



# Taskflow: A General-purpose Task-parallel Programming System

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<https://taskflow.github.io/>





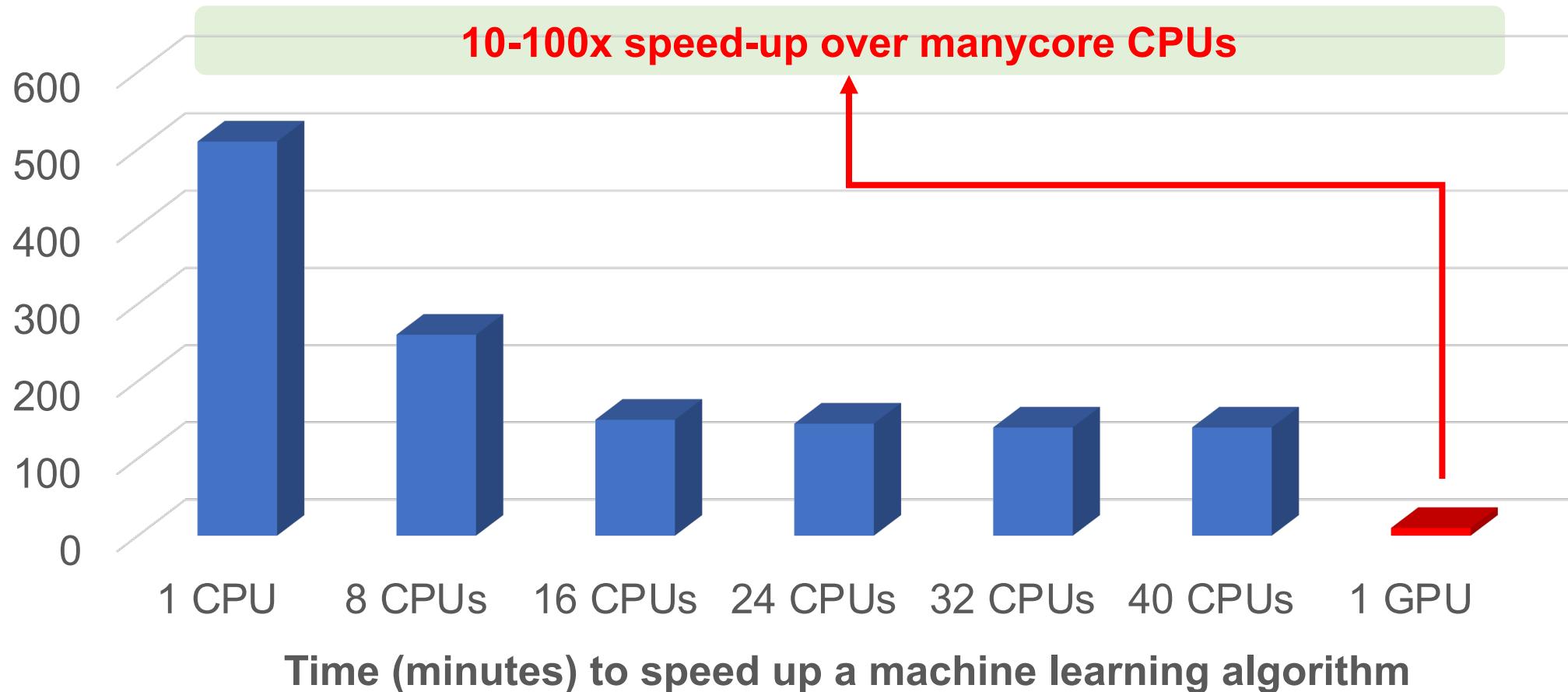
# Takeaways

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- Express your parallelism in the right way
- Program task graph parallelism using Taskflow
- Program dynamic task graph parallelism using Taskflow
- Overcome the scheduling challenges
- Demonstrate the efficiency of Taskflow in industrial application
- Conclude the talk

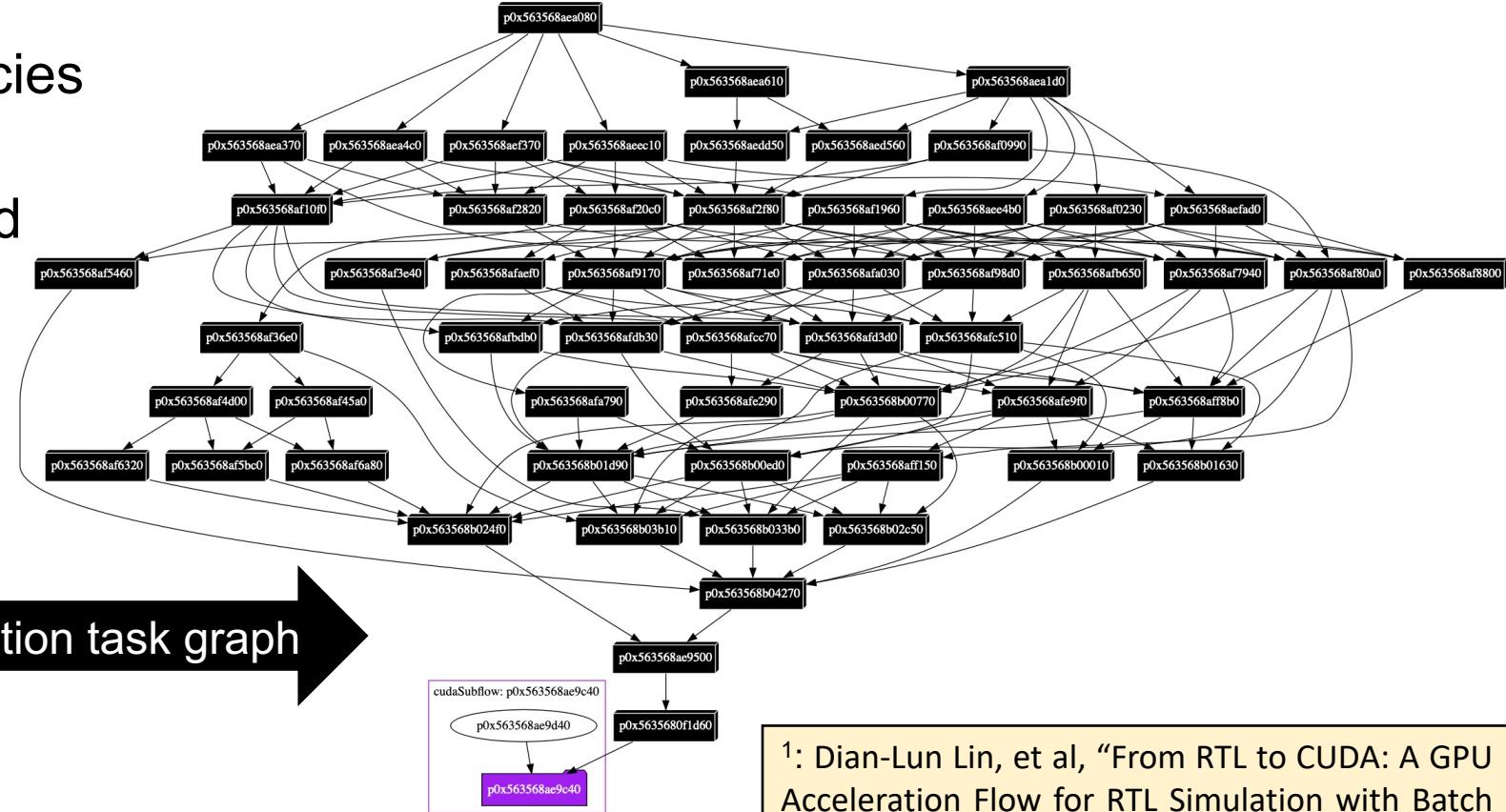
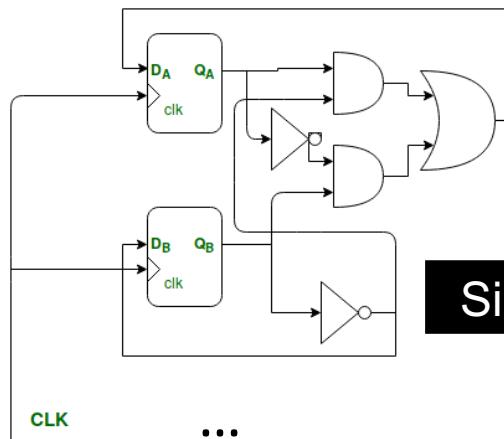
# Why Parallel Computing?

- Advances performance to a new level previously out of reach



# Today's Parallel Workload is Very Complex

- GPU-accelerated circuit simulation on a design of 500M gates<sup>1</sup>
  - >1000 kernels
  - >1000 dependencies
  - >500s to finish
  - >10hrs turnaround



<sup>1</sup>: Dian-Lun Lin, et al, "From RTL to CUDA: A GPU Acceleration Flow for RTL Simulation with Batch Stimulus," ACM ICPP, Bordeaux, France, 2022

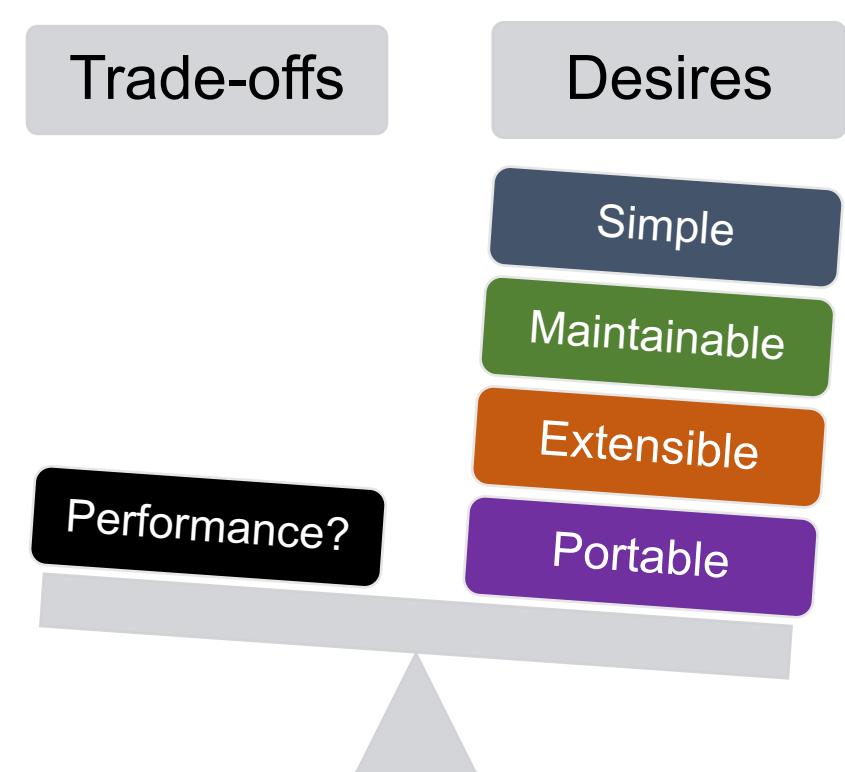
# Parallel Programming is Not Easy

- You need to deal with A LOT OF technical details

- Parallelism abstraction (software + hardware)
- Concurrency control
- Task and data race avoidance
- Dependency constraints
- Scheduling efficiencies (load balancing)
- Performance portability
- ...

- And, don't forget about trade-offs

- Performance vs Desires



# Need a High-level Programming Model

- From user's perspective, the biggest challenge is *transparency*
  - Parallelism abstraction, runtime optimization, load balancing, etc.
- Observing from the evolution of parallel programming:
  - **Task graph parallelism** (TGP) is the best model for future parallel arch
  - Capture programmers' intention in decomposing a parallel algorithm into a top-down task graph
  - Runtime can schedule dependent tasks across a large number of processing units (e.g., CPUs, GPUs)

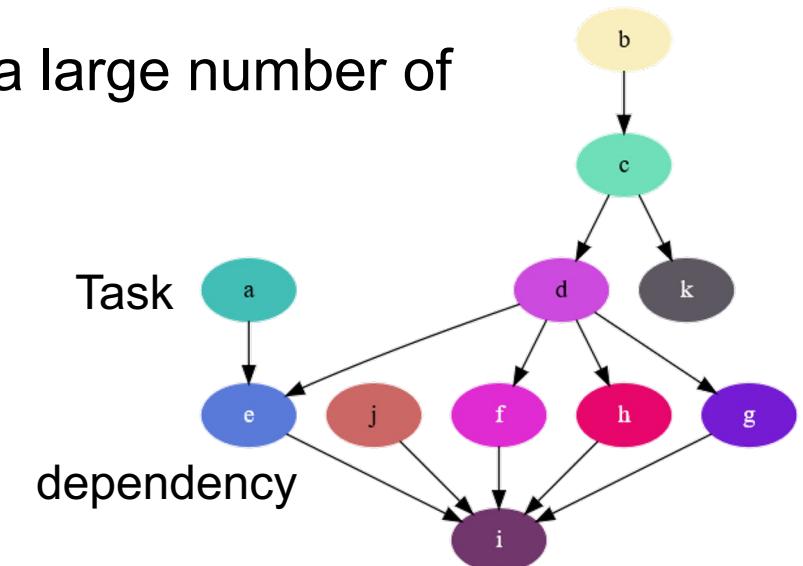
**OpenMP**



**StarPU**

**PaRSEC**

**kokkos**





# Takeaways

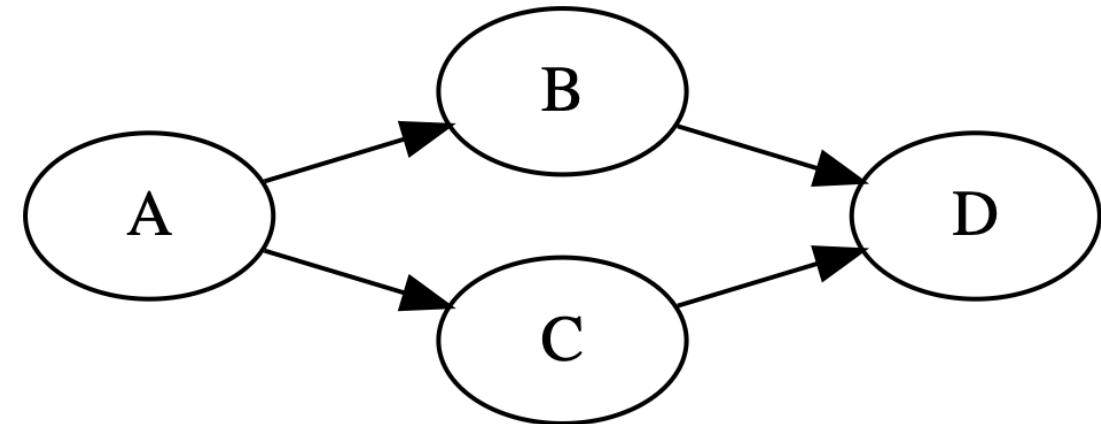
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# “Hello World” in Taskflow<sup>1</sup>

```
#include <taskflow/taskflow.hpp>
int main(){
    tf::Taskflow taskflow;
    tf::Executor executor;
    auto [A, B, C, D] = taskflow.emplace(
        [] () { std::cout << "TaskA\n"; },
        [] () { std::cout << "TaskB\n"; },
        [] () { std::cout << "TaskC\n"; },
        [] () { std::cout << "TaskD\n"; });
    A.precede(B, C);
    D.succeed(B, C);
    executor.run(taskflow).wait();
    return 0;
}
```

