# Goon

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# Base Template

```
#include <bits/stdc++.h>
typedef long long 11;
typedef __int128 ll1;
typedef long double ld;
typedef __float128 lld;
using namespace std;
signed main() {
     ios_base::sync_with_stdio(false);
     cin.tie(NULL);
     return 0:
```

# Range Queries

## Segment Tree

Allows for single element modification, and range sum queries, for any binary commutative operation.

```
template <typename T>
struct Segtree {
   //note that t[0] is not used
   int n;
   T* t;
   T uneut, qneut;
   //single element modification function
   function<T(T, T)> fmodify:
   //product of two elements for query and updating tree
   function<T(T, T)> fcombine;
   Segtree(int n, T updateNeutral, T queryNeutral,
        function<T(T, T)> fmodify, function<T(T, T)>
        fcombine) {
       this -> n = n;
       t = new T[2 * n];
       this -> fmodify = fmodify:
       this -> fcombine = fcombine;
       uneut = updateNeutral;
       aneut = quervNeutral:
       for(int i = 0; i < n; i++){
          t[i + n] = uneut;
       build():
   }
   void build() { // build the tree after manually assigning
       for (int i = n - 1; i > 0; i--) {
          t[i] = fcombine(t[i * 2], t[i * 2 + 1]);
   }
   void modify(int p, T value) { // set value at position p
       t[p] = fmodify(t[p], value);
       for (p /= 2; p > 0; p /= 2) {
          t[p] = fcombine(t[p * 2], t[p * 2 + 1]);
   T query(int 1, int r) { // sum on interval [1, r)
       T I_res = qneut, r_res = qneut;
       bool l_none = true, r_none = true;
       for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
           if (1 % 2 == 1) {
              if(1 none) {
                  l_none = false;
```

```
1_{res} = t[1];
              else {
                 l res = fcombine(l res. t[1]):
              1++;
          if (r % 2 == 1) {
              if(r_none) {
                 r_none = false;
                 r_res = t[r];
              else {
                  r_res = fcombine(t[r], r_res);
      if(l_none) {
          return r_res;
       if(r_none) {
          return l_res;
      return fcombine(l_res, r_res);
   T query(int ind) {
       return this->query(ind, ind + 1);
//useful examples:
  -- INCREMENT MODIFY, SUM QUERY --
   function<int(int, int)> fmodify = [](const int src, const
        int val) -> int{return src + val;};
   function<int(int, int)> fcombine = [](const int a, const
        int b) -> int{return a + b;};
   Segtree<int> segt(n, 0, 0, fmodify, fcombine);
  -- ASSIGNMENT MODIFY, MIN QUERY --
   function<int(int, int)> fmodify = [](const int src, const
        int val) -> int{return val:};
   function<int(int, int)> fcombine = [](const int a, const
        int b) -> int{return min(a, b);};
   Segtree<int> segt(n, 0, 1e9, fmodify, fcombine);
// -- INCREMENT MODIFY, MAX SUBARRAY SUM QUERY --
//subarray has to contain at least 1 element.
   struct seg {
      11 max_pfx, max_sfx, max_sum, sum;
       seg() {};
       seg(11 sum) {
          this->sum = sum;
       seg(ll max_pfx, ll max_sfx, ll max_sum, ll sum) {
           this->max_pfx = max_pfx;
           this->max_sfx = max_sfx;
          this->max_sum = max_sum;
          this->sum = sum;
   function<seg(seg, seg)> fmodify = [](const seg src, const
       seg val) -> seg{
       seg next;
       next.max_pfx = src.max_pfx + val.sum;
      next.max_sfx = src.max_sfx + val.sum;
      next.max_sum = src.max_sum + val.sum;
      next.sum = src.sum + val.sum;
      return next:
   function < seg(seg, seg) > fcombine = [](const seg lhs,
        const seg rhs) -> seg{
       seg next;
```

```
next.max_pfx = max(lhs.max_pfx, lhs.sum + rhs.max_pfx);
       next.max_sfx = max(rhs.max_sfx, rhs.sum + lhs.max_sfx);
       next.max_sum = max({lhs.max_sum, rhs.max_sum,
            lhs.max_sfx + rhs.max_pfx});
       next.sum = lhs.sum + rhs.sum;
       return next;
   Segtree<seg> segt(n, {0, 0, 0, 0}, {0, 0, 0}, fmodify,
  -- ASSIGNMENT MODIFY, XOR BASIS QUERY --
//returns the xor basis of a range. Refer to
    https://codeforces.com/blog/entry/68953
   struct seg{
       int basis[20];
       int nr_b = 0;
       seg(int val) {
          fill(basis, basis + 20, -1);
          nr b += basisAdd(val):
       seg() {
           fill(basis, basis + 20, -1):
       bool basisAdd(int val, int start = 0) {
          for(int i = start; i < 20; i++){</pre>
              if((val & 1 << i) == 0){
                  continue:
              if(basis[i] == -1){
                  basis[i] = val;
                  return true;
              val ^= basis[i];
          return false;
   };
   function < seg(seg, seg) > fmodify = [] (const seg src, const
        seg val) -> seg{
       return val:
   function < seg(seg, seg) > fcombine = [](const seg a, const
       seg b) -> seg{
if(a.nr b == 20){
          return a:
       if(b.nr_b == 20){
          return b:
       seg next;
       for(int i = 0; i < 20; i++){
          next.basis[i] = a.basis[i];
       next.nr_b = a.nr_b;
       for(int i = 19; i \ge 0; i--){
          if(b.basis[i] != -1){
              next.basisAdd(b.basis[i], i);
       return next;
   Segtree<seg> segt(n, {}, fmodify, fcombine);
```

