#### COP 5536 Advanced Data Structures

# University of Florida

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## Exam 1 Solution

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## Question 1

a) Implement a QUEUE with two STACKs having constant amortized cost for each QUEUE operation (6 points).

Name the two STACKs as Stack<sub>1</sub> and Stack<sub>2</sub>, we can implement the QUEUE as follows:

- ENQUEUE(x): PUSH x into Stack<sub>1</sub>
- DEQUEUE(x): If Stack<sub>2</sub> is not empty, then simply POP from Stack<sub>2</sub> and return the element. If Stack<sub>2</sub> is empty, POP all the elements of Stack<sub>1</sub>, PUSH them into Stack<sub>2</sub>, then POP from Stack<sub>2</sub> and return the result.
- b) Choose any two from the three methods to prove the amortized cost for each QUEUE operation is O(1) (4 points each).

### Aggregate method

Consider a sequence of n operations. The sequence of operations will involve at most n elements. The cost associated with each element will be at most 4 i.e. (pushed into  $Stack_1$ , popped from  $Stack_1$ , pushed to  $Stack_2$ , and popped from  $Stack_2$ ). Hence, the actual cost of n operations will be upper bounded by T(n) = 4 n. Hence, the amortized cost of each operation can be T(n)/n = 4 n / n = 4 = O(1).

# Accounting method

We guess that the amortized costs for ENQUEUE and DEQUEUE are 3 and 1. We show that the potential function P(n) satisfies P(n) - P(0) >= 0 for all n.

We have P(0) = 0. If an element is not popped, then it's only pushed twice and popped once. Thus, the cost of 3 is paid for by ENQUEUE operation. The cost for last pop operation is paid for by the DEQUEUE.

Note: Alternatively, we can set the costs for ENQUEUE and DEQUEUE as 4 and 0 respectively.