// live: <https://godbolt.org/z/j8hx3xnnx>



<sup>1</sup>: T.-W. Huang, et. al, “Taskflow: A Lightweight Parallel and Heterogeneous Task Graph Computing System,” *IEEE TPDS*, vol. 33, no. 6, pp. 1303-1320, June 2022



# Drop-in Integration

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- **Taskflow is header-only and written in completely standard C++**
  - ☺ No wrangle with installation

```
# clone the Taskflow project
~$ git clone https://github.com/taskflow/taskflow.git
~$ cd taskflow

# compile your program and tell it where to find Taskflow header files
~$ g++ -std=c++20 examples/simple.cpp -I ./ -O2 -pthread -o simple
~$ ./simple
```

TaskA

TaskC

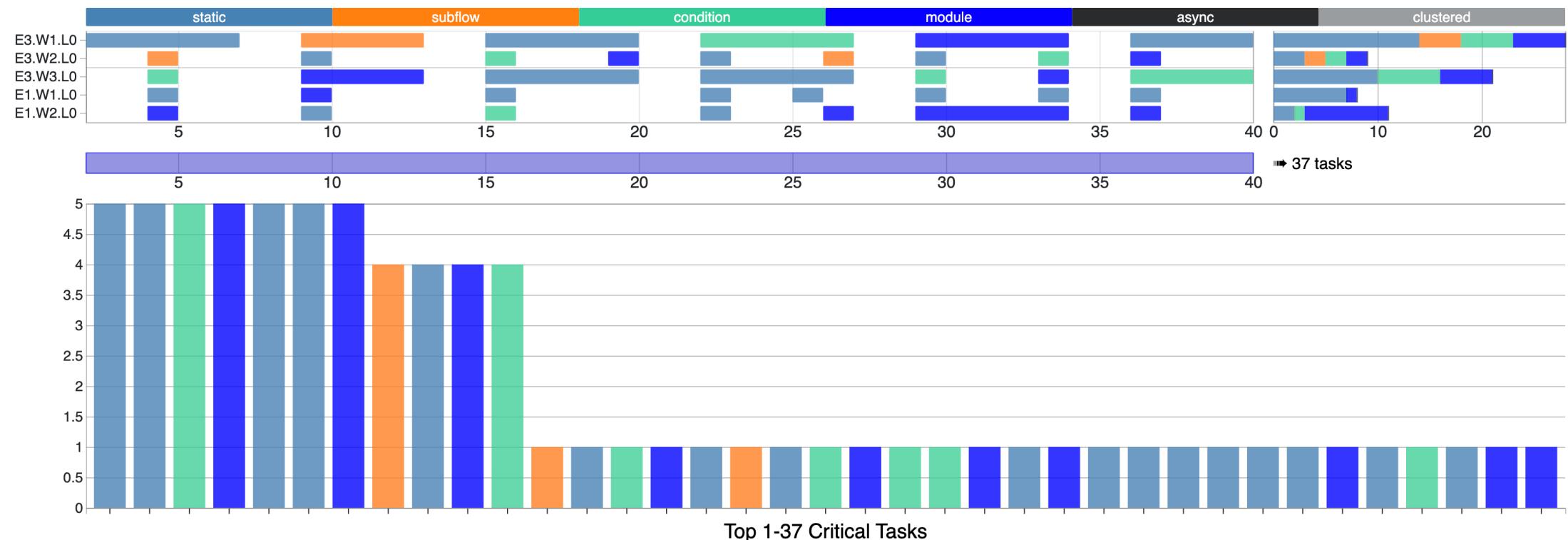
TaskB

TaskD



# Built-in Task Execution Visualizer

```
# run your program with the env variable TF_ENABLE_PROFILER enabled  
# to generate profile data and post it on https://taskflow.github.io/tfprof/  
~$ TF_ENABLE_PROFILER=simple.json ./simple
```

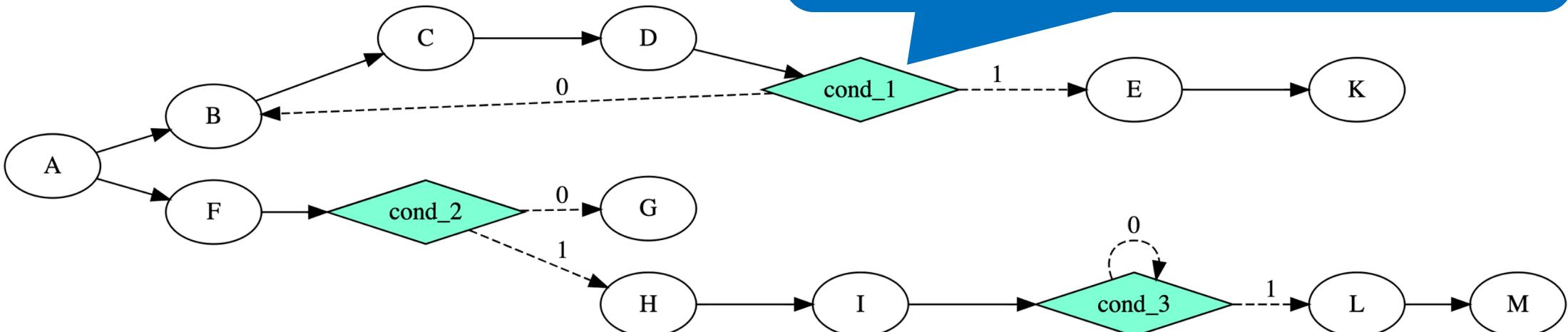


# Control Taskflow Graph Programming (CTFG)

// CTFG goes beyond the limitation of traditional DAG-based models

```
auto cond_1 = taskflow.emplace([](){ return run_B() ? 0 : 1; }); // 0: is the index of B
auto cond_2 = taskflow.emplace([](){ return run_G() ? 0 : 1; }); // 0: is the index of G
auto cond_3 = taskflow.emplace([](){ return loop() ? 0 : 1; }); // 0: is the index of cond_3
cond_1.precede(B, E);           // cycle
cond_2.precede(G, H);           // if-else
cond_3.precede(cond_3, L);      // loop
```

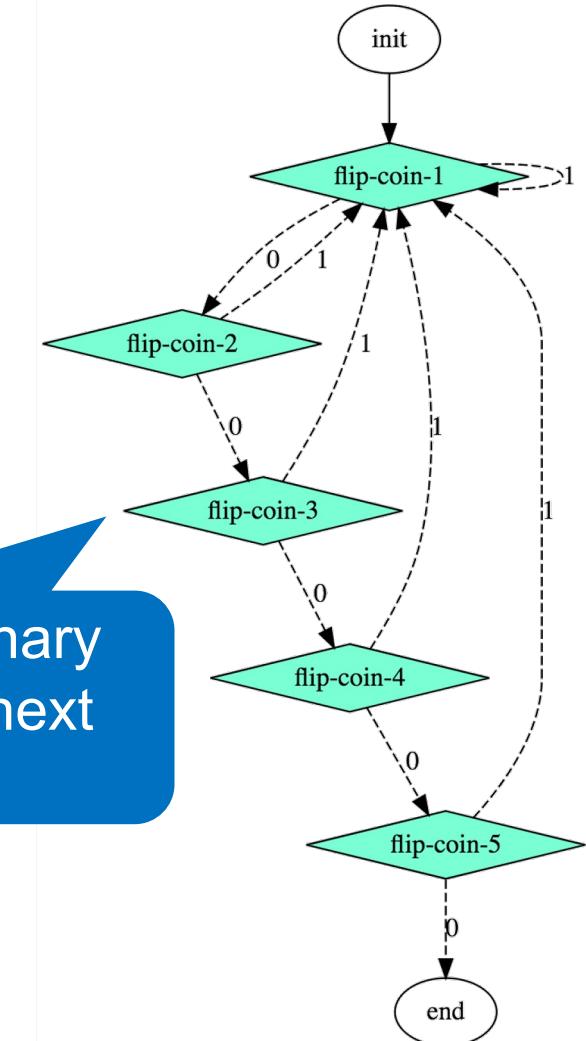
Very difficult for existing DAG-based systems to express an efficient overlap between tasks and control flow ...



# Non-deterministic Control Flow with CTFG

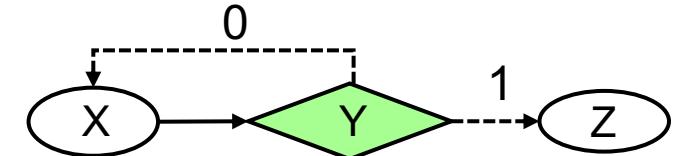
```
auto A = taskflow.emplace([&](){ } );
auto B = taskflow.emplace([&](){ return rand()%2; } );
auto C = taskflow.emplace([&](){ return rand()%2; } );
auto D = taskflow.emplace([&](){ return rand()%2; } );
auto E = taskflow.emplace([&](){ return rand()%2; } );
auto F = taskflow.emplace([&](){ return rand()%2; } );
auto G = taskflow.emplace([&]());
A.precede(B).name("init");
B.precede(C, B).name("flip-coin-1");
C.precede(D, B).name("flip-coin-2");
D.precede(E, B).name("flip-coin-3");
E.precede(F, B).name("flip-coin-4");
F.precede(G, B).name("flip-coin-5");
G.name("end");
```

Each task flips a binary coin to decide the next task to run



# Existing Frameworks on Control Flow?

- Most existing libraries are DAG-based
  - Do not anticipate conditional execution ...
- Unroll a task graph over fixed iterations
  - Task graph size becomes very large ...
- What about dynamic control flow?
  - Have no choice but resort to a client-side partition of the task graph
  - Synchronize the execution of partitioned task graphs around decision-making points
  - Lack end-to-end parallelism



```

tf::Taskflow G;
auto X = G.emplace([](){});
auto Y = G.emplace([](){
    return converged() ? 1 : 0;
});
cond.precede(Z, X);
executor.run(G).wait();
  
```

↓

```

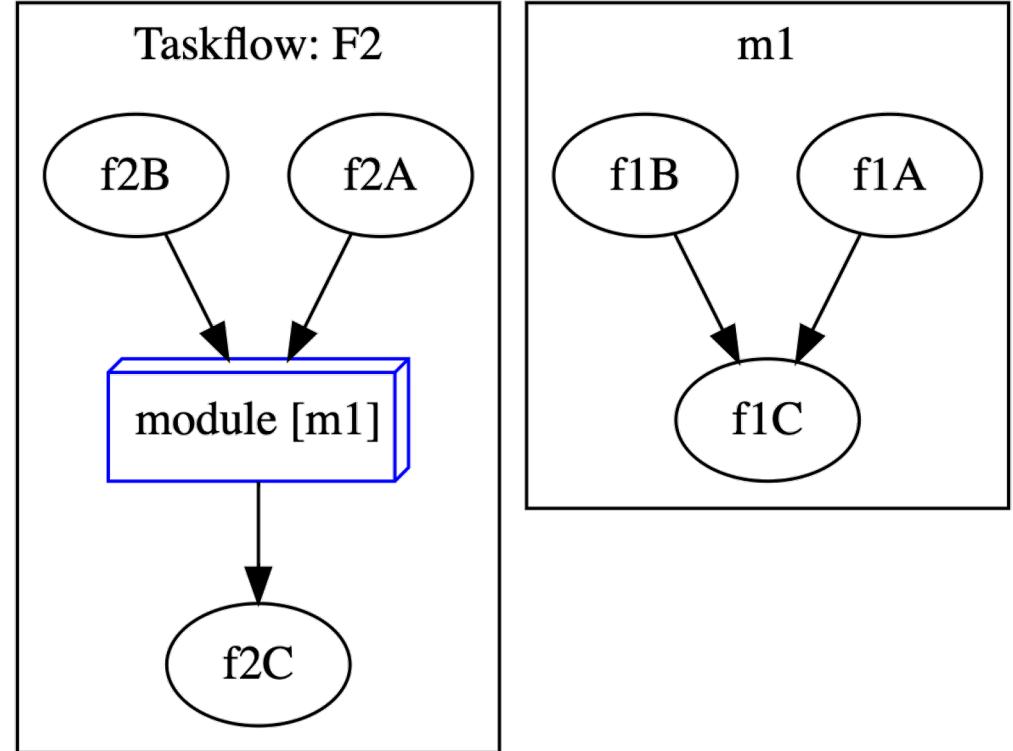
tbb::flow::graph X, Y, Z;
do {
    X.run();
    Y.run();
} while (!converged());
Z.run();
  
```

# Composable Tasking

```

tf::Taskflow m1, f2;
auto [f1A, f1B] = m1.emplace(
    []() { std::cout << "Task f1A\n"; },
    []() { std::cout << "Task f1B\n"; }
);
auto [f2A, f2B, f2C] = f2.emplace(
    []() { std::cout << "Task f2A\n"; },
    []() { std::cout << "Task f2B\n"; },
    []() { std::cout << "Task f2C\n"; }
);
auto f1_module_task = f2.composed_of(m1);
f1_module_task.succeed(f2A, f2B)
    .precede(f2C);

```



# Everything is Composable in Taskflow

- End-to-end parallelism in one graph
  - Task, dependency, control flow all together
  - Scheduling with whole-graph optimization
  - Efficiently overlap tasks with control flow
- Largely improved productivity!

Composition  
(HPDC'22, ICPP'22, HPEC'19)

Industrial use-case of productivity improvement using Taskflow

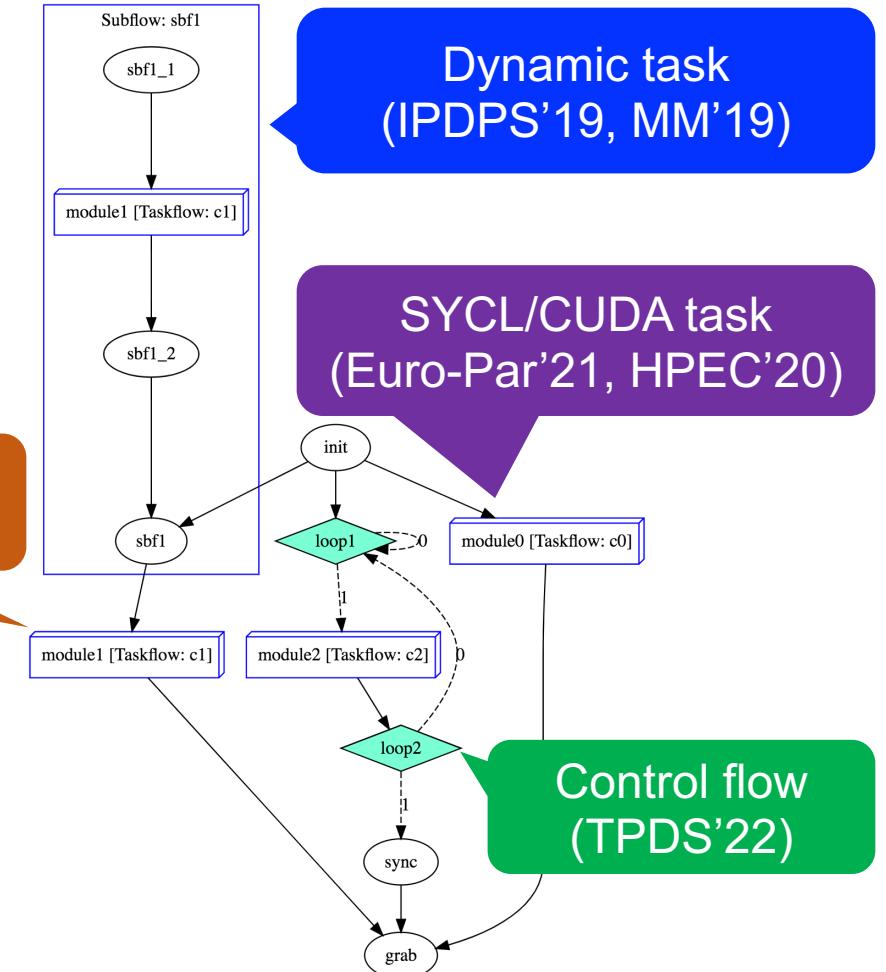
jcelerier ossia score

Reddit: <https://www.reddit.com/r/cpp/> [under taskflow]

I've migrated <https://ossia.io> from TBB flow graph to taskflow a couple weeks ago. Net +8% of throughput on the graph processing itself, and took only a couple hours to do the change! Also don't have to fight with building the TBB libraries for 30 different platforms and configurations since it's header only.

