

or

LinesRemaining
Δ Cursor lines!: \mathbb{N}
lines! = numLines - line line' = line column' = column

2. *UpKey*. This schema deals with what happens when the cursor is not on the top line of the display:

UpKeyNormal
Δ Cursor key?: KEY
key? = up line > 1 line' = line - 1 column' = column

The next schema deals with what happens when the cursor is on the top line of the display:

UpKeyAtTop
Δ Cursor key?: KEY
key? = up line = 1 line' = line - 1 line' = numLines

Note that the cursor has been defined to *wrap round* to the bottom line of the display. The full behaviour is given by:

$$\text{UpKey} == \text{UpKeyNormal} \vee \text{UpKeyAtTop}$$

3. *LeftKey*. The operation for moving left is given. It is easiest to deal first with what happens when the cursor is not at the far left of the display:

LeftKeyNormal
Δ Cursor key?: KEY
key? = left column > 1 column' = column - 1 line' = line

The next schema deals with the cursor's being at the left of a line other than the top line of the display. Note that the cursor wraps round to the start of the previous line:

LeftKeyAtStart
Δ Cursor key?: KEY
key? = left column = 1 column' = numColumns line > 1 line' = line - 1

Finally, a separate schema deals with the cursor being at the left of the top line. The cursor wraps round to the right of the bottom line:

LeftKeyAtTop
Δ Cursor key?: KEY
key? = left column = 1 column' = numColumns line = 1 line' = numLines

These schemas can be combined to form one schema which defines the response of the cursor to a left-move key in all initial positions of the cursor.

$\text{LeftKey} == \text{LeftKeyNormal} \vee \text{LeftKeyAtStart} \vee \text{LeftKeyAtTop}$

4. *NewDownKey*

NewDownKeyAtBottom _____	
ΔCursor	
key?:	KEY
key? = down	
line = numLines	
line' = numLines	
column' = column	

$\text{NewDownKey} \equiv \text{DownKeyNormal} \vee \text{NewDownKeyAtBottom}$

5. *NewRightKey*

NewRightKeyAtRight _____	
ΔCursor	
key?:	KEY
key? = right	
column = numColumns	
column' = numColumns	
line' = line	

$\text{NewRightKey} == \text{RightKeyNormal} \vee \text{NewRightKeyAtRight}$

6.

- (a) Yes, the Prime Minister must be a Member of Parliament because he or she is a member of the Cabinet and the Cabinet is a subset of the Members of Parliament.
- (b)

HP _____	
MPs:	PPERSON
Cabinet:	PPERSON
DPM, PM:	PERSON
Cabinet \subseteq MPs	
PM \in Cabinet	
DPM \in Cabinet	
DPM \neq PM	

- (c) The new Prime Minister may not be the same person as the old Prime Minister.
- (d) The new Prime Minister does not have to be chosen from the Cabinet.
- (e) The outgoing Prime Minister does not have to leave the Cabinet.
- (f) The outgoing Prime Minister may not leave the Cabinet.

7.

- (a) The members of the new Cabinet must all be MPs, to maintain the invariant that the Cabinet is a subset of the MPs.
- (b) *ChangeCabinet2* requires a complete change of personnel in the new Cabinet.
- (c) The error is that the PM is unchanged and is always a member of the Cabinet, so the Cabinet cannot change completely.

Chapter 7

1.

[PERSON] the set of all uniquely identifiable persons

Computer _____	
users, loggedIn:	PPERSON
loggedIn \subseteq users	

InitComputer _____	
Computer'	
loggedIn' = \emptyset	
users' = \emptyset	

RESPONSE ::=

OK | AlreadyAUser | NotAUser | LoggedIn | NotLoggedIn

2. Add user

AddUser ₀ _____	
$\Delta\text{Computer}$	
p?:	PERSON
p? \notin users	
users' = users \cup {p?}	
loggedIn' = loggedIn	

AddUserError
$\exists \text{Computer}$ $p?: \text{PERSON}$ $\text{reply!}: \text{RESPONSE}$
$p? \in \text{users}$ $\text{reply!} = \text{AlreadyAUser}$

AddUser ==
 $(\text{AddUser}_0 \wedge [\text{reply!}: \text{RESPONSE} \mid \text{reply!} = \text{OK}]) \vee$
 AddUserError

3.

