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UNIVERSAL GÖDEL STATEMENTS AND COMPUTABILITY OF INTELLIGENCE

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ABSTRACT. We show that there is a mathematical obstruction to complete Turing computability of intelligence. This obstruction can be circumvented only if human reasoning is fundamentally unsound, with the latter formally interpreted here as certain stable soundness. To this end, we develop in a specific setting an analogue of a Gödel statement for stably sound Turing machines, as defined here. This stably Gödel statement \mathcal{G} in addition has a certain universality with respect to a certain class of Turing machines / formal systems. Moreover, this \mathcal{G} is constructed explicitly, given the general form of our class of Turing machines. The main argument is implicitly based on set theory, in contrast to the partially meta-logical arguments of Gödel and Penrose, arguing for a similar obstruction. For this reason we are able to phrase our results as theorems.

In what follows we understand *human intelligence* very much like Turing in [2], as a black box which receives inputs and produces outputs. More specifically, this black box B is meant to be some system which contains a human subject. We do not care about what is happening inside B . So we are not directly concerned here with such intangible things as understanding, intuition, consciousness - the inner workings of human intelligence that are supposed as special. The only thing that concerns us is what output B produces given an input, not how it is produced. Given this *very* limited interpretation, the question that we are interested in is this:

Question 1. Can human intelligence be completely modelled by a Turing machine?

An informal definition of a Turing machine (see [1]) is as follows: it is an abstract machine which permits certain inputs, and produces outputs. The outputs are determined from the inputs by a fixed finite algorithm, defined in a certain precise sense. For a non-expert reader we point out that this “fixed” does not preclude the algorithm from “learning”,¹ it just means that how it “learns” is completely determined by the initial algorithm. In particular anything that can be computed by computers as we know them can be computed by a Turing machine. For our purposes the reader may simply understand a Turing machine as a digital computer with unbounded memory running some particular program. Unbounded memory is just a mathematical convenience. In specific arguments, also of the kind we make, we can work with non-explicitly bounded memory. Turing himself has started on a form of Question 1 in his “Computing machines and Intelligence”, [2], where he also informally outlined a possible obstruction to a yes answer coming from Gödel’s incompleteness theorem.

For the incompleteness theorem to have any relevance, we need some assumption on the soundness or consistency of human reasoning. Informally, a human is sound if whenever they asserts something in absolute faith, this something is indeed true. This requires context, as truth in general is undefinable. For our arguments later on the context will be in certain mathematical models. However, we cannot honestly hope for soundness, as even mathematicians are not on the surface sound at all times, they may assert mathematical untruths at various times, (but usually not in absolute faith). But we can certainly hope for some kind of fundamental soundness.

In this work we will formally interpret fundamental soundness as stable soundness. Essentially, our machine² B is now allowed to make corrections, and if a statement printed by B is never corrected, then this statement is true, if B has our stable soundness property. The negation of stably sound is stated as either stably unsound or not stably sound, synonymously. This stable soundness reflects our basic understanding of how science progresses. Of course even stable soundness needs idealizations to

¹In the sense of “machine learning”.

²Here we use the term machine as an abstraction for a process acting on inputs, but it need not be a computational process, in contrast to Turing machines.

make sense for humans. The human brain deteriorates, and eventually fails, so that either we idealize the human brain to never deteriorate, or B now refers not to an individual human but to the evolving scientific community.

Around the same time as Turing, Gödel argued for a no answer to Question 1, see [13, 310], relating the question to existence of absolutely undecidable problems, see also Feferman [8], and Koellner [15], [16] for a discussion. Since existence of absolutely undecidable problems is such a difficult and contentious issue, even if Gödel’s argument is in essence correct it is not completely compelling. Interestingly, for Gödel, fundamental unsoundness of human reasoning is not even a possibility, it does not seem to be stated in [13]. A more in depth analysis of Gödel vs Turing on computability and the mind appears for example in [5]. Later Lucas [12] and later again and more robustly Penrose [20] argued for a no answer based only on soundness, and by further elaborating the obstruction from the Gödel incompleteness theorem. Such an argument if correct would be much more compelling.

It should also be noted that for Penrose in particular, non-computability of intelligence is evidence for new physics, and he has specific and *very* intriguing proposals with Hameroff [11], on how this can take place in the human brain. Here is also a partial list of some partially related work on mathematical models of brain activity and or quantum collapse models: [14], [17], [9], [10].

From my perspective, one of the main issue with the Penrose argument concerns the soundness assumption, and we will review this in this introduction. Other authors, particularly Koellner [15], [16], argue that there are meta-logical issue even allowing for soundness. Following a technically very different approach to Penrose, we intend to completely resolve these issues here. In particular there should be no meta-logical issues because our results are theorems, and not just meta-theorems, cf. Section 6.

The following is a slightly informal version of our main Theorem 4.3, taking the view that our idealized human is represented by the evolving scientific community, or just H for short. The sole formal property that we postulate for H , what we call the stable Penrose property will be explained in Section 4. It is analogous to the arithmetic Penrose property we describe just below in the outline of the Penrose argument.

Theorem 0.1. *Either there are cognitively meaningful, absolutely non Turing computable processes in the human brain, or human beings are fundamentally unsound, meaning specifically that H is stably unsound. This theorem is indeed a mathematical fact,³ given our formalization of H and of stable soundness.*

By *absolutely* we mean in any physical model. Note that even existence of absolutely non Turing computable processes in nature is not known. For example we expect beyond reasonable doubt that solutions of fluid flow or N -body problems are generally non Turing computable, (over \mathbb{Z} , if not over \mathbb{R} cf. [3]), as modeled in essentially classical mechanics. But in a more physically accurate and fundamental model they may both become computable, possibly if the nature of the universe is ultimately discreet. It would be good to compare this theorem this with Deutch [7], where computability of any suitably finite and discreet physical system is conjectured. Although this is not immediately at odds with us, as the hypothesis of that conjecture may certainly not be satisfiable.

By strengthening the hypothesis of Theorem 0.1, from computability to provable computability as in Theorem 5.2, we can obtain more practical consequences. To the effect that not only is H stably unsound, but must in fact eventually stably assert $0 = 1$.

0.1. The Penrose argument. Following Lucas [12], Penrose has given variations of the argument for a no answer to Question 1 in his books [18], [19]. The final argument can be found in [20]. Which we now re-interpret in a language closer to our subsequent argument. The following argument (really just an outline) is more elaborate then what was originally proposed by Penrose, but this is because we make some additional things explicit.

³Specifically a theorem of set theory, although we keep set theory implicit as usual.

While this outline uses some of the language of formal systems, we will *not* use this language in our main argument, which is based purely on the language of Turing machines, and is much more elementary.

Suppose a human subject P is in contact with experimenter/operator E . The input strings that E gives P are pairs (Σ_T, n) for Σ_T specification of a Turing machines T , and $n \in \mathbb{N}$. The output $P(\Sigma_T, n)$ is a statement of arithmetic printed by P . (For now there is no requirement on truth.)

Given a Turing machine T let Θ_T be the statement:

$$(0.2) \quad T \text{ computes } P.$$

We ask that for each fixed T : $\{P(\Sigma_T, n)\}_n$ is the complete list of statements that P asserts to be true conditionally on Θ_T . We also ask that P satisfies what we call the arithmetic Penrose property. This says that P asserts for each T the statement I_T :

$$(0.3) \quad \Theta_T \implies T \text{ is sound,}$$

which means that $T(\Sigma_T, n)$ is true for each n . We ignore for now whether P making such an assertion is rational.

