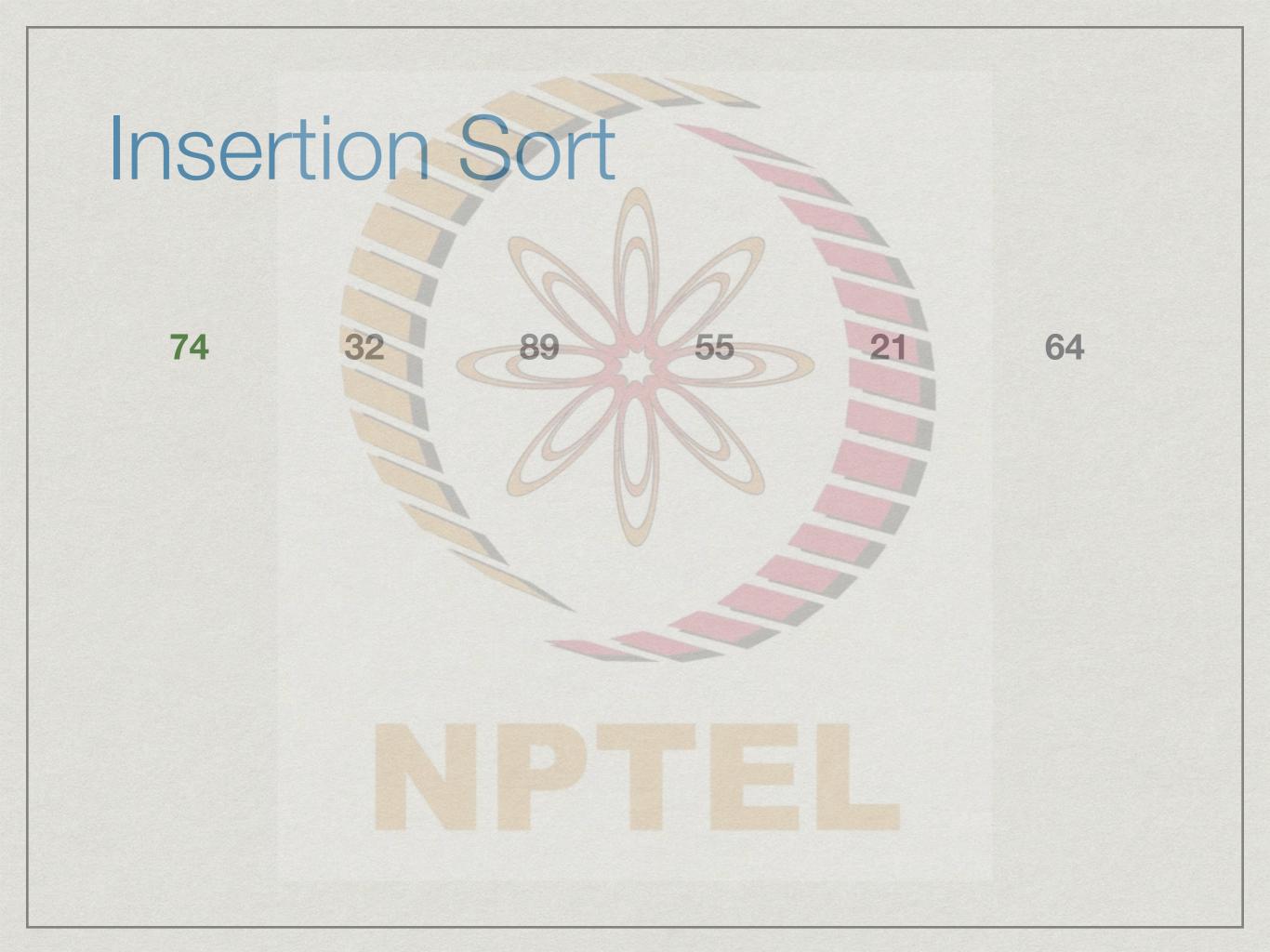
Insertion Sort

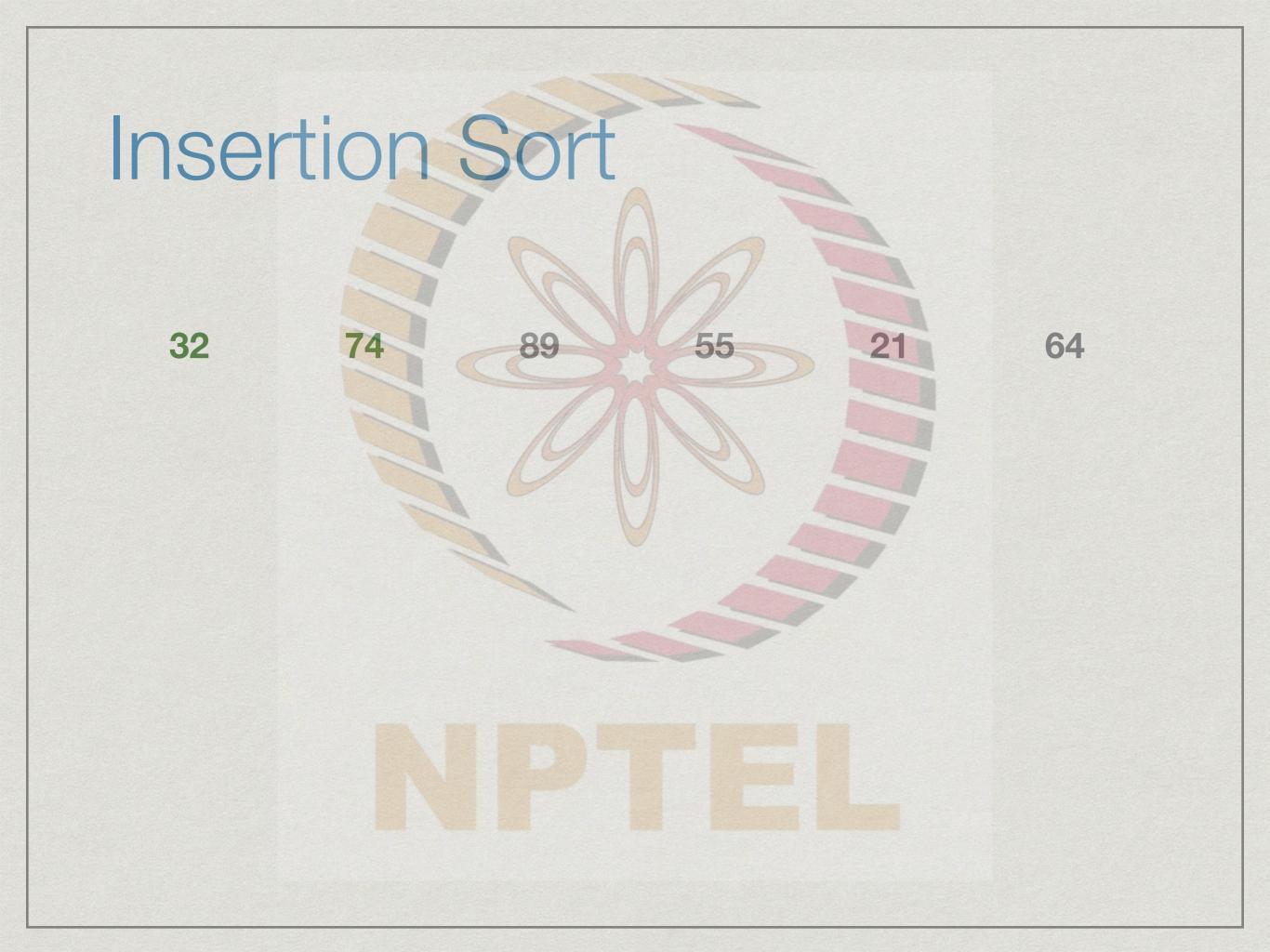
- * Start building a sorted sequence with one element
- * Pick up next unsorted element and insert it into its correct place in the already sorted sequence

