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ma:lifting_order_not_relevant⟩ **Lemma 1.** *Basically* $\ell_\Gamma^y[\ell_\Delta[\varphi]] = \ell_\Delta[\ell_\Gamma^y[\varphi]]$.

0.2 proof

Definition 2 ($\tau(\iota)$)**.** For an inference ι with $\sigma = \mathrm{mgu}(\iota)$, we define the infinite substitution $\tau(\iota)$ with $\mathrm{dom}(\tau(\iota)) = \mathrm{dom}(\sigma) \cup \{z_s \mid s\sigma \neq s\}$ as follows for a variable x:

define infinite substitutions properly and apply definition here

$$x\tau(\iota) = \begin{cases} x\sigma & x \text{ is a non-lifting variable} \\ z_{t\sigma} & x \text{ is a lifting variable } z_t \end{cases}$$

If the inference ι is clear from the context, we abbreviate $\tau(\iota)$ by τ. △

Definition 3 (Incremental lifting)**.** Let π be a resolution refutation of $\Gamma \cup \Delta$. We define $\mathrm{LI}(\pi)$ ($\mathrm{LI}_{\mathrm{cl}}(\pi)$) to be $\mathrm{LI}(\square)$ ($\mathrm{LI}_{\mathrm{cl}}(\square)$), where $\square$ is the empty clause derived in π.

Let C be a clause in π. For a literal λ in C, we denote the corresponding literal in $\mathrm{LI}_{\mathrm{cl}}(C)$ by λ_{LIcl}, which is exists by lemma 4.

We define $\mathrm{LI}(C)$ and $\mathrm{LI}_{\mathrm{cl}}(C)$ as follows:

Base case. If $C \in \Gamma$, $\mathrm{LI}(C) \stackrel{\mathrm{def}}{=} \bot$. If otherwise $C \in \Delta$, $\mathrm{LI}(C) \stackrel{\mathrm{def}}{=} \top$.

In any case, $\mathrm{LI}_{\mathrm{cl}}(C) \stackrel{\mathrm{def}}{=} \ell[C]$.

Resolution. If the clause C is the result of a resolution step ι of $C_1 : D \vee l$ and $C_2 : E \vee \neg l'$ using a unifier σ such that $l\sigma = l'\sigma$, then define $\mathrm{LI}(C)$ and $\mathrm{LI}_{\mathrm{cl}}(C)$ as follows:

$$\mathrm{LI}_{\mathrm{cl}}(C) \stackrel{\mathrm{def}}{=} \ell[(\mathrm{LI}_{\mathrm{cl}}(C_1)\backslash\{l_{\mathrm{LIcl}}\})\tau] \;\vee\; \ell[(\mathrm{LI}_{\mathrm{cl}}(C_2)\backslash\{l'_{\mathrm{LIcl}}\})\tau]$$

1. If l is Γ-colored: $\mathrm{LI}(C) \stackrel{\mathrm{def}}{=} \ell[\mathrm{LI}(C_1)\tau] \vee \ell[\mathrm{LI}(C_2)\tau]$
2. If l is Δ-colored: $\mathrm{LI}(C) \stackrel{\mathrm{def}}{=} \ell[\mathrm{LI}(C_1)\tau] \wedge \ell[\mathrm{LI}(C_2)\tau]$
3. If l is grey: $\mathrm{LI}(C) \stackrel{\mathrm{def}}{=} (\ell[l_{\mathrm{LIcl}}\tau] \wedge \ell[\mathrm{LI}(C_2)\tau]) \vee (\neg\ell[l'_{\mathrm{LIcl}}\tau] \wedge \ell[\mathrm{LI}(C_1)\tau])$

