

Qishi quiz 2

Group 2

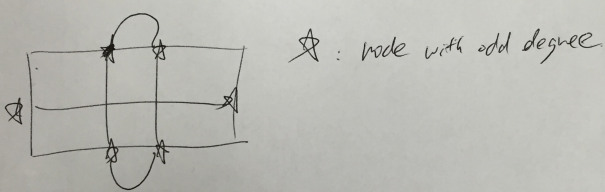
1. PROBLEM 1

According to wiki a χ^2 test is any statistical hypothesis test in which the sampling distribution of the test statistic is a χ^2 distribution when the null hypothesis is true. We can use it to test independence.

2. PROBLEM 2

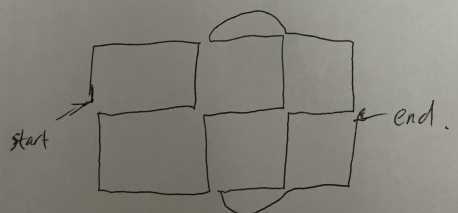
The shortest route to visit all edges has length 19 and it can be seen from the following:

Prob 2: There are six nodes in the graph whose degree is odd. So there is no Euler ~~circuit~~ circuit (i.e. no route visit ~~each~~ edges only once). But we ~~can~~ can connect 4 of them (and make them to two pair) and make the resulting ~~graph~~ graph as Euler ~~circuit~~ circuit.



☆ : node with odd degree

So the shortest route is the Euler circuit of the above graph.



The length of this route is 19.

3. PROBLEM 3

Let $Z = X + Y$. Since X, Y are iid $N(0, 1)$, Z is normal with mean 0 and variance 2. Moreover by symmetry we have $P(X|X + Y > 0) = P(Y|X + Y > 0)$. So $P(X|X + Y > 0) = \frac{1}{2}P(X + Y|X + Y > 0) = \frac{1}{2}P(Z|Z > 0) = \frac{1}{2} \cdot (\frac{1}{2\sqrt{\pi}} \int_0^\infty xe^{-\frac{x^2}{4}} dx) / P(Z > 0) = \frac{1}{2\sqrt{\pi}} \int_0^\infty xe^{-\frac{x^2}{4}} dx = \frac{1}{\sqrt{\pi}} \int_0^\infty e^{-\frac{x^2}{4}} d(\frac{x^2}{4}) = -\frac{1}{\sqrt{\pi}} e^{-y}|_0^\infty = \frac{1}{\sqrt{\pi}}$. So $P(X|X + Y > 0) = \frac{1}{\sqrt{\pi}}$.

4. PROBLEM 4

(1) If keep rolling for once, the payoff is 0 with probability $1/6$ and is $35 + i$ for $i = 2, \dots, 6$ with probability $1/6$ respectively. In this case the expected payoff is $\frac{1}{6} \cdot \sum_{i=2}^6 (35 + i) = 32.5$. Since by rolling once, the maximimul sum we can get is $35 + 6 = 41$ and $41 + 6 = 47 < 7^2 = 49$, we could still keep rolling without worrying about to get a square. After rolling twice the payoff become $32.5 + 3.5 = 36$ which is geater than 35. Hence we should keep rolling if the sum is 35.

(2)

5. PROBLEM 5

The two breaking points x and y are uniformly distributed on $[0, 1]$. Let c_1, c_2, c_3 be the length of each piece. Then $c_1 + c_2 + c_3 < 1$ and the area of volumn of it is $1/3! = 1/6$. So the joint distribution of them is $1/6$. Let $A = \min(X, Y)$, $B = \max(X, Y)$ and $C = \min(A, 1 - B, B - A)$. Let $F_C(a)$ be the cdf of C . Then $1 - F_C(a) = \text{Prob}(C \geq a) = P(a \leq X \leq 1, a \leq Y \leq 1 - a, |X - Y| \geq a) = 1 - (1 - 3a)^2$. So the pdf of C is $f_C(a) = F'_C(a) = 6(1 - 3a)$. So the expected length of the smallest piece is $\int_0^{1/3} 6a(1 - 3a) da = 1/9$.

