

ACTUATORS AND CYBERPHYSICAL SYSTEMS

DRAFT-MATTSSON-CORE-COAP-ACTUATORS
MATTSSON, FORNEHED, SELANDER, PALOMBINI
IRTF T2TPRG, OCT 31 2015

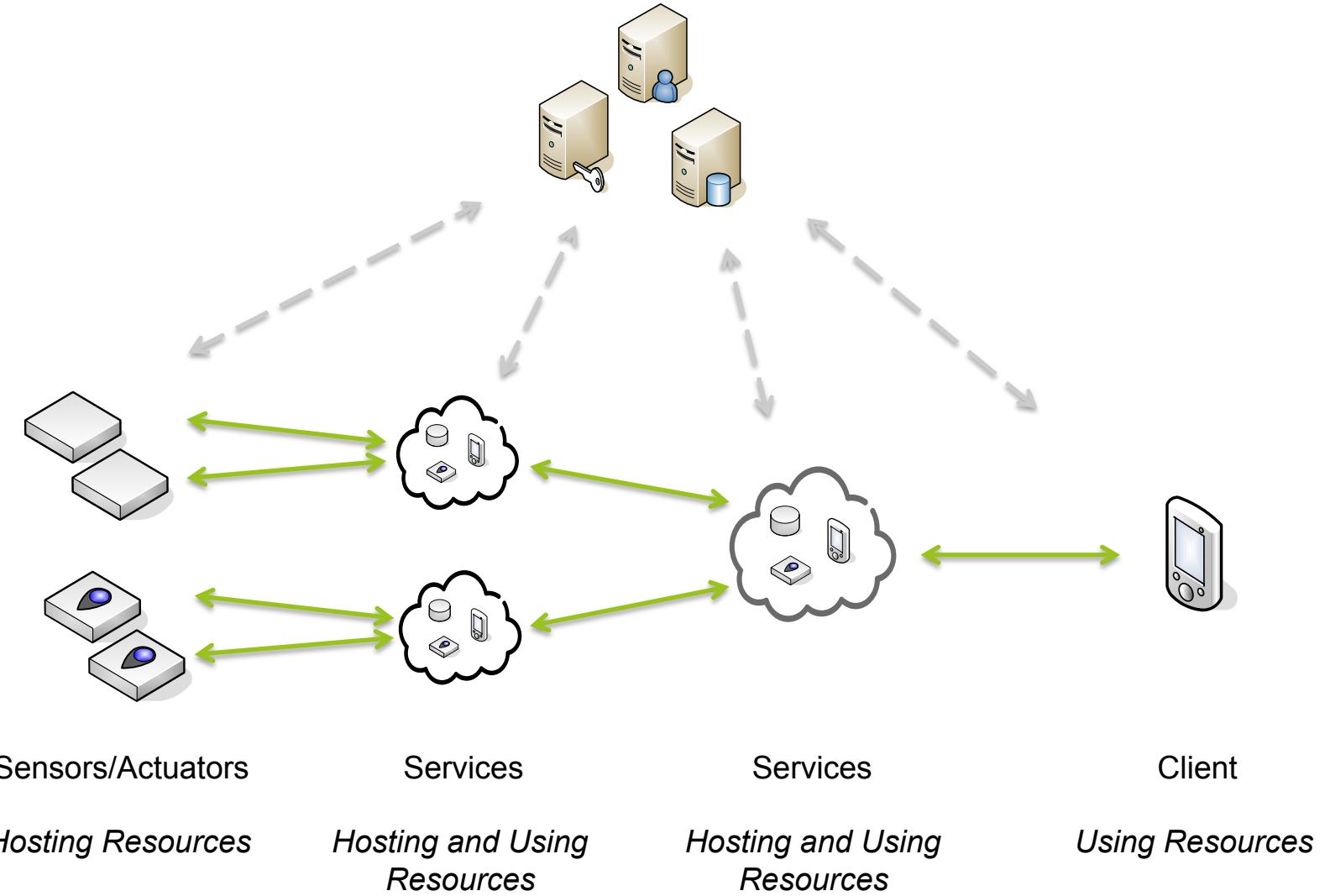


BACKGROUND

SECURITY AND PRIVACY FOR IOT SENSORS IS COMPLEX



- The *store-and-forward* and *publish-subscribe* architectures rely on middle boxes.
- As DTLS only offers hop-by-hop security, fully trusted intermediaries are necessary.
- e2e security mechanisms needed.



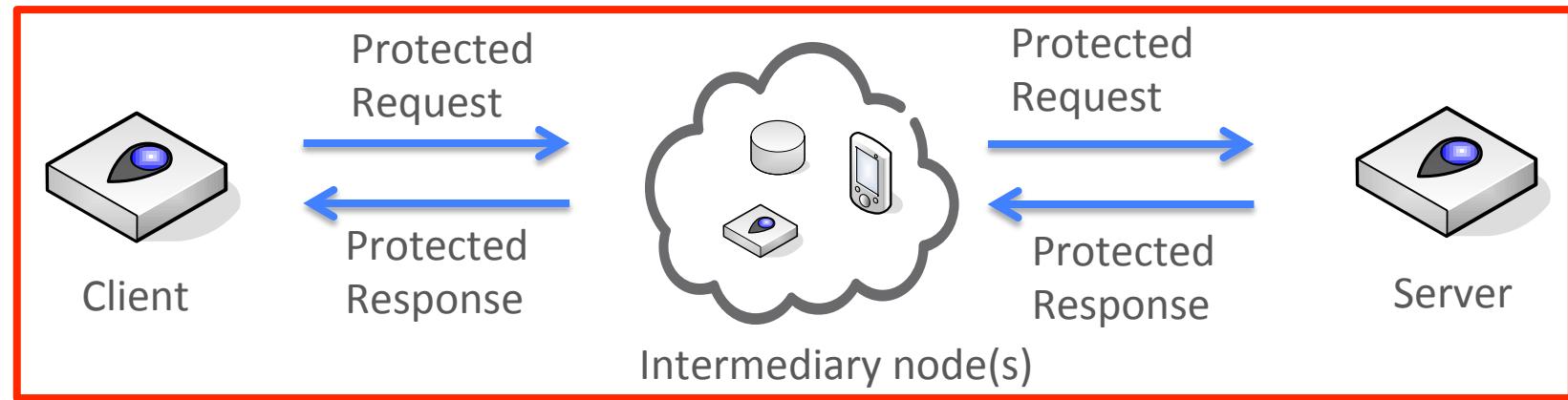
E2E SECURITY OVER MIDDLEBOXES



Object Security: More details in draft-selander-ace-object-security.