↑ 8 ↓ Reply Share Report Save Follow

ossia score





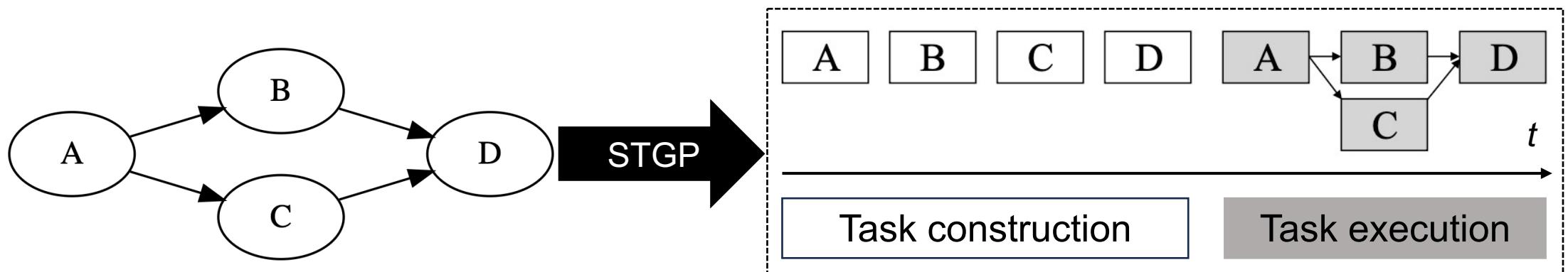
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# Static Task Graph Parallelism (STGP)

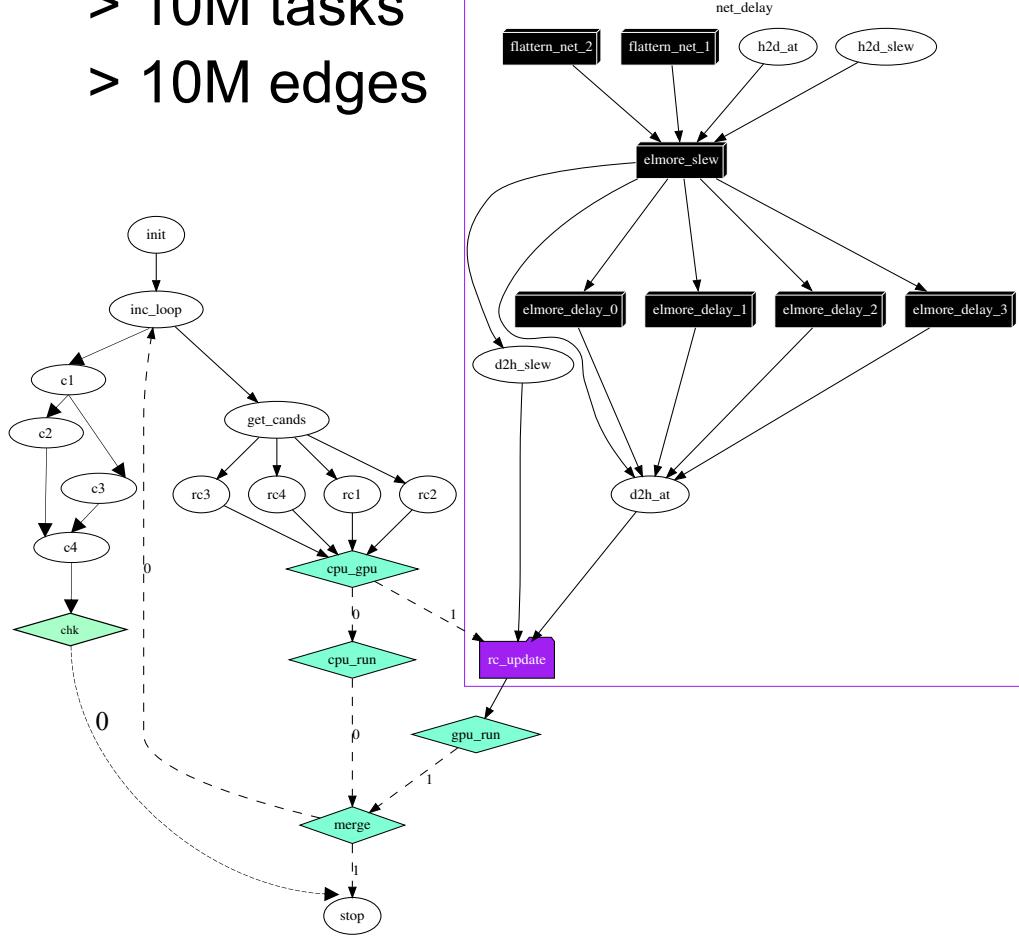
- In STGP, the graph structure must be known at compile time
  - Construct-and-run model – *Construct the task graph first and then run it*
- Lack of overlap between task construction and task execution
  - For large task graphs (e.g., multi-million tasks and dependencies), such an overlap can bring a significant performance advantage
- Lack of flexible and dynamic expressions of TGP
  - *Task graph structure cannot depend on runtime values or control-flow results*



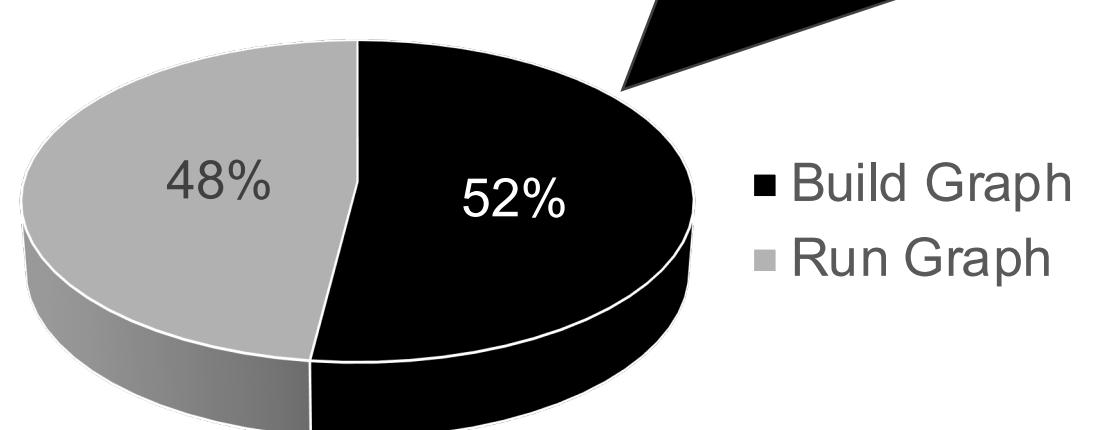
# Problem of STGP: Example #1

- Runtime breakdown of a task-parallel circuit analysis algorithm<sup>1</sup>

- > 10M tasks
- > 10M edges



Task graph construction time takes over 50% of the entire runtime



- Build Graph
- Run Graph

<sup>1</sup>: Tsung-Wei Huang, et al, "OpenTimer v2: A New Parallel Incremental Timing Analysis Engine," *IEEE TCAD*, 2022



# Problem of STGP: Example #2

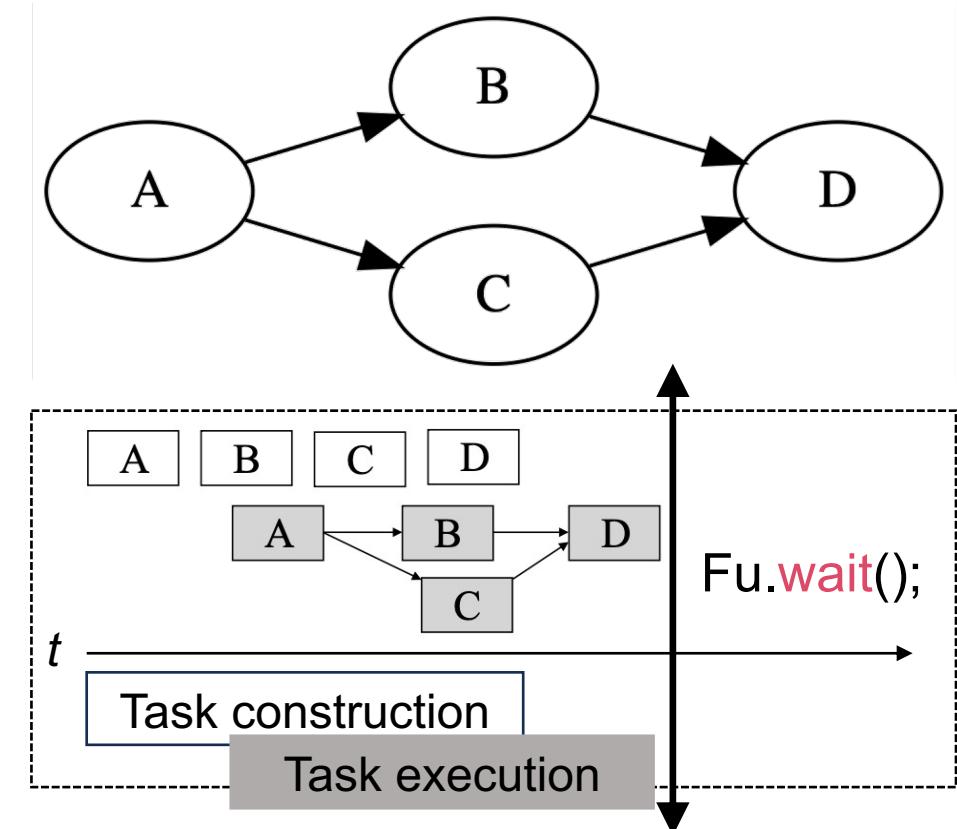
- TGP that depends on dynamic control-flow results

```
if (a == true) {  
    G1 = build_task_graph1();  
    if (b == true) {  
        G2 = build_task_graph2();  
        G1.precede(G2);  
        if (c == true) {  
            ... // another level of dynamic TGP  
        }  
    }  
    else {  
        G3 = build_task_graph3();  
        G3.precede(G1);  
    }  
}
```

```
G1 = build_task_graph1();  
G2 = build_task_graph2();  
if (G1.num_tasks() == 100) {  
    G1.precede(G2);  
}  
else {  
    G3 = build_task_graph3();  
    G2.precede(G1, G3);  
    if (G2.num_dependencies() >= 10) {  
        ... // another level of dynamic TGP  
    }  
}
```

# Dynamic TGP (DTGP) in Taskflow

```
// Live: https://godbolt.org/z/j76ThGbWK
tf::Executor executor;
auto A = executor.silent_dependent_async([](){
    std::cout << "TaskA\n";
});
auto B = executor.silent_dependent_async([](){
    std::cout << "TaskB\n";
}, A);
auto C = executor.silent_dependent_async([](){
    std::cout << "TaskC\n";
}, A);
auto [D, Fu] = executor.dependent_async([](){
    std::cout << "TaskD\n";
}, B, C); ←
Fu.wait();
```