# Lazy Segment Tree

}

Allows for range modification, and range sum query for any binary commutative operation.

```
template <typename T>
struct SegtreeLazy {
   public:
      T* t;
             //stores product of range
```

```
T* d; //lazy tree
                                                                       t[ind] = fcombine(t[1], t[r]);
                                                                                                                                         int val) -> int{return val;};
bool* upd; //marks whether or not a lazy change is here
T uneut, queut;
                                                                    //registers a lazy change llo this node
//single element modify
                                                                   void applv(int ind, T val) {
                                                                                                                                        int b) -> int{return a + b;};
function<T(T, T)> fmodify;
                                                                       upd[ind] = true;
                                                                       d[ind] = fmodify(d[ind], val);
//k element modify
function<T(T, T, int)> fmodifyk;
                                                                                                                                   -- INCREMENT MODIFY, MINIMUM QUERY --
                                                                   //applies lazy change to this node
//product of two elements for query
                                                                    //k is the amount of values that this node represents.
function<T(T, T)> fcombine;
                                                                   void push(int ind, int k) {
                                                                       if(!upd[ind]) {
SegtreeLazy(int maxSize, T updateNeutral, T
    queryNeutral, T initVal, function<T(T, T)>
                                                                           return:
     fmodify, function<T(T, T, int)> fmodifyk,
                                                                       t[ind] = fmodifyk(t[ind], d[ind], k);
                                                                                                                                        int b) -> int{return min(a, b);};
    function<T(T, T)> fcombine) {
                                                                       if(ind < n) {
   n = maxSize;
                                                                           int 1 = ind * 2:
   uneut = updateNeutral;
                                                                           int r = ind * 2 + 1;
   qneut = queryNeutral;
                                                                           apply(1, d[ind]);
                                                                                                                                   -- ASSIGNMENT MODIFY, MINIMUM QUERY --
                                                                           apply(r, d[ind]);
   this -> fmodify = fmodify;
   this -> fmodifyk = fmodifyk;
   this -> fcombine = fcombine:
                                                                       upd[ind] = false;
                                                                                                                                        int val) -> int{return val:}:
                                                                       d[ind] = uneut;
   //raise n to nearest pow 2
   int x = 1;
   while(x < n) {
                                                                   void _modify(int 1, int r, T val, int tl, int tr, int
                                                                                                                                        int b) -> int{return min(a, b);};
                                                                        ind) { `
       x <<= 1;
                                                                       if(1 == r){
   n = x:
                                                                           return;
   t = new T[n * 2];
                                                                       if(upd[ind]) {
   d = new T[n * 2];
                                                                           push(ind, tr - tl);
                                                                                                                                     and last indices of 0 and 1
   upd = new bool[n * 2];
                                                                                                                                // in the range that you queried.
                                                                       if(1 == t1 && r == tr) {
   //make sure to initialize values
                                                                           apply(ind, val);
                                                                                                                                    struct seg {
   for(int i = 0; i < n; i++){
                                                                           push(ind, tr - tl);
       t[i + n] = initVal;
                                                                                                                                       int size:
                                                                           return:
                                                                                                                                       bool has_one, has_zero;
   for(int i = n - 1; i > 0; i--){
                                                                       int mid = t1 + (tr - t1) / 2:
                                                                                                                                             end of closest 1
       t[i] = fcombine(t[i * 2], t[i * 2 + 1]);
                                                                       if(1 < mid) {
                                                                                                                                       int pfx_zero, sfx_zero;
                                                                           _{\text{modify}}(1, \min(r, \min), \text{ val}, \text{ tl}, \min, \text{ ind } * 2);
   for(int i = 0; i < n * 2; i++){
                                                                                                                                            the same value
       d[i] = uneut;
                                                                       if(r > mid) {
                                                                                                                                           this->size = size:
       upd[i] = false;
                                                                           _{modify(max(1, mid), r, val, mid, tr, ind * 2 +
                                                                                                                                           has_one = which;
                                                                                                                                           has zero = !which:
                                                                                                                                           pfx_one = 0;
                                                                                                                                           sfx_one = 1;
                                                                       combine(ind, tr - tl);
void modify(int 1, int r, T val) { //modifies the
                                                                                                                                           pfx_zero = 0;
     range [1, r)
                                                                                                                                           sfx_zero = 1;
    _{\text{modify}}(\hat{1}, r, val, 0, n, 1);
                                                                   T _query(int 1, int r, int tl, int tr, int ind) {
                                                                                                                                       seg(int size){
                                                                       if(I == r)
                                                                                                                                           this->size = size;
                                                                           return queut;
void modify(int ind, T val) { //modifies the range
                                                                                                                                           has one = false:
                                                                                                                                           has_zero = false;
     [ind, ind + 1)
                                                                       if(upd[ind]) {
    _modify(ind, ind + 1, val, 0, n, 1);
                                                                           push(ind, tr - tl);
                                                                                                                                       seg() {}
                                                                       if(l == tl && r == tr){
T query(int 1, int r) { //queries the range [1, r)
                                                                           return t[ind];
   return _query(1, r, 0, n, 1);
                                                                                                                                         seg val) -> seg{
                                                                       int mid = tl + (tr - tl) / 2:
                                                                                                                                       //set this element to 0 or 1
                                                                       T lans = qneut;
T query(int ind) { //queries the range [ind, ind + 1)
                                                                                                                                       seg next(src.size, val.has_one);
                                                                       T rans = qneut;
   return _query(ind, ind + 1, 0, n, 1);
                                                                                                                                       return next:
                                                                       if(1 < mid) {
                                                                           lans = _{query}(1, min(r, mid), tl, mid, ind * 2);
                                                                       if(r > mid) {
//calculates value of node based off of children
                                                                                                                                       //set the entire range to 0 or 1
                                                                           rans = _query(max(1, mid), r, mid, tr, ind * 2
//k is the amount of values that this node represents.
                                                                                                                                       seg next(src.size, val.has one):
void combine(int ind, int k) {
                                                                                                                                       return next:
   if(ind >= n){
       return;
                                                                       return fcombine(lans, rans);
                                                                                                                                        const seg rhs) -> seg{
                                                            };
   int 1 = ind * 2:
   int r = ind * 2' + 1;
   //make sure children are correct value before
                                                            //useful examples
        calculating
                                                               -- ASSIGNMENT MODIFY, SUM QUERY --
   push(1, k / 2);
                                                                                                                                       if(next.has_one) {
   push(r, k / 2);
                                                                function<int(int, int)> fmodify = [](const int src, const
```

```
function<int(int, int, int)> fmodifyk = [](const int src,
        const int val, const int k) -> int{return val * k;};
   function<int(int, int)> fcombine = [](const int a, const
   run_segt_tests(n, 0, 0, fmodify, fmodifyk, fcombine);
   function<int(int, int)> fmodify = [](const int src, const
        int val) -> int{return src + val:}:
    function<int(int, int, int)> fmodifyk = [](const int src,
        const int val, const int k) -> int{return src + val;};
   function<int(int, int)> fcombine = [](const int a, const
   run_segt_tests(n, 0, 1e9, fmodify, fmodifyk, fcombine);
   function<int(int, int)> fmodify = [](const int src, const
   function<int(int, int, int)> fmodifyk = [](const int src,
   const int val, const int k) -> int{return val;};
function<int(int, int)> fcombine = [](const int a, const
   run_segt_tests(n, 0, 1e9, fmodify, fmodifyk, fcombine);
// -- O1 ASSIGNMENT MODIFY, FIRST LAST INDEX O1 QUERY --
// when querying, returns a struct that tells you the first
       int pfx_one, sfx_one; //distance from beginning and
       seg(int size, bool which) { //sets the entire segment
    //assignment modify, range 'seg' query.
   function<seg(seg, seg)> fmodify = [](const seg src, const
   function<seg(seg, seg, int)> fmodifyk = [](const seg src,
        const seg val, const int k) -> seg{
   function<seg(seg, seg)> fcombine = [](const seg lhs,
       //combines lhs and rhs into one segment
       seg next(lhs.size + rhs.size, false);
       next.has_one = lhs.has_one || rhs.has_one;
       next.has_zero = lhs.has_zero || rhs.has_zero;
           next.pfx_one = lhs.has_one? lhs.pfx_one : lhs.size
```