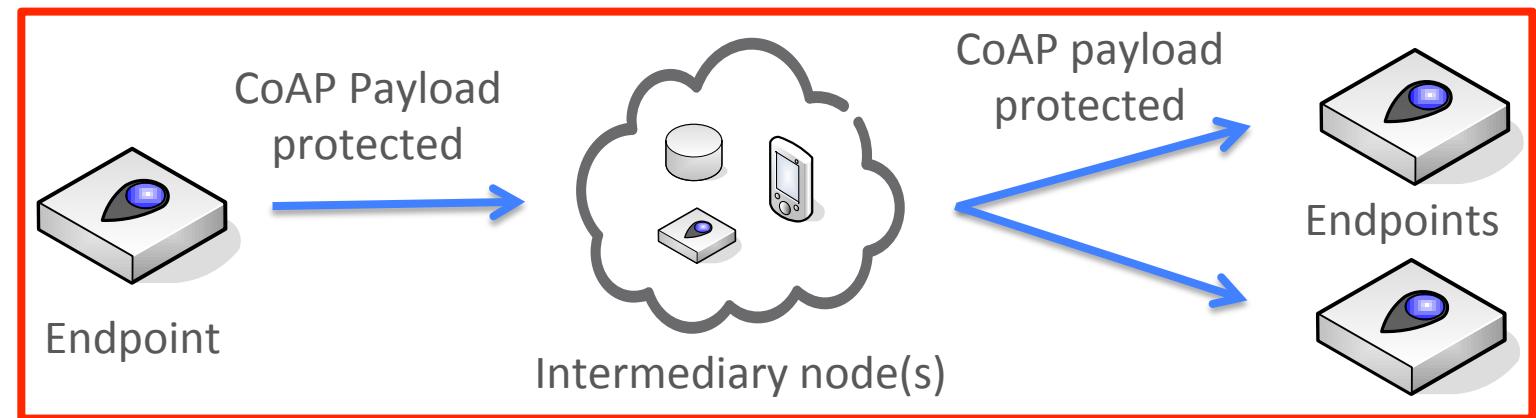
OSCOAP

- › CoAP Option that protects CoAP request-response
- › Supports store-and-forward



OSCON (COSE)

- › Protects CoAP Payload only
- › Supports publish-subscribe



REQUEST-RESPONSE MATCHING IN COAP, DTLS, AND OSCOAP

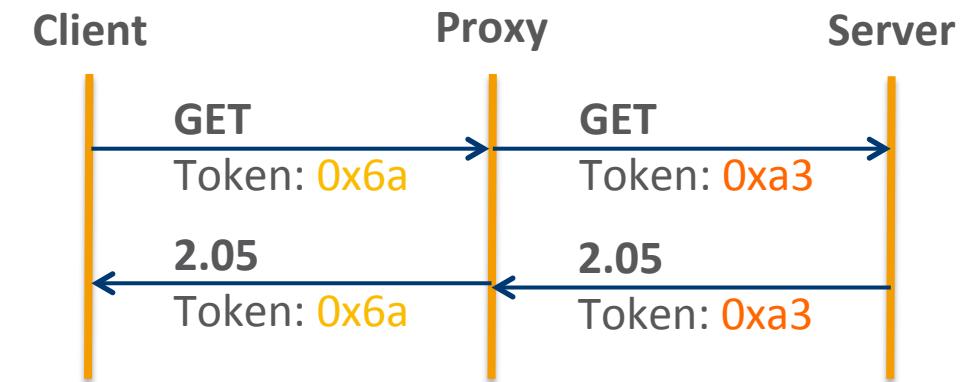


CoAP

- Response-request matching is done with hbh token

DTLS

- "Independent" sender and receiver contexts.
- No binding** between request and response.
- Responses can typically be **delivered out-of-order**.

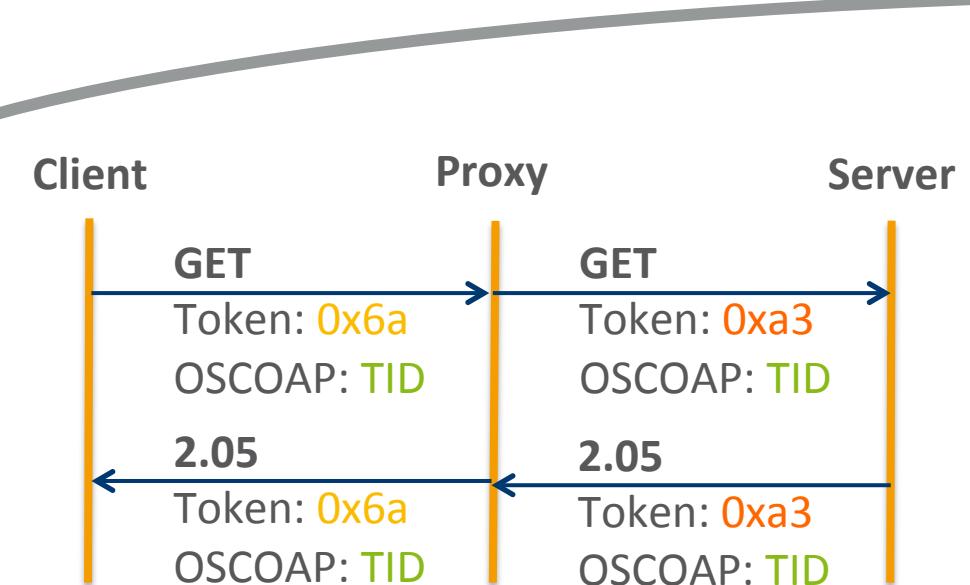


CoAP over DTLS

- CoAP RFC says that tokens currently "in use" SHOULD (not SHALL) be unique.
- Less strict** profiling of token when used with DTLS.
 - This leads to serious attacks.**

CoAP with OSCOAP

- e2e Transaction ID (TID) provides request-response binding.



ATTACKS

Assumptions:

- Attacks are performed by an on-path attacker, not a “trusted” middle box.
- CoAP is protected with a security protocol such as DTLS or OSCOAP.

BLOCK ATTACK

- Blocking sensor messages is just DoS.
- Blocking actuator messages results in the client losing information.
- Is the door locked or unlocked?
- Remedy: Inform user.