Factorisation. If the clause C is the result of a factorisation step ι of $C_1 : l \vee l' \vee D$ using a unifier σ such that $l\sigma = l'\sigma$, then $\mathrm{LI}(C) \stackrel{\mathrm{def}}{=} \ell[\mathrm{LI}(C_1)\tau]$ and $\mathrm{LI}_{\mathrm{cl}}(C) \stackrel{\mathrm{def}}{=} \ell[(\mathrm{LI}_{\mathrm{cl}}(C_1)\backslash\{l'_{\mathrm{LIcl}}\})\tau]$. △

vs_clause_plus_literals_equal⟩ **Lemma 4.** *Let C be a clause in a resolution refutation of $\Gamma \cup \Delta$.*

Then for every literal λ in C, there exists a literal λ_{LIcl} in $\mathrm{LI}_{\mathrm{cl}}(C)$ such that $\lambda_{\mathrm{LIcl}} = \ell[\lambda]$ and for resolved or factorised literals l and l' of a resolution or factorisation inference ι, we have that $\ell[l_{\mathrm{LIcl}}\tau] = \ell[l'_{\mathrm{LIcl}}\tau]$.

Proof. We proceed by induction.

Base case. For $C \in \Gamma \cup \Delta$, $\mathrm{LI}_{\mathrm{cl}}(C)$ is defined to be $\ell[C]$.

Resolution/Factorisation. Suppose the clause C is the result of a resolution or factorisation inference ι of the clauses $\bar{C}$ with $\sigma = \mathrm{mgu}(\iota)$.

For every literal in C, there exists a predecessor in a clause in $\bar{C}$. Let λ be a literal C_i with $C_i \in \bar{C}$, such that λ is not the predecessor of the literal being resolved or factorised in ι. Then $\lambda\sigma$ is occurs in C.

By the induction hypothesis, $\ell[\lambda]$ occurs in $\mathrm{LI}_{\mathrm{cl}}(C_i)$. The successor of $\ell[\lambda]$ in $\mathrm{LI}_{\mathrm{cl}}(C)$ is of the form $\ell[\ell[\lambda]\tau]$. But by Lemma 10, this is nothing else than $\ell[\lambda\tau]$. As no lifting variables occur in λ, we get by the definition of τ that $\ell[\lambda\tau] = \ell[\lambda\sigma]$.

Let l and l' be the resolved or factorised literals. In order to show that $\ell[l_{\mathrm{LIcl}}\tau] = \ell[l'_{\mathrm{LIcl}}\tau]$, consider that by the induction hypothesis, this is nothing else than $\ell[\ell[l]\tau] = \ell[\ell[l']\tau]$. But by applying the same argument as above, this is equivalent to $\ell[l\sigma] = \ell[l'\sigma]$, which is implied by $l\sigma = l'\sigma$. □

Definition 5. $\mathrm{LI}^{\Delta}(C)$ ($\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)$) for a clause C is defined as $\mathrm{LI}(C)$ ($\mathrm{LI}_{\mathrm{cl}}(C)$) with the difference that in its inductive definition, every lifting $\ell[\varphi]$ for a formula or term φ is replaced by a lifting of only the Δ-terms $\ell_{\Delta}[\varphi]$. △

Remark. Many results involving $\mathrm{LI}(C)$ ($\mathrm{LI}_{\mathrm{cl}}(C)$) are valid for $\mathrm{LI}^{\Delta}(C)$ ($\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)$) in a formulation which is adapted accordingly. This can easily be seen by the following proof idea:

Let $f_1, \dots, f_n$ be all Γ-colored function or constant symbols, c a fresh constant symbol and g be a fresh n-ary function symbol. Construct a formula $t = t$ such that $t = g(t_1, \dots, t_n)$, such that $t_i = f_i(c_1, \dots, c_m)$ for $1 \leq i \leq n$ where m is the arity of f_i and $c_j = c$ for $1 \leq j \leq m$. Let $\Delta' = \Delta$ and apply the desired result to the initial clause sets Γ and Δ'.

