

The Graph Algorithms Course

Pre-Requisite

- Basic Programming (we will code in C++)
- Basic - Intermediate Recursion
- Arrays, HM, LL, Trees

Outcome Expectation

- Course will be starting from very basics
but we will cover adv cf topics
- Beneficial with Internen perspective too.

Topics

(21 classes)

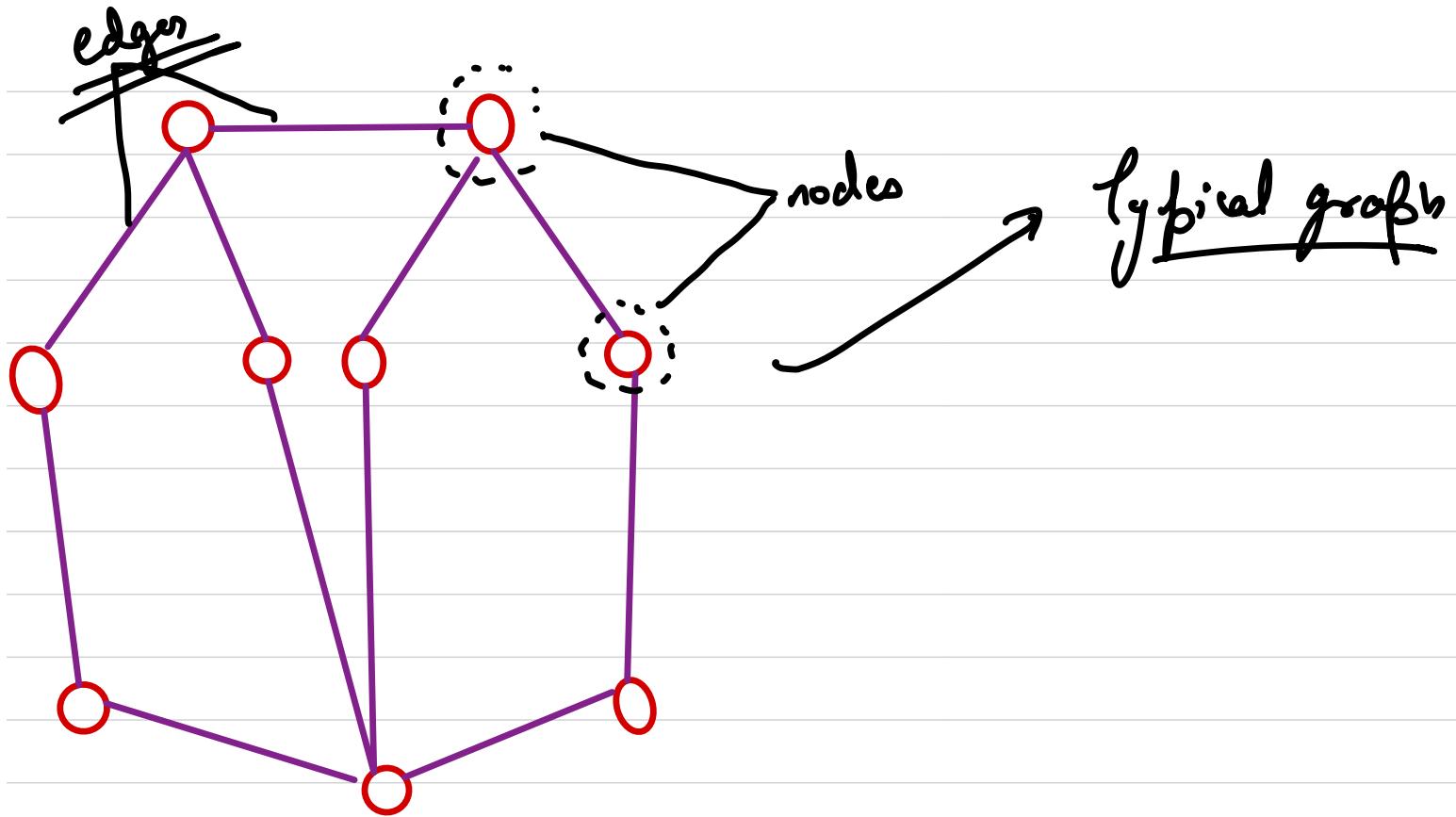
- Basic graphs & terminologies
- Traversal techniques Bfs / Dfs
- DP on graph (DP on Adj)
- Matrix Bfs / bfs , 0-1 Bfs dc
- Trees (n-ary) , DP on trees
- Shortest path, MST
- DSU
- SCC , Bridges / articulation point

Graphs



What is a graph ??

A graph is a collection of nodes and edges where each node might point to / connect to other nodes. The nodes represent real life entities and are connected by edges representing relationship between the nodes.



Applications of graph

Biology → PPI graph
→ Metabolic network graph

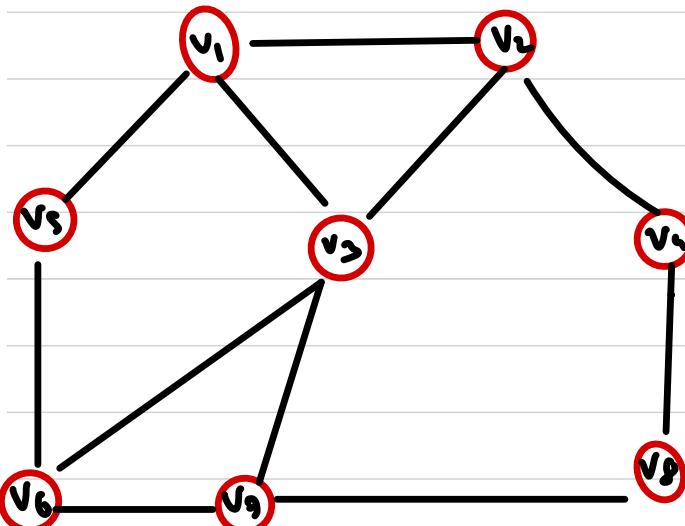
Electrical → Circuit net graph

Computer
Science → Shortest path (map)
→ gsm frequency assignment
→ routing mechanism
→ social media

formal definition of graph

set of vertices
 $V = \{v_1, v_2, v_3, \dots, v_n\}$

~~Edge set~~



$E = \{ \{v_1, v_2\}, \{v_2, v_3\}, \dots \}$ → unordered set

$$G = (V, E)$$

→ mathematical notation of graph

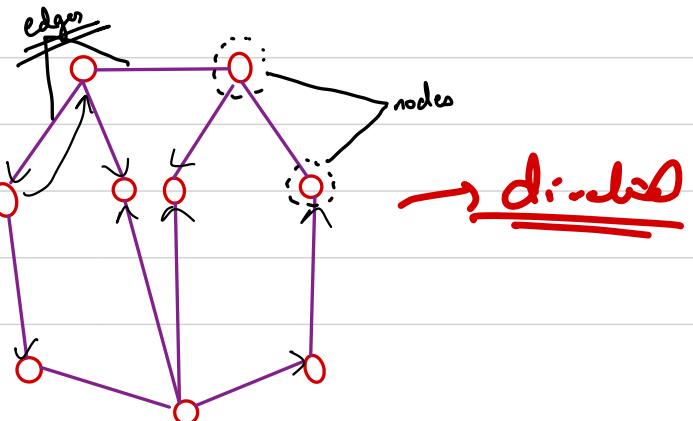
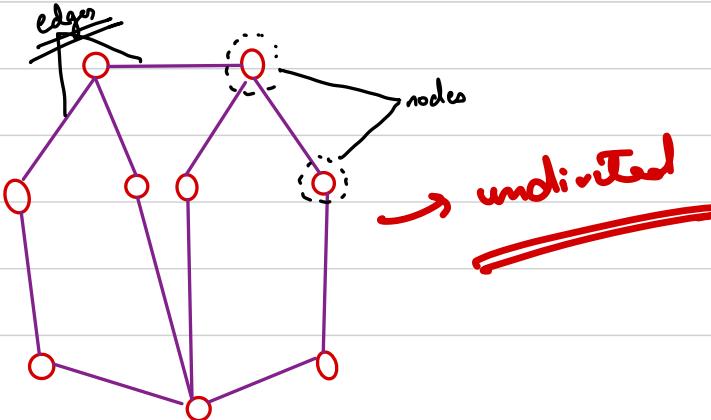
/unordered

A graph G is an ordered pair of a set V vertices & E , edges.

Types of graph

Based On direction

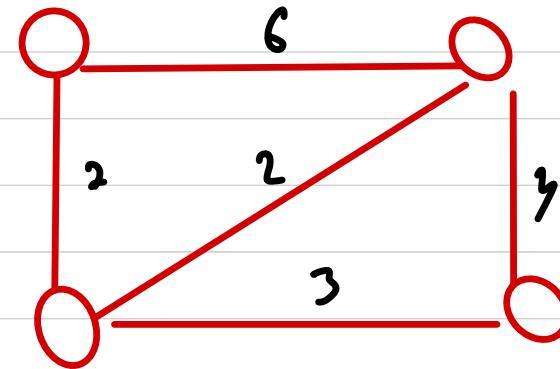
- **Directed** → water flow , road network,
- **Undirected** → facebook



Based on edge property

① Weighted

② Unweighted



Based on edge density

1) Sparse

2) Dense

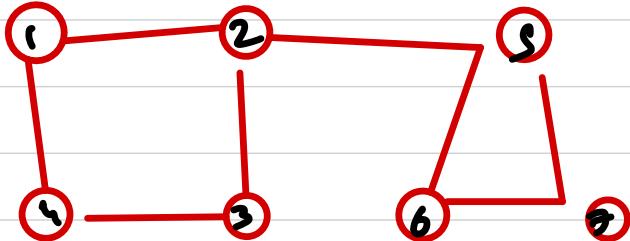
Representing graphs as a data structure

1) Adjacency Matrix

2) Adjacency List → adjacency set/map
3) Incidence Matrix ↗ edge list

adjacency matrix

$A_{ij} = \begin{cases} 1 & \rightarrow \text{if there is an edge} \\ & \text{from } i \text{ to } j \\ 0 & \rightarrow \text{else} \end{cases}$



	1	2	3	4	5	6	7
1	0	1	0	1	0	0	0
2	1	0	1	0	1	0	0
3	0	1	0	1	0	0	0
4	1	0	1	0	1	0	0
5	0	1	0	1	0	1	1
6	0	0	0	0	1	0	1
7	0	0	0	0	1	1	0

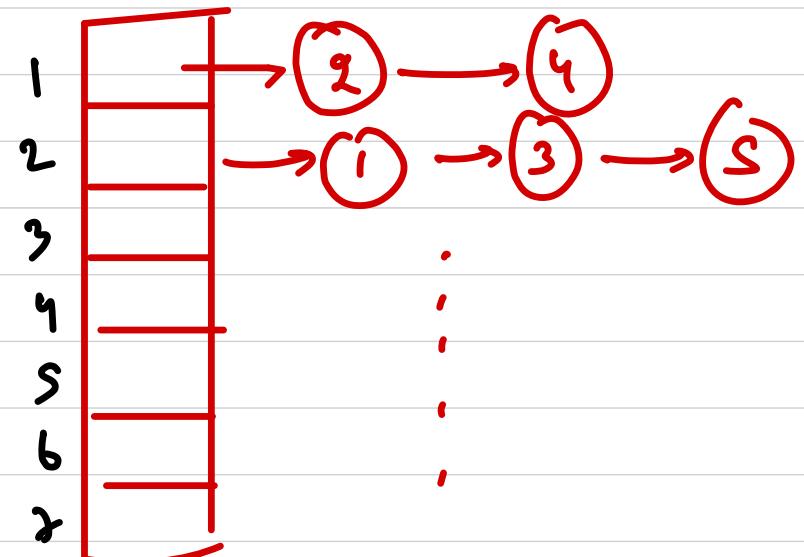
$\longrightarrow \underline{\underline{O(V^2)}}$

Nahui

adjacency list

$\{v, w\}$

Array of LL



Bucket array

$O(V+E)$ ~ space

adjacency map / set

array of hashmp

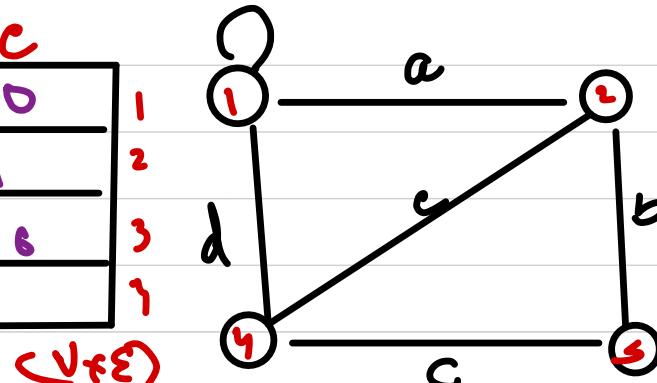
Incidence

Matrix

S_{min}

	a	b	c	d	e
1	1	0	0	1	0
2	1	1	0	0	1
3	0	1	1	0	0
4	0	0	1	1	1

$M =$



$(V+E)$

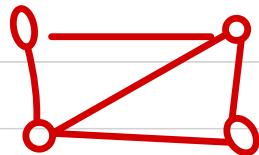
cols → edges
rows → vertices

$M_{ij} = \begin{cases} 1 & \text{if the } i^{\text{th}} \text{ vertex belongs to the } \\ & j^{\text{th}} \text{ edge} \\ 0 & \text{else} \end{cases}$

columns → $\underline{\underline{a}}$
rows → $\underline{\underline{dyne}}$

Q What is a degree ??

Degree of a vertex in a graph G , is the total no. of edges incident / associated with it.

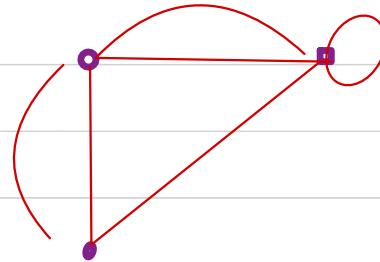


directed \rightarrow indegree / outdegree
graphs

outdegree \rightarrow no. of outgoing edges
indegree \rightarrow no. of incoming edges

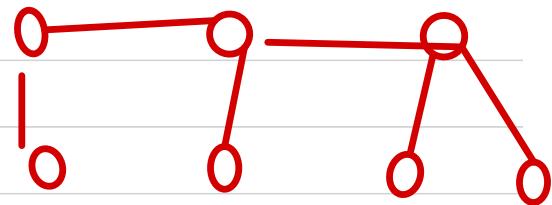
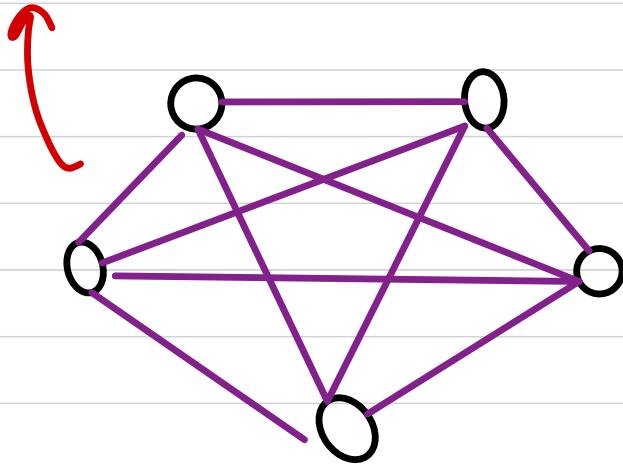
Multi graph → an undirected graph with multiple edges btw vertices & ^{e.g.} loops allowed.

I multi graph



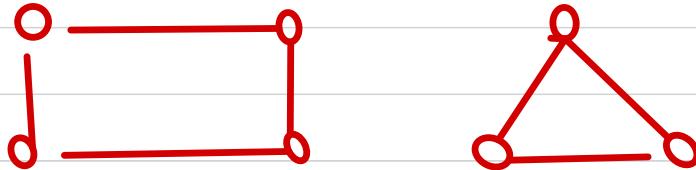
Simple graph → An undirected graph in which both multiple edges & loops are not allowed.

Complete graph → A graph in which every vertex is directly connected to every other vertex.



Connected graph → A graph in which there is some path between 2 vertices but not necessarily direct

Disconnected graph \rightarrow at least 2 vertices do not have
a path to every other vertex:



Component \rightarrow a subset of a disconnected/connected
graph which is connected.