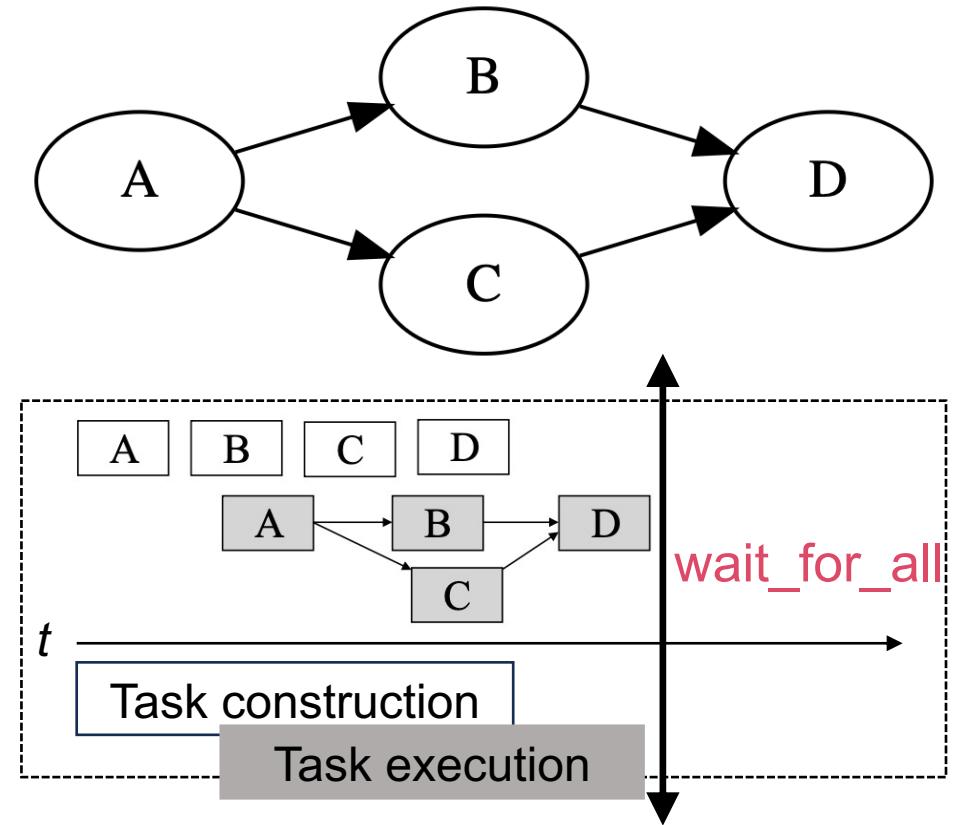


Specify arbitrary task dependencies using  
C++ variadic parameter pack

# silent\_dependent\_async is cheaper

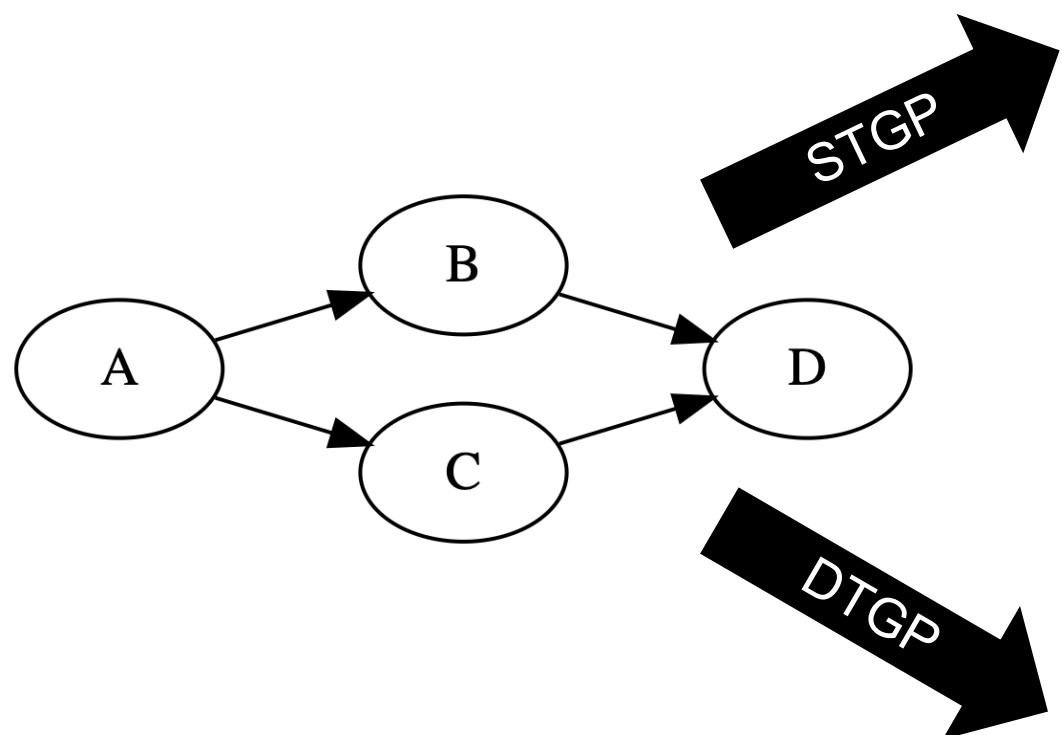
// Live: <https://godbolt.org/z/T87PrTarx>

```
tf::Executor executor;
auto A = executor.silent_dependent_async([](){
    std::cout << "TaskA\n";
});
auto B = executor.silent_dependent_async([](){
    std::cout << "TaskB\n";
}, A);
auto C = executor.silent_dependent_async([](){
    std::cout << "TaskC\n";
}, A);
auto D = executor.silent_dependent_async([](){
    std::cout << "TaskD\n";
}, B, C);
executor.wait_for_all();
```



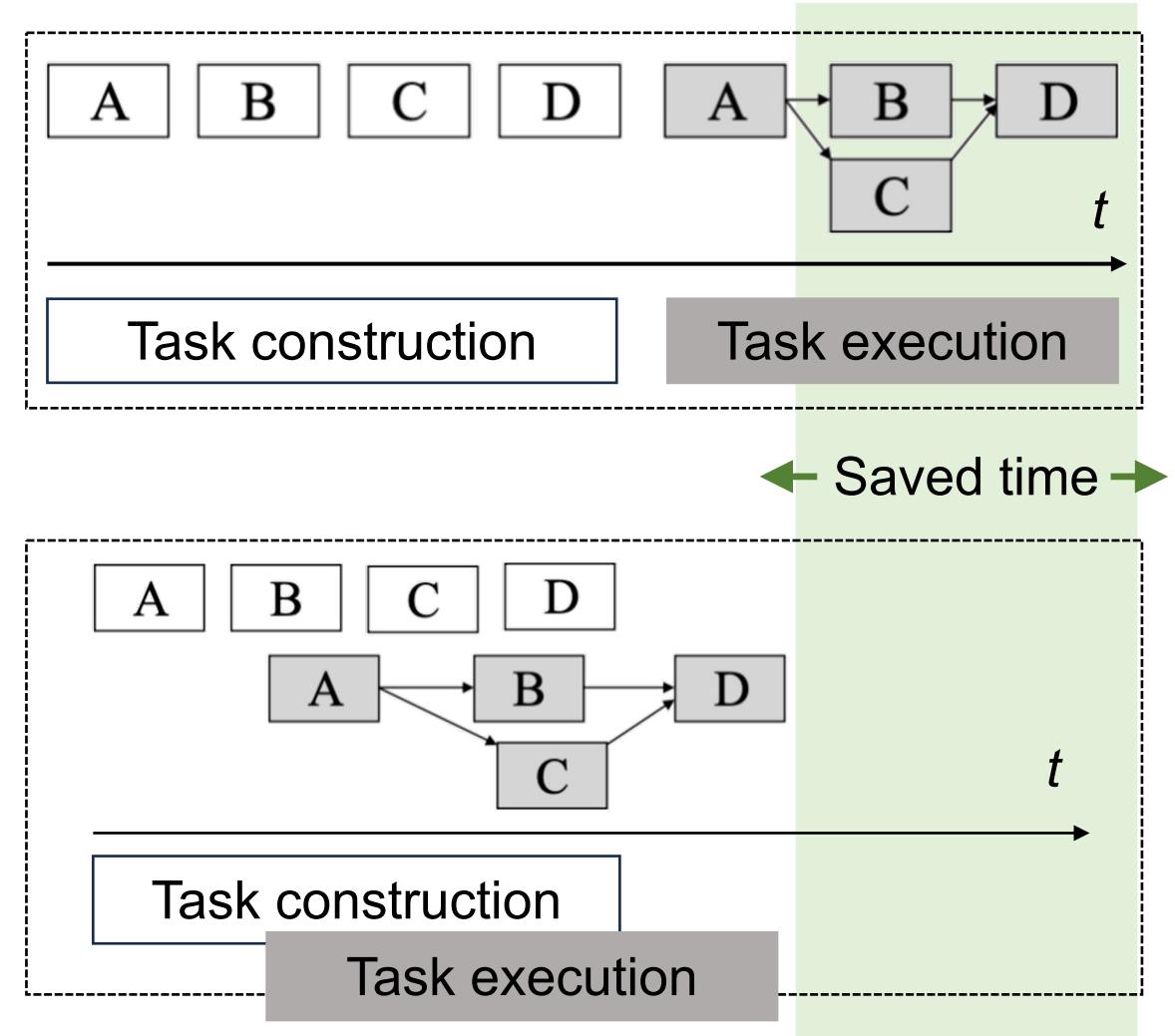
Block the caller until all tasks (A, B, C, and D) finish

# Comparison between STGP and DTGP



*STGP*

*DTGP*



# DTGP Needs a Correct Topological Order

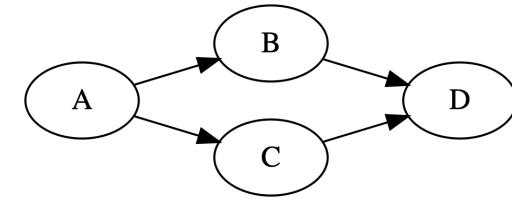
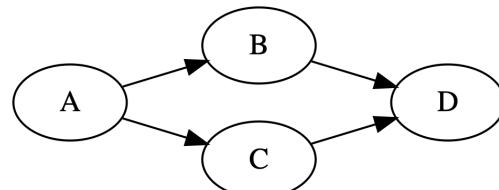
```
auto A = executor.silent_dependent_async([](){
    std::cout << "TaskA\n";
});

auto B = executor.silent_dependent_async([](){
    std::cout << "TaskB\n";
}, A);

auto C = executor.silent_dependent_async([](){
    std::cout << "TaskC\n";
}, A);

auto D = executor.silent_dependent_async([](){
    std::cout << "TaskD\n";
}, B, C);
```

Topological order #1: A→B→C→D



Topological order #2: A→C→B→D

```
auto A = executor.silent_dependent_async([](){
    std::cout << "TaskA\n";
});

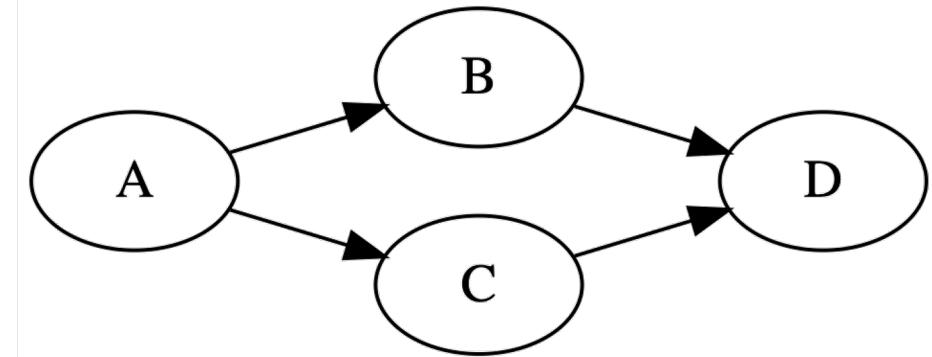
auto C = executor.silent_dependent_async([](){
    std::cout << "TaskC\n";
}, A);

auto B = executor.silent_dependent_async([](){
    std::cout << "TaskB\n";
}, A);

auto D = executor.silent_dependent_async([](){
    std::cout << "TaskD\n";
}, B, C);
```

# Incorrect Topological Order doesn't Work

```
tf::Executor executor;
auto A = executor.silent_dependent_async([](){
    std::cout << "TaskA\n";
});
auto D = executor.silent_dependent_async([](){
    std::cout << "TaskD\n";
}, B-is-unavailable-yet, C-is-unavailable-yet);
auto B = executor.silent_dependent_async([](){
    std::cout << "TaskB\n";
}, A);
auto C = executor.silent_dependent_async([](){
    std::cout << "TaskC\n";
}, A);
executor.wait_for_all();
```



An incorrect topological order ( $A \rightarrow D \rightarrow B \rightarrow C$ ) disallows us from expressing correct DTGP



# Variable Range of Task Dependencies

- Both methods accept a variable range of dependent tasks
  - Useful when the task dependencies come as a runtime variable

```
// Live: https://godbolt.org/z/6Pvco4KeE
std::vector<tf::AsyncTask> tasks = {
    executor.silent_dependent_async([](){ std::cout << "TaskA\n"; }),
    executor.silent_dependent_async([](){ std::cout << "TaskB\n"; }),
    executor.silent_dependent_async([](){ std::cout << "TaskC\n"; }),
    executor.silent_dependent_async([](){ std::cout << "TaskD\n"; })
};
// create a dependent-async tasks that depends on tasks, A, B, C, and D
executor.dependent_async([](){}, tasks.begin(), tasks.end());

// create a silent-dependent-async tasks that depends on tasks, A, B, C, and D
executor.silent_dependent_async([](){}, tasks.begin(), tasks.end());
```



# Existing DTGP Libraries: C++ Async

```
auto A = std::async([&](){ std::cout << "TaskA\n"; });
auto B = std::async([&](){
    A.get(); ←
    std::cout << "TaskB\n";
});
auto C = std::async([&](){
    A.get(); ←
    std::cout << "TaskC\n";
});
auto D = std::async([&](){
    B.get(); ←
    C.get(); ←
    std::cout << "TaskD\n";
}, B, C);
D.get();
```

Block until A completes

Block until A completes

Block until B and C complete



# Existing DTGP Libraries: OpenMP/Kokkos

## OpenMP “dependency clauses”

```
#omp parallel num_threads(hardware_concurrency())
{
    int A_B, A_C, B_D, C_D;
    #pragma omp task depend(out: A_B, A_C)
    {
        std::cout << "TaskA\n";
    }
    #pragma omp task depend(in: A_B; out: B_D)
    {
        std::cout << "TaskB\n";
    }
    #pragma omp task depend(in: A_C; out: C_D)
    {
        std::cout << "TaskC\n";
    }
    #pragma omp task depend(in: B_D, C_D)
    {
        std::cout << "TaskD\n";
    }
}
```

## Kokkos “task template”

```
struct A {
    template <class TeamMember> KOKKOS_INLINE_FUNCTION
    void operator()(TeamMember& member) { std::cout << "TaskA\n"; }
};

struct B {
    template <class TeamMember> KOKKOS_INLINE_FUNCTION
    void operator()(TeamMember& member) { std::cout << "TaskB\n"; }
};

struct C {
    template <class TeamMember> KOKKOS_INLINE_FUNCTION
    void operator()(TeamMember& member) { std::cout << "TaskC\n"; }
};

struct D {
    template <class TeamMember> KOKKOS_INLINE_FUNCTION
    void operator()(TeamMember& member) { std::cout << "TaskD\n"; }
};

auto scheduler = scheduler_type(/* ... */);
auto futA = Kokkos::host_spawn(Kokkos::TaskSingle(scheduler), A());
auto futB = Kokkos::host_spawn(Kokkos::TaskSingle(scheduler, futA), B());
auto futC = Kokkos::host_spawn(Kokkos::TaskSingle(scheduler, futA), C());
auto futD = Kokkos::host_spawn(
    Kokkos::TaskSingle(scheduler, when_all(futB, futC)), D()
);
... (more code to follow)
```



# Existing DTGP Libraries: PARSEC

```
int A(parsec_task_t* this_task) {
    int *out;
    unpack_args(this_task, &out);
    printf("TaskA\n");
    return APARSEC_HOOK_RETURN_DONE;
}
int B(parsec_task_t* this_task) {
    int *out, *in;
    unpack_args(this_task, &in, &out);
    printf("TaskB\n");
    return APARSEC_HOOK_RETURN_DONE;
}
int C(parsec_task_t* this_task) {
    int *in, *out;
    unpack_args(this_task, &in, &out);
    printf("TaskC\n");
    return APARSEC_HOOK_RETURN_DONE;
}
int D(parsec_task_t* this_task) {
    int *in1, *in2, *out;
    unpack_args(this_task, &in1, &in2, &out);
    printf("TaskD\n");
    return APARSEC_HOOK_RETURN_DONE;
}
```

```
int main() {
    /* additional boilerplate code to initialize PARSEC runtime
       environment (e.g., taskpool, dtd_tp, dependency data, etc.) */
    int *out;
    parsec_dtd_insert_task(A,
        tile_of_key(dependency, 0), INPUT, ARSEC_DTD_ARG_END
    );
    parsec_dtd_insert_task(B,
        tile_of_key(dependency, 0), INPUT,
        tile_of_key(dependency, 1), OUTPUT, PARSEC_DTD_ARG_END
    );
    parsec_dtd_insert_task(C,
        tile_of_key(dependency, 0), INPUT,
        tile_of_key(dependency, 2), OUTPUT, PARSEC_DTD_ARG_END
    );
    parsec_dtd_insert_task(D,
        tile_of_key(dependency, 1), INPUT,
        tile_of_key(dependency, 2), INPUT,
        tile_of_key(dependency, 3), OUTPUT, PARSEC_DTD_ARG_END
    );
    parsec_taskpool_wait();
}
```



# Limitations of Existing DTGP Libraries

---

## 1. Suffer from large verbosity from the ease-of-use standpoint

- The code does not explain itself – *I know this is subjective, but most applications just care how fast they can get things done ...*

## 2. Count on dataflow to express task dependencies

- Users need to explicitly define per-edge data to specify a task dependency
  - OpenMP's **in-out** dependency clauses
  - PARSEC's **INPUT-OUTPUT** keywords and **tile\_of\_key** constructs