# 2.3 Lazy Segment Tree (Duckling)

```
template<typename T, typename D>
struct Lazy {
   static constexpr T qn = 0; //stores the starting values
        at all nodes
   static constexpr D ln = 0;
   vector<T> v;
                   //stores values at each index we are
        querying for
   vector <D> lazy; //base, count of how many polynomials
        start at one at the beginning of this node
   //if OJ is not up to date, remove all occurrences of ln
   Lazy(int n = 0, \bar{T} def = qn) {
       this->n = n;
       this->size = 1;
       while(size < n) size *= 2;</pre>
       v.assign(size * 2, def);
       lazy.assign(size * 2, ln);
   bool isLeaf(int node) {
       return node >= size;
   T query_comb(T val1, T val2) {//update this depending on
       return val1 + val2:
   //how we combine lazy updates to lazy
   void lazy_comb(int node, D val) {//update this depending
        on update type. how do we merge the lazy changes?
       lazy[node] += val;
   void main_comb(int node, int size) {//update this
        depending on query type, how does the lazy value affect value at v for the query?
       v[node] += lazy[node];
   void push_lazy(int node, int size) {
       main_comb(node, size); //push lazy change to current
       if(!isLeaf(node)) {
           lazy_comb(node * 2, lazy[node]);
           lazy_comb(node * 2 + 1, lazy[node]);
       lazy[node] = ln;
   void update(int 1, int r, D val) {
       _update(1,0,size,l,r, val);
   void _update(int node, int currl, int currr, int
        &targetl, int &targetr, D val) {
       if (currl >= targetr || currr <= targetl) return;</pre>
       push_lazy(node, currr - currl);
       if(currl >= targetl && currr <= targetr) { //complete
           lazy_comb(node, val); //we apply the lazy change
                to this node, then update this node
```

```
} else { //partial overlap, should never be a leaf,
        otherwise it'd fall under previous categories
       int mid = (currl + currr) / 2;
       _update(node * 2, currl, mid, targetl, targetr,
        _update(node * 2 + 1, mid, currr, targetl,
            targetr, val);
       push_lazy(node * 2, (currr - currl) / 2);
       push_lazy(node * 2 + 1, (currr - currl) / 2);
       v[node] = query\_comb(v[node * 2], v[node * 2 + 1]);
T query(int 1, int r) {
   return _query(1,0,size,l,r);
T_query(int node, int currl, int currr, int &targetl,
    int &targetr) { //[1,r)
   if(currr <= targetl || currl >= targetr) return qn;
   push_lazy(node, currr-currl); //make pushes necessary
         before getting value, we always check for 2 cases
   if(currl >= targetl && currr <= targetr) { //complete</pre>
        overlap
       return v[node];
   } else {
       int mid = (currl + currr) / 2;
       return query_comb(
           _query(node * 2, currl, mid, targetl, targetr),
           _query(node * 2 + 1, mid, currr, targetl,
                targetr)
   }
}
```

# 3 Graphs

#### 3.1 DSU

};

```
//basic union find implementation with path compression
struct DSU {
    int N:
    vector<int> dsu;
    DSU(int n) {
       this->N = n;
       this->dsu = vector<int>(n, 0):
       for(int i = 0; i < n; i++){ //initialize roots</pre>
           dsu[i] = i;
    void dsu_init() {
       for(int i = 0; i < N; i++){</pre>
           dsu[i] = i;
    int find(int a) {
       if(dsu[a] == a) {
           return a;
       return dsu[a] = find(dsu[a]);
    //ret true if updated something
    bool unify(int a, int b) {
       int ra = find(a);
       int rb = find(b);
       if(ra == rb) {
           return false:
       dsu[rb] = ra;
       return true:
};
```

# 3.2 Bridges and Articulation Points

A bridge is an edge who's deletion will disconnect the graph. Articulation points are the same, but for nodes. The endpoints of a bridge are always articulation points (except in the case where an endpoint has degree 1), but it is possible for a non-bridge to have two articulation points as endpoints.

#### 3.2.1 Bridges

If edge (u, v) is a bridge, then the returned set should contain (u, v) and (v, u). As of 2/22/2024, this is untested.

```
set<pair<int, int>> find_bridges(int n, vector<vector<int>>&
   vector<bool> visited;
   vector<int> tin, low;
   int timer:
   timer = 0:
   visited.assign(n, false);
   tin.assign(n, -1);
   low.assign(n, -1);
   set<pair<int, int>> ans;
   function<void(int, int)> dfs = [&visited, &tin, &low,
        &timer, &dfs, &adj, &ans](int v, int p) -> void {
       visited[v] = true;
       tin[v] = low[v] = timer++;
      for (int to : adj[v]) {
          if (to == p) continue;
           if (visited[to]) {
              low[v] = min(low[v], tin[to]);
          } else {
              dfs(to, v);
              low[v] = min(low[v], low[to]);
              if (low[to] > tin[v]) {
                  ans.insert({v, to});
                  ans.insert({to, v});
      }
   };
   for (int i = 0; i < n; ++i) {</pre>
       if (!visited[i])
          dfs(i, -1);
   return ans;
```

#### 3.2.2 Articulation Points

Returns a boolean array, a[i] is true if node i is an articulation point.

```
vector<bool> find_articulation_points(int n,
    vector<vector<int>>& adj) {
   vector<bool> visited(n, false);
   vector<int> tin(n, -1);
   vector<int> low(n, -1);
   vector<bool> is_articulation_point(n, false);
   function<void(int, int)> dfs = [&visited, &tin, &low,
        &is_articulation_point, &timer, &adj, &dfs](int v,
        int p) -> void {
       visited[v] = true;
       tin[v] = low[v] = timer++;
       int children=0;
       for (int to : adj[v]) {
           if (to == p) continue;
           if (visited[to]) {
              low[v] = min(low[v], tin[to]);
           } else {
              dfs(to, v);
              low[v] = min(low[v], low[to]);
```

## 3.3 Tarjan SCC

Returns multiple lists of node ids, each list being a SCC.

```
vector<vector<int>>> find_scc(int n, vector<vector<int>>& adj)
   vector<vector<int>> adj_rev(n, vector<int>(0));
   vector<bool> used(n, false);
   vector<int> order(0);
   for(int i = 0; i < n; i++){</pre>
       for(int j = 0; j < adj[i].size(); j++){
   adj_rev[adj[i][j]].push_back(i);</pre>
   function < void (int) > dfs1 = [&used, &adj, &order,
        &dfs1](int v) -> void {
       used[v] = true;
       for (auto u : adj[v]) {
           if (!used[u]) {
               dfs1(u);
       order.push_back(v);
   for(int i = 0; i < n; i++){
       if(used[i]) {
           continue:
       dfs1(i);
   fill(used.begin(), used.end(), false);
   reverse(order.begin(), order.end());
   function < void (int, vector < int > k) > dfs2 = [kused,
        &adj_rev, &dfs2](int v, vector<int>& component) ->
        void {
       used[v] = true:
       component.push_back(v);
       for (auto u : adj_rev[v]) {
           if (!used[u]) {
               dfs2(u, component);
       }
   };
   vector<vector<int>> ans(0);
   for(int i = 0; i < n; i++){</pre>
       if(used[order[i]]){
           continue;
       vector<int> component(0);
       dfs2(order[i], component);
       ans.push_back(component);
   return ans:
```

#### 3.4 LCA

Lets you find the Lowest Common Ancestor of two nodes in a rooted tree in  $O(\log(n))$  time. Note that the segment tree

