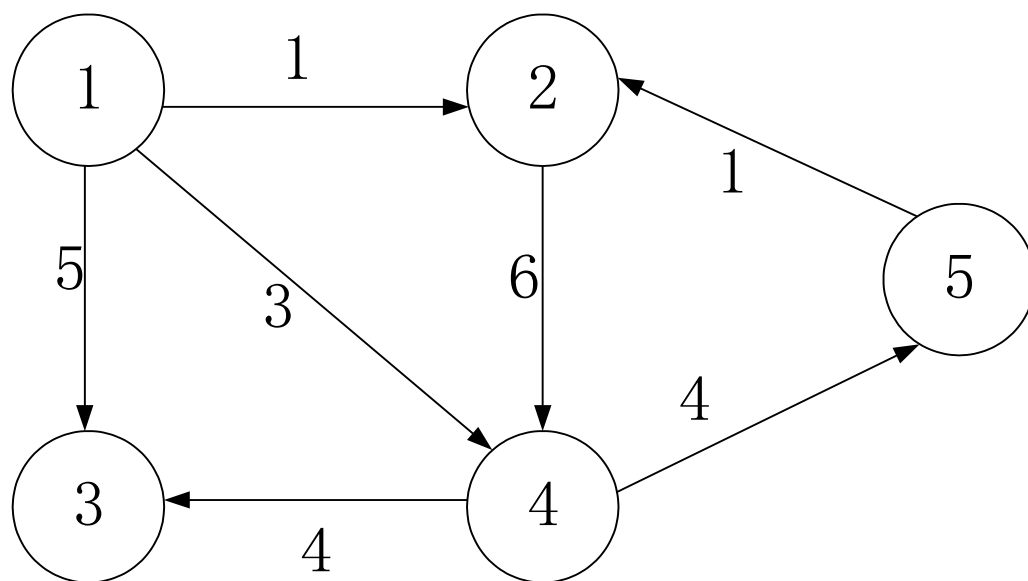


第十五章动态规划：最短路径一题（有向图见下面题）

第十六章回溯法：背包问题 $n=4, w=[2,7,3,5], p=[1,2,4,6], c=11$ ，用带限界函数的回溯法求解。

1 矩阵（回溯）

动态规划



解：先写出有向图对应的矩阵 A

$$A = \begin{pmatrix} \infty & 1 & 5 & 3 & \infty \\ \infty & \infty & \infty & 6 & \infty \\ \infty & \infty & \infty & \infty & \infty \\ \infty & \infty & 4 & \infty & 4 \\ \infty & 1 & \infty & \infty & \infty \end{pmatrix}$$

$$C^{(K)} = (C(i, j, k))$$

$$C(i, j, k) = \min \{ C(i, j, k-1), C(i, k, k-1) + C(k, j, k-1) \} (i \neq j)$$

其中 $i \in \{i | C(i, k, k-1) \neq \infty, i \neq k\}$,

(也就是在矩阵的第 k 列不为 ∞ 的那些元素对应的行号)

$$j \in \{j | C(k, j, k-1) \neq \infty, j \neq k\}$$

(也就是在矩阵的第 k 行不为 ∞ 的那些元素对应的列号)

$$C^{(0)} = \begin{pmatrix} 0 & 1 & 5 & 3 & \infty \\ \infty & 0 & \infty & 6 & \infty \\ \infty & \infty & 0 & \infty & \infty \\ \infty & \infty & 4 & 0 & 4 \\ \infty & 1 & \infty & \infty & 0 \end{pmatrix}$$

k=1 时: $i \in \phi, j \in \{2,3,4\}$

$$C^{(1)} = \begin{pmatrix} 0 & 1 & 5 & 3 & \infty \\ \infty & 0 & \infty & 6 & \infty \\ \infty & \infty & 0 & \infty & \infty \\ \infty & \infty & 4 & 0 & 4 \\ \infty & 1 & \infty & \infty & 0 \end{pmatrix}$$

k=2 时: $i \in \{1,5\}, j \in \{4\}$

$$C(1,4,2) = \min\{C(1,4,1), C(1,2,1)+C(2,4,1)\} = \{3, 1+6\} = 3$$

$$C(5,4,2) = \min\{C(5,4,1), C(5,2,1)+C(2,4,1)\} = \min\{\infty, 1+6\} = 7$$

$$C^{(2)} = \begin{pmatrix} 0 & 1 & 5 & 3 & \infty \\ \infty & 0 & \infty & 6 & \infty \\ \infty & \infty & 0 & \infty & \infty \\ \infty & \infty & 4 & 0 & 4 \\ \infty & 1 & \infty & 7 & 0 \end{pmatrix}$$

k=3 时: $i \in \{1,4\}, j \in \phi$

$$C^{(3)} = \begin{pmatrix} 0 & 1 & 5 & 3 & \infty \\ \infty & 0 & \infty & 6 & \infty \\ \infty & \infty & 0 & \infty & \infty \\ \infty & \infty & 4 & 0 & 4 \\ \infty & 1 & \infty & 7 & 0 \end{pmatrix}$$

