

Programmable Networks with Synthesis



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Network Misconfigurations are Common

Amazon server outage affects millions of companies and causes online chaos

MARCH 1, 2017 12:39PM

AMAZON'S giant servers crashed today causing chaos for millions of companies and people that use the cloud service and affected everything larger web sites to people's smarthomes and even library catalogues.

Amazon's massive AWS outage was caused by human error

One incorrect command and the whole internet suffers.

BY JASON DEL REY | @DELREY | MAR 2, 2017, 2:20PM EST

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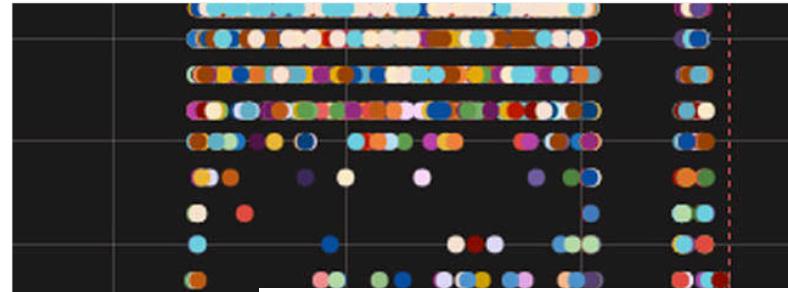


Sean Gallup / Getty

Amazon today blamed human error for the big AWS outage that took down a bunch of large Internet sites for several hours on Tuesday afternoon.

CloudFlare apologizes for Telia screwing you over

Unhappy about massive outage

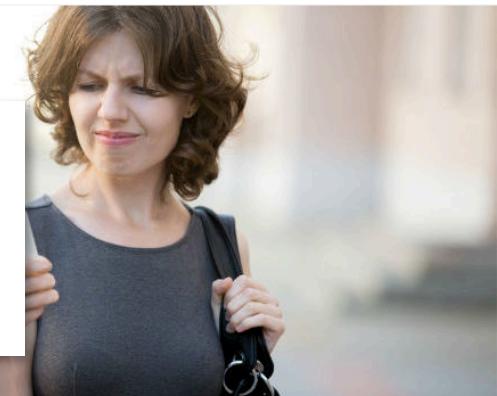


Networks

Level3 switch config blunder blamed for US-wide VoIP blackout

Network dropped calls because it was told to

21 Jun 2016 at 20:34, Kieren Mc



The summer of network misconfigurations

by Joanne Godfrey on August 11, 2016 in Application Connectivity Management, Firewall Change Management, Information Security, Risk Management and Vulnerabilities, Security Policy Management

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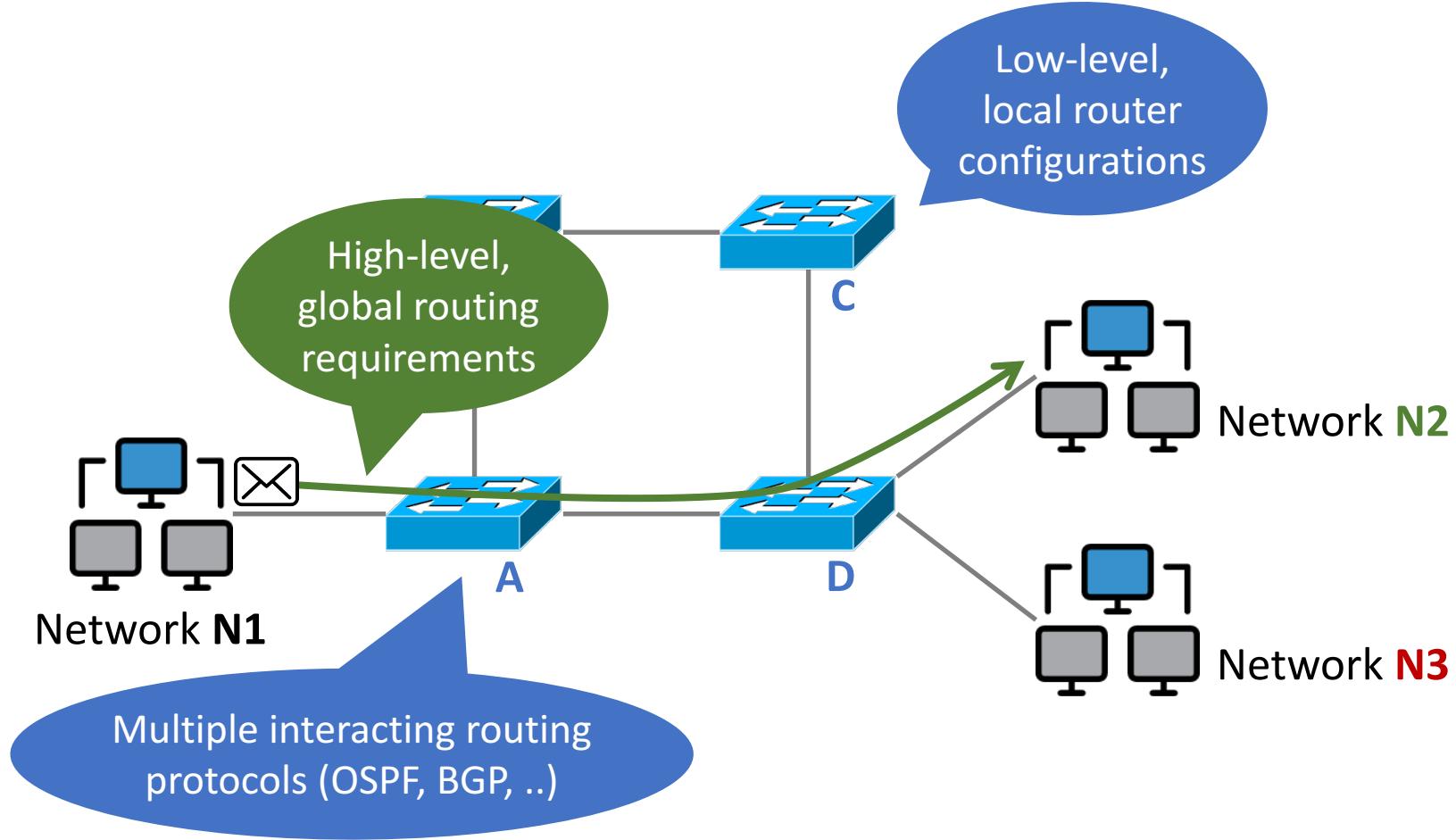
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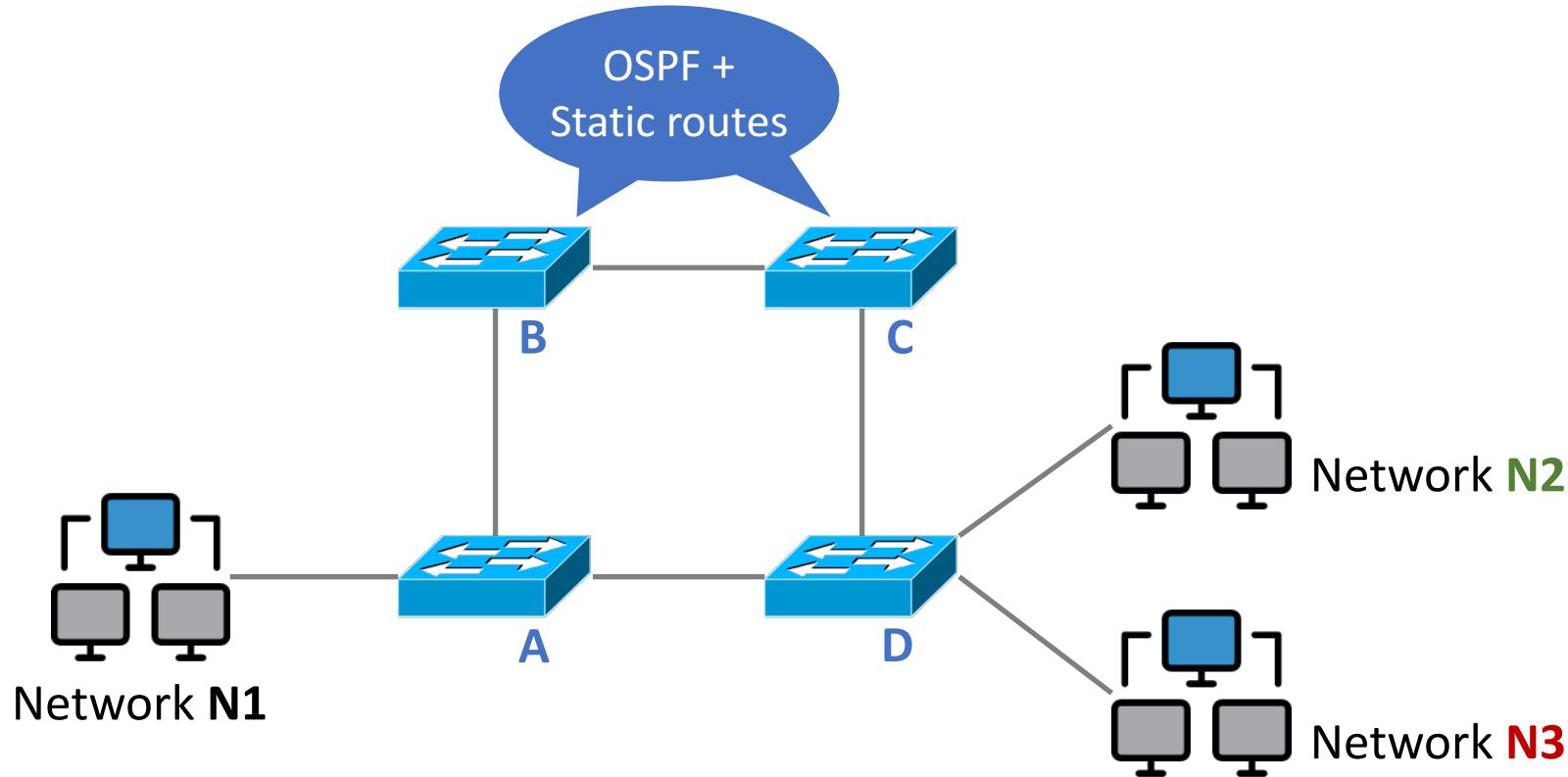
5 Oct 2016 at 11:33, Shaun Nichols

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What Makes Network Configuration Hard?



