

NPTEL MOOC, JAN-FEB 2015  
Week 4, Module 5

# DESIGN AND ANALYSIS OF ALGORITHMS

Minimum cost spanning trees

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# Example: Road network

- \* District hit by a cyclone, damaging the roads
- \* Government sets to work to restore the roads
- \* Priority is to ensure that all parts of the district can be reached
- \* What set of roads should be restored first?



# Spanning tree

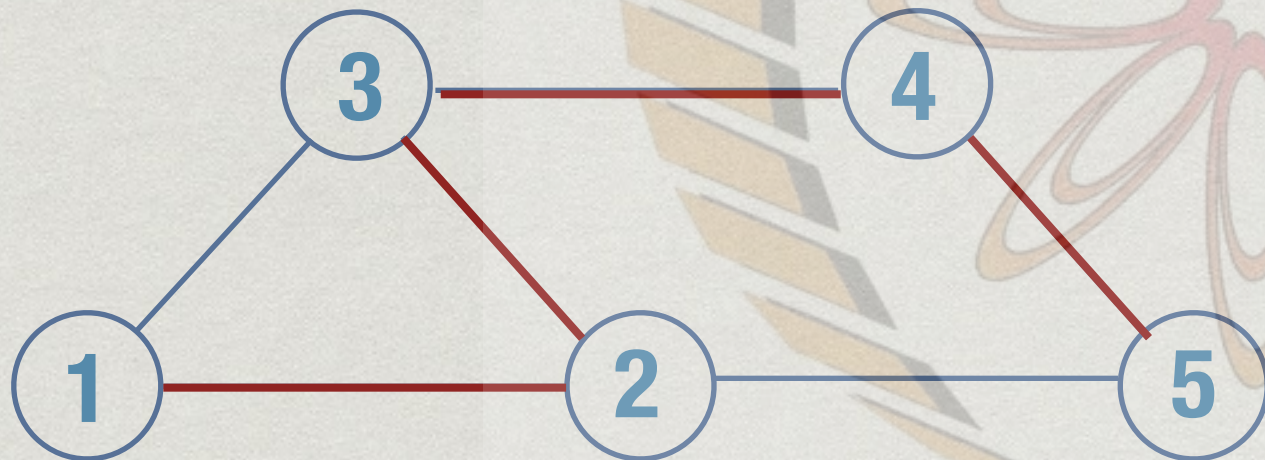


- \* Minimum connectivity:  
no loops
- \* Removing an edge  
from a loop cannot  
disconnect graph
- \* Connected acyclic  
graph — **tree**
- \* Spanning tree

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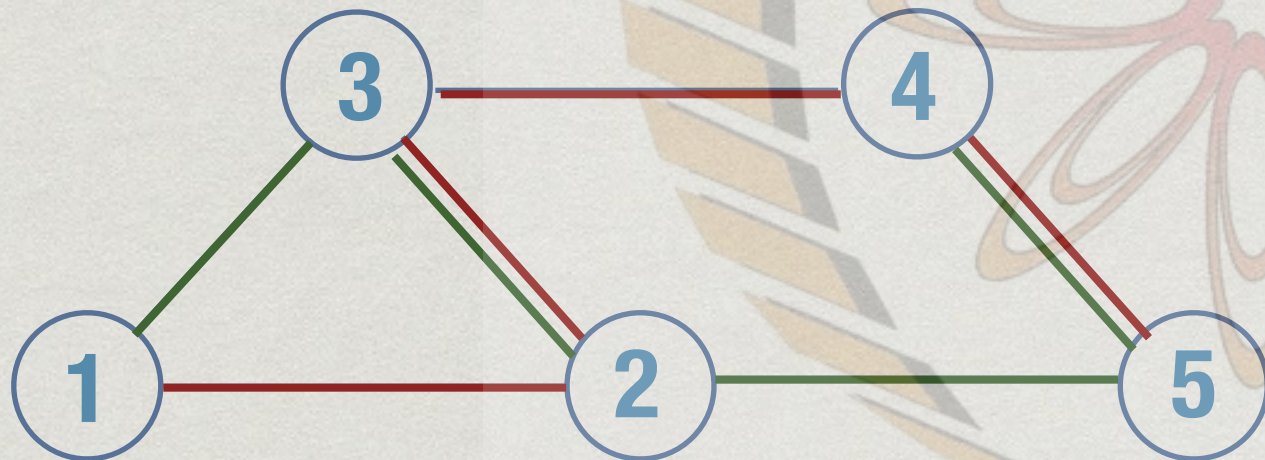
# Spanning tree