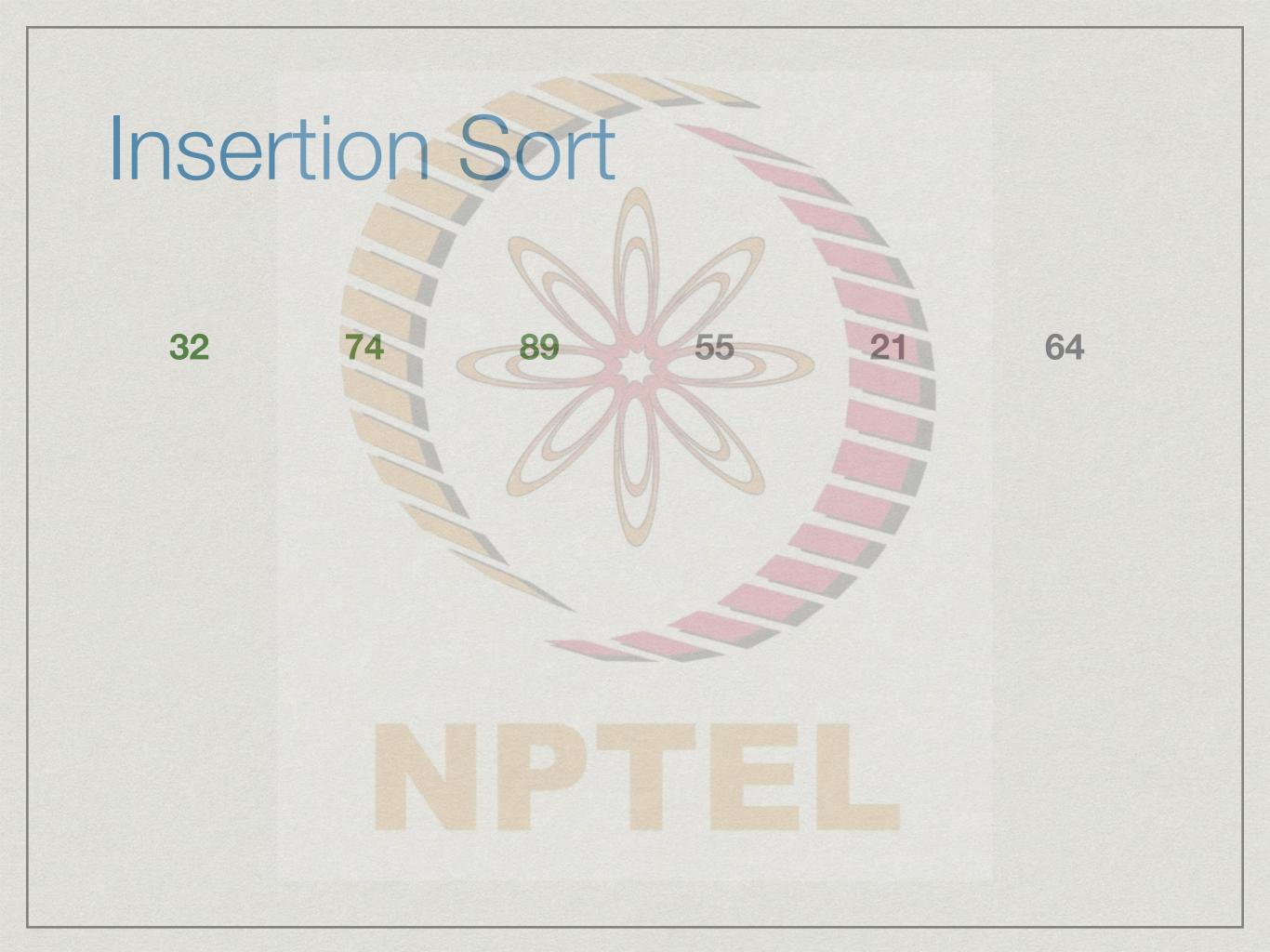
Insertion Sort

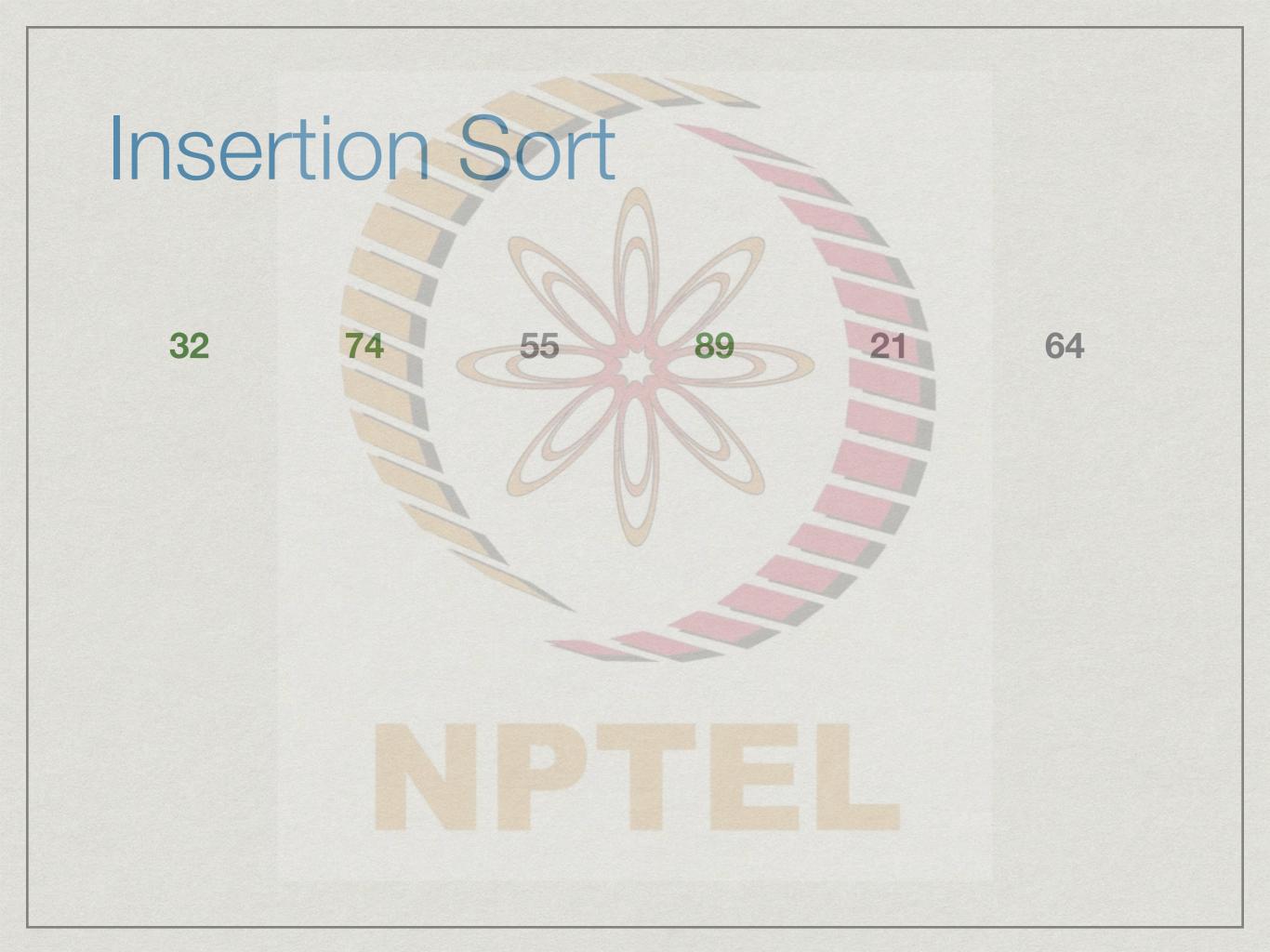
```
InsertionSort(A,n) // Sort A of size n
for (pos = 1; pos < n; pos++)
  // Build longer and longer sorted segments
  // In each iteration A[0]..A[pos-1] is already sorted
  // Move first element after sorted segment left
  // till it is in the correct place
   nextpos = pos
  while (nextpos > 0 &&
           A[nextpos] < A[nextpos-1])
      swap(A, nextpos, nextpos-1)
      nextpos = nextpos-1
```