Let then T_0 be a specified Turing machine, and suppose that E passes to P input of the form (Σ_{T_0}, n) . Now, as is well known⁴, the statements $\{T_0(\Sigma_{T_0}, n)\}_n$ ⁵ must be the complete list of provable statements in a certain formal system $\mathcal{F}(T_0)$ explicitly constructible given T_0 . Loosely, a formal system consists of a language: alphabet and grammar, a collection of sentences in this language understood as axioms, and finally a deductive system.

By construction $\mathcal{F}(T_0)$ would be sound if Θ_{T_0} and if I_{T_0} . In particular if $\Theta_{T_0} \wedge I_{T_0}$ then by the celebrated Gödel incompleteness theorem there would be a true Gödel statement $G(T_0)$ for this $\mathcal{F}(T_0)$, such that

$$T_0(\Sigma_{T_0}, n) \neq G(T_0), \quad \text{for all } n.$$

But P asserts I_{T_0} , hence he must assert by implication that

$$\Theta_{T_0} \implies G(T_0).$$

And so assuming P knows how to construct $G(T_0)$ then this statement must be in the list $\{P(\Sigma_{T_0}, n)\}_n$, and so in the list $\{T_0(\Sigma_{T_0}, n)\}_n$, so we would get a contradiction. So we conclude that either not Θ_{T_0} , that is P is not computed by T_0 or P is not consistent, but T_0 is arbitrary so we obtain an obstruction to computability of P .

Even if it was in essence correct the above argument is unsatisfactory because all it claims to prove is: either we are non-computable or inconsistent, which we appear to be anyway. Of course as we have argued we must talk of fundamental soundness, interpreted here as stable soundness. But then the argument cannot work exactly as above, since Gödel's theorem necessitates total consistency. This is a highly non-trivial issue, as while we can abstractly extract from a stably sound machine an absolutely sound machine, the latter may not be computable even if the former was, cf. Remark 3.2.

We will delve no further into critiquing the Lucas-Penrose argument. One such critique is given in Koellner [15], [16], see also Penrose [20], and Chalmers [4] for discussions of some issues. Note of course that our version of the Penrose argument is slightly different, and so the issues might be different.

So motivated by the discussion above, the ideal thing to do is to formally define stable soundness and construct a new type of Gödel statements, which works under this weaker hypothesis. This is actually what we will do, in the limited setting above. To this end, we reformulate the above idea using a more elementary approach, more heavily based in Turing machines. We first isolate a certain class of Turing machines that we name diagonalization machines. They print strings with a certain property C . As the name suggests, their behavior is related to the Cantor diagonalization argument. Next we explicitly construct a ‘‘Gödel string’’ \mathcal{G} which is universal for this whole class. This string \mathcal{G} has property C but cannot be printed by a Turing diagonalization machine. Crucially, this is then

⁴I don't know a standard reference but see for example [8].

⁵Strictly speaking after taking deductive closure.

extended to stable diagonalization machines, which print property C^6 strings only stably. Given this, our main result follows by an argument similar to the one in the outline above.

This is essentially as far as we can go in trying to outline the argument, as most of it just concerns the construction of the class of diagonalization machines and of \mathcal{G} , and this is hard to describe without details. However, technically the paper is mostly elementary, and should be widely readable in entirety.

1. SOME PRELIMINARIES

This section can be just skimmed on a first reading. Really what we are interested in is not Turing machines per se, but computations that can be simulated by Turing machine computations. These can for example be computations that a mathematician performs with paper and pencil, and indeed is the original motivation for Turing's specific model. However to introduce Turing computations we need Turing machines, here is our version which is a computationally equivalent, minor variation of Turing's original machine.

Definition 1.1. *A Turing machine M consists of:*

- *Three infinite (1-dimensional) tapes T_i, T_o, T_c , (input, output and computation) divided into discreet cells, next to each other. Each cell contains a symbol from some finite alphabet Γ . A special symbol $b \in \Gamma$ for blank, (the only symbol which may appear infinitely many often).*
- *Three heads H_i, H_o, H_c (pointing devices), H_i can read each cell in T_i to which it points, H_o, H_c can read/write each cell in T_o, T_c to which they point. The heads can then move left or right on the tape.*
- *A set of internal states Q , among these is "start" state q_0 . And a non-empty set $F \subset Q$ of final states.*
- *Input string Σ : the collection of symbols on the tape T_i , so that to the left and right of Σ there are only symbols b . We assume that in state q_0 H_i points to the beginning of the input string, and that the T_c, T_o have only b symbols.*
- *A finite set of instructions: I , that given the state q the machine is in currently, and given the symbols the heads are pointing to, tells M to do the following, the taken actions 1-3 below will be (jointly) called an **executed instruction set**, or just **step**:*
 - (1) *Replace symbols with another symbol in the cells to which the heads H_c, H_o point (or leave them).*
 - (2) *Move each head H_i, H_c, H_o left, right, or leave it in place, (independently).*
 - (3) *Change state q to another state or keep it.*
- *Output string Σ_{out} , the collection of symbols on the tape T_o , so that to the left and right of Σ_{out} there are only symbols b , when the machine state is final. When the internal state is one of the final states we ask that the instructions are to do nothing, so that these are frozen states.*

Definition 1.2. *A complete configuration of a Turing machine M or total state is the collection of all current symbols on the tapes, position of the heads, and current internal state. Given a total state s , $\delta(s)$ will denote the successor state of s , obtained by executing the instructions set of M on s , or in other words $\delta(s)$ is one step forward from s .*

So a Turing machine determines a special kind of function:

$$\delta^M : \mathcal{C}(M) \rightarrow \mathcal{C}(M),$$

where $\mathcal{C}(M)$ is the set of possible total states of M .

Definition 1.3. *A Turing computation, or computation sequence for M is a possibly not eventually constant sequence*

$$*M(\Sigma) := \{s_i\}_{i=0}^{i=\infty}$$

of total states of M , determined by the input Σ and M , with s_0 the initial configuration whose internal state is q_0 , and where $s_{i+1} = \delta(s_i)$. If elements of $\{s_i\}_{i=0}^{i=\infty}$ are eventually in some final machine state,

⁶The property is not exactly the same, it has to be suitably stabilized.

so that the sequence is eventually constant, then we say that the computation **halts**. In this case we denote by s_f the final configuration, so that the sequence is eventually constant with terms s_f . We define the **length** of a computation sequence to be the first occurrence of $n > 0$ s.t. $s_n = s_f$. For a given Turing computation $*M(\Sigma)$, we will write

$$*M(\Sigma) \rightarrow x,$$

if $*M(\Sigma)$ halts and x is the output string.

We write $M(\Sigma)$ for the output string of M , given the input string Σ , if the associated Turing computation $*M(\Sigma)$ halts.

Definition 1.4. Let *Strings* denote the set of all finite strings, including the empty string ϵ , of symbols in some fixed finite alphabet, with at least 2 elements, for example $\{0,1\}$. Given a partial function $f : \text{Strings} \rightarrow \text{Strings}$, that is a function defined on some subset of *Strings* - we say that a Turing machine M **computes** f if $*M(\Sigma) \rightarrow f(\Sigma)$, whenever $f(\Sigma)$ is defined.

So a Turing machine T itself determines a partial function, which is defined on all $\Sigma \in \text{Strings}$ s.t. $*T(\Sigma)$ halts, by $\Sigma \mapsto T(\Sigma)$. The following definition is purely for writing purposes.

