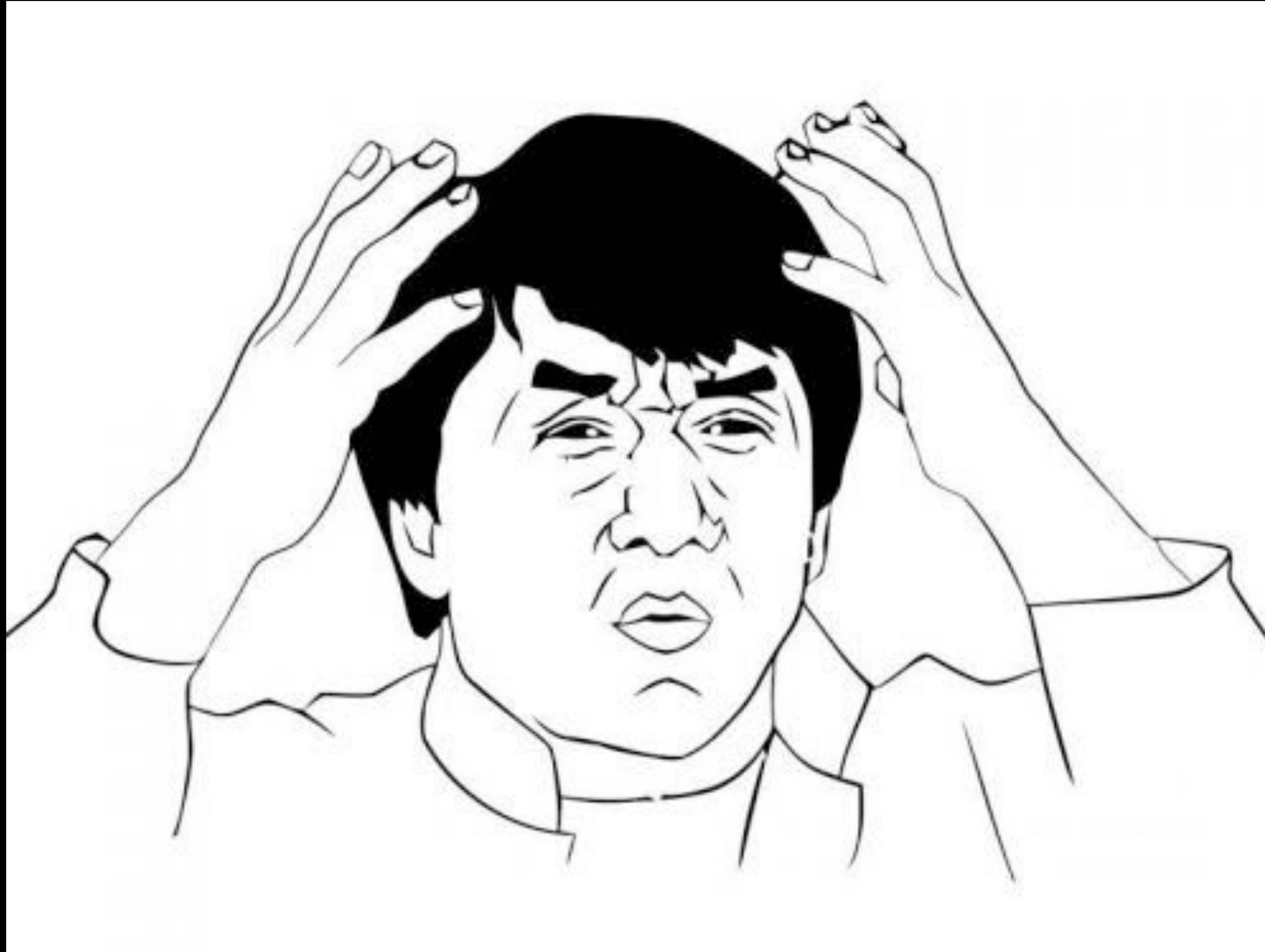


MySQL Lock



Deadlock & Slow query

MySQL Lock

- Isolation Level
 - Lock
 - MVCC

“isolation level is the setting that fine-tunes the balance between performance and reliability, consistency, and reproducibility of results when multiple transactions are making changes and performing queries at the same time.”