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# Spanning tree

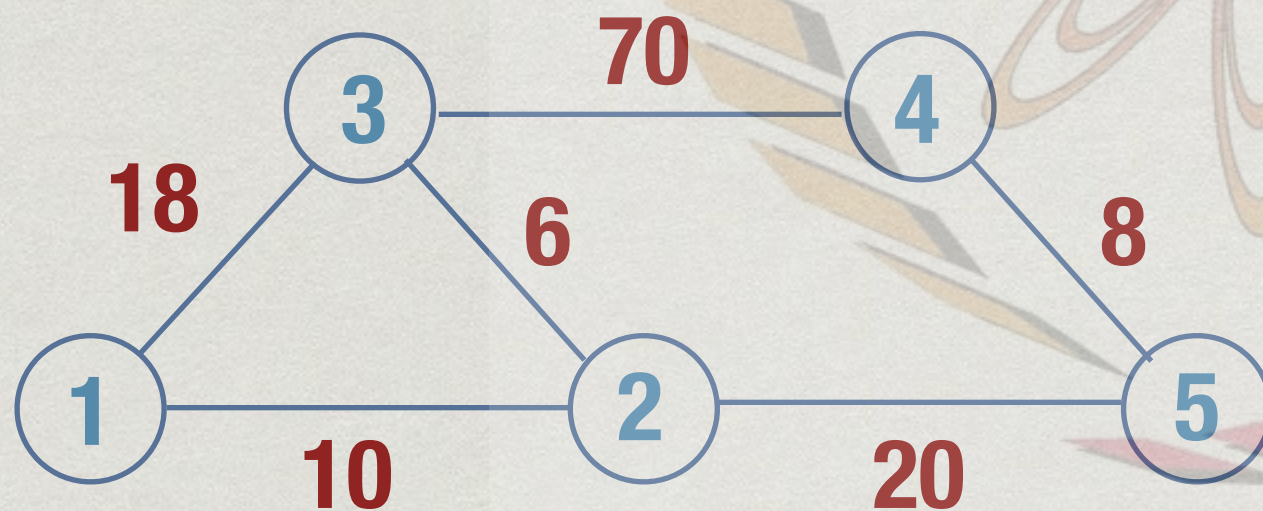


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# Spanning tree with costs



- \* Restoration of each road has a cost

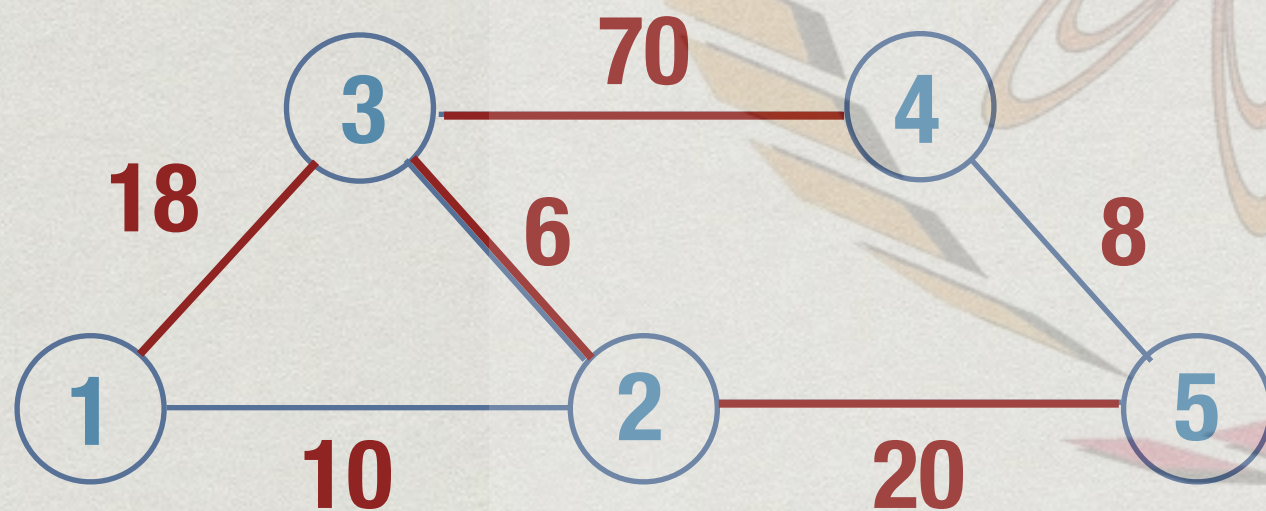
- \* Among the different spanning trees, choose the one with minimum cost

