Excercise 4 Implementing a centralized agent

Group №: 90 Kyle Gerard, Yann Bolliger

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1 Solution Representation

1.1 Variables

Instead of representing a vehicle's journey as a sequence of tasks, we chose to represent it as a sequence of pickup and delivery actions. Each task $t \in \mathcal{T}$ has one action of both types.

$$\mathcal{P} = \{pickup(t) : t \in \mathcal{T}\}, \ \mathcal{D} = \{delivery(t) : t \in \mathcal{T}\}, \ \mathcal{A} = \mathcal{P} \cup \mathcal{D}\}$$

This accounts for the fact that a vehicle can carry multiple tasks at a time if there are two pickups in a row. The following variables define the first pickup of each vehicle (where \mathcal{V} is the set vehicles).

$$\forall v \in \mathcal{V}: firstPickup(v) \in \mathcal{P} \cup \{\text{null}\}\$$

If the variable is null this means the vehicle does not accomplish any actions. All subsequent actions of a vehicle are defined by the next set of variables:

$$\forall a \in \mathcal{A} : nextAction(a) \in \mathcal{A} \cup \{null\}$$

where again the null signifies that a vehicle has no further actions to perform. We will define two other sets of variables which will clarify the former:

$$\forall a \in \mathcal{A} : vehicle(a) \in \mathcal{V}; \ \forall a \in \mathcal{A} : time(a) \in \mathbb{N}$$

The vehicle variables define which vehicle carries out a certain action. This can be derived from the firstPickup action at the start of each action chain defined by nextAction. (For example if firstPickup(v) = a, nextAction(a) = b then vehicle(a) = vehicle(b) = v.)

The second variable can also be derived from the action chains. It simply gives the rank of each action in the chain (for example if firstPickup(v) = a, nextAction(a) = b then time(a) = 1, time(b) = 2).

1.2 Constraints

As explained before, the action chain for each vehicle define the time and the vehicle for each action:

$$firstPickup(v) = a \Rightarrow time(a) = 1; \ nextAction(b) = c \Rightarrow time(c) = time(b) + 1$$

 $firstPickup(v) = a \Rightarrow vehicle(a) = v; \ nextAction(b) = c \Rightarrow vehicle(c) = vehicle(b)$

Additionally, the same vehicle must *pickup* and *deliver* a task. It has to pickup the task before it delivers it, of course, and **each task must be picked up and delivered**:

$$\forall a \in \mathcal{A} : nextAction(a) \neq a, nextAction(a) = null \Rightarrow a \in \mathcal{D}$$

$$\forall t \in \mathcal{T} : vehicle(pickup(t)) = vehicle(delivery(t)) \in \mathcal{V}$$
$$\forall t \in \mathcal{T} : time(pickup(t)) < time(delivery(t))$$

$$\forall t \in \mathcal{T} \exists \{pickup(t), delivery(t)\} \subset \mathcal{A} \text{ and } \forall a \in \mathcal{A} \exists v \in \mathcal{V} : vehicle(a) = v$$

Last but not least, at each time τ a vehicle v can never carry more weight than its capacity.

$$carriedTasks(\tau, v) = \{t \in \mathcal{T} : vehicle(pickup(t)) = v \land time(pickup(t)) < \tau \land time(delivery(t)) > \tau\};$$

$$\forall \tau \in \mathbb{N} : \sum_{t \in carriedTasks(\tau,v)} weight(t) \leq capacity(v)$$

- 1.3 Objective function
- 2 Stochastic optimization
- 2.1 Initial solution
- 2.2 Generating neighbours
- 2.3 Stochastic optimization algorithm
- 3 Results
- 3.1 Experiment 1: Model parameters
- 3.1.1 Setting
- 3.1.2 Observations
- 3.2 Experiment 2: Different configurations
- 3.2.1 Setting
- 3.2.2 Observations