FAKULTÄT FÜR !NFORMATIK
Faculty of Informatics

Diplomarbeitspräsentation

Interpolation in First-Order Logic with Equality

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Craig Interpolation

Theorem (Craig). *Let A and B be first-order formulas such that $\models A \supset B$. Then there is an interpolant I for A and B such that:*

- $\models A \supset I$
- $\models I \supset B$
- $\mathbf{Lang}(I) \subseteq \mathbf{Lang}(A) \cap \mathbf{Lang}(B)$

$$\begin{tikzcd} A \arrow[rr, Rightarrow] \arrow[dr, Rightarrow] & & B \\ & I \arrow[ur, Rightarrow] & \end{tikzcd}$$

$\Rightarrow$ Interpolants give a concise logical summary of the implication

Applications of Craig Interpolation

Theoretical:
- Proof of Beth's Definability Theorem

Practical:
- Program analysis: Detect loop invariants
- Model checking: Overapproximate set of reachable states

Aim and Scope of the Thesis

Give comprehensive account of existing techniques and extend them:
- Model-theoretic proof
- Reduction to first-order logic without equality
- Interpolant extraction from resolution proofs

Model-theoretic proof

- Non-constructive proof:
  - Let T_A and $T_{\neg B}$ be theories extending A and $\neg B$
  - Build model from maximal consistent intersection of T_A and $T_{\neg B}$ (assuming the non-existence of interpolants)
    $\Rightarrow A \wedge \neg B$ satisfiable
- Related to Robinson's Joint Consistency Theorem

Reduction to first-order logic without equality [?]

Translate equality and function symbols:

$$\begin{aligned} \big(P(c)\big)^* &\equiv \exists x(C(x) \wedge P(x)) \\ \big(P(f(c))\big)^* &\equiv \exists x(\exists y(C(y) \wedge F(y,x)) \wedge P(x)) \\ \big(s = t\big)^* &\equiv E(s,t) \end{aligned}$$

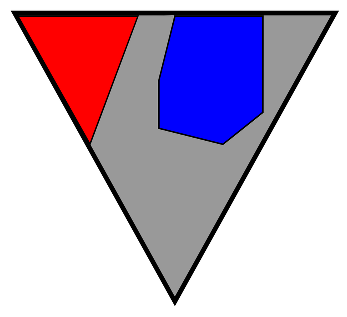
Add theory of equality:

$$\varphi \quad \to \quad T_E \supset \varphi^*$$

$\Rightarrow$ Then calculate interpolant in reduced logic

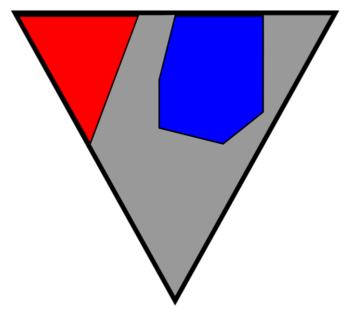
Interpolant extraction from proofs in two phases [?]

Proof:



$\Downarrow$ *Extract propositional interpolant structure from proof*

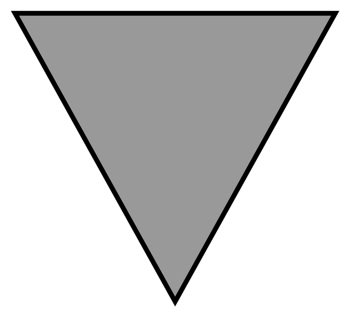
Propositional Interpolant:



$\ldots Q(f(c), c) \ldots$

$\Downarrow$ *Replace colored function and constant symbols*

Prenex First-Order Interpolant:



$\exists x_3 \forall x_5 \ldots Q(x_5, x_3) \ldots$

Interpolant extraction from proofs in one phase

- Replace colored terms during extraction

$\ldots \forall x_5 \ldots Q(x_5, c) \ldots$

$\exists x_3 \ldots \forall x_5 \ldots Q(x_5, x_3) \ldots$

Contributions

- We introduced the one phase-approach.
- We showed that the number of quantifier alternations in the interpolant essentially corresponds to the number of color alternations in terms.

References

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