- \* Minimum cost spanning tree



# Spanning tree with costs

**Cost = 114**



- \* Restoration of each road has a cost

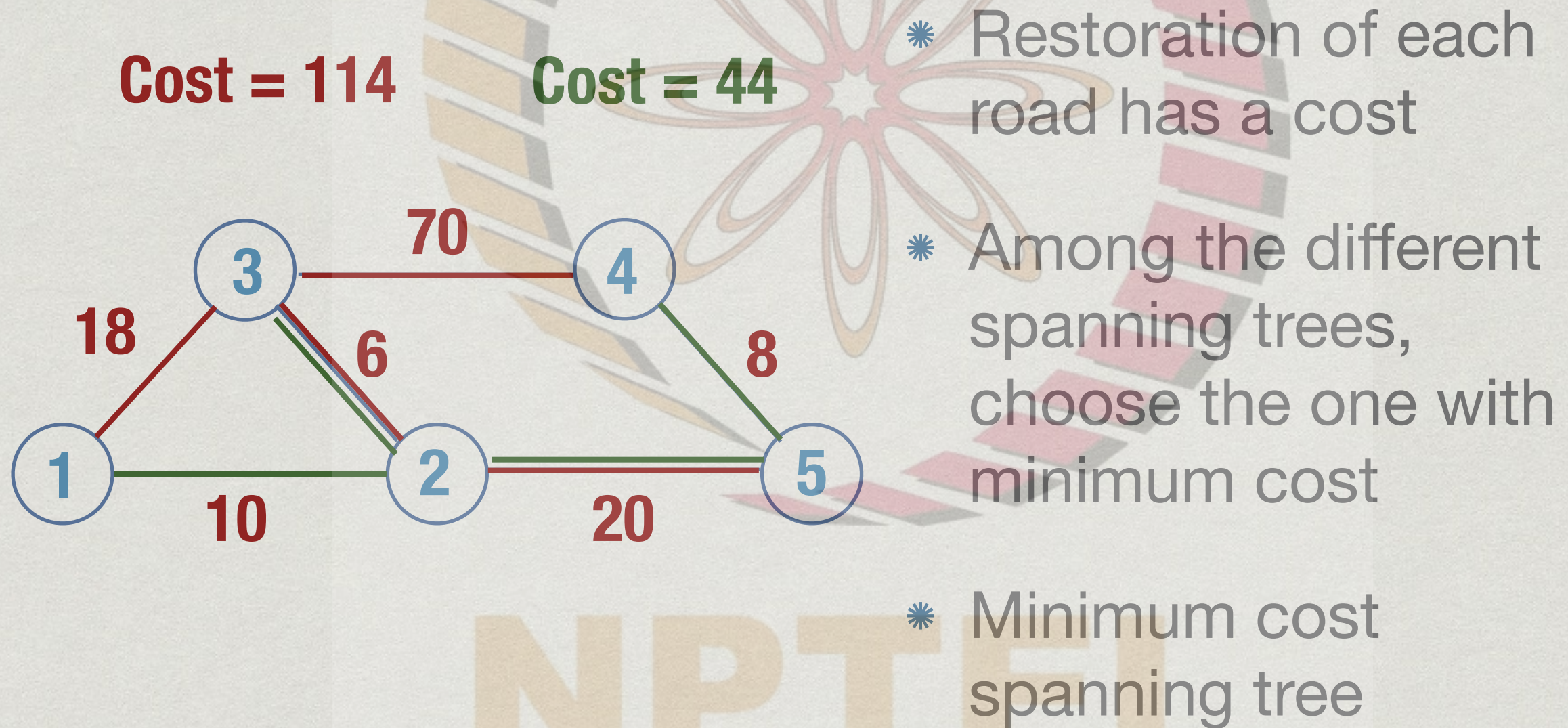
- \* Among the different spanning trees, choose the one with minimum cost