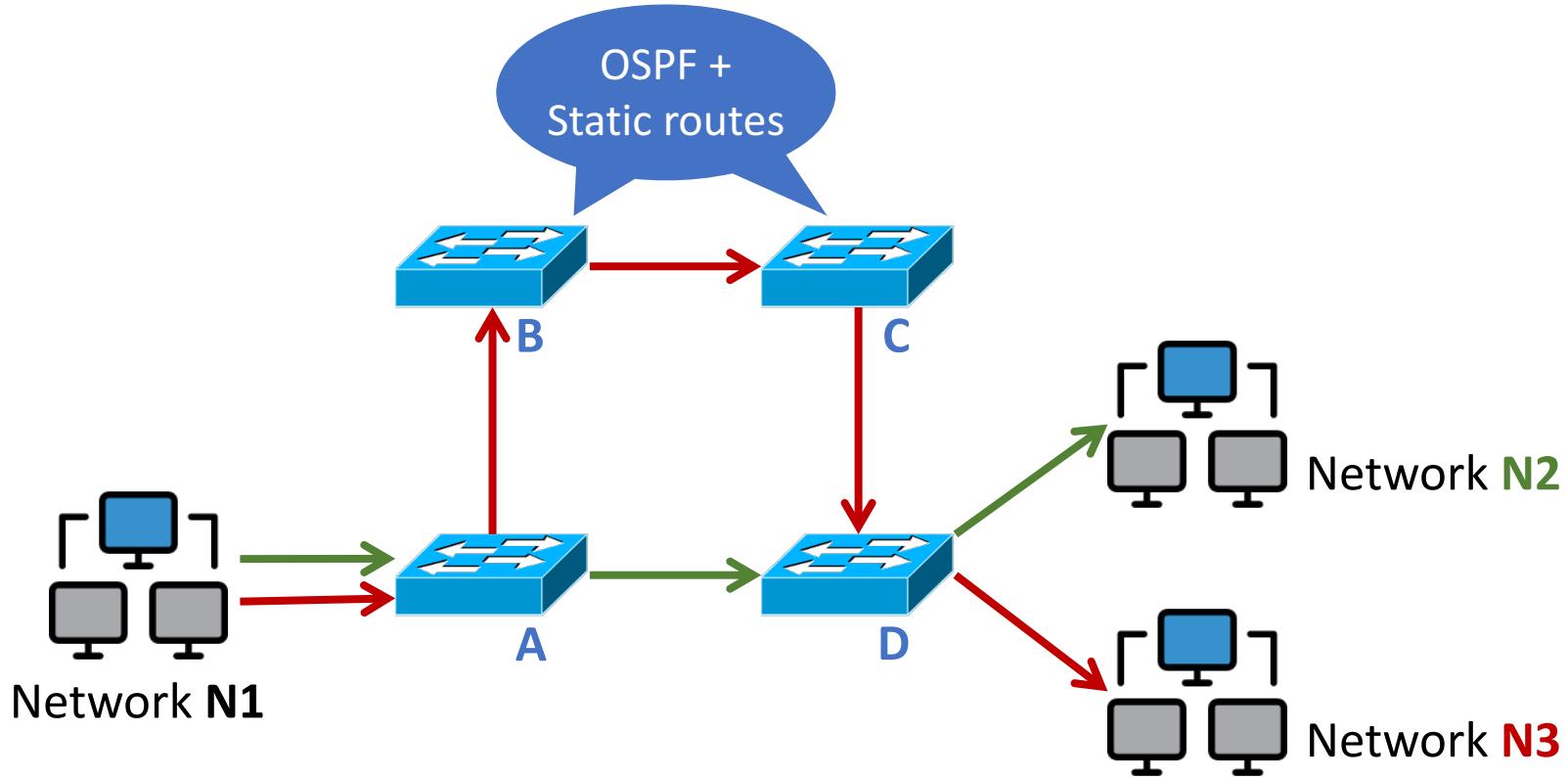
Example



R1: Packets from **N1** to **N2** must follow the path **A → D**

R2: Packets from **N1** to **N3** must follow the path **A → B → C → D**

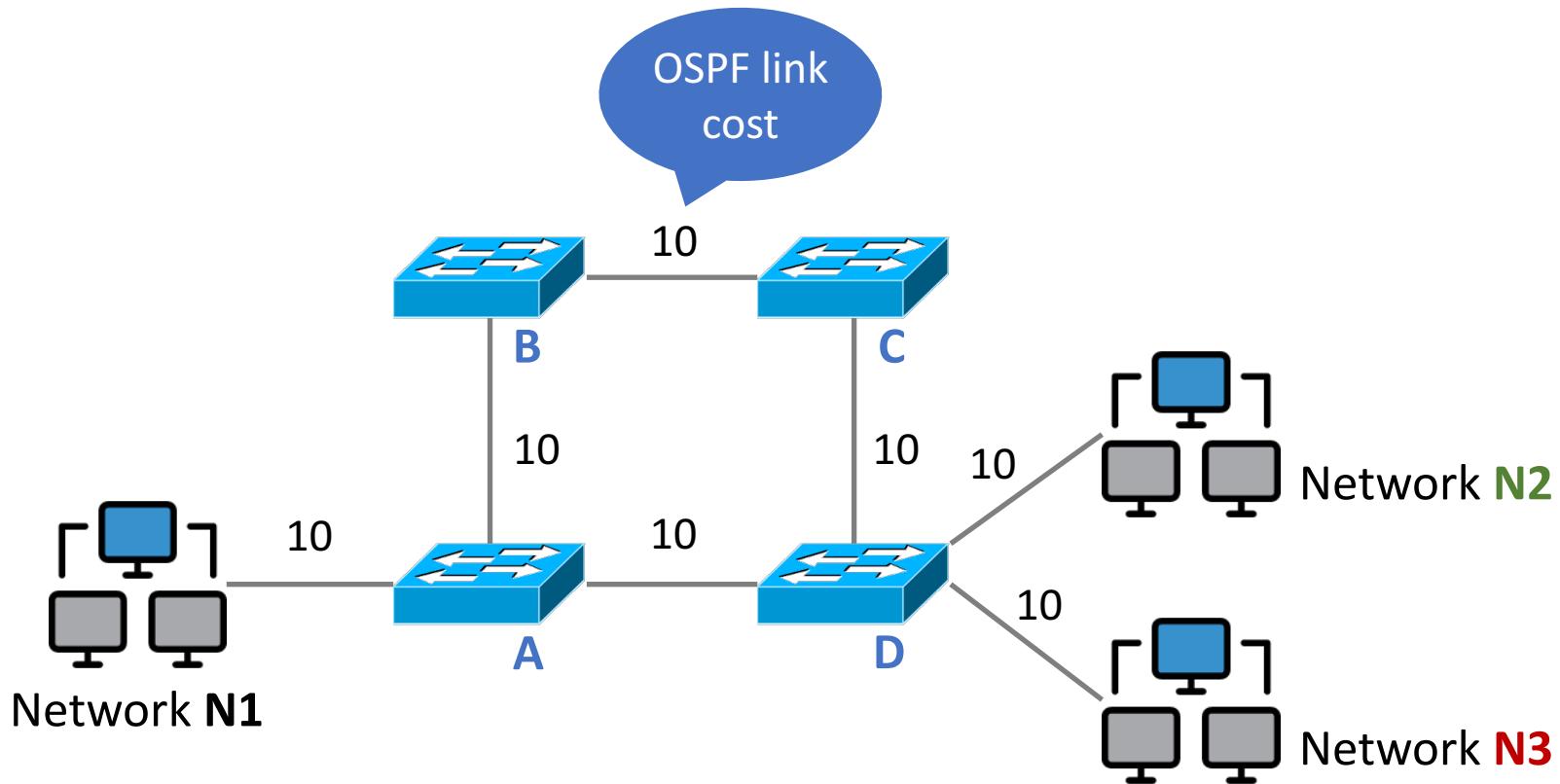
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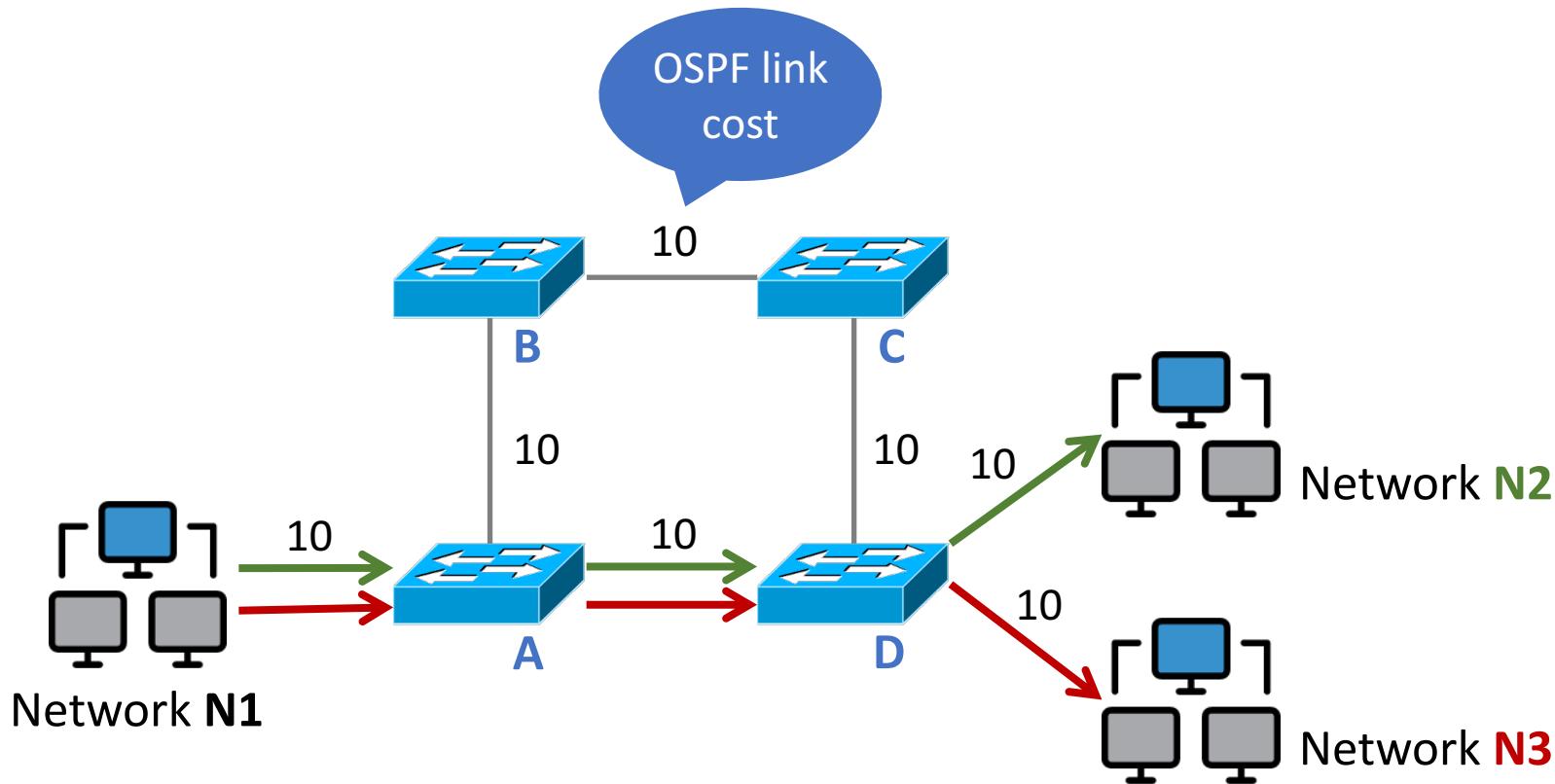
Example



R1: Packets from **N1** to **N2** must follow the path **A → D**

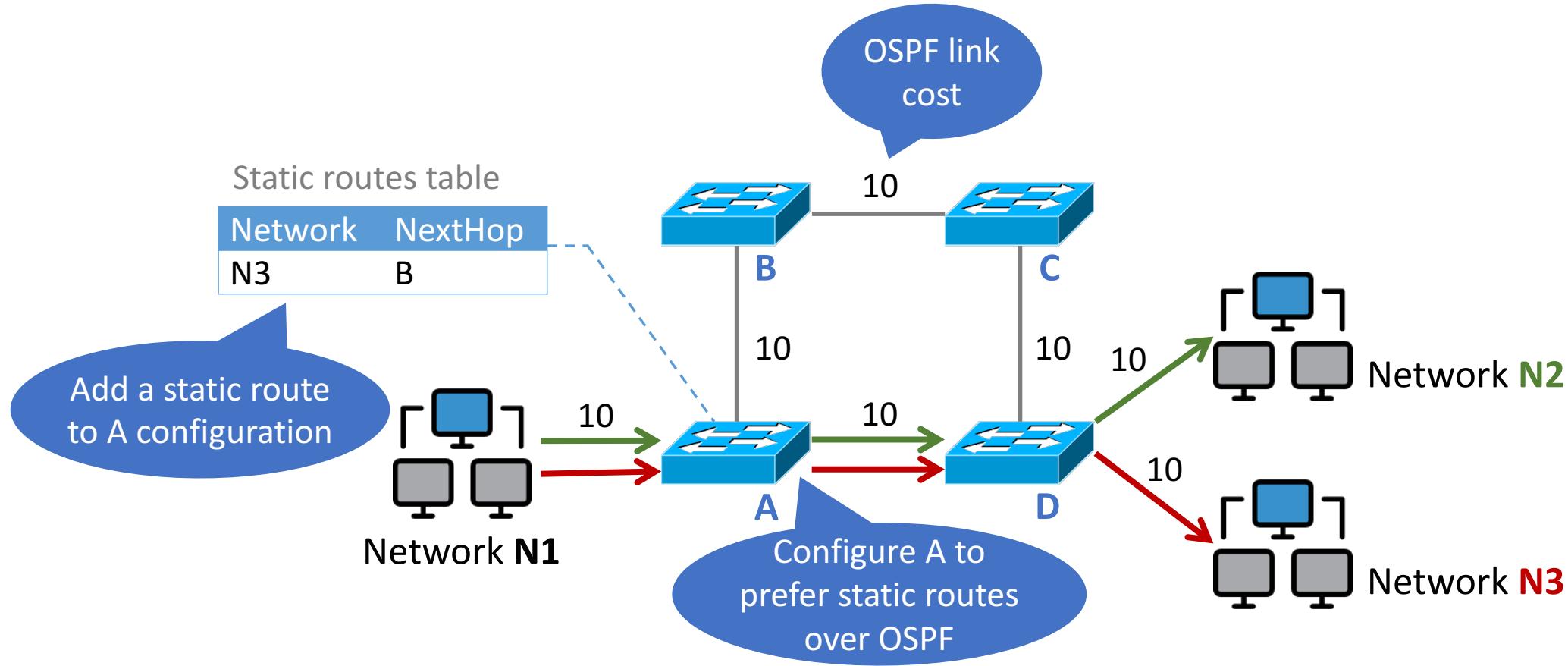
R2: Packets from **N1** to **N3** must follow the path **A → B → C → D**

Example



- ✓ R1: Packets from **N1** to **N2** must follow the path **A → D**
- ✗ R2: Packets from **N1** to **N3** must follow the path **A → B → C → D**