## 3. Rely on complex data structure to schedule tasks

- Libomp implements a lock-based hash table to schedule tasks
  - Key: address of the input and output data of a task
  - Value: a list of tasks accessing that input and output data
- Frequent access to this data structure results in large runtime overhead



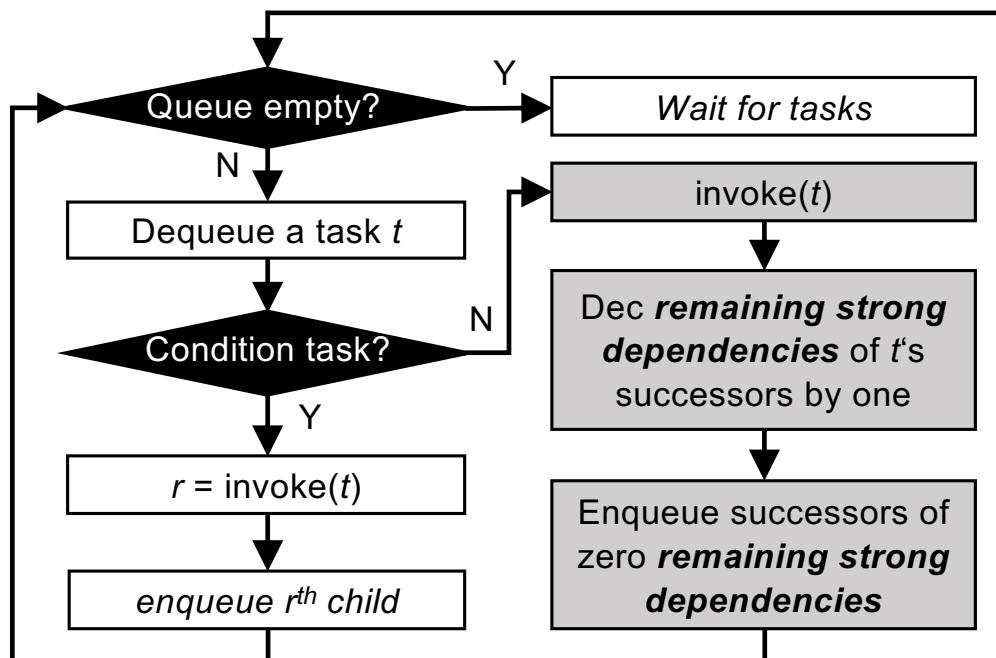
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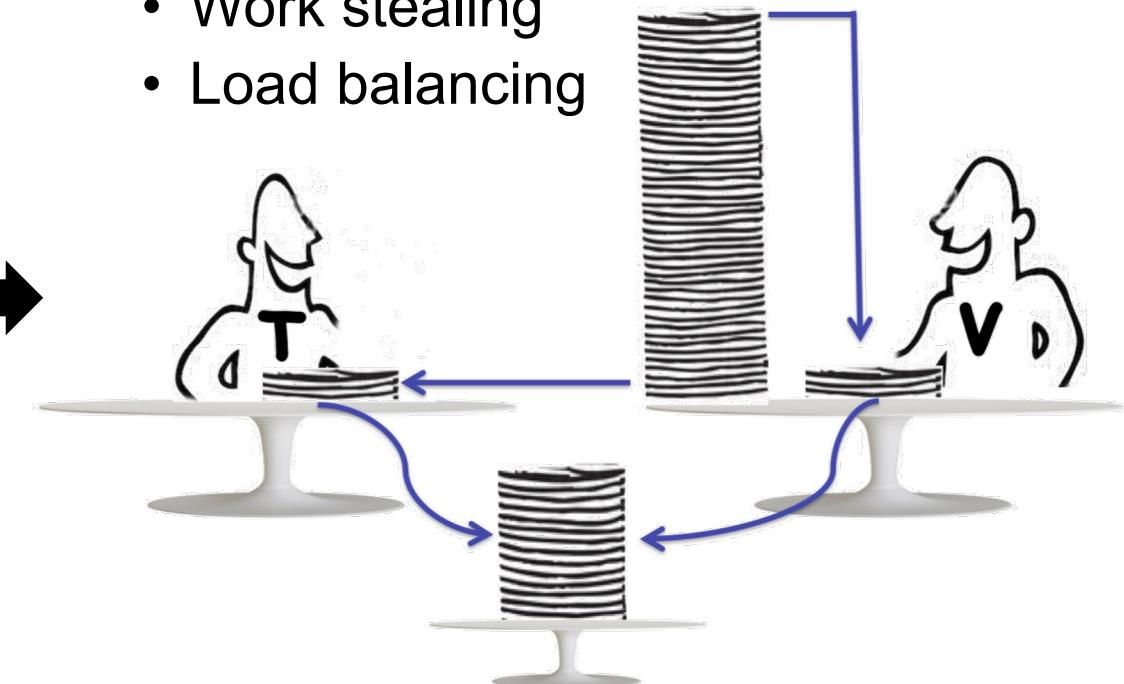
# STGP Scheduling Algorithm

- Task-level scheduling



- Worker-level scheduling

- Work stealing
- Load balancing

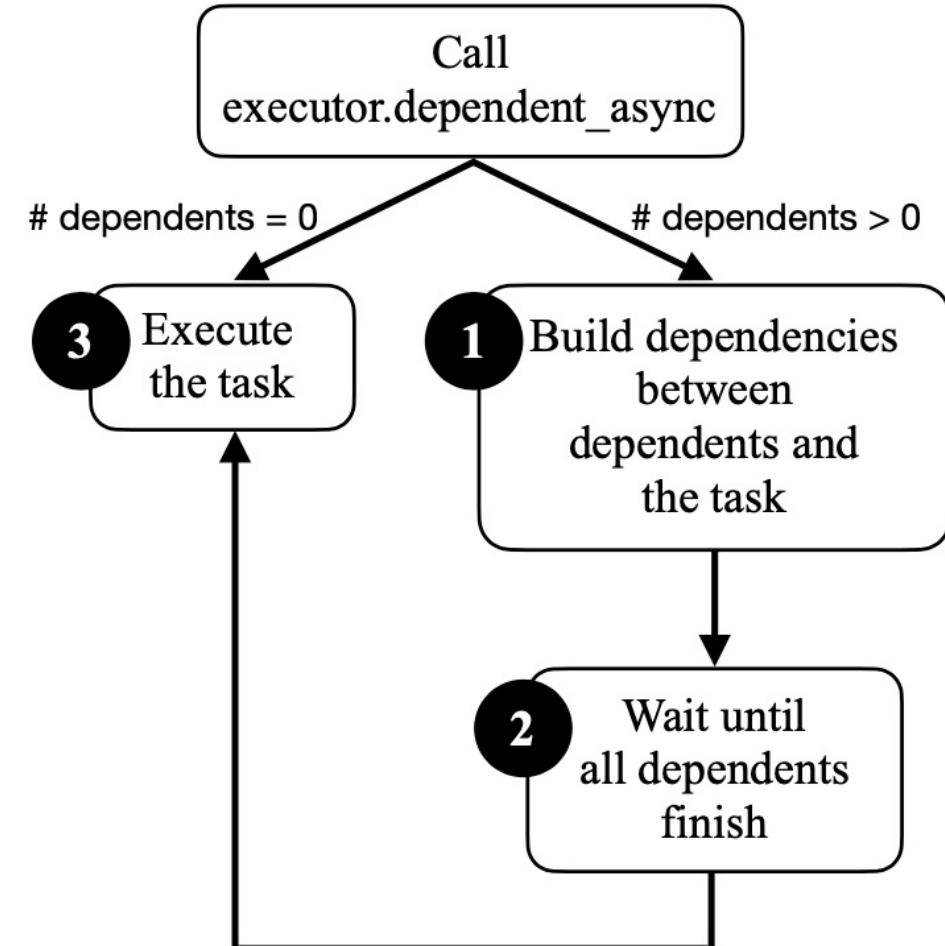


**Key results:** schedule tasks with in-graph control flow with a **strong balance** between the number of active workers and dynamically generated tasks – *low latency, energy efficient, and high throughput*



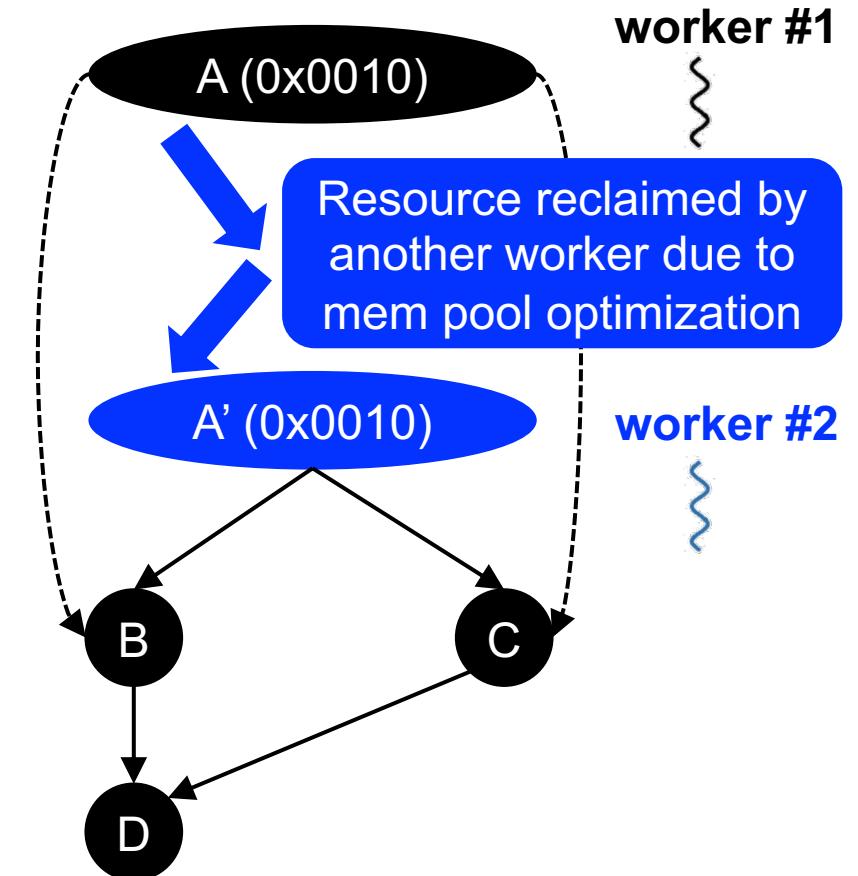
# DTGP Scheduling Algorithm

- The algorithm has three parts:
  - Build dependencies
  - Wait for dependents to finish
  - Execute the task
- Three key scheduling challenges:
  1. **ABA** – a specified task dependency must exist correctly
  2. **Data race** – multiple threads may simultaneously modify the dependency structure of a task
  3. **Synchronization** – application can issue a global synchronization at any time to wait for all tasks to finish



# Solving Challenge #1: ABA Problem<sup>1</sup>

```
tf::Executor executor;
auto A = executor.silent_dependent_async([](){
    std::cout << "TaskA\n";
});
auto B = executor.silent_dependent_async([](){
    std::cout << "TaskB\n";
}, A);
auto C = executor.silent_dependent_async([](){
    std::cout << "TaskC\n";
}, A);
auto D = executor.silent_dependent_async([](){
    std::cout << "TaskD\n";
}, B, C);
executor.wait_for_all();
```



<sup>1</sup>: ABA Problem: [https://en.wikipedia.org/wiki/ABA\\_problem](https://en.wikipedia.org/wiki/ABA_problem)



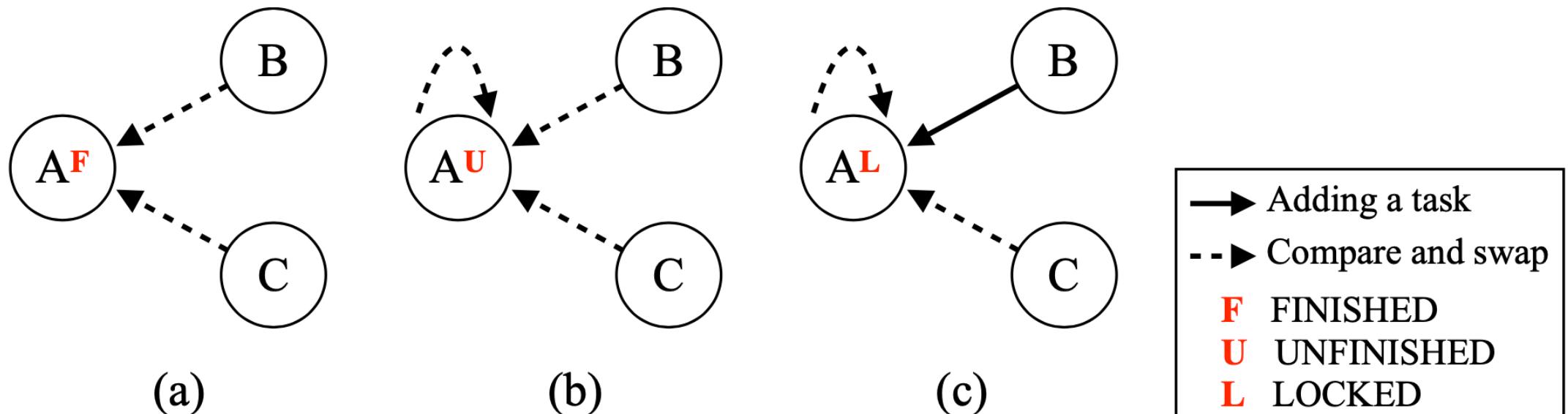
# Retain Shared Ownership of Every Task

```
tf::Executor executor;
tf::AsyncTask A = executor.silent_dependent_async([](){
    std::cout << "TaskA\n";
});
tf::AsyncTask B = executor.silent_dependent_async([](){
    std::cout << "TaskB\n";
}, A); ←
tf::AsyncTask C = executor.silent_dependent_async([](){
    std::cout << "TaskC\n";
}, A);
tf::AsyncTask D = executor.silent_dependent_async([](){
    std::cout << "TaskD\n";
}, B, C);
executor.wait_for_all();
```

tf::AsyncTask acts like  
a std::shared\_ptr to  
ensure a task is alive  
when being used

# Solving Challenge #2: Data Race

- Both B and C want to add themselves to the successors of A
  - In the meantime, A may want to remove its successor
- Apply compare-and-swap (CAS) to enable exclusive access
  - As a result, constructing a dynamic task graph can be completely thread-safe





# Solving Challenge #3: Synchronization

- Application can issue a global synchronization at any time

- executor.wait\_for\_all();

```
tf::Executor executor;
auto A = executor.silent_dependent_async([](){});  
auto B = executor.silent_dependent_async([](){}, A);
executor.wait_for_all(); // wait for A and B to finish
```

```
auto C = executor.silent_dependent_async([](){}, A);
auto D = executor.silent_dependent_async([](){}, B, C);
executor.wait_for_all(); // wait for C and D to finish
```

```
...
```

```
executor.wait_for_all(); // wait for other tasks to finish
```

```
// lock-based solution
std::unique_lock lock(mutex);
cv.wait(lock, [&](){
    return num_tasks == 0;
});
```

```
// atomic wait-based solution
auto n = num_tasks.load();
while(n != 0) {
    num_tasks.wait(n);
    n = num_tasks.load();
});
```