- \* Minimum cost spanning tree

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# Spanning tree with costs





# Facts about trees

**Definition:** A tree is a connected acyclic graph

**Fact 1:** A tree on  $n$  vertices has exactly  $n-1$  edges

- \* Start with a tree and delete edges
- \* Initially one single component
- \* Deleting an edge must split a component into two
- \* After  $n-1$  edge deletions,  $n$  components, each an isolated vertex



# Facts about trees

**Fact 2:** Adding an edge to a tree must create a cycle

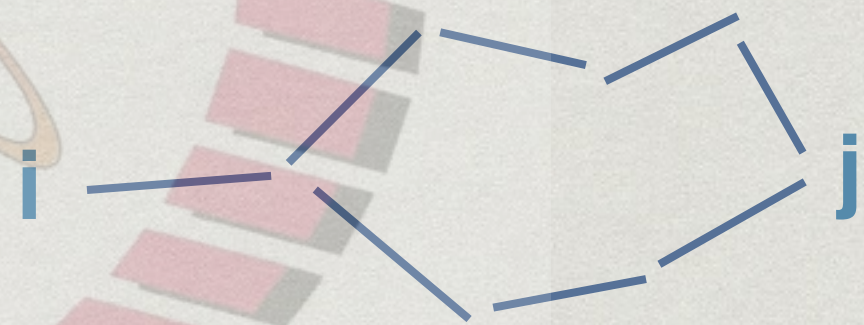
- \* Suppose we add an edge  $(i,j)$
- \* Tree is connected, so there is already a path  $p$  from  $i$  to  $j$
- \* New edge  $(i,j)$  plus path  $p$  creates a cycle



# Facts about trees

**Fact 3:** In a tree, every pair of nodes is connected by a unique path

- \* If there are two paths from  $i$  to  $j$ , there must be a cycle





# Facts about trees

Any two of the following facts about a graph  $G$  implies the third

- \*  $G$  is connected
- \*  $G$  is acyclic
- \*  $G$  has  $n-1$  edges

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# Building a minimum cost spanning trees

