Simplifying Mathematical Notations in Recurrence Equations

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1 Introduction

In the book "Introduction to Algorithms" written by CLRS, section 4.4, the following recurrence equation is given:

$$T(n) = \begin{cases} \Theta(1), & \text{if } n = 1\\ 3T(\lfloor n/4 \rfloor) + \Theta(n^2), & \text{otherwise} \end{cases}$$
 (1)

which later on was simplified as

$$T(n) = \begin{cases} \Theta(1), & \text{if } n = 1\\ 3T(\lfloor n/4 \rfloor) + cn^2, & \text{otherwise} \end{cases}$$
 (2)

we will argue that even if $\Theta(n^2)$ is a set, and $c*n^2$ is a scalar and they generally cannot be swapped one with the other, in this case it can be done and will help the equation to be more tractable.

2 Proof

Let's consider the general case only, so we have

$$3T(|n/4|) + \Theta(n^2)$$

which expanded is

$$3T(|n/4|) + f(n)$$
 where $f(n) \in \Theta(n^2)$

$$f(n)\in\Theta(n^2)$$
 means that
$$\exists~c_1>0,c_2>0,n_0>0~\text{such that}~c_1*n^2\leq f(n)\leq c_2*n^2~\forall n\geq n_0$$

So we have

$$3T(\lfloor n/4 \rfloor) + c_1 * n^2 \le 3T(\lfloor n/4 \rfloor) + f(n) \le 3T(\lfloor n/4 \rfloor) + c_2 * n^2 \forall n \ge n_0$$

$$= \underbrace{3T(\lfloor n/4 \rfloor) + c_1 * n^2}_{\alpha} \le T(n) \le \underbrace{3T(\lfloor n/4 \rfloor) + c_2 * n^2}_{\beta} \forall n \ge n_0$$

So if we resolve α only, we'll have a lower bound for T(n). If we resolve β only, we'll have an upper bound for T(n).

But α and β are asymptotically equivalent, that means that T(n) is asymptotycally equivalent to α and β too.

So we just need to resolve either α or β to obtain the asymptotic value of T(n), so we can just resolve for a generic c:

$$T(n) = \begin{cases} \Theta(1), & \text{if } n = 1\\ 3T(\lfloor n/4 \rfloor) + cn^2, & \text{otherwise} \end{cases}$$
 (3)

since using a generic constant won't change the asymptotic value.

3 Conclusions

The concept described in this document and the relative proof can be easily extended to any similar recurrence, and generalized.