used in this LCA is slightly modified.

```
struct LCA {
   struct Segtree {
       //note that t[0] is not used
       int n:
       int* t;
       int* node_id;
       int uneut, queut;
       //single element modification function
       function<int(int, int)> fmodify;
       //product of two elements for query and updating tree
       function<int(int, int)> fcombine;
       Segtree() {
           //do nothing
       Segtree(int n, int updateNeutral, int queryNeutral,
            function<int(int, int)> fmodify,
            function<int(int, int)> fcombine) {
           this -> n = n;
           t = new int[2 * n]
           node_id = new int[2 * n];
           this -> fmodify = fmodify:
           this -> fcombine = fcombine;
           uneut = updateNeutral:
           qneut = queryNeutral;
           for(int i = 0; i < 2 * n; i++){
              t[i] = uneut;
              node_id[i] = 0;
       void build() { // build the tree after manually
            assigning the values.
           for (int i = n - 1; i > 0; i--) {
              t[i] = fcombine(t[i * 2], t[i * 2 + 1]);
              node_id[i] = t[i] == t[i * 2]? node_id[i * 2] :
                   node_id[i * 2 + 1];
          }
       }
       int query(int 1, int r) { // least deep node on
            interval [1, r)
           int min_depth = 1e9;
           int res = -1:
           for (1 += n, r += n; 1 < r; 1 /= 2, r /= 2) {
              if (1 % 2 == 1) {
                  if(t[1] < min_depth) {</pre>
                      res = node_id[1];
                      min_depth = t[1];
                  1++;
              if (r % 2 == 1) {
                  if(t[r] < min_depth) {</pre>
                      res = node_id[r];
                      min_depth = t[r];
              }
           return res;
   };
   int n;
   vector<vector<int>> edges;
   vector<int> depth; //distance of each node from the root
    vector<int> left_occ, right_occ; //leftmost and rightmost
        occurrences for each node in the euler tour
```

```
//single assignment modify, range min query
//stores the euler tour of the tree to compute lca
Segtree segt;
void euler_tour(int cur, int p, vector<int>& ret) {
   left_occ[cur] = ret.size();
   ret.push_back(cur);
   for(int i = 0; i < edges[cur].size(); i++){
       int next = edges[cur][i];
       if(next == p){
           continue:
       euler_tour(next, cur, ret);
       ret.push_back(cur);
   right_occ[cur] = ret.size();
void find_depth(int cur, int p) {
   for(int i = 0; i < edges[cur].size(); i++){</pre>
       int next = edges[cur][i];
       if(next == p){
           continue;
       depth[next] = depth[cur] + 1;
       find_depth(next, cur);
}
void init(int n. int root, vector<vector<int>>& edges) {
   this->n = n;
   this->root = root;
   this->edges = edges;
   this->depth = vector<int>(n, 0);
   find_depth(root, -1);
   vector<int> tour(0);
   this->left_occ = vector<int>(n, -1);
   this->right_occ = vector<int>(n, -1);
   euler_tour(root, -1, tour);
   function<int(int, int)> fmodify = [](const int src,
        const int val) -> int{return val;};
   function<int(int, int)> fcombine = [](const int a,
        const int b) -> int{return min(a, b);};
   this->segt = Segtree(tour.size(), 0, 1e9, fmodify,
        fcombine);
   for(int i = 0; i < tour.size(); i++){</pre>
       segt.node_id[i + tour.size()] = tour[i];
       segt.t[i + tour.size()] = depth[tour[i]];
   segt.build();
}
//adjacency list constructor
LCA(int n, int root, vector<vector<int>> edges) {
   init(n, root, edges);
//parent list constructor
//if node i is the root, then parents[i] must equal -1
LCA(int n, vector<int> parents) {
   int root = -1;
   vector<vector<int>> edges(n, vector<int>(0));
   for(int i = 0; i < n; i++){
       if(parents[i] == -1){
           root = i:
       edges[parents[i]].push_back(i);
       edges[i].push_back(parents[i]);
   init(n, root, edges);
int lca(int a, int b) {
   int l = min(left_occ[a], left_occ[b]);
   int r = max(right_occ[a], right_occ[b]);
   int lc = segt.query(1, r);
```

```
return lc;
}
int dist(int a, int b) {
   int lc = lca(a, b);
   return depth[a] + depth[b] - 2 * depth[lc];
}
};
```

# 3.5 Centroid Decomposition

A Centroid of a tree is a node such that when the tree is rooted at it, no other nodes have a subtree of size greater than  $\frac{N}{2}$ . Restructures the tree such that the maximum distance from root to leaf is  $O(\log(n))$ . Useful for solving problems like answering queries that ask what's the closest white node in a white black tree online, with color updates.

```
struct CentroidDecomp {
   CentroidDecomp() {
       //yay
   vector<bool> vis:
   vector<int> centroid_parent;
   vector<int> size; //size of subtree in original tree
   vector<vector<int>> edges;
   int find size(int cur. int p = -1) {
       if(vis[cur]) {
           return 0:
       size[cur] = 1;
       for(int i = 0; i < edges[cur].size(); i++){</pre>
           int next = edges[cur][i];
           if(next != p){
    size[cur] += find_size(next, cur);
       return size[cur];
   int find_centroid(int cur, int p, int sub_size) {
       for(int i = 0; i < edges[cur].size(); i++){</pre>
           int next = edges[cur][i];
if(next == p){
           if(!vis[next] && size[next] > sub_size / 2) {
               return find_centroid(next, cur, sub_size);
       return cur;
   void init_centroid(int cur, int p = -1) {
       find_size(cur);
       int centroid = find_centroid(cur, -1, size[cur]);
       vis[centroid] = true;
       centroid_parent[centroid] = p;
       for(int i = 0; i < edges[centroid].size(); i++){</pre>
           int next = edges[centroid][i];
           if(!vis[next]){
               init_centroid(next, centroid);
   //returns an array 'a' where the parent of node i is a[i].
   //if i is the root, then a[i] = -1.
   vector<int> calc_centroid_decomp(int n,
        vector<vector<int>>& adj_list) {
```

```
edges = adj_list;
  vis = vector<bool>(n, false);
  centroid_parent = vector<int>(n, -1);
  size = vector<int>(n, -1);
  init_centroid(0);
  return centroid_parent;
}
};
```

#### 3.5.1 Find Centroids

Takes in adjacency list, returns a list of the centroids of the tree.

```
vector<int> findCentroid(const vector<vector<int>> &g) {
   int n = g.size();
   vector<int> centroid;
   vector<int> sz(n);
   function<void (int, int)> dfs = [&](int u, int prev) {
      sz[u] = 1;
      bool is_centroid = true;
      for (auto v : g[u]) if (v != prev) {
            dfs(v, u);
            sz[u] += sz[v];
            if (sz[v] > n / 2) is_centroid = false;
      }
      if (n - sz[u] > n / 2) is_centroid = false;
      if (is_centroid) centroid.push_back(u);
    };
    dfs(0, -1);
    return centroid;
}
```

