ECE 350
Real-time
Operating
Systems

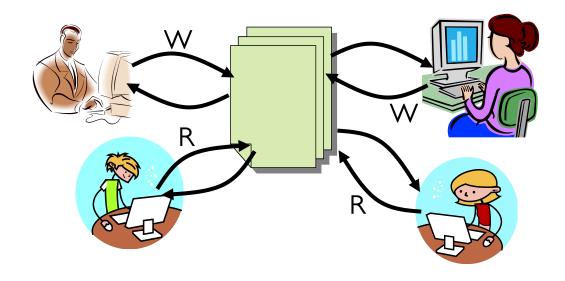


Tutorial Reader/Writer Lock

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Readers/Writers Lock



- Motivation: consider shared database with two classes of users
 - Readers: never modify database
 - Writers: read and modify database
- Database can have many readers at the same time
- But there can be only one writer active at a time

Properties of Readers/Writers Lock

- Common variant of mutual exclusion
 - One writer at a time, if no readers
 - Many readers, if no writer

Thread 2 Thread I	Writer	Reader
Writer	NO!	NO!
Reader	NO!	OK!

- Correctness constraints
 - Readers can read when no writers
 - Writers can read/write when no readers or writers
 - Only one thread manipulates state of the lock at a time

Readers/Writers Lock Class

```
class ReadWriteLock {
   private:
      Lock lock;
                             // needed to change state vars
      CV okToRead
                             // CV for readers
      CV okToWrite;
                             // CV for writers
      int AW = 0;
                             // # of active writers
      int AR = 0;
                             // # of active readers
      int WW = 0;
                             // # of waiting writers
                             // # of waiting readers
      int WR = 0;
   public:
      void acquireRL();
      void releaseRL();
      void acquireWL();
      void releaseWL();
```

Readers/Writers Lock Design Pattern

```
read() {
    lock.acquireRL();

    // Read shared state

    lock.releaseRL();
}

write() {
    lock.acquireWL();

// Read/write shared state

lock.releaseWL();
}
```

Readers/Writers Lock Implementation

```
acquireRL() {
   lock.acquire();
                                // Need lock to change state vars
  while (AW + WW > 0) { // Is it safe to read?
                               // No! add to # of waiting readers
     WR++;
     okToRead.wait(&lock);
                                // Wait on condition variable
                                   No longer waiting
     WR - -;
  AR++:
                                 // Now we are active again
   lock.release();
                          Why release lock
                               here?
releaseRL() {
   lock.acquire();
                                // No longer active
  AR - - :
  if (AR == 0 \&\& WW > 0)
                             // If no active reader,
     okToWrite.signal();
                              // wake up waiting writer
   lock.release();
```

Lock is acquired to change internal state of R/W lock

Readers/Writers Lock Implementation (cont.)

```
acquireWL() {
   lock.acquire();
  while (AW + AR > 0) {
                                // is it safe to write?
                                // No! add to # of waiting writers
     WW++;
                                   Wait on condition variable
     okToWrite.wait(&lock);
                                   No longer waiting
     WW--:
  AW++;
                                // Now we are active again
   lock.release();
releaseWL() {
   lock.acquire();
                                // No longer active
  AW--;
  if (WW > 0) {
                                // Give priority to writers
     okToWrite.signal();
                              // Wake up a waiting writer
   } else if (WR > 0) {
                                // If there are waiting readers,
                                    wake them all up
     okToRead.broadcast();
   lock.release();
```

• Consider following sequence of arrivals: R1, R2, W1, R3

```
read() {
   lock.acquire();
   while (AW + WW > 0) {
      WR++;
      okToRead.wait(&lock);
      WR - -:
   AR++;
   lock.release();
   // Read
   lock.acquire();
   AR - - :
   if (AR == 0 \&\& WW > 0)
      okToWrite.signal();
   lock.release();
}
```

```
write() {
   lock.acquire();
   while (AW + AR > 0) {
      WW++;
      okToWrite.wait(&lock);
      WW--:
   AW++;
   lock.release();
   // Read and Write
   lock.acquire();
   AW--:
   if (WW > 0) {
      okToWrite.signal();
   } else if (WR > 0) {
      okToRead.broadcast();
   lock.release();
```

• RI comes - AR = 0, WR = 0, AW = 0, WW = 0

```
read()
    lock.acquire();
   while (AW + WW > 0) {
       WR++;
       okToRead.wait(&lock);
       WR - -:
   AR++;
   lock.release();
   // Read
   lock.acquire();
   AR - - ;
   if (AR == 0 \&\& WW > 0)
       okToWrite.signal();
    lock.release();
```