- <https://dev.mysql.com/doc/refman/8.0/en/innodb-transaction-isolation-levels.html>

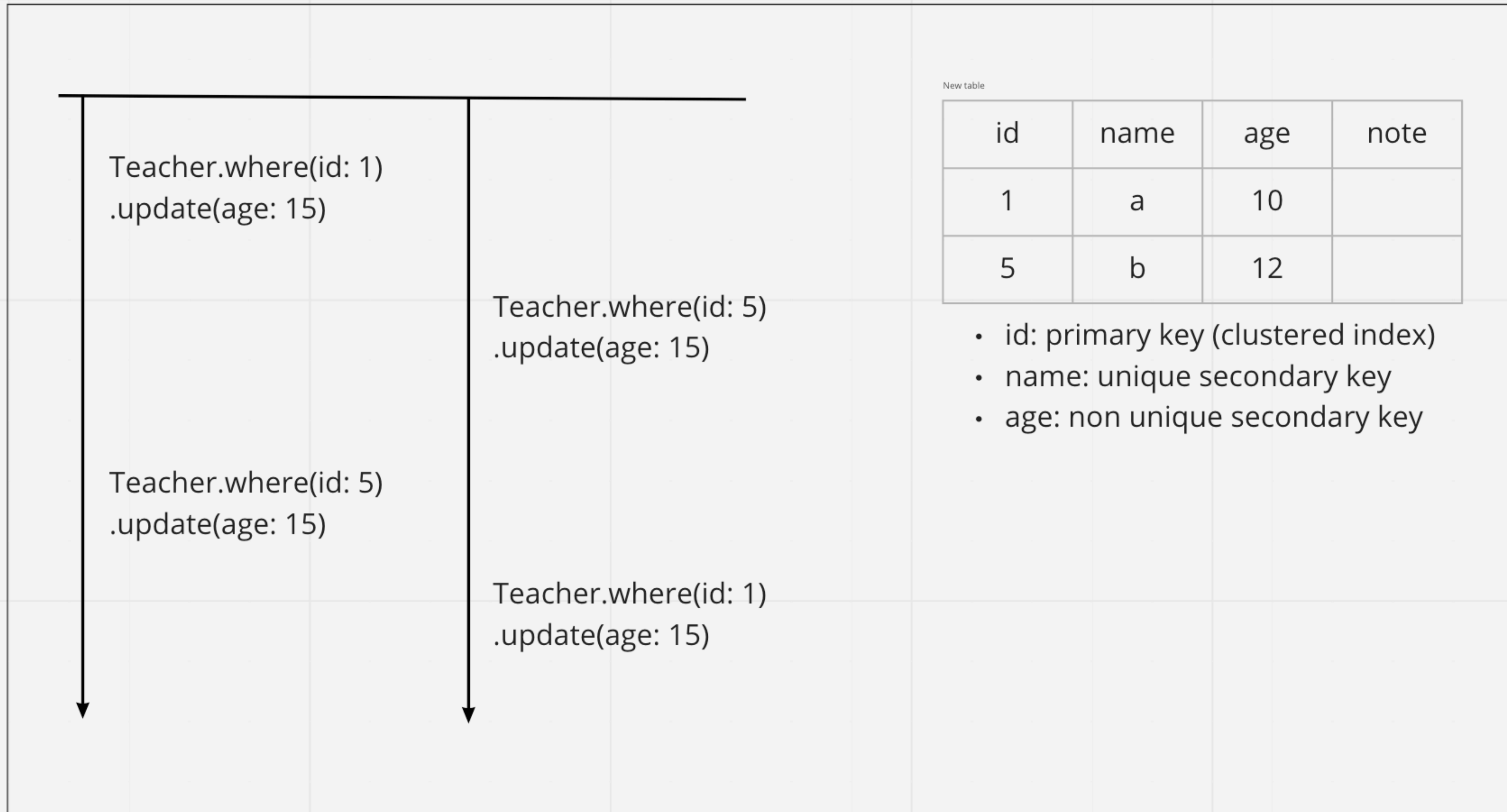
Agenda

- MySQL Lock
 - Lock and Index
 - Deadlock
 - Deadlock Error Message
 - Slow Query
- Real Case
- Conclusion