## Two natural strategies

- \* Start with smallest edge and grow it into a tree

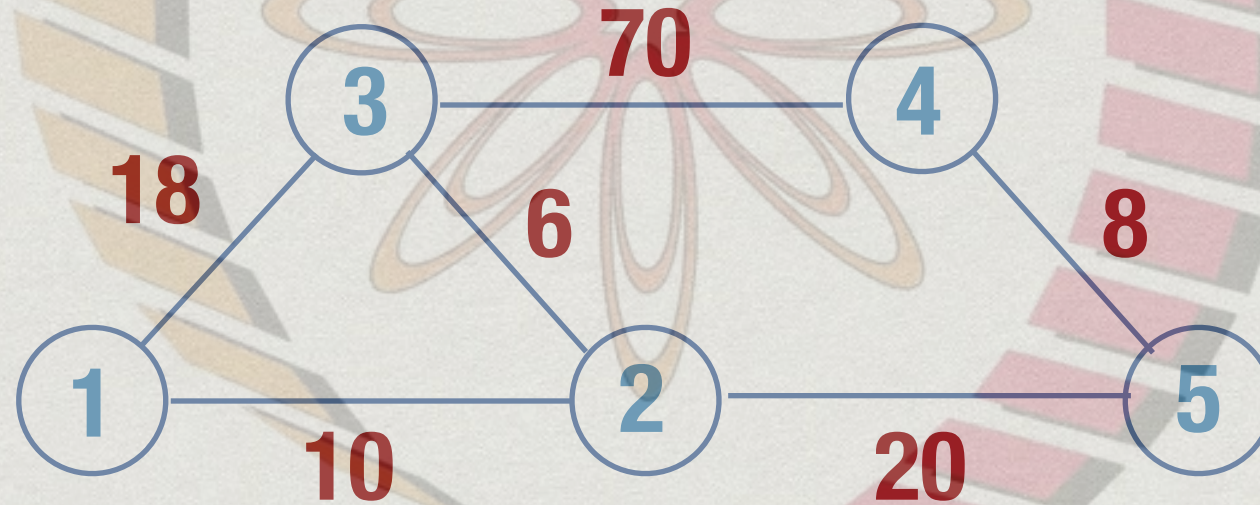
### Prim's Algorithm

- \* Scan edges in ascending order of cost and connect components to form a tree

### Kruskal's Algorithm



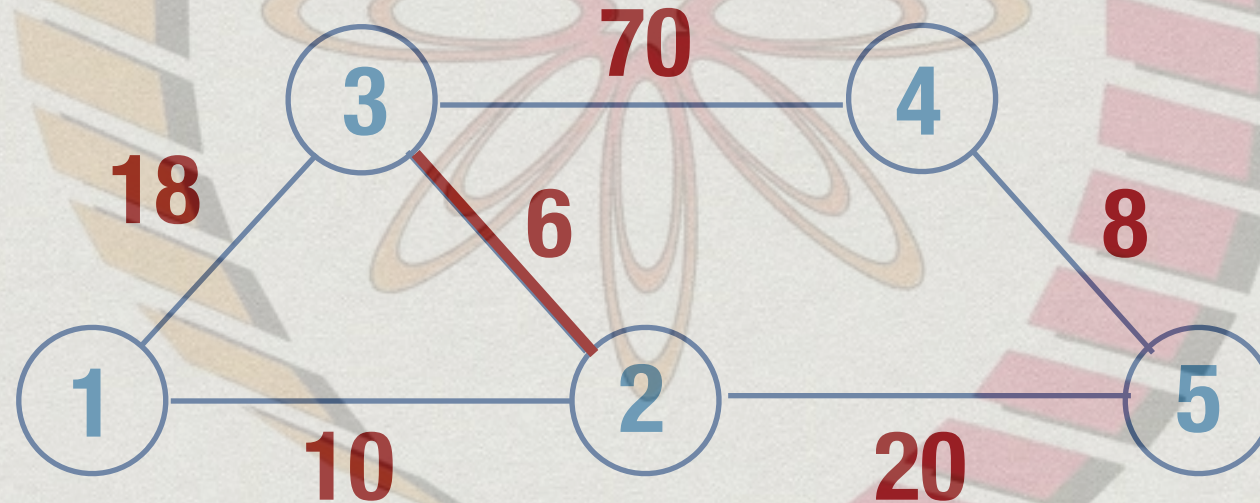
# Prim's algorithm



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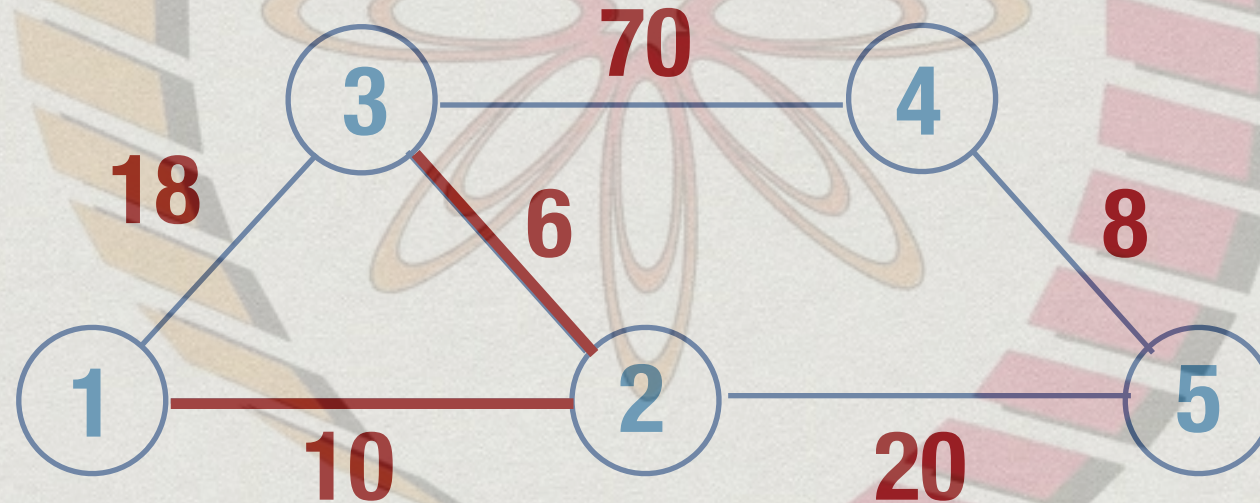
# Prim's algorithm