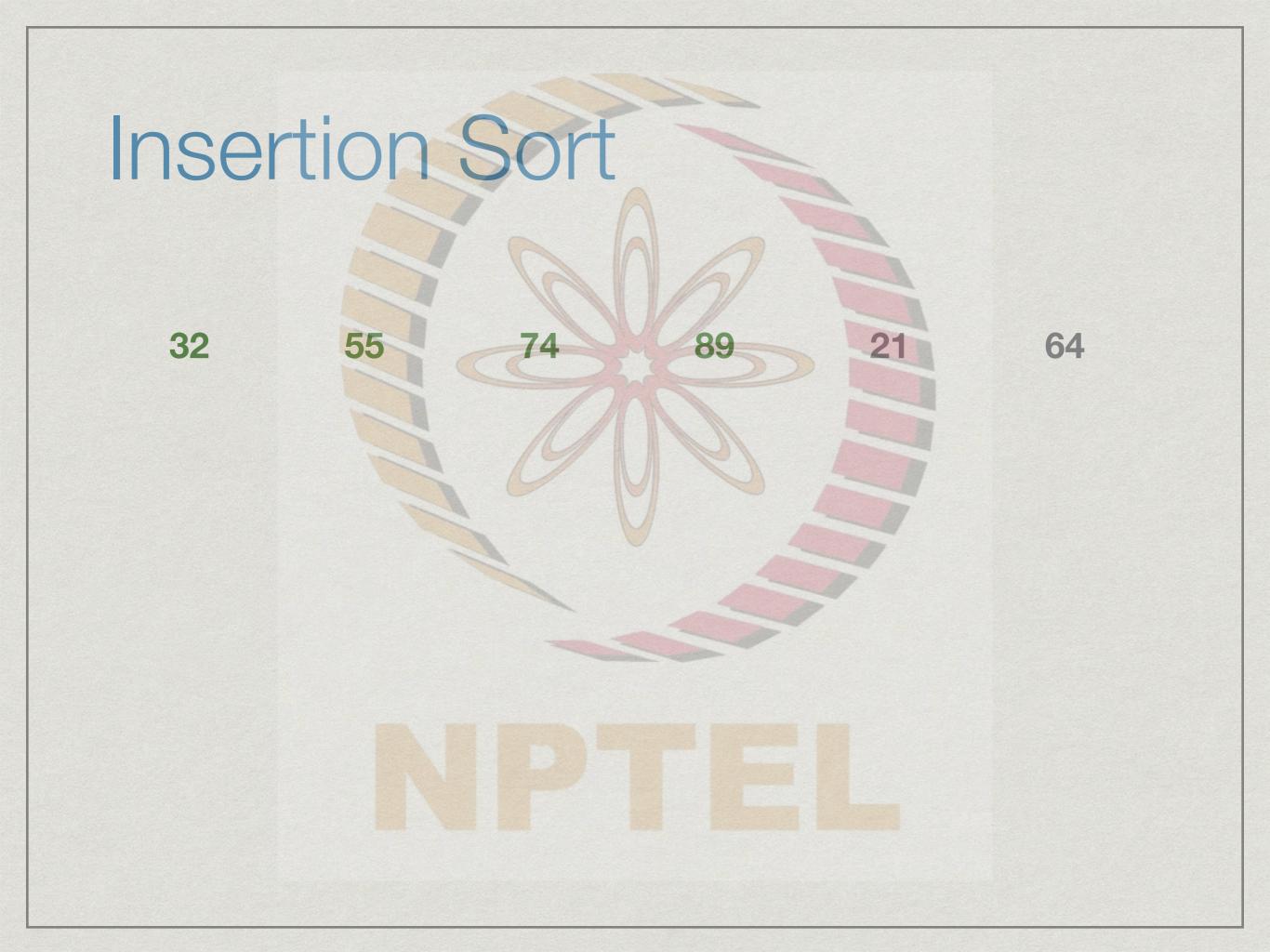


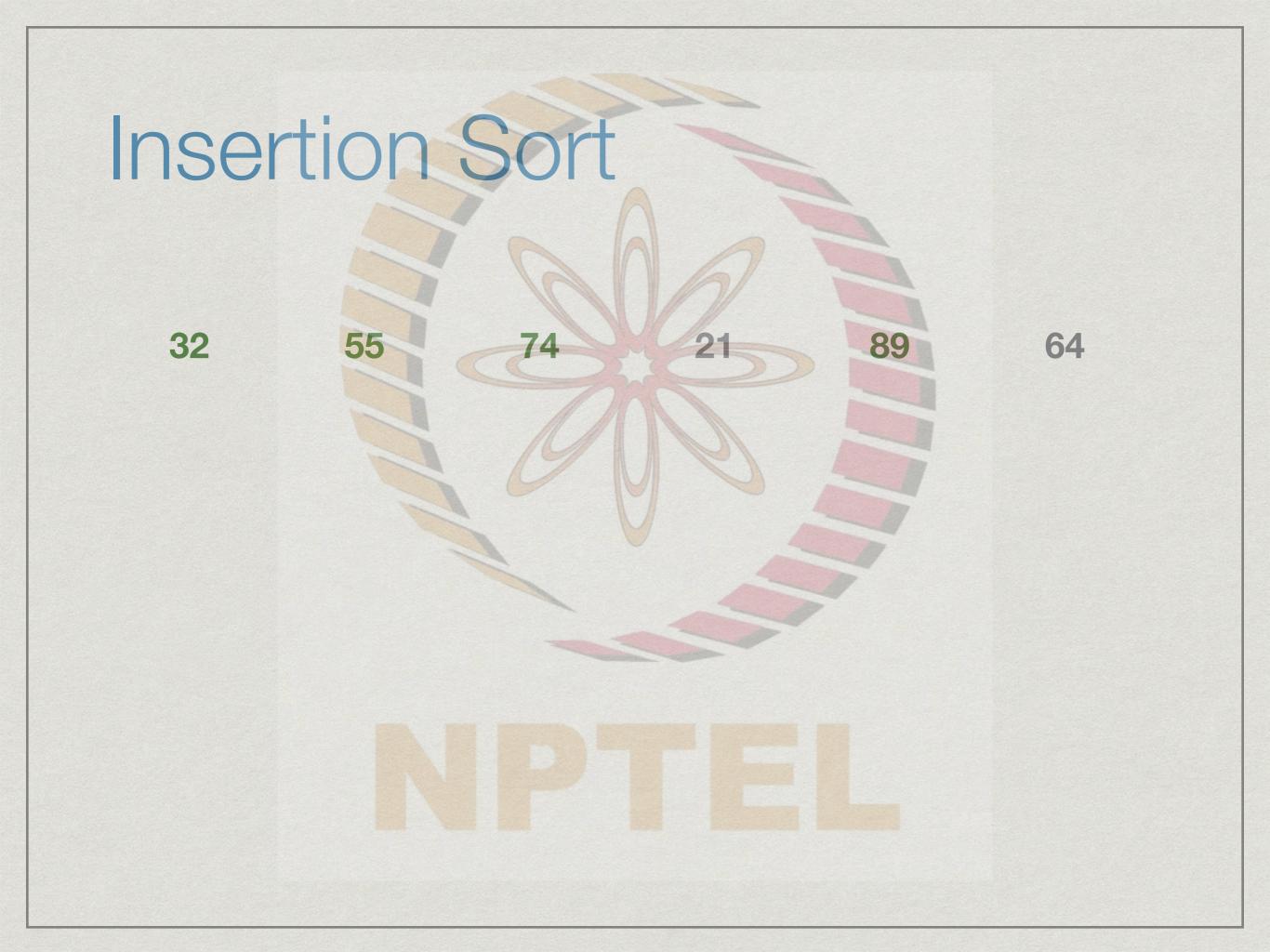


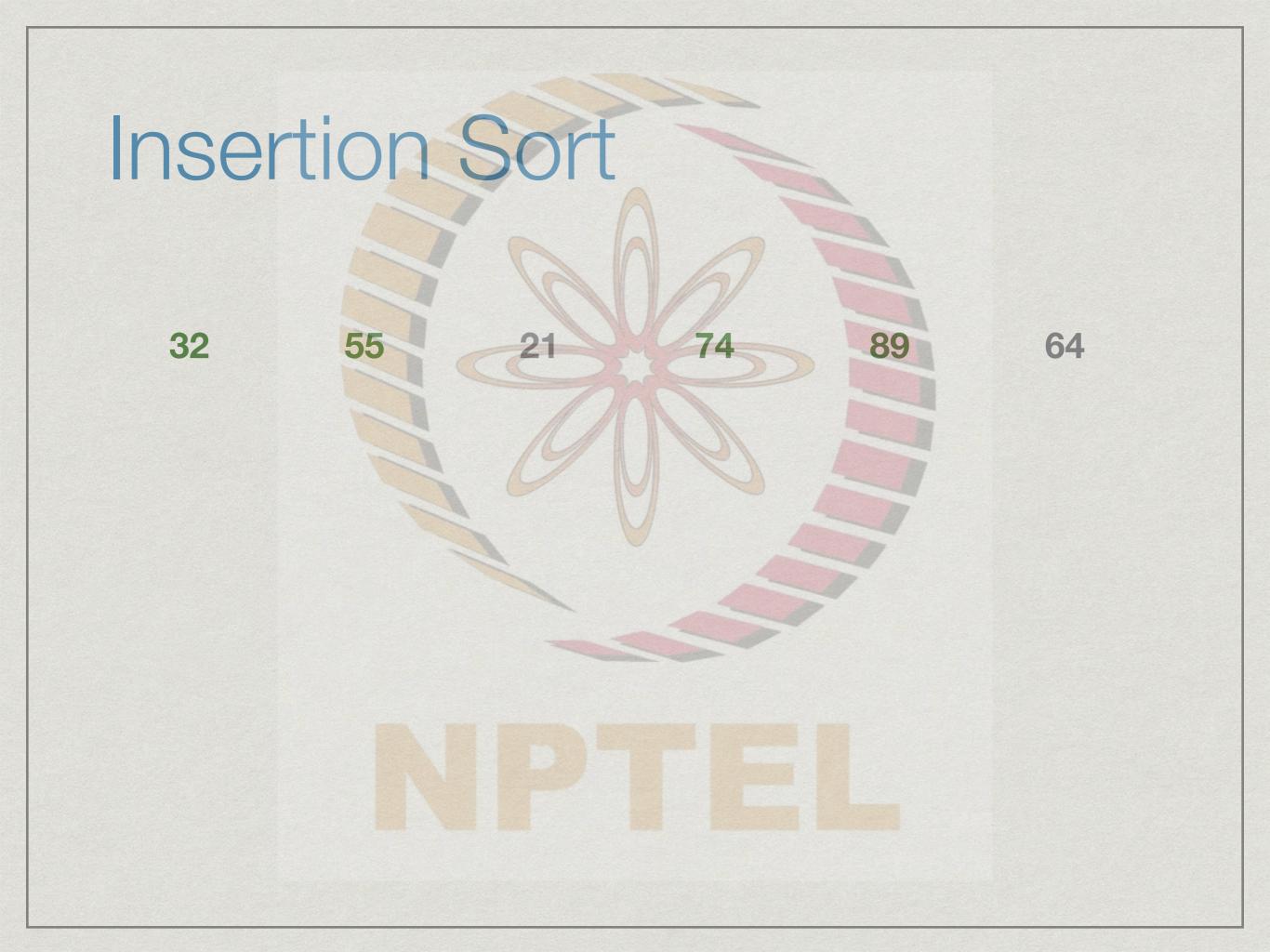


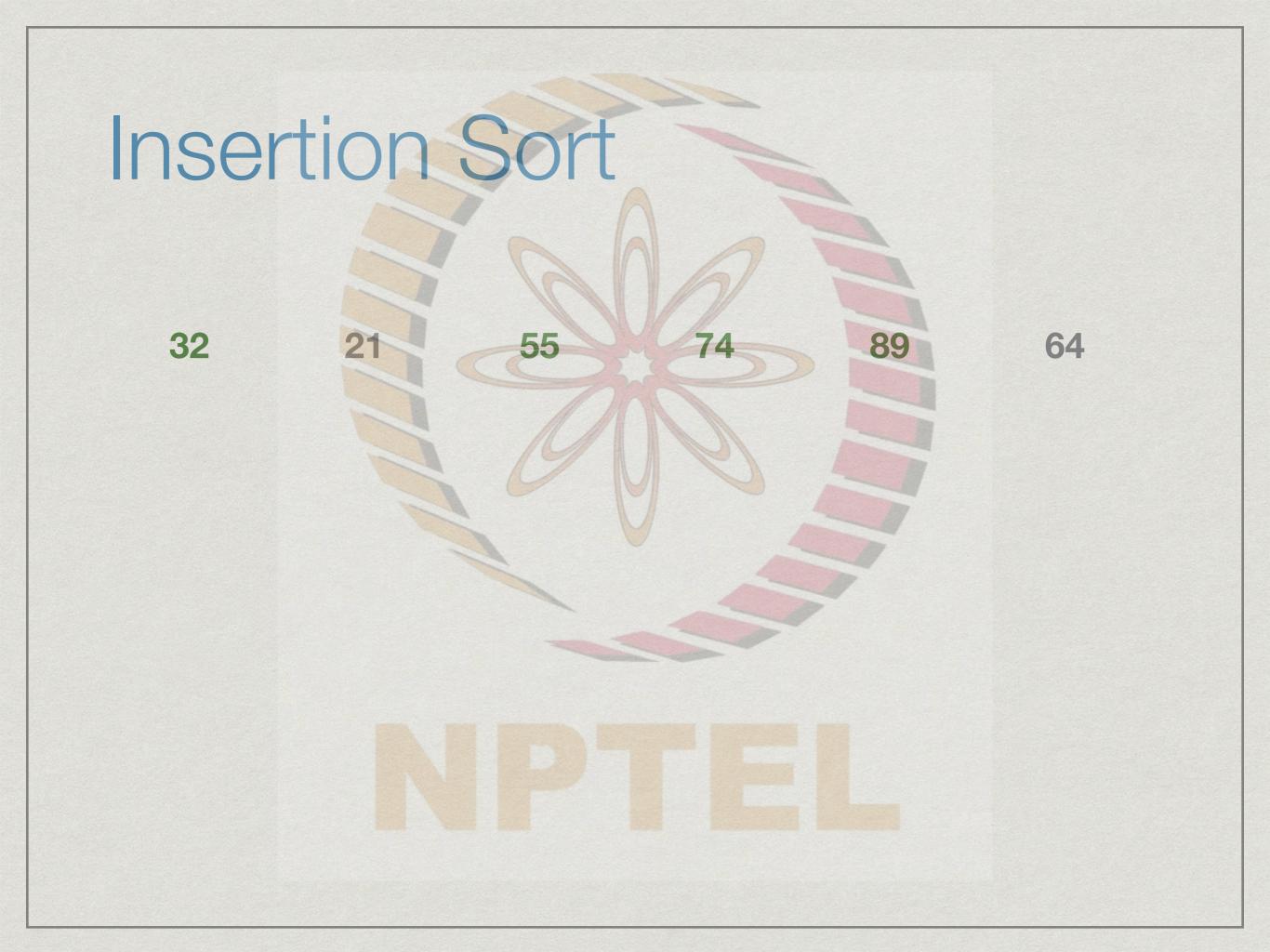


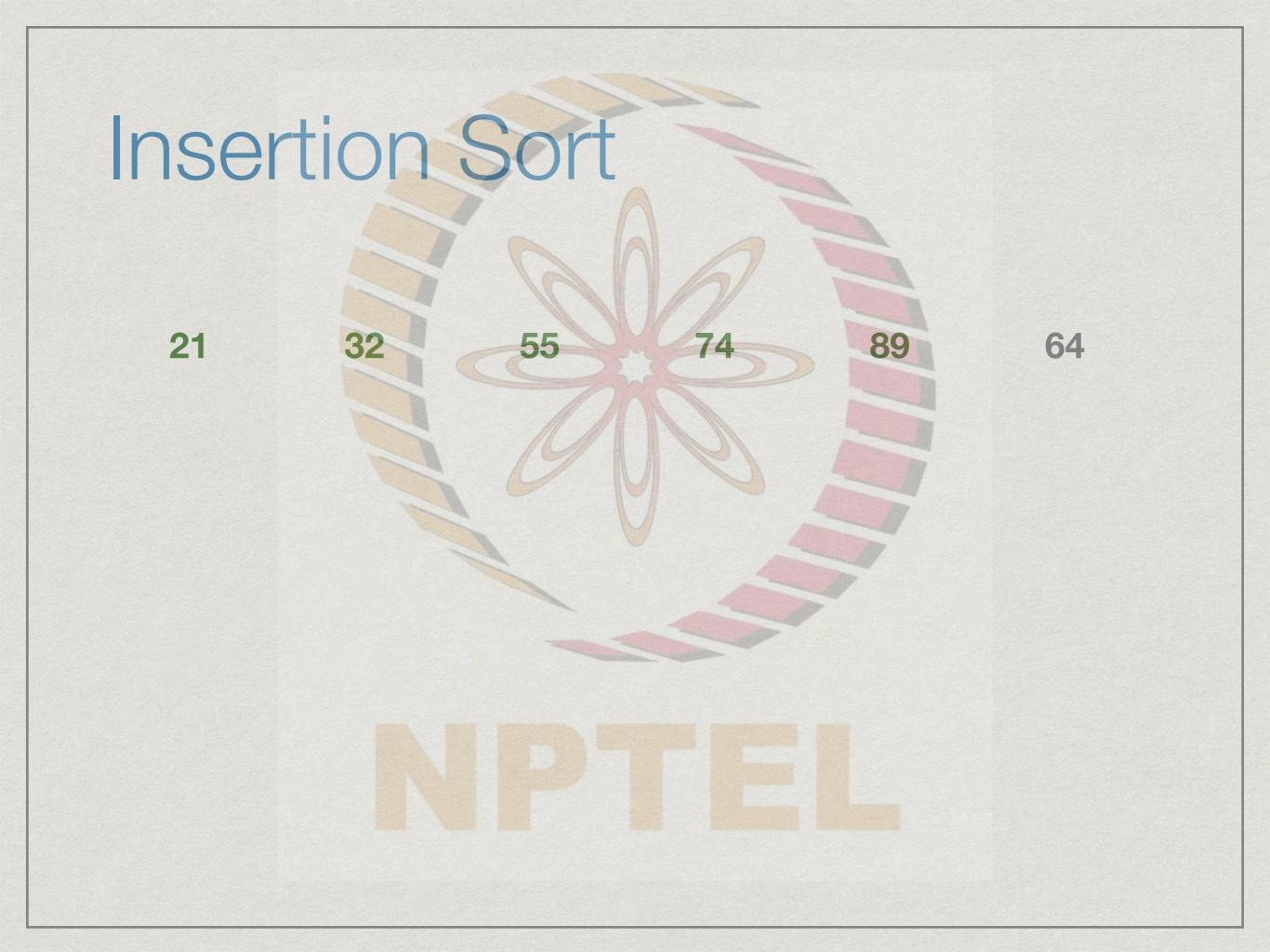


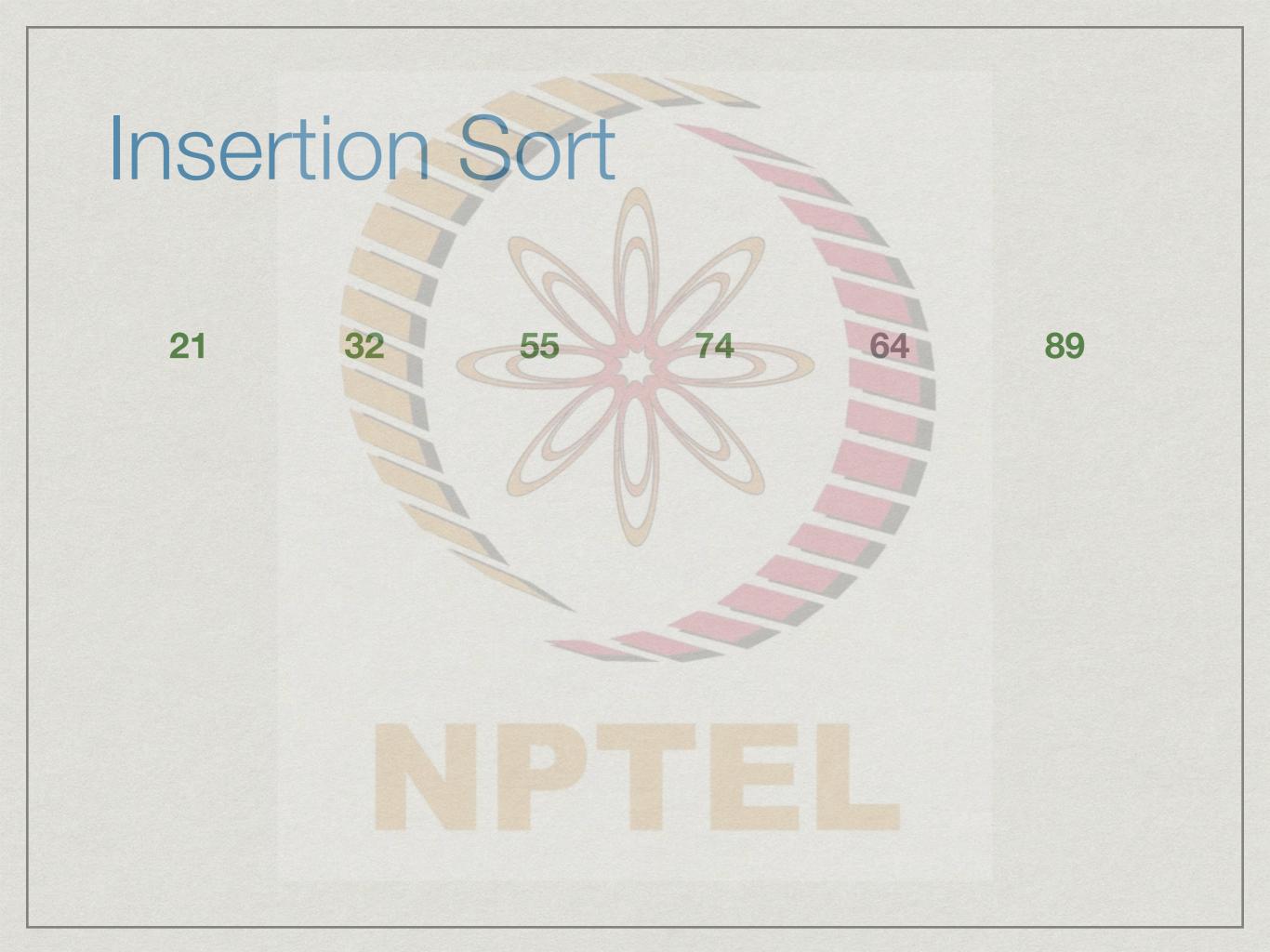


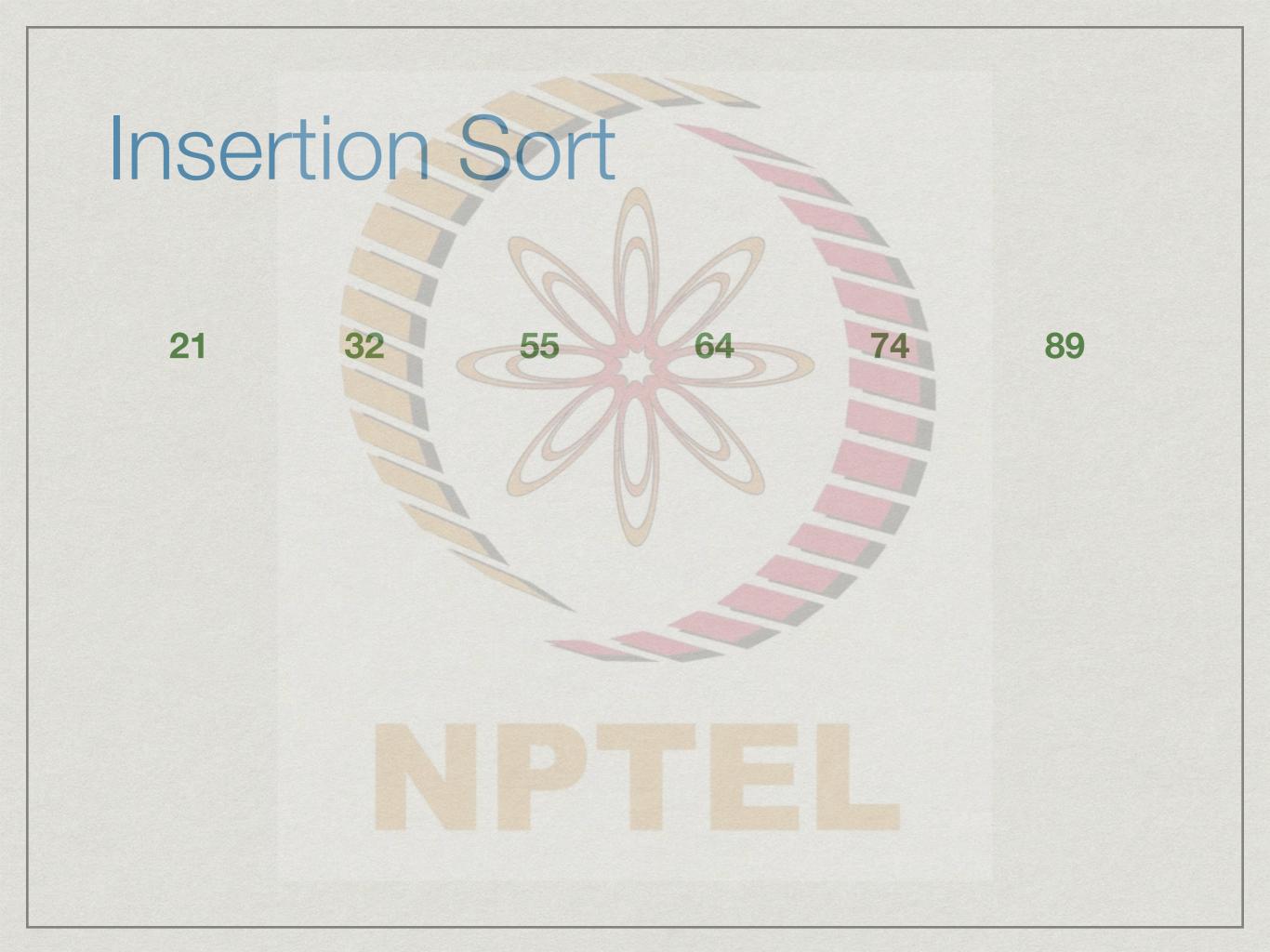












Analysis of Insertion Sort