# Lock-free Scheduling Algorithm<sup>1</sup>

---

**Algorithm 1** dependent\_async(callable, deps)

```
1: Create a future
2: num_deps  $\leftarrow$  sizeof(deps)
3: task  $\leftarrow$  initialize_task(callable, num_deps, future)
4: for all dep  $\in$  deps do
5:   process_dependent(task, dep, num_deps)
6: end for
7: if num_deps == 0 then
8:   schedule_async_task(task)
9: end if
10: return (task, future)
```

---

**Algorithm 2** process\_dependent(task, dep, num\_deps)

```
1: dep_state  $\leftarrow$  dep.state
2: target_state  $\leftarrow$  UNFINISHED
3: if dep_state.CAS(target_state, LOCKED) then
4:   dep.successors.push(task)
5:   dep_state  $\leftarrow$  UNFINISHED
6: else if target_state == FINISHED then
7:   num_deps  $\leftarrow$  AtomDec(task.join_counter)
8: else
9:   goto line 2
10: end if
```

---

**Algorithm 3** schedule\_async\_task(task)

```
1: target_state  $\leftarrow$  UNFINISHED
2: while not task.state.CAS(target_state, FINISHED)
   do
3:   target_state  $\leftarrow$  UNFINISHED
4: end while
5: Invoke(task.callable)
6: for all successor  $\in$  task.successors do
7:   if AtomDec(successor.join_counter) == 0 then
8:     schedule_async_task(successor)
9:   end if
10: end for
11: if AtomDec(task.ref_count) == 0 then
12:   Delete task
13: end if
```

---

<sup>1</sup>: Cheng-Hsiang Chiu, et. al, "Programming Dynamic Task Parallelism for Heterogeneous EDA Algorithms," IEEE/ACM ICCAD, CA, 2023



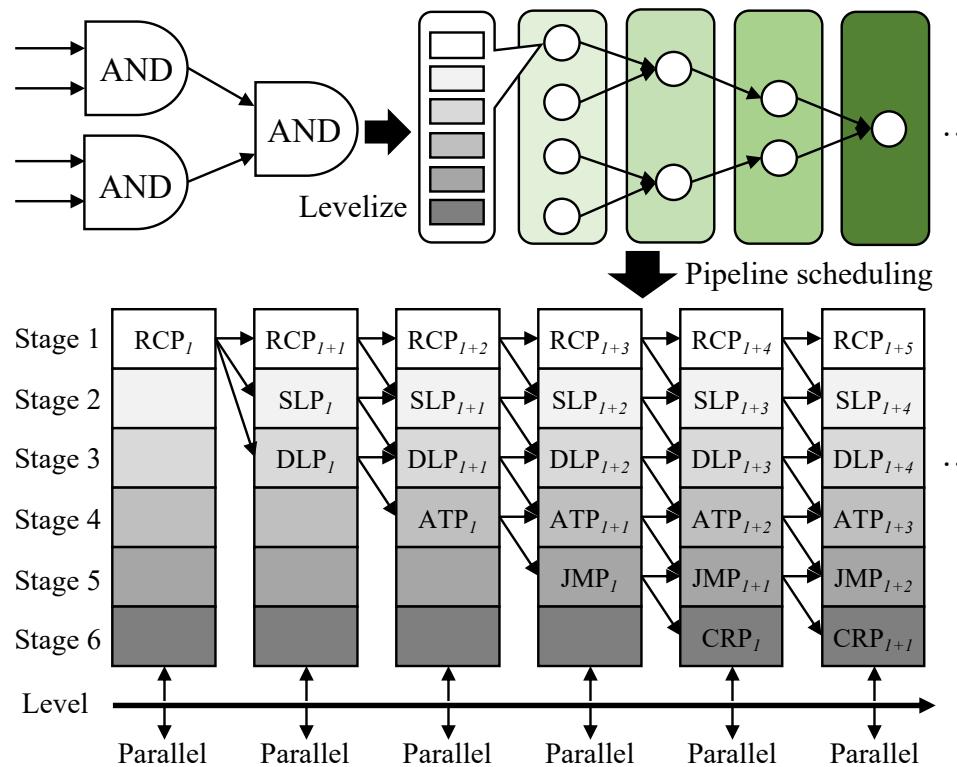
# Takeaways

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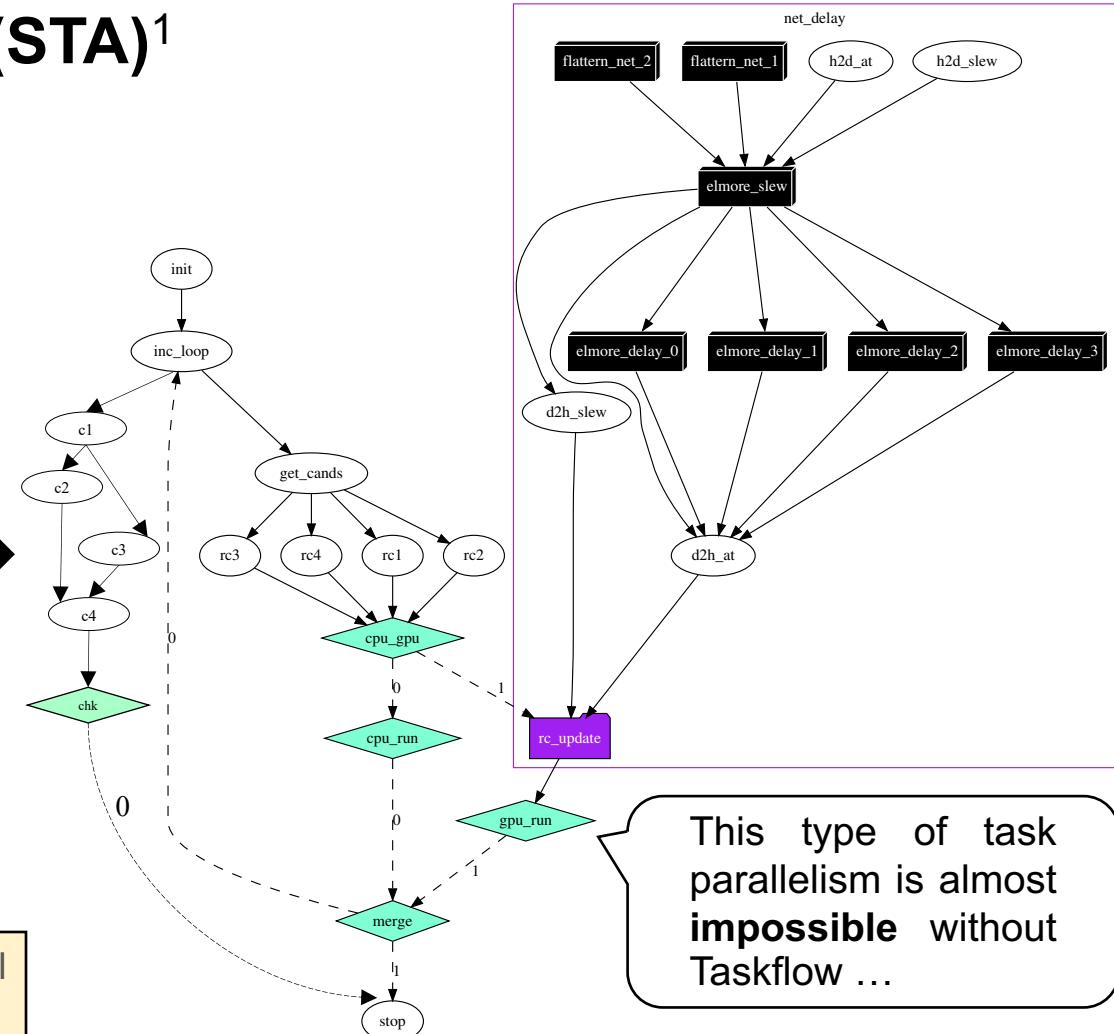
- Express your parallelism in the right way
- Program task graph parallelism using Taskflow
- Program dynamic task graph parallelism using Taskflow
- Overcome the scheduling challenges
- Demonstrate the efficiency of Taskflow in industrial application
- Conclude the talk

# Experimental Results – STGP

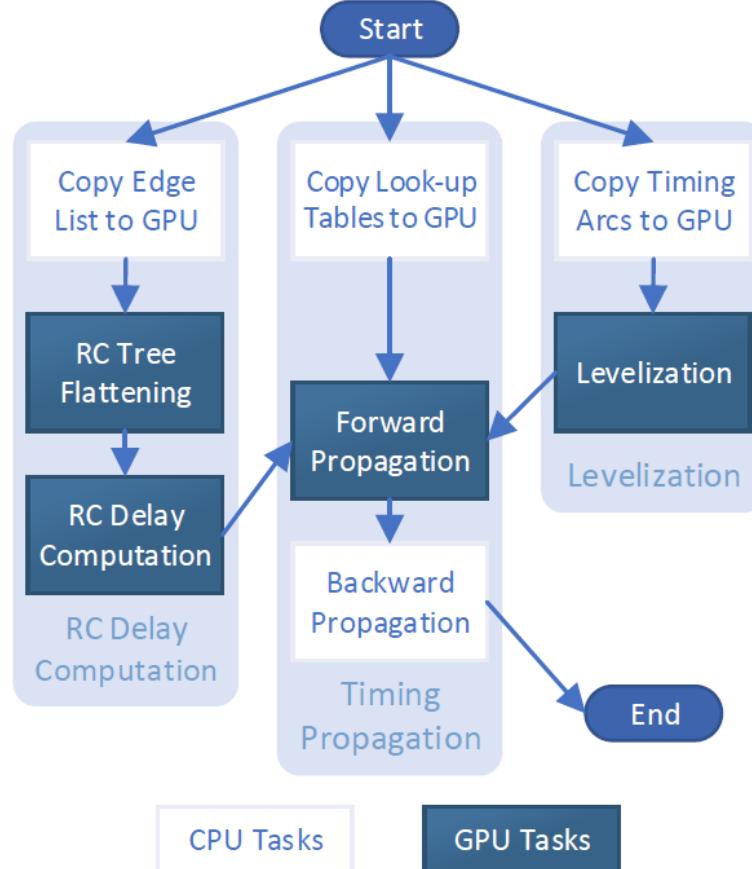
- Parallelize static timing analysis (STA)<sup>1</sup>



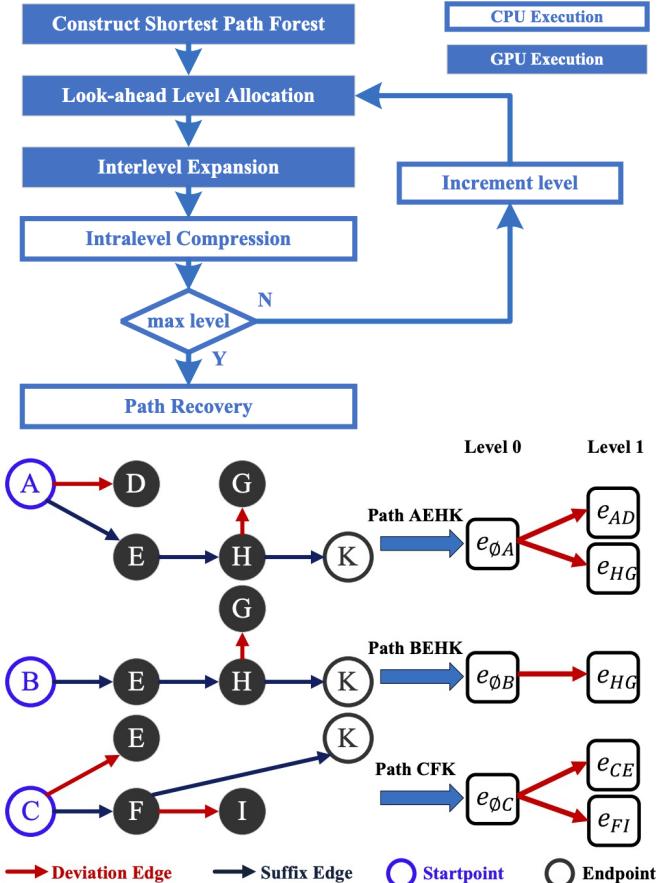
<sup>1</sup>: Tsung-Wei Huang, et al, "OpenTimer v2: A New Parallel Incremental Timing Analysis Engine," *IEEE TCAD*, 2022



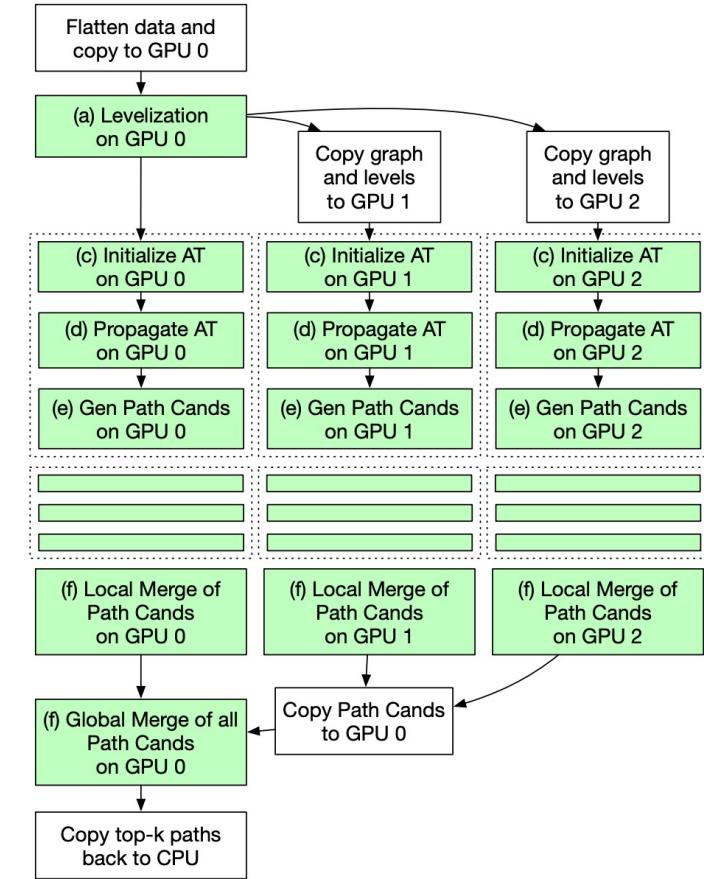
# Task-parallel STA Alg with CPU and GPU



GPU-based graph analysis (ICCAD'20)



GPU-based path analysis (DAC'21)



GPU-based CPPR (ICCAD'21)



