

# 67865 - Mathematical Tools In Computer Science

February 20, 2024

## Topics:

The course is divided into 3 mini courses.

1. Probability. (Weeks 1-3)
  - (a) Rehearsal
  - (b) Useful Tools In research and CS
  - (c) week 4 = break
2. Advanced topics in linear algebra. (Weeks 5-7/8)
3. Optimization. (Week 8/9 - 11)

## Technicalities:

- Exercise every week.
- Final grade is the average of the best (n-3) exercises.
- Every week of miluim is excemption from 1 exercise. At least 4 exercises must be submitted.
- There is no exam, and there are no quizzes. :)
- The topics are going to be quite similar to previous years, so notes apply. If Amit does anything different than previous years, he will publish notes.

## Part I

# Probability

## 1 Rehearsal

### Today's Topics:

- Probability spaces
- Conditional Probability
- Independence
- Random Variables
- Expectation

### 1.1 Probability Spaces (Rehearsal)

**Definition 1.** (Probability Space/Merhav Histabrut) A probability space is a set  $\Omega$  (is finite or countable - in this course) and a function  $Pr : \Omega \rightarrow [0, 1]$  s.t.  $\sum_{\omega \in \Omega} Pr(\omega) = 1$

### Terminology

- $\omega \in \Omega$  is called an elementary/basic event.
- $A \subseteq \Omega$  is called an Event.
- $Pr$  is extended to events by  $Pr(A) = \sum_{\omega \in A} Pr(\omega)$
- $Pr$  is uniform if for any  $\omega \in \Omega$  it holds that  $Pr(\omega) = \frac{1}{|\Omega|}$

**Example.** Throwing a die:  $\Omega = \{1, 2, \dots, 6\}$ ,  $Pr(i) = \frac{1}{6}$   
Throwing a pair of dice:  $\Omega = \{1, \dots, 6\}^2$ ,  $Pr((i, j)) = \frac{1}{36}$

### 1.2 Conditional Probability (Rehearsal)

**Definition 2.** (Conditional Probability) Let  $A, B \subseteq \Omega$ , s.t.  $Pr(A) > 0$ . Define the conditional probability of  $B$  given  $A$ :  
$$Pr(B|A) = \frac{Pr(A \cap B)}{Pr(A)}$$

### 1.3 Independence (Rehersal)

**Definition 3.** (Independence)  $A, B \subseteq \Omega$  are independent if  $Pr(A \cap B) = Pr(A) \cdot Pr(B)$ .

Note: If  $Pr(A) > 0$ , then  $A, B$  are independent  $\iff P(B|A) = P(B)$

### 1.4 Random Variables (Rehersal)

**Definition 4.** (Random Variable) A random variable is a function  $X : \Omega \rightarrow \mathbb{R}$ .

#### Terminology

- If  $Range(X) \subseteq \{0, 1\}$  then  $X$  is called Indicator random variable.
- $(X + Y)(\omega) = X(\omega) + Y(\omega)$
- $(X \cdot Y)(\omega) = X(\omega) \cdot Y(\omega)$
- $(aX)(\omega) = aX(\omega)$
- $P(X = a) = P(\{\omega \in \Omega : X(\omega) = a\})$
- $P(X \leq Y) = P(\{\omega \in \Omega : X(\omega) \leq Y(\omega)\})$
- $P(X = 7 : Y = 3) = P(\{\omega \in \Omega : X(\omega) = 7\} | \{\omega \in \Omega : Y(\omega) = 3\})$
- $\vdots$

**Fact.** If  $X$  is an indicator R.V.  $\mathbb{E}(X) = Pr(X = 1)$ .

**Definition 5.** (Independent Random Variables) Random Variables  $X, Y$  are independent if for any  $a, b \in \mathbb{R}$  the events  $\{\omega \in \Omega : X(\omega) = a\}$  and  $\{\omega \in \Omega : Y(\omega) = b\}$  are independent.

### 1.5 Expectation (Rehersal)

**Definition 6.** (Expectation) The expectation of a random variable  $X$  is  $\mathbb{E}(X) = \sum_{\omega \in \Omega} X(\omega) \cdot Pr(X = \omega)$ .

**Example.**

$$\Omega = \{1, 2, 3\} \times \{1, 2, 3\}. Pr((i, j)) = \frac{1}{9}.$$

$$X((i, j)) = i + j$$

$$\mathbb{E}(X) = \frac{1}{9}(X((1, 1)) + X((1, 2)) + \dots + X((3, 3))) = 4$$

**Lemma.** (Linearity Of Expectation) Let  $X, Y$  be random variables,  $a, b \in \mathbb{E}$ . then:

$$\mathbb{E}(aX + bY) = a\mathbb{E}(X) + b\mathbb{E}(Y)$$

**Example.** The expected number of fixed points in a random permutation.

- A fixed point in a permutation  $\pi : [n] \rightarrow [n]$  (1:1 function) is  $i \in [n]$  s.t.  $\pi(i) = i$ .
- $S_n = \{(\pi : [n] \rightarrow [n]) : \pi \text{ is a permutation}\}$

$$Pr(\pi) = \frac{1}{n!}$$

$X(\pi) = |\{i : \pi(i) = i\}|$  (number of fixed points of  $\pi$ ) This is indeed a random variable.  $X : S_n \rightarrow \mathbb{R}$ .