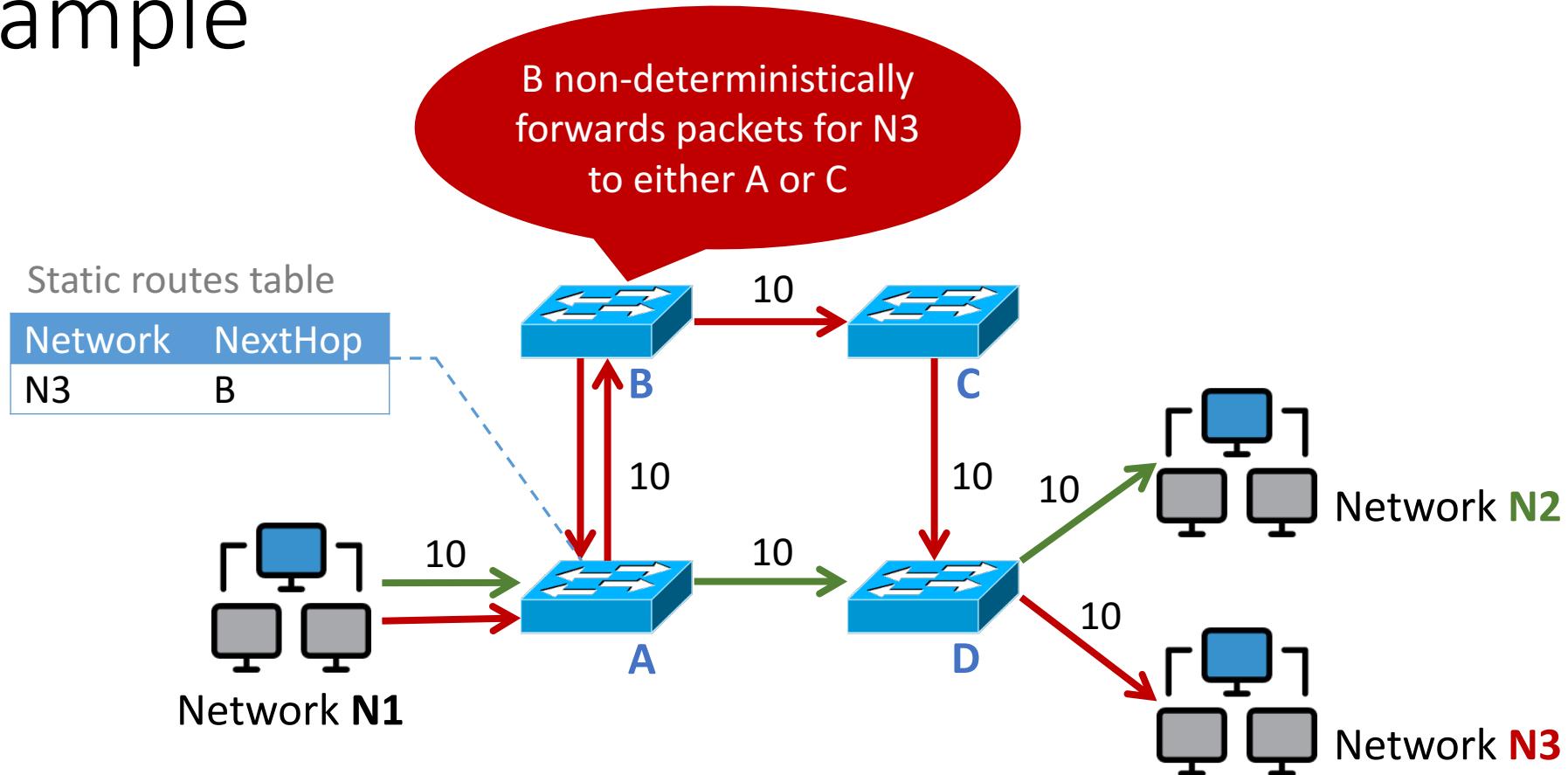
Example



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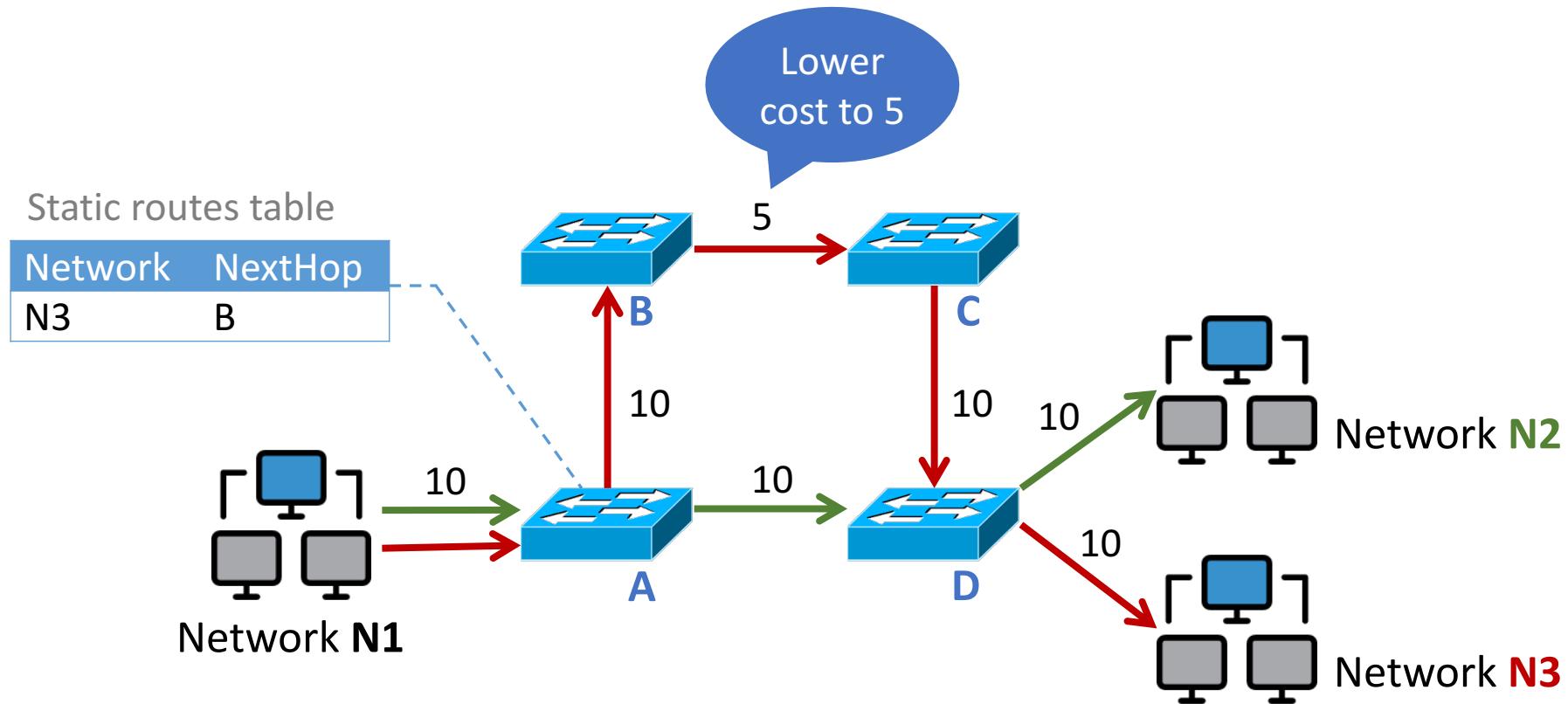
Example



✓ R1: Packets from **N1** to **N2** must follow the path **A** → **D**

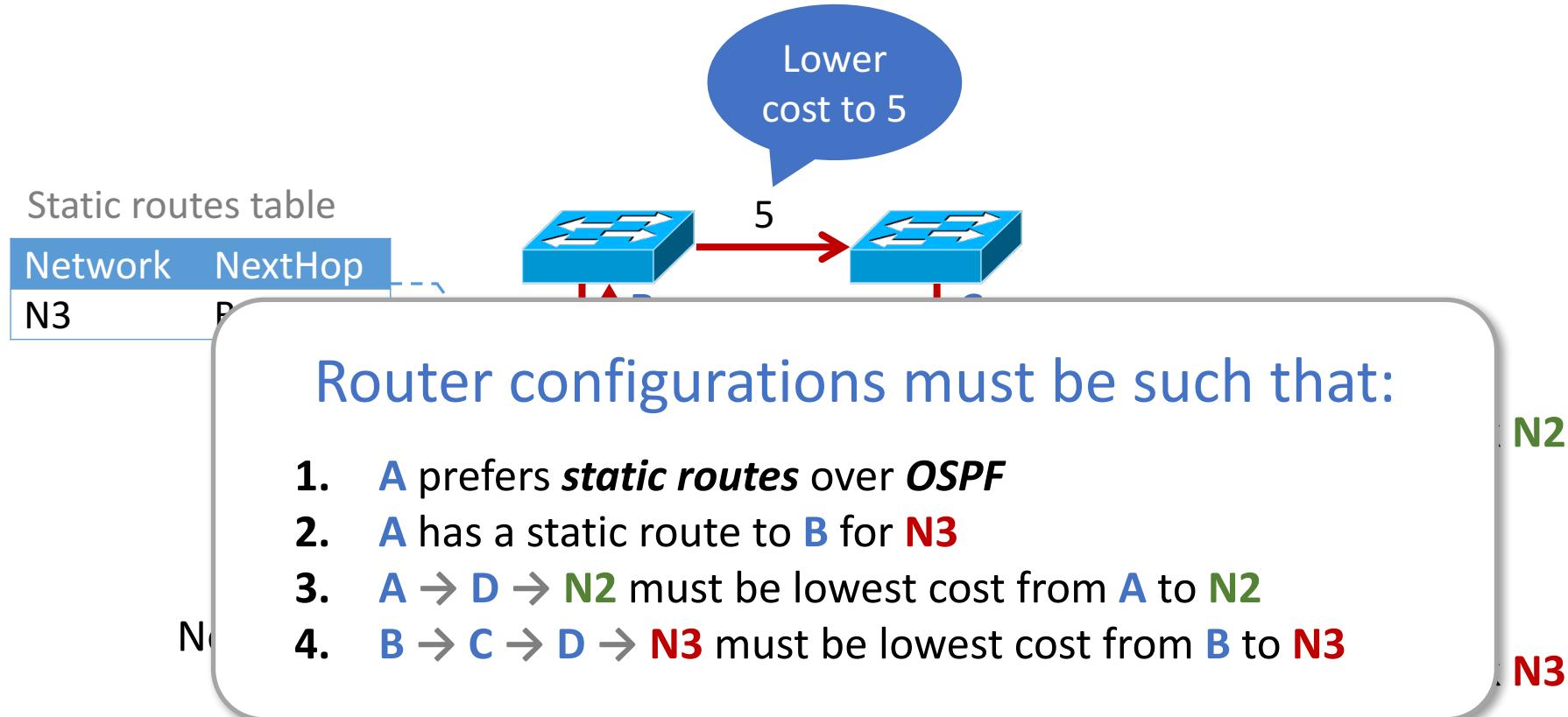
✗ R2: Packets from **N1** to **N3** must follow the path **A** → **B** → **C** → **D**

Example



- ✓ R1: Packets from **N1** to **N2** must follow the path **A → D**
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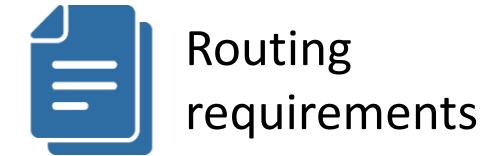
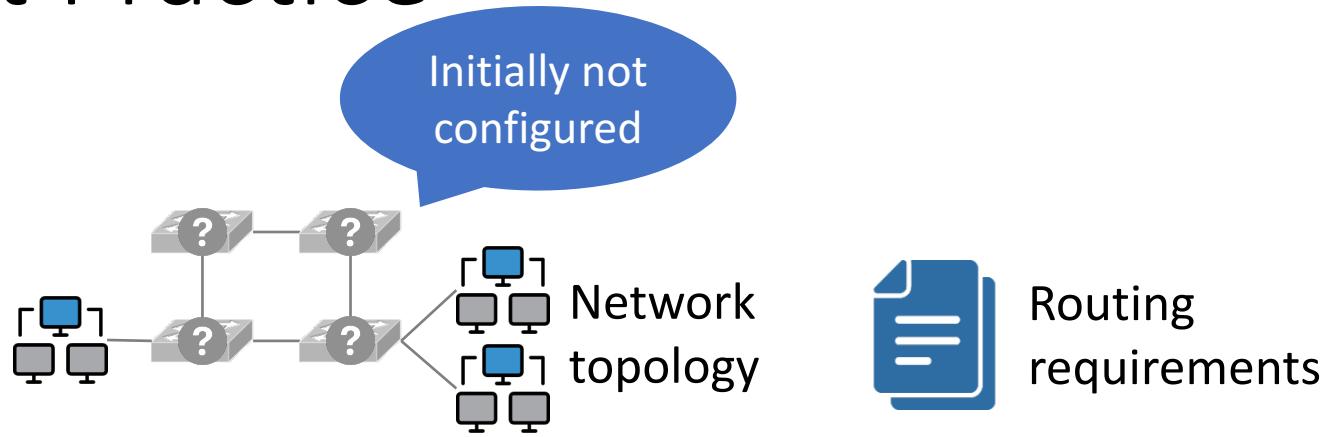
Example