Definition 1.5. Given Turing computations (for possibly distinct Turing machines) $*T_1(\Sigma_1)$, $*T_2(\Sigma_2)$ we say that they are **equivalent** if they both halt with the same output string or both do not halt. We write $T_1(\Sigma_1) = T_2(\Sigma_2)$ if $*T_1(\Sigma_1)$, $*T_2(\Sigma_2)$ both halt with the same value.

In practice we will allow our Turing machine T to reject some elements of *Strings* as valid input. We may formalize this by asking that there is a special final machine state q_{reject} , so that $T(\Sigma)$ halts with q_{reject} for

$$\Sigma \notin I \subset \text{Strings},$$

where I is some set of all valid, that is T -**permissible** input strings. We do not ask that for $\Sigma \in I$ $*T(\Sigma)$ halts. If $*T(\Sigma)$ does halt then we will say that Σ is T -**acceptable**. It will be convenient to forget q_{reject} and instead write

$$T : I \rightarrow O,$$

where $I \subset \text{Strings}$ is understood as the subset of all T -permissible strings, or just **input set** and O is the set output strings or **output set**.

We will sometimes use abstract sets to refer to input and output sets. However, these are understood to be subsets of *Strings* under some implicit, *fixed* encoding. Concretely an **encoding** of A is an injective set map $i : A \rightarrow \text{Strings}$. For example if the input set is Strings^2 , we may encode it as a subset of *Strings* as follows. The encoding string of $\Sigma \in \text{Strings}^2$ will be of the type: “this string encodes an element Strings^2 , whose components are Σ_1 and Σ_2 .” In particular the sets of integers \mathbb{N}, \mathbb{Z} , which we use often, will under some encoding correspond to subsets of *Strings*. Indeed this abstracting of sets from their encoding in *Strings* is partly what computer languages do. The fixing of the encoding can be understood as fixing the computer language.

The above will allow us to work with a set \mathcal{T} of Turing machines, with abstract sets of inputs and outputs implicitly encoded as subsets of *Strings* as above. Note \mathcal{T} itself has an induced encoding. Of course, concretely \mathcal{T} is nothing more then the set of Turing machines, with a distinguished final state called q_{reject} .

Definition 1.6. We say that a Turing machine T computes a partial function $f : I \rightarrow J$, if I is contained in the set of permissible inputs of T and $*T(\Sigma) \rightarrow f(\Sigma)$, whenever $f(\Sigma)$ is defined, for $\Sigma \in I$.

Given Turing machines

$$M_1 : I \rightarrow O, M_2 : J \rightarrow P,$$

we may naturally **compose** them to get a Turing machine $M_2 \circ M_1 : C \rightarrow P$, for $C = M_1^{-1}(O \cap J)$, ($O \cap J$ is understood as intersection of subsets of *Strings*). C can be empty in which case this is a Turing machine which rejects all input. Let us not elaborate further.

1.1. Join of Turing machines. Our Turing machine of Definition 1.1 is a multi-tape enhancement of a more basic notion of a Turing machine with a single tape, but we need to iterate this further.

We replace a single tape by tapes T^1, \dots, T^n in parallel, which we denote by $(T^1 \dots T^n)$ and call this n -tape. The head H on the n -tape has components H^i pointing on the corresponding tape T^i . When moving a head we move all of its components separately. A string of symbols on $(T^1 \dots T^n)$ is an n -string, formally just an element $\Sigma \in \text{Strings}^n$, with i 'th component of Σ specifying a string of symbols on T^i . The blank symbol b is the symbol (b^1, \dots, b^n) with b^i blank symbols of T^i .

Given Turing machines M^1, M^2 we can construct what we call a **join** $M^1 \star M^2$, which is roughly a Turing machine where we alternate the operations of M^1, M^2 . In what follows symbols with superscript 1, 2 denote the corresponding objects of M^1 , respectively M^2 , cf. Definition 1.1.

$M^1 \star M^2$ has three 2-tapes:

$$(T_i^1 T_i^2), (T_c^1 T_c^2), (T_o^1 T_o^2),$$

three heads H_i, H_c, H_o which have component heads H_i^j, H_c^j, H_o^j , $j = 1, 2$. It has machine states:

$$Q_{M^1 \star M^2} = Q^1 \times Q^2 \times (\mathbb{Z}_2 = \{0, 1\}),$$

with initial state $(q_0^1, q_0^2, 0)$ and final states:

$$F_{M^1 \star M^2} = F^1 \times Q^2 \times \{1\} \sqcup Q^1 \times F^2 \times \{0\}.$$

Clearly we have a natural splitting

$$\mathcal{C}(M^1 \star M^2) = \mathcal{C}(M^1) \times \mathcal{C}(M^2) \times \mathbb{Z}_2.$$

In terms of this splitting we define the transition function

$$\delta^{M^1 \star M^2} : \mathcal{C}(M^1 \star M^2) \rightarrow \mathcal{C}(M^1 \star M^2),$$

for our Turing machine $M^1 \star M^2$ by:

$$\begin{aligned} \delta^{M^1 \star M^2}(s^1, s^2, 0) &= (\delta^{M^1}(s^1), s^2, 1), \\ \delta^{M^1 \star M^2}(s^1, s^2, 1) &= (s^1, \delta^{M^2}(s^2), 0). \end{aligned}$$

Or, concretely this means the following. Given machine state $q = (q^1, q^2, 0)$ and the symbols

$$(\sigma_i^1 \sigma_i^2), (\sigma_c^1 \sigma_c^2), (\sigma_o^1 \sigma_o^2)$$

to which the heads H_i, H_c, H_o are currently pointing, we first check instructions in I^1 for $q^1, \sigma_i^1, \sigma_c^1, \sigma_o^1$, and given those instructions as step 1 execute:

- (1) Replace symbols σ_c^1, σ_o^1 to which the head components H_c^1, H_o^1 point, or leave them unchanged, while leaving unchanged the symbols to which H_i^1, H_o^1 point.
- (2) Move each head component H_i^1, H_c^1, H_o^1 left, right, or leave it in place, (independently). (The second components of the heads are unchanged.)
- (3) Change the first component of q to another machine state in Q^1 or keep it, based on the instruction in I^1 . Leave the second component of q unchanged. The third component of q is changed to 1.

Then likewise given machine state $q = (q^1, q^2, 1)$, we check instructions in I^2 for $q^2, \sigma_i^2, \sigma_c^2, \sigma_o^2$ and given those instructions as step 2 execute:

- (1) Replace symbols σ_c^2, σ_o^2 to which the head components H_c^2, H_o^2 point, or leave them unchanged, while leaving unchanged the symbols to which H_i^2, H_o^2 point.
- (2) Move each head component H_i^2, H_c^2, H_o^2 left, right, or leave it in place.
- (3) Change the second component of q to another or keep it, based on instruction in I^2 . Leave the first component unchanged, and change the third component of q to 0.

1.1.1. Input. The input for $M^1 \star M^2$ is a 2-string or in other words pair (Σ_1, Σ_2) , with Σ_1 an input string for M^1 , and Σ_2 an input string for M^2 .

1.1.2. *Output.* The output for

$$*M^1 \star M^2(\Sigma_1, \Sigma_2)$$

is defined as follows. If this computation halts then the 2-tape $(T_o^1 T_o^2)$ contains a 2-string, bounded by b symbols, with T_o^1 component Σ_o^1 and T_o^2 component Σ_o^2 . Then the output $M^1 \star M^2(\Sigma_1, \Sigma_2)$ is defined to be Σ_o^1 if the final state is of the form $(q_f, q, 1)$ for q_f final, or Σ_o^2 if the final state is of the form $(q, q_f, 0)$, for q_f likewise final.

