Kernel 里面对抢占的关闭和打开,是依赖 preempt_count 这个变量的,而不是某个开关。 preempt count 是个 u64 变量,属于 struct thread info 结构体。

以 arch\arm64\include\asm\thread_info.h 为例:

```
* low level task data that entry.S needs immediate access to.
struct thread_info {
                                  /* low level flags */
#endif
    union
               preempt_count; /* 0 => preemptible, <0 => bug */
struct {
#ifdef CONFIG_CPU_BIG_ENDIAN
u32 need_resched;
           u32 count
#else
           u32 count;
u32 need_resched;
#endif
        ] preempt;
=0 可以被抢占, >0 不能被抢占。
```

读取这个值可以用函数:

```
static inline int preempt_count(void)
   return READ_ONCE(current_thread_info()->preempt.count);
```

不同的场景,抢占的值是不一样的。

```
reserved bits
                      bit21
                                  bit20 bit19-bit16 bit15-bit8
                                                                  bit7-bit0
                PREEMPT_ACTIVE NMI HARDIRQ
* PREEMPT_MASK:
* SOFTIRQ_MASK:
                     0x0000ff00
* HARDIRQ_MASK:
     NMI_MASK:
* PREEMPT_ACTIVE:
                  0x00200000
#define PREEMPT BITS
#define SOFTIRQ_BITS
#define HARDIRQ_BITS
#define NMI BITS
```

普通场景(PREEMPT_MASK)打开和关闭抢占:

arch\arm64\include\asm\preempt.h

```
\label{eq:count_dec_and_test())} $$ if (unlikely(preempt_count_dec_and_test())) \land \\ preempt_schedule(); \land \\ $$ while (0) $$
static _always_inline bool __preempt_count_dec_and_test(void)
       * Because of load-store architectures cannot do per-cpu atomic
* operations; we cannot use PREEMPT_NEED_RESCHED because it might get
      return ! --*preempt_count_ptr() && tif_need_resched();
static _always_inline volatile int *preempt_count_ptr(void)
      return &current_thread_info()->preempt_count;
```

可以看到此处只是 preempt_count++或者--软中断场景 (SOFTIRQ_MASK) 打开和关闭抢占:

Kernel\softirq.c

可以看到此处增减的幅度是 SOFTIRQ_DISABLE_OFFSET