#path \Rightarrow A path p_n is a graph whose vertices can be arranged in some sequence such that



$$V = \{v_1, v_2, v_3, v_n\}$$

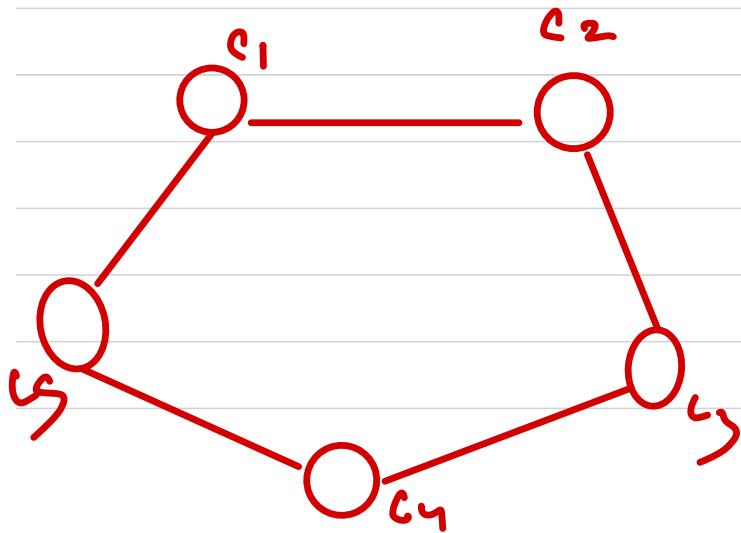
edge set of the graph is

$$E = \{v_i v_{i+1} \mid i \in [1, n-1]\}$$



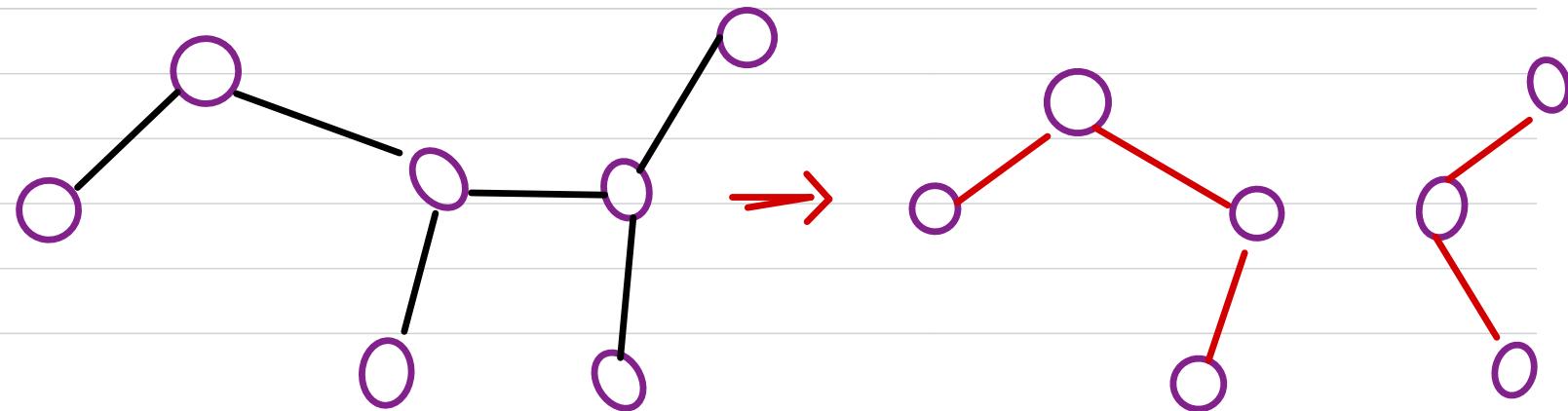
Cycle \rightarrow A cycle C_n is a graph whose vertices can be arranged in a cyclic sequence. such that edge set is

$$E = \{v_i; v_{i+1} \neq i \in \{1, n-1\}\} \cup \{\underline{v_i, v_1}\}$$

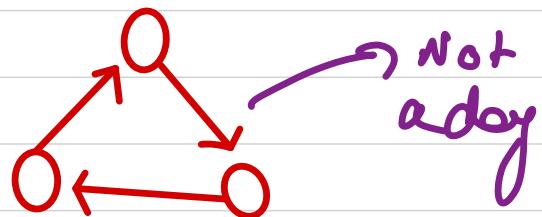


TREE → Tree is a connected graph with no cycles

forest → If we remove an edge from tree we get a forest which is collection of trees



DAG (directed acyclic graph) \rightarrow



Not
acyclic

How to read graphs

As graphs are non-linear, we need some mechanism to read graphs.

Graph traversal

→ Depth first

→ Breadth first



Depth first search

TC
?

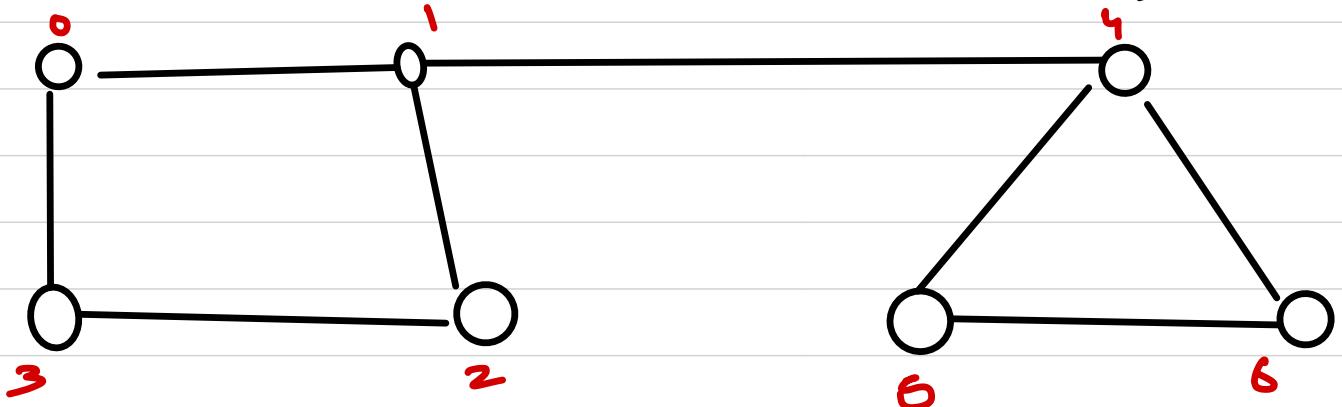
$O(V+E)$

* Motivation problem:

~~Q.~~ Given a graph calc all paths between two vertices

OR

~~Q.~~ Given a graph check whether there is a path between 2 vertices.



is there a path from 0 - 5

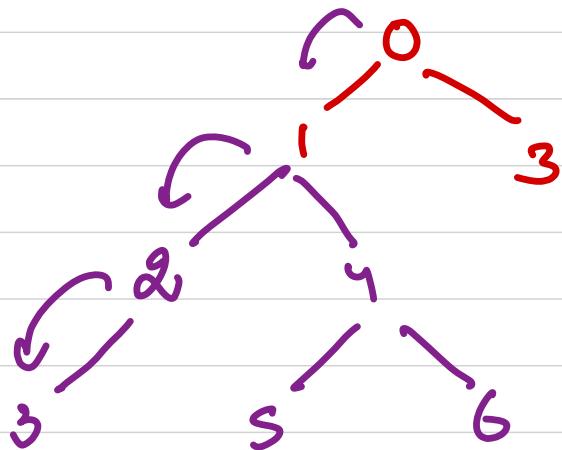
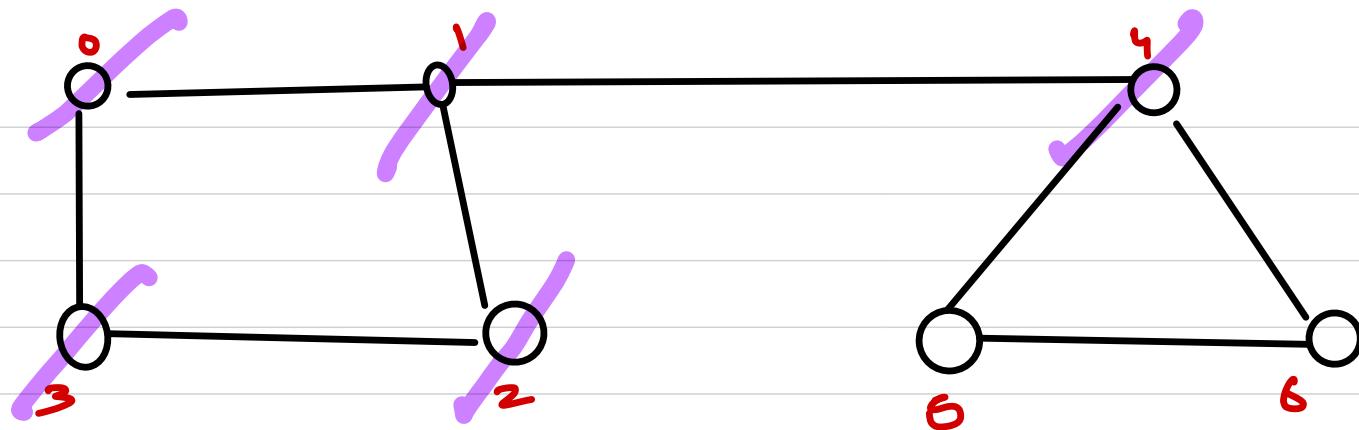
What is the most easiest query for this? ?

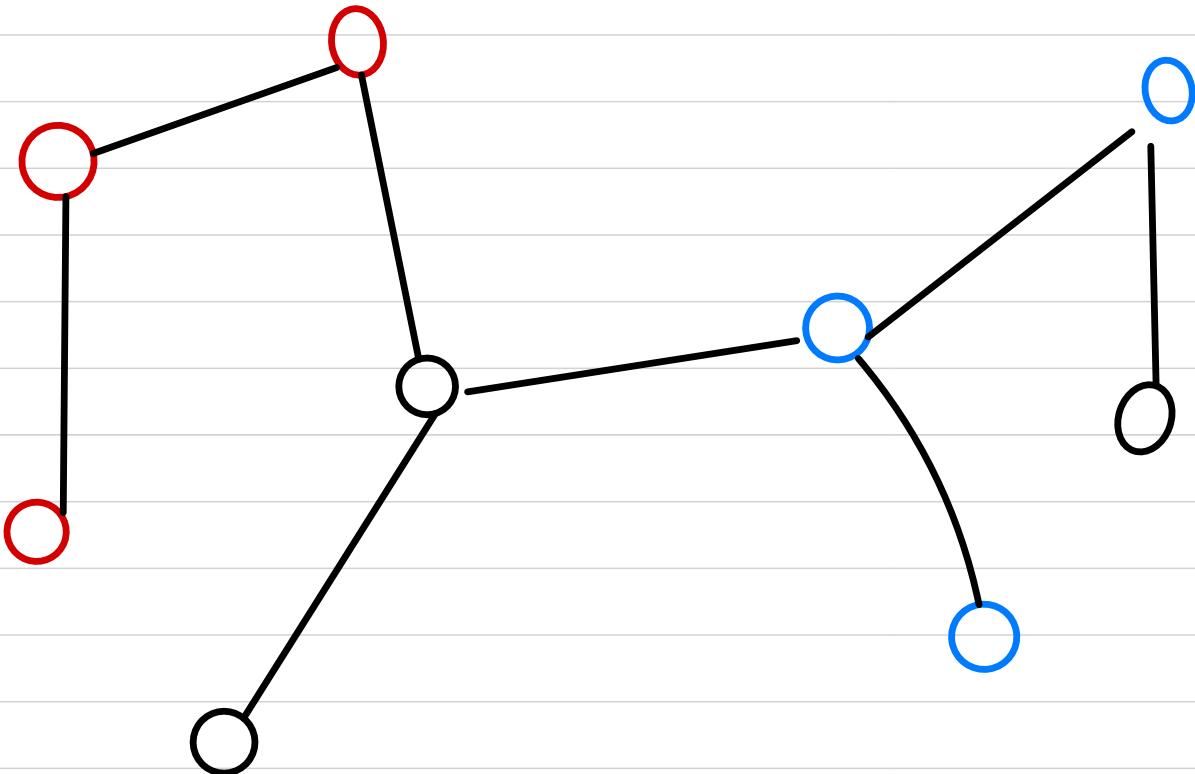
→ neighbours

$$f(u, v) = \begin{cases} f(n_1, v) \\ \text{or} \\ f(n_2, v) \\ \text{or} \\ f(n_3, v) \\ \vdots \end{cases}$$

whether there is
a path from
u to v.

where $(n_1, n_2, n_3, \dots) \in$ neighbors of u.





→ total red
total blue

Proof that a tree with n nodes has $n-1$

~~edges.~~

P MI $\rightarrow f(n)$
no. of edges for
 n nodes

$$f(1) \rightarrow 0$$

$$f(2) \rightarrow 1$$

To prove $\rightarrow f(n+1)$

We assume function works for $f(1) \rightarrow n-1$

Total no. of edges \rightarrow no. of edges required for $(n+1)^{th}$ node
+ $(n-1)$ edges

Every node that will be added to a tree requires

only 1 node.

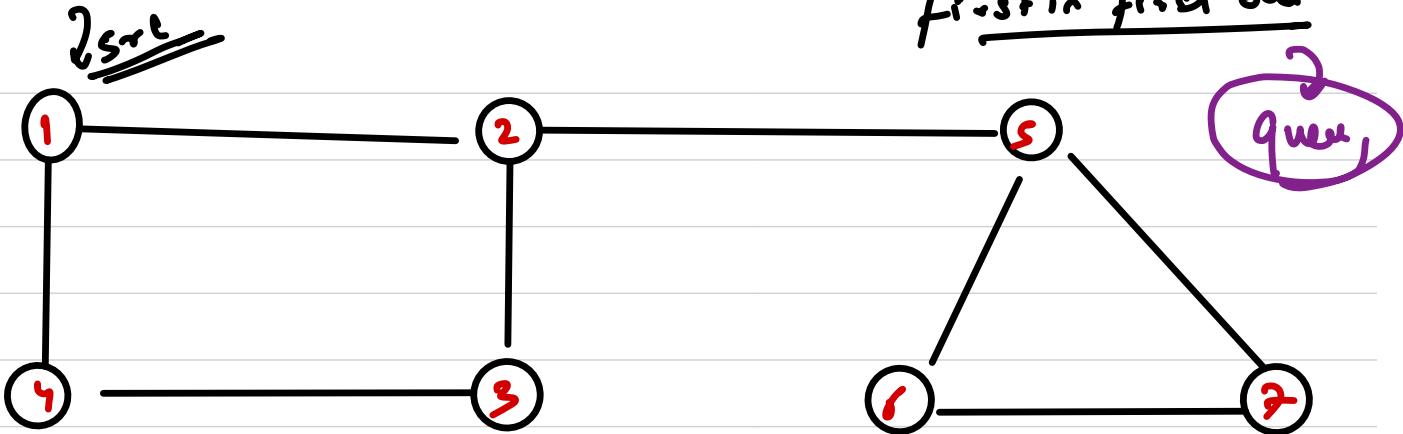
$$\boxed{f(n+1) = (n-1) + 1 = n}$$

equivalent
to level order
traversed

Breadth First Search



Say we have an arbitrary src node then all the nodes
shortest (by edges)
with same distance from src node are said to be
on the same level.



1, 2, 4, 3, 5, 6, 7

Breadth first
traversal

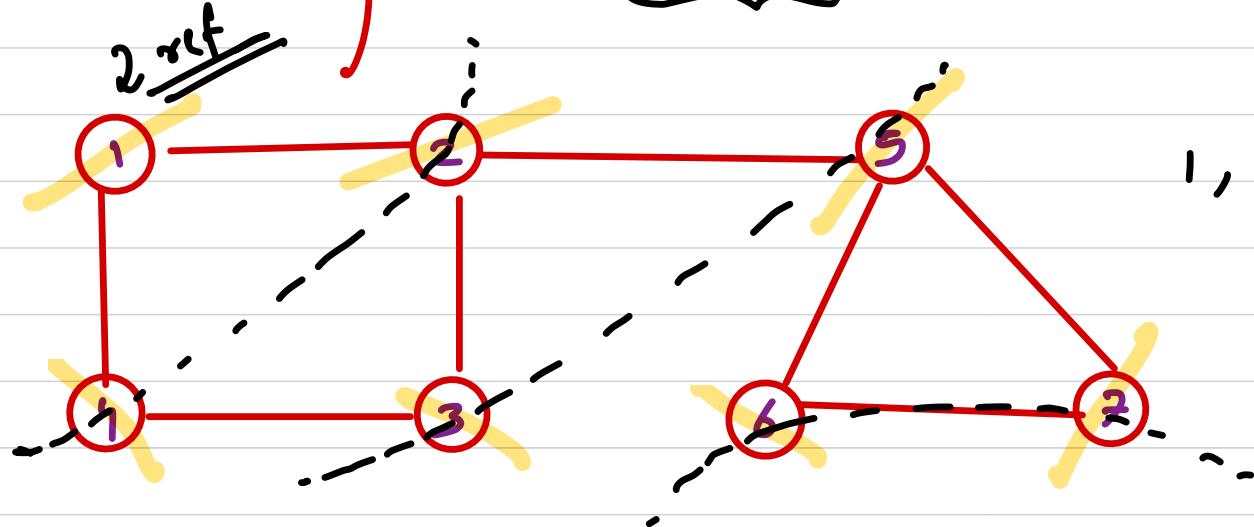


→ queue

unreliable

Breadth First Search

$f(f)$



1, 4, 2, 3, 5, 6, 7



1, 2, 4, 5, 3,
6, 7

X X X X X X X X

5 queen

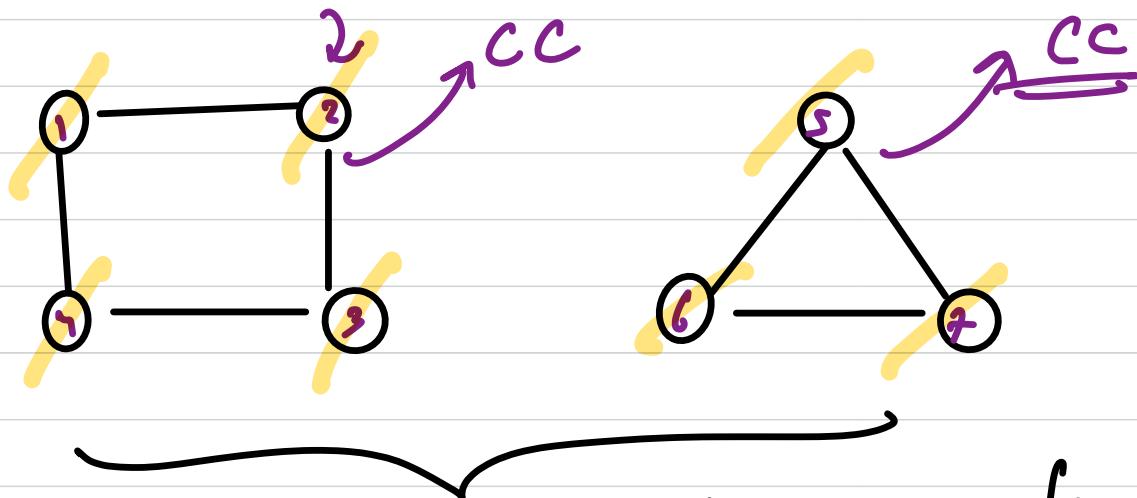
Time $\rightarrow O(V+E)$
Space $\rightarrow O(V)$

So in breadth first search, when we travel from γ of node \rightarrow any node then it always consider shortest path.

Connected

Components

\rightarrow Dfs / Bfs



Disconnected
graph

$vis = \{ \}$

iterate on all
possibly vertices

```
for (i = 1; i <= 8; i++)  
if (!vis[i])  
dfs(i)
```

So the no. of times we will call dfs will be
the count of CC.

Journey to the moon

pair of ost id $\propto a_1, a_2 \geq$ Then a_1, a_2 are
from same Country

Choose $\rightarrow 2$

Say, total no. of ast = n

$${}^n C_2 \rightarrow \frac{n!}{2!(n-2)!} \Rightarrow \frac{n(n-1)}{2}$$

? ?