- * Inserting a new value in sorted segment of length k requires upto k steps in the worst case
- * In each iteration, sorted segment in which to insert increased by 1
- * $t(n) = 1 + 2 + ... + n-1 = n(n-1)/2 = O(n^2)$

Recursive formulation

- * To sort A[0..n-1]
 - * Recursively sort A[0..n-2]
 - * Insert A[n-1] into A[0..n-2]
- * Base case: n = 1

Insertion Sort, recursive

```
InsertionSort(A,k) // Sort A[0..k-1]
  if (k == 1)
    return;
  InsertionSort(A, k-1);
  Insert(A, k-1);
  return;
Insert(A, j) // Insert A[j] into A[0..j-1]
 pos = j;
  while (pos > 0 && A[pos] < A[pos-1])
    swap(A,pos,pos-1);
    pos = pos-1;
```

Recurrence

- * t(n), time to run insertion sort on length n
 - * Time t(n-1) to sort segment A[0] to A[n-2]
 - * n-1 steps to insert A[n-1] in sorted segment

* Recurrence

*
$$t(n) = n-1 + t(n-1)$$

 $t(1) = 1$

*
$$t(n) = n-1 + t(n-1) = n-1 + ((n-2) + t(n-2)) = ... = (n-1) + (n-2) + ... + 1 = n(n-1)/2 = O(n^2)$$

O(n²) sorting algorithms

- * Selection sort and insertion sort are both O(n²)
- * So is bubble sort, which we will not discuss here
- * O(n²) sorting is infeasible for n over 10000
- * Among O(n²) sorts, insertion sort is usually better than selection sort and both are better than bubble sort
 - * What happens when we apply insertion sort to an already sorted list?