# Algorithm Library

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## 1 多项式

#### 1.1 FFT - tourist

```
/* copy from tourist */
1
2
   namespace FFT {
       typedef double dbl;
3
4
        struct num {
5
6
            dbl x, y;
7
            num() \{ x = y = 0; \}
            num(dbl x, dbl y) : x(x), y(y) \{ \}
8
9
        };
10
        inline num operator+(num a, num b) { return num(a.x + b.x, a.y + b.y); }
11
        inline num operator—(num \ a, num \ b) \{ return num(a.x - b.x, a.y - b.y); \}
12
        inline num operator*(num a, num b) { return num(a.x * b.x - a.y * b.y, a
13
           .x * b.y + a.y * b.x);
        inline num conj(num a) { return num(a.x, -a.y); }
14
15
16
       int base = 1;
17
        vector < num > roots = \{ \{0, 0\}, \{1, 0\} \};
        vector < int > rev = \{ 0, 1 \};
18
19
        const dbl PI = a cosl(-1.0);
20
21
       void ensure_base(int nbase) {
22
23
            if (nbase <= base) {
24
                return;
25
            }
            rev.resize(1 << nbase);
26
27
            for (int i = 0; i < (1 << nbase); i++) {
                rev[i] = (rev[i >> 1] >> 1) + ((i & 1) << (nbase - 1));
28
29
            roots.resize(1 << nbase);
30
31
            while (base < nbase) {
                dbl \ angle = 2 * PI / (1 << (base + 1));
32
                for (int i = 1 \ll (base - 1); i \ll (1 \ll base); i++) {
33
                     roots[i \ll 1] = roots[i];
34
                     dbl \ angle_i = angle * (2 * i + 1 - (1 << base));
35
36
                     roots[(i \ll 1) + 1] = num(cos(angle_i), sin(angle_i));
37
38
                base++;
39
```

```
40
        }
41
42
        void fft (vector <num>& a, int n = -1) {
            if (n == -1) {
43
                 n = a.size();
44
45
            assert((n & (n - 1)) == 0);
46
            int zeros = __builtin_ctz(n);
47
            ensure_base(zeros);
48
            int shift = base - zeros;
49
50
            for (int i = 0; i < n; i++) {
                 if (i < (rev[i] >> shift)) {
51
                     swap(a[i], a[rev[i] >> shift]);
52
53
                 }
54
            for (int k = 1; k < n; k <<= 1) {
55
                 for (int i = 0; i < n; i += 2 * k) {
56
                     for (int j = 0; j < k; j++) {
57
                         num z = a[i + j + k] * roots[j + k];
58
                          a[i + j + k] = a[i + j] - z;
59
                         a[i + j] = a[i + j] + z;
60
61
                     }
                }
62
            }
63
64
        }
65
66
        vector < num> fa, fb;
67
        vector < long long > multiply (vector < int > & a, vector < int > & b) {
68
            int need = a.size() + b.size() - 1;
69
70
            int nbase = 1;
71
            while ((1 \ll \text{nbase}) < \text{need}) \text{ nbase}++;
72
            ensure_base(nbase);
            int sz = 1 \ll nbase;
73
            if (sz > (int) fa. size()) 
74
75
                 fa.resize(sz);
76
            for (int i = 0; i < sz; i++) {
77
                 int x = (i < (int)a.size() ? a[i] : 0);
78
                 int y = (i < (int)b.size() ? b[i] : 0);
79
80
                 fa[i] = num(x, y);
81
82
            fft (fa, sz);
```

```
83
             num r(0, -0.25 / (sz >> 1));
             for (int i = 0; i \le (sz >> 1); i++) {
84
85
                 int j = (sz - i) & (sz - 1);
                 num z = (fa[j] * fa[j] - conj(fa[i] * fa[i])) * r;
86
                 if (i != j) {
87
                      fa[j] = (fa[i] * fa[i] - conj(fa[j] * fa[j])) * r;
88
89
                 fa[i] = z;
90
91
             }
             for (int i = 0; i < (sz >> 1); i++) {
92
93
                 num A0 = (fa[i] + fa[i + (sz >> 1)]) * num(0.5, 0);
                 num A1 = (fa[i] - fa[i + (sz >> 1)]) * num(0.5, 0) * roots[(sz)]
94
                    >> 1) + i];
95
                 fa[i] = A0 + A1 * num(0, 1);
96
             fft(fa, sz \gg 1);
97
             vector < long long > res (need);
98
             for (int i = 0; i < need; i++) {
99
                 if (i \% 2 == 0) {
100
101
                      res[i] = fa[i >> 1].x + 0.5;
102
                 } else {
                      res[i] = fa[i >> 1].y + 0.5;
103
104
                 }
105
             }
106
             return res;
107
        }
108
         vector<long long> square(const vector<int>& a) {
109
             int need = a.size() + a.size() - 1;
110
111