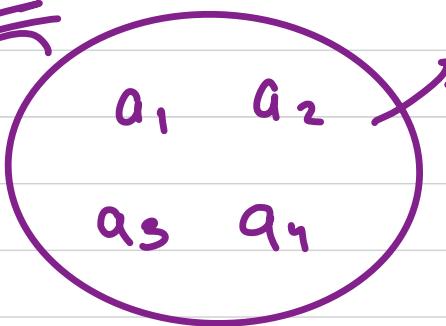
$\langle a_1, a_2 \rangle$

$\langle a_2, a_3 \rangle$

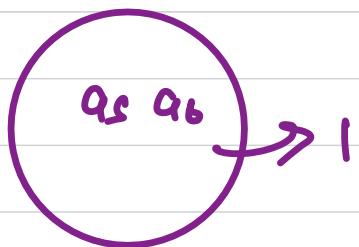
$\langle a_3, a_1 \rangle$

$\langle a_5, a_6 \rangle$

gack



$${}^4 C_2 \rightarrow \frac{4 \times 3}{2} = 6$$



$${}^6 C_2 \rightarrow \frac{6 \times 5}{2} = 15$$

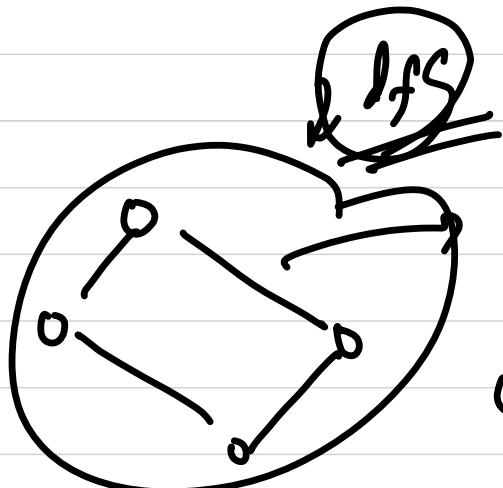
$$15 - 6 - 1 \\ \Rightarrow 8$$

Permutation Swaps

[11, 21, 31]
 ^ ^ ^
 1 2 3

: d 4 1 2 , 7
A > 1 3 2 4
B > 1 4 2 3

1, 2
2, 3



(3, 4) ← index

those indices
are in the same
connected
component

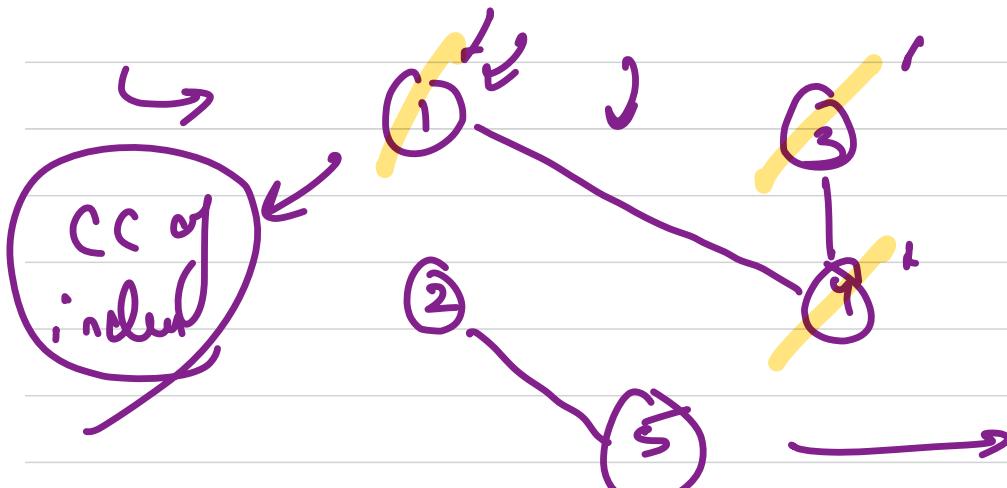
idn 1 2 , 1 5



(2, 3)

B \rightarrow 1 5 2 3 4

(1, 4)
(1, 5)



relabel

$P \rightarrow [1, 2, 3]$

$q \rightarrow [1, 3, 2]$

$P \rightarrow [q, s]$
 $q \rightarrow [s, r]$

$i d_1$ 1 2 3 4 5

$(2, 5)$

4 1 3 4 3 2

$(3, 4)$

5 1 4 2 3 5

$(1, 4)$

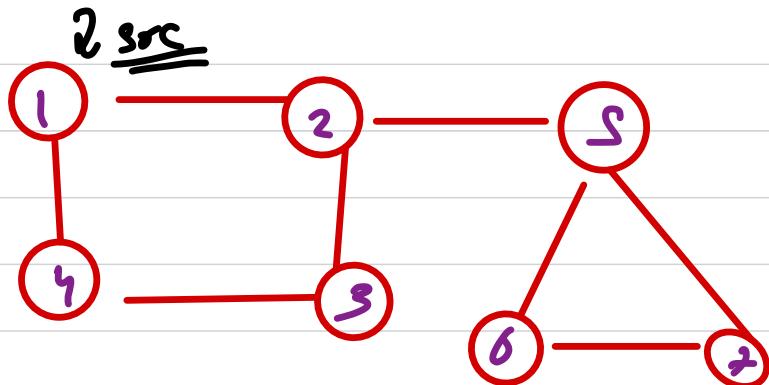


$\rho \rightarrow [1, 3, 4]$

$q \rightarrow [1, \cancel{2}, \cancel{2}, 3]$



~~Q~~ Given an ^{unweighted undirected} graph, and a source node calc the length of the shortest path to all the other nodes from the source node.



→

0	1	2	1	2	3	3
1	2	3	4	5	6	7

↑ output

BFS

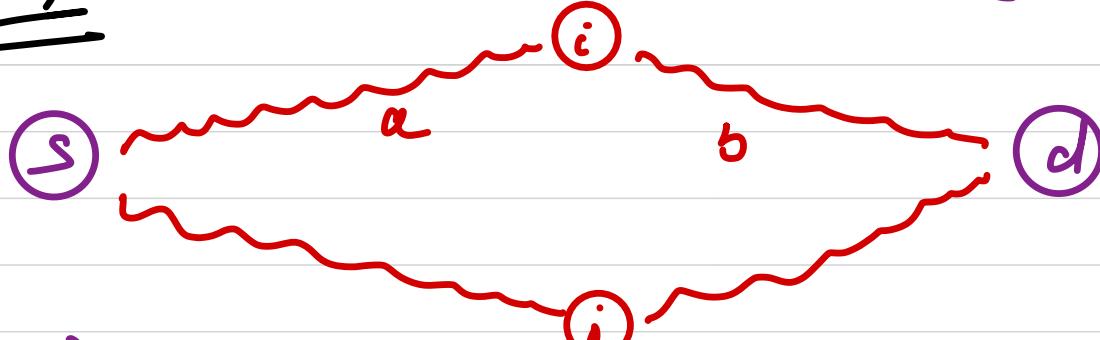
BFS

Bfs traverse the nodes based on shortest

path

Q Given a graph, a src and a destination,
find all the nodes which are a part of
at least 1 shortest path.

O(v + e)



$\text{dist} \rightarrow (s, d) \rightarrow x$

$s[i] + d[i] = x$

$a + b = x$
 $b = x - a$

Rotten

Oranges

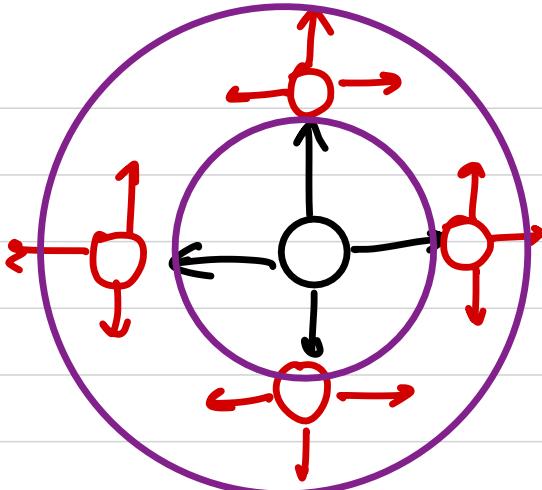
6	2	X2	Y2	2
2	Y2	X2	0	0
0	21	0	X2	2
0	0	2	2	2
0	0	0	0	0

(2D array)

m x n

-1 → when no
other oranges
can rot a
orange

$t = t'$ → all oranges which
are rotten will offset their
neighbours at same time



BFS

queue

vis

(i, j)

Multi Source BFS

$O(nm)$

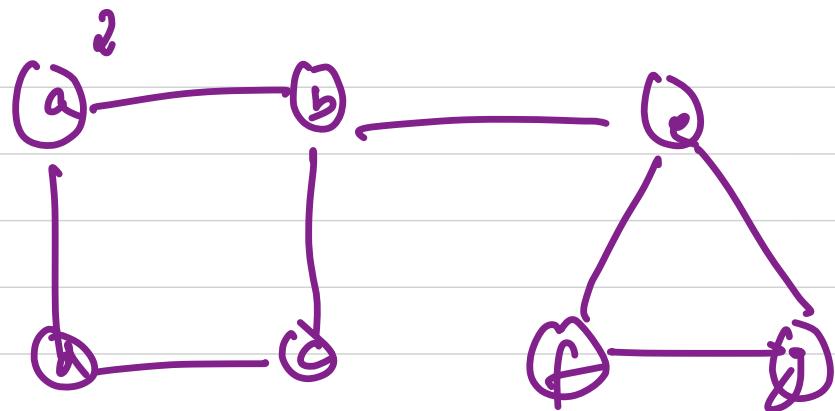
$(i+1, j) \quad (i-1, j) \quad (i, j-1) \quad (i, j+1)$

all this will be marked
if there is a fresh orange

$\begin{array}{c} \uparrow \\ \leftarrow 2 \rightarrow \\ \downarrow \end{array}$

$\begin{array}{c} 1 \\ \leftarrow 2 \rightarrow \\ \downarrow \end{array}$

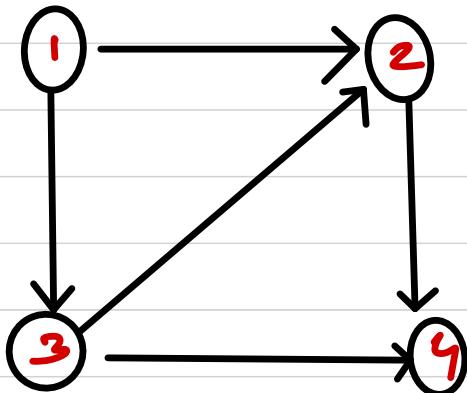
$\begin{array}{c} 9 \\ \leftarrow 2 \rightarrow \end{array}$



~~graph LR; a((a)) --- b((b)); a --- d((d)); b --- c((c)); b --- a; c --- f((f)); c --- e((e)); d --- e; d --- f; e --- f; e --- c; g((g)) --- c~~
fig

Longest Path

Dynamic Prog.



→ The longest directed path can start from any node.

$f(v)$

1D DP
=

returns the length

of the longest path

that starts with v.

when $n_1, n_2, n_3, \dots, n_m$ are
the neighbours of v

ans $\max(f(1), f(2), f(3), \dots, f(n))$

$1 + \max$

$f(n_1)$

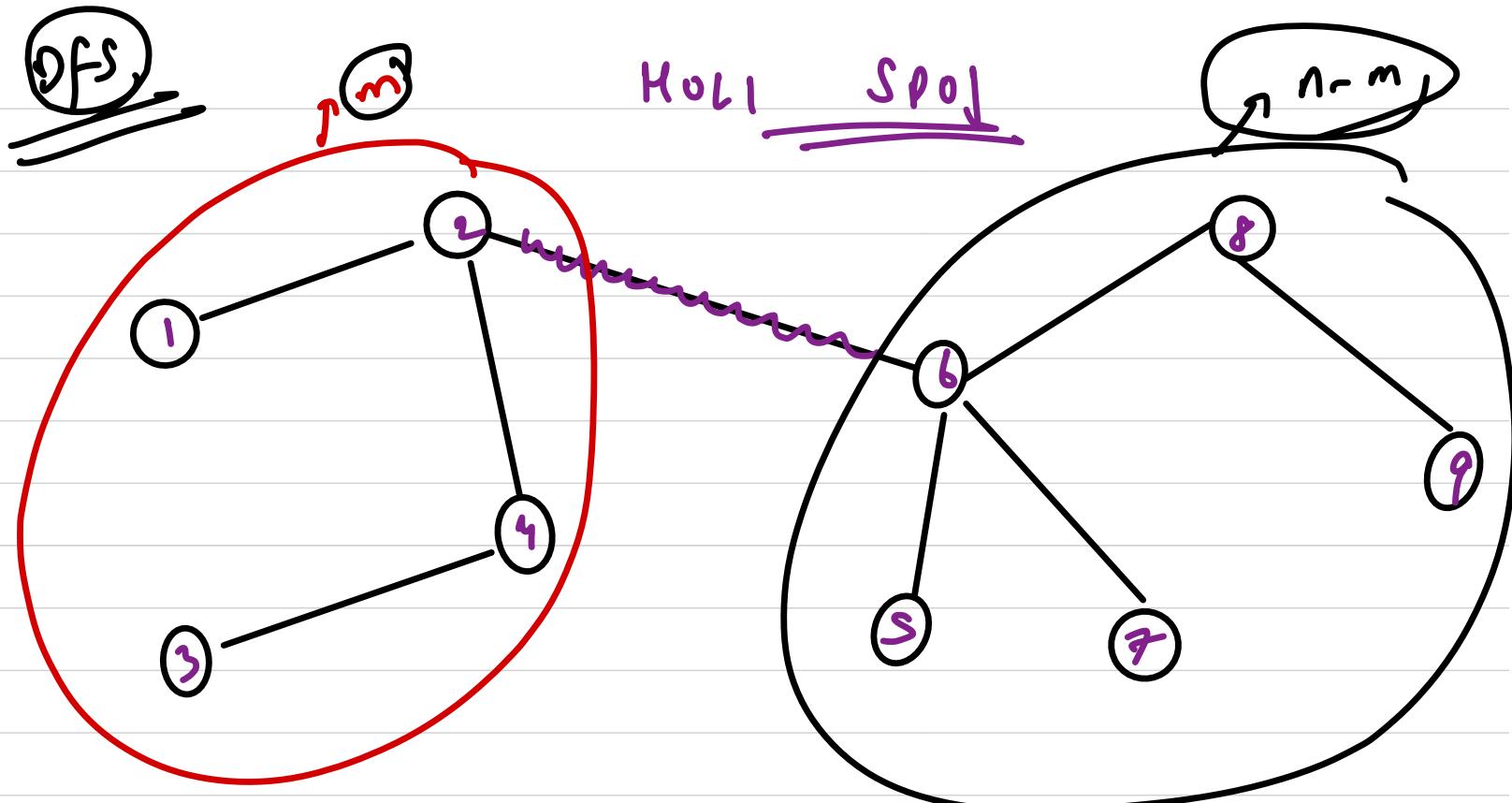
$f(n_2)$

$f(n_3)$

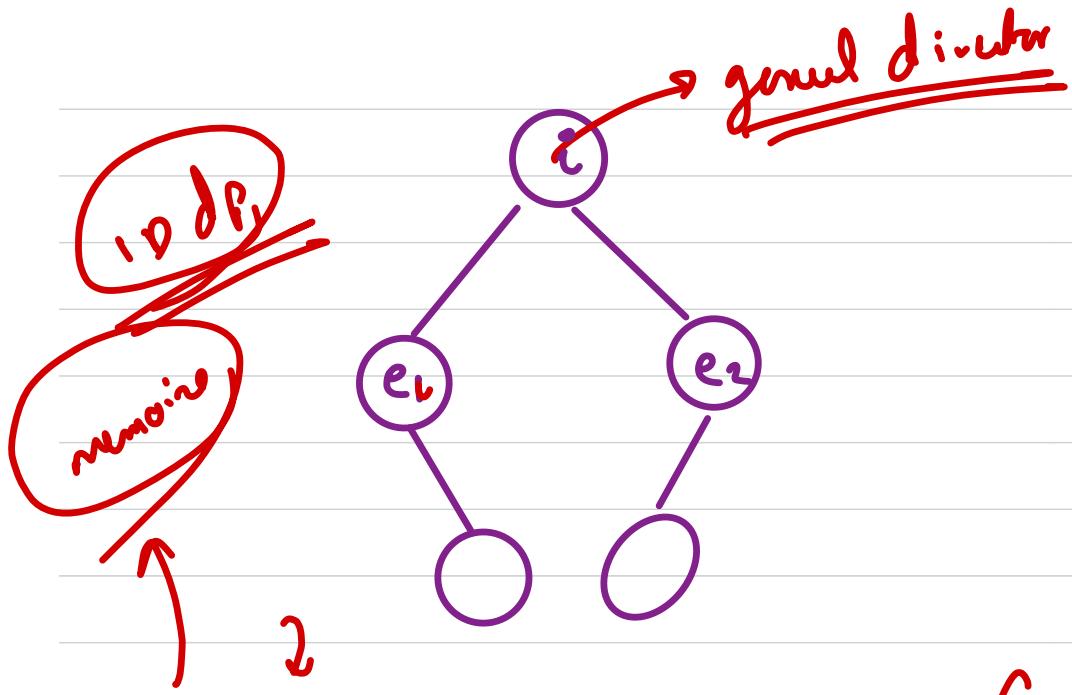
⋮

$f(n_m)$

Base
case
length
0

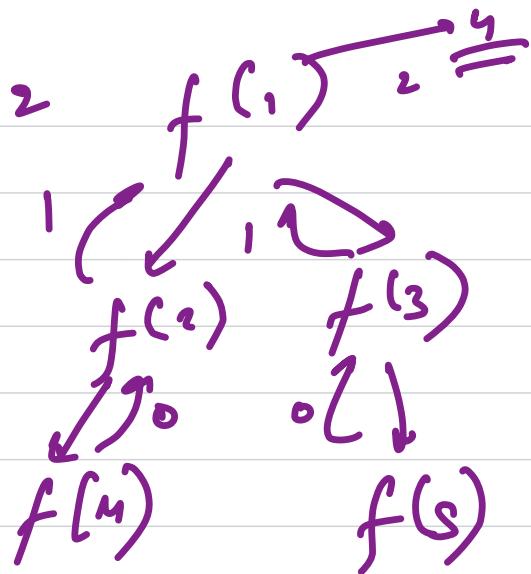
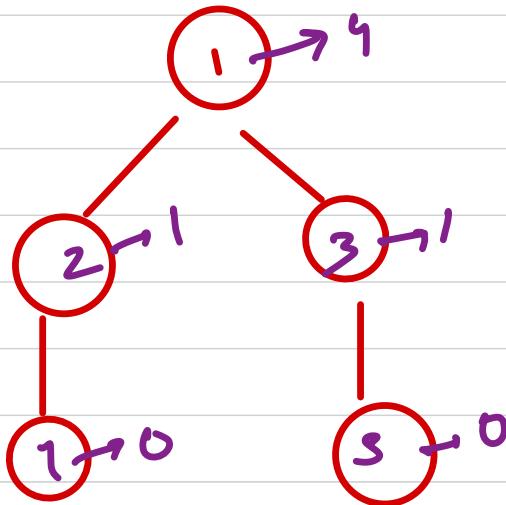


$$\underbrace{g(n-m, m)}_{\text{ }} \times \overrightarrow{\text{wt}}$$

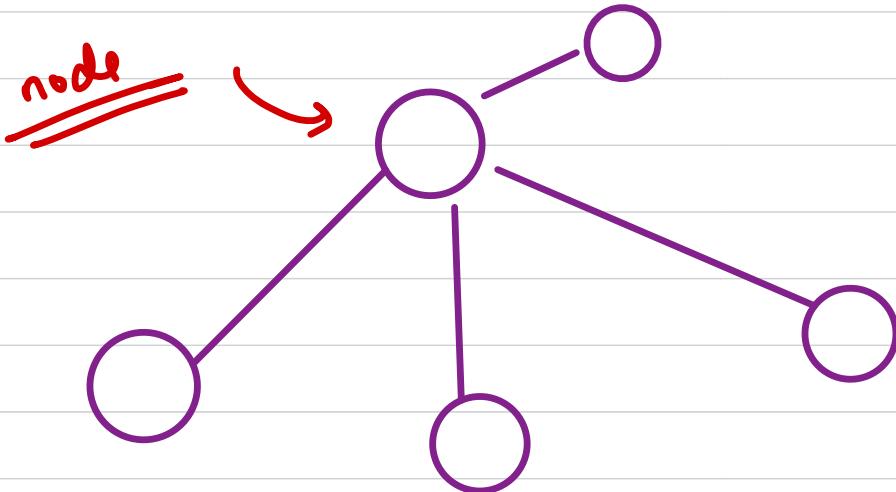


$$f(i) = \sum (1 + f(e_i))$$

↓
 no. of subordinates
 under employee i



Independent Set → Attoder



color $\begin{matrix} \text{red} \\ \text{white} \end{matrix}$ $\begin{matrix} \text{B} \\ \text{W} \end{matrix}$

