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ESCUELA DE INGENIERÍA
DEPARTAMENTO DE CIENCIAS DE LA COMPUTACIÓN

IIC2343 – Arquitectura de Computadores

Ayudantía 2

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Preguntas

1. Assembly

- a) Cargue los valores 5 y 17 en los registros \$s1 y \$s2. Considere registros de 1 byte.
- b) Utilizando la arquitectura ISA MEEP MEEP, programe una multiplicación en *assembly* entre dos números guardados en memoria. Considere que los números a multiplicar están en las direcciones 8 y 9, mientras que el resultado desea guardarse en la dirección 13.
- c) Utilizando la arquitectura ISA MEEP MEEP, programe en *assembly* un código que le permita sumar N números desde el número P_0 en adelante. Asuma que N y P_0 están guardados en la memoria RAM en las direcciones 13 y 21, respectivamente. Guarde el valor final en la dirección 4 de la memoria. Comente si lo encuentra necesario.

2. Anexos

MIPS assembly language				
Category	Instruction	Example	Meaning	Comments
Arithmetic	add	add \$s1,\$s2,\$s3	$\$s1 = \$s2 + \$s3$	Three register operands
	subtract	sub \$s1,\$s2,\$s3	$\$s1 = \$s2 - \$s3$	Three register operands
	add immediate	addi \$s1,\$s2,20	$\$s1 = \$s2 + 20$	Used to add constants
Data transfer	load word	lw \$s1,20(\$s2)	$\$s1 = \text{Memory}[\$s2 + 20]$	Word from memory to register
	store word	sw \$s1,20(\$s2)	$\text{Memory}[\$s2 + 20] = \$s1$	Word from register to memory
	load half	lh \$s1,20(\$s2)	$\$s1 = \text{Memory}[\$s2 + 20]$	Halfword memory to register
	load half unsigned	lhu \$s1,20(\$s2)	$\$s1 = \text{Memory}[\$s2 + 20]$	Halfword memory to register
	store half	sh \$s1,20(\$s2)	$\text{Memory}[\$s2 + 20] = \$s1$	Halfword register to memory
	load byte	lb \$s1,20(\$s2)	$\$s1 = \text{Memory}[\$s2 + 20]$	Byte from memory to register
	load byte unsigned	lbu \$s1,20(\$s2)	$\$s1 = \text{Memory}[\$s2 + 20]$	Byte from memory to register
	store byte	sb \$s1,20(\$s2)	$\text{Memory}[\$s2 + 20] = \$s1$	Byte from register to memory
	load linked word	ll \$s1,20(\$s2)	$\$s1 = \text{Memory}[\$s2 + 20]$	Load word as 1st half of atomic swap
	store condition. word	sc \$s1,20(\$s2)	$\text{Memory}[\$s2 + 20] = \$s1; \$s1 = 0 \text{ or } 1$	Store word as 2nd half of atomic swap
Logical	load upper immed.	lui \$s1,20	$\$s1 = 20 * 2^{16}$	Loads constant in upper 16 bits
	and	and \$s1,\$s2,\$s3	$\$s1 = \$s2 \& \$s3$	Three reg. operands; bit-by-bit AND
	or	or \$s1,\$s2,\$s3	$\$s1 = \$s2 \$s3$	Three reg. operands; bit-by-bit OR
	nor	nor \$s1,\$s2,\$s3	$\$s1 = \sim (\$s2 \$s3)$	Three reg. operands; bit-by-bit NOR
	and immediate	andi \$s1,\$s2,20	$\$s1 = \$s2 \& 20$	Bit-by-bit AND reg with constant
	or immediate	ori \$s1,\$s2,20	$\$s1 = \$s2 20$	Bit-by-bit OR reg with constant
	shift left logical	sll \$s1,\$s2,10	$\$s1 = \$s2 \ll 10$	Shift left by constant
Conditional branch	shift right logical	srl \$s1,\$s2,10	$\$s1 = \$s2 \gg 10$	Shift right by constant
	branch on equal	beq \$s1,\$s2,25	if ($\$s1 == \$s2$) go to PC + 4 + 100	Equal test; PC-relative branch
	branch on not equal	bne \$s1,\$s2,25	if ($\$s1 \neq \$s2$) go to PC + 4 + 100	Not equal test; PC-relative
	set on less than	slt \$s1,\$s2,\$s3	if ($\$s2 < \$s3$) $\$s1 = 1$; else $\$s1 = 0$	Compare less than; for beq, bne
	set on less than unsigned	sltu \$s1,\$s2,\$s3	if ($\$s2 < \$s3$) $\$s1 = 1$; else $\$s1 = 0$	Compare less than unsigned
	set less than immediate	slti \$s1,\$s2,20	if ($\$s2 < 20$) $\$s1 = 1$; else $\$s1 = 0$	Compare less than constant
	set less than immediate unsigned	sltiu \$s1,\$s2,20	if ($\$s2 < 20$) $\$s1 = 1$; else $\$s1 = 0$	Compare less than constant unsigned
Unconditional jump	jump	j 2500	go to 10000	Jump to target address
	jump register	jr \$ra	go to \$ra	For switch, procedure return
	jump and link	jal 2500	$\$ra = PC + 4$; go to 10000	For procedure call

Figura 1: ISA de un computador MIPS

SAMPLE INSTRUCTION	MEANING
ADD \$S1, \$S2, \$S3	$\$S1 = \$S2 + \$S3$
SUB \$S1, \$S2, \$S3	$\$S1 = \$S2 - \$S3$
ADDi \$S1, \$S2, Lit	$\$S1 = \$S2 + Lit$
AND \$S1, \$S2, \$S3	$\$S1 = \$S2 \wedge \$S3$
OR \$S1, \$S2, \$S3	$\$S1 = \$S2 \vee \$S3$
ANDi \$S1, \$S2, Lit	$\$S1 = \$S2 \wedge Lit$
ORi \$S1, \$S2, Lit	$\$S1 = \$S2 \vee Lit$
SL \$S1, \$S2, Lit	$\$S1 = \$S2 \ll Lit$
SRL \$S1, \$S2, Lit	$\$S1 = \$S2 \gg Lit$
BEQ \$S1, \$S2, Lit	if ($\$S1 == \$S2$) go to PC + Lit
BNE \$S1, \$S2, Lit	if ($\$S1 != \$S2$) go to PC + Lit
J Lit	go to Lit
LB \$S1, Lit(\$S2)	$\$S1 = Memory[\$S2 + Lit]$
SB \$S1, Lit(\$S2)	$Memory[\$S2 + Lit] = \$S1$
MOV \$S1, \$S2	$\$S1 = \$S2$
MOVi \$S1, Lit	$\$S1 = Lit$

* REGISTROS DE 1 BYTE, MEMORIA BYTE-INDEXED

Figura 2: ISA de un computador MEEP MEEP