Precondition

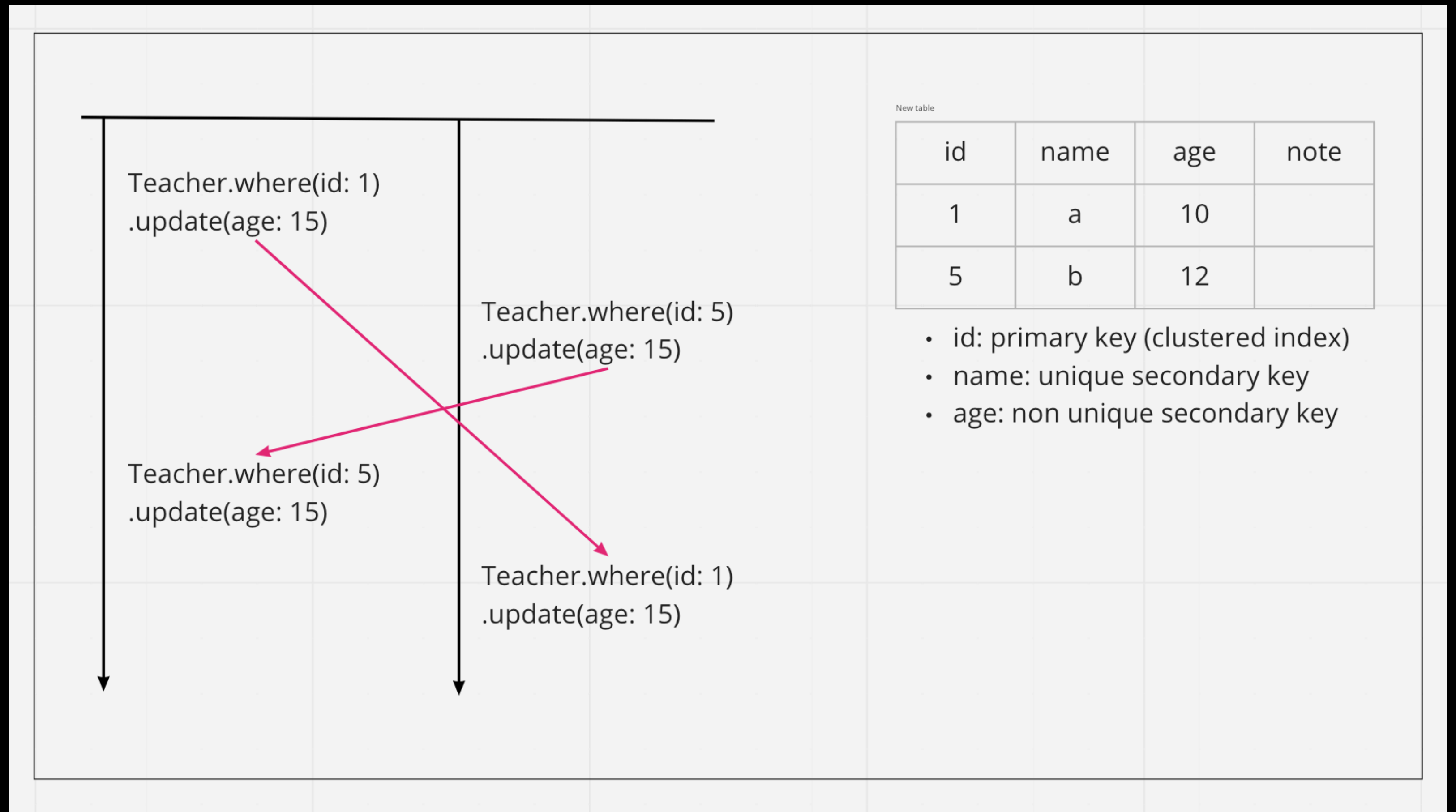
- MySQL 5.7
- Repeatable Read
- Table Teachers
 - Indice:
 - Primary key: id
 - Unique Index: name
 - Secondary Index: age