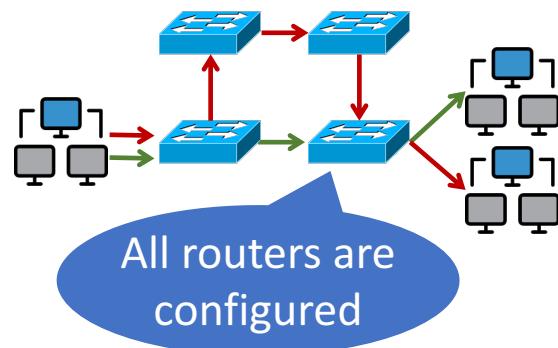
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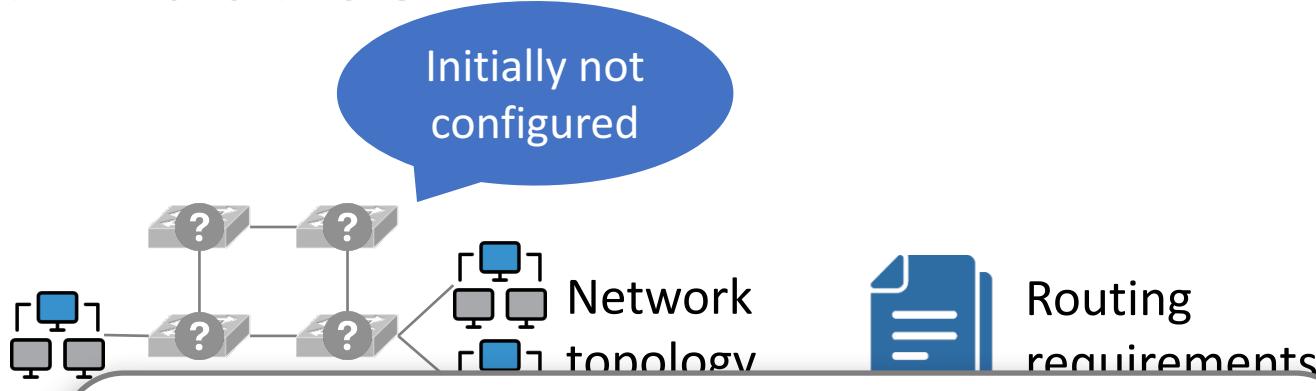
Current Practice



Operators manually
configure each router

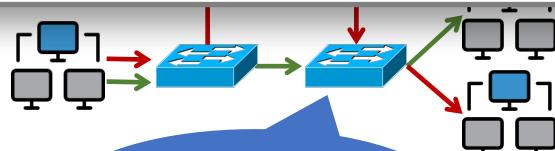


Current Practice



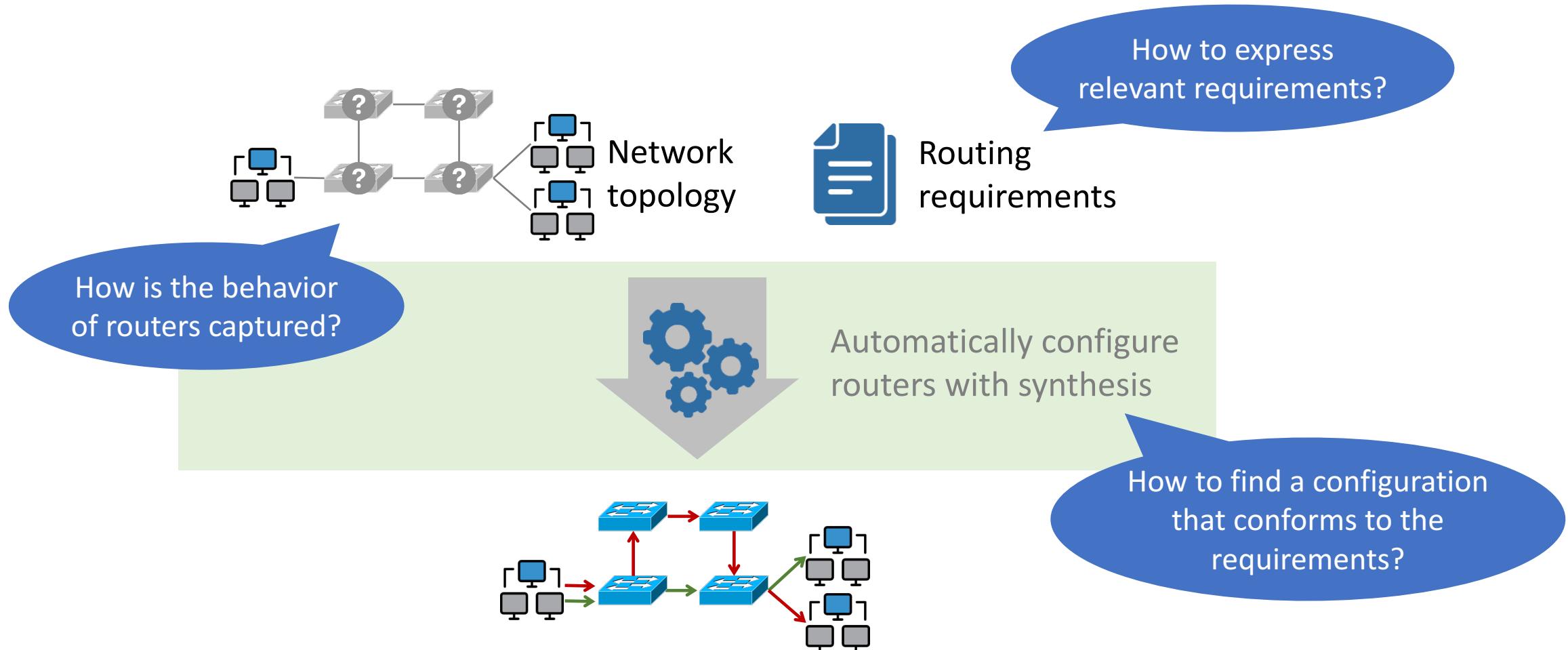
Problems and Challenges

- *Diversity* in protocol expressiveness
- Protocol *dependencies*
- No *correctness* guarantees



All routers are configured

Wanted: Programmable Networks with Synthesis



Programmable Networks with *Synthesis*: Dimensions

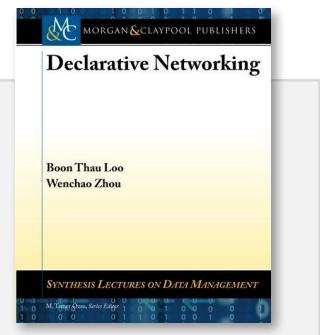
Deployment scenarios	Datacenter Incremental	ISP Enterprise
Routing protocols	Deterministic OSPF BGP Static routes MPLS	Probabilistic ECMP Gossip
Requirements	Paths Isolation Reachability Waypointing	Failures Congestion
Synthesis Techniques	Enumerative learning Symbolic execution (SyNET**)	Probabilistic Constraint solving CEGIS

**SyNET: <http://synt.net.ethz.ch>

Capturing Network Behavior



Key idea: Express routing protocols, along with their dependencies, in *stratified Datalog*



Datalog Example

Input

```
parent(bob, alice)  
parent(carol, alice)
```

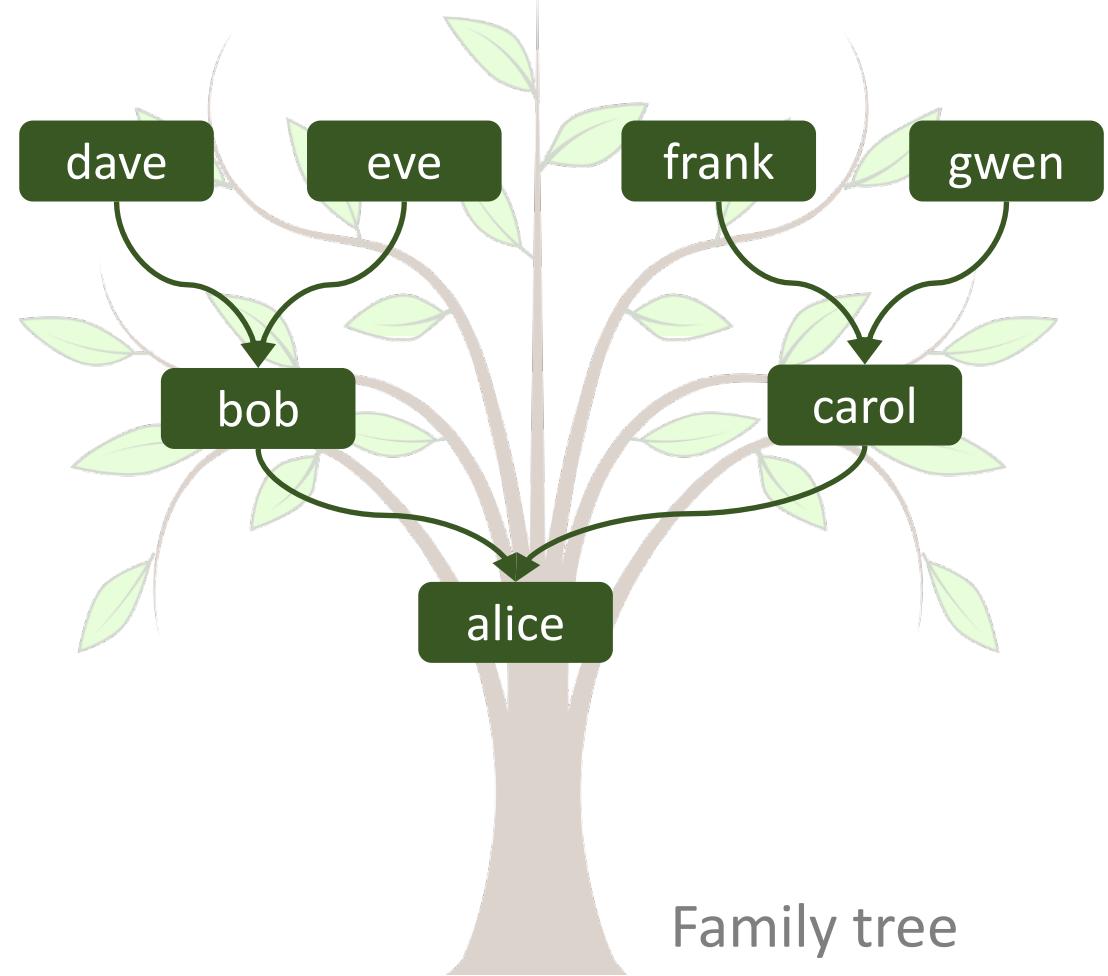
...

Program

```
anc(X, Y) ← parent(X, Y)  
anc(X, Y) ← parent(X, Z), anc(Z, Y)
```

Query

```
anc(dave, alice)?
```



How is this related to networks?

Datalog Syntax (1/2)

To define a Datalog program, we need:

Constants: $\mathcal{C} = \{alice, bob, carol, \dots\}$

Variables: $\mathcal{V} = \{X, Y, \dots\}$

Predicates: $\mathcal{P} = \{parent, anc, \dots\}$

The sets above can be used to construct the “atoms” of a Datalog program:

Ground atoms: $A_{\mathcal{P}(\mathcal{C})} = \{p(t_1, \dots, t_n) \mid p \in \mathcal{P}, \forall 0 \leq i \leq n. t_i \in \mathcal{C}\}$

Atoms: $A_{\mathcal{P}(\mathcal{C}, \mathcal{V})} = \{p(t_1, \dots, t_n) \mid p \in \mathcal{P}, \forall 0 \leq i \leq n. t_i \in \mathcal{C} \cup \mathcal{V}\}$

Example:
 $parent(dave, alice)$

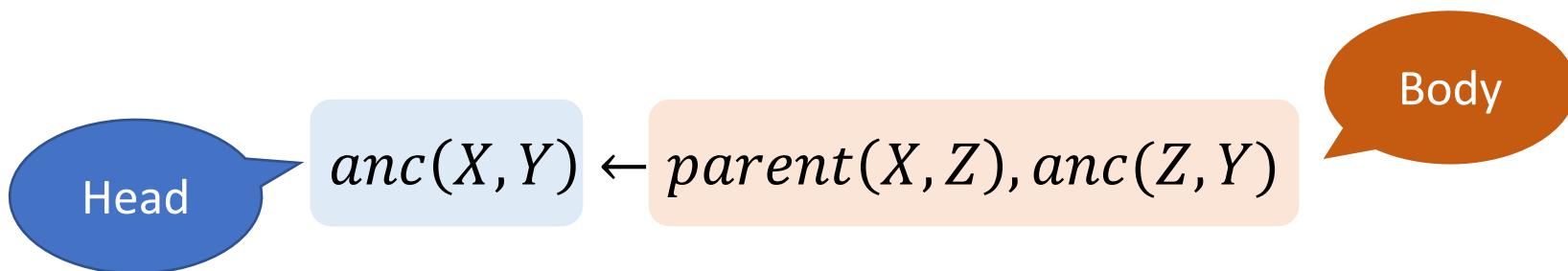
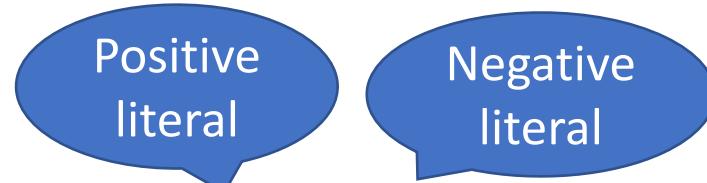
Example:
 $parent(X, Y)$

Datalog Syntax (1/2)

A **Datalog program** is a set of rules of the form

$$a \leftarrow l_1, \dots, l_n$$

where a is an atom and l_1, \dots, l_n are literals of the form a or $\neg a$



A Datalog program is **well-formed** if for any rule in the program, all variables that appear in the head also appear in the body

Is this rule well-formed $anc(X, Y) \leftarrow parent(Y, Z)$?