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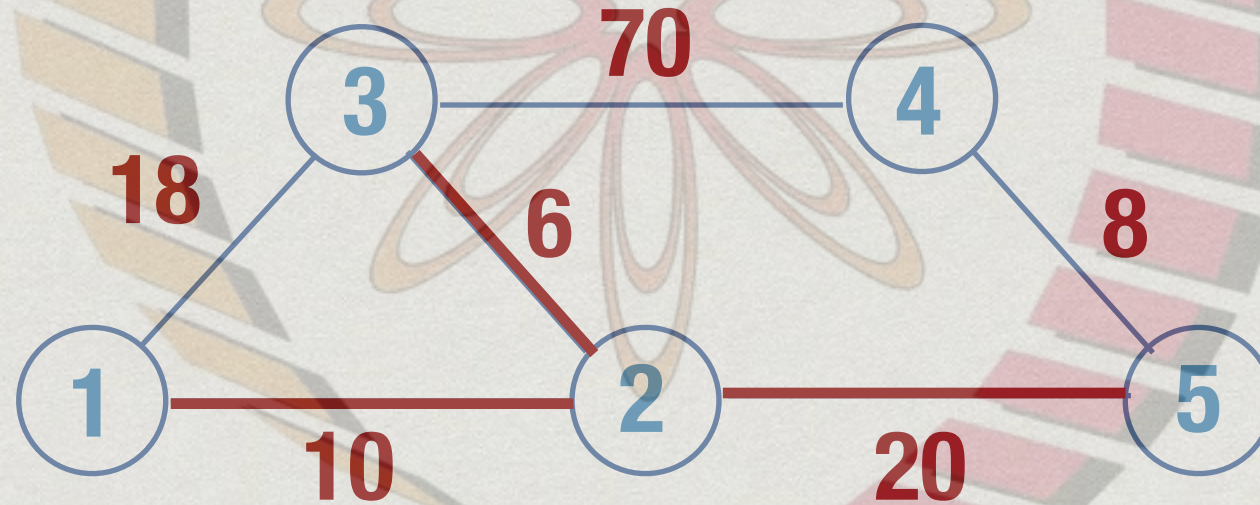