X Warm Up 1



Deadlock

- Mutual Exclusive
- Hold and Wait
- No Preemption
- Circular Wait



How to Debug DeadLock

- Deadlock Log
 - innodb deadlock detect: 是否偵測 deadlock
 - innodb print all deadlocks: 如果要全部 log 到 error.log 需要打開 flag
 - 否則只能看到最後一筆 deadlock
> SHOW ENGINE INNODB STATUS

How to Debug DeadLock

2022-04-21 10:36:45 0x40de538700

*** (1) TRANSACTION:

TRANSACTION 8115, ACTIVE 3 sec starting index read

mysql tables in use 1, locked 1

LOCK WAIT 3 lock struct(s), heap size 1136, 2 row lock(s), undo log entries 1

MySQL thread id 827, OS thread handle 278607652608, query id 11483 172.17.0.1 root updating

UPDATE `teachers` SET `teachers`.`age` = 6 WHERE `teachers`.`id` = 1

*** (1) WAITING FOR THIS LOCK TO BE GRANTED:

RECORD LOCKS space id 227 page no 3 n bits 72 index PRIMARY of table

`test`.`teachers` trx id 8115 **lock_mode X locks rec but not gap waiting**

Record lock, heap no 2 PHYSICAL RECORD: n_fields 6; compact format; info bits 0

0: len 8; hex 8000000000000001; asc ;;

1: len 6; hex 000000001fb2; asc ;;

2: len 7; hex 53000001cc295d; asc S)];;

3: len 3; hex 616161; asc aaa;;

4: len 4; hex 8000000f; asc ;;

5: SQL NULL;

*** (2) TRANSACTION:

TRANSACTION 8114, ACTIVE 3 sec starting index read

mysql tables in use 1, locked 1

3 lock struct(s), heap size 1136, 2 row lock(s), undo log entries 1

MySQL thread id 828, OS thread handle 278607922944, query id 11482 172.17.0.1 root updating

UPDATE `teachers` SET `teachers`.`age` = 10 WHERE `teachers`.`id` = 5

*** (2) HOLDS THE LOCK(S):

RECORD LOCKS space id 227 page no 3 n bits 72 index PRIMARY of table `test`.`teachers`

trx id 8114 lock_mode X locks rec but not gap

Record lock, heap no 2 PHYSICAL RECORD: n_fields 6; compact format; info bits 0

0: len 8; hex 8000000000000001; asc ;;

1: len 6; hex 000000001fb2; asc ;;

2: len 7; hex 53000001cc295d; asc S)];;

3: len 3; hex 616161; asc aaa;;

4: len 4; hex 8000000f; asc ;;

5: SQL NULL;

*** (2) WAITING FOR THIS LOCK TO BE GRANTED:

RECORD LOCKS space id 227 page no 3 n bits 72 index PRIMARY of table `test`.`teachers`

trx id 8114 **lock_mode X locks rec but not gap waiting**

Record lock, heap no 3 PHYSICAL RECORD: n_fields 6; compact format; info bits 0

0: len 8; hex 8000000000000005; asc ;;

1: len 6; hex 000000001fb3; asc ;;

2: len 7; hex 54000001ca2ac9; asc T *];;

3: len 3; hex 626262; asc bbb;;

4: len 4; hex 80000010; asc ;;

5: SQL NULL;

*** WE ROLL BACK TRANSACTION (2)

X Warm Up 2

```
Teacher.where('id >  
1').order(id: :asc)  
.update_all(age: 10)
```

```
Teacher.where('id >  
1').order(id: :desc)  
.update_all(age: 10)
```

New table

id	name	age	note
1	a	10	
5	b	12	

- id: primary key (clustered index)
- name: unique secondary key
- age: non unique secondary key

Row Lock - One By One

- Row locks are acquired one by one
 - ORDER BY
 - Index ordering (MySQL 8.0+)
- Follow index orders

“Transaction an UPDATE statement includes an ORDER BY clause, the **rows are updated in the order** specified by the clause.”