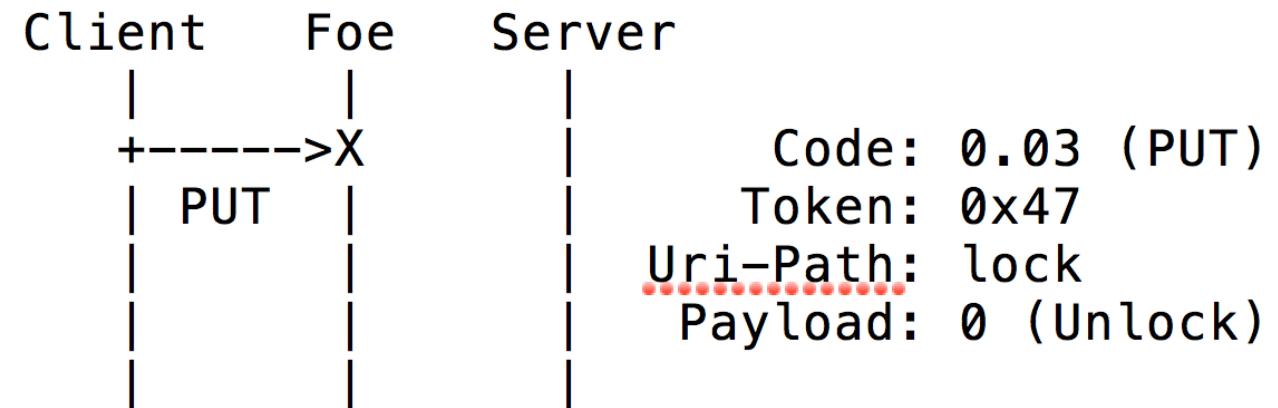


Figure 1: Blocking a Request

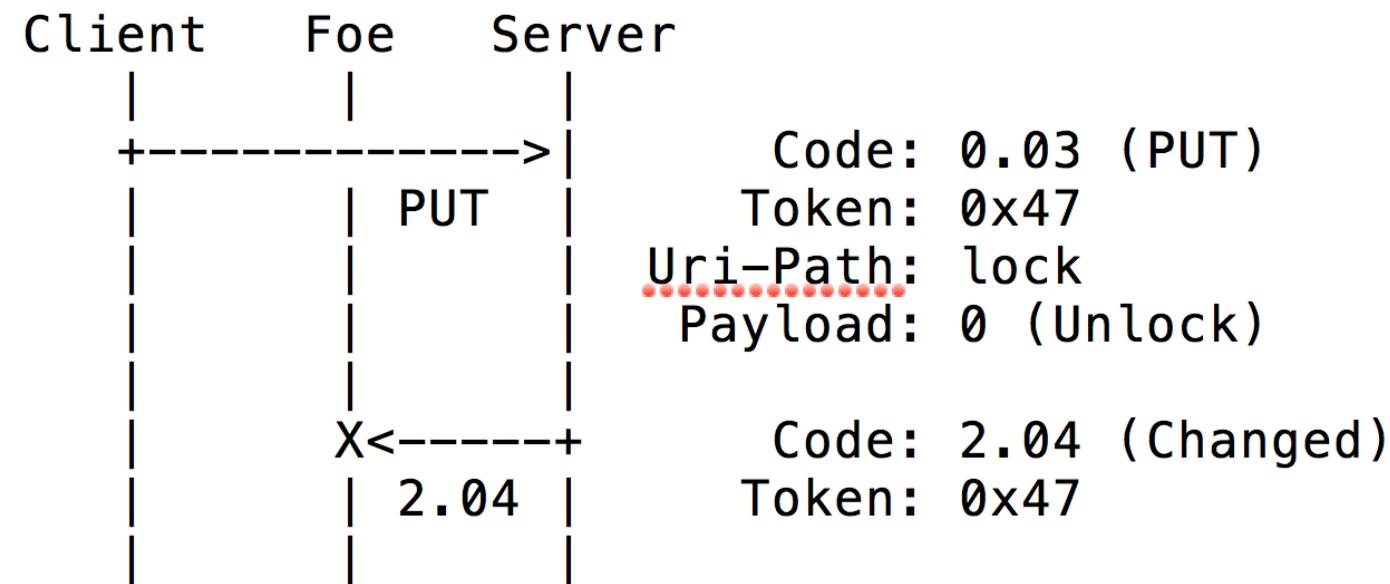


Figure 2: Blocking a Response

DELAY ATTACK (1/2)



- An attacker can delay the delivery of any packet (request or response) by a chosen amount of time.
- The attacker can control the replay window.
- Attacker can e.g. unlock a door at a given time.

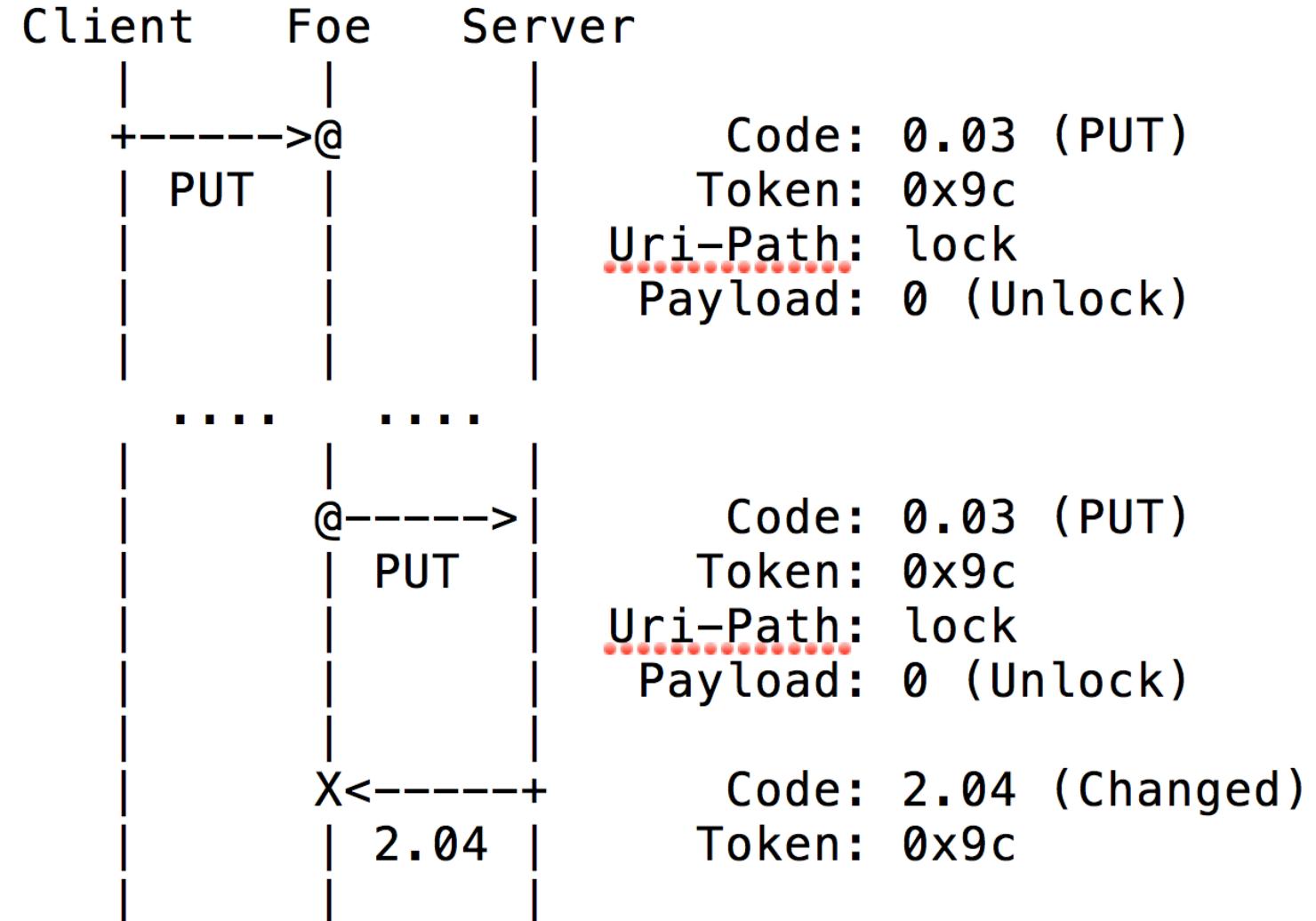


Figure 3: Delaying a Request

DELAY ATTACK (2/2)