# Prim's algorithm



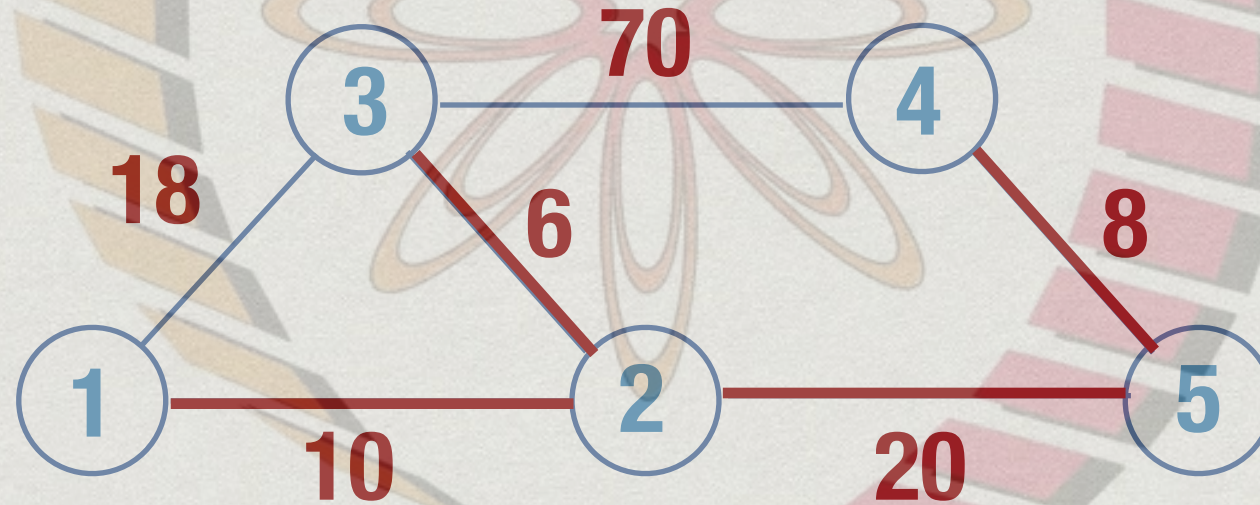


# Prim's algorithm



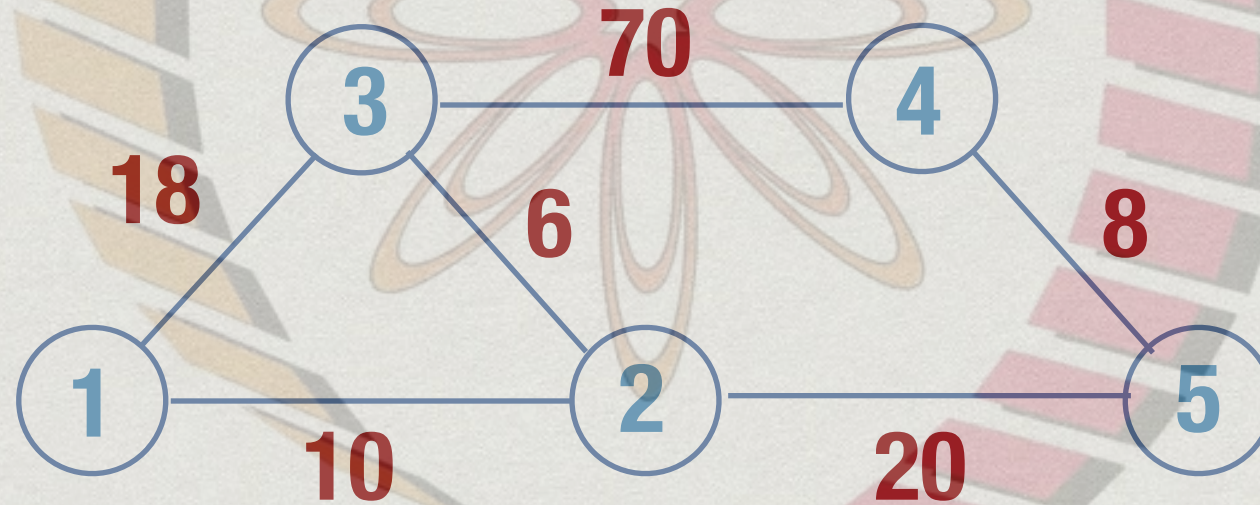


# Prim's algorithm





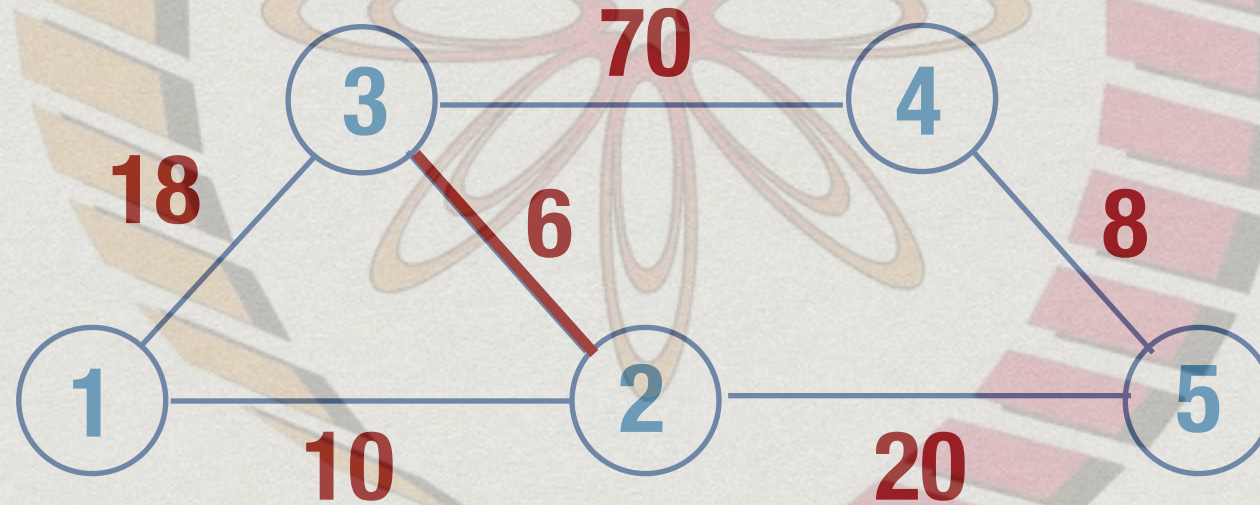
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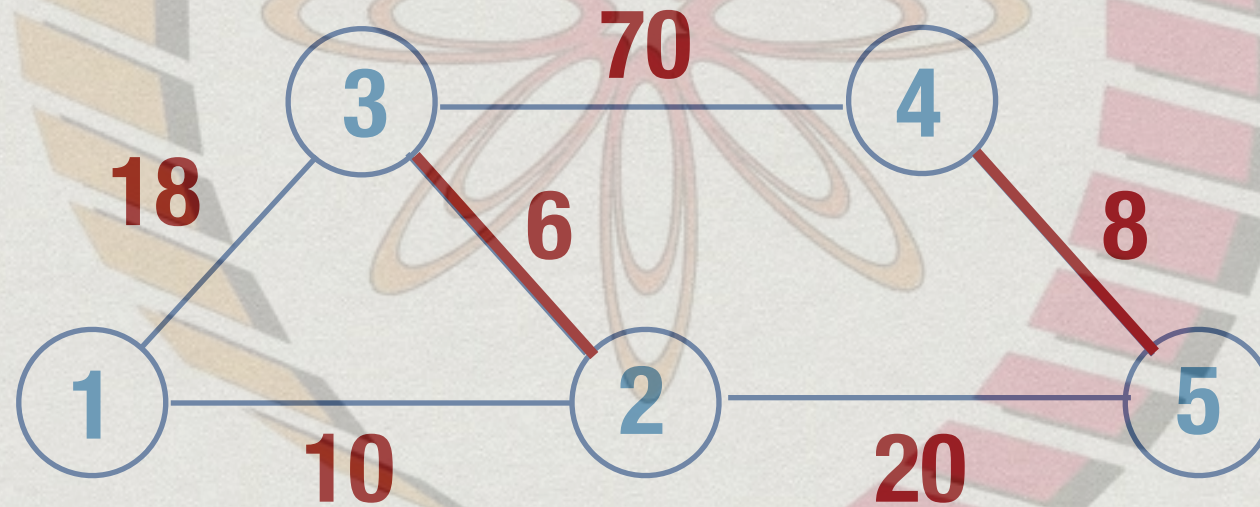