             int nbase = 1;
             while ((1 \ll \text{nbase}) < \text{need}) \text{ nbase}++;
112
113
             ensure_base(nbase);
114
             int sz = 1 \ll nbase;
             if ((sz >> 1) > (int) fa. size()) {
115
                 fa.resize(sz \gg 1);
116
117
             for (int i = 0; i < (sz >> 1); i++) {
118
119
                 int x = (2 * i < (int) a. size() ? a[2 * i] : 0);
                 int y = (2 * i + 1 < (int) a. size() ? a[2 * i + 1] : 0);
120
121
                 fa[i] = num(x, y);
122
123
             fft(fa, sz \gg 1);
             num r(1.0 / (sz >> 1), 0.0);
124
```

```
125
             for (int i = 0; i \le (sz >> 2); i++) {
                  int j = ((sz >> 1) - i) & ((sz >> 1) - 1);
126
127
                 num fe = (fa[i] + conj(fa[j])) * num(0.5, 0);
                 num fo = (fa[i] - conj(fa[j])) * num(0, -0.5);
128
                 num \ aux = fe * fe + fo * fo * roots[(sz >> 1) + i] * roots[(sz >> 1) + i]
129
                     >> 1) + i];
                 num tmp = fe * fo;
130
131
                  fa[i] = r * (conj(aux) + num(0, 2) * conj(tmp));
                  fa[j] = r * (aux + num(0, 2) * tmp);
132
             }
133
134
             fft(fa, sz \gg 1);
135
             vector < long long > res (need);
             for (int i = 0; i < need; i++) {
136
137
                  if (i \% 2 == 0) {
                      res[i] = fa[i >> 1].x + 0.5;
138
139
                 } else {
                      res[i] = fa[i >> 1].y + 0.5;
140
141
142
143
             return res;
        }
144
145
         vector<int> multiply_mod(vector<int>& a, vector<int>& b, int m, int eq =
146
             0) {}
147
             int need = a.size() + b.size() - 1;
148
             int nbase = 0;
149
             while ((1 \ll \text{nbase}) < \text{need}) \text{ nbase}++;
             ensure base(nbase);
150
             int sz = 1 \ll nbase;
151
             if (sz > (int) fa.size())  {
152
                  fa.resize(sz);
153
154
155
             for (int i = 0; i < (int)a.size(); i++) {
                 int x = (a[i] \% m + m) \% m;
156
                  fa[i] = num(x & ((1 << 15) - 1), x >> 15);
157
158
             }
             fill(fa.begin() + a.size(), fa.begin() + sz, num{0, 0});
159
160
             fft (fa, sz);
             if (sz > (int) fb.size()) 
161
                  fb.resize(sz);
162
163
164
             if (eq) {
                 copy(fa.begin(), fa.begin() + sz, fb.begin());
165
```

```
166
             } else {
167
                 for (int i = 0; i < (int)b.size(); i++) {
168
                     int x = (b[i] \% m + m) \% m;
                     fb [i] = num(x & ((1 << 15) - 1), x >> 15);
169
                 }
170
171
                 fill(fb.begin() + b.size(), fb.begin() + sz, num{0, 0});
172
                 fft (fb, sz);
173
174
             dbl ratio = 0.25 / sz;
             num r2(0, -1);
175
176
             num r3(ratio, 0);
177
             num r4(0, -ratio);
             num r5(0, 1);
178
179
             for (int i = 0; i \le (sz >> 1); i++) {
                 int j = (sz - i) & (sz - 1);
180
181
                 num a1 = (fa[i] + conj(fa[j]));
                 num a2 = (fa[i] - conj(fa[j])) * r2;
182
                 num b1 = (fb[i] + conj(fb[j])) * r3;
183
                 num b2 = (fb[i] - conj(fb[j])) * r4;
184
185
                 if (i != j) {
186
                     num c1 = (fa[j] + conj(fa[i]));
187
                     num c2 = (fa[j] - conj(fa[i])) * r2;
                     num d1 = (fb[j] + conj(fb[i])) * r3;
188
                     num d2 = (fb[j] - conj(fb[i])) * r4;
189
190
                     fa[i] = c1 * d1 + c2 * d2 * r5;
191
                     fb[i] = c1 * d2 + c2 * d1;
192
                 fa[j] = a1 * b1 + a2 * b2 * r5;
193
                 fb[j] = a1 * b2 + a2 * b1;
194
195
196
             fft (fa, sz);
197
             fft (fb, sz);
198
             vector < int > res (need);
             for (int i = 0; i < need; i++) {
199
                 long long aa = fa[i].x + 0.5;
200
201
                 long long bb = fb[i].x + 0.5;
202
                 long long cc = fa[i].y + 0.5;
203
                 res[i] = (aa + ((bb \% m) \ll 15) + ((cc \% m) \ll 30)) \% m;
204
205
             return res;
206
        }
207
208
        vector < int > square_mod (vector < int > & a, int m) {
```

```
209 | return multiply_mod(a, a, m, 1);
210 | };
```