## 3.6 Heavy-Light Decomposition

 $O(log(n)^2)$  modify and query over any path in the tree. LCA, and Lazy Segment Tree not included.

```
template <typename T>
struct HLD {
   vector<vector<int>> edges:
   vector<bool> toParentHeavy
   vector<bool> hasOutHeavy;
   vector<int> parent;
vector<int> subtreeSize;
   SegtreeLazy<T> segt;
   vector<int> segEndInd; //stores the index at which this
        heavy path ends.
   vector<int> segBeginInd; //stores the index at which this
        heavy path begins.
   vector<int> segParent; //what is the parent node of this
        heavy path
   vector<int> segPos; //stores the index of each node
        within the segment tree.
   vector<T> maxSegCache; //stores results to segments which
        take up the entire heavy path
   void calcSubtreeSize(int cur, int p = -1) {
       parent[cur] = p;
subtreeSize[cur] = 1;
       for(int i = 0; i < edges[cur].size(); i++){</pre>
           int next = edges[cur][i];
           if(next == p){
               continue;
           calcSubtreeSize(next, cur);
           subtreeSize[cur] += subtreeSize[next];
   void calcHLD(int cur, int p = -1) {
       for(int i = 0; i < edges[cur].size(); i++){</pre>
```

```
int next = edges[cur][i];
        if(next == p){
            continue:
        if(subtreeSize[next] > subtreeSize[cur] / 2) {
            hasOutHeavy[cur] = true;
            toParentHeavy[next] = true;
        calcHLD(next, cur);
HLD(int n, int root, vector<vector<int>> adjList, T
    updateNeutral, T queryNeutral, T initVal,
     function<T(T, T)> fmodify, function<T(T, T, int)>
     fmodifyk, function\langle T(T, T) \rangle fcombine) {
    this->lca = LCA(n, root, adjList);
    this->edges = adjList;
    this->parent = vector<int>(n, -1);
    this->subtreeSize = vector<int>(n, 0);
    this->toParentHeavy = vector<bool>(n, false);
   this->hasOutHeavy = vector<bool>(n, false);
this->calcSubtreeSize(root);
    this->calcHLD(root);
    //create the segment tree needed to do the range
    this->segt = SegtreeLazy<T>(n, updateNeutral,
         queryNeutral, initVal, fmodify, fmodifyk,
    this->segBeginInd = vector<int>(n, -1);
    this->segEndInd = vector<int>(n, -1);
    this->segParent = vector<int>(n, -1);
    this->segPos = vector<int>(n, -1):
    //find the positions of the nodes in the segment tree.
    int posPtr = 0;
    for(int i = 0; i < n; i++){
        if(this->hasOutHeavy[i]) {
            //we want to have each heavy path be contiguous
                 in the segment tree, so we want to start
                 at the beginning.
            continue;
        int begin = posPtr;
int cur = i;
        vector<int> heavyPath(0);
        heavyPath.push_back(cur);
        this->segPos[cur] = posPtr ++;
        while(toParentHeavy[cur]) {
            cur = parent[cur];
            heavyPath.push_back(cur);
            this->segPos[cur] = posPtr ++;
        cur = parent[cur];
        for(int j = 0; j < heavyPath.size(); j++){</pre>
           this->segBeginInd[heavyPath[j]] = begin;
this->segEndInd[heavyPath[j]] = posPtr;
this->segParent[heavyPath[j]] = cur;
    //compute max cache values
    this->maxSegCache = vector<T>(n);
    for(int i = 0; i < n; i++){</pre>
        if(this->segPos[i] == this->segBeginInd[i]) {
            int begin = this->segBeginInd[i];
            int end = this->segEndInd[i];
            this->maxSegCache[begin] =
                 this->segt.query(begin, end);
void modify(int a, int b, T val) {
    int _lca = this->lca.lca(a, b);
    _modify(a, _lca, val);
   _modify(b, _lca, val);
```

```
this->modify(_lca, val);
   void modify(int a, T val) {
       this->segt.modify(this->segPos[a], val);
       //update cache
       int begin = this->segBeginInd[a];
       int end = this->segEndInd[a];
       this->maxSegCache[begin] = this->segt.query(begin,
   T query(int a, int b) {
       int _lca = this->lca.lca(a, b);
       T ret = this->segt.qneut;
       ret = this->segt.fcombine(ret, _query(a, _lca));
       ret = this->segt.fcombine(ret, _query(b, _lca));
       ret = this->segt.fcombine(ret, this->query(_lca));
       return ret;
   T query(int a) {
       return this->segt.query(this->segPos[a]);
       void _modify(int a, int _lca, int val) {
           //while a and _lca aren't in the same heavy path
           while(this->segEndInd[a] != this->segEndInd[_lca])
               //modify until the end of the segment a belongs
               this->segt.modify(this->segPos[a],
                    this->segEndInd[a], val);
               a = this->segParent[a]:
               //update cache
               int begin = this->segBeginInd[a];
               int end = this->segEndInd[a];
               this->maxSegCache[begin] =
                    this->segt.query(begin, end);
           //a and _lca are in the same heavy path. Now, just
modify the segment from a to _lca, not
                including _lca.
           this->segt.modify(this->segPos[a],
                this->segPos(_lca], val);
       T _query(int a, int _lca) {
           Tret = this->segt.qneut;
           while(this->segEndInd[a] != this->segEndInd[_lca])
               //see if we can use the cache
               if(this->segBeginInd[a] == this->segPos[a]) {
                  //use the cache
                  ret = this->segt.fcombine(ret,
                       this->maxSegCache[this->segBeginInd[a]]);
              }
               else {
                  ret = this->segt.fcombine(ret,
                       this->segt.query(this->segPos[a],
                       this->segEndInd[a]));
               a = this->segParent[a];
           ret = this->segt.fcombine(ret,
                this->segt.query(this->segPos[a],
                this->segPos[_lca]));
           return ret;
};
```

#### 3.7 Min Cost Flow

Given some amount of flow, what's the minimum cost required to achieve that flow? Can also be used to compute max flow if max\_flow = INF. Slightly modified MCMF template from KACTL.

```
#include <bits/extc++.h>
#define all(x) begin(x), end(x)
#define sz(x) (int) (x).size()
#define rep(i, a, b) for(int i = a; i < (b); i++)
typedef pair<int, int> pii;
typedef vector<11> VL;
const 11 INF = numeric limits<11>::max() / 4;
struct MCMF {
    struct edge {
        int from, to, rev;
        11 cap, cost, flow;
    vector<vector<edge>> ed;
    vector<int> seen;
    vector<ll> dist, pi;
    vector<edge*> par;
    MCMF(int N) : N(N), ed(N), seen(N), dist(N), pi(N),
         par(N) {}
    void addEdge(int from, int to, ll cap, ll cost) {
       if (from == to) return;
        ed[from].push_back(edge{ from,to,sz(ed[to]),cap,cost,0
        ed[to].push_back(edge{
             to,from,sz(ed[from])-1,0,-cost,0 });
    void path(int s) {
       fill(all(seen), 0);
fill(all(dist), INF);
        dist[s] = 0; 11 di;
        __gnu_pbds::priority_queue<pair<11, int>> q;
vector<decltype(q)::point_iterator> its(N);
        q.push({ 0, s });
        while (!q.empty()) {
            s = q.top().second; q.pop();
seen[s] = 1; di = dist[s] + pi[s];
            for (edge& e : ed[s]) if (!seen[e.to]) {
                11 val = di - pi[e.to] + e.cost;
                if (e.cap - e.flow > 0 && val < dist[e.to]) {</pre>
                    dist[e.to] = val;
par[e.to] = &e;
                    if (its[e.to] == q.end())
                        its[e.to] = q.push({ -dist[e.to], e.to
                        q.modify(its[e.to], { -dist[e.to], e.to
                }
           }
        rep(i,0,N) pi[i] = min(pi[i] + dist[i], INF);
    pair<11, 11> maxflow(int s, int t, 11 max_flow = INF) {
       11 totflow = 0, totcost = 0;
while (path(s), seen[t] && totflow < max_flow) {</pre>
            11 fl = max_flow - totflow;
            for (edge* x = par[t]; x; x = par[x->from])
                fl = min(fl, x->cap - x->flow);
            totflow += fl;
            for (edge* x = par[t]; x; x = par[x->from]) {
   x->flow += fl;
                ed[x->to][x->rev].flow -= fl;
```