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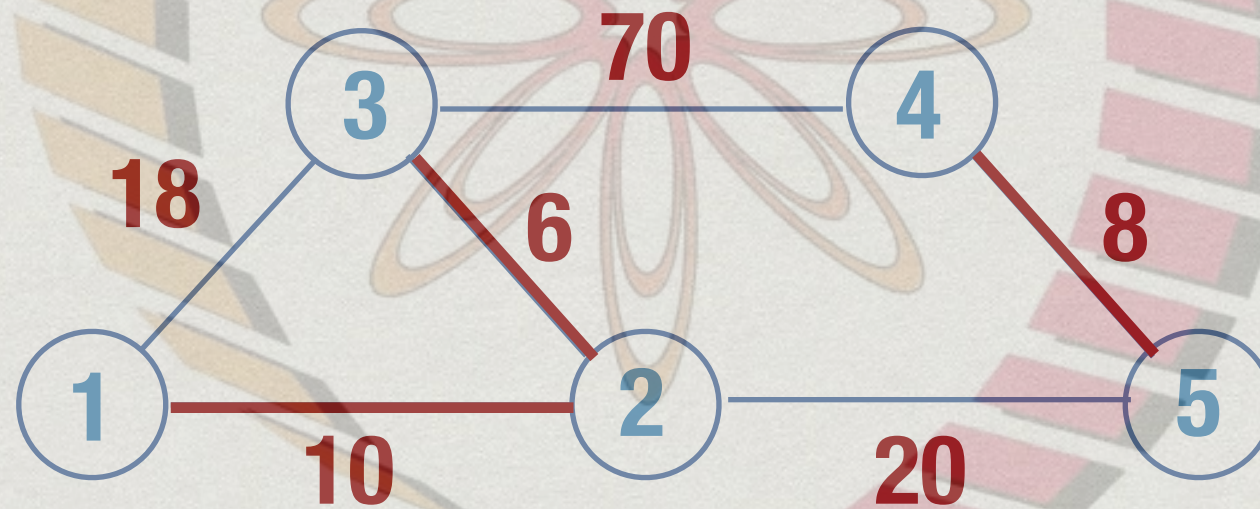
# Kruskal's algorithm



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# Kruskal's algorithm



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# Kruskal's algorithm