• RI comes - AR = 0, WR = 0, AW = 0, WW = 0

```
read() {
    lock.acquire();
    white (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR--;
    }
    AR++;
    lock.release();

    // Read

    lock.acquire();
    AR--;
    if (AR == 0 && WW > 0)
        okToWrite.signal();
    lock.release();
}
```

• RI comes — AR = 0, WR = 0, AW = 0, WW = 0

```
read() {
    lock.acquire();
    while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR--;
    }
    AR++;
    lock.release();

    // Read

    lock.acquire();
    AR--;
    if (AR == 0 && WW > 0)
        okToWrite.signal();
    lock.release();
}
```

```
• RI comes — AR = 1, WR = 0, AW = 0, WW = 0
  read() {
     lock.acquire();
     while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR - -:
     AR++;
     lock.release();
     // Read
     lock.acquire();
     AR--;
     if (AR == 0 \&\& WW > 0)
        okToWrite.signal();
     lock.release();
```

```
• RI comes - AR = 1, WR = 0, AW = 0, WW = 0
  read() {
     lock.acquire();
     while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR - -:
     AR++:
     lock.release();
     // Read
     lock.acquire();
     AR--;
     if (AR == 0 \&\& WW > 0)
        okToWrite.signal();
     lock.release();
```

```
• RI comes — AR = 1, WR = 0, AW = 0, WW = 0
  read() {
     lock.acquire();
     while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR - -:
     AR++;
     lock.release();
        Read
     lock.acquire();
     AR - - ;
     if (AR == 0 \&\& WW > 0)
        okToWrite.signal();
     lock.release();
```

• R2 comes - AR = 1, WR = 0, AW = 0, WW = 0

```
read()
    lock.acquire();
   while (AW + WW > 0) {
       WR++;
       okToRead.wait(&lock);
       WR - -:
   AR++;
   lock.release();
   // Read
   lock.acquire();
   AR - - ;
   if (AR == 0 \&\& WW > 0)
       okToWrite.signal();
    lock.release();
```

• R2 comes — AR = 1, WR = 0, AW = 0, WW = 0

```
read() {
   lock.acquire();
   while (AW + WW > 0) {
      WR++;
      okToRead.wait(&lock);
      WR - -:
   AR++;
   lock.release();
   // Read
   lock.acquire();
   AR--;
   if (AR == 0 \&\& WW > 0)
      okToWrite.signal();
   lock.release();
```

• R2 comes — AR = 1, WR = 0, AW = 0, WW = 0 read() { lock.acquire(): while $(\overline{AW} + \overline{WW} > 0)$ WR++; okToRead.wait(&lock); WR - -: AR++; lock.release(); // Read lock.acquire(); AR--; if (AR == 0 && WW > 0)okToWrite.signal(); lock.release();

```
• R2 comes — AR = ^{2}, WR = ^{0}, AW = ^{0}, WW = ^{0}
  read() {
     lock.acquire();
     while (AW + WW > 0) {
         WR++;
         okToRead.wait(&lock);
         WR - -:
     AR++;
      lock.release();
     // Read
     lock.acquire();
     AR--;
     if (AR == 0 \&\& WW > 0)
         okToWrite.signal();
      lock.release();
```

```
• R2 comes — AR = 2, WR = 0, AW = 0, WW = 0
  read() {
     lock.acquire();
     while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR - -:
     AR++:
     lock.release();
     // Read
     lock.acquire();
     AR--;
     if (AR == 0 \&\& WW > 0)
        okToWrite.signal();
     lock.release();
```

```
• R2 comes — AR = 2, WR = 0, AW = 0, WW = 0
  read() {
     lock.acquire();
     while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR - -:
     AR++;
     lock.release();
        Read
     lock.acquire();
     AR--;
     if (AR == 0 \&\& WW > 0)
        okToWrite.signal();
     lock.release():
```

Assume readers take a while to access database Situation: Locks released, only AR is non-zero

• WI comes — AR = 2, WR = 0, AW = 0, WW = 0

```
write()
   lock.acquire();
   while (AW + AR > 0) {
      WW++;
      okToWrite.wait(&lock);
      WW--:
   AW++;
   lock.release();
   // Read and Write
   lock.acquire();
   AW--;
   if (WW > 0) {
      okToWrite.signal();
   } else if (WR > 0) {
      okToRead.broadcast();
   lock.release();
```