k=4 时: $i \in \{1,2,5\}, j \in \{3,5\}$

$$C(1,3,4) = \min\{C(1,3,4), C(1,4,3)+C(4,3,3)\} = \{5, 3+4\} = 5$$

$$C(1,5,4) = \min\{C(1,5,4), C(1,4,3)+C(4,5,3)\} = \{\infty, 3+4\} = 7$$

$$C(2,3,4) = \min\{C(2,3,3), C(2,4,3)+C(4,3,3)\} = \{\infty, 6+4\} = 10$$

$$C(2,5,4) = \min\{C(2,5,3), C(2,4,3)+C(4,5,3)\} = \{\infty, 6+4\} = 10$$

$$C(5,3,4) = \min\{C(5,3,3), C(5,4,3)+C(4,3,3)\} = \{\infty, 7+4\} = 11$$

$$C^{(4)} = \begin{pmatrix} 0 & 1 & 5 & 3 & 7 \\ \infty & 0 & 10 & 6 & 10 \\ \infty & \infty & 0 & \infty & \infty \\ \infty & \infty & 4 & 0 & 4 \\ \infty & 1 & 11 & 7 & 0 \end{pmatrix}$$

k=5 时: $i \in \{1,2,4\}, j \in \{2,3,4\}$

$$C(1,2,5) = \min\{C(1,2,4), C(1,5,4)+C(5,2,4)\} = \{5, 7+1\} = 5$$

$$C(1,3,5) = \min\{C(1,3,4), C(1,5,4)+C(5,3,4)\} = \{5, 7+11\} = 5$$

$$C(1,4,5) = \min\{C(1,4,4), C(1,5,4)+C(5,4,4)\} = \{3, 7+7\} = 3$$

$$C(2,3,5) = \min\{C(2,3,4), C(2,5,4)+C(5,3,4)\} = \{10, 1+11\} = 10$$

$$C(2,4,5) = \min\{C(2,4,4), C(2,5,4)+C(5,4,4)\} = \{6, 1+7\} = 6$$

$$C(4,2,5) = \min\{C(4,2,4), C(4,5,4)+C(5,2,4)\} = \{\infty, 4+1\} = 5$$

$$C(4,3,5) = \min\{C(4,3,4), C(4,5,4)+C(5,3,4)\} = \{4, 4+11\} = 4$$

$$C^{(5)} = \begin{pmatrix} 0 & 1 & 5 & 3 & \infty \\ \infty & 0 & 10 & 6 & 10 \\ \infty & \infty & 0 & \infty & \infty \\ \infty & 5 & 4 & 0 & 4 \\ \infty & 1 & 11 & 7 & 0 \end{pmatrix}$$

2 背包问题 $n=4, w=[2,7,3,5], p=[1,2,4,6], c=11$, 用带限界函数的回溯法求解。

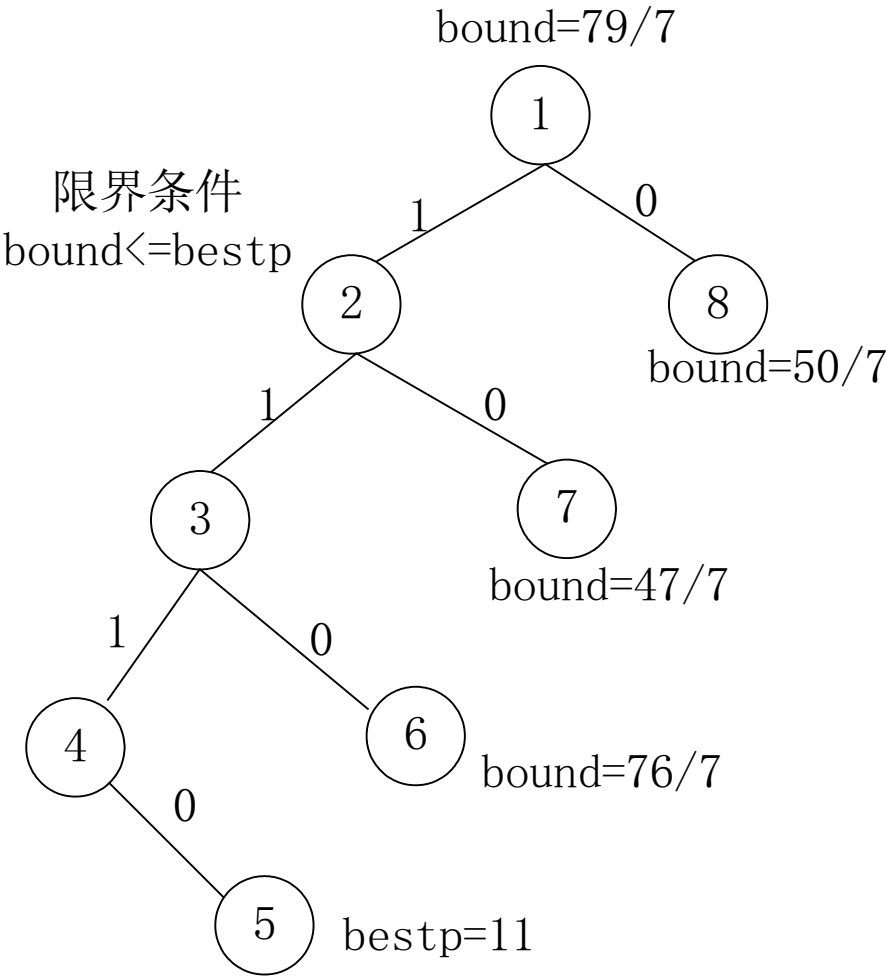
答: 效益密度为 $[0.5, 0.286, 1.33, 1.2]$, 按密度排列为 $[3, 4, 1, 2]$

$w'=[3, 5, 2, 7], p'=[4, 6, 1, 2]$

限界方法 1: $cp+r \leq bestp$, 则停止生成右子树。cp 为当前已得到的效益值, r 为尚为考虑的物品的效益值之和。

限界方法 2: 定义 $bound = cp + \text{对其余物品的贪心解效益值}$ 。如 $bound <$

=bestp 则停止产生右子树



解 为 $X =$

[1,0,1,1]