Similarly we can compute the expected length of the longest piece. Let $D = \max(A, 1 - B, B - a)$. Then the cdf for D is $F_D(a) = \text{Prob}(D \leq a) = \text{Prob}(A \leq a, 1 - B \leq a, B - A \leq a)$. So $F_D(a) = (3a - 1)^2$ if $1/3 \leq a \leq 1/2$ and $F_D(a) = 1 - 3(1 - a)^2$ if $1/2 \leq a \leq 1$. So the pdf of D is $f_D(a) = 6(3a - 1)$ if $1/3 \leq a \leq 1/2$ and $f_D(a) = 6(1 - a)$ if $1/2 \leq a \leq 1$. So the expected length of the longest piece is $\int_{1/3}^{1/2} 6a(3a - 1) da + \int_{1/2}^1 6a(1 - a) da = 11/18$.

Hence, the expected length of the middle-sized piece is $1 - 1/9 - 11/18 = 5/18$.

6. PROBLEM 6

(1) For $i = 1, \dots, N$, define random variable $I_i = 1$ if person i choose his/her own hat and $I_i = 0$ otherwise. Then $E(I_i) = \text{Prob}(I_i = 1) = 1/N$. Then $Y = \sum_i I_i$. So $E(Y) = E(\sum_i I_i) = \sum_i E(I_i) = N \cdot 1/N = 1$.

(2) We have $E(Y^2) = E((\sum_i I_i) \cdot (\sum_j I_j)) = E(\sum_{i,j=1}^N I_i I_j) = \sum_{i,j=1}^N E(I_i I_j) = \sum_{i,j=1, i \neq j}^N \text{Prob}(I_i = 1, I_j = 1) + \sum_i \text{Prob}(I_i = 1) = \sum_{i,j=1, i \neq j}^N \frac{1}{N(N-1)} + \sum_i \text{Prob}(I_i = 1) \frac{1}{N} = 1 + 1 = 2$. So the variance of Y is $\text{var}(Y) = E(Y^2) - (E(Y))^2 = 2 - 1^2 = 1$.

(3) Let X_N be the random variable counting the nuber of people picking their own hats in the first round. Since from (1) we know that on average there are one person who picks his/her own hat, we might guess that be the number of rounds that are run are $R(N) = N$. In fact, it is easy to see that it is true for $N = 1$. For general N , we could prove it by

induction. In fact, suppose that it holds up to $N - 1$. Then

$$\begin{aligned}
E(R(N)) &= \sum_{i=0}^N E(R(N)|X_N = i) \text{Prob}(X_N = i) \\
&= \sum_{i=0}^N (1 + E(R(N - i)) \text{Prob}(X_N = i)) \\
&= \sum_{i=0}^N \text{Prob}(X_N = i) + E(R(N)) \text{Prob}(X_N = 0) + \sum_{i=1}^N (N - i) \text{Prob}(X_N = i) \\
&= 1 + E(R(N)) \text{Prob}(X_N = 0) + N(1 - \text{Prob}(X_N = 0)) - E(X_N) \\
&= E(R(N)) \text{Prob}(X_N = 0) + N(1 - \text{Prob}(X_N = 0)),
\end{aligned}$$

where the third equality follows from induction hypothesis and the last equality follows from (1) that $E(X_N) = 1$. So $E(R(N))(1 - \text{Prob}(X_N = 0)) = N(1 - \text{Prob}(X_N = 0))$ which implies $E(R(N)) = N$ since $P(X_N = 0) \neq 1$. This finishes the induction process.

(4) Conditioning on X_N gives

$$\begin{aligned}
E(S(N)) &= \sum_{i=0}^N E(S(N)|X_N = i) \text{Prob}(X_N = i) \\
&= \sum_{i=0}^N (N + E(S(N - i)) \text{Prob}(X_N = i)) \\
&= N + \sum_{i=0}^N E(S(N - i) \text{Prob}(X_N = i)).
\end{aligned}$$

This implies

$$E(S(N)) = N + E(S(N - X_N)).$$

Because of this relation, it is natural to assume that $E(S(N))$ is a polynomial function of N . But if there is exactly only one match in each round, we would have $\sum_{i=1}^N i = N(N + 1)/2$ selections. So it would be natural to assume $E(S(N))$ is a quadratic polynomial function of N . So suppose that $E(S(N)) = aN^2 + bN$ and plug this into the above equation. We get $aN^2 + bN = N + E(a(N - X_N)^2 + b(N - X_N))$. Since $E(X_N) = \text{var}(X_N) = 1$, we have $aN^2 + bN = aN^2 + (b - 2a + 1)N + 2a - b$. So $a = 1/2$ and $b = 1$. This implies $E(S(N)) = N + N^2/2$. To prove it rigorously, we could use induction. In the case of $N = 2$, it is true because the number of rounds is geometric distribution with parameter $1/2$. Suppose that it is true up to $N - 1$. Then