• WI comes — AR = 2, WR = 0, AW = 0, WW = 0

```
write() {
   lock.acquire();
   while (AW + AR > 0) {
      WW++;
      okToWrite.wait(&lock);
      WW--:
   AW++;
   lock.release();
   // Read and Write
   lock.acquire();
   AW--;
   if (WW > 0) {
      okToWrite.signal();
   } else if (WR > 0) {
      okToRead.broadcast();
   lock.release();
```

```
• WI comes — AR = 2, WR = 0, AW = 0, WW = 0
  write() {
     lock.acquire():
     while (AW + AR > 0)
        WW++;
        okToWrite.wait(&lock);
        WW--:
     AW++;
     lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
     lock.release();
```

```
• WI comes – AR = 2, WR = 0, AW = 0, WW = \mathbf{1}
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++;
        okToWrite.wait(&lock);
        WW - - :
     AW++;
     lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
      lock.release();
```

```
• WI comes — AR = 2, WR = 0, AW = 0, WW = 1
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++:
        okToWrite.wait(&lock);
        WW--:
     AW++;
     lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
     lock.release();
```

WI cannot start because of readers, so it releases lock and goes to sleep

• R3 comes — AR = 2, WR = 0, AW = 0, WW = 1

```
read()
    lock.acquire();
   while (AW + WW > 0) {
       WR++;
       okToRead.wait(&lock);
       WR - -:
   AR++;
   lock.release();
   // Read
   lock.acquire();
   AR - - ;
   if (AR == 0 \&\& WW > 0)
       okToWrite.signal();
    lock.release();
```

• R3 comes — AR = 2, WR = 0, AW = 0, WW = 1

```
read() {
    lock.acquire();
    while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR--;
    }
    AR++;
    lock.release();

    // Read

    lock.acquire();
    AR--;
    if (AR == 0 && WW > 0)
        okToWrite.signal();
    lock.release();
}
```

• R3 comes — AR = 2, WR = 0, AW = 0, WW = 1 read() { lock.acquire(): while $(\overline{AW} + \overline{WW} > 0)$ WR++; okToRead.wait(&lock); WR - -: AR++; lock.release(); // Read lock.acquire(); AR--; if (AR == 0 && WW > 0)okToWrite.signal(); lock.release();

```
• R3 comes — AR = 2, WR = \frac{1}{1}, AW = \frac{1}{1}, AW = \frac{1}{1}
  read() {
      lock.acquire();
      while (AW + WW > 0) {
         WR++;
         okToRead.wait(&lock);
         WR - -:
      AR++;
      lock.release();
      // Read
      lock.acquire();
      AR--;
      if (AR == 0 \&\& WW > 0)
         okToWrite.signal();
      lock.release();
```

• R3 comes — AR = 2, WR = 1, AW = 0, WW = 1

```
read() {
   lock.acquire();
   while (AW + WW > 0) {
      WR++:
      okToRead.wait(&lock);
      WR - - ;
   AR++;
   lock.release();
   // Read
   lock.acquire();
   AR--;
   if (AR == 0 \&\& WW > 0)
      okToWrite.signal();
   lock.release();
```

Status:

- R1 and R2 still reading
- WI and R3 waiting on okToWrite and okToRead, respectively

• R2 is done reading -AR = 2, WR = 1, AW = 0, WW = 1read() { lock.acquire(); while (AW + WW > 0) { WR++; okToRead.wait(&lock); WR - -: AR++; lock.release(); // Read lock.acquire(); AR - - ; if (AR == 0 && WW > 0)okToWrite.signal(); lock.release();

• R2 is done reading -AR = 1, WR = 1, AW = 0, WW = 1 read() { lock.acquire(); while (AW + WW > 0) { WR++; okToRead.wait(&lock); WR - -: AR++; lock.release(); // Read lock.acquire(): AR - -; if (AR == 0 && WW > 0) okToWrite.signal(); lock.release();

• R2 is done reading -AR = 1, WR = 1, AW = 0, WW = 1 read() { lock.acquire(); while (AW + WW > 0) { WR++; okToRead.wait(&lock); WR - -: AR++; lock.release(); // Read lock.acquire(); AR - - : if (AR == 0 && WW > 0)okloWrite.signal();

lock.release();