Under this construction, every originally Γ-colored symbol is now grey, which implies that $\mathrm{LI}(C) = \mathrm{LI}^{\Delta}(C)$ as well as $\mathrm{LI}_{\mathrm{cl}}(C) = \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)$. But $\Delta \models \varphi \Leftrightarrow \Delta' \models \varphi$ for any formula φ. △

⟨lemma:no_colored_terms⟩ **Lemma 6.** *Let C be a clause of a resolution refutation π of $\Gamma \cup \Delta$. $\mathrm{LI}(C)$ and $\mathrm{LI}_{\mathrm{cl}}(C)$ do not contain colored symbols. $\mathrm{LI}^{\Delta}(C)$ and $\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)$ do not contain Δ-colored symbols.*

Proof. For $\mathrm{LI}(C)$ and $\mathrm{LI}_{\mathrm{cl}}(C)$, consider the following: In the base case of the inductive definitions of $\mathrm{LI}(C)$ and $\mathrm{LI}_{\mathrm{cl}}(C)$, no colored symbols occur. In the inductive steps, any colored symbol which is added by τ to intermediary formulas is lifted.

For $\mathrm{LI}^{\Delta}(C)$ and $\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)$, a similar argument goes through by reading colored as Δ-colored. □

⟨lemma:substitute_and_lift⟩ **Lemma 7.** *Let σ be a substitution and F a formula without Φ-colored terms such that for a set of formulas Ψ, $\Psi \models F$. Then $\Psi \models \ell_{\Phi}[F\sigma]$.*

Proof. $\ell_{\Phi}[F\sigma]$ is an instance of F: σ substitutes variables either for terms not containing Φ-colored symbols or by terms containing Φ-colored symbols. For the first kind, the lifting has no effect. For the latter, the lifting only replaces subterms of the terms introduced by the substitution by a lifting variable such that the original structure of F remains invariant as it by assumption does not contain colored terms. □

⟨lemma:gamma_proves_pide⟩ **Lemma 8.** *For a clause C in a resolution refutation of $\Gamma \cup \Delta$, $\Gamma \models \mathrm{LI}^{\Delta}(C) \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)$.*

Proof. We proceed by induction of the strengthening $\Gamma \models \mathrm{LI}^{\Delta}(C) \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_{\Gamma})$.

Base case. For $C \in \Gamma$, $\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_\Gamma) = \ell_\Delta[C] = C$. Hence $\Gamma \models \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_\Gamma)$.

For $C \in \Delta$, $\mathrm{LI}^{\Delta}(C) = \top$, so $\Gamma \models \mathrm{LI}^{\Delta}(C)$.

Resolution. Suppose the clause C is the result of a resolution step ι of $C_1 : D \vee l$ and $C_2 : E \vee \neg l'$ with $\sigma = \mathrm{mgu}(\iota)$.

We define the following abbreviations:

$\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma)^* = \mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma \backslash \{l_{\mathrm{LIcl}^\Delta}\})$

$\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_2)_\Gamma)^* = \mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_2)_\Gamma \backslash \{\neg l'_{\mathrm{LIcl}^\Delta}\})$

Hence the induction hypothesis can be stated as follows:

$\Gamma \models \mathrm{LI}^{\Delta}(C_1) \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma)^* \vee (l_{\mathrm{LIcl}^\Delta})_\Gamma$

$\Gamma \models \mathrm{LI}^{\Delta}(C_2) \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_2)_\Gamma)^* \vee \neg(l'_{\mathrm{LIcl}^\Delta})_\Gamma$

By Lemma 6, $\mathrm{LI}^{\Delta}(C_i)$ and $\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_i)$ for $i \in \{1, 2\}$ do not contain Δ-colored terms. Hence we are able to apply Lemma 7 in order to obtain

$\Gamma \overset{(\circ)}{\models} \ell_\Delta[\mathrm{LI}^{\Delta}(C_1)\tau] \vee \ell_\Delta[\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma)^*\tau] \vee \ell_\Delta[(l_{\mathrm{LIcl}^\Delta})_\Gamma\tau]$

$\Gamma \overset{(*)}{\models} \ell_\Delta[\mathrm{LI}^{\Delta}(C_2)\tau] \vee \ell_\Delta[\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_2)_\Gamma)^*\tau] \vee \neg\ell_\Delta[(l'_{\mathrm{LIcl}^\Delta})_\Gamma\tau]$

By Lemma 4, we obtain that $\ell_\Delta[l_{\mathrm{LIcl}^\Delta}\tau] = \ell_\Delta[l'_{\mathrm{LIcl}^\Delta}\tau]$.