#### 3.8 2SAT

Given a boolean expression like  $(a \lor b) \land (\neg b \lor c) \land (c \land \neg a)$ , figures out whether there is a valid assignment to the variables, a, b, c. If a valid assignment exists, can produce an example. Depends on find\_scc to work.

```
struct TSAT {
   int n:
   vector<vector<int>>> node_id; //{false, true}
   vector<vector<int>> c;
   TSAT(int n) {
       this->n = n;
       this->node_id = vector<vector<int>>(n, {0, 0});
       for(int i = 0; i < n; i++){</pre>
          node_id[i][0] = ptr ++;
node_id[i][1] = ptr ++;
       this->c = vector<vector<int>>(n * 2, vector<int>(0));
   //clears all implications
   void clear() {
       this->c = vector<vector<int>>(n * 2, vector<int>(0));
   //a being a_state implies b being b_state
   void imply(int a, bool a_state, int b, bool b_state) {
       int a_id = node_id[a][a_state];
       int b_id = node_id[b][b_state];
       c[a_id].push_back(b_id);
                                         //positive
       c[b_id ^ 1].push_back(a_id ^ 1); //contrapositive
   //forces a to be state
   void set(int a, int state) {
                                         //if a is !state.
       imply(a, !state, a, state);
            then we have a contradiction, so a must be state.
   //at least one of a has to be a_state or b has to be
   void addOR(int a, bool a_state, int b, bool b_state) {
       imply(a, !a_state, b, b_state); //if a is not good,
            then b has to be
   //exactly one of a has to be a_state or b has to be
   void addXOR(int a, bool a_state, int b, bool b_state) {
       imply(a, !a_state, b, b_state); //normal or
```

```
imply(a, a_state, b, !b_state); //if a is good, then
        b can't be
//either both a and b are good, or both a and b are not
     good
void addXNOR(int a, bool a_state, int b, bool b_state) {
   imply(a, a_state, b, b_state); //if a is good, then
        b has to be good
    imply(b, b_state, a, a_state); //if b is good, then
        a has to be good
}
//a and b both have to be good
//equivalent to set(a, a_state); set(b, b_state);
void addAND(int a, bool a_state, int b, bool b_state) {
    imply(a, a_state, b, b_state); //if a is good, then
        b has to be good
    imply(b, b_state, a, a_state); //if b is good, then
        a has to be good
    set(a, a_state);
                                     //make sure that a is
        good
//if a solution exists, returns a possible configuration
     of the variables.
//otherwise, returns an empty vector
vector<bool> generateSolution() {
    //first, split into sccs.
    vector<vector<int>> scc = find_scc(n * 2, c);
    //check for contradictions, eg if a and !a are in the
        same scc
    vector<int> node_scc(n * 2);
    for(int i = 0; i < scc.size(); i++){</pre>
       for(int j = 0; j < scc[i].size(); j++){</pre>
           int id = scc[i][j];
           node_scc[id] = i;
    for(int i = 0; i < n; i++){
       if(node\_scc[i * 2 + 0] == node\_scc[i * 2 + 1]){
          return {}:
    //otherwise, a solution always exists
    vector<bool> v(n, false);
    vector<bool> ans(n, false);
    //toposort scc
    vector<vector<int>> scc_c(scc.size(), vector<int>(0));
    for(int i = 0; i < node_scc.size(); i++){</pre>
       int cur = i:
       int cur_scc = node_scc[i];
       for(int j = 0; j < this->c[cur].size(); j++){
           int next = this->c[cur][j];
           int next_scc = node_scc[next];
           if(next_scc != cur_scc) {
              scc_c[cur_scc].push_back(next_scc);
   vector<int> scc_indeg(scc.size(), 0);
    for(int i = 0; i < scc_c.size(); i++){</pre>
       for(int j = 0; j < scc_c[i].size(); j++){</pre>
           int next_scc = scc_c[i][j];
           scc_indeg[next_scc] ++;
    queue<int> q;
    for(int i = 0; i < scc_indeg.size(); i++){</pre>
       if(scc_indeg[i] == 0){
           q.push(i);
    vector<int> toporder:
    while (q.size()^{\dagger}!=0) {
       int cur = q.front();
       q.pop();
       toporder.push back(cur):
```

```
for(int i = 0; i < scc_c[cur].size(); i++){
    int next = scc_c[cur][i];
    scc_indeg[next] --;
    if(scc_indeg[next] == 0){
        q.push(next);
    }
}

//assign the answers in reverse topological order
for(int i = toporder.size() - 1; i >= 0; i--){
    int cur_scc = toporder[i];
    for(int j = 0; j < scc[cur_scc].size(); j++){
        int cur = scc[cur_scc][j];
        bool state = cur % 2;
        cur /= 2;
        if(v[cur]) {
            break;
        }
        v[cur] = true;
        ans[cur] = state;
    }
}

bool solutionExists() {
    return generateSolution().size() != 0;
}</pre>
```

# 4 Strings

};

#### 4.1 Suffix Tree

Equivalent to a trie that contains all suffixes of a string S, but only uses  $O(\log(A)|S|)$  memory, where A is the size of the alphabet. Can find the number of occurrences of T in S, compute the suffix array of S, and compute the Longest Common Prefix between two indices in S.