1.2. **Universal Turing machines.** It will be convenient to refer to the universal Turing machine

$$U : \mathcal{T} \times \text{Strings} \rightarrow \text{Strings},$$

for \mathcal{T} the set of Turing machines as already indicated above. This universal Turing machine already appears in Turing's [1]. It permits as input a pair (T, Σ) for T an encoding of a Turing machine and Σ input to this T . It can be partially characterized by the property that for every Turing machine T and string Σ we have:

$$*T(\Sigma) \text{ is equivalent to } *U(T, \Sigma).$$

1.3. **Notation.** In what follows \mathbb{Z} is the set of all integers and \mathbb{N} non-negative integers. We will sometimes specify a Turing machine simply by specifying a function

$$T : I \rightarrow O,$$

with the full data of the underlying Turing machine being implicitly specified, in a way that should be clear from context. When we intend to suppress dependence of a variable V on some parameter p we often write $V = V(p)$, this equality is then an equality of notation not of mathematical objects.

2. PRELIMINARY SETUP FOR THE PROOF OF THEOREM 0.1

This section can be understood to be a warm up, as we will not yet work with stable soundness. But most of this will carry on to the more technical setup of Section 3.

Definition 2.1. A **machine**⁷ will be a synonym for a partial function $A : I \rightarrow O$, with I, O abstract sets with a fixed, prescribed encoding as subsets of *Strings*, (cf. *Preliminaries*).

\mathcal{M} will denote the set of machines. Given a Turing machine $T : I \rightarrow O$, we have an associated machine $fog(T)$ by forgetting all structure except the structure of a partial function. \mathcal{T} will denote the set of machines, which in addition have the structure of a Turing machine. So we have a forgetful map $fog : \mathcal{T} \rightarrow \mathcal{M}$.

2.1. **Diagonalization machines.** There is a well known connection between Turing machines and formal systems to which we already alluded in Section 0.1. So Gödel statements can already be interpreted in Turing machine language as certain Gödel strings. But we will be aiming to construct, in a specific setting relevant to our goals, a more flexible and in a certain sense universal (for our class of Turing machines) such Gödel string \mathcal{G} . Extending this construction to more general classes of Turing machines / formal systems would be very interesting, but at the moment it is not clear what that would entail.

To make this \mathcal{G} exceptionally simple we will need to formulate some specific properties for our machines, which will require a bit of setup. We denote by $\mathcal{T}_{\mathbb{Z}} \subset \mathcal{T}$ the subset of Turing machines of the type:

$$X : (S_X \times \mathbb{N} \subset \text{Strings} \times \mathbb{N}) \rightarrow \mathbb{Z}.$$

In other words, the input set of $X \in \mathcal{T}_{\mathbb{Z}}$ is of the form $S_X \times \mathbb{N}$, for $S_X \subset \text{Strings}$, and the output set of X is \mathbb{Z} .

Let $\mathcal{O} \subset \mathcal{T}_{\mathbb{Z}} \times \text{Strings}$ consist of $(X, \Sigma) \in \mathcal{T}_{\mathbb{Z}} \times \text{Strings}$ with $\Sigma \in S_X$, defined as above. And set

$$\mathcal{O}' := \mathcal{O} \times \mathbb{N} \subset \mathcal{T}_{\mathbb{Z}} \times \text{Strings} \times \mathbb{N}.$$

⁷For some authors and in some of the writing of Turing and Gödel “machine” is synonymous with Turing machine. For us the term machine is just abstraction for a process.

Let

$$D_1 : \mathbb{Z} \sqcup \{\infty\} \rightarrow \mathbb{Z},$$

be a fixed Turing machine which satisfies

$$(2.2) \quad D_1(x) = x + 1 \text{ if } x \in \mathbb{Z} \subset \mathbb{Z} \sqcup \{\infty\}$$

$$(2.3) \quad D_1(\infty) = 1.$$

Here $\{\infty\}$ is the one point set containing the element ∞ , which is just a particular distinguished symbol, also implicitly encoded as an element of *Strings*, s.t. $\{\infty\} \cap \mathbb{Z} = \emptyset$, where the intersection is taken in *Strings*. In what follows we sometimes understand D_1 as an element of $\mathcal{T}_{\mathbb{Z}}$, denoting the Turing machine:

$$(2.4) \quad (x, m) \mapsto D_1(x),$$

for all $(x, m) \in (\mathbb{Z} \sqcup \{\infty\}) \times \mathbb{N}$.

We need one more Turing machine.

Definition 2.5. *We say that a Turing machine*

$$R : D \supset \mathcal{O}' \rightarrow \mathbb{Z} \sqcup \{\infty\},$$

has property G if the following is satisfied:

- *R halts on the entire \mathcal{O}' , that is \mathcal{O}' is contained in the set of R-acceptable strings.*
- *$R(X, \Sigma, m) \neq \infty \implies R(X, \Sigma, m) = X(\Sigma, m)$, for $(\Sigma, m) \in S_X \times \mathbb{N}$, and $X \in \mathcal{T}_{\mathbb{Z}}$.*
- *$\forall m : R(D_1, \infty, m) \neq \infty$, and so $\forall m : R(D_1, \infty, m) = 1$, by the previous property.*

Lemma 2.6. *There is a Turing machine R satisfying property G.*

Proof. Let W_n be some Turing machine $W_n : \{\epsilon\} \rightarrow \{\infty\}$, for $\epsilon \in \text{Strings}$ the empty string. So as a function it is not very interesting since the input and output sets are singletons. We ask that the length of $*W_n(\epsilon)$ is $n > 0$, (cf. Preliminaries). Let R_n be the Turing machine, specified as

$$R_n(Z) := W_n \star U(\epsilon, Z),$$

in the language of the join operation described in Section 1, for $Z \in \text{Strings}$, and for U the universal Turing machine. Clearly R_n always halts, although it may halt with machine state q_{reject} . Moreover by construction every $Z = (X, \Sigma, m) \in \mathcal{O}' \subset \text{Strings}$ is permitted. Additionally, for $(X, \Sigma, m) \in \mathcal{O}'$,

$$R_n(X, \Sigma, m) \neq \infty \implies R_n(X, \Sigma, m) = X(\Sigma, m),$$

in particular every $(X, \Sigma, m) \in \mathcal{O}'$ is R_n -acceptable. As a function $\mathbb{Z} \sqcup \{\infty\} \rightarrow \mathbb{Z}$, D_1 is completely determined but it could have various implementations as a Turing machine, so that the length l_m of $*D_1(\infty, m)$ depends on this implementation. Clearly we may assume that $\forall m : l = l_m$ for some l , by definition of D_1 as an element of $\mathcal{T}_{\mathbb{Z}}$, as in (2.4). We then ask that $n_0 > l$ is fixed. Then by construction we get:

$$\forall m : R_{n_0}(D_1, \infty, m) = D_1(\infty, m) = 1.$$

So set $R := R_{n_0}$, and this gives the desired Turing machine. Note that the domain $D \subset \mathcal{T} \times \text{Strings}$ of R -permissible strings is not explicitly determined by our construction, as we cannot tell without additional information when a general Z is rejected by R . We can only say that $D \supset \mathcal{O}'$. \square

Define \mathcal{M}_0 to be the set of machines whose input set is $\mathcal{I} = \mathcal{T} \times \mathbb{N}$ and whose output set is *Strings*. That is

$$\mathcal{M}_0 := \{M \in \mathcal{M} \mid M : \mathcal{T} \times \mathbb{N} \rightarrow \text{Strings}\}.$$

We set

$$\mathcal{T}_0 := \{T \in \mathcal{T} \mid fog(T) \in \mathcal{M}_0\},$$

and we set $\mathcal{I}_0 := \mathcal{T}_0 \times \mathbb{N}$. Given $M \in \mathcal{M}_0$ and $M' \in \mathcal{T}_0$ let $\Theta_{M, M'}$ be the statement:

$$(2.7) \quad M \text{ is computed by } M'.$$

For each $M \in \mathcal{M}_0$, we define a machine:

$$\begin{aligned} \widetilde{M} : \mathcal{I} &\rightarrow \text{Strings} \times \mathbb{N} \\ (2.8) \quad \widetilde{M}(B, m) &= (M(B, m), m), \end{aligned}$$

which is naturally a Turing machine when M is a Turing machine.