• R2 is done reading -AR = 1, WR = 1, AW = 0, WW = 1 read() { lock.acquire(); while (AW + WW > 0) { WR++; okToRead.wait(&lock); WR - -: AR++; lock.release(); // Read lock.acquire(); AR--; if (AR == 0 && WW > 0)

okToWrite.signal():

lock.release();

• RI is done reading -AR = 1, WR = 1, AW = 0, WW = 1read() { lock.acquire(); while (AW + WW > 0) { WR++; okToRead.wait(&lock); WR - -: AR++; lock.release(); // Read lock.acquire(); AR - - ; if (AR == 0 && WW > 0)okToWrite.signal(); lock.release();

• RI is done reading -AR = 0, WR = 1, AW = 0, WW = 1 read() { lock.acquire(); while (AW + WW > 0) { WR++; okToRead.wait(&lock); WR - -: AR++; lock.release(); // Read lock.acquire(): AR - -; if (AR == 0 && WW > 0) okToWrite.signal(); lock.release();

• RI is done reading -AR = 0, WR = 1, AW = 0, WW = 1

```
read() {
    lock.acquire();
    while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR--;
    }
    AR++;
    lock.release();

    // Read

    lock.acquire();
    AR--:
    if (AR == 0 && WW > 0)
        okToWrite.signal();
    lock.release();
}
```

• RI is done reading -AR = 0, WR = 1, AW = 0, WW = 1read() { lock.acquire(); while (AW + WW > 0) { WR++; okToRead.wait(&lock); WR - -: AR++; lock.release(); All readers are finished, R1 signals waiting writer - note, R3 is still waiting // Read lock.acquire(); AR - - : if (AR == 0 && WW > 0)okToWrite.signal(); lock.release();

• RI is done reading -AR = 0, WR = 1, AW = 0, WW = 1

```
read() {
    lock.acquire();
    while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR--;
    }
    AR++;
    lock.release();

    // Read

    lock.acquire();
    AR--;
    if (AR == 0 && WW > 0)
        okToWrite.signal():
    lock.release();
}
```

```
• WI gets a signal - AR = 0, WR = 1, AW = 0, WW = 1
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++:
                                    WI gets signal from RI
        okToWrite.wait(&lock);
        WW--:
     AW++;
     lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
     ĺock.release();
```

```
• WI gets a signal - AR = 0, WR = 1, AW = 0, WW = 0
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++;
        okToWrite.wait(&lock):
        WW--;
     AW++;
     lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
     lock.release();
```

```
• WI gets a signal - AR = 0, WR = 1, AW = 1, WW = 0
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++;
        okToWrite.wait(&lock);
        WW--:
     AW++;
      lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
      lock.release();
```

```
• WI gets a signal - AR = 0, WR = 1, AW = 1, WW = 0
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++;
        okToWrite.wait(&lock);
        WW--:
     AW++:
     lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
     lock.release();
```

```
• WI gets a signal - AR = 0, WR = 1, AW = 1, WW = 0
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++;
        okToWrite.wait(&lock);
        WW--:
     AW++;
     lock.release();
        Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
     lock.release();
```

```
• WI is done — AR = 0, WR = 1, AW = 1, WW = 0
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++;
        okToWrite.wait(&lock);
        WW--:
     AW++;
     lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
     lock.release();
```

```
• WI is done - AR = 0, WR = 1, AW = 0, WW = 0
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++;
        okToWrite.wait(&lock);
        WW--:
     AW++;
     lock.release();
     // Read and Write
     lock.acouire():
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
     lock.release();
```

```
• WI is done - AR = 0, WR = 1, AW = 0, WW = 0
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++;
        okToWrite.wait(&lock);
        WW--:
     AW++;
     lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0)
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
     lock.release();
```

```
• WI is done - AR = 0, WR = 1, AW = 0, WW = 0
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++;
        okToWrite.wait(&lock);
        WW--:
     AW++;
     lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal():
       else if (WR > 0)
        okToRead.broadcast();
     lock.release();
```

```
• WI is done - AR = 0, WR = 1, AW = 0, WW = 0
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++;
        okToWrite.wait(&lock);
        WW--:
     AW++;
     lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
                                      No waiting writer, so
        okToRead.broadcast();
                                         only signal R3
     lock.release();
```

```
• WI is done - AR = 0, WR = 1, AW = 0, WW = 0
  write() {
     lock.acquire();
     while (AW + AR > 0) {
        WW++;
        okToWrite.wait(&lock);
        WW--:
     AW++;
     lock.release();
     // Read and Write
     lock.acquire();
     AW--;
     if (WW > 0) {
        okToWrite.signal();
     } else if (WR > 0) {
        okToRead.broadcast();
     lock.release();
```

```
• R3 gets a signal - AR = 0, WR = 1, AW = 0, WW = 0
  read() {
     lock.acquire();
     while (AW + WW > 0) {
        WR++:
        okToRead.wait(&lock);
                                        R3 gets signal from W3
        WR--:
     AR++;
     lock.release();
     // Read
     lock.acquire();
     AR--;
     if (AR == 0 \&\& WW > 0)
        okToWrite.signal();
     lock.release();
```

