

## 1 Polynomial Practice

$$\underline{(x^2+1)(x^2+1)}$$

- Note 8** (a) If  $f$  and  $g$  are non-zero real polynomials, how many real roots do the following polynomials have at least? How many can they have at most? (Your answer may depend on the degrees of  $f$  and  $g$ .)

(i)  $f + g$

(ii)  $f \cdot g$

(iii)  $f/g$ , assuming that  $f/g$  is a polynomial

(ii)  $\underline{(x^2+1)(x^2+1)}$  zero

degree(f) + degree(g) at most

(i) if  $f+g=0$  int

(iii) at least zero

$f+g$  is odd, at least one

$f+g = x^2+1$ , no one more  $\max(d(f), d(g))$  at most  $\deg(f) - \deg(g)$

- (b) Now let  $f$  and  $g$  be polynomials over  $GF(p)$

- (i) We say a polynomial  $f = 0$  if  $\forall x, f(x) = 0$ . Show that if  $f \cdot g = 0$ , it is not always true that either  $f = 0$  or  $g = 0$ .

- (ii) How many  $f$  of degree exactly  $d < p$  are there such that  $f(0) = a$  for some fixed  $a \in \{0, 1, \dots, p-1\}$ ?

(i)  $p=2$   $f=x$   $g=x-1$   $\rightarrow 0, 1$   $\rightarrow$  Fermat's little theorem

$$x^{p-1} \equiv 1 \pmod{p}$$

(ii)  $p^{d-1} \rightarrow f(x) = \sum_{k=0}^d c_k x^k$

$$x \neq 0$$

- (c) Find a polynomial  $f$  over  $GF(5)$  that satisfies  $f(0) = 1, f(2) = 2, f(4) = 0$ . How many such polynomials of degree at most 4 are there?

$c_0 = a$ ,  $c_d \neq 0$ ,  $(p-1)p^{d-1}$

degree : 3  $\leftarrow \frac{2(x-2)(x-4)}{(x-2)(x-4)} \leftarrow$  在 GF 下的变形

$$\Delta_1(x) = \frac{(x-2)(x-4)}{(c_0-2)(0-4)} = \frac{(x-2)(x-4)}{4}$$

$$f(x) = \frac{(x-4)(x-4)}{8} - \frac{4x(x-4)}{8}$$

$$= \frac{(-3x-2)(x-4)}{8}$$

$$\Delta_2(x) = \frac{x(x-4)}{2 \neq -2} = -\frac{x(x-4)}{4}$$

$$\Delta_4(x) = \frac{x(x-2)}{4 \neq 2} = \frac{x(x-2)}{2}$$

5个解,  $\sum_{k=0}^d c_k x^k$   $c_d \neq 0$ ,  $c_0 = 1$ ,  $2(x^2 - 6x + 8) + 2x^2 - 8x$   
 $2x^2 - 12x + 16 + 2x^2 - 8x$

$$\underline{4x^2 - 20x + 16} \equiv \underline{4x^2 + 1}$$

## 2 Lagrange Interpolation in Finite Fields

**Note 8** Find a unique polynomial  $p(x)$  of degree at most 2 that passes through points  $(-1, 3)$ ,  $(0, 1)$ , and  $(1, 2)$  in modulo 5 arithmetic using the Lagrange interpolation.

- (a) Find  $p_{-1}(x)$  where  $p_{-1}(0) \equiv p_{-1}(1) \equiv 0 \pmod{5}$  and  $p_{-1}(-1) \equiv 1 \pmod{5}$ .

$$P_{-1}(x) = k(x)(x-1) \quad P_{-1}(-1) \equiv 2k \equiv 1 \pmod{5}$$

$$[(1-0) * (-1-1)]^{-1} = 2^{-1} \quad \frac{2*3-5}{2-3} = 1$$

$$P_{-1}(x) = 3x(x-1) = 3x^2 - 3x$$

- (b) Find  $p_0(x)$  where  $p_0(-1) \equiv p_0(1) \equiv 0 \pmod{5}$  and  $p_0(0) \equiv 1 \pmod{5}$ .

