

Complexity & Efficiency

Motivation

Large Scale

- Healthcare.gov
- Amazon S3
- Fortnite

Smaller Scale

- Stock data analysis

Example - Search (Count # of comparisons)

1	3	5	6	7	8	11	13	17	19	20
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Sequential Search
Best: 1 test
Worst: n tests

Binary Search
Best: 1 test
Worst: 1 data item: 1 test
2 data items: 2 tests
3 data items: 3 tests
4 data items: 4 tests
 2^k data items: $k+1$ tests

Trend? \rightarrow # of tests \rightarrow n tests $\rightarrow \log_2 n$

Given n items what is the power of 2 to?

How do I write down ~~the~~ $2^? = n$

$$? = \log_2 n$$

With n items; # of tests is $\log_2 n + 1$

If n is not a power of 2? $\text{floor}(\log_2 n) + 1$

Show google drive.

How do we communicate this? Big-Oh
Write big oh for Binary / Sequential search

Big - Oh

Informal Idea! Remove the constants, leave largest growth term

(Compare with n^2, n)

$O(1)$	constant	
$O(\log n)$	logarithmic	
$O(n)$	linear	
$O(n \log n)$		
$O(n^2)$	quadratic	
$O(n^3)$	cubic	Ex: Matrix Multiplication
$O(2^n)$	exponential	Ex: Fibonacci (actually the constant is smaller), generate subsets
$O(n!)$	factorial	Ex: TSP

~~Proof~~

$$5n^2 + \frac{n}{3} \text{ is } O(n^2)$$

Means for some value $c > 0$, $c \cdot n^2 \geq 5n^2 + \frac{n}{3}$ once n is big enough

Formal Definition: $f(n)$ is $O(g(n))$ means for some $c > 0$

$$c \cdot g(n) \geq f(n) \text{ for } n \geq K$$

$$c \cdot n^2 \geq 5n^2 + \frac{n}{3}$$

$$c \geq 5 + \frac{1}{3n}$$

$$\text{Choose } c=6, K=1$$

Show $f(n)$ is not $O(g(n))$?

#1 $C \cdot n \geq n^2$

n^2 is to $O(n)$

$C \geq n$
constant goes to infinity

#2 $C \cdot \log(n) \geq n$ $\log n$ is not $O(\log(n))$

$C \geq \frac{n}{\log(n)}$ $\xrightarrow{\frac{dn}{dn}}$ $\lim_{n \rightarrow \infty} \frac{n}{\log n} = \lim_{n \rightarrow \infty} \frac{1}{\frac{1}{n}} = n$

Stacks

Idea: How do I tell if parentheses are balanced?

$$2 + (7 \cdot 9) / (4 + (6 + 2))$$

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Data Structure : Stack

Stacks

Characteristics :
- allows duplicates
- Last in, First out (LIFO) order

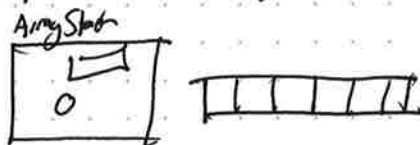
Operations :
push (T item)
pop() T
boolean isEmpty

Ex: Braces and Parentheses

$$\{2 + (7 \cdot 9) / (4 + [6 + 2])\} + 9$$

push { push { pop } pop } pop

Implement with array

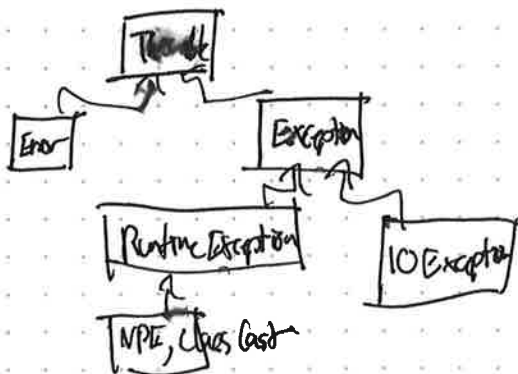


Do array stack from Array Ring

Exceptions

try throw notation
create our own exception

Checked Exception
Try / catch



Queues

Characteristics:

- Allows duplicates
- First in, First out order FIFO

Applications?

Linked Queue

void enqueue add at rear
T dequeue
T getFront
is Empty

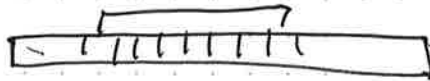
Trick: use Front and Rear pointers

Array Queue

#1 Shift

#2 Track Front

#3 Wrap around



2 versions of wrapping

~~rear~~ (FR)

```
if ((rear + 1) == max) {  
    rear = 0  
} else {  
    rear++  
}
```

rear = (rear + 1) % max

→

How do we tell if circular queue is full?

Leave an empty spot ✓

Store a placeholder

Keep a counter

Why do we have to define ~~to~~ CarStack / CarQueue?

Stack ~~is Queue~~ ^{if} ~~are~~ old and not written as interfaces
Queue uses weird terminology

add/offer, remove/poll