Semantics of *Positive* Datalog Programs

Each Datalog program P is associated with a consequence operator T_P :

Interpretations: $\mathcal{I} = 2^{A_{\mathcal{P}}(\mathcal{C})}$ (an interpretation I is a set of ground atoms)

Substitutions: $\sigma: \mathcal{V} \rightarrow \mathcal{C}$ (a substitution maps variables to constants)

Consequence operator: $T_P: \mathcal{I} \rightarrow \mathcal{I}$

$$T_P(I) = \{\sigma(a) \mid a \leftarrow l_1, \dots, l_n \in P, \forall i \in [1 \dots n]. I \vdash \sigma(l_i)\}$$

where: $I \vdash l_i$ if $l_i = a$ and $a \in I$

$I \vdash l_i$ if $l_i = \neg a$ and $a \notin I$

A Datalog program P is **positive** if the negation operator does not appear in its rules

Is T_P monotone if P is a positive Datalog program?

The **semantics** of a positive Datalog program P is given by the least-fixed point of T_P

How can we compute the least-fixed point of T_P ?

Example

Compute the least-fixed point of the following Datalog program:

```
p(a) ← q(X)
q(b) ← r(a), p(b)
p(b) ← r(a)
```

defined over the signature $\mathcal{C} = \{a, b, c\}$, $\mathcal{V} = \{X\}$, and $\mathcal{P} = \{p, q, r\}$

Datalog *Inputs*

We can split the set \mathcal{P} of predicates into

1. *input predicates*: predicates that do not appear in the head of rules, and
2. *output predicates*: all remaining predicates

Which are the input/output predicates of this program?

$$\begin{aligned} \textit{anc}(X, Y) &\leftarrow \textit{parent}(X, Y) \\ \textit{anc}(X, Y) &\leftarrow \textit{parent}(X, Z), \textit{anc}(Z, Y) \end{aligned}$$

An *input* for a program P is an interpretation I that contains only atoms constructed using input predicates

The *semantics* of a positive Datalog program P given an *input* I for P is given by the smallest fixed point of T_P that contains I . Let's denote this by $\llbracket P \rrbracket_I$.

How can we compute $\llbracket P \rrbracket_I$?

Datalog and Negation

What should be the semantics of this program?

$$\begin{aligned} p(X) &\leftarrow \neg q(X) \\ q(X) &\leftarrow p(X) \\ p(X) &\leftarrow r(X) \end{aligned}$$

Problem: For a Datalog program P with negation, the consequence operator T_P is not guaranteed to be monotone!

What about the semantics of this program?

$$\begin{aligned} p(X) &\leftarrow \neg q(X) \\ q(X) &\leftarrow r(X) \end{aligned}$$

This Datalog
program is called
“stratified”

Semantics of Stratified Datalog

A **Datalog** program P is **stratified** if its rules can be partitioned into sets P_1, \dots, P_n called strata, such that:

1. for every predicate p , all rules with p in their heads are in one stratum P_i
2. if a predicate symbol p occurs in a positive literal in P_i , then all rules with p in their heads are in a stratum P_j with $j \leq i$
3. if a predicate symbol p occurs in a negative literal in P_i , then all rules with p in their heads are in a stratum P_j with $j < i$

What is an example of a Datalog program that is/is not stratified?

The semantics of a stratified Datalog program P , with strata P_1, \dots, P_n , and an input I for P , is given by the fixed-point M_n where

$M_0 = I$, and $M_i = \llbracket P_i \rrbracket_{M_{i-1}}$, for $i \in [1, n]$.

Is M_n unique for any stratified Datalog program? What if we partition the rules into different partitions?

Encoding Network Behavior in stratified Datalog

Datalog (2/3 Graph Reachability)

Input

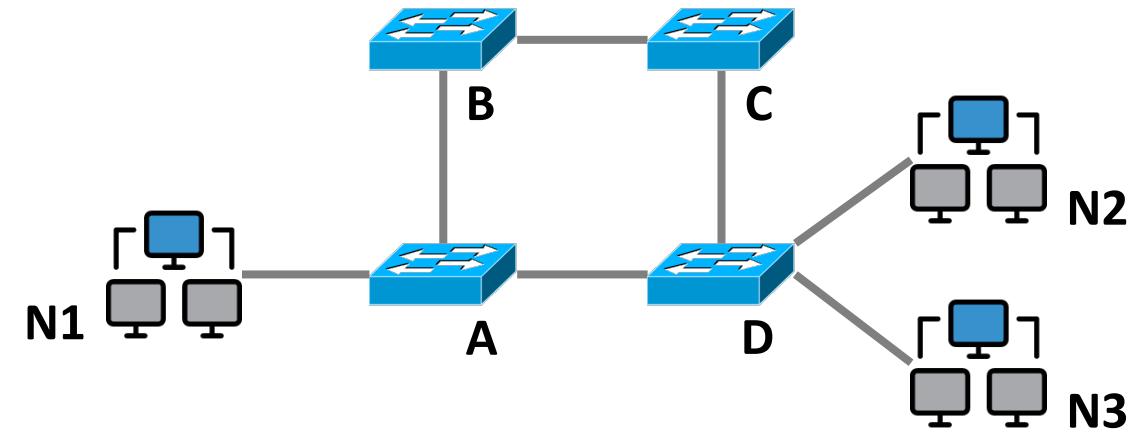
$link(n1, a)$
 $link(a, b)$
...

Program

$path(X, Y) \leftarrow link(X, Y)$
 $path(X, Y) \leftarrow link(X, Z), path(Z, Y)$

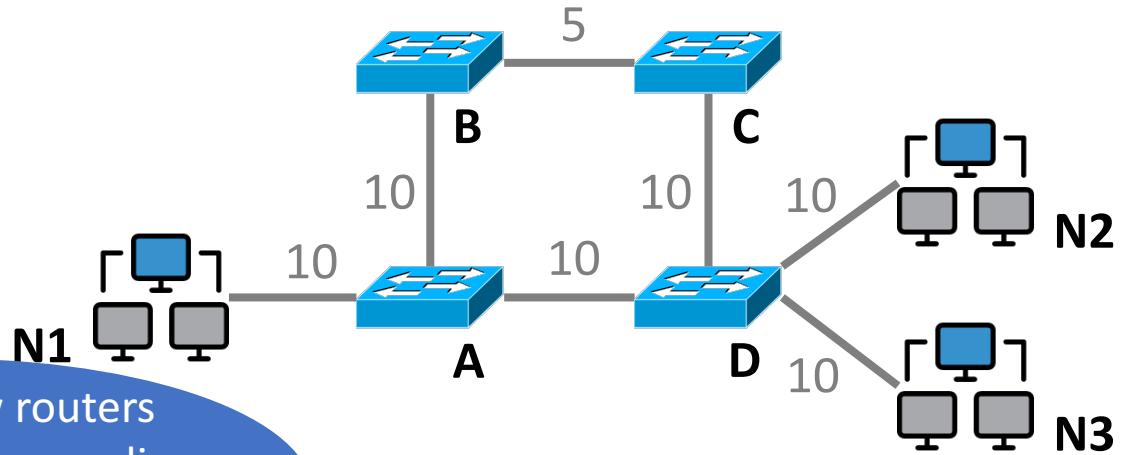
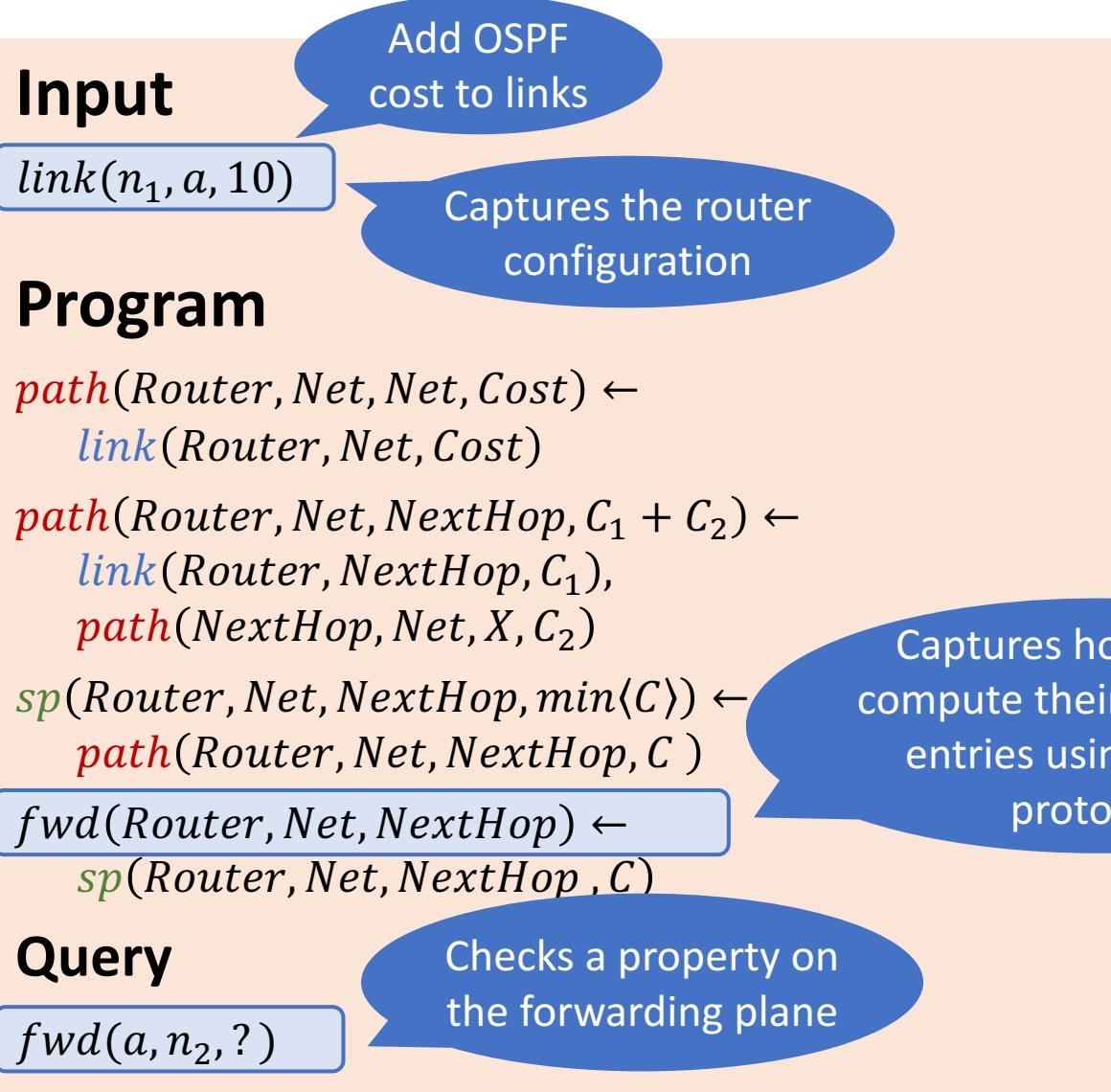
Query

$path(n1, n2)?$



Can we capture the network's forwarding plane?