- **Reordering:** If a non-zero replay window is used, the attacker can let the client interact with the actuator before delivering the delayed request.
- **Remedy:** New CoAP option:
The Repeat Option

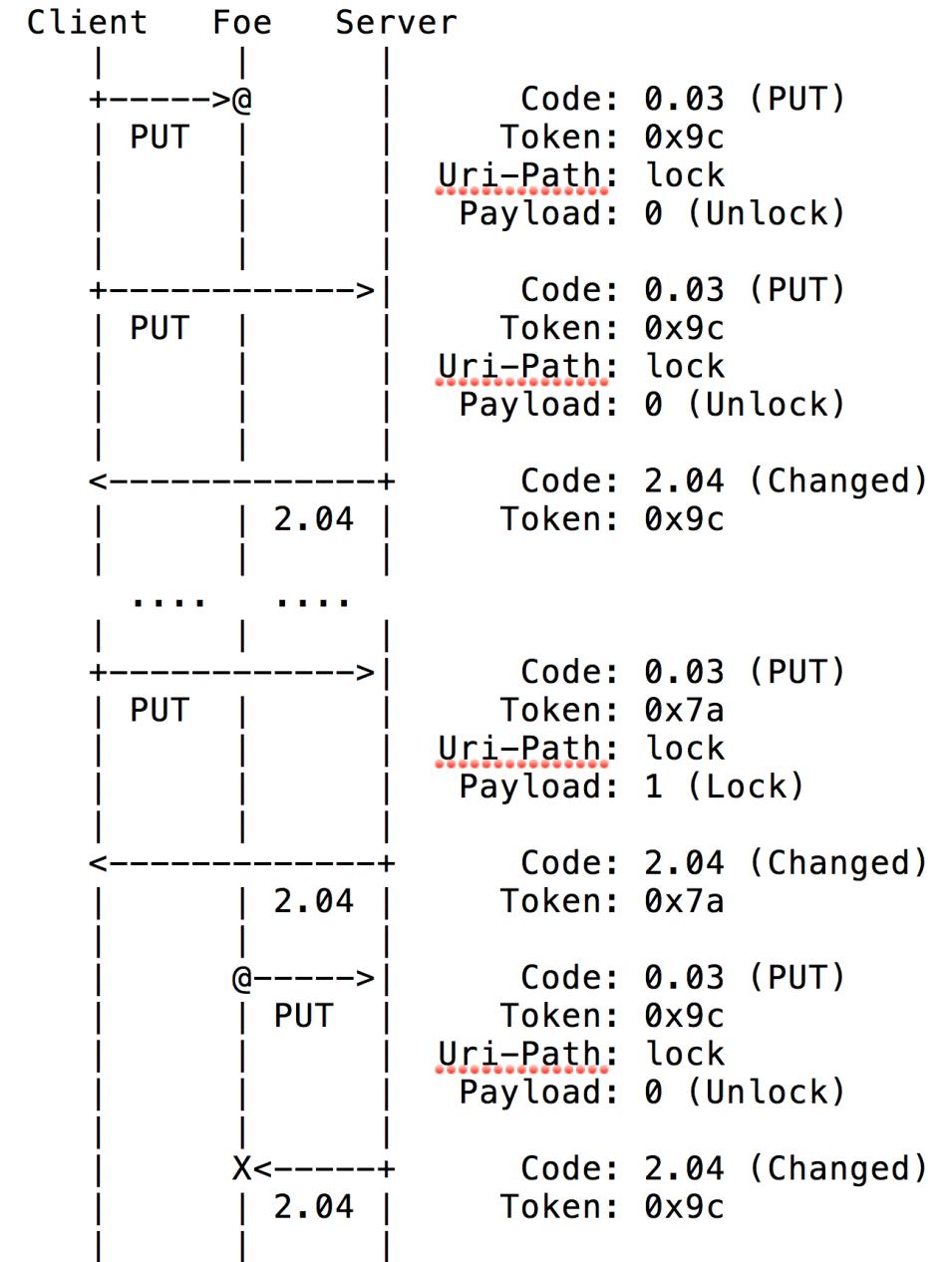


Figure 4: Delaying Request with Reordering

MISMATCH ATTACK (1/3)

- Possible if CoAP is protected with DTLS (but not OSCOAP).
- The client needs to reuse the same token.
- The attacker can still control the replay windows.
- Attacker can e.g. fool client to believe a door has been locked.