Question: What is the expected value of  $X$ .

For every  $i \in [n]$ , we define  $X_i(\pi) = \begin{cases} 1 & \pi(i) = i \\ 0 & \text{otherwise} \end{cases}$ .

So  $X = \sum_{i=1}^n X_i$

$$\Rightarrow \mathbb{E}(X) \stackrel{*1}{=} \sum_{i=1}^n \mathbb{E}(X_i) \stackrel{*2}{=} \sum_{i=1}^n Pr(X_i = 1) = \sum_{i=1}^n \frac{1}{n} = 1$$

Where  $*1$  is from linearity of expectation and  $*2$  is from  $X_i$ 's being indicator random variables.

## 1.6 Inequalities (Rehearsal)

### 1.6.1 Union Bound

**Lemma.** (Union Bound) If  $A, \dots, A_k \subseteq \Omega$  then:

$$Pr\left(\bigcup_{i=1}^k A_i\right) \leq \sum_{i=1}^k Pr(A_i)$$

### 1.6.2 Markov

**Theorem.** (Markov Inequality) Let  $X : \Omega \rightarrow \mathbb{R}_+$  (NON NEGATIVE!) be a R.V. then for any  $c > 0$ :

$$Pr(X \geq c\mathbb{E}(X)) \leq \frac{1}{c} \iff Pr(X \geq c) \leq \frac{\mathbb{E}(X)}{c}$$

Proof.

$$\begin{aligned}
\mathbb{E}[X] &\stackrel{def}{=} \sum_{\omega \in \Omega} Pr(\omega) \cdot X(\omega) \\
&\stackrel{*_1}{\geq} \sum_{\substack{\omega \in \Omega \\ X(\omega) \geq c\mathbb{E}[X]}} Pr(\omega) \cdot X(\omega) \\
&\geq \sum_{\substack{\omega \in \Omega \\ X(\omega) \geq c\mathbb{E}[X]}} Pr(\omega) \cdot c\mathbb{E}[X] \\
&= c\mathbb{E}[X] \cdot \sum_{\substack{\omega \in \Omega \\ X(\omega) \geq c\mathbb{E}[X]}} Pr(\omega) \\
&\stackrel{def}{=} c\mathbb{E}[X] \cdot Pr(X \geq c\mathbb{E}[X])
\end{aligned}$$

\*<sub>1</sub> is because  $X$  is non negative.

$$\Rightarrow Pr(X \geq c\mathbb{E}(X)) \leq \frac{1}{c}$$

□

### 1.6.3 Chebyshev

**Definition 7.** (Variance and standard deviation) The Variance of R.V.  $X$  is:

$$Var(X) = \mathbb{E}(X - \mathbb{E}(X))^2 = \mathbb{E}(X^2) - \mathbb{E}(X)^2$$

The standard deviation of  $X$  is  $\sigma = \sqrt{Var(X)}$ .

**Fact.**

$$\begin{aligned}
Var(X) &= \mathbb{E}(X - \mathbb{E}(X))^2 \\
&= \mathbb{E}(X^2 - 2X\mathbb{E}(X) + \mathbb{E}(X)^2) \\
&= \mathbb{E}(X^2) - \mathbb{E}(2X\mathbb{E}(X)) + \mathbb{E}(X)^2 \\
&= \mathbb{E}(X^2) - 2\mathbb{E}(X)\mathbb{E}(X) + \mathbb{E}(X)^2 \\
&= \mathbb{E}(X^2) - \mathbb{E}(X)^2
\end{aligned}$$

**Fact.**  $Var(X) \leq \mathbb{E}(X^2)$  ( $\leftarrow$  called second moment)

**Theorem.** (Chebychev inequality) Let  $X$  be a R.V. then:

$$Pr\left(|X - \mathbb{E}(X)| \geq c\sqrt{Var(X)}\right) \leq \frac{1}{c^2} \iff Pr(|X - \mathbb{E}(X)| \geq c) \leq \frac{Var(X)}{c^2}$$

Proof. Denote  $Y(\omega) = (X - \mathbb{E}(X))^2$ . By markov inequality:

$$Pr(Y \geq c^2 \mathbb{E}(Y)) \leq \frac{1}{c^2}$$

$$\begin{aligned} Pr(Y \geq c^2 \mathbb{E}(Y)) &= Pr((X - \mathbb{E}(X))^2 \geq c^2 Var(X)) \\ &= Pr(X - \mathbb{E}(X) \geq c\sqrt{Var(X)}) \end{aligned}$$

□

## 2 Ramsey Theory & Probabilistic Method

### 2.1 Ramsey Theory

Ramsey theory tries to find the largest possible “island” of “order” in a “chaos”.