- <https://dev.mysql.com/doc/refman/5.7/en/update.html>

“A *key_part* specification can end with ASC or DESC. These keywords are permitted for future extensions for specifying ascending or descending index value storage. Currently, they are parsed but **ignored**”

- <https://dev.mysql.com/doc/refman/5.7/en/create-index.html>

* Warm Up 2 - 1

update teachers use
index (id_desc) set
age=5 where id > 1

update teachers use
index (id_asc) set
age=5 where id > 1

New table

id	name	age	note
1	a	10	
5	b	12	

- id: primary key (clustered index)
- name: unique secondary key
- age: non unique secondary key
- id_asc: 'create unique index id_asc on teachers (id ASC)'
- id_desc: 'create unique index id_desc on teachers (id DESC)'

X Warm Up 3

Teacher.where(id:
6).update_all(age:
10)

Teacher.create(id: 6)

Teacher.where(id:
10).update_all(age:
10)

Teacher.create(id: 10)

New table

id	name	age	note
1	a	10	
5	b	12	

- id: primary key (clustered index)
- name: unique secondary key
- age: non unique secondary key

○ Warm Up 3 - 1

Teacher.where(id:
6).update_all(age:
10)

Teacher.create(id: 6)

Teacher.where(id:
10).update_all(age:
10)

Teacher.create(id: 10)

New table

id	name	age	note
1	a	10	
8	b	12	

- id: primary key (clustered index)
- name: unique secondary key
- age: non unique secondary key

MySQL Lock

- Lock Mode
 - Share (S) / Exclusive (X)
- Lock Type
 - Record Lock (Table / Row)
 - Intention Lock (Table Level)
 - Auto-Incr Lock
 - Gap Lock (Prevent Insert)
 - Insert Intention Lock
 - Next-Key Lock

	id	name	age	note
(Negative infinity, 1]	1	Aaa	10	
(1, 5]	5	Bob	12	
(5, positive infinity]				

MySQL Lock

LATEST DETECTED DEADLOCK

2022-04-22 10:36:54 277165582080

*** (1) TRANSACTION:

TRANSACTION 1872, ACTIVE 3 sec inserting

mysql tables in use 1, locked 1

LOCK WAIT 3 lock struct(s), heap size 1128, 2 row lock(s)

MySQL thread id 25, OS thread handle 278780622592, query id 203

172.17.0.1 root update

INSERT INTO `teachers` (`id`) VALUES (6)

*** (1) HOLDS THE LOCK(S):

RECORD LOCKS space id 3 page no 4 n bits 72 index PRIMARY of table `test`.`teachers` trx id 1872 **lock_mode X**

Record lock, heap no 1 PHYSICAL RECORD: n_fields 1; compact format; info bits 0

0: len 8; hex 73757072656d756d; asc supremum;;

*** (1) WAITING FOR THIS LOCK TO BE GRANTED:

RECORD LOCKS space id 3 page no 4 n bits 72 index PRIMARY of table `test`.`teachers` trx id 1872 **lock_mode X insert intention waiting**

Record lock, heap no 1 PHYSICAL RECORD: n_fields 1; compact format; info bits 0

0: len 8; hex 73757072656d756d; asc **supremum**;;

*** (2) TRANSACTION:

TRANSACTION 1873, ACTIVE 3 sec inserting

mysql tables in use 1, locked 1

LOCK WAIT 3 lock struct(s), heap size 1128, 2 row lock(s)

MySQL thread id 26, OS thread handle 278783792896, query id 204

172.17.0.1 root update

INSERT INTO `teachers` (`id`) VALUES (10)