In what follows when we write $M'(M', m)$ we mean $M'(\Sigma_{M'}, m)$ for $\Sigma_{M'}$ the string encoding of the specification of the Turing machine M' . So we conflate the notation for the Turing machine and its string specification.

Definition 2.9. For $M \in \mathcal{M}_0$, $M' \in \mathcal{T}_0$, an abstract string $O \in \text{Strings}$ is said to have **property** $C = C(M, M')$ if:

$$\begin{aligned} \Theta_{M, M'} \implies \forall m : (&*M'(M', m) \text{ does not halt}) \vee (M'(M', m) \notin \mathcal{O}) \\ \vee (M'(M', m) \in \mathcal{O}, O \in \mathcal{O} \text{ and } &X(\Sigma, m) = D_1 \circ R \circ \widetilde{M}'(M', m), \text{ where } (X, \Sigma) = O), \end{aligned}$$

and where \widetilde{M}' is determined by M' as in (2.8).

At a glance, this is a somewhat complicated property, but essentially it just says that if $\Theta_{M, M'}$ then for all m “ $O \neq M'(M', m)$ ” unless either $*M'(M', m)$ does not halt, or the output does not have the right (data) type, or $R(O, m) = \infty$. Thus the string O with property $C(M, M')$ is “diagonal” in a certain sense, where by “diagonal” we mean that something analogous to Cantor’s diagonalization is happening, but we will not elaborate.

Remark 2.10. The fact that data types get intricated is perhaps not surprising. On one hand there is a well known correspondence, the Curry-Howard correspondence [6], between proof theory in logic and type theory in computer science, and on the other hand we are doing something at least loosely related to Gödel incompleteness, but in the language of Turing machines.

Definition 2.11. We say that $M \in \mathcal{M}_0$ is **C-sound**, or is a **diagonalization machine**, if for each $(M', m) \in \mathcal{I}_0$, with $M(M', m) = O$ defined, O has property $C(M, M')$. We say that M is C-sound on M' if the list $\{M(M', m)\}_m$ has only elements with property $C(M, M')$.

Define a C-sound $M' \in \mathcal{T}_0$ analogously.

Definition 2.12. If M, M' as above are C-sound we will say that $\text{sound}(M)$, $\text{sound}(M')$ hold. If M is C-sound on M' we say that $\text{sound}(M, M')$ holds.

Example 1. A trivially C-sound machine M is one for which

$$M(M', m) = (D_1 \circ R \circ \widetilde{M}', M')$$

for every $(M', m) \in \mathcal{I}$. As $(D_1 \circ R \circ \widetilde{M}', M')$ automatically has property $C(M, M')$ for each $M' \in \mathcal{T}_0$. In general, for any $M \in \mathcal{M}_0$, $M' \in \mathcal{T}_0$ the list of all strings O with property $C(M, M')$ is always infinite, as by this example there is at least one such string $(D_1 \circ R \circ \widetilde{M}', M')$, which can then be modified to produce infinitely many such strings.

Theorem 2.13. If $\text{sound}(M, M') \wedge \Theta_{M, M'}$ then

$$\forall m : M(M', m) \neq (D_1, \infty).$$

On the other hand, if $\text{sound}(M, M')$ then the string

$$\mathcal{G} := (D_1, \infty) \in \mathcal{O}$$

has property $C(M, M')$. In particular if $\text{sound}(M)$ then \mathcal{G} has property $C(M, M')$ for all M' .

So given any C-sound $M \in \mathcal{M}_0$ there is a certain string \mathcal{G} with property $C(M, M')$ for all M' , such that for each M' if $\Theta_{M, M'}$ then

$$\mathcal{G} \neq M(M', m),$$

for all m . This ‘‘Gödel string’’ \mathcal{G} is what we are going to use further on. What makes \mathcal{G} particularly suitable for our application, is that it is independent of the particulars of M , all that is needed is $\mathcal{M} \in \mathcal{M}_0$ and is C -sound. So \mathcal{G} is in a sense universal.

Proof. Suppose not and let M'_0 be such that $\Theta_{M,M'_0} \wedge \text{sound}(M, M'_0)$ and such that

$$M(M'_0, m_0) = \mathcal{G} \text{ for some } m_0,$$

so that \mathcal{G} has property $C(M, M')$. Set $I = (M'_0, m_0)$ then we have that:

$$1 = D_1(\infty, m_0),$$

$$D_1(\infty, m_0) = D_1 \circ R \circ \widetilde{M}'(I), \text{ by } \mathcal{G} \text{ having property } C(M, M'), \text{ and by } *M'(I) \rightarrow \mathcal{G} \in \mathcal{O} \text{ since } \Theta_{M,M'},$$

$$D_1 \circ R \circ \widetilde{M}'(I) = D_1 \circ R(D_1, \infty, m_0) \quad \text{by } M'(I) = \mathcal{G},$$

$$D_1 \circ R(D_1, \infty, m_0) = 2 \quad \text{by property } G \text{ of } R \text{ and by (2.2),}$$

$$1 = 2.$$

So we obtain a contradiction.

We now verify the second part of the theorem. Given $M' \in \mathcal{T}_0$, we show that:

$$(2.14) \quad \forall m : \left(\text{sound}(M, M') \wedge (M'(I) \in \mathcal{O}) \wedge \Theta_{M,M'} \implies R(\widetilde{M}'(I)) = \infty \right),$$

where $I = (M', m)$. Suppose otherwise that for some m_0 and $I_0 = (M', m_0)$ we have:

$$\text{sound}(M, M') \wedge (*M'(I_0) \text{ halts}) \wedge (M'(I_0) \in \mathcal{O}) \wedge \Theta_{M,M'} \wedge (R(\widetilde{M}'(I_0)) \neq \infty).$$

So we have:

$$(2.15) \quad *M'(I_0) \rightarrow (X, \Sigma) \in \mathcal{O},$$

for some (X, Σ) having property $C(M, M')$. And so, since R is defined on all of \mathcal{O}' :

$$R(\widetilde{M}'(I_0)) = R(X, \Sigma, m_0) = X(\Sigma, m_0) = x \in \mathbb{Z}, \text{ for some } x,$$

by Property G of R and by $R(\widetilde{M}'(I_0)) \neq \infty$.

Then we get:

$$x = X(\Sigma, m_0) = D_1 \circ R \circ \widetilde{M}'(I_0) = D_1(x) = x + 1$$

by (X, Σ) having property $C(M, M')$, and by (2.15). So we get a contradiction and (2.14) follows. \square

2.2. A system with a human subject S as a machine in \mathcal{M}_0 . Let S be a human subject, in a controlled environment, in communication with an experimenter/operator E that as input passes to S elements of $\mathcal{I} = \mathcal{T} \times \mathbb{N}$. Here **controlled environment** means primarily that no information i.e. stimulus, that is not explicitly controlled by E and that is usable by S , passes to S while he is in this environment. This condition is only for simplicity, so long as we know in principle, or can compute in principle, what ‘‘input’’ our S receives, it doesn’t matter what kind of environment he is in. For practical purposes S has in his environment a general purpose digital computer with arbitrarily, as necessary, expendable memory, (in other words a universal Turing machine).