#### 1.2 形式幂级数

### 2 数论

#### 2.1 简单的防爆模板

```
1
   namespace SimpleMod {
2
        constexpr int MOD = (int)1e9 + 7;
3
        inline int norm(long long a) { return (a % MOD + MOD) % MOD; }
        inline int add(int a, int b) { return a + b >= MOD ? a + b - MOD : a + b
4
5
        inline int sub(int a, int b) { return a - b < 0 ? a - b + MOD : a - b; }
        inline int mul(int a, int b) { return (int)((long long)a * b % MOD); }
6
7
        inline int powmod(int a, long long b) {
            int res = 1;
8
            while (b > 0) {
9
10
                 if (b \& 1) res = mul(res, a);
                a = mul(a, a);
11
12
                b >>= 1;
13
            return res;
14
15
16
        inline int inv(int a) {
            a \%= MOD;
17
            \quad \textbf{if} \quad (a < 0) \quad a \ +\!\!= MOD;
18
            int b = MOD, u = 0, v = 1;
19
20
            while (a) {
21
                int t = b / a;
22
                b = t * a; swap(a, b);
                u = t * v; swap(u, v);
23
24
            }
            assert(b == 1);
25
            if (u < 0) u += MOD;
26
            return u;
27
28
        }
29
```

#### 2.2 筛法

#### 2.2.1 线性素数筛

```
vector < bool > is Prime; // true 表示非素数 false 表示是素数
1
   vector<int> prime; // 保存素数
2
   int sieve(int n) {
3
       isPrime.resize(n + 1, false);
4
       isPrime[0] = isPrime[1] = true;
5
6
       for (int i = 2; i \le n; i++) {
7
           if (!isPrime[i]) prime.emplace_back(i);
8
           for (int j = 0; j < (int) prime.size() && prime[j] * i \le n; ++j) {
               isPrime[prime[j] * i] = true;
9
               if (!(i % prime[j])) break;
10
           }
11
12
13
       return (int)prime.size();
14
```

- 2.3 欧拉定理
- 2.4 欧拉函数
- 2.5 中国剩余定理
- 2.5.1 CRT
- 2.5.2 EXCRT
- 2.6 BSGS
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- 4.3 小球盒子模型
- 4.4 斯特林数
- 5 博弈论
- 6 其他数学
- 6.1 蔡勒公式
- 7 图论
- 7.1 并查集

```
struct dsu {
private:
    // number of nodes
    int n;
    // root node: -1 * component size
    // otherwise: parent
    std::vector<int> pa;
```

```
public:
8
         dsu\left( \begin{array}{cccc} \textbf{int} & n_{-} = \ 0 \end{array} \right) \ : \ n\left( n_{-} \right) \, , \ pa\left( n_{-}, \ -1 \right) \ \left\{ \right\}
9
10
         // find node x's parent
         int find(int x) {
11
12
              return pa[x] < 0 ? x : pa[x] = find(pa[x]);
13
         }
         // merge node x and node y
14
         // if x and y had already in the same component, return false, otherwise
15
              return true
         // Implement (union by size) + (path compression)
16
         bool unite(int x, int y) {
17
18
              int xr = find(x), yr = find(y);
              if (xr != yr) {
19
                   if (-pa[xr] < -pa[yr]) std::swap(xr, yr);
20
                   pa[xr] += pa[yr];
21
                   pa[yr] = xr; // y \rightarrow x
22
23
                   return true;
24
              return false;
25
26
         }
         // size of the connected component that contains the vertex x
27
         int size(int x) {
28
              return -pa[find(x)];
29
30
         }
31
    };
```

- 7.2 最小树形图
- 7.3 最近公共祖先
- 7.4 最大流
- 7.5 全局最小割
- 7.6 二分图最大权匹配
- 7.7 一般图最大匹配
- 7.8 2-sat
- 7.9 最大团
- 8 字符串
- 8.1 KMP
- 8.2 Manacher
- 8.3 Trie
- 8.4 01-Trie
- 9 计算几何