$$\begin{aligned}
E(S(N)) &= N + E(S(N)) \text{Prob}(X_N = 0) + \sum_{i=1}^N E(S(N - i) \text{Prob}(X_N = i)) \\
&= N + E(S(N)) \text{Prob}(X_N = 0) + \sum_{i=1}^N ((N - i)^2/2 + (N - i)) \text{Prob}(X_N = i) \\
&= N + E(S(N)) \text{Prob}(X_N = 0) + (N^2/2 + N)(1 - \text{Prob}(X_N = 0)) - (N + 1)E(X_N) \\
&\quad + E(X_N^2)/2
\end{aligned}$$

where the second equality follows from induction hypothesis. Plugging $E(X_N) = 1$, $E(X_N^2) = 2$ into the above equation gives $E(S(N)) = N + N^2/2$, as desired.

(5) Let H_i be the random variable counting the number of hats chosen by person i . Then $\sum_i H_i = S(N)$. Hence $E(\sum_i H_i) = E(S(N))$. By symmetry, H_i has the same mean, so $E(H_i) = S(N)/N = 1 + n/2$. So the expected number of false selections made by one of the N people is $E(H_i - 1) = n/2$.

7. PROBLEM 7

The upper bound is 1 and the lower bound is 0.02. Let ι be a column vector every element of which is 1. Let x be a $n \times k$ matrix and let $S(x)$ be the subspace generated by the column vectors of x . Then $P_x = x(x^T x)^{-1}x^T$ is the projection matrix projecting onto $S(x)$ and $M_x = I - P_x = I - x(x^T x)^{-1}x^T$ is the annihilator space projecting onto the orthogonal complement subspace $S^\perp(X)$. Then the (centered) R^2 is

$$R^2 = \frac{\|P_x M_\iota y\|^2}{\|M_\iota y\|^2},$$

where $M_\iota y = (I - P_\iota)y = y - \iota(\iota^T \iota)^{-1}\iota^T y = y - \bar{y}\iota$ is the centered version of y and $\bar{y} = \frac{1}{n} \sum_{t=1}^n y_t$ is the mean of y . Here we assume that ι is contained in the span $S(x)$ of the regressors. So in this problem $S(x)$ is $S(x_1, \iota)$ or $S(x_2, \iota)$ or $S(x_1, x_2, \iota)$.

Upper bound: If $M_\iota y \in S(x_1, x_2)$, the vector space generated by the column vectors of x_1 and x_2 , then $\|P_{(x_1, x_2)} M_\iota y\|^2 = \|M_\iota y\|^2 \Rightarrow R^2 = 1$. Since $R^2 \leq 1$, the upper bound of R^2 is 1. Geometrically, this means that $M_\iota y$ belongs to the subspace span by (the column vectors of) x_1 and x_2 , so the model is perfectly fit and $R^2 = 1$.

Lower bound: First we show that $R_{M_3}^2 \geq R_{M_1}^2, R_{M_2}^2$, where we use subscripts to differentiate R^2 from different models. Without loss of generality, we just need to prove the case for M_2 . In fact, since both $P_{x_1, x_2, \iota}$ and $P_{x_2, \iota}$ are symmetric, so is $P_{x_1, x_2, \iota} - P_{x_2, \iota}$. Moreover, since $S(x_2, \iota) \subset S(x_1, x_2, \iota)$, we have $P_{x_1, x_2, \iota} P_{x_2, \iota} = P_{x_2, \iota} P_{x_1, x_2, \iota} = P_{x_2, \iota}$. So $(P_{x_1, x_2, \iota} - P_{x_2, \iota})^2 = P_{x_1, x_2, \iota}^2 - P_{x_1, x_2, \iota} P_{x_2, \iota} - P_{x_2, \iota} P_{x_1, x_2, \iota} + P_{x_2, \iota}^2 = P_{x_1, x_2, \iota} - P_{x_2, \iota}$ as both of $P_{x_1, x_2, \iota}$ and $P_{x_2, \iota}$ are idempotents. So $P_{x_1, x_2, \iota} - P_{x_2, \iota}$ is an orthogonal projection matrix. Hence $\|P_{x_1, x_2, \iota} M_\iota y\|^2 - \|P_{x_2, \iota} M_\iota y\|^2 = (M_\iota y)^T (P_{x_1, x_2, \iota} - P_{x_2, \iota}) M_\iota y = \|(P_{x_1, x_2, \iota} - P_{x_2, \iota}) M_\iota y\|^2 \geq 0 \Rightarrow \frac{\|P_{x_1, x_2, \iota} M_\iota y\|^2}{\|M_\iota y\|^2} \geq \frac{\|P_{x_2, \iota} M_\iota y\|^2}{\|M_\iota y\|^2}$. So $R_{M_3}^2 \geq R_{M_2}^2 = 0.02$. But this equality can be attained. In fact, if $\|P_{x_1, x_2, \iota} M_\iota y\|^2 = \|P_{x_2, \iota} M_\iota y\|^2$, then $R_{M_3}^2 = 0.02$. Geometrically, to do the OLS for model M_3 , it amounts to find a vector v inside $S(x_1, x_2, \iota)$ such that $\|M_\iota y - v\|^2$ is minimized. In our case that $R_{M_3}^2 = 0.02$ this minimization is exactly achieved by $P_{x_2, \iota} M_\iota y$.

