

undirected edges (from  $\mathcal{M}$ ) are to be interpreted as two directed edges.

$$E(C) = \mathcal{A}(C) \cup \mathcal{M}(C)$$

$$V(C) = V(E(C))$$

$$G(C) = (V(C), E(C))$$

color of component is color of some term in it (all the same)

per resolution step: oppositely colored components are not unifiable

## Components

nodes: max col term occurrences and variables in grey occurrences.

1. components initially: for every variable, all grey occurrences and all colored occurrences
2. resolution: components of  $C_1$  and  $C_2$  are carried over, some are merged.

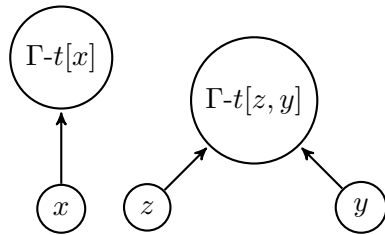
$$l\sigma = l'\sigma$$

For each max col term  $t$  in  $l\sigma$ : merge component of  $t$  and  $t'$ .

quantifier ordering: Build  $\mathcal{A}(C)$ , which is the condensation of  $G(C)$ . If in the condensation there is a path from a node containing a term containing  $u_i$  to a node containing term containing  $u_j$ , then  $u_i <_{\mathcal{A}(C)} u_j$ .

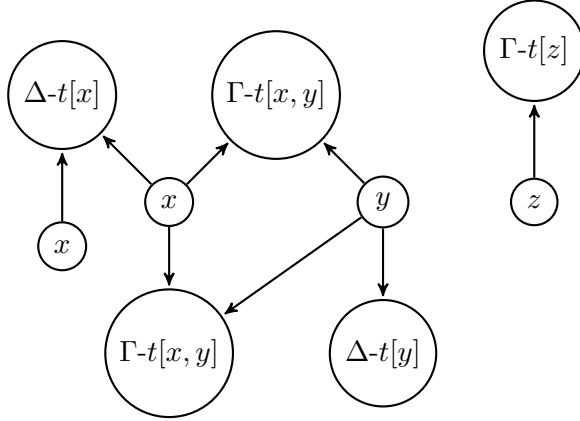
## graph components visualised

initially



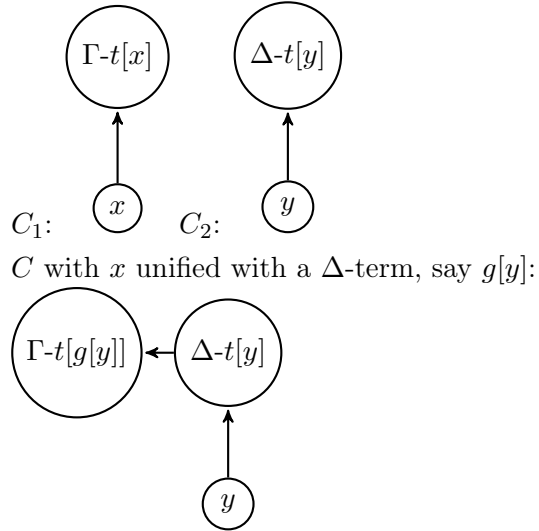
Note: initially, all colored terms are in one component

in the derivation, single color



Note:  $\Gamma$ - and  $\Delta$ -terms can not be merged (unified). All other combinations are possible.

in the derivation



proofs

**Conjecture 1** (Lemma 1). *continue with other proof before fixing this*

*If a variable has a grey and a colored occurrence, there is an arrow between the component containing the grey occurrence and the component containing the colored occurrence (in the condensation graph).*

*Proof.* Induction start is easy.

Induction step: Suppose the statement holds  $C_1$  and  $C_2$ .

Note that no new grey components can be added, just merged. Hence it suffices to show that the component of an arbitrary in  $C$  newly added colored occurrence of a variable, say  $x$ , has an arrow starting from a grey component.

No component is ever added.  $x$  is not unified as otherwise it would not exist anymore (the lemma statement requires the variable to occur). A new colored occurrence of  $x$  can be created by either putting  $x$  into a colored term or by a colored occurrence of  $x$  in the codomain of the unifier.

1. Putting  $x$  into a colored term. Then there is some  $\gamma[y]$  with  $y\sigma = t[x]$ . In the easy case,  $y$  is just unified with  $t[x]$ . Let  $\hat{y}$  be the occurrence of  $y$  in the resolved literal which causes a change of  $y$  in the unification algorithm and  $\hat{t}[x]$  the corresponding term at the same position in the other resolved literal.

Then the component of  $\hat{y}$  is merged with the component of  $\hat{t}[x]$ .

Afterwards, we have some other component of  $x$  as well. This could be:

- a) in the same clause as  $\hat{t}[x]$ .

Then distinguish on the the shape of  $\hat{t}[x]$ :

- Either it is grey, then  $\hat{y}$  is grey as well and we have an arrow by the induction hypothesis from  $\hat{y}$  to  $\gamma[y]$ .
- Otherwise it is colored. Then by the induction hypothesis, as there exists a grey component of  $x$  in this clause, there is an arrow to  $\hat{t}[x]$ . By some Lemma yet to define, there either is a merge arrow between  $\gamma[y]$  and  $\hat{y}$ , which is also a colored term, or there is a grey occurrence of  $y$  with arrows to the two colored occurrences. in the first case, we are done, and in the second ???

TODO: ICI

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“induction”  
until we  
hit a term  
contain-  
ing  $x$

TODO: ICI

TODO: ICI

TODO: ICI

?? as some var which is unified to  $x$  ???

- b) in the same clause as  $\hat{y}$  in the form of a component which is called  $y$  in  $C_i$  for some  $i$ . By the subst, it's now  $x$ . Then we have an arrow by the induction hypothesis from the component of  $y$  to the component of  $\gamma[y]$ .

TODO: what if  $y$  is substituted by a colored term containing  $x$ ?

□

**Conjecture 2** (Conjecture 4). *Suppose in  $\text{AI}^\Delta(C)$  a maximal  $\Gamma$ -term  $\gamma_j[z_i]$  contains a lifting variable  $z_i$ . Then  $z_i <_{\mathcal{A}(C)} z_j$ .*

## random notes

1. if two variable-nodes in the condensation are connected when disregarding the arrow direction, they occur in the same clause.