

## LING 570 – HW2

### Q1 - Part (1)

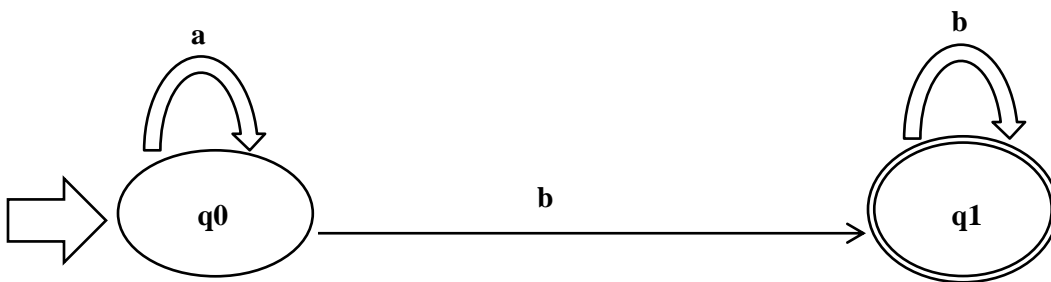
Consider the following DFA of  $D1 = (\Sigma, Q, q_0, F, \delta)$

$$\Sigma = \{a, b\}$$

$$Q = \{q_0, q_1\}$$

$$F = \{q_1\}$$

$$\delta = \left\{ \begin{array}{l} q_0 \times a \rightarrow q_0 \\ q_0 \times b \rightarrow q_1 \\ q_1 \times b \rightarrow q_1 \end{array} \right\}$$



The language accepted by D1 is  $\{ab, aab, abb, abbb, bbb, \dots\}$  or  $a^*b^+$

We can create the a regular grammar that generates the same language with the following mappings;

<u>DFA</u>	<u>Regular Grammar</u>
$\Sigma$	$\Sigma$
$q_0$	$S$
$\delta$	$P$
$Q$	$N$

So by substitution into a grammar  $G = (N, \Sigma, P, S)$

$$N = \{q_0, q_1\}$$

$$\Sigma = \{a, b\}$$

$$P = \left\{ \begin{array}{l} S \rightarrow a S \\ S \rightarrow b q_1 \\ q_1 \rightarrow b q_1 \\ q_1 \rightarrow \epsilon \end{array} \right\}$$

We can see that the same language  $L(G) = a^*b^+$  is also accepted by G.

### Q1 - Part (2)

Consider the grammar  $G = (\{q1, q2, q3\}, \{a, b\}, P, q1)$

$$P = \left\{ \begin{array}{ll} S & \rightarrow a S \\ S & \rightarrow \epsilon S \\ S & \rightarrow \epsilon q3 \\ S & \rightarrow b q2 \\ q2 & \rightarrow a q2 \\ q2 & \rightarrow a q3 \\ q2 & \rightarrow b q3 \\ q3 & \rightarrow a S \end{array} \right\}$$

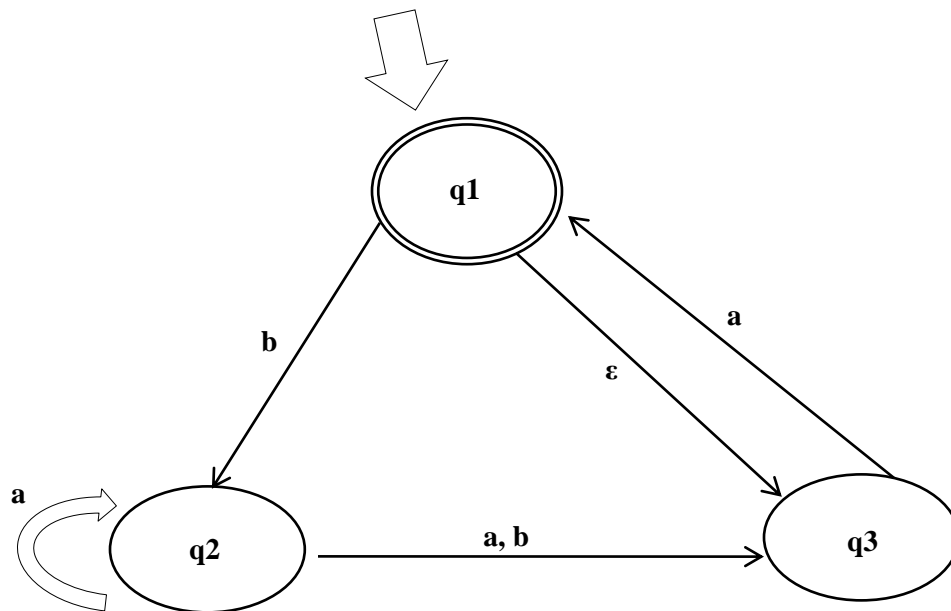
where  $L(G) = \{ \epsilon, a, baba, baa, \dots \}$ .

We can also map the grammar to a NFA of  $N1 = (Q, \Sigma, \delta, q0, F)$

<u>Regular Grammar</u>	<u>NFA</u>
$\{a, b\}$	$\Sigma$
$q1$	$F$
$P$	$\delta$
$\{q1, q2, q3\},$	$Q$

The mapping gives us  $N1 = (\{q1, q2, q3\}, \{a, b\}, \delta, q1, q1)$ .

Based on the transition function of  $P = \delta$  we can construct a NFA for  $N1$  below. In addition, the NFA also takes into account of the grammar showing transition leading to more than one state and transition happening without reading any inputs ( $\epsilon$ ).



## **Q2 - Part (2)**

Both commands below produced the same outcome of:

```
(0 -> 0 "they" : "PRO" / 1) (1 -> 1 "can" : "AUX" / 0.99) (2 -> 2 "fish" : "NOUN" / 0.7)
0.693
```

1. `carmel -k 1 fsa7 wfst1`
2. `cat wfst1_test | carmel -k 1 -sli wfst1`

These commands take in a FSA (fsa7) or a word-string sequence (wfst1\_test) as input, pass them through a weighted finite state transducer (wfst1) to produce a weighted FSA path that shows the top 1 most likely path consisting of tag sequences based on the input word-string. The ranking of the most likely path is based on the combined conditional probabilistic values for each pair of output tag and input word.

## **Q2 - Part (3)**

The command produced the following output:

```
(0 -> 0 "they" : "PRO" / 1) (1 -> 1 "can" : "AUX" / 0.99) (2 -> 2 "fish" : "NOUN" / 0.7)
0.693
(0 -> 0 "fish" : "NOUN" / 0.7) (1 -> 1 "they" : "PRO" / 1) (2 -> 2 "can" : "AUX" / 0.99)
0.693
```

The `-b` option accepts batch composition of input word-string sequences. The outcome is similar to what were described in Part (2) and that is output the result of the top 1 most likely path for each given input sequence. If there are 3 input sequences of word-string then the output would also consists of 3 lines, with each line output in the same order as the input sequence.

*End of HW2 – submitted by Wee Teck Tan*

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