RemoveUser ₀
$\Delta \text{Computer}$ $p?: \text{PERSON}$
$p? \in \text{users}$ $p? \notin \text{loggedIn}$ $\text{users}' = \text{users} \setminus \{p?\}$ $\text{loggedIn}' = \text{loggedIn}$

RemoveUserError
$\exists \text{Computer}$ $p?: \text{PERSON}$ $\text{reply!}: \text{RESPONSE}$
$(p? \notin \text{users} \wedge \text{reply!} = \text{NotAUser})$ \vee $(p? \in \text{users} \wedge p? \in \text{loggedIn} \wedge \text{reply!} = \text{LoggedIn})$

RemoveUser ==
 $(\text{RemoveUser}_0 \wedge [\text{reply!}: \text{RESPONSE} \mid \text{reply!} = \text{OK}]) \vee$
 RemoveUserError

4. Log in

Login ₀
$\Delta \text{Computer}$ $p?: \text{PERSON}$
$p? \in \text{users}$ $p? \notin \text{loggedIn}$ $\text{loggedIn}' = \text{loggedIn} \cup \{p?\}$ $\text{users}' = \text{users}$

LoginError
$\exists \text{Computer}$ $p?: \text{PERSON}$ $\text{reply!}: \text{RESPONSE}$
$(p? \notin \text{users} \wedge \text{reply!} = \text{NotAUser})$ \vee $(p? \in \text{users} \wedge p? \in \text{loggedIn} \wedge \text{reply!} = \text{LoggedIn})$

Login ==
 $(\text{Login}_0 \wedge [\text{reply!}: \text{RESPONSE} \mid \text{reply!} = \text{OK}]) \vee \text{LoginError}$

5. Log out

Logout ₀
$\Delta \text{Computer}$ $p?: \text{PERSON}$
$p? \in \text{users}$ $p? \in \text{loggedIn}$ $\text{loggedIn}' = \text{loggedIn} \setminus \{p?\}$ $\text{users}' = \text{users}$

LogoutError

\exists Computer

$p?:$ PERSON

$reply!:$ RESPONSE

$(p? \notin \text{users} \wedge$
 $reply! = \text{NotAUser})$

\vee

$(p? \in \text{users} \wedge$
 $p? \notin \text{loggedIn} \wedge$
 $reply = \text{NotLoggedIn})$

Logout ==

$(\text{Logout}_0 \wedge [\text{reply!}: \text{RESPONSE} \mid \text{reply!} = \text{OK}]) \vee$
LogoutError

Chapter 8

1. $\text{loggedIn} \subseteq \text{users}$
2. $\forall i: \mathbb{Z} \cdot i * i \geq 0$
3. $\exists n: \mathbb{Z} \cdot n * n = n$
4. $\{n: \mathbb{N} \mid (\forall m: \mathbb{N} \mid m \neq 1 \wedge m \neq n \cdot n \bmod m \neq 0) \cdot n\}$

Chapter 9

1. Latin: LANGUAGE
Latin \notin ran speaks
2. $\# \text{ speaks } \{\text{Switzerland}\} = 4$
3. EU: PCOUNTRY
speaksInEU: COUNTRY \leftrightarrow LANGUAGE
speaksInEU = EU \triangleleft speaks
4. grandParent: PERSON \leftrightarrow PERSON
grandParent = parent ; parent

5.

firstCousin: PERSON \leftrightarrow PERSON

firstCousin = (grandParent ; grandParent[~]) \ sibling

6. Students are either from EU or overseas, but not both. Students study and teachers teach. Only offered modules can be studied. Those modules that are taught are studied.

7.

studies $\{\{p\}\}$

8.

$\#(\text{teaches}\{\{p\}\})$

9. Inverse of *studies* relates modules to persons studying them.

10. The composition relates students to the teachers who teach modules the students study.

11.

$(\text{studies} ; \text{teaches}^{\sim})\{\{p\}\}$

12.

$\#((\text{studies} ; \text{teaches}^{\sim})\{\{p\}\} \cap (\text{studies} ; \text{teaches}^{\sim})\{\{q\}\})$

13.

inter \triangleleft studies

14.

$((\text{teaches} ; \text{studies}^{\sim})\{\{p\}\} \cap (\text{teaches} ; \text{studies}^{\sim})\{\{q\}\}) \triangleright \text{inter} \neq \emptyset$

15.

(a)

delegates \subseteq dom speaks

(b)

ran speaks \cap official $\neq \emptyset$

(c)

$\exists \text{ lang: LANGUAGE} \cdot$

$(\forall \text{ del: PERSON} \mid \text{del} \in \text{delegates} \cdot \text{del speaks lang})$

(d)

$\exists \text{ del: PERSON} \cdot \text{del} \in \text{delegates} \cdot$

$(\exists \text{ lang: LANGUAGE} \cdot \text{del speaks lang} \wedge$

$(\forall \text{ otherDel: PERSON} \mid \text{otherDel} \in \text{delegates} \setminus \{\text{del}\} \cdot$
 $\neg(\text{otherDel speaks lang}))$