Datalog (3/3 Shortest-path Routing)



Routing Requirements as Datalog Queries

Paths

Packets for traffic class TC must follow the path

$r_1 \rightarrow \dots \rightarrow r_n$

$fwd(r_1, tc, r_2) \wedge \dots \wedge fwd(r_{n-1}, tc, r_n)$

Routing Requirements as Datalog Queries

Paths

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$r_1 \rightarrow \dots \rightarrow r_n$

$$fwd(r_1, tc, r_2) \wedge \dots \wedge fwd(r_{n-1}, tc, r_n)$$

Traffic isolation

The paths for two distinct traffic classes tc_1 and tc_2 do not share links in the same direction

$$\forall R_1, R_2. fwd(R_1, tc_1, R_2) \Rightarrow \neg fwd(R_1, tc_2, R_2)$$

Routing Requirements as Datalog Queries

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$$\forall R_1, R_2. fwd(R_1, tc_1, R_2) \Rightarrow \neg fwd(R_1, tc_2, R_2)$$

Reachability

Packets for traffic class tc can reach router r_2 from router r_1

$$reach(r_1, tc, r_2)$$

Routing Requirements as Datalog Queries

Paths

Packets for traffic class TC must follow the path

$r_1 \rightarrow \dots \rightarrow r_n$

$$fwd(r_1, tc, r_2) \wedge \dots \wedge fwd(r_{n-1}, tc, r_n)$$

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Reachability

Packets for traffic class tc can reach router r_2 from router r_1

$$reach(r_1, tc, r_2)$$

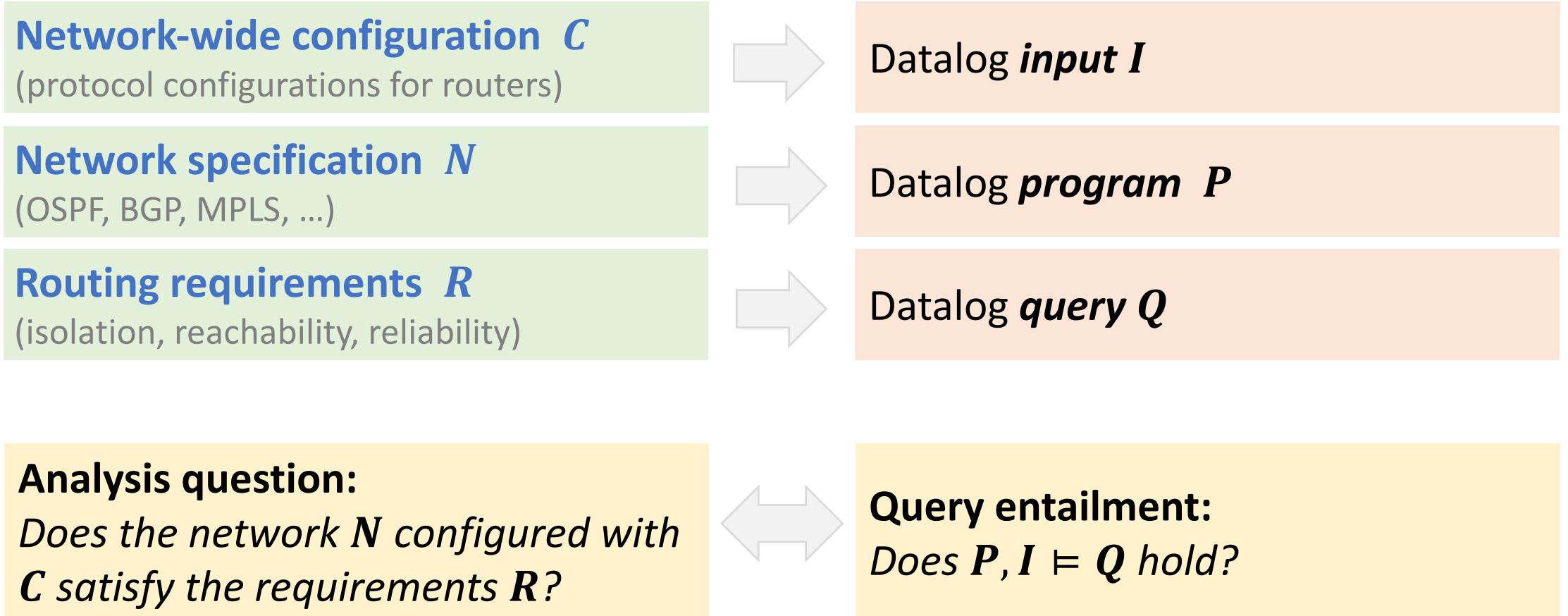
Loop-freeness

The forwarding plane has no loops

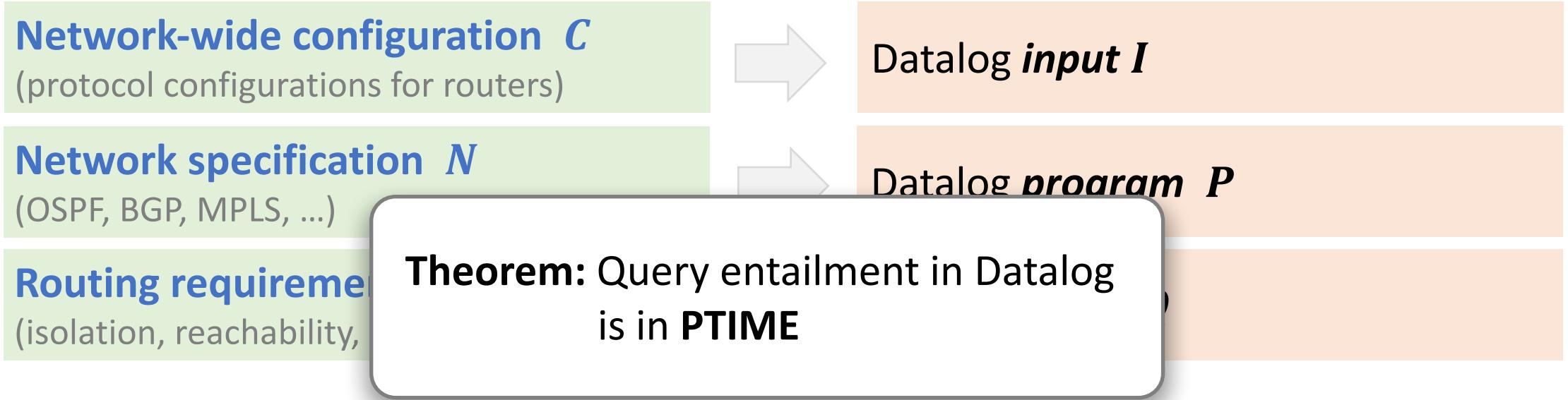
$$\forall TC, R. \neg reach(R, TC, R)$$

Analysis of Network Configurations in Datalog

Analysis of Network Configurations in Datalog



Analysis of Network Configurations in Datalog

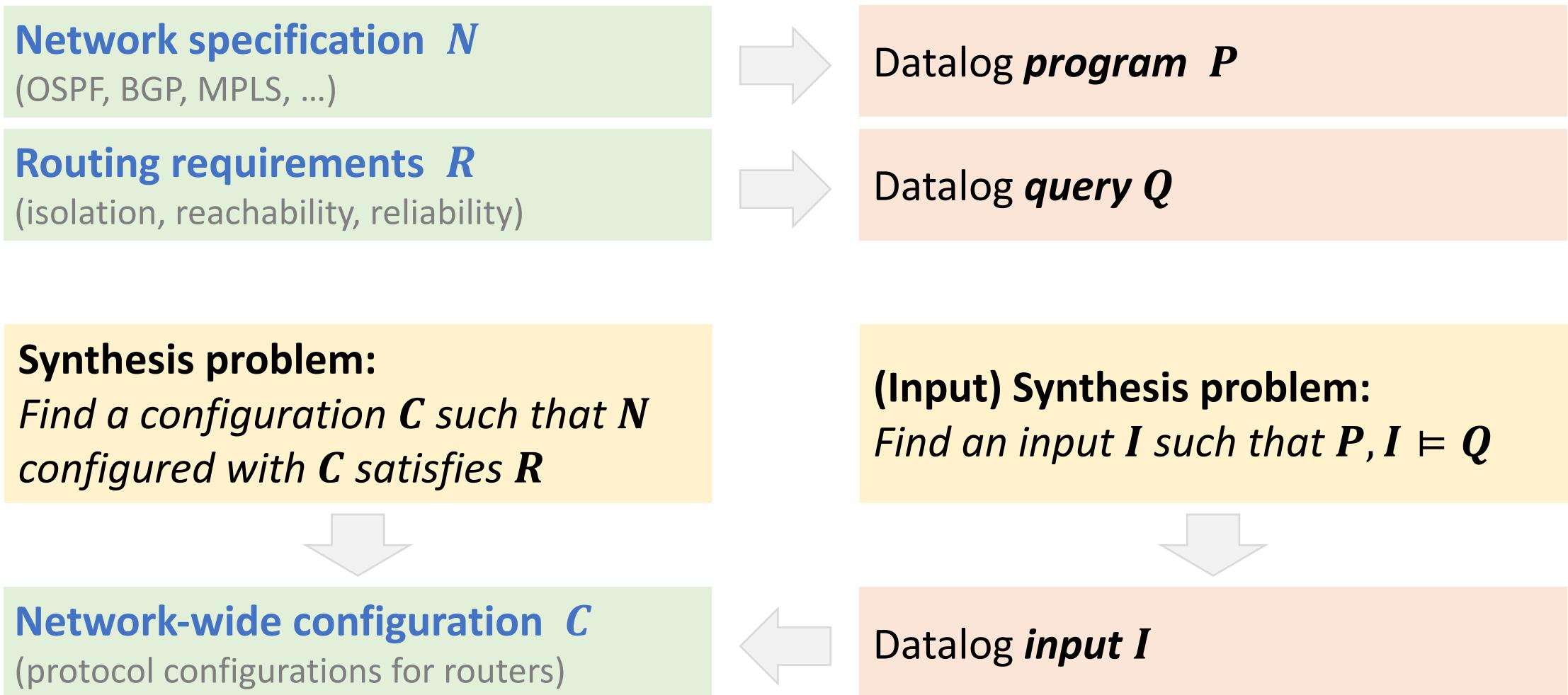


Analysis question:
Does the network N configured with C satisfy the requirements R ?

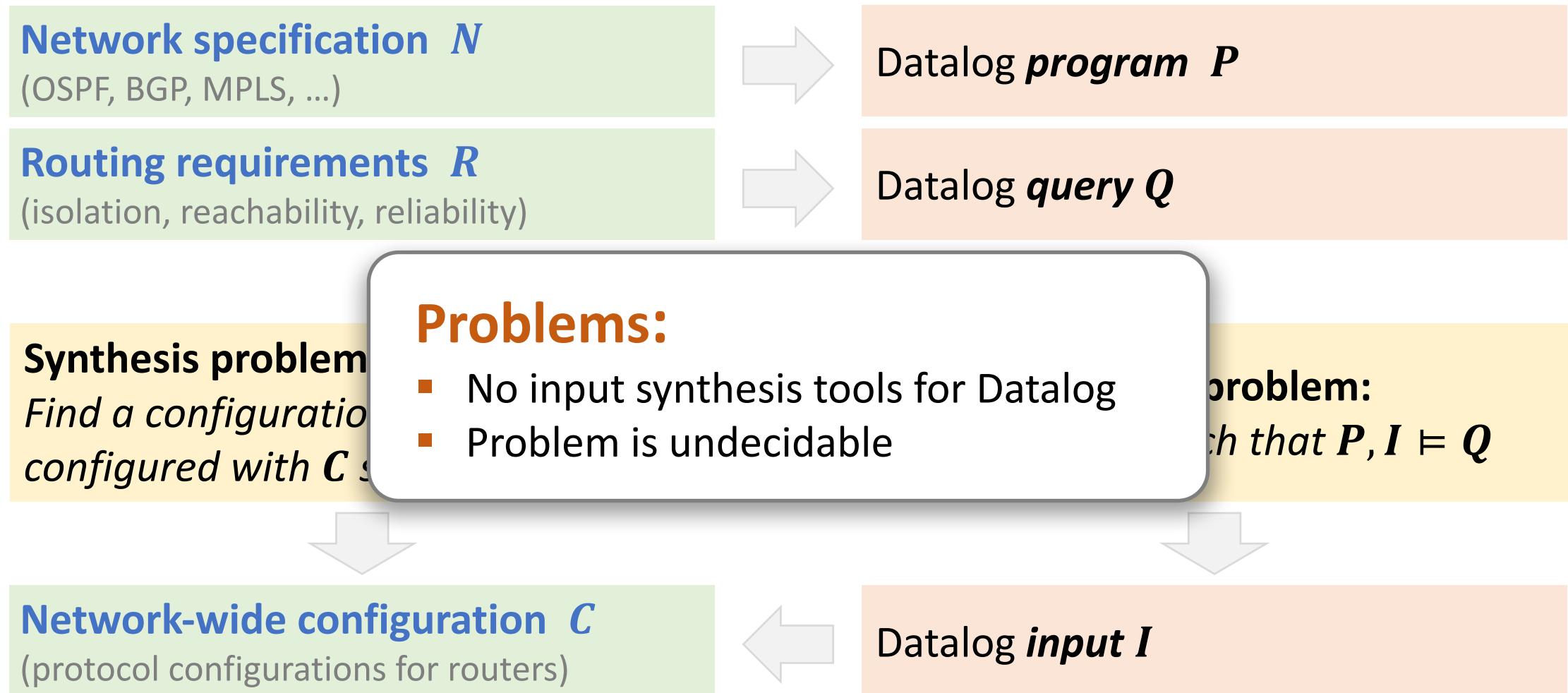
Query entailment:
Does $P, I \models Q$ hold?