Now we distinguish cases based on the color of the resolved literal:

- Suppose that l is Γ-colored. Then as $\ell_\Delta[l_{\mathrm{LIcl}^\Delta}\tau] = \ell_\Delta[l'_{\mathrm{LIcl}^\Delta}\tau]$, we can perform a resolution step on $(\circ)$ and $(*)$, which gives that $\Gamma \models \ell_\Delta[\mathrm{LI}^{\Delta}(C_1)\tau] \vee \ell_\Delta[\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma)^*\tau] \vee \ell_\Delta[\mathrm{LI}^{\Delta}(C_2)\tau] \vee \ell_\Delta[\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_2)_\Gamma)^*\tau]$. This however is nothing else than $\Gamma \models \mathrm{LI}^{\Delta}(C) \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)$.
- Suppose that l is Δ-colored. Then $(\circ)$ and $(*)$ simply to the following:

  $\Gamma \models \ell_\Delta[\mathrm{LI}^{\Delta}(C_1)\tau] \vee \ell_\Delta[\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma)^*\tau]$

  $\Gamma \models \ell_\Delta[\mathrm{LI}^{\Delta}(C_2)\tau] \vee \ell_\Delta[\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_2)_\Gamma)^*\tau]$

  These however imply that $\Gamma \models \mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma)^* \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_2)_\Gamma)^* \vee (\ell_\Delta[\mathrm{LI}^{\Delta}(C_1)\tau] \wedge \ell_\Delta[\mathrm{LI}^{\Delta}(C_2)\tau])$, which is nothing else than $\Gamma \models \mathrm{LI}^{\Delta}(C) \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)$.
- Suppose that l is grey. Suppose that M is a model of Γ such that $M \not\models \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)$, i.e. $M \not\models \ell_\Delta[\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma)^*\tau] \vee \ell_\Delta[\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_2)_\Gamma)^*\tau]$. Then $M \models \ell_\Delta[\mathrm{LI}^{\Delta}(C_1)\tau] \vee \ell_\Delta[l_{\mathrm{LIcl}^\Delta}\tau]$ as well as $M \models \ell_\Delta[\mathrm{LI}^{\Delta}(C_2)\tau] \vee \neg\ell_\Delta[l'_{\mathrm{LIcl}^\Delta}\tau]$.

  Due to $\ell_\Delta[l_{\mathrm{LIcl}^\Delta}\tau] = \ell_\Delta[l'_{\mathrm{LIcl}^\Delta}\tau]$, we obtain that

  $M \models (\ell_\Delta[l_{\mathrm{LIcl}^\Delta}\tau] \wedge \ell_\Delta[\mathrm{LI}^{\Delta}(C_2)\tau]) \vee (\neg\ell_\Delta[l'_{\mathrm{LIcl}^\Delta}\tau] \wedge \ell_\Delta[\mathrm{LI}^{\Delta}(C_1)\tau])$, which is nothing else than $M \models \mathrm{LI}^{\Delta}(C)$.

Factorisation. Suppose the clause C is the result of a factorisation inference ι of $C_1 : l \vee l' \vee D$ with $\sigma = \mathrm{mgu}(\iota)$.