Black → cast
White → no cast

No, two adjacent vertices should be black

$f(\text{node}, c) \rightarrow$ these two parameters can
uniquely identify our state

$c \rightarrow \text{constraint}$

$f(\text{node}, 0) \rightarrow$ # of ways to color subtree
rooted at node

as $C \rightarrow 0$ then node can be w or b

$f(\text{node}, 1) \rightarrow$ # of ways to color subtree rooted
at node and node is under constraint

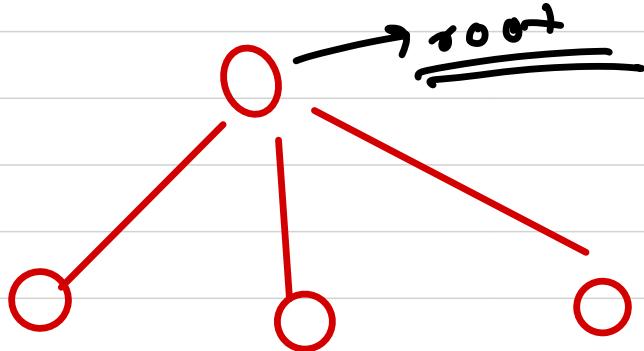
$$\begin{cases} f(\text{node}, 0) = \pi f(\text{child}, 0) + \pi f(\text{child}, 1) \\ f(\text{node}, 1) = \pi f(\text{child}, 0) \end{cases}$$

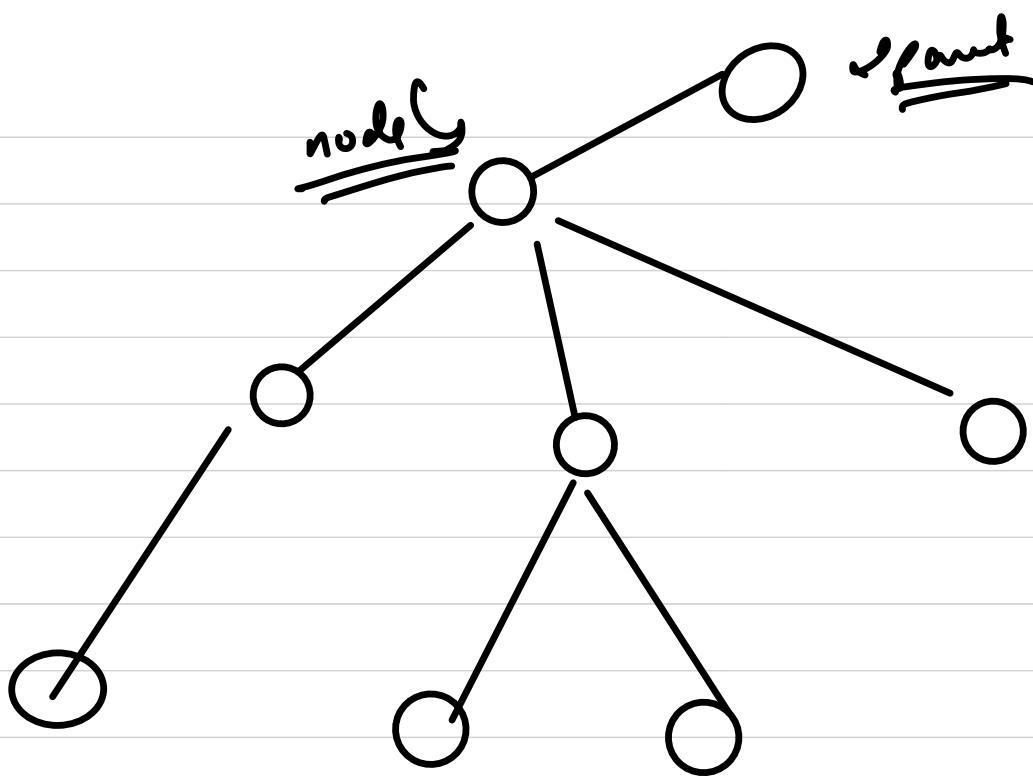
final ans $\rightarrow f(\text{root}, 0)$

LeetCode

Binary Tree Cameras

Brute force





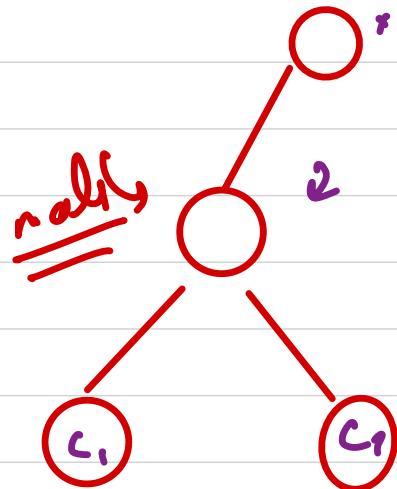
$f(\text{node}, \text{cam}, \text{par cam})$

↓

return the min number of cans

to cover the subtree rooted at node.

Final ans $\rightarrow \min(f(\text{root}, 1, 0), f(\text{root}, 0, 0))$



for current node , we decided to have
or can the parent is not have a car

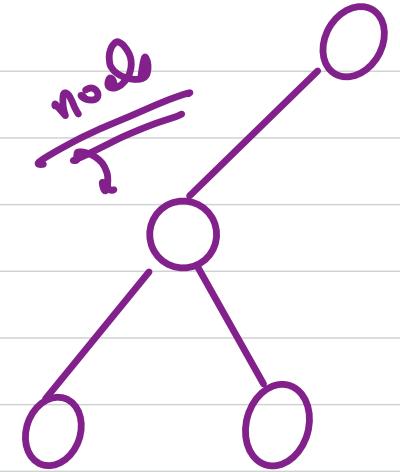
$$f(\text{node}, 1, 0)$$

for your cars

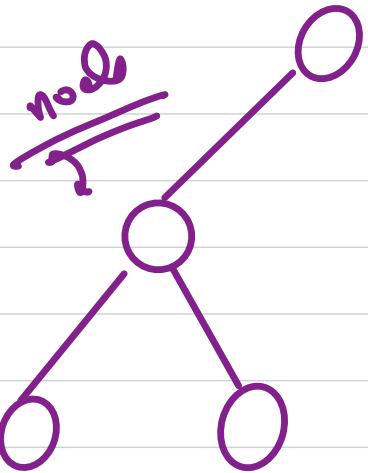
$$f(\text{node}, 1, 0) = 1 + \min(f(c_0, 1, 1), f(c_0, 0, 1)) + \\ \min(f(c_1, 1, 1), f(c_1, 0, 1))$$

when node isn't key can but parent
has it

$$f(\text{node}, 0, 1)$$



$$\begin{aligned} f(\text{node}, 0, 1) = \min & (f(c_1, 0, 0), f(c_1, 1, 0)) + \\ & \min (f(c_2, 0, 0), f(c_2, 1, 0)) \end{aligned}$$



$$f(\text{node}, 0, 0) =$$

$$f(\text{node}, 0, 0)$$

so dp

one of the children are bounded to have a can.

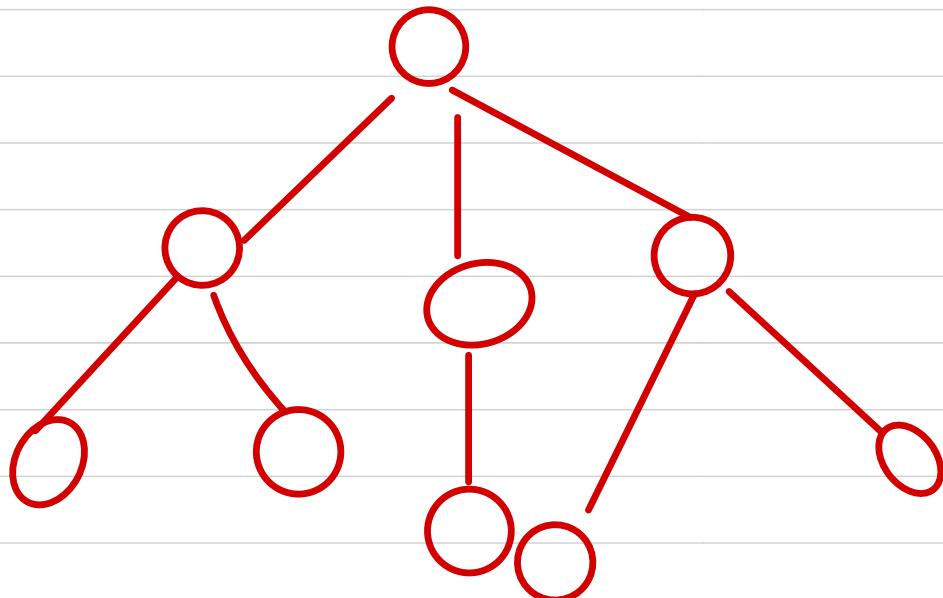
$$\begin{cases}
 f(c_1, 1, 0) + \min(f(c_2, 1, 0), \\ f(c_2, 0, 0)) \\
 f(c_2, 1, 0) + \min(f(c_1, 1, 0), \\ f(c_1, 0, 0))
 \end{cases}$$

ans = ∞

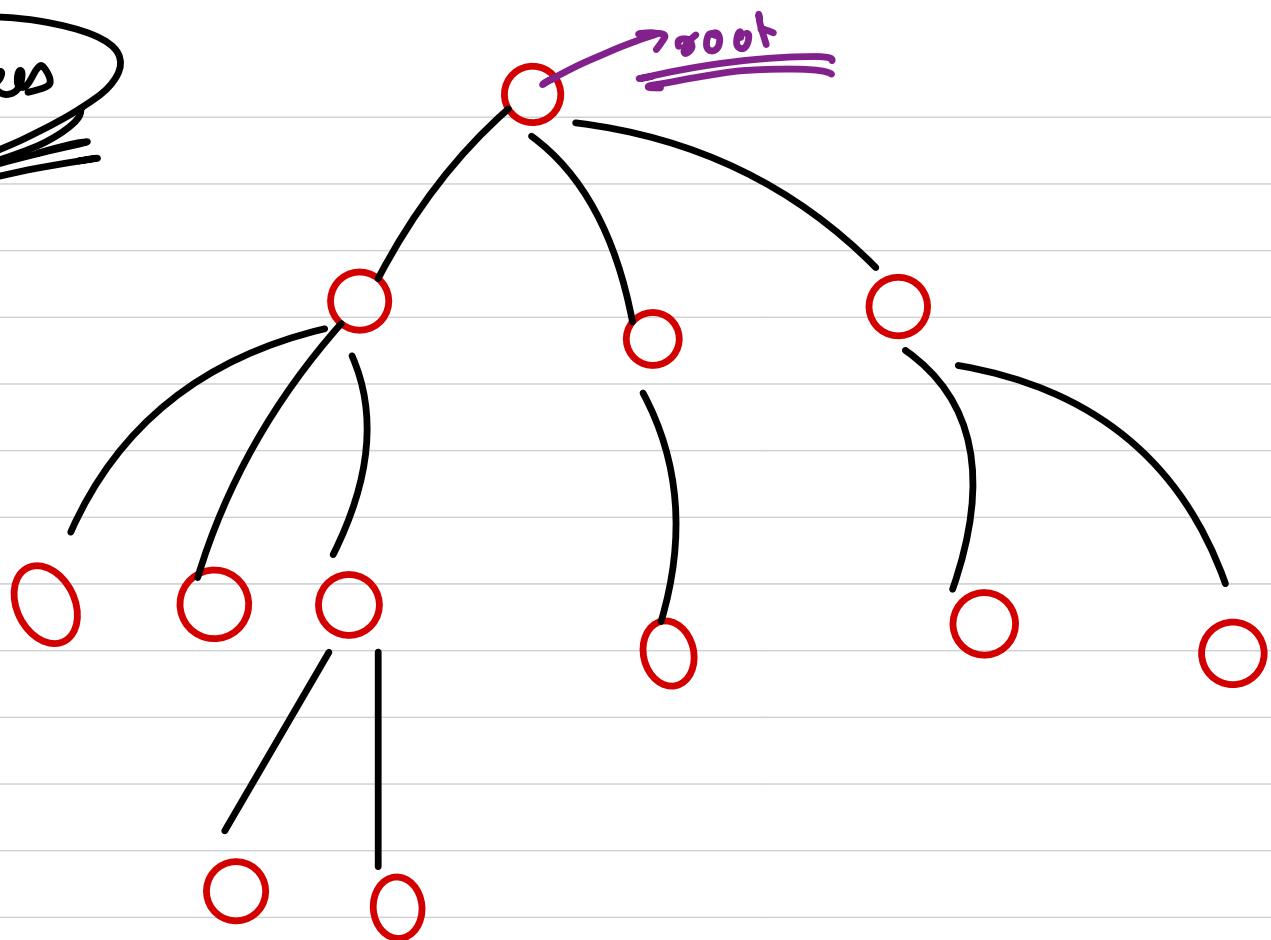
fam tree

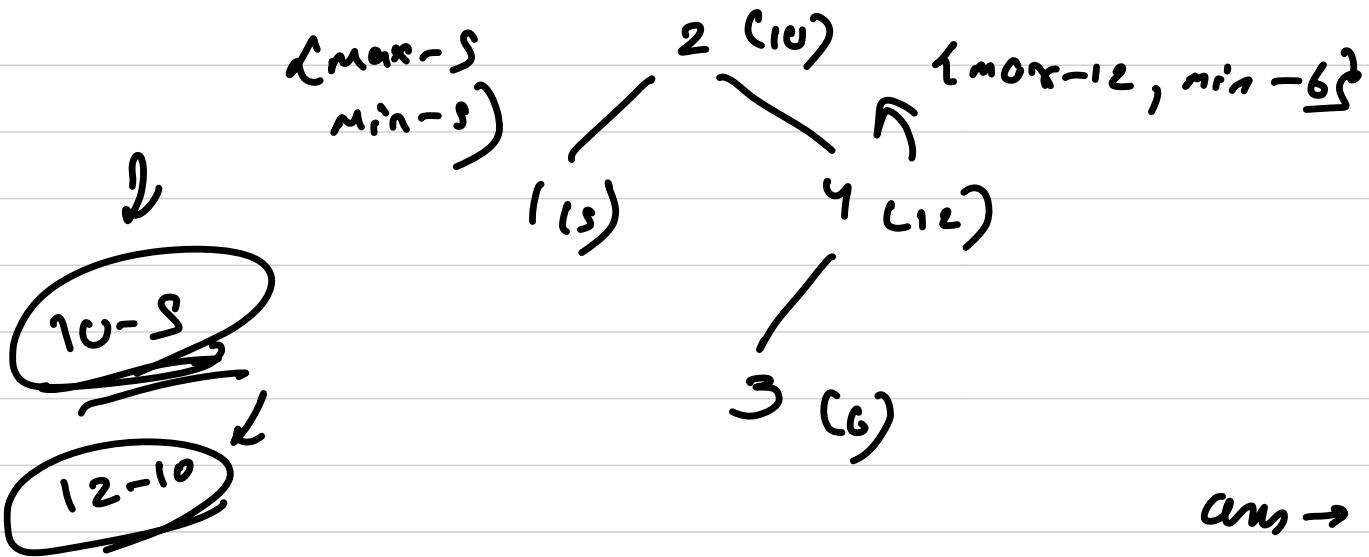
(Codichef)

max
min



Dynam trees





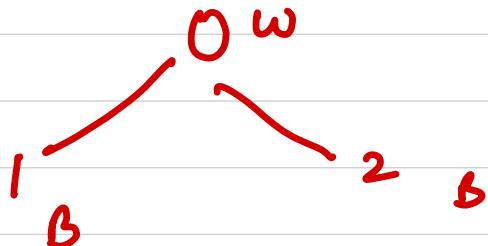
~~ans → -05~~
~~6~~
 = =

#applemans and tree → codforces

Divide this tree into arbitrary k components

such that each component has exactly one black

node-



$$0 \leq k \leq n$$

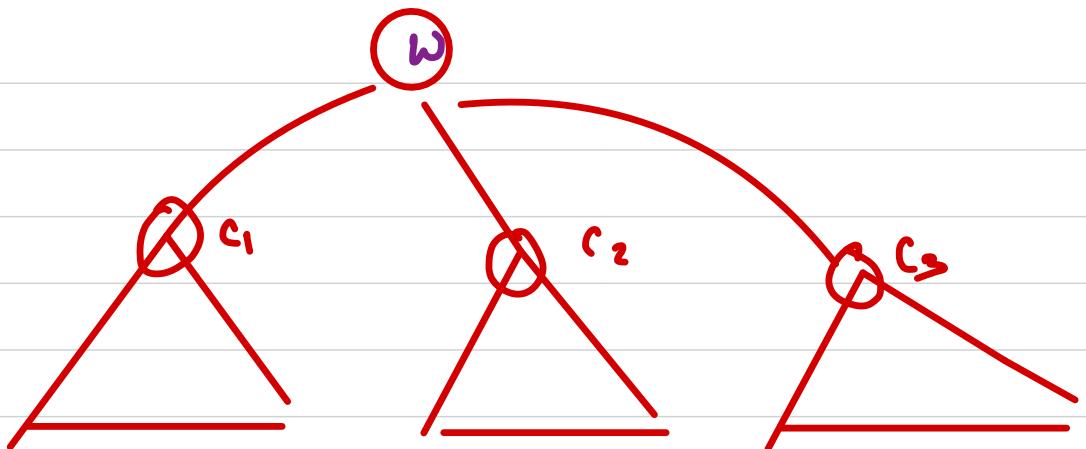
tree →
 (n^m)

every node will have different choices
based on colors of node inSubtree

$f(\text{node}, c)$ \downarrow
count
of black
n-nodes

$f(\text{node}, 0) \rightarrow$ # of ways subtree rooted at node
has no black node vertex