Network-wide Configuration Synthesis

Network-wide Configuration Synthesis



Network-wide Configuration Synthesis



Input Synthesis for Datalog



Key idea: Reduce to solving SMT constraints

Input Synthesis for Positive Datalog (First Attempt)

Datalog program P

```
path(X, Y) ← link(X, Y)  
path(X, Y) ← link(X, Z), path(Z, Y)
```

Datalog query Q

```
path(a, c) ∧ ¬link(a, c)
```



Generate SMT constraints

```
∀X, Y. path(X, Y) ⊜ link(X, Y)  
∀X, Y. path(X, Y) ⊜ ∃Z. link(X, Z) ∧ path(Z, Y)  
path(a, c) ∧ ¬link(a, c)
```

Easy encoding...
unfortunately it
does not work

SMT Constraints ψ



Solve SMT constraints

```
path(a, c)
```

Model $M \models \psi$



Derive input (by project on
input predicates)

Datalog input I

\emptyset

Unfortunately,
we get $P, I \not\models Q$

Input Synthesis for Positive Datalog (1/2)

Datalog program P

$$\begin{aligned} path(X, Y) &\leftarrow link(X, Y) \\ path(X, Y) &\leftarrow link(X, Z), path(Z, Y) \end{aligned}$$

Datalog query Q

$$path(a, c) \wedge \neg link(a, c)$$

path(a, c) is a
positive query

Generate SMT constraints

$$\forall X, Y. path_1(X, Y) \Leftrightarrow link(X, Y)$$
$$\forall X, Y. path_2(X, Y) \Leftrightarrow \left(link(X, Y) \vee \left(\exists Z. (link(X, Z) \wedge path_1(Z, Y)) \right) \right)$$
$$path_2(a, c) \wedge \neg link(a, c)$$

Bounded unrolling
for positive queries

SMT Constraints ψ

Solve SMT constraints

$$link(a, b), link(b, c),$$
$$path_1(a, b), path_1(b, c)$$
$$path_2(a, b), path_2(b, c), path_2(a, c)$$

Model $M \models \psi$

Derive input

$$link(a, b), link(b, c)$$

Datalog input I

Input Synthesis for Positive Datalog (2/2)

Datalog program P

$$\begin{aligned} path(X, Y) &\leftarrow link(X, Y) \\ path(X, Y) &\leftarrow link(X, Z), path(Z, Y) \end{aligned}$$

Datalog query Q

$$\neg path(a, c) \wedge \neg link(a, c)$$

*path(a, c) is a
negative query*

Generate SMT constraints

$$\begin{aligned} \forall X, Y. path(X, Y) &\Leftarrow link(X, Y) \\ \forall X, Y. path(X, Y) &\Leftarrow \exists Z. link(X, Z) \wedge path(Z, Y) \\ \neg path(a, c) \wedge \neg link(a, c) \end{aligned}$$

No unrolling for
negative queries

SMT constraints ψ

Solve SMT constraints

{}

Model $M \models \psi$

Derive input

{}

Datalog input I

Input Synthesis for Positive Datalog (2/2)

Datalog program P

$$\begin{aligned} path(X, Y) &\leftarrow link(X, Y) \\ path(X, Y) &\leftarrow link(X, Z), path(Z, Y) \end{aligned}$$

Datalog query Q

$$\neg path(a, c) \wedge \neg link(a, c)$$

path(a, c) is a
negative query



Generate SMT constraints


$$\begin{aligned} \forall X, Y. path(X, Y) &\Leftarrow link(X, Y) \\ \forall X, Y. path(X, Y) &\Leftarrow \exists Z. link(X, Z) \wedge path(Z, Y) \\ \neg path(a, c) \wedge \neg link(a, c) \end{aligned}$$

No unrolling for
negative queries

SMT Constraints ψ



Solve SMT constraints

{}

Model $M \models \psi$



Derive input

?

Datalog input I

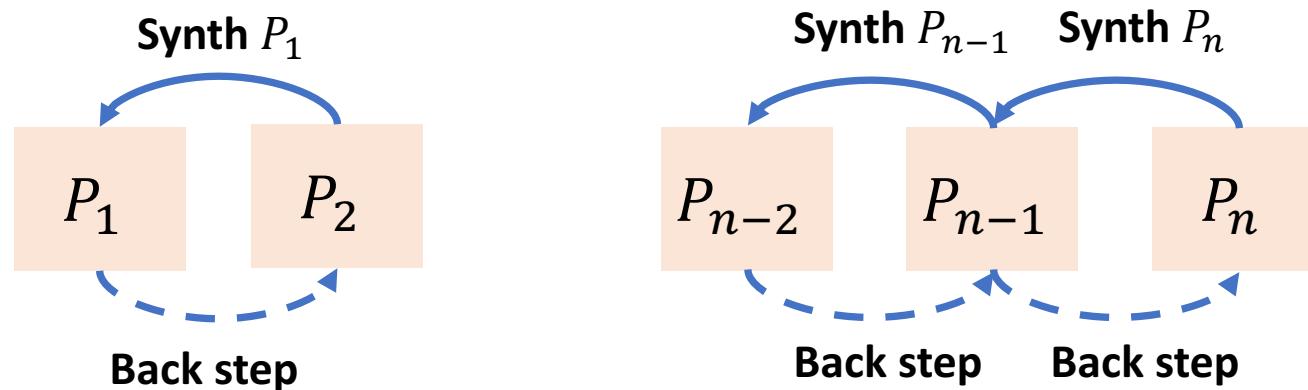
Summary:

- Unroll rules for positive queries, do not unrolling rules for negative queries
- Combine both kinds of constraints for constraints that contain both positive/negative queries.

Input Synthesis for *Stratified* Datalog

Suppose we have a program P with strata P_1, \dots, P_n , and a query Q .

High-level idea:

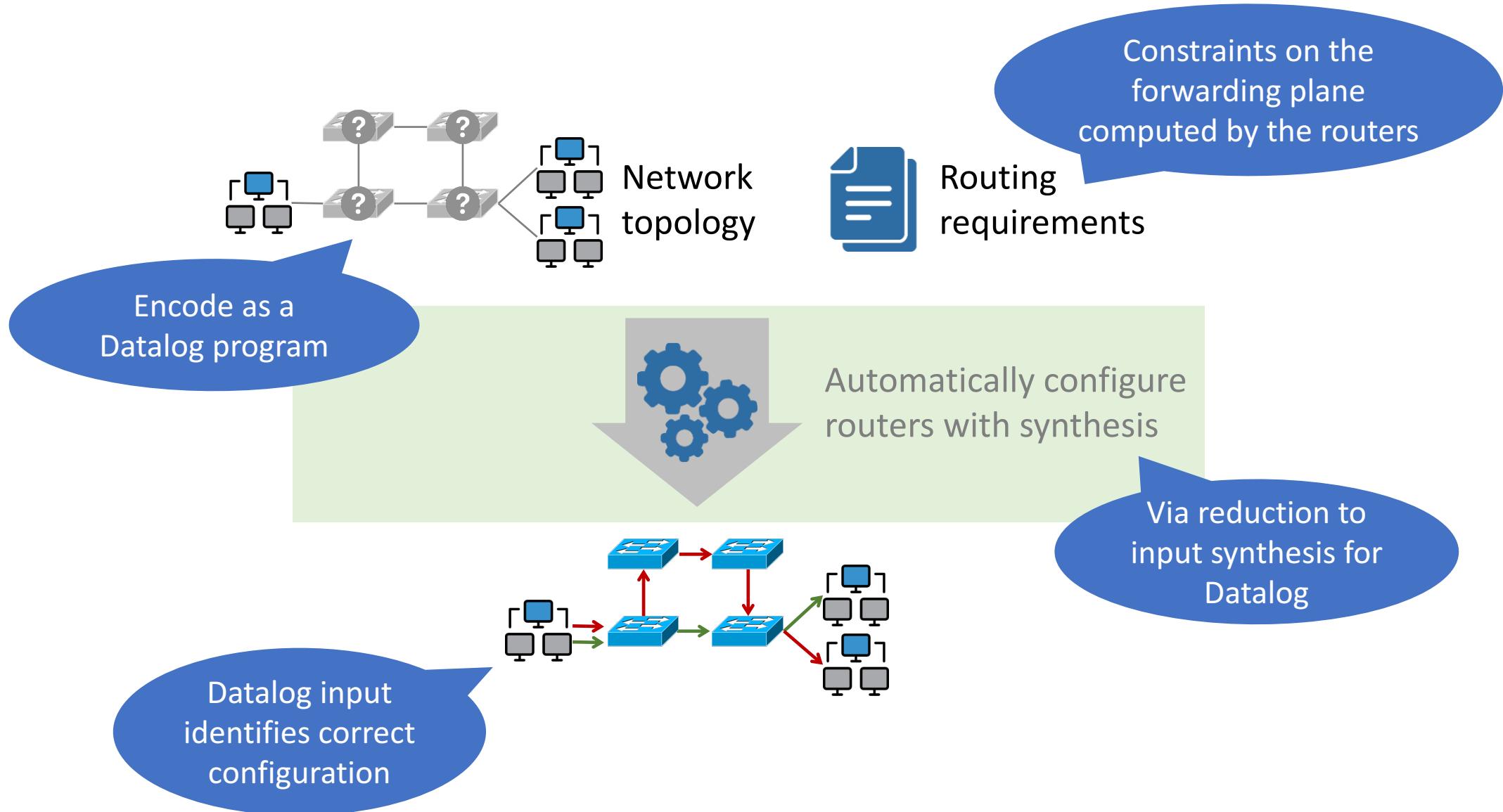


Synth P_n : Compute input I_n for stratum P_n such that $\llbracket P_n \rrbracket_{I_n}$ satisfies Q .

Synth $P_{n-1} \dots P_1$: Compute input I_i for stratum P_i such that $\llbracket P_i \rrbracket_{I_i}$ produces the input I_{i+1} synthesized by the previous step

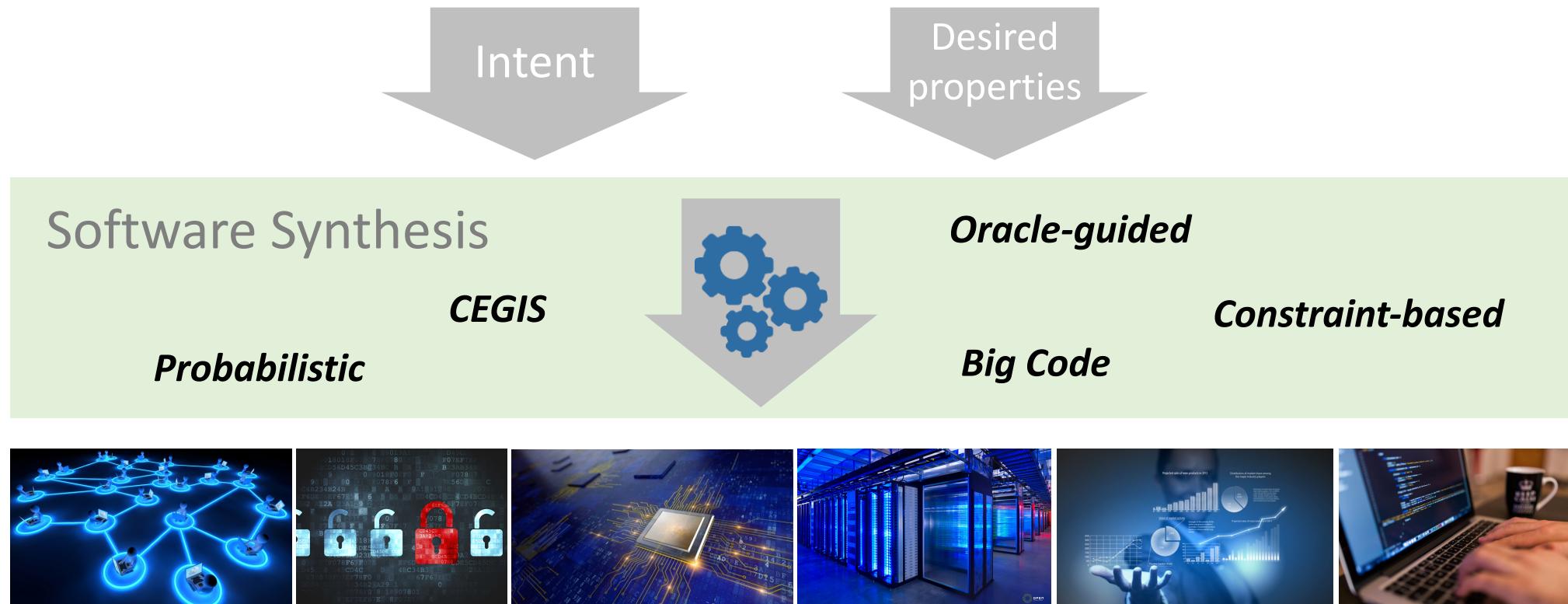
Back step: Backtrack to step **Synth P_i** if the step **Synth P_{i-1}** returns **unsat**

Network-wide Configuration Synthesis



Software Synthesis @ SRL

Develop new synthesis techniques to solve practical system challenges



Computer networks

Security and privacy

Modern architectures

Datacenters

Data science

End-user programming

Application domains

Implementation

Implementation

The SyNET system (<http://synet.ethz.ch>)