Make sure to modify the terminator character if '\$' is used as a character in the input. Terminator character is to 'flush' the buffered changes so that all suffixes are in the tree.

```
struct SuffixTree {
       struct SuffixNode {
           //l, r: left and right boundaries [1, r) of the
                edge that leads to this node.
           //parent: index of parent node
           //link: index of link node
           int index:
           int 1, r;
           SuffixNode* parent;
SuffixNode* link;
           map<char, SuffixNode*> children;
           SuffixNode(int index, int 1 = 0, int r = 0,
                SuffixNode* parent = nullptr) : index{index},
                1{1}, r{r}, parent{parent}, link{nullptr} {}
           int len() {return r - 1;}
           SuffixNode* get_child(char c) {
               if(children.find(c) == children.end()) {
                  return nullptr;
              return children[c];
           void set_child(char c, SuffixNode* ptr) {
               if(children.find(c) == children.end()) {
                  children.insert({c, ptr});
               children[c] = ptr;
```

```
};
   vector<char> chars:
   vector<SuffixNode*> nodes:
   SuffixTree(string s) {
       //add terminator char
       s.push_back('$');
       //build tree
       this->n = s.size():
       for(int i = 0; i < s.size(); i++){</pre>
           this->add_char(s[i]);
       //calculate useful information
       this->calc leaf cnt():
       this->calc_suf_arr();
       this->calc_lcp();
   //runs in O(|s|) time.
   bool contains_string(string s) {
       return count_occurrences(s);
   //each leaf node corresponds to a suffix, so it
        suffices to see how many leaf nodes there
    //are in the subtree corresponding to s.
   int count_occurrences(string s) {
       int i = 0:
       SuffixNode* node = this->nodes[0];
       while(i != s.size()) {
           SuffixNode* child = node -> get_child(s[i]);
           if(child == nullptr) {
               return 0;
           node = child:
           for(int j = child -> 1; j < min(child -> r,
                (int) this->chars.size()) && i < s.size();</pre>
               if(this->chars[j] != s[i]) {
                  return 0;
              i ++:
       return this->leaf_cnt[node -> index];
   vector<int> get_suffix_array() {
       return this->suf_arr;
   int get_lcp(int a, int b) {
       if(a == b){
           return this->n - a - 1;
       int a_ind = this->suf_to_suf_ind[a];
       int b_ind = this->suf_to_suf_ind[b];
       if(a_ind > b_ind) {
           swap(a_ind, b_ind);
       return this->lcp_rmq.query(a_ind, b_ind);
private:
   Segtree<int> lcp_rmq;
   vector<int> suf_arr; //suffix array
   vector<int> suf_to_suf_ind; //maps suffix indices to
        their locations in the suffix array.
   vector<int> leaf_cnt; //number of leaves in subtree
   vector<int> lcp; //longest common prefix between
adjacent suffixes in suffix array.
   //uses kasai's algorithm to compute lcp array in O(n).
   //lcp stands for longest common prefix.
   void calc_lcp() {
```

```
int k = 0;
   int n = this -> n;
   vector<int> lcp(n,0);
   vector<int> rank(n,0);
   for(int i = 0; i < n; i++) {</pre>
       rank[this->suf_arr[i]] = i;
   for(int i = 0; i < n; i++, k? k-- : 0) {
       if(rank[i]==n-1) {
          k = 0;
          continue;
       int j = this->suf_arr[rank[i] + 1];
       while(i + k < n && j + k < n && this->chars[i +
           k] == this -> chars[j + k]) {
       lcp[rank[i]] = k;
   this->lcp = lcp;
   //create lcp_rmq segtree
   function<int(int, int)> fmodify = [](const int
        src, const int val) -> int{return val;};
   function<int(int, int)> fcombine = [](const int a,
        const int b) -> int{return min(a, b);};
   this->lcp_rmq = Segtree<int>(n, 0, 1e9, fmodify,
        fcombine);
   for(int i = 0; i < n; i++){
       this->lcp_rmq.t[n + i] = lcp[i];
   this->lcp_rmq.build();
void calc_leaf_cnt() {
   vector<int> ans(this->nodes.size(), 0);
   function<int(SuffixNode*)> dfs = [&ans,
        &dfs](SuffixNode* cur) -> int {
       int cnt = cur -> children.size() == 0;
       for(auto i = cur -> children.begin(); i != cur
            -> children.end(); i++){
          cnt += dfs(i -> second);
       ans[cur -> index] = cnt;
      return cnt:
   dfs(this -> nodes[0]);
   this->leaf_cnt = ans;
//do greedy dfs on the tree.
//Ordering can be changed by switching the comparator
    in the ordered map in the node struct.
void calc_suf_arr() {
   vector<int> ans(this->chars.size(), 0);
   int ind = 0:
   int n = this \rightarrow n:
   function < void (SuffixNode*, int) > dfs = [&ans,
        &dfs, &ind, &n](SuffixNode* cur, int dist) ->
        void {
       dist += cur -> len():
       if(cur -> children.size() == 0){
          ans[ind++] = n - dist;
       for(auto i = cur -> children.begin(); i != cur
            -> children.end(); i++){
          dfs(i -> second, dist);
   dfs(this->nodes[0], 0);
   this->suf_arr = ans;
   //compute mapping from indices to suffix array
   this->suf_to_suf_ind = vector<int>(n, 0);
   for(int i = 0; i < n; i++){
       this->suf_to_suf_ind[this->suf_arr[i]] = i;
```

```
struct SuffixState {
    SuffixNode* v;
    SuffixState(SuffixNode* v, int pos) : v{v},
        pos{pos} {}
};
SuffixState ptr = SuffixState(nullptr, 0);
//runs in amortized O(1) time
void add char(char c) {
    this->chars.push_back(c);
    this->tree_extend((int) this->chars.size() - 1);
void tree_extend(int pos) {
    if(pos == 0){
       nodes.push_back(new SuffixNode(0));
       ptr = SuffixState(nodes[0], 0);
    while(true) {
       SuffixState nptr = go(ptr, pos, pos + 1);
       if(nptr.v != nullptr) {
           ptr = nptr;
           return;
       SuffixNode* mid = split(ptr);
       SuffixNode* leaf = new SuffixNode(nodes.size(),
            pos, this -> n, mid);
       nodes.push_back(leaf);
       mid -> set_child(chars[pos], leaf);
       ptr.v = get_link(mid);
       ptr.pos = ptr.v -> len();
       if(mid == nodes[0]) {
           break;
SuffixState go(SuffixState st, int 1, int r) {
   while(l < r) {
       if(st.pos == st.v -> len()) {
           st = SuffixState(st.v ->
                get_child(chars[1]), 0);
           if(st.v == nullptr) {
              return st;
       }
       else |
           if(chars[st.v -> 1 + st.pos] != chars[1]) {
              return SuffixState(nullptr, -1);
           if(r - 1 < st.v \rightarrow len() - st.pos) {
              return SuffixState(st.v, st.pos + r -
           1 += st.v -> len() - st.pos;
           st.pos = st.v \rightarrow len();
    return st:
SuffixNode* split(SuffixState st) {
    if(st.pos == st.v -> len()) {
       return st.v;
    if(st.pos == 0){
       return st.v -> parent;
    SuffixNode* par = st.v -> parent;
    SuffixNode* new_node = new
        SuffixNode(nodes.size(), st.v -> 1, st.v -> 1
        + st.pos, par);
    nodes.push_back(new_node);
   par -> set_child(chars[st.v -> 1], new_node);
   new_node -> set_child(chars[st.v -> 1 + st.pos],
```

```
st.v);
st.v -> parent = new_node;
st.v -> 1 += st.pos;
return new_node;
}

SuffixNode* get_link(SuffixNode* v) {
    if(v -> link != nullptr) {
        return v -> link;
    }
    if(v -> parent == nullptr) {
        return nodes[0];
    }
    SuffixNode* to = get_link(v -> parent);
    v -> link = split(go(SuffixState(to, to -> len()),
        v -> 1 + (v -> parent == nodes[0]), v -> r));
    return v -> link;
};
```

## 5 Maths

 $11 \mod = 1e9 + 7;$ 

#### 5.1 Modular Arithmetic

REMEMBER: power function is power(11, 11), not pow(11, 11).

```
vector<ll> fac;
map<pair<ll, l1>, l1> nckdp;
11 add(ll a, ll b) {
    11 \text{ ret} = a + b:
    while(ret >= mod) {
       ret -= mod;
    return ret;
ll sub(ll a, ll b) {
    11 \text{ ans} = a - b;
    while (ans < 0) {
       ans += mod;
    return ans;
11 mul(ll a, ll b) {
    return (a * b) % mod;
ll power(ll a, ll b) {
    11 \text{ ans} = 1;
    11 p = a;
    while(b != 0){
       if(b \% 2 == 1){
            ans = mul(ans, p);
       p = mul(p, p);
       b /= 2:
    return ans:
ll divide(ll a, ll b){
    return mul(a, power(b, mod - 2));
11 gcd(ll a, ll b){
    if(b == 0){
       return a;
    return gcd(b, a % b);
void fac_init() {
    fac = vector<ll>(1e6, 1);
```

```
for(int i = 2; i < fac.size(); i++){</pre>
       fac[i] = mul(fac[i - 1], i);
ll nck(ll n, ll k) {
   if(nckdp.find({n, k}) != nckdp.end()) {
       return nckdp.find({n, k}) -> second;
   ll ans = divide(fac[n], mul(fac[k], fac[sub(n, k)]));
   nckdp.insert({{n, k}, ans});
   return ans;
//true if odd, false if even.
bool nck_parity(ll n, ll k) {
   return (n \& (n - k)) == 0;
11 catalan(ll n){
   return sub(nck(2 * n, n), nck(2 * n, n + 1));
//cantor pairing function, uniquely maps a pair of integers
     back to the set of integers.
ll cantor(ll a, ll b, ll m) {
   return ((a + b) * (a + b + 1) / 2 + b) \% m;
ll extended_euclidean(ll a, ll b, ll& x, ll& y) {
   x = 1, y = 0;
ll x1 = 0, y1 = 1, a1 = a, b1 = b;
while (b1) {
       11 q = a1 / b1;
       tie(x, x1) = make_tuple(x1, x - q * x1);
       tie(y, y1) = make_tuple(y1, y - q * y1);
       tie(a1, b1) = make_tuple(b1, a1 - q * b1);
   return a1;
//modular inverse of a for any mod m.
//if -1 is returned, then there is no solution.
ll mod_inv(ll a, ll m) {
   11 x, y;
   11 g = extended_euclidean(a, m, x, y);
   if (g != 1) {
       return -1:
   else {
       x = (x \% m + m) \% m;
       return x:
//only works when all modulo is coprime.
//if you want to do this with non-coprime modulos, then you
     need to factor all of the modulos,
//and resolve the factors independently; converting them back
//it is not guaranteed that there is a solution if the
     modulos are not coprime.
ll chinese_remainder_theorem(vector<ll>& modulo, vector<ll>&
     remainder) {
   if(modulo.size() != remainder.size()) {
       return -1;
   for(int i = 0: i < modulo.size(): i++){</pre>
       M *= modulo[i];
   11 \text{ solution} = 0;
   for(int i = 0; i < modulo.size(); i++){</pre>
       11 a_i = remainder[i];
       11 M_i = M / modulo[i];
       11 N_i = mod_inv(M_i, modulo[i]);
       solution = (solution + a_i * M_i % M * N_i) % M;
   return solution;
```