We introduce the abbreviation $\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma)^* = \mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)\Gamma \backslash \{l_{\mathrm{LIcl}^\Delta}, \neg l'_{\mathrm{LIcl}^\Delta}\})$ and express the induction hypothesis as follows:

$\Gamma \models \mathrm{LI}^{\Delta}(C_1) \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma)^* \vee (l_{\mathrm{LIcl}^\Delta})_\Gamma \vee \neg(l'_{\mathrm{LIcl}^\Delta})_\Gamma$

By Lemma 6, $\mathrm{LI}^{\Delta}(C_i)$ and $\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_i)$ for $i \in \{1,2\}$ do not contain Δ-colored terms. Hence we are able to apply Lemma 7 in order to obtain

$$\Gamma \overset{(\circledast)}{\vDash} \ell_\Delta[\mathrm{LI}^{\Delta}(C_1)\tau] \vee \ell_\Delta[\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma)^*\tau] \vee \ell_\Delta[(l_{\mathrm{LIcl}^\Delta})_\Gamma\tau] \vee \neg\ell_\Delta[(l'_{\mathrm{LIcl}^\Delta})_\Gamma\tau]$$

As by Lemma 4 we get that $\ell_\Delta[l_{\mathrm{LIcl}^\Delta}\tau] = \ell_\Delta[l'_{\mathrm{LIcl}^\Delta}\tau]$, we can perform a factorisation step on ($\circledast$) to obtain that $\Gamma \vDash \ell_\Delta[\mathrm{LI}^{\Delta}(C_1)\tau] \vee \ell_\Delta[\mathrm{LI}^{\Delta}_{\mathrm{cl}}((C_1)_\Gamma)^*\tau] \vee \ell_\Delta[(l_{\mathrm{LIcl}^\Delta})_\Gamma\tau]$. But this is nothing else than $\Gamma \vDash \mathrm{LI}^{\Delta}(C) \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_\Gamma)$. $\square$

?⟨def:arrow_quantifier_block⟩? **Definition 9** (Quantifier block)**.** Let C be a clause in a resolution refutation π of $\Gamma \cup \Delta$ and $\bar{x}$ be the Δ-lifting variables and $\bar{y}$ the Γ-lifting variables occurring in $\mathrm{LI}(C)$ and $\mathrm{LI}_{\mathrm{cl}}(C)$. $Q(C)$ denotes an arrangement of the elements of $\{\forall x_t \mid x_t \in \bar{x}\} \cup \{\exists y_t \mid y_t \in \bar{y}\}$ such that for two lifting variable z_s and z_r, if s is a subterm of r, then z_s is listed before z_r. We denote $Q(\square)$ by $Q(\pi)$. $\triangle$

⟨lemma:lifting_tau_commute⟩ **Conjectured Lemma 10.** $\ell[\ell[\varphi]\tau] = \ell[\varphi\tau]$.

Proof. proof by induction.
Supp constant: done.
Supp grey function: apply to children.
supp variable: $\ell[\ell[x]\tau] = \ell[x\tau]$
supp lft var: $\ell[\ell[z_t]\tau] = \ell[z_t\tau]$
supp col term t
$\ell[\ell[t]\tau] = \ell[z_t\tau] = \ell[z_{t\sigma}] = z_{t\sigma} = \ell[t\sigma] = \ell[t\tau]$
$\square$

⟨lemma:gamma_lifted_lide⟩ **Lemma 11.** *For a clause C of a resolution refutation of $\Gamma \cup \Delta$, $\ell_\Gamma[\mathrm{LI}^{\Delta}(C) \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)] = \mathrm{LI}(C) \vee \mathrm{LI}_{\mathrm{cl}}(C)$.*

Proof. Base case.

LI^{Δ}: easy.

$\mathrm{LI}^{\Delta}_{\mathrm{cl}}$: By Lemma 1, $\ell_\Gamma[\ell_\Delta[C]] = \ell[C]$

Resolution.

IH:

$\ell_\Gamma[\mathrm{LI}^{\Delta}(C_1) \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_1)] = \mathrm{LI}(C_1) \vee \mathrm{LI}_{\mathrm{cl}}(C_1)$.

$\ell_\Gamma[\mathrm{LI}^{\Delta}(C_2) \vee \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_2)] = \mathrm{LI}(C_2) \vee \mathrm{LI}_{\mathrm{cl}}(C_2)$.