*** (2) HOLDS THE LOCK(S):

RECORD LOCKS space id 3 page no 4 n bits 72 index PRIMARY of table `test`.`teachers` trx id 1873 **lock_mode X**

Record lock, heap no 1 PHYSICAL RECORD: n_fields 1; compact format; info bits 0

0: len 8; hex 73757072656d756d; asc supremum;;

*** (2) WAITING FOR THIS LOCK TO BE GRANTED:

RECORD LOCKS space id 3 page no 4 n bits 72 index PRIMARY of table `test`.`teachers` trx id 1873 lock_mode **X insert intention**

waiting

Record lock, heap no 1 PHYSICAL RECORD: n_fields 1; compact format; info bits 0

0: len 8; hex 73757072656d756d; asc **supremum**;;

*** WE ROLL BACK TRANSACTION (2)

Range Update

update teachers use
index (PRIMARY) set
note="123" where id
>= 3 and age = 40

block!

Teacher.where(id:
3).update_all(age:
15)

Teacher.create(id: 7,
name: 'blocked!')

New table

id	name	age	note
2	a	10	
3	b	12	
4	c	13	
100	d	15	

- id: primary key (clustered index)
- name: unique secondary key
- age: non unique secondary key

Range Update

- Lock every rows even it is not satisfied where condition

“When using the default REPEATABLE READ isolation level, the first UPDATE acquires an x-lock on each row that it reads and **does not release any of them:**”

- https://dev.mysql.com/doc/refman/8.0/en/innodb-transaction-isolation-levels.html#isolevel_repeatable-read

MySQL Lock - Clustered Index

- Hit => Row Lock
- Not Hit => Next Key Lock
- Range Search => Lock every row it read