We suppose that upon receiving any $I \in \mathcal{I}$, as a string in his computer, after possibly using his computer in some way, S instructs his computer to print after some indeterminate time a string $S(I)$. We are not actually assuming that $S(I)$ is defined on every I , (although this would likely be a safe assumption). So S in our language also denotes a machine:

$$S : \mathcal{I} \rightarrow \text{Strings},$$

which we suppose satisfies the condition that for any fixed $T \in \mathcal{T}_0$

$$\{S(T, m)\}_m$$

is the complete list of strings that S asserts to have property $C(S, T)$. (While being a part of the system above.) Of course in our argument we will not actually need S to list infinitely many strings.

Definition 2.16. We say that S the human subject is **computable** if the corresponding machine S above is computable.

Additional condition for S . Let S be a subject as above, which additionally satisfies what we call the *Penrose property*. We ask that the physical S asserts that they are sound, which entails in this case that they assert $\text{sound}(S)$ for S the above machine. And we ask that S is aware of Theorem 2.13, so that as a consequence S asserts that \mathcal{G} has property $C(S, T)$, for all $T \in \mathcal{T}_0$. Or still more concretely: for all T , $\mathcal{G} = S(T, m)$ for some m . This condition is preliminary, since asserting soundness is at least on the surface irrational, and we formally treat fundamental soundness only in the next section.

Theorem 2.17.

$$S \text{ is computable} \implies \neg \text{sound}(S).$$

In fact we prove more, for any $S' \in \mathcal{T}_0$:

$$\Theta_{S, S'} \implies \neg \text{sound}(S, S').$$

This partly formalizes Theorem 0.1, to completely formalize it we must wait till the following sections.

Proof. Suppose $\Theta_{S, S'}$ for some $S' \in \mathcal{T}_0$. Suppose in addition $\text{sound}(S, S')$. Then by Theorem 2.13

$$S(S', m) \neq \mathcal{G}$$

for any m , which contradicts the Penrose property above. \square

3. FUNDAMENTAL SOUNDNESS AS STABLE SOUNDNESS

Imagine a machine P which sequentially prints statements of arithmetic, which it asserts are true, but so that P can also delete a printed statement, if P decided the statement to be untrue. We say that P is stably sound if any printed statement by P that survives to infinity is in fact true. More formally, for each $n \in \mathbb{N}$, $P(n)$ will correspond to an operation denoted by the string $(\Sigma, +)$ or $(\Sigma, -)$ meaning add Σ to the list or remove Σ from list, respectively, where Σ is a statement of arithmetic. If there is an n_0 with $P(n_0) = (\Sigma, +)$, s.t. there is no $m > n_0$ with $P(m) = (\Sigma, -)$, then Σ is called *P -stable* and we say that P **prints Σ stably**.

Definition 3.1. We say that P is **stably sound** if every P -stable Σ is true.

Remark 3.2. Given a stably sound P , we may construct from it a sound machine P^s simply by enumerating, in order, all the P -stable Σ . However this is of no usefulness in our context, as in general P^s may not be computable even if P is computable. Explicit examples of this sort can be constructed by hand. We can in fact construct a Turing machine A , whose stabilization A^s enumerates every Diophantine equation with no integer solution, and this set is well known to be undecidable, [1].

We now translate this to our setting. The crucial point of our Gödel string is that it will still function in this stable soundness context. Let \mathcal{M}^\pm denote the set of machines

$$M : \mathcal{I} = \mathcal{T} \times \mathbb{N} \rightarrow \text{Strings} \times \{\pm\},$$

where $\{\pm\}$ is the set containing two symbols $+$, $-$, likewise implicitly encoded as a subset of *Strings*. We set

$$\mathcal{T}^\pm := \{T \in \mathcal{T} \mid \text{fog}(T) \in \mathcal{M}^\pm\}.$$

Definition 3.3. For $M \in \mathcal{M}^\pm$, and for $(T, m) \in \mathcal{I}$, we say that an abstract $O \in \text{Strings}$ is (M, T) -**stable**, and that M **prints O T -stably** if there exists an $m \in \mathbb{N}$ s.t. $M(T, m) = (O, +)$ and there is no $k > m$ s.t. $M(T, k) = (O, -)$. When $T \in \mathcal{T}^\pm$ and $\text{fog}(T) = M$, instead of writing (M, T) -stable we just write T -stable.

Let

$$pr : Strings \times \{\pm\} \rightarrow Strings,$$

be the natural projection. For each $M \in \mathcal{M}^\pm$, we define a machine:

$$\widetilde{M} : \mathcal{I} \rightarrow Strings \times \mathbb{N},$$

$$(3.4) \quad \widetilde{M}(T, m) = (pr \circ M(T, m), m),$$

which is naturally a Turing machine when M is a Turing machine.

In what follows $\mathcal{O} \subset \mathcal{T}_{\mathbb{Z}} \times Strings$ is as before.

Definition 3.5. For $M \in \mathcal{M}^\pm$, $M' \in \mathcal{T}^\pm$, an abstract string $O \in Strings$ is said to have property $sC = sC(M, M')$ if:

$$\begin{aligned} \Theta_{M, M'} \implies & \forall m : (*M'(M', m) \text{ does not halt}) \vee (pr \circ M'(M', m) \notin \mathcal{O}) \vee (pr \circ M'(M', m) \text{ is not } M'\text{-stable}) \\ & \vee (pr \circ M'(M', m) \in \mathcal{O}, O \in \mathcal{O} \text{ and } X(\Sigma, m) = D_1 \circ R \circ \widetilde{M}'(M', m), \text{ where } (X, \Sigma) = O), \end{aligned}$$

for \widetilde{M}' determined by M' as in (3.4).

Definition 3.6. We say that $M \in \mathcal{M}^\pm$ is **stably C-sound** on M' , and we write that $s\text{-sound}(M, M')$ holds, if every (M, M') -stable O has property $sC(M, M')$. We say that M is **stably C-sound** if it is stably C-sound on all M' , and in this case we write that $s\text{-sound}(M)$ holds.

Example 2. As before an example of a trivially stably C-sound machine M is one for which

$$M(M', m) = (D_1 \circ R \circ \widetilde{M}', M', +)$$

for every $(M', m) \in \mathcal{I}$.

Theorem 3.7. If $s\text{-sound}(M, M') \wedge \Theta_{M, M'}$ for some $M' \in \mathcal{T}^\pm$ then

$$(O \text{ is } (M, M')\text{-stable}) \implies O \neq \mathcal{G},$$

where

$$\mathcal{G} := (D_1, \infty) \in \mathcal{O}.$$

On the other hand,

$$(3.8) \quad \forall M' \in \mathcal{T}^\pm : s\text{-sound}(M, M') \implies \mathcal{G} \text{ has property } sC(M, M').$$

In particular if $s\text{-sound}(M)$ then \mathcal{G} has property $sC(M, M')$ for all M' .

Proof. This is mostly analogous to the proof of Theorem 2.13. Suppose not and let M' be such that $s\text{-sound}(M, M') \wedge \Theta_{M, M'}$ and such that for some m_0 :

$$M(M', m_0) = (\mathcal{G}, +) \text{ and } \mathcal{G} \text{ is } (M, M')\text{-stable},$$

with \mathcal{G} consequently having property $sC(M, M')$, by $s\text{-sound}(M, M')$.