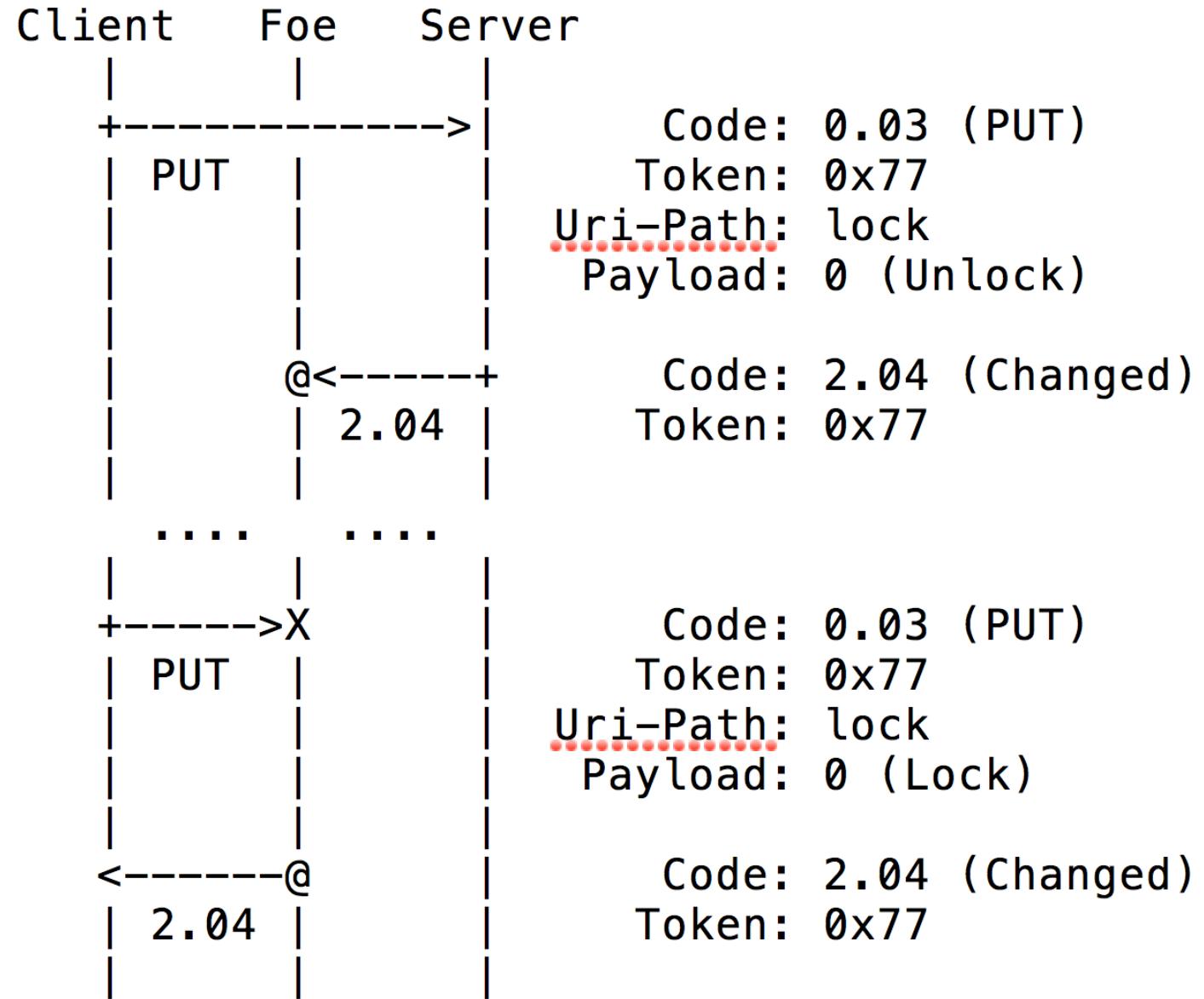


Figure 5: Mismatching Response to PUT

MISMATCH ATTACK (2/3)

- Attack is also valid for sensors.
Also here with serious consequences.
- Attacker can e.g. fool client to believe a door is locked.

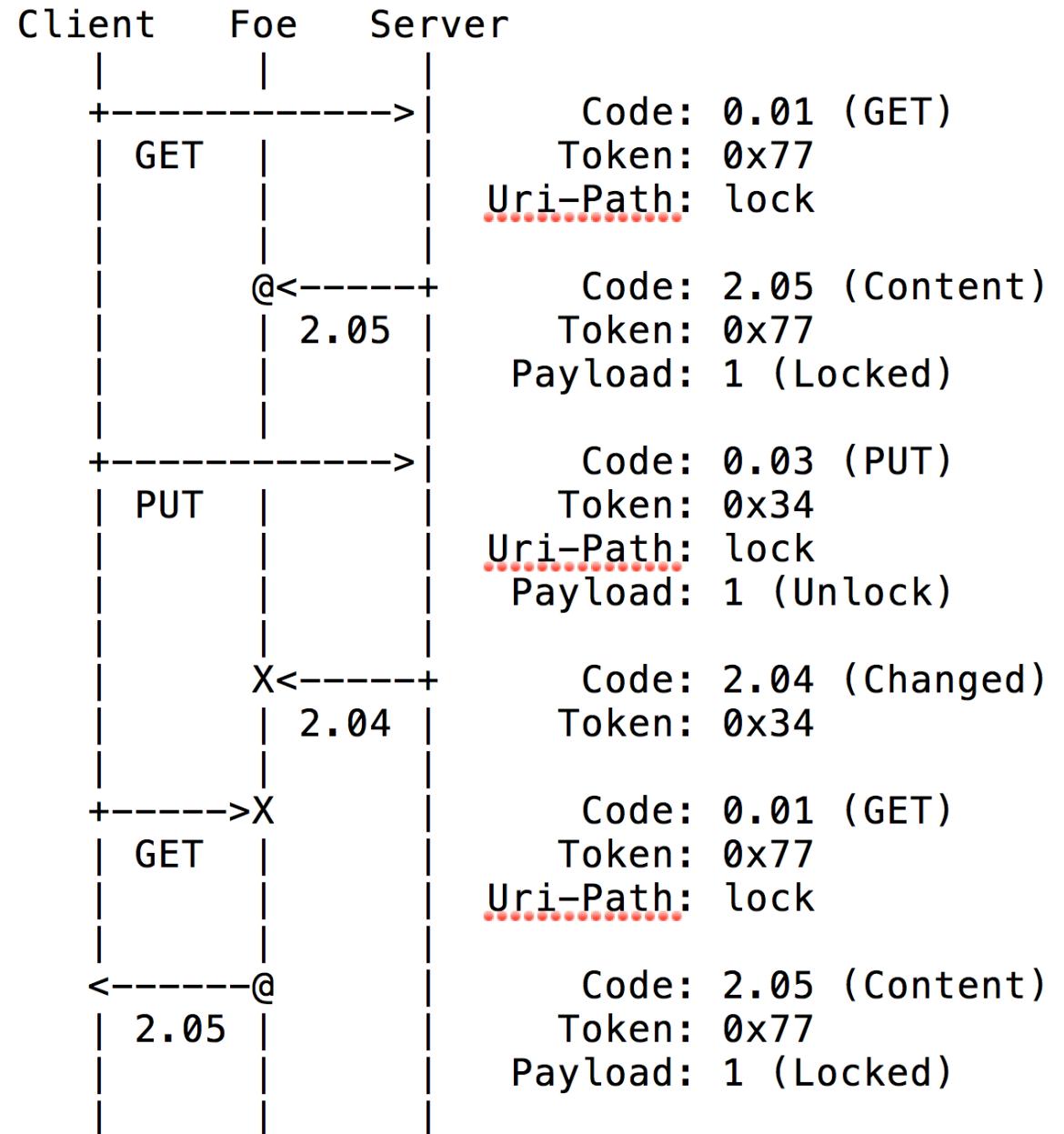


Figure 6: Mismatching Response to GET

MISMATCH ATTACK (3/3)



- Attacker can even mix responses from different resources if the resources share the same DTLS connection **on some part of the path.**
 - Located behind a common gateway
 - Served by the same CoAP proxy.
- Example: Attacker makes client believe the house is on fire.
- Remedy: Never reuse tokens (which the client has not received responses to).