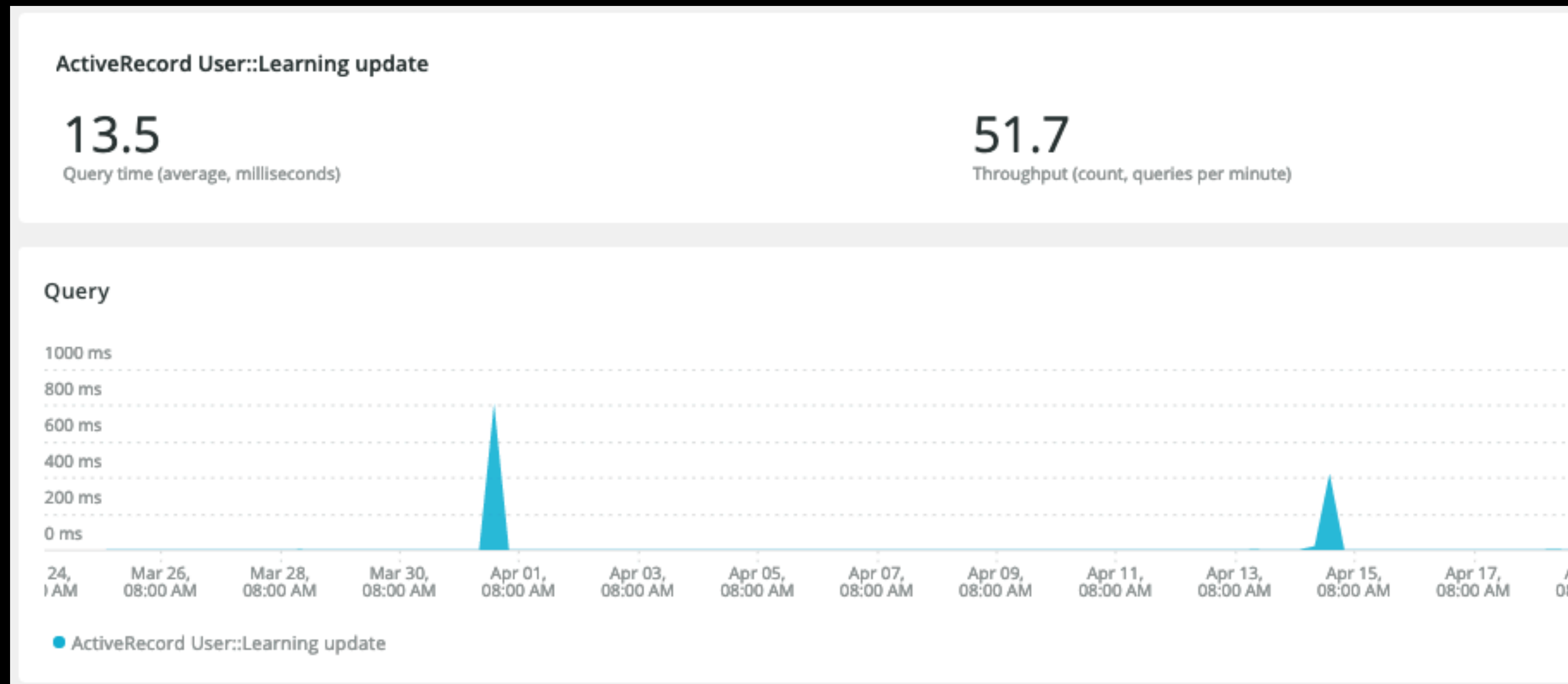
Real Case 1

- Hourly::CheckAppointmentRequestExpired
 - 預約 Appointment 時，會需要老師 confirm，如果超過時間沒有 confirm 則更新 Appointment 狀態為 expired
- Metrics: Appointment Update
- Metrics: Appointment Create

```
# Query_time: 15.574314 Lock_time: 0.000079 Rows_sent: 0 Rows_examined: 7484706
SET timestamp=1651940578;
UPDATE `appointments` SET `appointments`.`state` = 10, `appointments`.`state_change_time` =
'2022-05-07 16:22:43' WHERE `appointments`.`deleted_at` IS NULL AND ((state = 0) and
(confirm_expired_time <= '2022-05-07 16:20:02.597431'));
```

Real Case 2

- Monthly::DividedRatio
- metrics reference



```
User::Learning.update_all('intrested_ratio = intrested_ratio / 2')  
Locale::LanguageTagSearch.update_all('count = count / 2')
```

Not Unique Secondary Index

Teacher.where(age:
10).update_all(note:
'special')

Teacher.create(nam
e: 'a', age: 9)

New table

id	name	age	note
2	b	10	
3	c	12	
4	d	13	

- id: primary key (clustered index)
- name: unique secondary key
- age: non unique secondary key

Secondary Index

- Type
 - Unique => Almost Like Clustered Index
 - Not Unique => require gap lock even hit
- Keep primary key reference

“For other search conditions, and for non-unique indexes, InnoDB locks the index range scanned, **using gap locks** or next-key locks to block insertions by other sessions into the gaps covered by the range.”

- <https://dev.mysql.com/doc/refman/8.0/en/innodb-locks-set.html>

How To Solve

- 批次更新 (find_in_batch)
- 篩選出 clustered id 再更新，避免多餘的 gap lock

```
Ids = Collection.where(conditions).pluck(:id)
```

```
Collection.where(id: ids).update_all(...)
```

- 改成 read-committed
- 檢查 ORM 產生的 SQL

Why not set Read-Committed as default?

- Controversial

“In general I think good practice is to **use READ COMMITTED isolation mode as default** and change to REPEATABLE READ for those applications or transactions which require it.”

- <https://www.percona.com/blog/2015/01/14/mysql-performance-implications-of-innodb-isolation-modes/>

“THE GENERAL RULE could be the following :

- if your queries and transactions are short : use rather the **default REPEATABLE-READ** mode!
 - if your queries are long and **reading a lot of data which** are likely to be modified by other transactions in parallel : then **use the READ-COMMITTED mode**
- ”

- <https://dev.mysql.com/blog-archive/performance-impact-of-innodb-transaction-isolation-modes-in-mysql-5-7/>