$\mathrm{LI}^{\Delta}_{\mathrm{cl}}$:

$\ell_\Gamma[\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_1)] = \mathrm{LI}_{\mathrm{cl}}(C_1)$

$\ell_\Delta[\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_1)\tau] \subseteq \mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)$

$\ell[\mathrm{LI}_{\mathrm{cl}}(C_1)\tau] \subseteq \mathrm{LI}_{\mathrm{cl}}(C)$

to show: $\ell^y_\Gamma[\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C)] = \mathrm{LI}_{\mathrm{cl}}(C)$

$\ell[\ell_\Gamma[\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_1)]\tau] = \ell[\mathrm{LI}_{\mathrm{cl}}(C_1)\tau]$ IH + same op on both sides

new lemma above

$\ell[\ell_\Gamma[\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_1)]\tau] = \ell[\mathrm{LI}^{\Delta}_{\mathrm{cl}}(C_1)\tau]$

LI^{Δ}:

- Supp Γ:
  IH: $\ell_\Gamma[\mathrm{LI}^\Delta(C_1)] = \mathrm{LI}(C_1)$
  hence also: $\ell[\mathrm{LI}^\Delta(C_1)] = \mathrm{LI}(C_1)$ (by lemma: no Δ-terms in ...)
  $+\ \tau$: $\ell[\mathrm{LI}^\Delta(C_1)]\tau = \mathrm{LI}(C_1)\tau$
  $+\ \ell$: $\ell[\ell[\mathrm{LI}^\Delta(C_1)]\tau] = \ell[\mathrm{LI}(C_1)\tau]$
  by new lemma $\ell[\mathrm{LI}^\Delta(C_1)\tau] = \ell[\mathrm{LI}(C_1)\tau]$
  hence by Lemma 1, $\ell_\Gamma[\ell_\Delta[\mathrm{LI}^\Delta(C_1)\tau]] \subseteq \mathrm{LI}^\Delta(C)$
  hence $\ell_\Gamma[\mathrm{LI}^\Delta(C)] \subseteq \mathrm{LI}^\Delta(C)$

Factorisation.

□

Lemma 12. *For a clause C of a resolution refutation of $\Gamma \cup \Delta$, $\Gamma \models Q(C)(\mathrm{LI}(C) \vee \mathrm{LI}_{\mathrm{cl}}(C))$.*

Proof. By Lemma 11 $\ell_\Gamma[\mathrm{LI}^\Delta(C) \vee \mathrm{LI}^\Delta_{\mathrm{cl}}(C)] = \mathrm{LI}(C) \vee \mathrm{LI}_{\mathrm{cl}}(C)$.

By Lemma 8, $\Gamma \models \mathrm{LI}^\Delta(C) \vee \mathrm{LI}^\Delta_{\mathrm{cl}}(C)$. Hence the terms in $\mathrm{LI}^\Delta(C) \vee \mathrm{LI}^\Delta_{\mathrm{cl}}(C)$ provide witness terms for the Γ-lifting variables in $\mathrm{LI}(C) \vee \mathrm{LI}_{\mathrm{cl}}(C)$, which are existentially quantified in $Q(C)(\mathrm{LI}(C) \vee \mathrm{LI}_{\mathrm{cl}}(C))$.

Furthermore, the ordering imposed on the quantifiers in $Q(C)$ implies that if a Δ-lifting variable x_s occurs in a witness term for a Γ-lifting variable y_r, y_r is quantified in the scope of the quantifier of x_s as s is a subterm of r. This however ensures that the witness terms are valid. □

?⟨lemma:li_symmetry⟩? **Lemma 13.** *symmetry: $Q(C)(\mathrm{LI}(C)) \Leftrightarrow Q(\hat{C})(\mathrm{LI}(\hat{C}))$.*

Proof. todo: copy from other pdf □

Theorem 14. *same as other pdf*