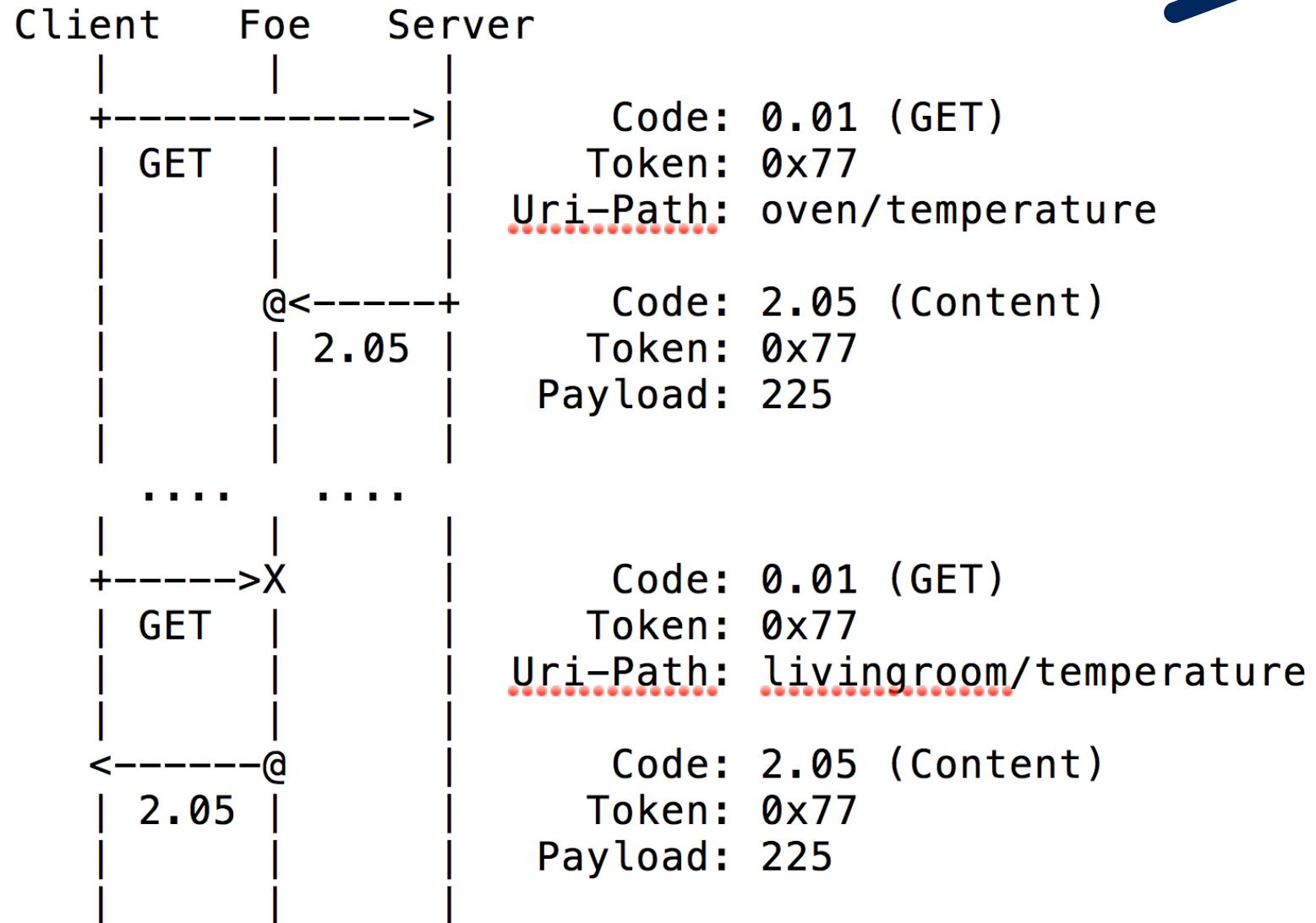


Figure 7: Mismatching Response from other resource

RELAY ATTACK



- Common with deployments where actuator actions are triggered automatically based on proximity.
- E.g. a car (the client) constantly polling for the car key (the server) and unlocking both doors and engine as soon as the car key responds.
- An attacker relaying the CoAP messages out-of-band, **completely breaks the security**.
- Unfortunately already happening in practice.
- Remedy: Don't use automatic car locks.

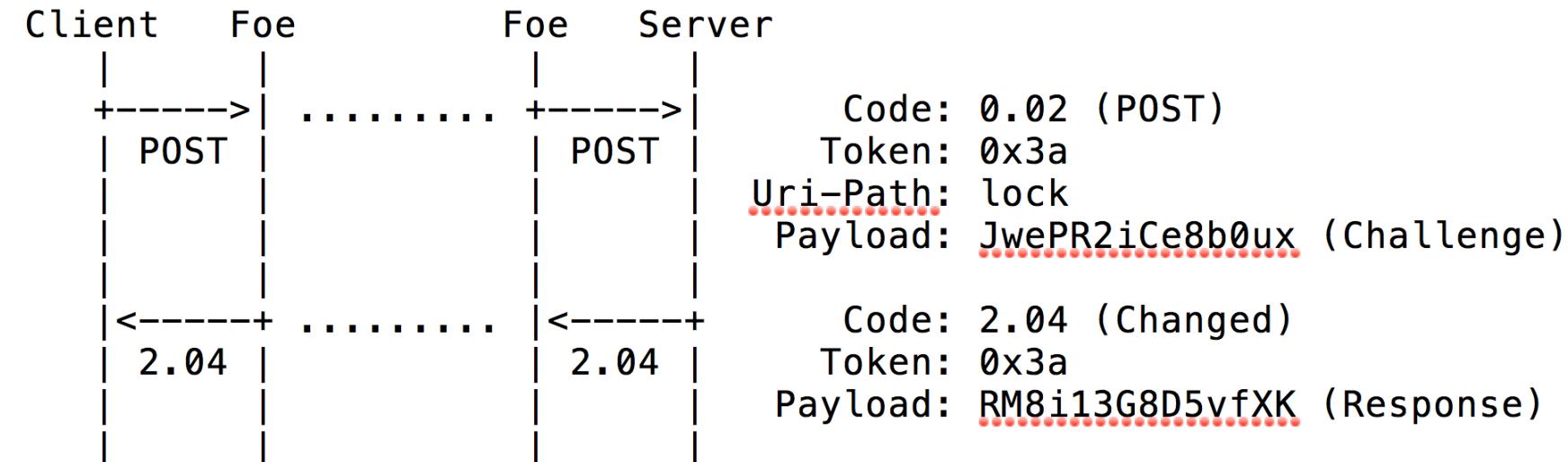


Figure 8: Relay Attack (the client is the actuator)



THE REPEAT OPTION

REPEAT OPTION



- New e2e challenge-response mechanism for CoAP, binding a request to an earlier response.
- The challenge (for the client) is simply to echo the value.
- Enables the server to verify the freshness of a request:
 $(t_1 - t_0) = \text{RTT} < \text{Threshold}$, thus mitigating delay attacks.