- Written in **Python** ($\approx 4K$ lines of code)
- Protocols encoded in **stratified Datalog** (≈ 100 rules)
- Uses the **Z3** constraint solver. Relies on linear integer arithmetic theories (**LIA**).
- Outputs **CISCO** configurations
- Supports **BGP**, **OSPF**, and **static routes**

Network-specific optimizations

- Partial evaluator for Datalog
- Protocol-specific constraints

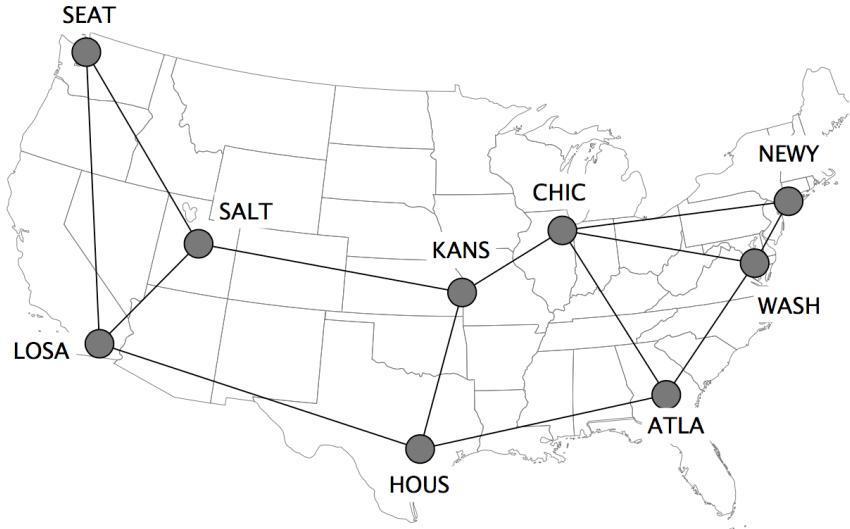
```
! A snippet from router A
interface f0/1
    ip address 10.0.0.2 255.255.255.254
    ip ospf cost 10
    description "To B"
interface f0/0
    ip address 10.0.0.0 255.255.255.254
    ip ospf cost 65530
    description "To C"
interface f1/0
    ip address 10.0.0.4 255.255.255.254
    ip ospf cost 65530
    description "To D"
!
```

Sample **CISCO** configuration
output by SyNET

Experiments



Experiment



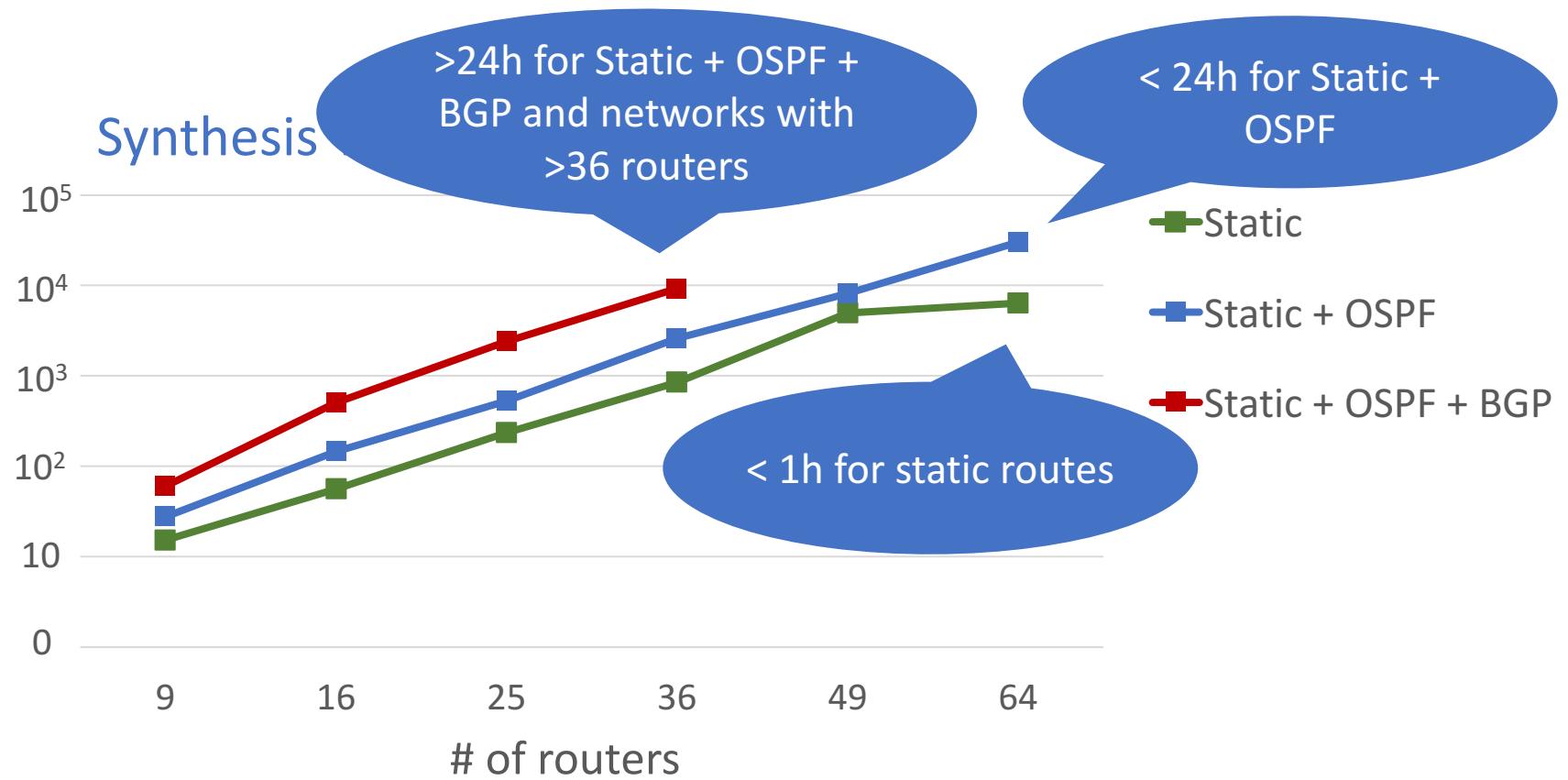
US-based network connecting major universities and research institutes

Protocols / # Traffic classes	1 class	5 classes	10 classes
Static	1.3s	2.0s	4.0s
Static + OSPF	9.0s	21.3s	49.3s
Static + OSPF + BGP	13.3s	22.7s	1m19.7s

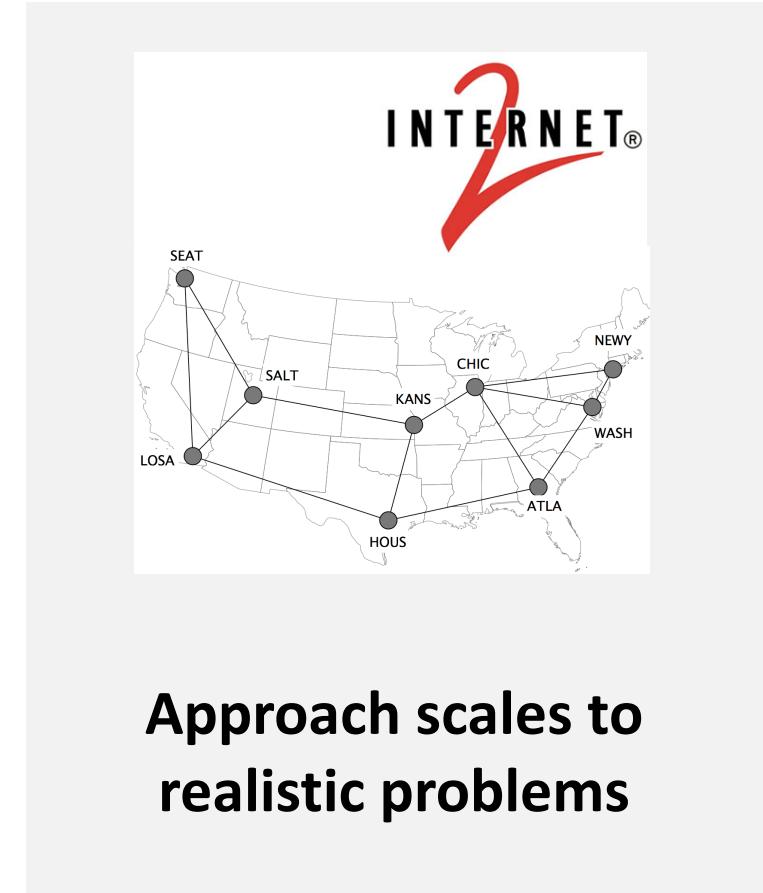
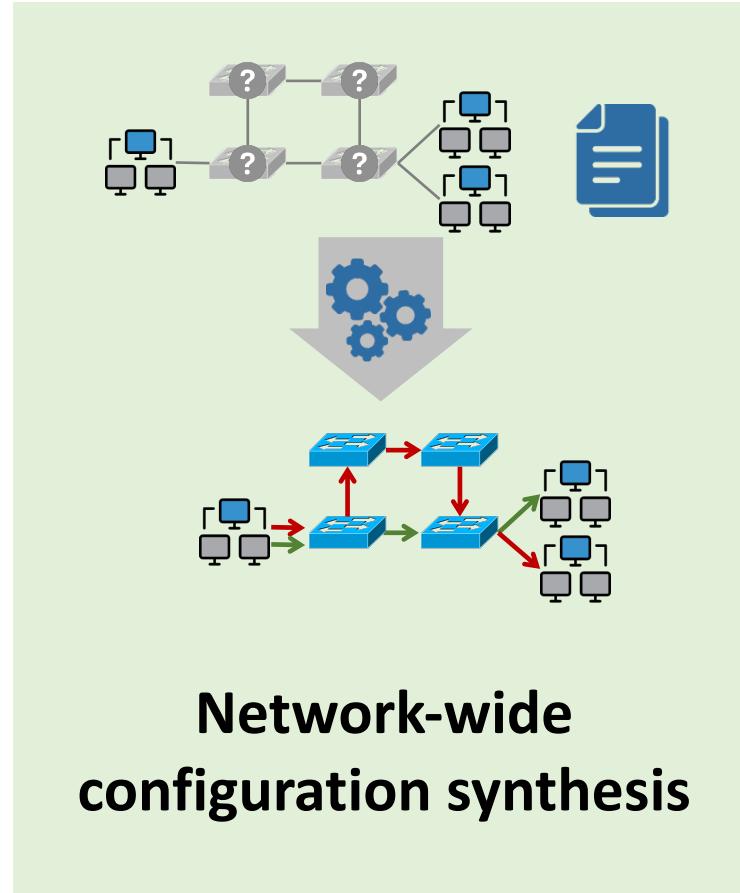
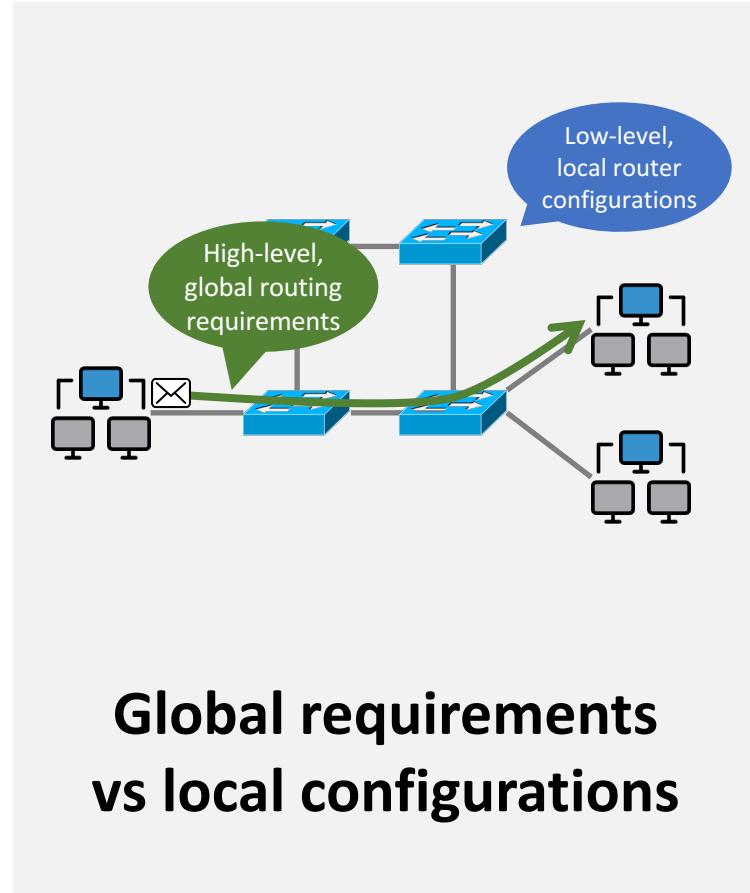
Synthesis Times

Scalability Experiment

- Grid topologies with up to 64 routers
- Requirements for 10 traffic classes



Summary: Programmable Networks with Synthesis



For more details read this paper: <https://arxiv.org/abs/1611.02537>