• R3 gets a signal - AR = 0, WR = 0, AW = 0, WW = 0read() { lock.acquire(); while (AW + WW > 0) { WR++; okToRead.wait(&lock): WR--; AR++: lock.release(); // Read lock.acquire(); AR--; if (AR == 0 && WW > 0)okToWrite.signal(); lock.release();

```
• R3 gets a signal - AR = 1, WR = 0, AW = 0, WW = 0
  read() {
     lock.acquire();
     while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR - -:
     AR++;
     lock.release();
     // Read
     lock.acquire();
     AR--;
     if (AR == 0 \&\& WW > 0)
        okToWrite.signal();
     lock.release();
```

• R3 finishes — AR = 0, WR = 0, AW = 0, WW = 0 read() { lock.acquire(); while (AW + WW > 0) { WR++; okToRead.wait(&lock); WR - -: AR++; lock.release(); // Read lock.acquire(); AR--; if (AR == 0 && WW > 0)okToWrite.signal(): lock.release();

```
read() {
   lock.acquire();
   while (AW + WW > 0) {
      WR++:
      okToRead.wait(&lock);
      WR - -:
   AR++:
   lock.release();
                            What if we
                         remove this line?
   // Read
   lock.acquire();
   AR - -:
   if (AR == 0 \&\& WW > 0)
      oklowrite.signal();
   lock.release();
```

It works but it's inefficient, writer wakes up and goes to sleep again when it's not save to write

```
write() {
   lock.acquire();
   while (AW + AR > 0) {
      WW++:
      okToWrite.wait(&lock);
      WW--:
   AW++:
   lock.release();
   // Read and Write
   lock.acquire();
   AW--:
   if (WW > 0) {
      okToWrite.signal();
   } else if (WR > 0) {
      okToRead.broadcast();
   lock.release();
```

```
read() {
   lock.acquire();
   while (AW + WW > 0) {
      WR++:
      okToRead.wait(&lock);
      WR - -:
   AR++:
                         What if we turn
   lock.release();
                            signal() to
   // Read
                           broadcast()?
   lock.acquire();
   AR - - :
   if (AR == 0 \&\& WW > 0)
      okToWrite.broadcast();
   lock.release();
```

```
write() {
   lock.acquire();
   while (AW + AR > 0) {
      WW++:
      okToWrite.wait(&lock);
      WW--:
   AW++:
   lock.release();
   // Read and Write
   lock.acquire();
   AW--:
   if (WW > 0) {
      okToWrite.signal();
   } else if (WR > 0) {
      okToRead.broadcast();
   lock.release();
```

It works but it's inefficient to wake up all writers only for one to becomes active

```
read() {
    lock.acquire();
    while (AW + WW > 0) {
        WR++;
        okToContinue.wait(&lock);
        WR--;
    }
    AR++;
    lock.release();

    // Read

    lock.acquire();
    AR--;
    if (AR == 0 && WW > 0)
        okToContinue.signal();
    lock.release();
}
```

What if we turn **okToWrite** and **okToRead** into **okContinue**?

Signal could be delivered to wrong thread (reader) and get waisted!

```
write() {
   lock.acquire();
   while (AW + AR > 0) {
      WW++:
      okToContinue.wait(&lock);
      WW--:
   AW++:
   lock.release();
   // Read and Write
   lock.acquire();
   AW--:
   if (WW > 0) {
      okToContinue.signal();
   } else if (WR > 0) {
      okToContinue.broadcast();
   lock.release();
```

```
read() {
   lock.acquire();
   while (AW + WW > 0) {
      WR++:
      okToContinue.wait(&lock);
      WR - -:
   AR++:
   lock.release(); Does changing signal()
                   to broadcast() solve the
   // Read
                           problem?
   lock.acquire();
   AR - - :
   if (AR == 0 \&\& WW > 0)
      okToContinue.broadcast();
   lock.release();
```

```
write() {
   lock.acquire();
   while (AW + AR > 0) {
      WW++:
      okToContinue.wait(&lock);
      WW--:
   AW++:
   lock.release();
   // Read and Write
   lock.acquire();
   AW--:
   if (WW > 0) {
      okToContinue.broadcast();
   } else if (WR > 0) {
      okToContinue.broadcast();
   lock.release();
```