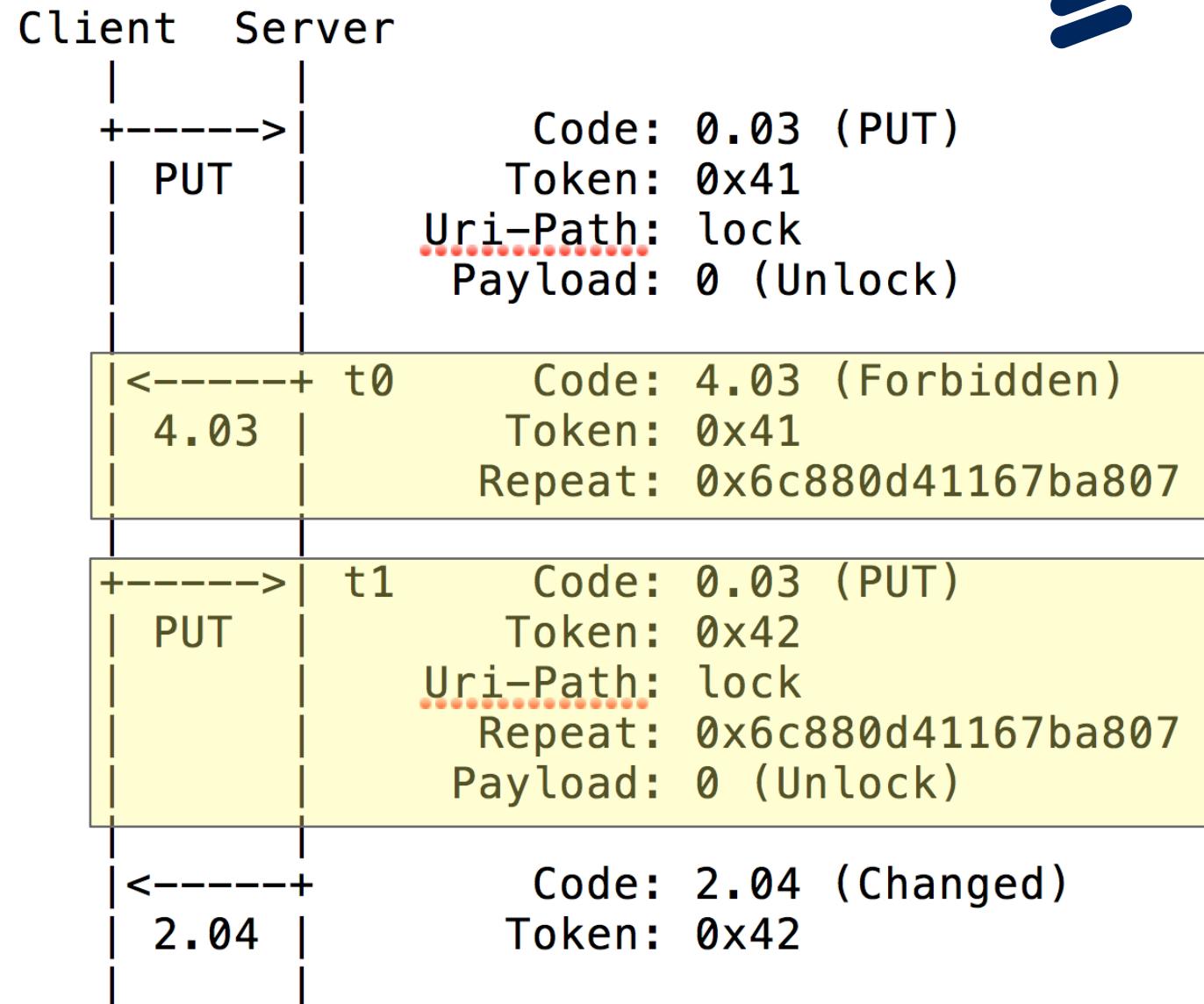


Figure 9: The Repeat Option

REPEAT OPTION



- Gives the server freshness guarantee independently of what the client does. But takes two round-trips and server has to save the option value and time t_0 .
- The server may simply send the current time t_0 . The client may then estimate the current time in the servers timescale when sending future requests (i.e. not echoing). This reduces round-trips and server state, but has security problems.
- The server may instead of a pseudorandom value send an CCM-encrypted timestamp:

option value = AEAD(key, t_0 , NULL)
or AEAD(key, t_0 , parts of request)

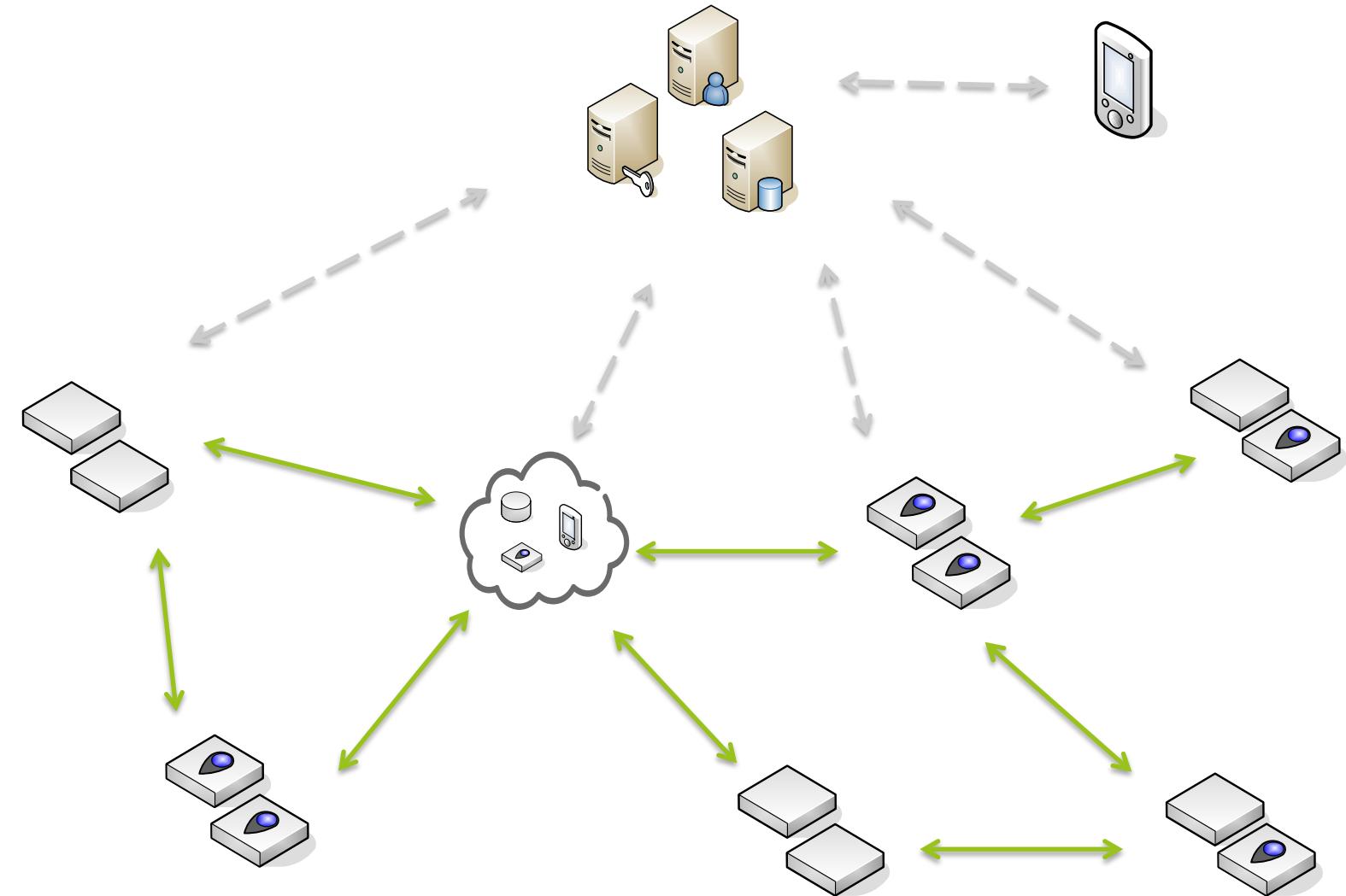
This does not reduce the number of round-trips but eliminates server state.