$f(\text{node}, 1) \rightarrow$ # of ways for subtree rooted at node
& has 1 black vertex \equiv



$$f(\text{node}, i) = \pi(f(c_1, 0) + f(c_1, 1))$$

$$f(\sim \text{node}, 0) = 0$$

$$f(\text{node}, 0) = \pi(f(c_1, 0) + f(c_1, 1))$$

$$f(\text{node}, 1) = \sum f(c_{1,1}) \times \frac{f(\text{node}, 0)}{f(c_{1,0}) + f(c_{1,1})}$$

~~Qn~~ You're given a set of inequalities .
~~Eg~~ $(a < b \quad c > d \quad b < d \dots)$

You've to find that whether for any set of int values for the variables, can be solve the inequality .

Dependency Resolution

- ↳ $a < b$, $c > d$, $b < d$
- ↳ package of a software
- ↳ course structure

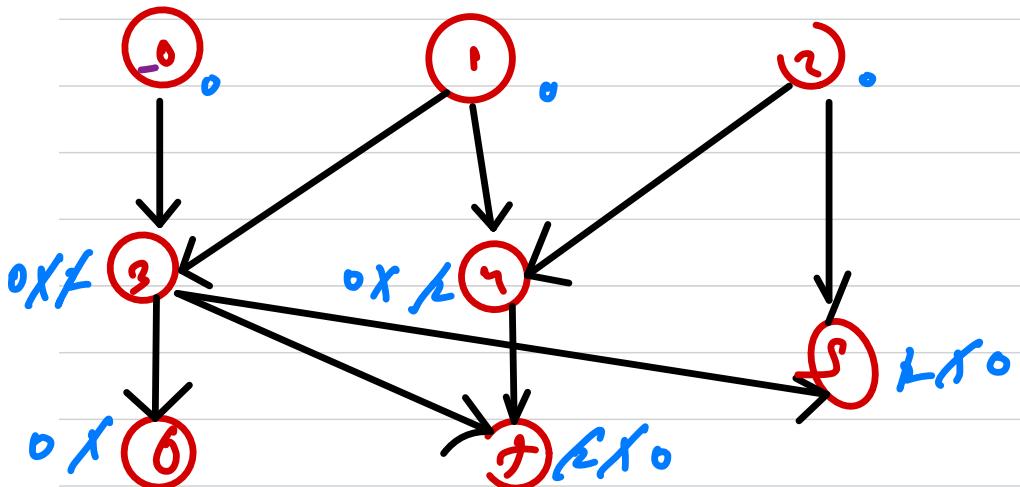
→ directed graph

Can be visualized
in the form of
a graph.

In order to resolve the dependency graph we can
use an algorithm called as topological sorting

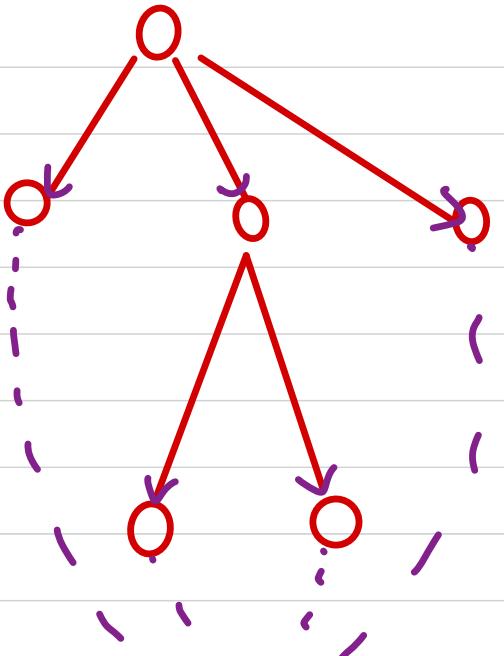
Kahn's algorithm \rightarrow .BFS based topological sort

queue



5, 2, 3, 1, 8, 2, 4, 6

0, 1, 2, 3, 4, 5, 6, 7



0, 1, 2, 3

→ fix the value
here i.e Node N.
0, 1, 2, 3

$f(\text{node}, \text{xor})$

~~Ques~~ Given a grid with char [L, R, U, D]

on each cell of grid. You're standing at top left of the grid. From each cell you can go in the direction written on the cell for ex L → Left.

Decide if it is possible to reach the bottom right corner or not. Dimension → $M \times N$

Explain Space \Rightarrow O(1)

$m, n \leq 10^3$

↳ There is only one case for not ready the bottom right.

↳ if there exist a cycle



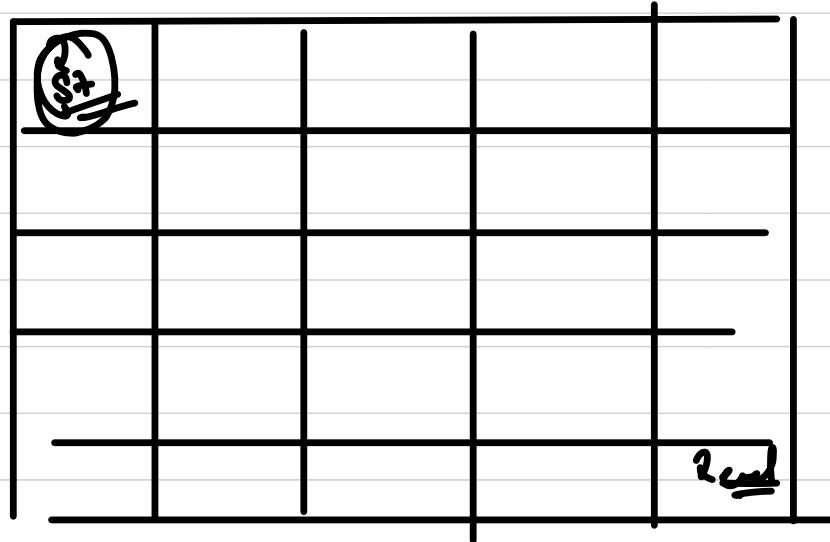
$1 \rightarrow 1x$

$2 \rightarrow 2x$

Say we can swap the characters of the cell with any other then we want. Cost of changing char is 1 unit.

Calc the min cost to reach bottom right.

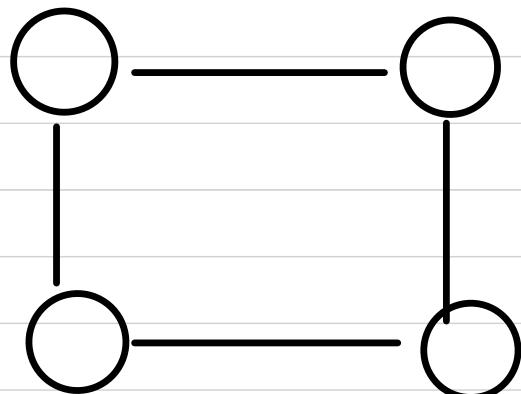
I cant



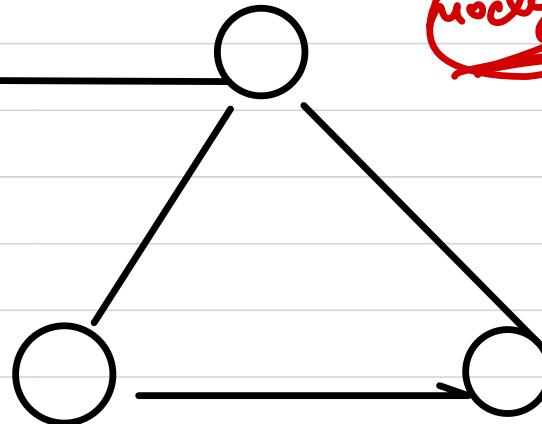
on this weighted graph calc the shortest path

Given a graph & a src and a dest node
Find the shortest path from src \rightarrow dest

? src

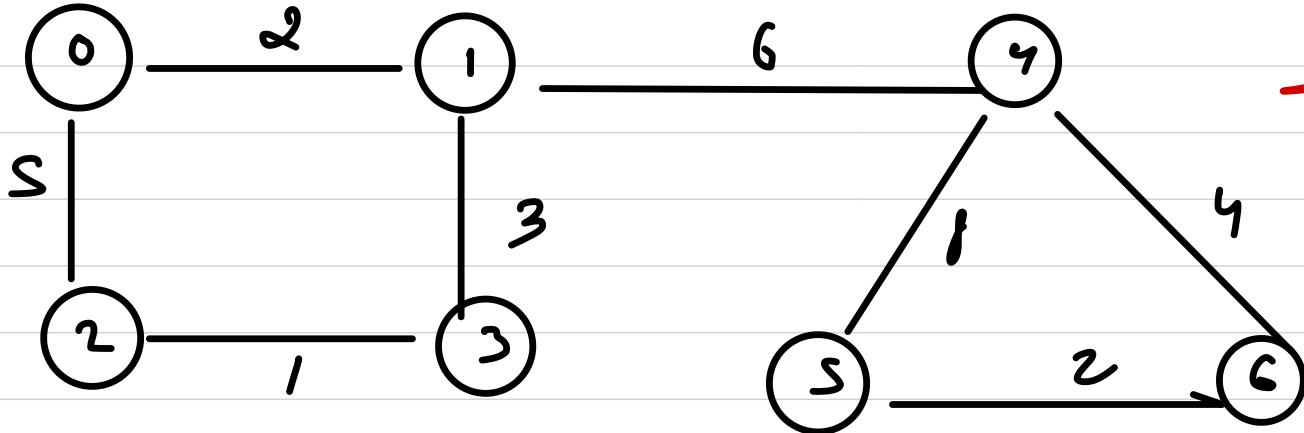


BFS
→ modify



short

28-0



Dijkstra's

Priority Queue



node	parent	wt
0	-	0
1	0	2
2	0	5
3	0	6
4	0	8
5	0	9
6	0	10

Joint

Joint

0	2	1	3	4	5	6
---	---	---	---	---	---	---

Strongly Connected Components

directed graph



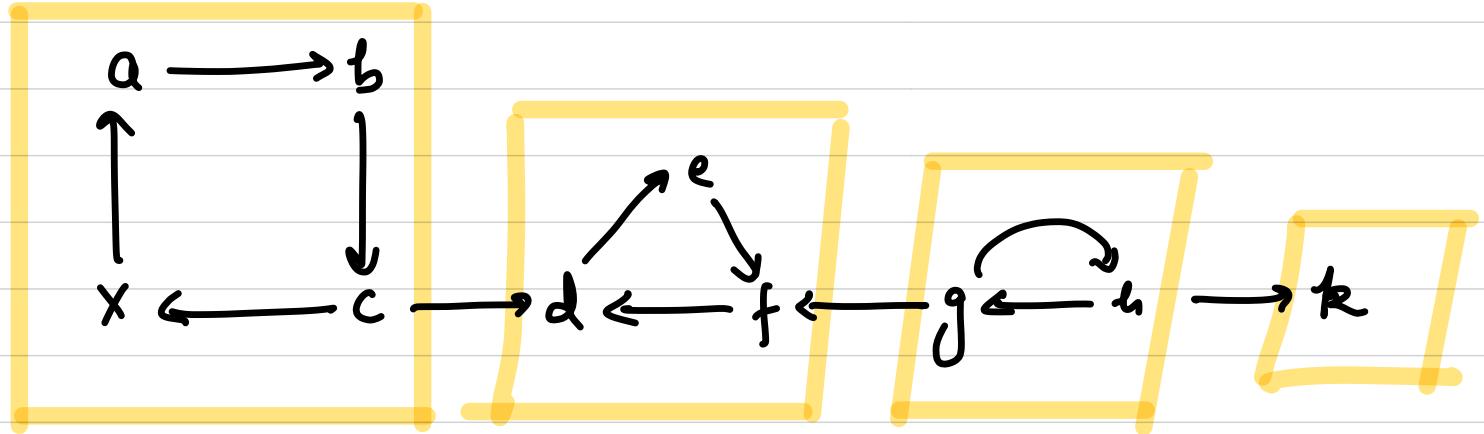
We have a path

from a to b

but not vice-versa

Within a component in a SCC

every vertex should have a path
to other vertex in a directed graph



How can we find SCC ??



for finding normal cc in an undirected graph,
we used to take help of dfs / bfs

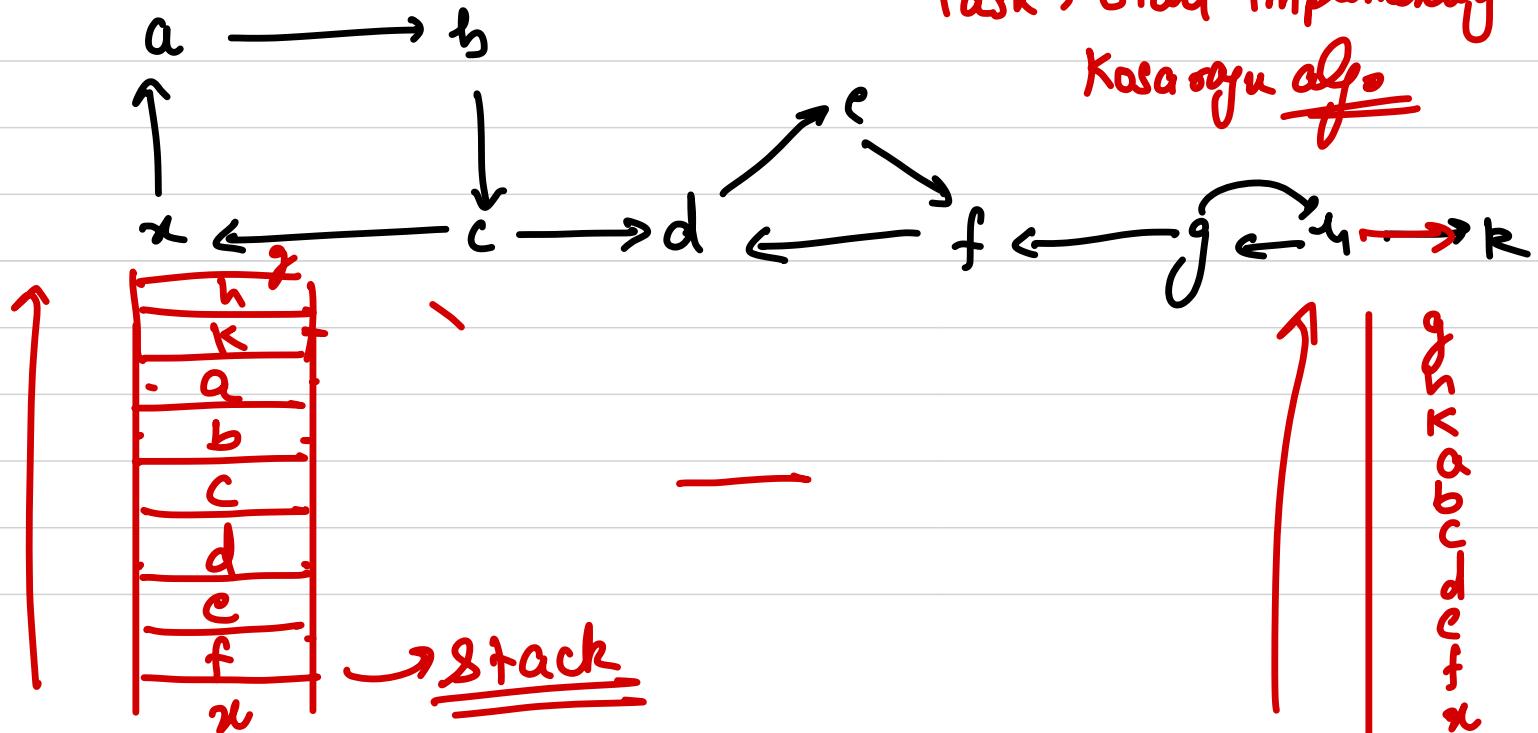
Kosaraju's

Algorithm

visited

a, b, c, x, d, e, f
g, h, k

Task → Start Implementing
Kosaraju algo



We will start bfs/dfs from an unvisited node.

