

Chapter 17: Basics of Agda

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November 12, 2025



What is Agda?

- A **dependently typed** programming language
 - Implemented in Haskell
- A **proof assistant**: propositions-as-types
 - Types: propositions
 - Terms (functional programs): proofs

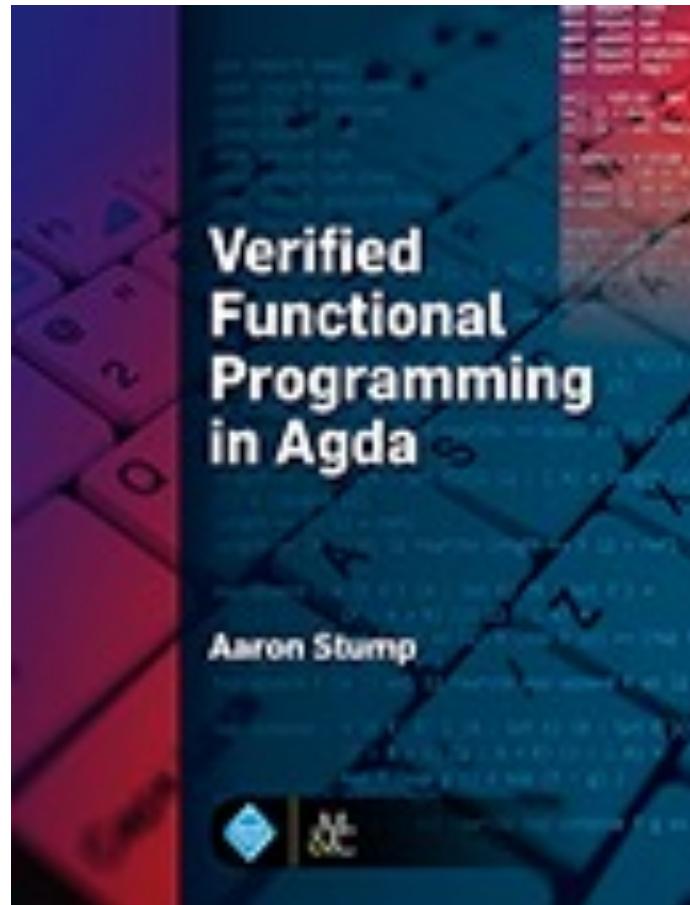


Installation

- There are several ways to install Agda:
 - Using a released **source package** from Hackage
 - Using a **binary package** prepared for your platform
 - Using the development version from the **Git repository**
- More information is available at
<https://agda.readthedocs.io/en/v2.6.2.2/getting-started/installation.html>



Reference



<https://dl.acm.org/doi/book/10.1145/2841316>

Agda source: <http://svn.divms.uiowa.edu/repos/clc/projects/agda/ial-releases/1.2>



17.1 Functional Programming with the Booleans



Declaring the Datatype of Booleans

bool.agda

```
module bool where

-- datatypes

data B : Set where
  tt : B
  ff : B
```

基本命令

| | |
|-------------|---------|
| Load: | C-c C-l |
| Type Check: | C-c C-d |
| Compute: | C-c C-n |



Defining Boolean Operations

```
-- not
~_ : B → B
~ tt = ff
~ ff = tt

-- and
_&&_ : B → B → B
tt && b = b
ff && b = ff

-- or
_||_ : B → B → B
tt || b = tt
ff || b = b
```

注：类型不可省略。



Defining Boolean Operations

```
if_then_else_ : ∀ {ℓ} {A : Set ℓ}
              → ℒ → A → A → A
if tt then y else z = y
if ff then y else z = z

_xor_ : ℒ → ℒ → ℒ
tt xor ff = tt
ff xor tt = tt
tt xor tt = ff
ff xor ff = ff

_nor_ : ℒ → ℒ → ℒ
x nor y = ~ (x || y)
```



17.2 Constructive Proof



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Curry-Howard Isomorphism

- Formulas (Properties) as Types

```
~~ tt ≡ tt
```

- Proof as Programs

```
~~tt : ~~ tt ≡ tt
~~tt = refl
```

```
~~ff : ~~ ff ≡ ff
~~ff = refl
```

definitionally equal.

```
data _≡_ {ℓ} {A : Set ℓ} (x : A) : A → Set ℓ where
  refl : x ≡ x
```



Proving Theorems using Pattern Matching

```
~~-elim : ∀ (b : ℂ) → ~ ~ b ≡ b
~~-elim tt = ~tt
~~-elim ff = ~ff
```

```
&&-idem : ∀ (b : ℂ) → b && b ≡ b
&&-idem tt = refl
&&-idem ff = refl
```



Implicit Arguments

```
&&-idem : ∀ (b : ℂ) → b && b ≡ b  
&&-idem tt = refl  
&&-idem ff = refl
```



自动推导 ℂ

```
&&-idem : ∀ (b) → b && b ≡ b  
&&-idem tt = refl  
&&-idem ff = refl
```



implicit argument b

```
&&-idem : ∀ {b} → b && b ≡ b  
&&-idem{tt} = refl  
&&-idem{ff} = refl
```



Implicit Arguments

```
&&-idem-tt : tt && tt ≡ tt  
&&-idem-tt = &&-idem
```

=

```
&&-idem-tt : tt && tt ≡ tt  
&&-idem-tt = &&-idem{tt}
```



Theorems with Hypotheses

```
| |≡ff2 : ∀ {b1 b2} → b1 || b2 ≡ ff → b2 ≡ ff
| |≡ff2 {tt} ()
| |≡ff2 {ff}{tt} ()
| |≡ff2 {ff}{ff} p = refl
```

absurd pattern (): the case being considered is impossible.

In-Class Exercise: Prove the following.

```
| |≡ff1 : ∀ {b1 b2} → b1 || b2 ≡ ff → ff ≡ b1
```



Matching on Equality Proofs

```
||-cong1 : ∀ {b1 b1' b2} →  
          b1 ≡ b1' → b1 || b2 ≡ b1' || b2  
||-cong1 refl = refl
```

引入等式

定义性相等

=

```
||-cong1 : ∀ {b1 b1' b2} →  
          b1 ≡ b1' → b1 || b2 ≡ b1' || b2  
||-cong1{b1}{.b1'}{b2} refl = refl
```

dot pattern



The rewrite Directive

```
||-cong2 : ∀ {b1 b2 b2'} →  
          b2 ≡ b2' → b1 || b2 ≡ b1 || b2'  
||-cong2 p rewrite p = refl
```

rewrite p: allow a more complicated equation proved by p



Other Examples

Polymorphic theorem

```
ite-same : ∀{ℓ}{A : Set ℓ} →  
          ∀(b : B) (x : A) →  
          (if b then x else x) ≡ x  
ite-same tt x = refl  
ite-same ff x = refl
```

```
B-contra : ff ≡ tt → ∀{ℓ} {P : Set ℓ} → P  
B-contra ()
```



Homework

17.1. Define a datatype day, which is similar to the B datatype but has one constructor for each day of the week.

17.2. Using the day datatype from the previous problem, define a function nextday of type day \rightarrow day, which given a day of the week will return the next day of the week.

17.3. Give a proof of the following formula:

```
ite-arg : ∀{ℓ ℓ'}{A : Set ℓ}{B : Set ℓ'} →  
          (f : A → B)(b : B)(x y : A) →  
          (f (if b then x else y)) ≡ (if b then f x else f y)
```