**Definition 8.** A Coloring of a graph  $G = (V, E)$  is a function  $c : E \rightarrow \{Red(R), Black(B)\}$ .

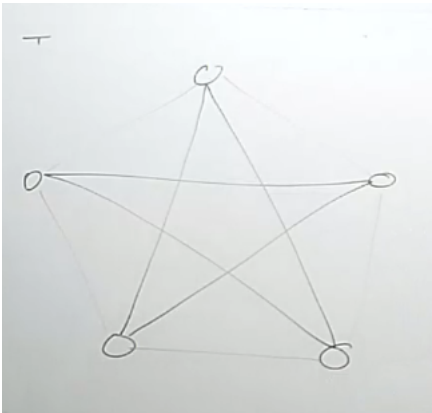
**Definition 9.**  $S \subseteq V$  is Monochromatic w.r.t coloring  $c$ , if all edges whose both vertices are in  $S$  are in the same colors.

**Definition 10.** The Ramsey Number  $R(k)$  is the minimal number  $R$  s.t. for each coloring of  $K_R$  (The complete graph of  $R$  vertices) there is a monochromatic set  $S \subseteq V$  of size  $k$ .

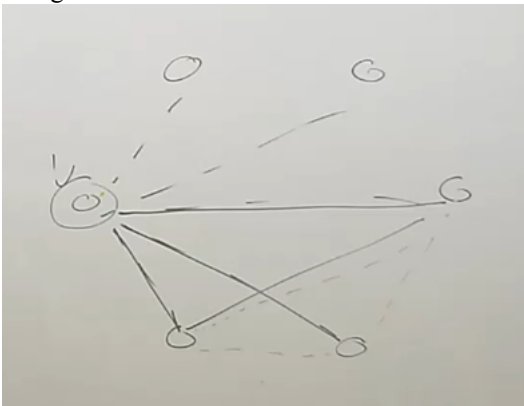
Claim.  $R(2) = 2$  (trivial)

Claim.  $R(3) = 6$

Proof. first, we can see that  $R(3) > 5$  because of the following counter example.



Now we will show that  $R(3) \leq 6$ . We look at a verticie. 3 of it's edges have to be either black or red. assume they are red, now look at the 3 edges connecting the other vrtices of the red edges. if they are all black we won. if one of them is red there is a red triangle and we won.



□

Claim.  $R(4) = 18$  (hard, discovered 1978)

Claim.  $R(5) = \text{Unknown}$  (between 43 and 48)

**Theorem.** (Ramsey)

$$2^{\frac{k}{2}-1} < R(k) \leq 2^{2k}$$

Proof. Upper bound in the tirgul: Generalization of  $R(3) \leq 6$ .

Lower Bound: Using the **probabilstic method**.

Denote  $N = 2^{\frac{k}{2}-1}$ . Let  $C$  be a random coloring of  $K_N$  (The complete graph of  $N$  vertices). Denote  $A$  the event that there is a mono set of size  $k$ :  $A = \{c : E \rightarrow \{R, B\} : \text{there is a mono set of size } k\}$

is is enough to show that  $Pr(A) < 1$ .

For  $S \subseteq V$  s.t.  $|S| = k$ , Denote  $A_S = \{c : S \text{ is mono w.r.t } c\}$ .

As such  $A = \bigcup_{S \subseteq V, |S|=k} A_S$ .

$Pr(A_s) = 2^{-\binom{k}{2}+1}$  because there are  $\binom{k}{2}$  edges formed in  $S$ . and they all need to be the same color. the  $+1$  is because they can all be either red or black.

By the union bound:

$$\begin{aligned}
 Pr(A) &\leq \sum_{S \subseteq V, |S|=k} Pr(A_s) \\
 &= \sum_{S \subseteq V, |S|=k} 2^{-\binom{k}{2}+1} \\
 &= \binom{N}{k} \cdot 2^{-\binom{k}{2}+1} \\
 &\leq N^k \cdot 2^{-\binom{k}{2}+1} \\
 &= 2^{(\frac{k}{2}-1)k} \cdot 2^{-\binom{k}{2}+1} \\
 &= 2^{\frac{k^2}{2}-k-\frac{k(k-1)}{2}+1} \\
 &= 2^{-\frac{k}{2}+1} \\
 &\stackrel{*}{<} 1
 \end{aligned}$$

\* is for  $k \geq 3$ , for  $k = 2, 1$  the case is trivial.

□