If we have touched all possible nodes in a go

then we restart with any other unvisited if

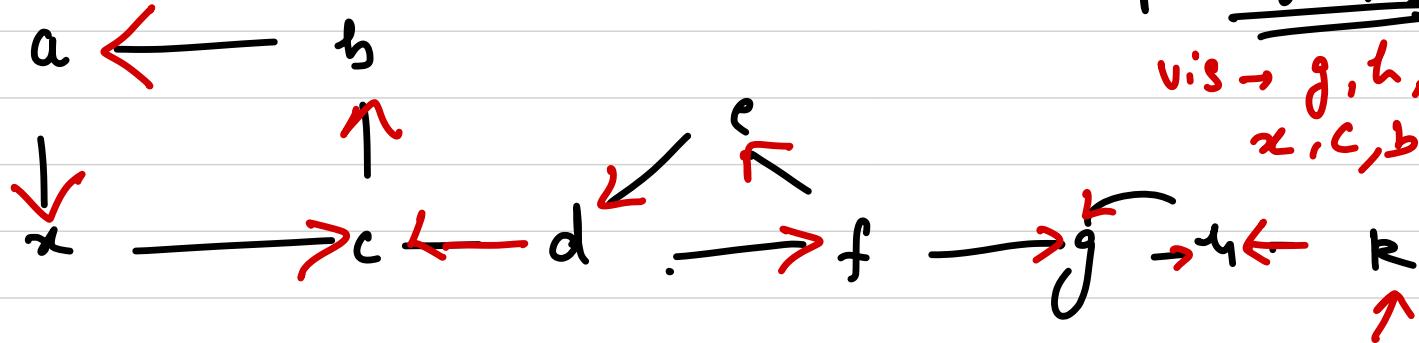
it is left.

Do dfs & store vertices in the stack as per their finish order.

In the phase 2, take transpose of the graph

Start dfs on the transposed graph from the node at the top of stack:

2 transposed graphs



v.i.s → g, h, K, a,
z, c, b, d, f, i

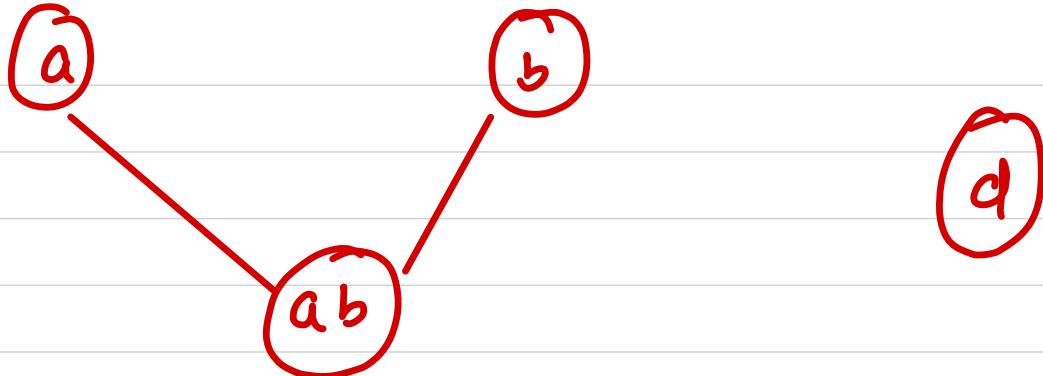
gh

K

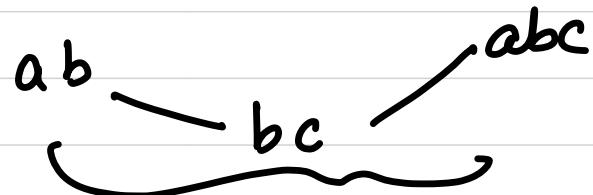
a z c b

d f c

=>



2₁ ofs / Bfs / Dsu



one component

Same component

10^5

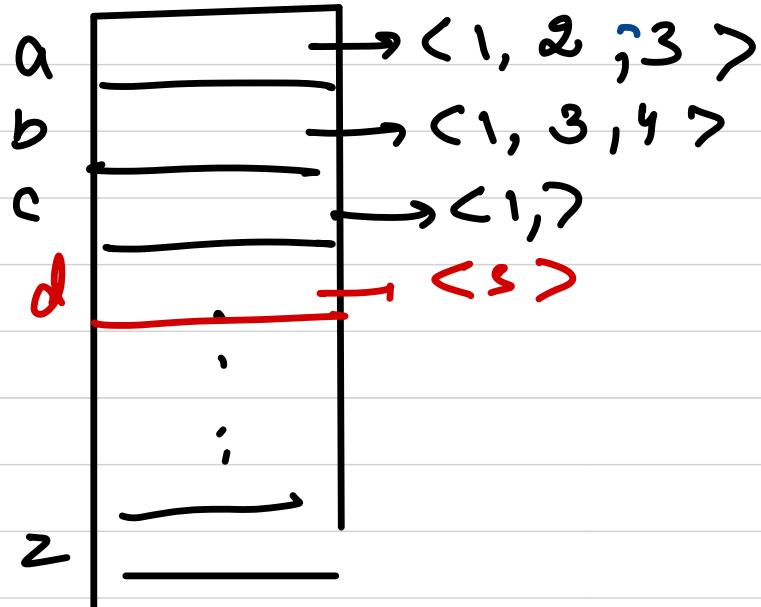
$a b c$

$a b$

b

$a b$

d



4 3 2 1

5

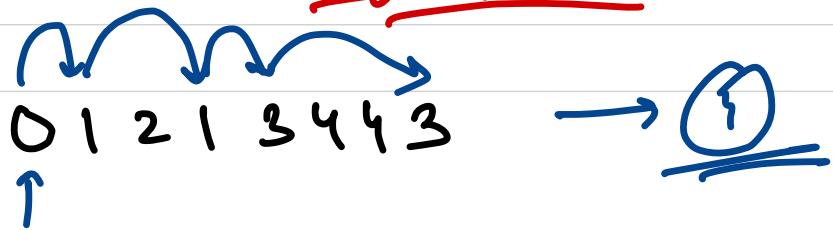
Q7 Given a string of digits of length N. (NS^{10^5})

You're standing at 0^{th} index & want to reach $(N-1)^{\text{th}}$ index. For each jump you take a cost of 1 unit.

And from any i^{th} digit, you can go to $i-1$ or $i+1$

or any other digit S_j such that $S_i = S_j$

Find the min jumps.



0 1 2 1 3 4 4 3

worst case → you need to touch
every digit only
once

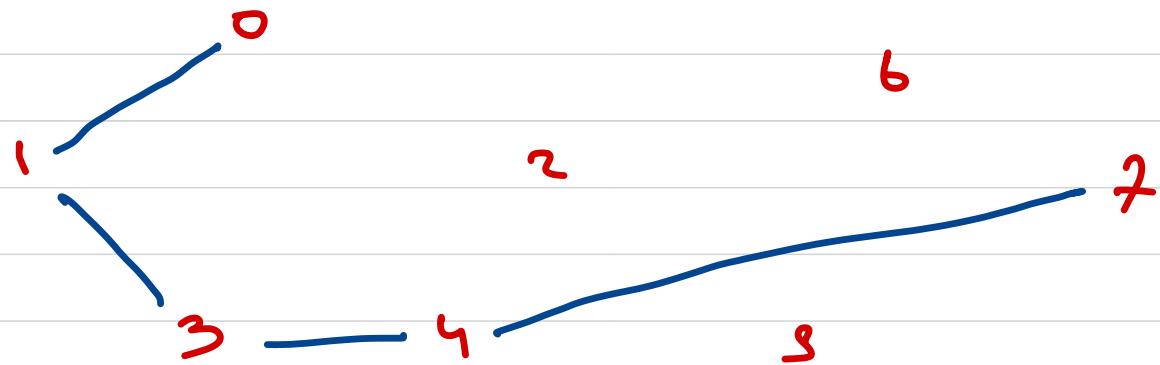
→ 19

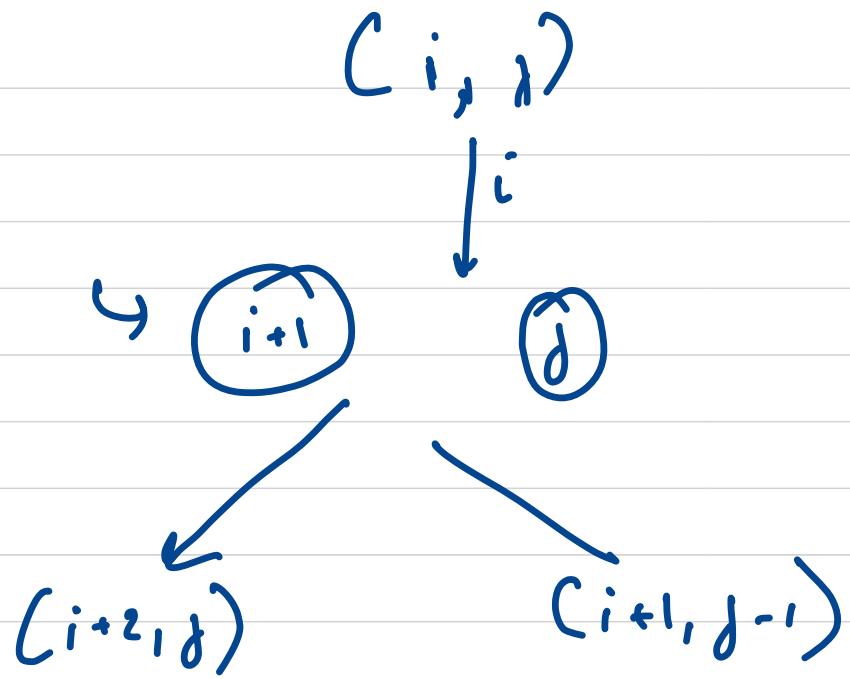
0 0 1 0 0 2 2 2 3 3 4 5 6 6 0 2 3 6 8 9

0 1 2 1 3 4 4 3

0 1 2 3 7 8 6 2

0 → 1 → 2





Disjoint Set Union

Qn Lets say we want to create clusters

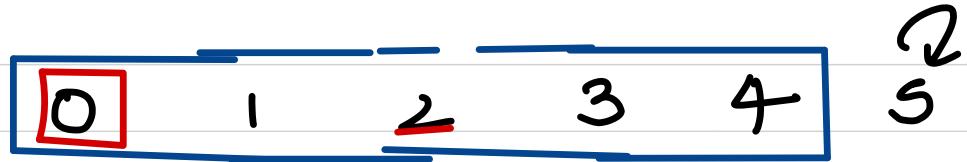
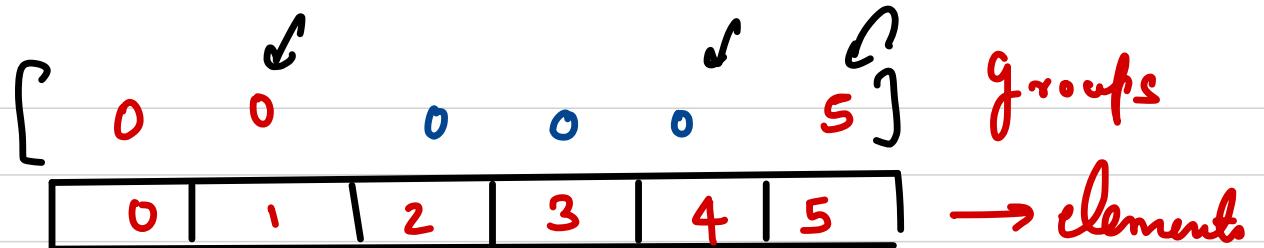
We have a set of elements and we need to group them. Sometimes you might be asked to return the group of any element belongs to.

To uniquely identify a group, we will pick any element of the group & name it the leader/representative / parent of the group.

union (a,b) → Adds b to the group of a or

vice - versa.

get(x)/find(x) → To what group the element
x belongs to.



union(0,1)

union(2,3)

union(3,4)
union(1,3)

find(1) → 0

find(5) → 5

find(4) → 0

int find (x) {
 ↗ $O(1)$

return par[x];

}

void union (a,b) {

a = find (a);
 ↗ $O(n)$

b = find (b);

for i in 1..n {
 if par[i] == b
 par[i] = a

,

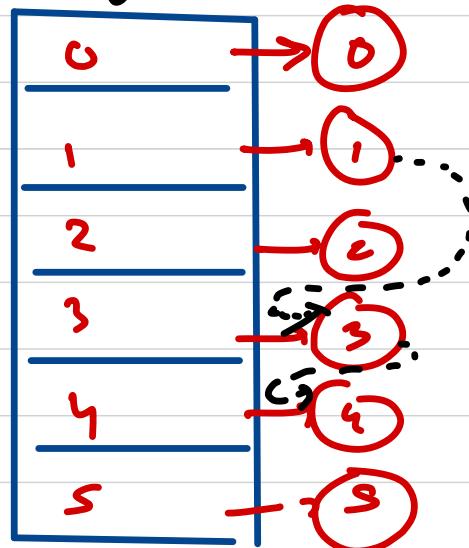
3

Can we
optimize?

How about storing the grouping mechanism such that we represent it as vector of vector or

vector of ll

head tail



union(1,3)

doubly list

find / get $\rightarrow O(n)$

union $\rightarrow \underline{O(1)}$

The core problem is, if we add n elements
then the operation of updating parent array
for each element is $O(n)$

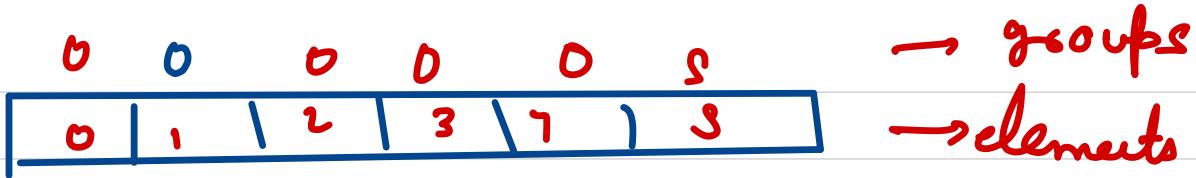
Amortized

$$\frac{1 \times O(n)}{n} \rightarrow \underline{\underline{O(n)}}$$

Somehow optimize updating the parent.

Now let's say we will not union $b \rightarrow a$ directly

We can maintain sizes of the group, and we will always add smaller group to the bigger one.



union(0,1)

ops $\rightarrow \underline{\underline{1}}$

union(2,3)

union(1,0)

union(2,1)

What is the benefit

1

1

:

1

2

2

2

:

2

4

4

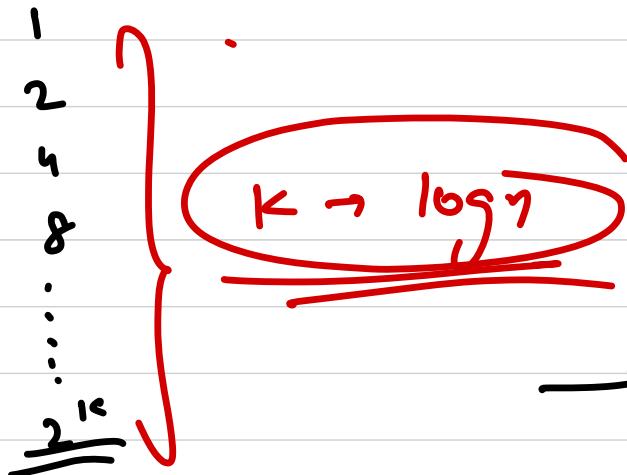
4

8

8

union
by
size

$$\cancel{\frac{1}{2} \times (1 \log n)} \rightarrow \cancel{O(\log n)}$$

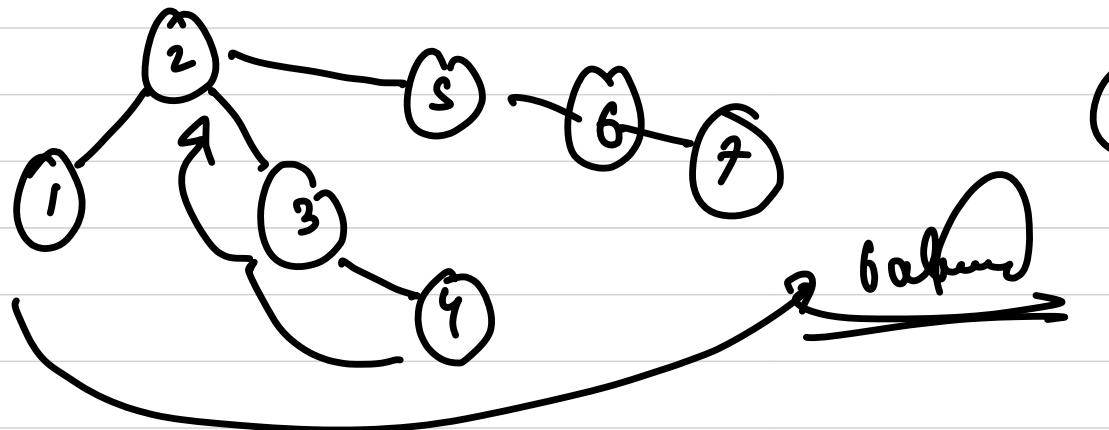
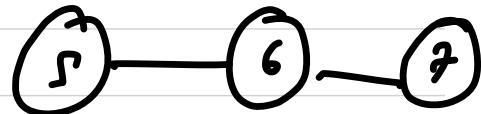


$$1 \times \frac{n}{2} + 2 \times \frac{n}{4} + 4 \times \frac{n}{8} \dots$$

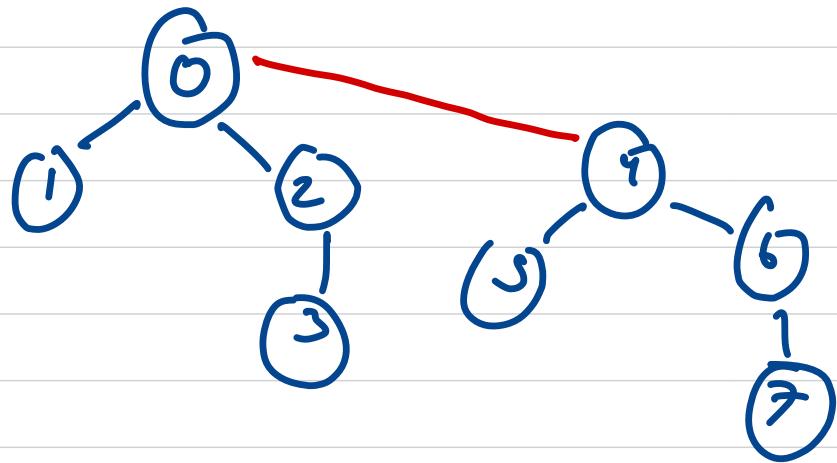
$$\xrightarrow{\text{union}} \underline{\log n}$$