So the range of R^2 of Model 3 is $[0.02, 1]$.

8. PROBLEM 8

$1/n$ is a rational number so it can be written as binary number with finite sum: $1/n = \sum_{i=1}^m p_i 2^{-i}$ for $p_i = 0, 1$. Then the function `biasedCoin()` can be implement (by using the function `fairCoin()`) in the following ways:

```
int biasedCoin(int [] p, int m) {
    for (int i = 0; i < m; i++) {
        int s = fairCoin();
```

```

        if (s < p[i]) {
            return 1;
        }
        else if (s > p[i]) {
            return 0;
        }
    }
    return 0;
}

```

The algorithm is explained as follows: when we coss a fair coin, the heads count as 1 and the tails count as 0. Let $s_i \in \{0, 1\}$ be the result of i -th toss starting with $i = 1$. After each toss, we compare s_i with p_i . If $s_i < p_i$ (resp. $s_i > p_i$), the `baisedCoin()` return a head (resp. tail) and the tossing stops. If $s_i = p_i$, we continue tossing the fair coin. If we end up with $s_n = p_n$, `baisedCoin()` return a tail. This will give a function which return heads with probabily $1/n$.

9. PROBLEM 9

Let A_i be the event that the i -th bridge is weak and the previus $i - 1$ bridges are all strong (i.e i is the first strong bridge), then A_i are disjoint and $P(\cup_{i=1}^{11} A_i) = 1$, where A_{11} means that all the bridges are strong. Moreover, we have $Prob(A_i) = \frac{1}{2^i}$ for $i \leq 10$ and $Prob(A_{11}) = \frac{1}{2^{10}}$. Define random variables I_i as $I_i = 1$ on A_i and $I_i = 0$, otherwise. Let X be the random variable counting how many brdges the man has to cross. Then $X|_{A_i} = i - 1 + 10 = i + 9$ for $i \leq 10$ and $X|_{A_{11}} = 10$. So $E(X) = \sum_{i=1}^{10} (i + 9) \cdot \frac{1}{2^i} + 10 \cdot \frac{1}{2^{10}}$. But $\sum_{i=1}^{10} \frac{i}{2^i} = 2 \cdot \sum_{i=1}^{10} \frac{i}{2^i} - \sum_{i=1}^{10} \frac{i}{2^i} = 1 + \sum_{i=1}^9 \frac{i+1}{2^i} - \sum_{i=1}^{10} \frac{i}{2^i} = 1 + \sum_{i=1}^9 \frac{1}{2^i} - \frac{10}{2^{10}}$. So $E(X) = 1 + \sum_{i=1}^9 \frac{1}{2^i} - \frac{10}{2^{10}} + 9 \cdot \sum_{i=1}^{10} \frac{1}{2^i} + \frac{10}{2^{10}} = 1 + \frac{9}{2^{10}} + 10 \cdot \sum_{i=1}^9 \frac{1}{2^i} = 1 + \frac{9}{2^{10}} + 10 \cdot (\frac{\frac{1}{2} - \frac{1}{2^{10}}}{1 - \frac{1}{2}}) = 1 + \frac{9}{2^{10}} + 10 - \frac{10}{2^9} = 11 - \frac{11}{2^{10}}$. So to arrive the 10th island, on average the man has to cross $11 - \frac{11}{2^{10}}$ bridges.

10. PROBLEM 10

It is a reference to a constant pointer to a constant integer.