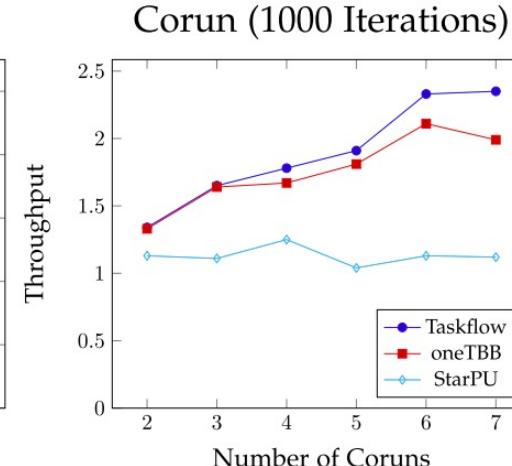
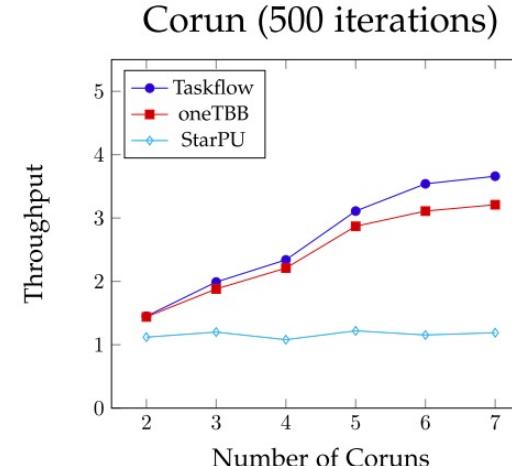
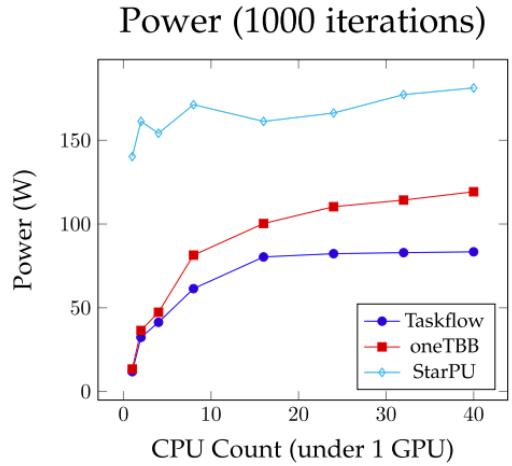
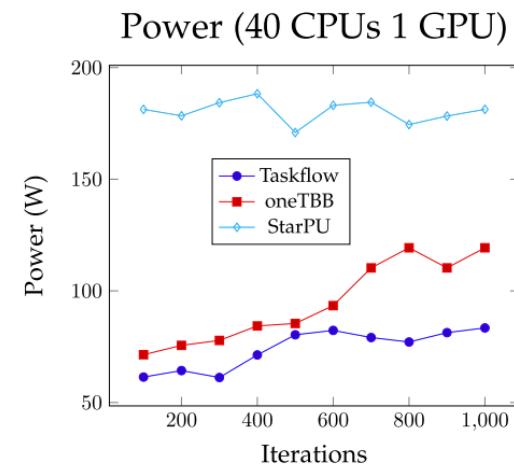
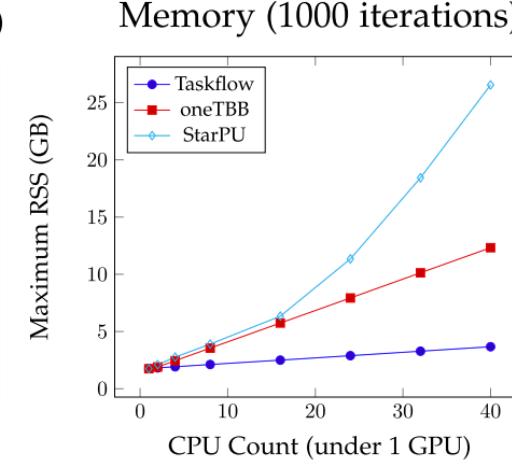
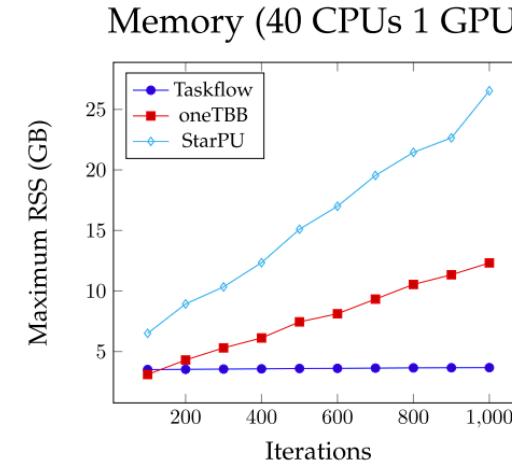
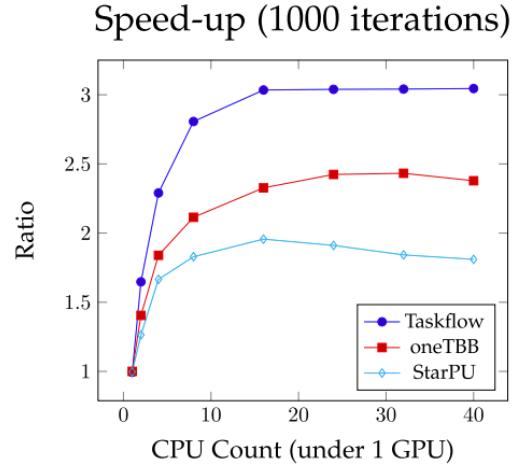
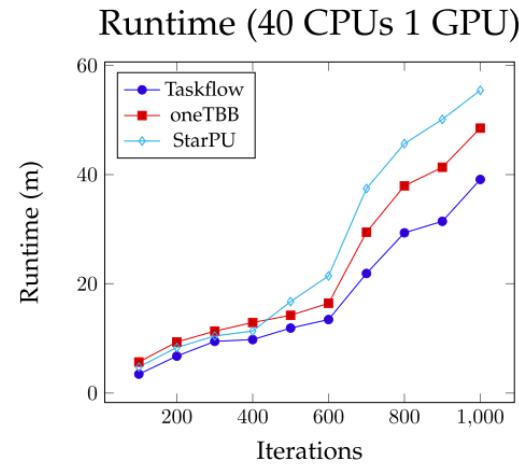
# Task-parallel Path Generation Algorithm<sup>1</sup>

- Applied Taskflow to accelerate path-based analysis on GPU
  - Ex: leon3mp (1.6M gates): **611x speed-up** over 1 CPU (**44x** over 40 CPUs)
  - **Best paper award** in ACM TAU 2021

Benchmark	#Pins	#Gates	#Arcs	OpenTimer Runtime	Our Algorithm #MDL=10		Our Algorithm #MDL=15		Our Algorithm #MDL=20	
					Runtime	Speed-up	Runtime	Speed-up	Runtime	Speed-up
leon2	4328255	1616399	7984262	2875783	4708.36	611×	5295.49ms	543×	5413.84	531×
leon3mp	3376821	1247725	6277562	1217886	5520.85	221×	7091.79ms	172×	8182.84	149×
netcard	3999174	1496719	7404006	752188	2050.60	367×	2475.90ms	304×	2484.08	303×
vga_lcd	397809	139529	756631	53204	682.94	77.9×	683.04ms	77.9×	706.16	75.3×
vga_lcd_iccad	679258	259067	1243041	66582	720.40	92.4×	754.35ms	88.3×	766.29	86.9×
b19_iccad	782914	255278	1576198	402645	2144.67	188×	2948.94ms	137×	3483.05	116×
des_perf_ispd	371587	138878	697145	24120	763.79	31.6×	766.31ms	31.5×	780.56	30.9×
edit_dist_ispd	416609	147650	799167	614043	1818.49	338×	2475.12ms	248×	2900.14	212×
mgc_edit_dist	450354	161692	852615	694014	1463.61	474×	1485.65ms	467×	1493.90	465×
mgc_matric_mult	492568	171282	948154	214980	994.67	216×	1075.90ms	200×	1113.26	193×

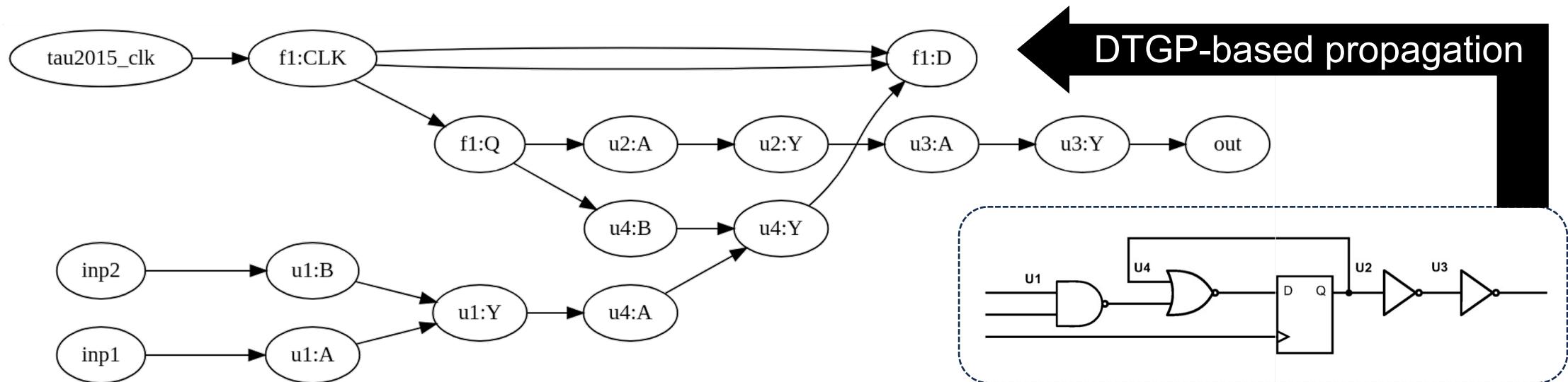
<sup>1</sup>: Guannan Guo, Tsung-Wei Huang, Yibo Lin, and Martin Wong, "GPU-accelerated Path-based Timing Analysis," *IEEE/ACM Design Automation Conference (DAC)*, CA, 2021

# Comparison with Existing TGP Systems



# Experimental Results – DTGP

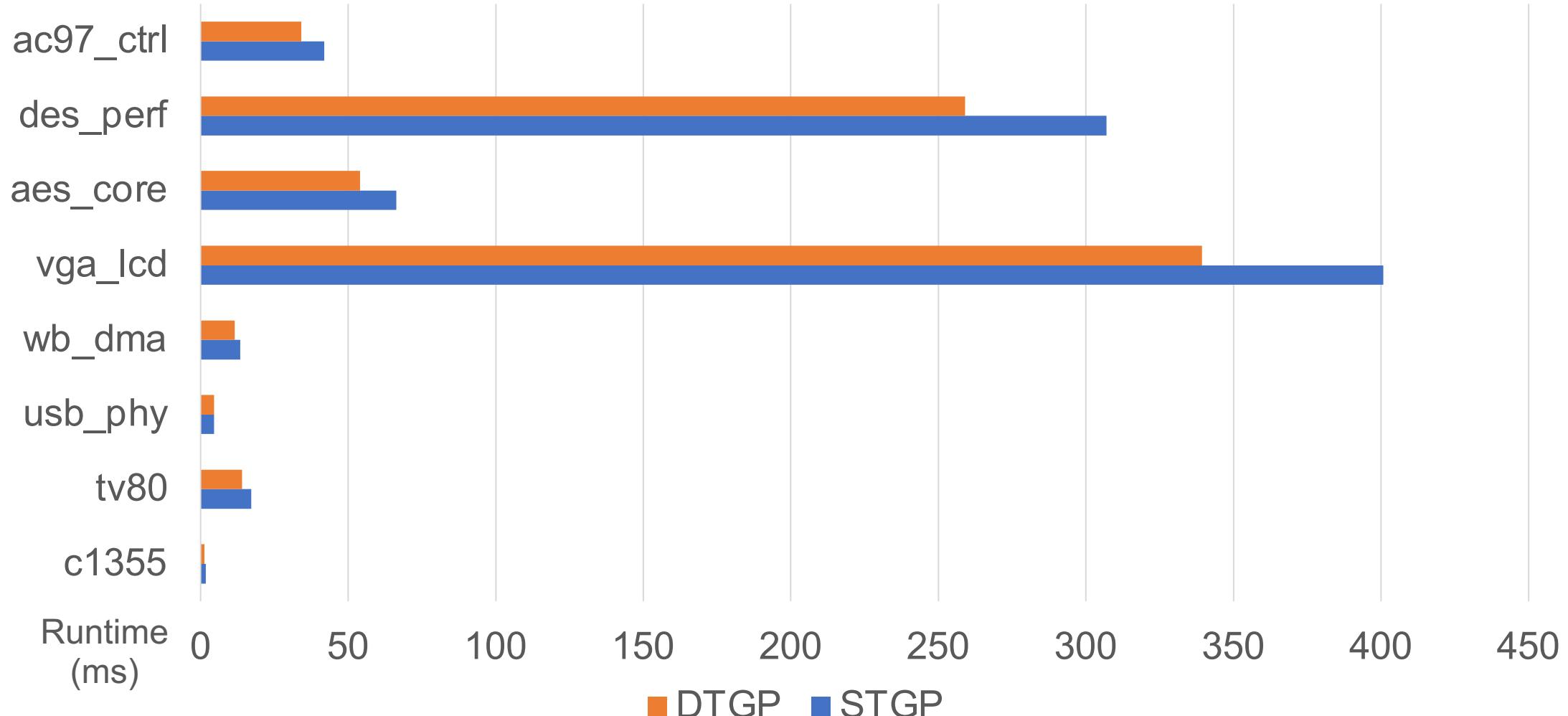
- Evaluated on a real-world static timing analysis application<sup>1</sup>
  - Formulated the timing propagation algorithm into a dynamic task graph
  - Ex (below): a task graph for a full-timing propagation on a five-gate circuit
    - Large circuits can compose millions of tasks and dependencies



<sup>1</sup>: T.-W. Huang, et. al, "OpenTimer v2: A New Parallel Incremental Timing Analysis Engine," *IEEE TCAD*, vol. 40, no. 4, pp. 776-789, April 2021