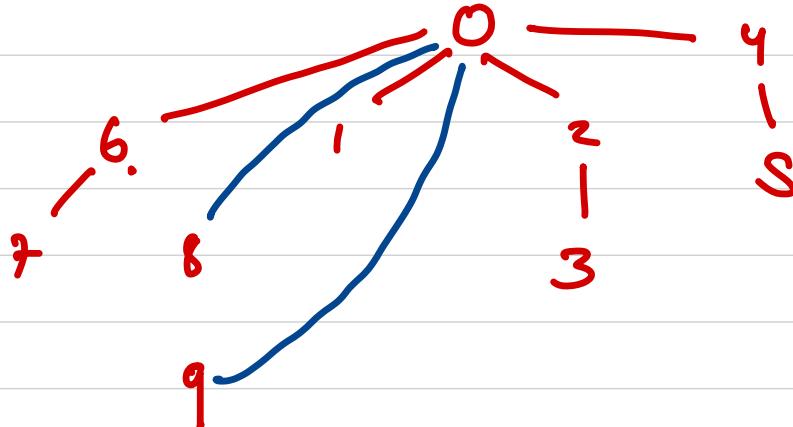
$$\text{get} \rightarrow \underline{\log n}$$

curve
obj
sw



semic





1
2
4
8
16
 \vdots
 2^k

$\log \gamma$

versus
alpha

$\log^2 \gamma$

path
compression

union by
rank

~~Size~~

$\log^* n \rightarrow$ extremely slow growing funcⁿ

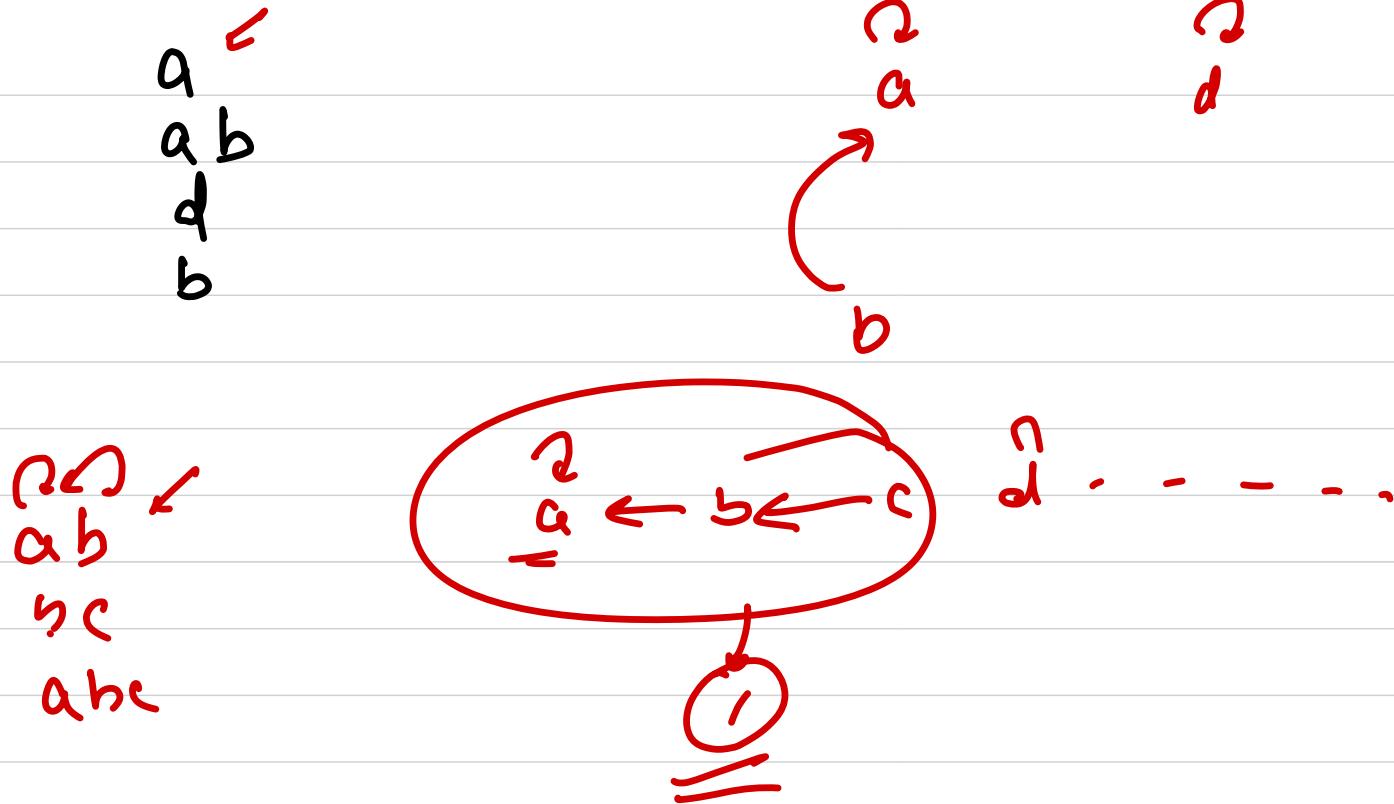
$\log^* n$ no. of steps we need to take

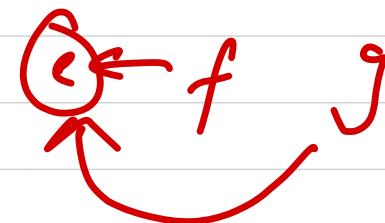
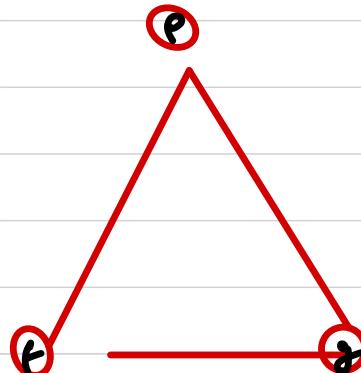
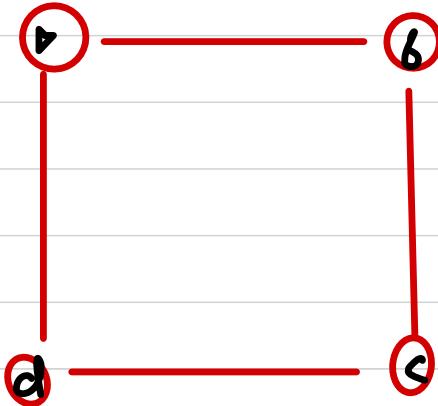
\log_2 on the value of n to make it

smaller than 1.

$$n \xrightarrow{\underline{\underline{2^{16}}}}$$

$$\log 2^{16} \rightarrow \log(16) \rightarrow \log(4) \log(2) \rightarrow 1$$





~~Q3~~ Given, n persons standing in a row at pos $\{1-n\}$. We can do 2 ops -

① $-x \rightarrow$ remove the person at pos x .

② $?x \rightarrow$ find the nearest person to the right of x that is still present (means not removed)

$$\boxed{n \leq 10^6}$$
$$q \leq 10^6$$

$$n \Rightarrow q = 3$$

$$\begin{array}{r} -2 \\ -3 \\ \hline ?1 \end{array}$$

$$\rightarrow \underline{\underline{1}}$$

~~1 X X 4 5~~



? 1

- 2

{ 1

union(x, y)



- x



union (x, x+1)

? x



get (x+1)

Q: Given an undirected graph, Some ops are given as -

remove a b → remove the edge between a and b vertex

get a b → tell if a & b are in same component.

after all the given no edges will be left,

$$v \leq 5 \times 10^4$$

$$e \leq 10^5$$

$$q \leq 1.5 \times 10^4$$



a — d
b — c

?
f
j

Offline
query

Answer a d

run d c

run f e

get b c — no

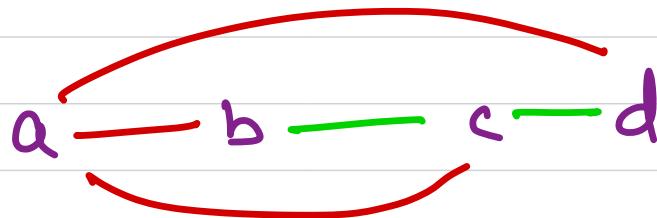
Answer c b
get b c — no

get a f — no

d. Given a graph with v nodes, & e edges , every edge has a color either red or green. Given src & dest , find the shortest path such that path starts from red edge & ends at green edge . & we can switch from a red to green edge only once.

src - a
dest - d

$$\begin{matrix} v \leq 10^3 \\ e \leq 10^3 \end{matrix}$$





Do BFS from both sides



first do bfs from src - dest with only red edges

second do bfs from dest - src with only green edges

src

o

b

c

d

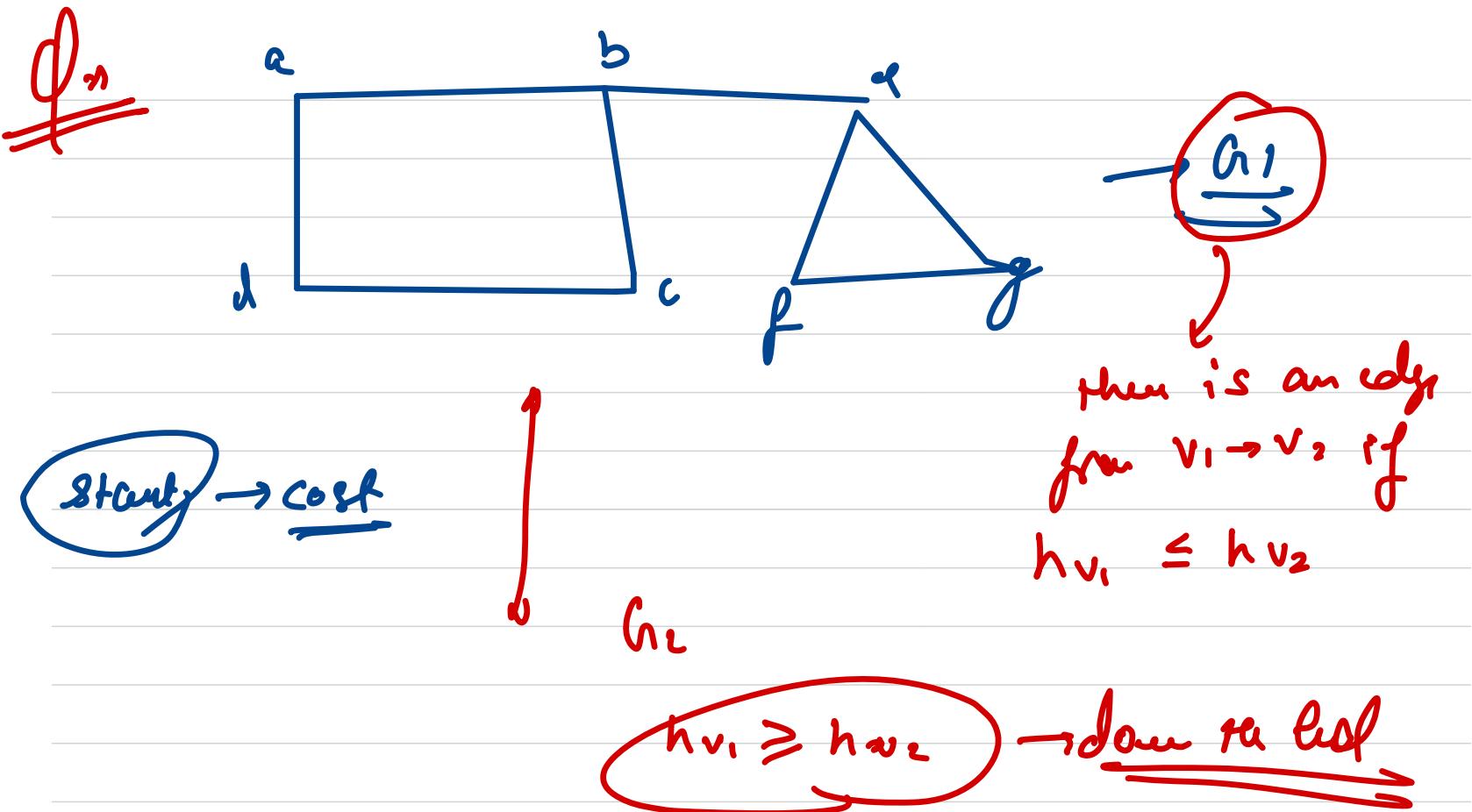
des

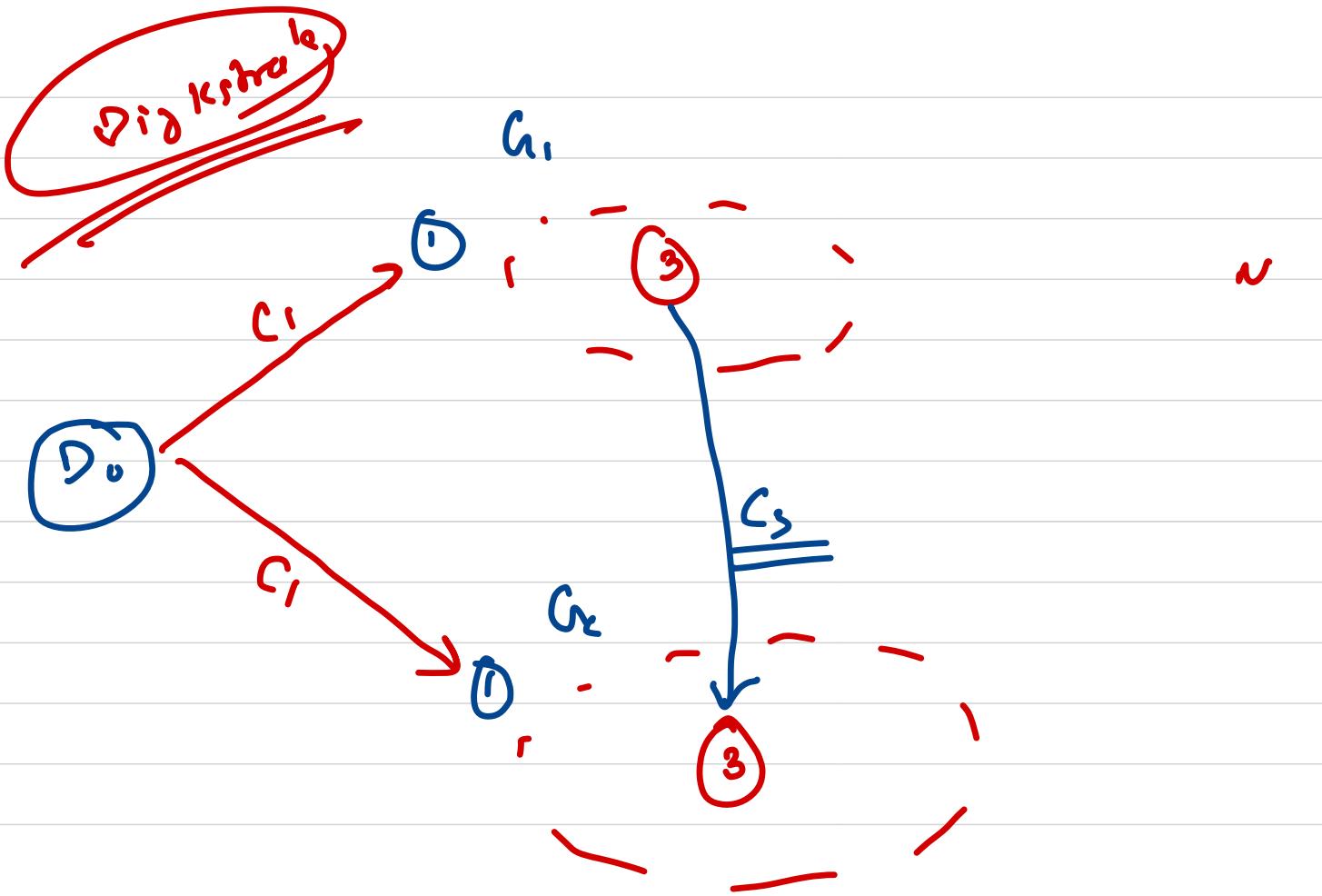
o

a — s — d

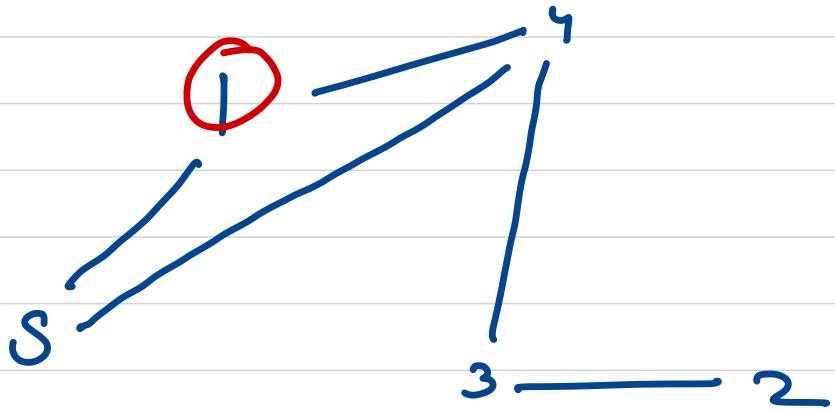
result

~~sd~~

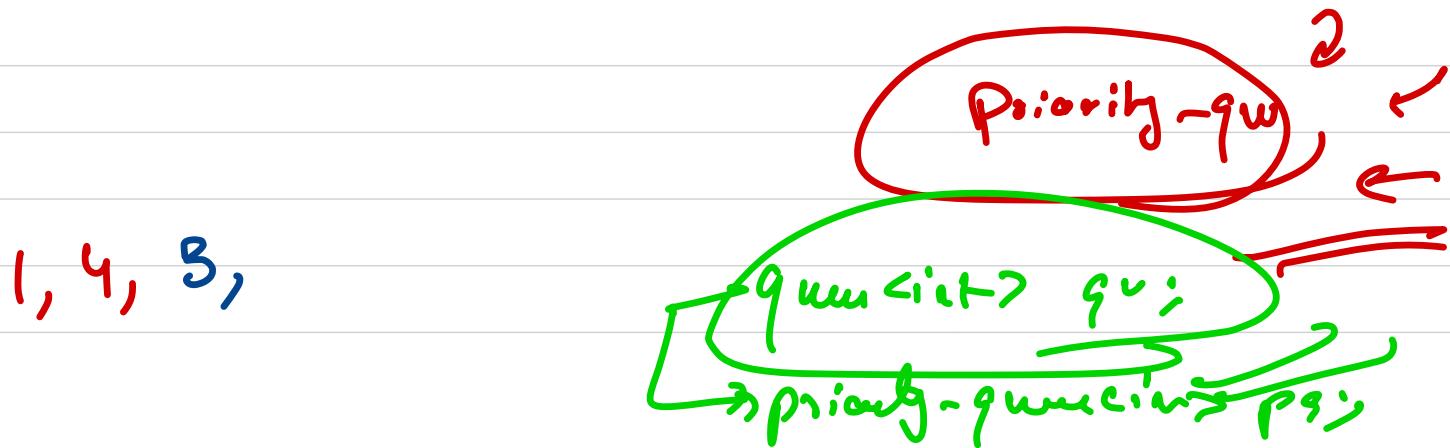




BFS



We have to take
child of sneeze
child first



[wrt, wrt, er, eff, fft]

w r t
w o f

w r t
e r

w r t
c t t

f r t
e t t

t f

c. w x

t → f c → e $\omega \rightarrow e$
 $\omega \rightarrow x$
 $x \rightarrow t$

w e r t f

[w_ot, w_rf, l_ev, e_tf, s_f⁺⁺]

w r f
w r f
t r f

t → f

w r f
w r f

w → s

w o t
e r

w → e

e r
e t t
e r
s f t t

O(n³)

insert

w . r

· f e ·

w e s t f

Topo

Qn Given a matrix M of binary values · find
if we can have a 1D array B such that