CYBERPHYSICAL

CYBERPHYSICAL SYSTEMS



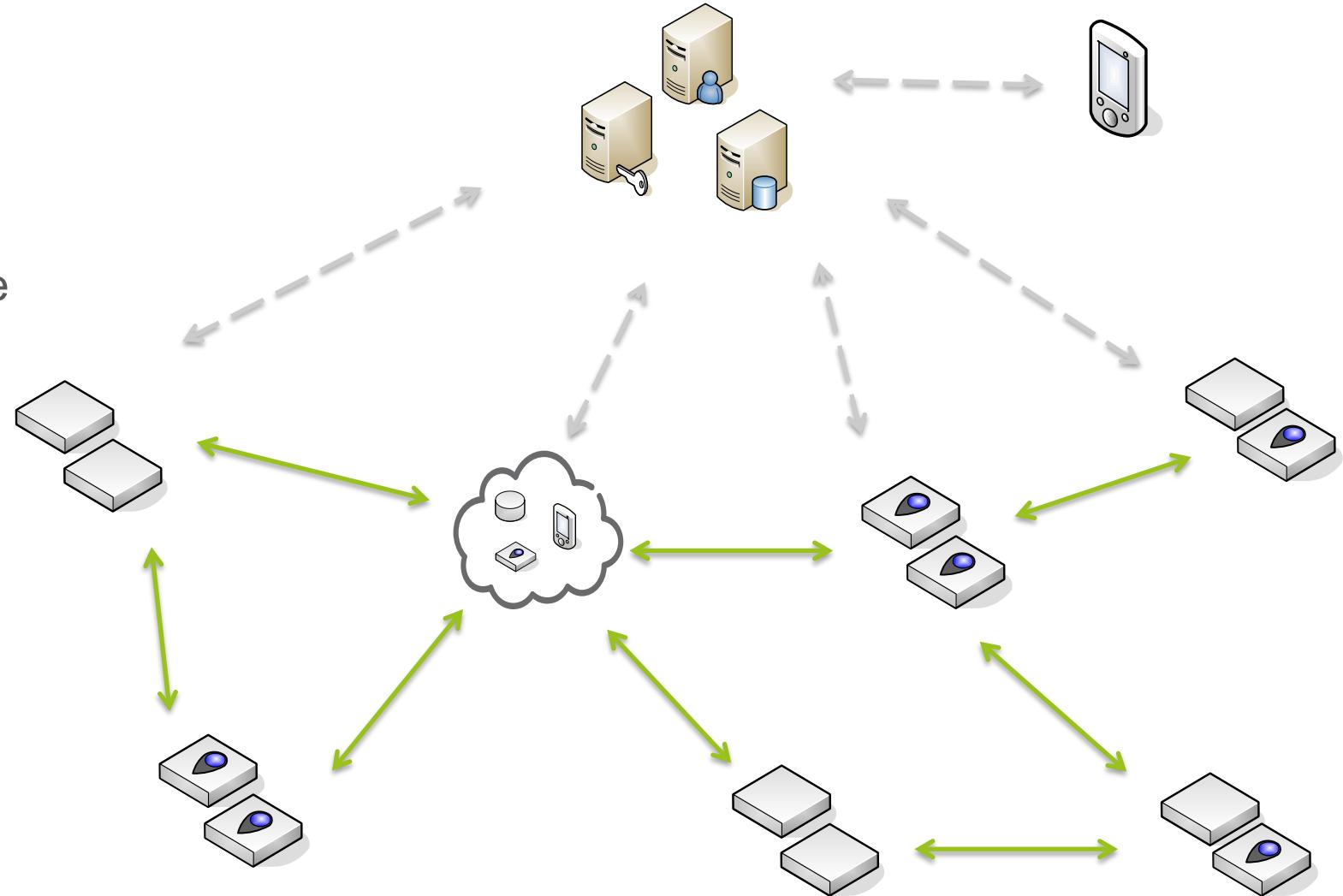
- Security and privacy for sensors is complex
- Security for actuators is worse
- Security for privacy for cyberphysical systems where actuators are controlled by sensors is much worse.



CYBERPHYSICAL SYSTEMS



- Attackers can as before compromise nodes, block, delay, and relay messages.
- Attackers can also break sensors or actuators, influence sensors, move sensors, change time, and selectively block, delay or relay new programming...
- How to get general such systems secure and reliable?
- Which layers?
- A job for T2TRG?



CONCLUSIONS

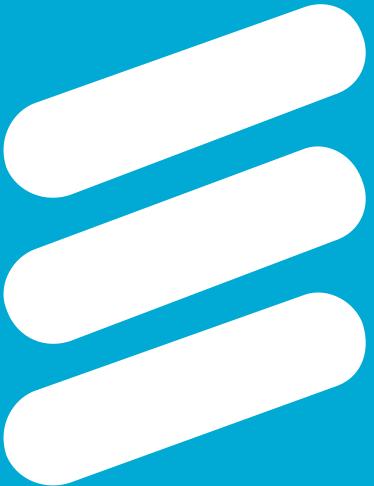
CONCLUSIONS

IoT systems needs more then integrity for individual messages:

- **Message bundle integrity.** Binding for example a response to a request, or an authentication handshake together. Lack of this has led to various attacks on e.g. TLS. Responses should always be bound to requests.
- **Time integrity**, binding a message or a bundle to a particular time. Standard to do in authentication protocols using timestamps or challenge-response, but missed when sending things over security protocols.
- **Space integrity**, binding a message or a message flow to a particular location. Is done in for e.g. HDCP 2.0 but needed also for cyberphysical systems.

Adding TLS or DTLS is not enough. Security analysis needed.





ERICSSON