Yes, but it's inefficient to wake up all threads for only one to becomes active

```
read() {
    lock.acquire();
    while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR--;
    }
    AR++;
    lock.release();

    // Read

    lock.acquire();
    AR--;
    if (AR == 0 && WW > 0)
        okToWrite.signal();
    lock.release();
}
```

Can readers starve?

Yes: writers take priority

```
write() {
   lock.acquire();
   while (AW + AR > 0) {
      WW++:
      okToWrite.wait(&lock);
      WW--:
   AW++:
   lock.release();
   // Read and Write
   lock.acquire();
   AW--:
   if (WW > 0) {
      okToWrite.signal();
   } else if (WR > 0) {
      okToRead.broadcast();
   lock.release();
```

```
read() {
    lock.acquire();
    while (AW + WW > 0) {
        WR++;
        okToRead.wait(&lock);
        WR--;
    }
    AR++;
    lock.release();

    // Read

    lock.acquire();
    AR--;
    if (AR == 0 && WW > 0)
        okToWrite.signal();
    lock.release();
}
```

Can writers starve?

Yes: a waiting writer may not be able to proceed if another writer slips in between signal and wakeup

```
write() {
   lock.acquire();
   while (AW + AR > 0) {
      WW++:
      okToWrite.wait(&lock);
      WW--:
   AW++:
   lock.release();
   // Read and Write
   lock.acquire();
   AW--:
   if (WW > 0) {
      okToWrite.signal();
   } else if (WR > 0) {
      okToRead.broadcast();
   lock.release();
```

Readers/Writers Lock Without Writer Starvation (Take 1)

```
check also for waiting writers

acquireWL() {
    lock.acquire();
    while (AW + AR + WW > 0) {
        WW++;
        okToWrite.wait(&lock);
        WW--;
    }
    AW++;
    lock.release();
}
```

- Does this work?
 - No! If there **WW** is more than zero, then no waiting writer can successfully proceed

Readers/Writers Lock Without Writer Starvation (Take 2)

Idea: keep track of writers' waiting order, allow writer with longest waiting time to proceed

Does this work?

```
numWriters = 0;
                                                 releaseWL() {
                  No, Signal can wake up a
                                                    lock.acquire();
nextToGo = 1;
                  wrong writer and get waisted!
                                                    AW--:
acquireWL() {
                                                    nextToGo++:
                                                    if (WW > 0) {
   lock.acquire();
   mvPos = numWriters++;
                                                       okToWrite.signal();
                                                      else if (WR > 0) {
   while (AW + AR > 0 | |
            myPos > nextToGo) {
                                                       okToRead.broadcast();
      WW++;
                                                    lock.release();
      okToWrite.wait(&lock);
      WW--:
                     Inefficient solution is to change
   AW++:
                       signal() to broadcast()
   lock.release();
```

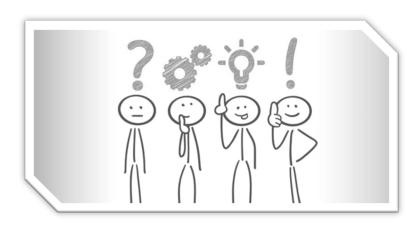
Readers/Writers Lock Without Writer Starvation (Take 3)

Idea for efficient solution: have separate CV for each waiting writer and put CV's in ordered queue

```
numWriters = 0;
nextToGo = 1;
acquireWL() {
   lock.acquire();
   myPos = numWriters++;
   myCV = new CV();
   queue.enqueue(myCV);
   while (AW + AR > 0 | |
            mvPos > nextToGo) {
      WW++:
      myCV.wait(&lock);
      WW--:
   AW++;
   queue.dequeue();
   lock.release();
```

```
releaseWL() {
   lock.acquire();
   AW--:
   nextToGo++;
   if (WW > 0) {
      queue.head().signal();
   } else if (WR > 0) {
      okToRead.broadcast();
   lock.release();
}
releaseRL() {
   lock.acquire();
   AR - - ;
   if (AR == 0 \&\& WW > 0)
      queue.head().signal();
   lock.release();
```

Questions?



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