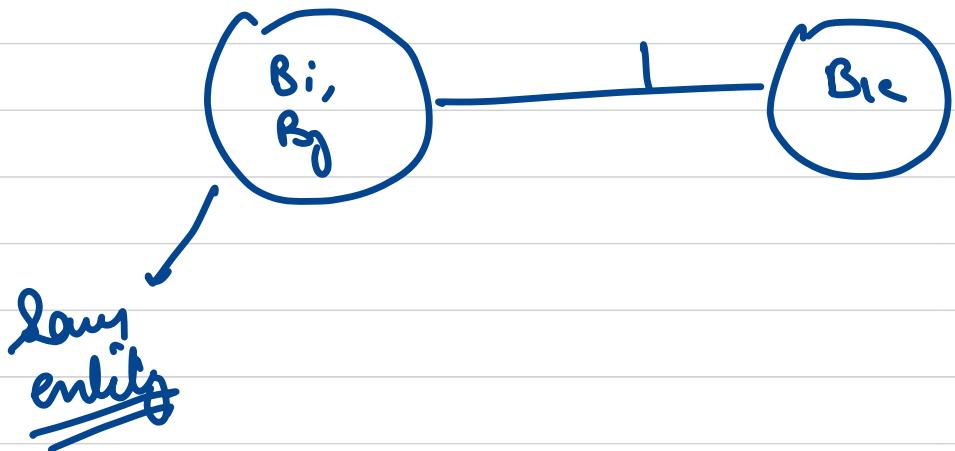
$$M_{:,j} = \{B_i - B_0\}$$

$$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

$$\boxed{3 \ 1 \ 4}$$

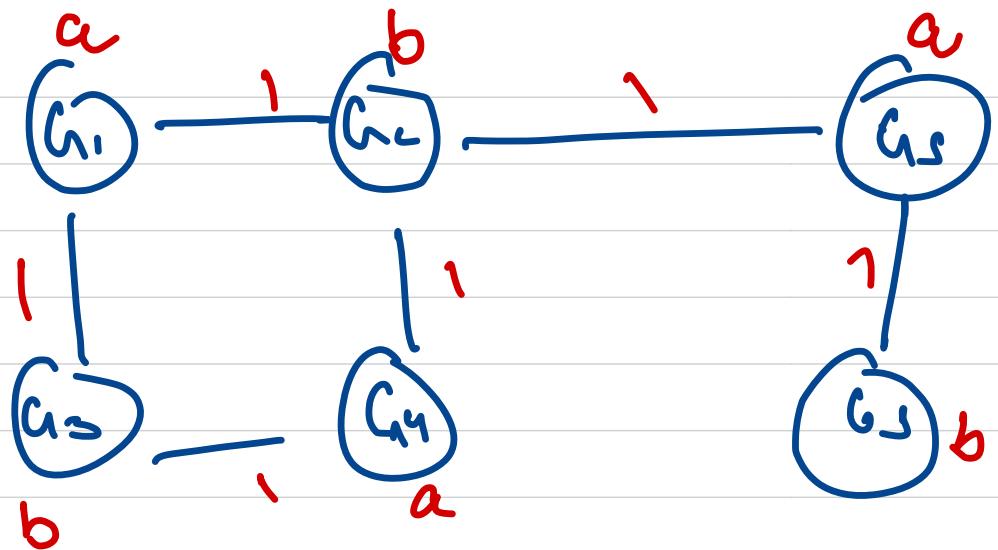
$$\begin{pmatrix} 1 \\ 1 \\ 0 \end{pmatrix}$$

if $M_{ij} = 0$ then $\underline{\underline{B_i = B_j}}$



$$\begin{aligned}M_{ik} &= 1 \\B_i &\neq B_{ik}\end{aligned}$$

$M_{ij} = 1$



Zioloczy

Qn Given a tree, select 2 nodes, we need to satisfy two conditions

- ① Subtrees rooted at the 2 selected nodes must be disjoint.
- ② Sum of node values \rightarrow max

$f(\text{node}, c)$

}

max value you

can get by consider

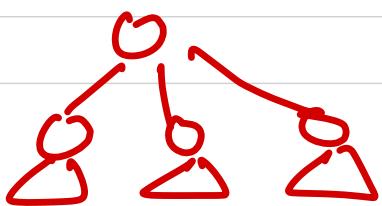
subtrees of node &

taking c no of nodes

$f(n, 1)$ $f(n, 2)$

↳ Sum of nodes

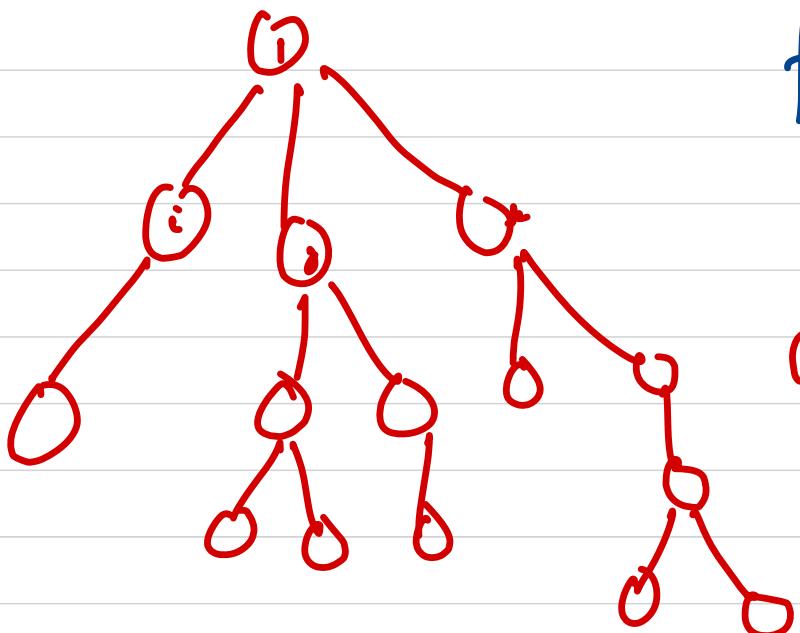
$$f(n, 1) = \max \left\{ \begin{array}{l} \text{sum of nodes}(n) \\ f(c_i, 1) \end{array} \right.$$

$$f(n, 2) = \max \left\{ \begin{array}{l} f(c_1, 2) \\ \text{or} \\ \text{calc } f(c_i, 1) \rightarrow \max, 2^{\text{nd}} \text{ min} \end{array} \right.$$


Base

$$f(u, 1) = au \quad u \rightarrow 1 \text{ by } u = 1$$

$$f(v, e) = -\infty$$

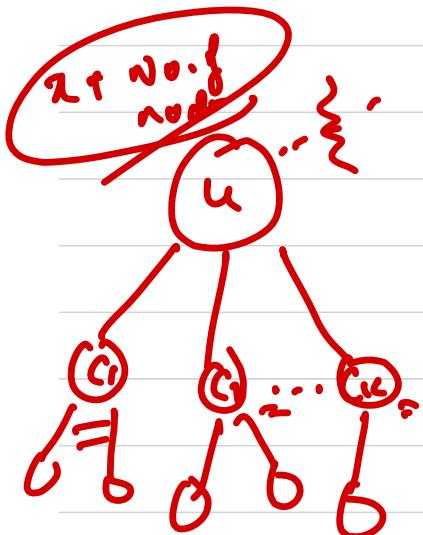


$f(i)$: sum of distance
 $(1, i) + i$

$g(i) =$

$f(n)$

$$g(u) = \sum_{i=1}^k g(c_i) + \text{nodes}(c_i)$$



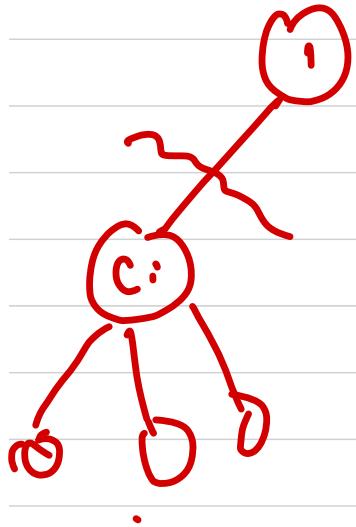
$$f(c_1) = g(c_1)$$

nodes(n) → count of nodes of subtree rooted at n

Rerooty

$g(c_i)$

$f(c_i) ??$



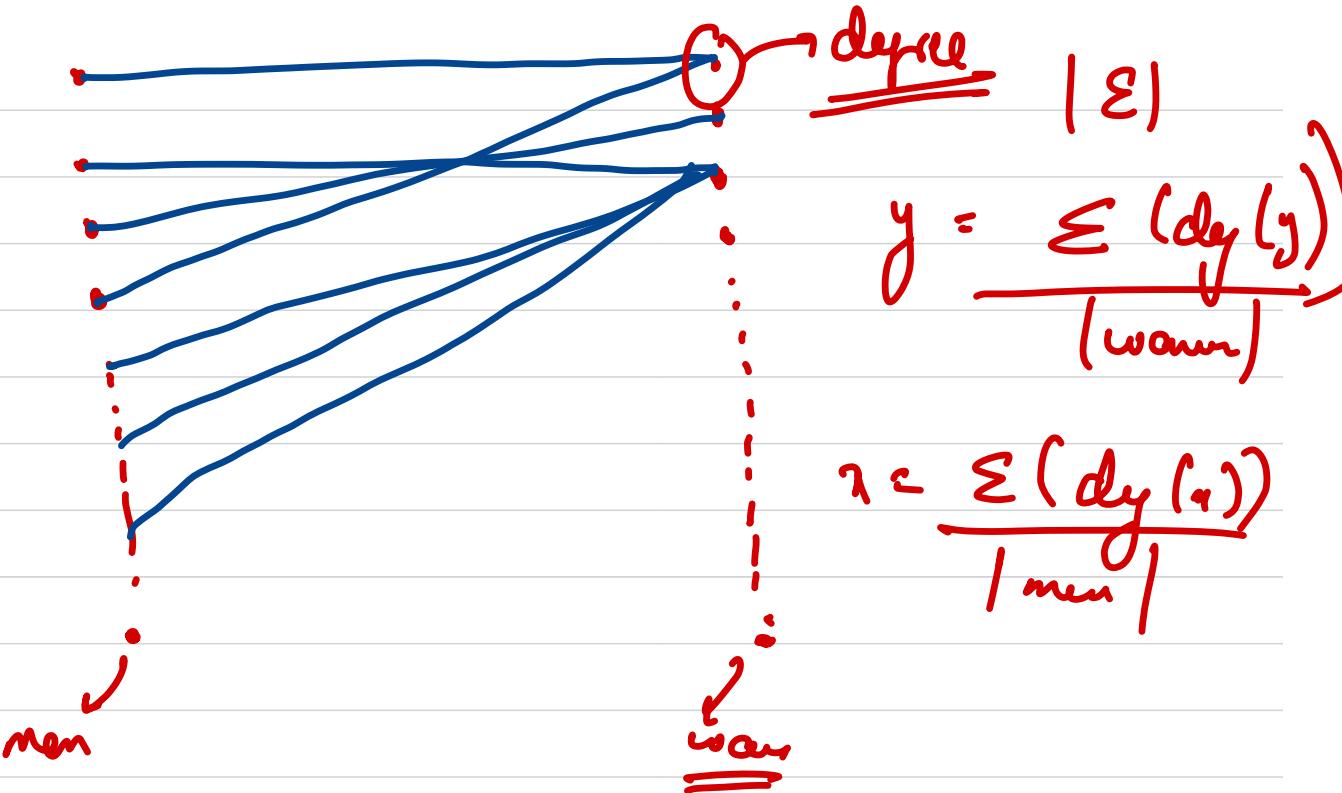
$p(l, c_i)$

$$b(l, c_i) = f(l) - \left(\overline{g(c_i)} + \overline{\text{node}(c_i)} \right)$$

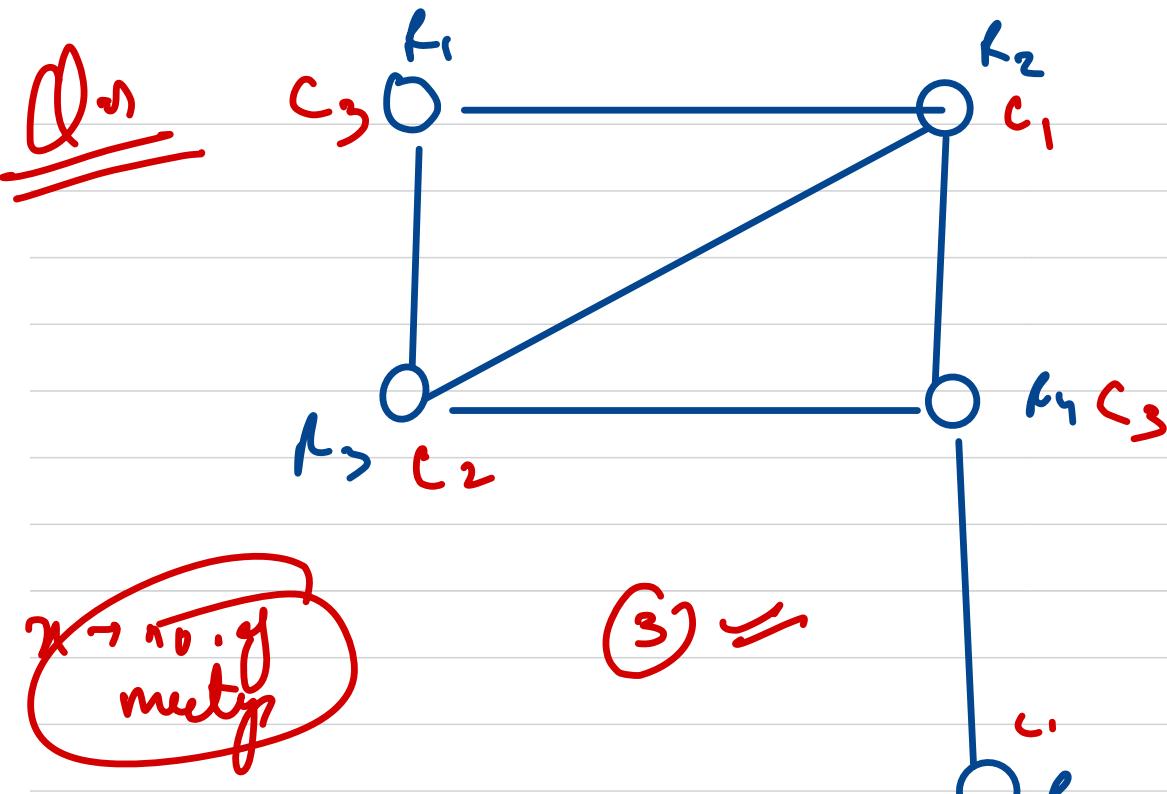
$$f(c_i) = g(c_i) + p(l, c_i) + \overline{n - \text{node}(c_i)}$$

Q. Let's say we have a gender dist i.e (male female count) which represent no. of male pop & no. of female pop. Say x is a no. that denotes avg no. of opp gender partners for men & y is denoting avg no. of opp gender partners for women. find the value of α/y

ans



$$x_y = \frac{\sum (\text{deg}(x))}{|\text{men}|} \times \frac{|\text{ways}|}{\sum (\text{deg}(y))} \Rightarrow \frac{|\text{ways}|}{|\text{men}|}$$



Range of meeting times

S-2 pm
2-9 pm
4-6 pm
11-1 am
5-5 am

graph color

~~edges → rule~~

meeting

rule → room

edge → new book

rooms can't share

parallel edges

wrap up the meetings

What is graph coloring problem?

Given a graph G , and k colors, assign a color to each node such that adjacent nodes get different colors.

Minimum no. of colors for which the above statement holds true is called as chromatic no. of graph $\chi(G)$

$\chi(G)$ ← How to do this ??

Theorem \rightarrow Given a graph G , if it has a acyclic,
then the base algo uses at most $d+1$ colors
for G .

I.S assume true for $P(n)$

B.C. \rightarrow $n=1$ \circlearrowleft $(0+1) \rightarrow 1$ color

Proof $\rightarrow P(n+1)$

$G \rightarrow (V, E) \rightarrow (n+1)$ nodes graph

$\hookrightarrow D$ is the max degree in \underline{G} .

$v_1 v_2 v_3 \dots \dots v_{n-1} v_n v_{n+1}$

Remove $\underline{v_{n+1}}$ & all corresponding edges we get

a graph $\underline{\underline{G'}} \leq D$

To color $\underline{\underline{G'}}$, acc to $P(n)$ we need $d+1$ colors
at most

dⁿ

Given a graph, check if it 2 colorable.

If the graph is bipartite or not ??

==