11. PROBLEM 11

```

/**
 * Definition for singly-linked list.
 * struct ListNode {
 *     int val;
 *     ListNode *next;
 *     ListNode(int x) : val(x), next(NULL) {}
 * };

```

```

    */
class Solution {
public:
    void deleteNode(ListNode* node) {
        node->val = node->next->val;
        node->next = node->next->next;
    }
};

```

12. PROBLEM 12

The .hh file for the matrix class is as follows:

```

class Matrix {

private:
    int numRows;
    int numCols;
    // matriarray to store the data in Matrix
    int *matrixarray;

public:
    // Constructors
    Matrix(); // default constructor
    Matrix(int sizeRow, int sizeCol);
    Matrix(Matrix &a); // copy constructor

    // Destructor
    ~Matrix();

    // Mutator methods
    void setelem(int row, int col, int elem);
    void add (Matrix &a);
    void subtract (Matrix &a);
    bool equals (Matrix &a);

    // Accessor methods
    int getrows();
    int getcols();
    int getelem(int row, int col);
};

```

The .cpp file for the matrix class is as follows:

```
#include "Matrix.hh"
#include <cassert>
using namespace std;

// The default constructor creates a 0 by 0 matrix.
Matrix::Matrix() {
    numRows = 0;
    numCols = 0;
    matrixarray = new int[0];
}

// Copy constructor
Matrix::Matrix(Matrix &a) {
    numRows = a.getrows();
    numCols = a.getcols();
    matrixarray = new int[numRows * numCols];
    for (int i = 0; i < numRows; i++) {
        for (int j = 0; j < numCols; j++) {
            matrixarray[i * numCols + j] =
                a.matrixarray[i * numCols + j];
        }
    }
}

// Initialize matrix of size sizeRow by sizeCol
Matrix::Matrix(int sizeRow, int sizeCol) {
    matrixarray = new int[sizeRow * sizeCol];
    numRows = sizeRow;
    numCols = sizeCol;
    for (int i = 0; i < sizeRow; i++) {
        for (int j = 0; j < sizeCol; j++) {
            matrixarray[i * sizeCol + j] = 0;
        }
    }
}
```

```

// Get the rows tall of the matrix
int Matrix::getrows () {
    return numRows;
}

// Get the columns wide of the matrix
int Matrix::getcols () {
    return numCols;
}

// Set the element at row-th row, col-th column as elem
void Matrix::setelem (int row, int col, int elem) {
    assert ((row >= 0) && (row < numRows));
    assert ((col >= 0) && (col < numCols));
    matrixarray[row * numCols + col] = elem;
}

// Get the element at row-th row, col-th column
int Matrix::getelem (int row, int col) {
    assert ((row >= 0) && (row < numRows));
    assert ((col >= 0) && (col < numCols));
    return matrixarray[row * numCols + col];
}

// Add two matrices
void Matrix::add (Matrix &a) {
    assert(numRows == a.getrows());
    assert(numCols == a.getcols());
    for (int i = 0; i < numRows; i++) {
        for (int j = 0; j < numCols; j++) {
            matrixarray[i * numCols + j] =
                matrixarray[i * numCols + j]
                + a.matrixarray[i * numCols + j];
        }
    }
}

```



```

// Subtract a from another matrix
void Matrix::subtract (Matrix &a) {
    assert(numRows == a.getrows());
    assert(numCols == a.getcols());
    for (int i = 0; i < numRows; i++) {
        for (int j = 0; j < numCols; j++) {
            matrixarray[i * numCols + j] =
                matrixarray[i * numCols + j]
                - a.matrixarray[i * numCols + j];
        }
    }
}

// Test for equality
bool Matrix::equals (Matrix &a) {
    if (numRows != a.getrows() || numCols != a.getcols()){
        return false;
    }
    for (int i = 0; i < numRows; i++) {
        for (int j = 0; j < numCols; j++) {
            if ((matrixarray[i * numCols + j])
                != (a.matrixarray[i * numCols + j])) {
                return false;
            }
        }
    }
    return true;
}

// Release the memory
Matrix::~Matrix() {
    delete [] matrixarray;
}

```

13. PROBLEM 13

No, constructor of a class CANNOT be virtual. Here is a quote from Bjarne Stroustrup's C++ Style and Technique FAQ: "*virtual call is a mechanism to get work done given partial*

information. In particular, "virtual" allows us to call a function knowing only any interfaces and not the exact type of the object. To create an object you need complete information. In particular, you need to know the exact type of what you want to create. Consequently, a "call to a constructor" cannot be virtual."