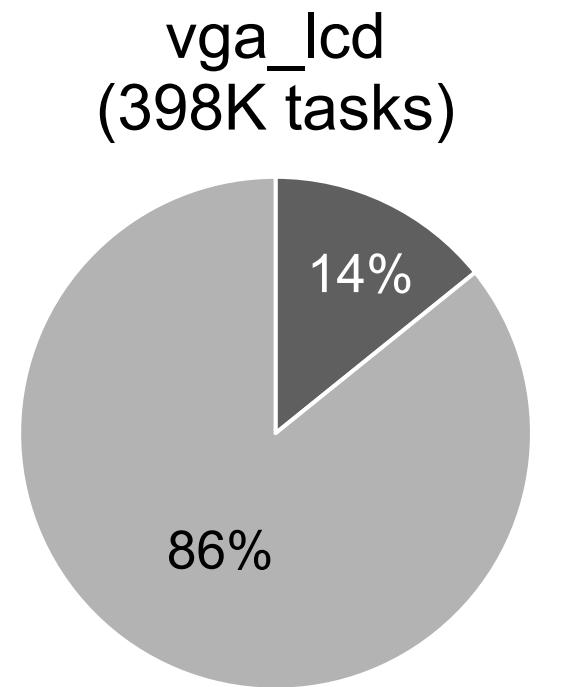
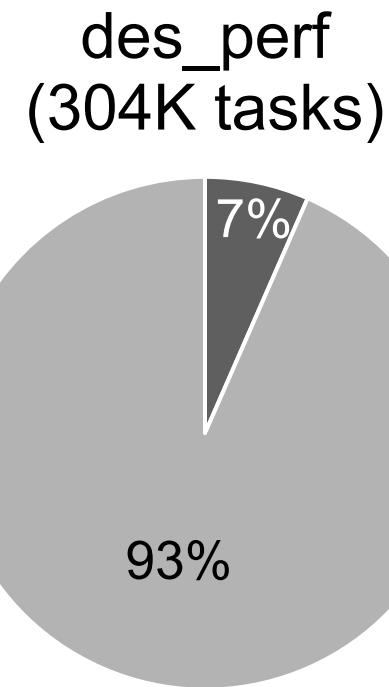
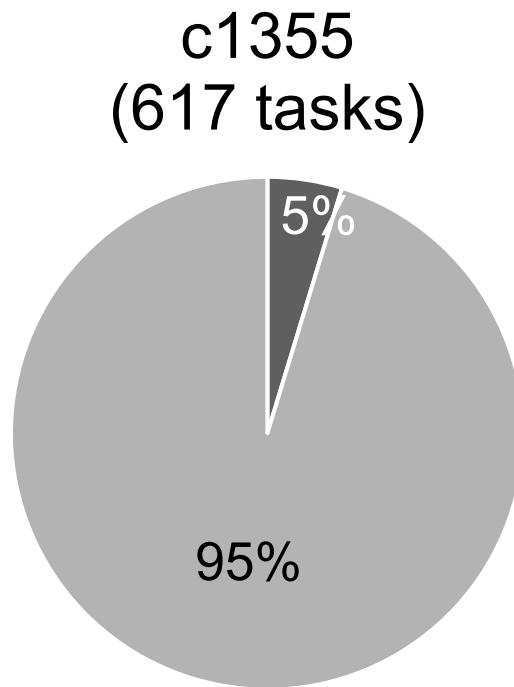


# Runtime Comparison: STGP vs DTGP



# Runtime Breakdown of STGP

- Graph construction time increases as the circuit size increases



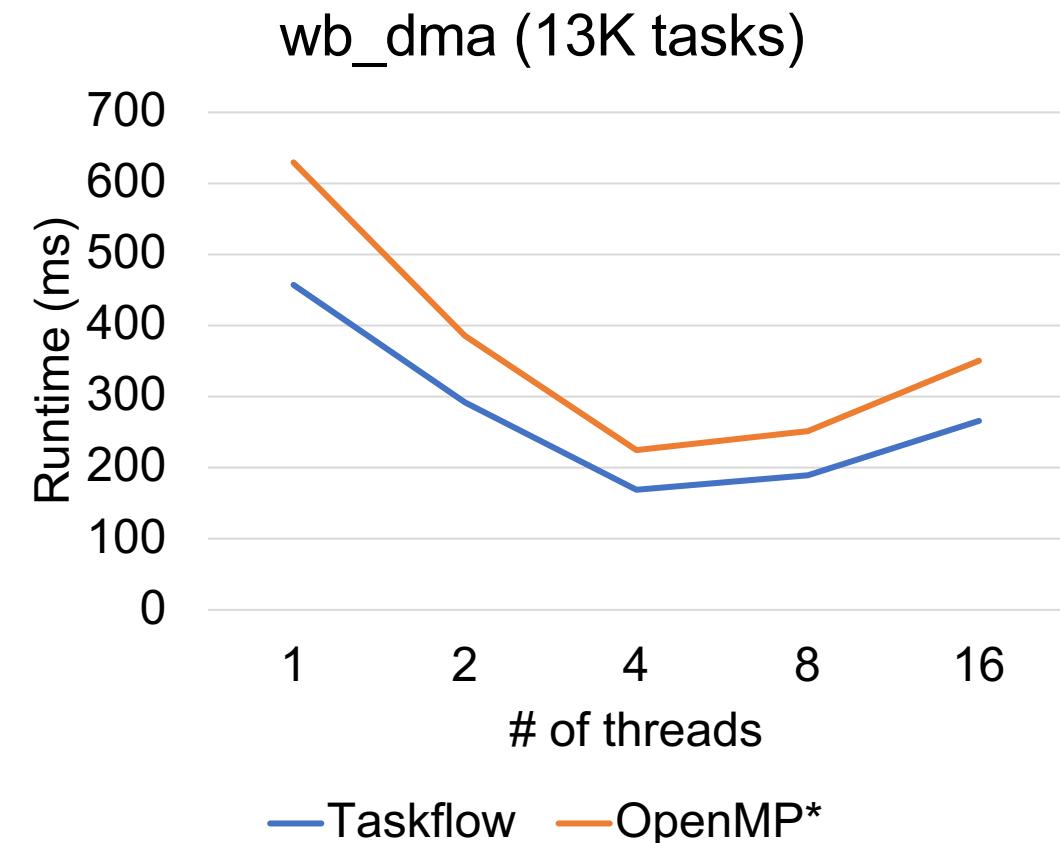
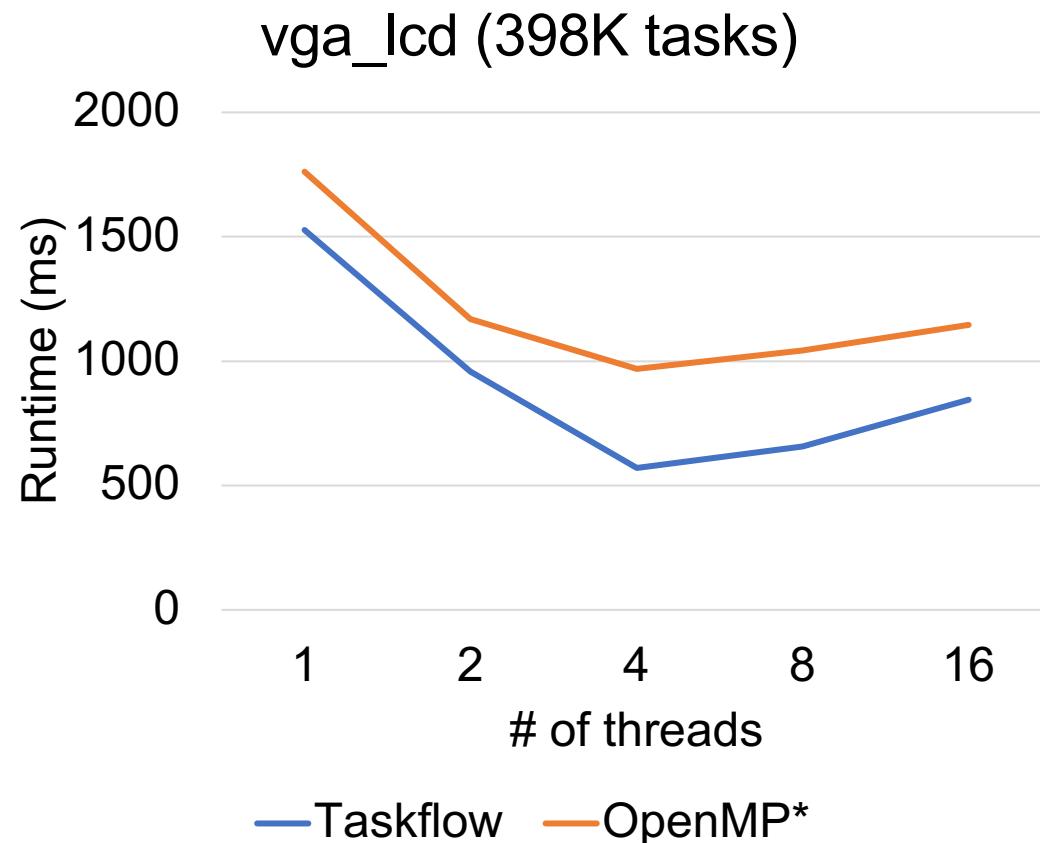
■ Build Graph ■ Run Graph

■ Build Graph ■ Run Graph

■ Build Graph ■ Run Graph

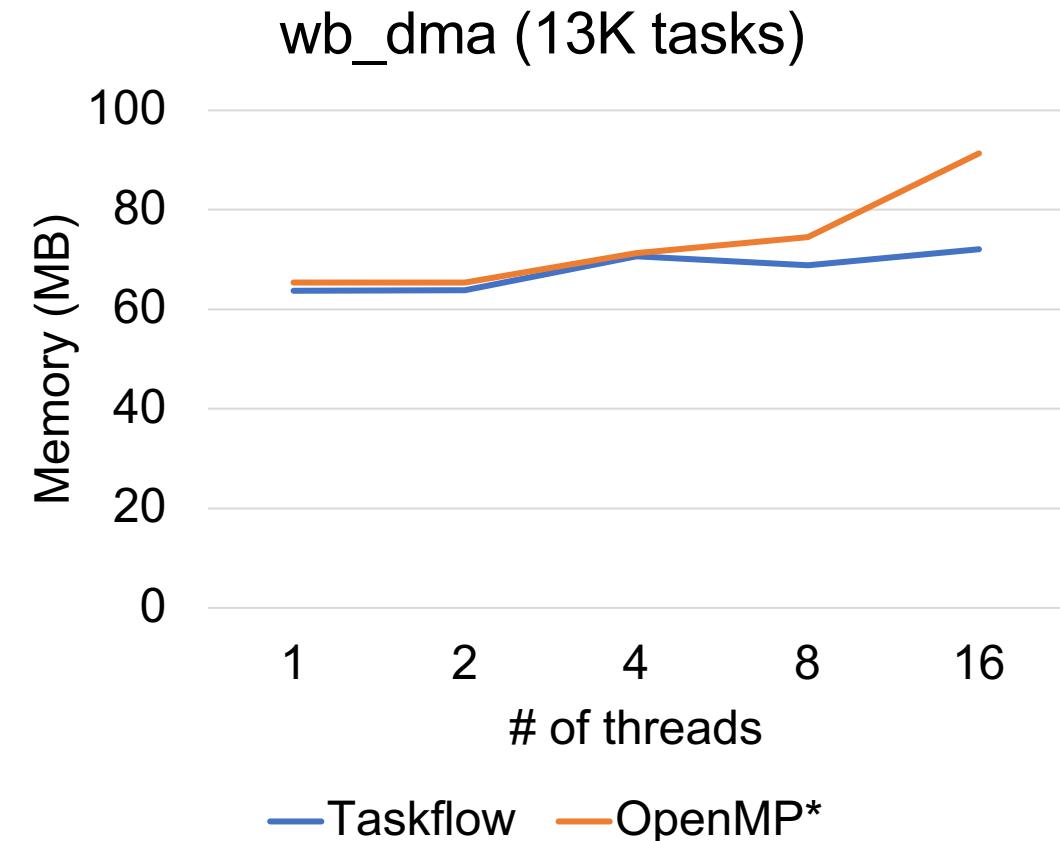
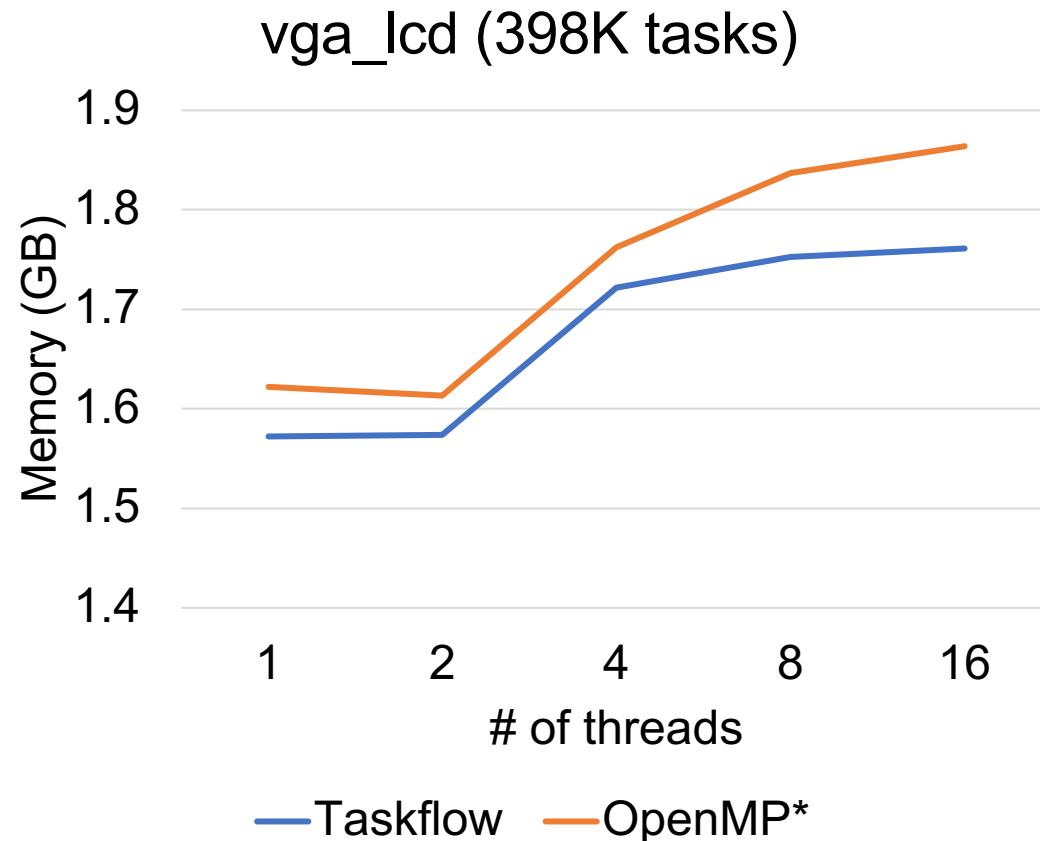
# Runtime Comparison with OpenMP

- Taskflow scales better than OpenMP with increasing # threads



# Memory Comparison with OpenMP\*

- Taskflow scales better than OpenMP with increasing # threads





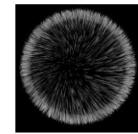
# Conclusion

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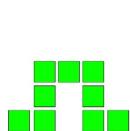
- Expressed your parallelism in the right way
- Programmed task graph parallelism using Taskflow
- Programmed dynamic task graph parallelism using Taskflow
- Overcame the scheduling challenges
- Showcased the efficiency of Taskflow in industrial application
- **Concluding the talk**



# Thank You for using Taskflow!



Explosion





# Thank you for Sponsoring Taskflow!

---



Google Summer of Code





# Questions?



Taskflow: <https://taskflow.github.io>

## Static task graph parallelism

```
// live: https://godbolt.org/z/j8hx3xnnx
tf::Taskflow taskflow;
tf::Executor executor;
auto [A, B, C, D] = taskflow.emplace(
    []() { std::cout << "TaskA\n"; },
    []() { std::cout << "TaskB\n"; },
    []() { std::cout << "TaskC\n"; },
    []() { std::cout << "TaskD\n"; });
A.precede(B, C);
D.succeed(B, C);
executor.run(taskflow).wait();
return 0;
```

## Dynamic task graph parallelism

```
// Live: https://godbolt.org/z/T87PrTax
tf::Executor executor;
auto A = executor.silent_dependent_async([](){
    std::cout << "TaskA\n";
});
auto B = executor.silent_dependent_async([](){
    std::cout << "TaskB\n";
}, A);
auto C = executor.silent_dependent_async([](){
    std::cout << "TaskC\n";
}, A);
auto D = executor.silent_dependent_async([](){
    std::cout << "TaskD\n";
}, B, C);
executor.wait_for_all();
```