If we set $I = (M', m_0)$, then by \mathcal{G} having property $sC(M, M')$, by $*M'(I) \rightarrow (\mathcal{G}, +)$, $\mathcal{G} \in \mathcal{O}$ since $\Theta_{M, M'}$ and by \mathcal{G} being M' -stable as \mathcal{G} is (M, M') -stable:

$$D_1(\infty, m_0) = D_1 \circ R \circ \widetilde{M}'(I).$$

On the other hand:

$$\begin{aligned} D_1 \circ R \circ \widetilde{M}'(I) &= D_1 \circ R(D_1, \infty, m_0) \quad \text{by } M'(I) = (\mathcal{G}, +), \\ D_1 \circ R(D_1, \infty, m_0) &= 2 \quad \text{by property } G \text{ of } R \text{ and by (2.2),} \\ D_1(\infty, m_0) &= 1, \\ 1 &= 2. \end{aligned}$$

So we obtain a contradiction.

We now verify the second part of the theorem. Given any $M' \in \mathcal{T}_0$, for any $m \in \mathbb{N}$, setting $I := (M', m)$ we show that:

$$(3.9) \quad s\text{-sound}(M, M') \wedge (pr \circ M'(I) \in \mathcal{O}) \wedge (pr \circ M'(I) \text{ is } M'\text{-stable}) \wedge \Theta_{M, M'} \implies R(\widetilde{M}'(I)) = \infty.$$

Suppose otherwise that for some m_0 and $I_0 := (M', m_0)$ we have:

$$s\text{-sound}(M, M') \wedge (*M'(I_0) \text{ halts}) \wedge (pr \circ M'(I_0) \in \mathcal{O}) \wedge (pr \circ M'(I_0) \text{ is } M'\text{-stable}) \wedge \Theta_{M, M'} \wedge (R(\widetilde{M}'(I_0)) \neq \infty).$$

Then by the above condition we get:

$$(3.10) \quad *M'(I_0) \rightarrow (O, +), \text{ or } *M'(I_0) \rightarrow (O, -),$$

for some $O = (X, \Sigma) \in \mathcal{O}$, which is (M, M') -stable, and with property $sC(M, M')$. We can of course guarantee that there is some m'_0 with $M'(M', m'_0) = (O, +)$, but we arranged the details so that this is not necessary.

Since R is defined on all of \mathcal{O}' we get:

$$R(\widetilde{M}'(I_0)) = R(O, m_0) = X(\Sigma, m_0) = x \in \mathbb{Z}, \text{ for some } x,$$

by Property G of R and by $R(\widetilde{M}'(I_0)) \neq \infty$. Then we have:

$$x = X(\Sigma, m_0) = D_1 \circ R \circ \widetilde{M}'(I_0) = D_1(x) = x + 1,$$

by (X, Σ) having property $sC(M, M')$, and by (3.10). So we get a contradiction and (3.9) follows. Our conclusion readily follows. \square

4. A SYSTEM WITH A HUMAN SUBJECT S AS A MACHINE IN \mathcal{M}^\pm

Let S be a human subject in a controlled environment as before. For the moment S is not in any way idealized. We may then suppose, as in Section 2.2, that S determines an element of \mathcal{M}^\pm :

$$S : \mathcal{I} \rightarrow \text{Strings} \times \{\pm\}.$$

Once again S is only a partial function. We suppose S satisfies the condition that for any fixed $T \in \mathcal{T}^\pm$ the list $\{S(T, m)\}_m$ is complete, in the sense that if the physical S , while part of the environment above, eventually stably asserts that some O has property $sC(S, T)$ then

$$S(T, m) = (O, +)$$

for some m , and that moreover O is (S, T) -stable.

And conversely if O is (S, T) -stable, then S stably asserts that O has property $sC(S, T)$. Here, “stably asserts”, analogously to previous usage means, means that S is never to change their mind on this.

As S now denotes two things: the human subject and the corresponding machine, we will say **physical** S when we want to clarify that we are talking of the actual human.

Definition 4.1. *As before, we say that the physical S is **computable** if the corresponding machine S above is computable.*

Definition 4.2. *We say that $M \in \mathcal{M}^\pm$ has the **stable Penrose property** if \mathcal{G} is (M, T) -stable for every $T \in \mathcal{T}^\pm$. In particular if the physical S stably asserts $s\text{-sound}(S)$ and is aware of the second half of Theorem 3.7 then the associated $S \in \mathcal{M}^\pm$ has the stable Penrose property, by the condition above.*

We suppose from now on that S has the stable Penrose property. We further justify this property below.

The following theorem is an immediate consequence of Theorem 3.7.

Theorem 4.3. *Let $M \in \mathcal{M}^\pm$ have the stable Penrose property then*

$$M \text{ is computable} \implies \neg s - \text{sound}(M).$$

In fact, for any $M' \in \mathcal{T}^\pm$:

$$\Theta_{M,M'} \implies \neg s - \text{sound}(M, M').$$

In particular if our physical S is computable, they cannot be fundamentally sound, specifically meaning stably sound.

4.1. Justifying the stable Penrose property. It remains to explain however how the stable Penrose property is justified. The interpretation of this for a given S is that such a human is simply asserting that the list of certain mathematical statements, that they assert to be true, converges in the exact sense above to a list of things which are actually true. For example I assert in absolute faith L : 5 is an odd number. This statement L is likely stably on my list, unless I would have lost my sanity and hence would no longer be me. On the other hand, the statement $0 = 1$ should not be stably on my list.

To make further sense of this, we could idealize our S so that their brain is not subject to deterioration, and so that S aware of this. But then we have to understand the implications of our theorem for idealized S , for non-idealized S . In principle this shouldn't be a problem, since we can imagine a process of idealization that does not magically introduce non-computability in the limit, where the was none before. But this discussion is outside our scope.

Or we may say more concretely that the idealized human is represented by “the evolving scientific community” H , as we have already mentioned in the introduction. The fact that it is “evolving” because its members change (death/birth) presents no problems. If each individual human is Turing computable, then so is this H itself understood as machine:

$$H : \mathcal{I} \rightarrow \text{Strings} \times \{\pm\},$$

where the output $H(I)$ is determined by, for example, majority consensus. H believes and must believe if science is to exist that scientific or at least mathematical progress converges on truth. So it is natural to postulate that H would (eventually) stably assert $s - \text{sound}(H)$, or more directly that H has the stable Penrose property. Then we may apply Theorem 4.3 above to this H , which then formalizes Theorem 0.1.

Proof of Theorem 4.3. Suppose $\Theta_{M,M'}$ for some $M' \in \mathcal{T}^\pm$. Suppose in addition $s - \text{sound}(M, M')$. Then by Theorem 3.7 for all m s.t. $pr \circ M(M', m)$ is (M, M') -stable:

$$M(M', m) \neq \mathcal{G},$$

but this contradicts the stable Penrose property. \square

5. FORMAL SYSTEM INTERPRETATION

Let us fix our interpretation of an idealized human as H , from before, postulated to have the stable Penrose property. As before we will say physical H when we want to emphasize that we are talking of the physical thing representing the partial function H . Theorem 4.3, allows us to conclude that if H is computable then H not stably sound. The one string \mathcal{G} that H is guaranteed to stably print, which expresses unsoundness of H , while elementary is slightly esoteric. Can we see more clearly that H is unsound? Yes, but we need stronger assumptions, and some language of formal systems. This section can be safely omitted as it is only of secondary interest.