14. PROBLEM 14

Yes, but the non-virtual function was not redefined in any derived class.

15. PROBLEM 15

The algorithm find palindromes by expanding from the centers i and $(i, i + 1)$. The time complexity is $O(N^2)$.

```
def expandAroundCenter(str , c1 , c2):
    left = c1
    right = c2
    length = len(str)

    while(left >= 0 && right <=length-1 && str[left] == str[right]):
        left --
        right++

    return str[left+1:right]

def longestPalindrome(str):
    length = len(str)
    if (length==0): return ""

    longest = str[0:1]
    for i in range(length):
        p1 = expandAroundCenter(str , i , i)
        if (len(p1) > len(longest)): longest = p1

        p2 = expandAroundCenter(str , i , i+1)
        if (len(p2) > len(longest)): longest = p2

    return longest
```

16. PROBLEM 16

We can first reverse the sentence and then reverse each words. The time complexity is $O(N^2)$.

```

def reverseWords(word):
    length = len(word)
    if (length < 2): return word
    left = 0
    right = length-1

    while (left < right):
        tmp = word[left]
        word[right] = word[left]
        word[left] = tmp

    return word

def reverseSentences(sentence):
    res = []
    reverseWords(sentence)

    pt1 = 0
    pt2 = 0

    while (pt2 != len(sentence)):
        if (sentence[pt2] != " "):
            pt2++
        sentence[pt1,pt2-1] = reverseWords(sentence[pt1,pt2-1])
        pt1 = pt2

    return sentence

```

17. PROBLEM 17

Use dynamic programming. For each i , find the max profits $\text{profit1}[i]$ and $\text{profits2}[i]$ with one transection from $\text{prices}[0]$ to $\text{prices}[i]$ and from $\text{prices}[i+1]$ to $\text{prices}[n-1]$ respectively. Find the maximal $\text{profit1}[i] + \text{profit2}[i]$ from 0 to $n-1$. The time complexity is $O(N)$.

```

def maxProfit(self, prices):
    length=len(prices)
    if length < 2: return 0
    profit1=[0 for i in range(length)]
    profit2=[0 for i in range(length)]

```

```

# find the maximal profits with one transaction
# in sub list prices[0],...,price[i]
minP=prices[0]
for i in range(1,length):
    minP=min(minP, prices[i])
    profit1[i]=max(profit1[i-1],prices[i]-minP)

# find the maximal profits with one transaction
# in sub list prices[i+1],...,price[n-1]
maxP=prices[length-1]
for i in range(length-2,1,-1):
    maxP=max(maxP, prices[i])
    profit2[i]=max(profit2[i+1],maxP-prices[i])

res=0
for i in range(length):
    if profit1[i]+profit2[i]>res: res=profit1[i]+profit2[i]
return res

```

18. PROBLEM 18

Form a graph such that each node represent each person and if two persons are friends, then there is an edge connecting the corresponding two nodes. Then this problem actually asks to find a Hamiltonian Path in athis graph, that is a path that visits each vertex exactly once. We could use back-tracking algorithm to find such path. The input is a N by N matrix `graph[N][N]` which is the adjacency matrix representation of the graph. A value `graph[i][j]` is 1 if there is a direct edge from i to j (i.e. person i and person j are friends), otherwise `graph[i][j]` is 0.

```

#function to check if the vertex v can be
# added at index 'pos'
def isSafe(v,graph,path,pos):
    #Check if this vertex is an adjacent vertex of the
    # previously added vertex
    if (graph[path[pos-1]][v] == 0):
        return False

# Check if the vertex has already been included.
for i in range(pos):
    if (path[i]==v):
        return False

```

```

    return True

# A recursive function to solve hamiltonian cycle
# problem
def hamPath(graph, path, pos):
    if (pos == N):
# if there is an edge from the last included vertex
# to the first node
        if (graph[path[pos-1]][path[0]] == 1):
            return True
        else:
            return False
# Try different vertices as a next candidate in
# Hamiltonian Cycle.
    for v in range(1,V):
        if (isSafe(v,graph,path,pos)):
            path[pos] = v
            if (hamPath(graph,path,pos+1) == True):
                return True

            path[pos] = -1

    return False

#This function solves the Hamiltonian Cycle problem
# using Backtracking.

def BookProb(graph, N)
    path = N*[-1]

    path[0] = 0
    if (hamPath(graph,path,1) == False):
        print "No solution"
        return False

    for i in range(N):

```

```
        print(path[i])  
        print(" ")  
  
    return True
```