$$P_0(x) = k(x+1)(x-1) \quad -k \equiv 1 \pmod{5} \quad \rightarrow k = -1$$

$$P_0(x) = -(x^2 - 1) = 1 - x^2 \quad (x+1)(x-1) \quad (-1)^{-1}$$

- (c) Find  $p_1(x)$  where  $p_1(-1) \equiv p_1(0) \equiv 0 \pmod{5}$  and  $p_1(1) \equiv 1 \pmod{5}$ .

$$P_1(x) = k(x+1)x \quad 2k \equiv 1 \pmod{5} \quad k = 3$$

$$P_1(x) = 3x(x+1) = 3x^2 + 3x$$

- (d) Construct  $p(x)$  using a linear combination of  $p_{-1}(x)$ ,  $p_0(x)$ , and  $p_1(x)$ .

$$\begin{aligned} P(x) &= \underbrace{3P_{-1}(x)}_{\equiv 3(3x^2 - 3x)} + P_0(x) + 2P_1(x) \\ &\equiv 3(3x^2 - 3x) + 1 - x^2 + 2(3x^2 + 3x) \\ &\equiv (4x^2 - 3x + 1) \equiv 4x^2 - 3x + 1 \end{aligned}$$

## 3 Secrets in the United Nations

**Note 8** A vault in the United Nations can be opened with a secret combination  $s \in \mathbb{Z}$ . In only two situations should this vault be opened: (i) all 193 member countries must agree, or (ii) at least 55 countries, plus the U.N. Secretary-General, must agree.

- (a) Propose a scheme that gives private information to the Secretary-General and all 193 member countries so that the secret combination  $s$  can only be recovered under either one of the two specified conditions.

choose a prime  $P > 193$  work on  $\text{GF}(P)$   $P > 193 + 138$

$P(x)$  of degree 192

$P(0) = \text{secret}$   $P(1) \dots P(193)$  give to 193 member countries.

138  $P(194) - P(332)$  gives to Secretary-General

- (b) The General Assembly of the UN decides to add an extra level of security: each of the 193 member countries has a delegation of 12 representatives, all of whom must agree in order for that country to help open the vault. Propose a scheme that adds this new feature. The scheme should give private information to the Secretary-General and to each representative of each country.

for  $p(1)$  to  $p(193)$   
 $g(1)$  of degree  $\leq 11 \dots$  work on  $\text{GF}(p)$   
let  $p(n) = g(0) \rightarrow$  give  $g(1), g(2) \dots g(12)$  to  
(2 representatives)

## 4 To The Moon!

### Note 8

A secret number  $s$  is required to launch a rocket, and Alice distributed the values  $(1, p(1)), (2, p(2)), \dots, (n+1, p(n+1))$  of a degree  $n$  polynomial  $p$  to a group of  $\$GME$  holders  $\text{Bob}_1, \dots, \text{Bob}_{n+1}$ . As usual, she chose  $p$  such that  $p(0) = s$ .  $\text{Bob}_1$  through  $\text{Bob}_{n+1}$  now gather to jointly discover the secret. However,  $\text{Bob}_1$  is secretly a partner at Melvin Capital and already knows  $s$ , and wants to sabotage  $\text{Bob}_2, \dots, \text{Bob}_{n+1}$ , making them believe that the secret is in fact some  $s' \neq s$ . How could he achieve this? In other words, what value should he report (in terms of variables known in the problem, such as  $s'$ ,  $s$  or  $y_1$ ) in order to make the others believe that the secret is  $s'$ ?

$$p(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + s$$

$\times$

$$p(1) = a_n + a_{n-1} + \dots + a_1 + s$$

$\downarrow$

$$p(1) - s + s'$$

$$s = y_1 \Delta_1(0) + \sum_{k=2}^{n+1} y_k \Delta_k(0)$$

$$s' = y'_1 \Delta_1(0) + \sum_{k=2}^{n+1} y'_k \Delta_k(0)$$

$$s - s' = \Delta_1(0) (y_1 - y'_1)$$

$$y'_1 = (\Delta_1(0))^{-1} (s' - s) + y_1$$

$\nearrow$

$\Delta_1(x)$  Root  $\xrightarrow{2-n+1 \rightarrow n+1 \text{ root}}$  so  $\Delta_1(0) \neq 0$

$\downarrow$  degree,  $n$