For simplicity we will base everything of standard set theory \mathcal{ST} (Zermelo-Fraenkel axioms). Turing machines, and arithmetic are assumed to be naturally formalized in \mathcal{ST} . In what follows, for a statement L , $\mathcal{F} \vdash L$ means that L is provable in the formal system \mathcal{F} .

Let \mathcal{A} denote the set of sentences of arithmetic, as formalized by \mathcal{ST} . Let

$$P : \mathbb{N} \rightarrow \mathcal{A} \times \{+, -\},$$

be a machine associated to scientific community H , analogously to the previous discussion, and as in the preamble to Section 3.

Definition 5.1. We will say that the physical H is **captured by a formal system** $\mathcal{F} \supset \mathcal{ST}$ if the following are satisfied:

(1) For any $T \in \mathcal{T}^\pm$:

$$(O \text{ is } (H, T)\text{-stable}) \iff \mathcal{F} \vdash (O \text{ has property } sC(H, T)).$$

(2)

$$(A \in \mathcal{A} \text{ is } P\text{-stable}) \iff \mathcal{F} \vdash A.$$

We will suppose from now on that the physical H accepts \mathcal{ST} as sound, meaning more specifically that H stably asserts its theorems as true, where the meaning of this is as in previous sense. In this case a formal system \mathcal{F} in the language of set theory, that captures H , always exists. We may simply define $\mathcal{F}(H)$ capturing H by adding to \mathcal{ST} as axioms the P -stable statements, and add as an axiom, for each O (H, T) -stable, the statement:

$$O \text{ has property } sC(H, T).$$

The resulting formal system $\mathcal{F}(H)$ is in fact minimal, in the sense that any other \mathcal{F} which captures H contains $\mathcal{F}(H)$.⁸ Let $Con(H)$ denote the statement:

$$\exists \mathcal{F} : (\mathcal{F} \supset \mathcal{ST} \text{ s.t. } \mathcal{F} \text{ captures } H) \wedge (\mathcal{F} \text{ is consistent}).$$

This is of course equivalent to the condition that $\mathcal{F}(H)$ is consistent by the above discussion.

Theorem 5.2. Let H be as above then:

$$\mathcal{ST} \vdash (\exists H' \in \mathcal{T}^\pm : \Theta_{H, H'}) \implies \neg Con(H).$$

In other words if H is provably computable then they are not consistent.

Note that the statement K :

$$\mathcal{ST} \vdash (\exists H' \in \mathcal{T}^\pm : \Theta_{H, H'})$$

does not mean that the physical H can prove the statement M :

$$(5.3) \quad \exists H' \in \mathcal{T}^\pm : \Theta_{H, H'}$$

in the practical sense. It just means that after the term H in the statement $\Theta_{H, H'}$ has been suitably interpreted in set theory \mathcal{ST} , M is provable in \mathcal{ST} . But a set theoretic, in other words mathematical, interpretation of the term H may not be practically attainable. At the least this necessitates detailed knowledge of the physics and biology underlying our humans. And even if this interpretation was attainable, H may not be clever enough to find the proof of M , again in the practical sense. Also note that $\neg Con(H)$ expresses *fundamental* inconsistency of H , as we only take stable assertions of H above.

Example 3. Suppose that H is provably computable as in the hypothesis of the theorem above. Then $\mathcal{F}(H)$ captures H and $\neg Con(\mathcal{F}(H))$ by the theorem, and so $\mathcal{F}(H)$ proves $0 = 1$. Since $\mathcal{F}(H)$ captures H the physical H must *stably* assert $0 = 1$.

Proof of Theorem 5.2. Let \mathcal{F} capture H as above. By the proof of the second part of Theorem 4.3:

$$\mathcal{ST} \vdash ((\exists H' \in \mathcal{T}^\pm : \Theta_{H, H'}) \implies L),$$

where L is:

$$\begin{aligned} \exists m \in \mathbb{N}, \exists H' \in \mathcal{T}^\pm : (pr \circ H(H', m) \text{ is defined and is } (H, H')\text{-stable}) \wedge \\ \neg(pr \circ H(H', m) \text{ has property } sC(H, H')). \end{aligned}$$

So:

$$(\mathcal{ST} \vdash M) \implies (\mathcal{ST} \vdash L) \implies (\mathcal{F} \vdash L),$$

where M is as in (5.3).

⁸ $\mathcal{F}(H)$ may not however be effectively axiomatized, cf. Remark 3.2, but this is not relevant to our present discussion.

On the other hand, by the assumption that H is captured by \mathcal{F} :

$$\mathcal{F} \vdash (\forall m \in \mathbb{N}, \forall H' \in \mathcal{T}^\pm : (pr \circ H(H', m) \text{ is defined and is } (H, H')\text{-stable}) \implies (pr \circ H(H', m) \text{ has property } sC(H, H'))).$$

So $\mathcal{F} \vdash \neg L$. □

6. RELATIONSHIP WITH THE PENROSE ARGUMENT AND ITS PREVIOUS CRITICISM

The most lucid criticism of the Penrose argument known to me appears in Koellner [15], [16]. Although our argument is very different it is worthwhile to consider it in light of that existing criticism. Brushing aside for the moment concerns such as soundness of human reasoning, the reason the argument of Penrose is susceptible criticism, is that it is only a meta-argument, meaning principally that it partly relies on interpretation of truth and knowledge. In other words, it is not written as a proof in some formal system, even implied one. And there may be logical subtleties to interpretation of truth that a non-logician such as myself or Penrose may simply miss.

So what of the main argument of this paper? Here the situation is rather different. With the possible exception of the inessential Section 5, where we touch on some logic, our theorems are in fact theorems of set theory, just classical Zermelo-Fraenkel set theory. For example our Gödel string \mathcal{G} for a stably C -sound M is proved in set theory to have property sC . In contrast, in [20] Penrose asserts the truth of a certain analogously used Gödel statement by a meta-proof. And this is exactly what is contested by Koellner.

This is not to say that there are no issues of interpretation in this paper. One must interpret our definition of stable soundness, and more to the point the stable Penrose property as it applies to actual human beings. We have of course already partly addressed this. For the moment my position is that the idealized human being of this paper is represented by the world scientific community H (considered as evolving). As already discussed, it makes perfect sense for H to have the stable Penrose property, so under the postulate that H really does have it, our theorem can be interpreted to directly apply to this H .

It is of course always the case that we must interpret mathematical theorems when applied to the real world. What one looks for is whether there is any meaningful physical obstruction to carrying out the necessary idealization in principle. In our specific case I see no such obstruction. Of course if the universe and humanity must eventually go extinct then our idealized humans cannot even in principle exist. But to me this is not a meaningful obstruction. The potential mortality of the universe is very unlikely to have any causal relation with computability of intelligence. So we can imagine an eternal universe, an immortal H and run the argument, then translate to our universe.

7. CONCLUDING REMARK

While it can be argued that humans are not sound, it would be very difficult to argue that we are not stably sound. Scientists operate on the unshakeable faith that scientific progress converges on truth. And our interpretation above of this convergence as stable soundness is very simple and natural. Thus our results put a very serious obstruction to computability of intelligence.

In addition, at least under the stronger hypothesis of Theorem 5.2, stable unsoundness is testable/observable, at least in principle. For if H is provably computable as in Example 3 then H must eventually *stably* assert $0 = 1$. If we knew in addition that a given H' computes H then as H' is a Turing machine we can simulate it on a powerful computer and see if such non-sense statements really do appear. Given our basic understanding of humanity, such a possibility seems too ridiculous.

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