## 5.2 Primes, Factors, Divisors

```
vector<int> lp; //lowest prime factor
vector<int> pr; //prime list
void prime_sieve(int n) {
   l\bar{p} = vector < int > (n + 1);
   pr = vector<int>(0);
    for(int i = 2; i <= n; i++) {
       if(lp[i] == 0) {
          lp[i] = i;
           pr.push_back(i);
       for (int j = 0; i * pr[j] <= n; j++) {
           lp[i * pr[j]] = pr[j];
           if (pr[j] == lp[i]) {
       }
   }
vector<int> find_prime_factors(int val) {
   vector<int> factors(0);
   while(val != 1) {
       factors.push_back(lp[val]);
       val /= lp[val];
   return factors;
void find_divisors_helper(vector<int>& p, vector<int>& c, int
    ind, int val, vector<int>& ans) {
   if(ind == p.size()) {
       ans.push_back(val);
   for(int i = 0; i <= c[ind]; i++){</pre>
       find_divisors_helper(p, c, ind + 1, val, ans);
       val *= p[ind];
vector<int> find_divisors(int val) {
   vector<int> factors = find_prime_factors(val);
   map<int, int> m;
   vector<int> p(0);
   vector<int> c(0);
   for(int i = 0; i < factors.size(); i++){</pre>
       int next = factors[i];
       if(m.find(next) == m.end()) {
          p.push_back(next);
           c.push_back(0);
           m.insert({next, m.size()});
       int ind = m[next];
       c[ind] ++;
    vector<int> div(0);
   find_divisors_helper(p, c, 0, 1, div);
   return div:
```

#### 5.3 FFT

For our purposes, allows us to multiply two polynomials of length N in O(Nlog(N)) time. Useful when we can convert a problem into polynomial multiplication.

```
const double PI = acos(-1);
void fft(vector<complex<ld>> & a, bool invert) {
   int n = a.size():
   if (n == 1) {
       return:
    vector<complex<ld>> a0(n / 2), a1(n / 2);
   for (int i = 0; 2 * i < n; i++) {
       a0[i] = a[2*i];
       a1[i] = a[2*i+1];
   fft(a0, invert);
   fft(a1, invert);
   double ang = 2 * PI / n * (invert ? -1 : 1);
    complex<ld> w(1), wn(cos(ang), sin(ang));
   for (int i = 0; 2 * i < n; i++) {
       a[i] = a0[i] + w * a1[i];
       a[i + n/2] = a0[i] - w * a1[i];
       if (invert) {
          a[i] /= 2;
          a[i + n/2] /= 2;
       \bar{w} = wn:
//given two polynomials, returns the product.
//if the two polynomials sizes arent same or powers of 2,
     this handles that.
vector<int> fft_multiply(vector<int> const& a, vector<int>
    const& b) {
    vector<complex<ld>>> fa(a.begin(), a.end()), fb(b.begin(),
        b.end());
    int n = 1:
    while (n < a.size() + b.size()) {</pre>
       n < < 1;
   fa.resize(n);
   fb.resize(n):
   fft(fa, false);
   fft(fb, false);
   for (int i = 0; i < n; i++) {</pre>
       fa[i] *= fb[i]:
   fft(fa, true);
    vector<int> result(n);
   for (int i = 0; i < n; i++) {
       result[i] = round(fa[i].real());
    return result;
```

# 5.4 Geometry

```
ld pi = acos(-1);
ld epsilon = 1e-9;

struct vec2 {
    ld x, y;
    vec2() {this->x = 0; this->y = 0;}
    vec2(ld x, ld y) {this->x = x; this->y = y;}
};

vec2 add(vec2 a, vec2 b){
```

```
vec2 ret;
                                                                      vec2 v2 = sub(t2, t3);
                                                                                                                                     }
   ret.x = a.x + b.x;
                                                                      return abs(cross(v1, v2) / 2.0);
                                                                                                                                      //assuming that the density of the polygon is uniform, the
   ret.v = a.v + b.v;
   return ret;
                                                                                                                                          centroid is the center of mass.
                                                                   //returns the distance along the ray from ray_a to the
                                                                                                                                      //winding direction matters...
                                                                                                                                      vec2 polygon_centroid(vector<vec2>& poly) {
                                                                       nearest point on the circle.
vec2 sub(vec2 a, vec2 b) {
                                                                                                                                         vec2 c = vec2();
                                                                  ld rayCircleIntersect(vec2 ray_a, vec2 ray_b, vec2 center, ld
   vec2 ret:
                                                                                                                                          for(int i = 0; i < poly.size(); i++){</pre>
                                                                       radius) {
   ret.x = a.x - b.x:
                                                                                                                                             vec2 v0 = poly[i];
                                                                      vec2 ray_dir = normalize(sub(ray_b, ray_a));
   ret.y = a.y - b.y;
                                                                      vec2 to_center = sub(center, ray_a);
                                                                                                                                             vec2 v1 = poly[(i + 1) \% poly.size()];
   return ret:
                                                                      vec2 center_proj = add(ray_a, mul(ray_dir, dot(ray_dir,
                                                                                                                                             ld p = cross(v0, v1);
                                                                                                                                             c.x' += (v0.x + v1.x) * p;
                                                                           to_center)));
                                                                      ld center_proj_len = length(sub(center, center_proj));
                                                                                                                                             c.y += (v0.y + v1.y) * p;
ld cross(vec2 a, vec2 b) {
                                                                      //radius^2 = center_proj_len^2 + int_depth^2
   return a.x * b.y - a.y * b.x;
                                                                       //int_depth = sqrt(radius^2 - center_proj_len^2)
                                                                                                                                         ld area = polygon_area(poly);
                                                                      ld int_depth = sqrt(radius * radius - center_proj_len *
                                                                                                                                         c.x /= (6.0 * area);
ld dot(vec2 a, vec2 b) {
                                                                           center_proj_len);
                                                                                                                                         c.y /= (6.0 * area);
   return a.x * b.x + a.y * b.y;
                                                                                                                                         return c;
                                                                      return dot(ray_dir, to_center) - int_depth;
ld length(vec2 a) {
                                                                                                                                      //i believe this gives in CCW order, have to verify though.
                                                                   //sector area of circle
   return sqrt(a.x * a.x + a.y * a.y);
                                                                                                                                      vector<vec2> convex hull(vector<vec2> a, bool
                                                                   ld sector_area(ld theta, ld radius) {
                                                                                                                                           include_collinear = false) {
                                                                      return radius * radius * pi * ((theta) / (2.0 * pi));
                                                                                                                                          function<int(vec2, vec2, vec2)> orientation = [](vec2 a,
ld distance(vec2 a, vec2 b){
                                                                                                                                              vec2 b, vec2 c) -> int {
   return length(sub(a, b));
                                                                  ld chord_area(ld theta, ld radius) {
                                                                                                                                             1d v = a.x*(b.y-c.y)+b.x*(c.y-a.y)+c.x*(a.y-b.y);
                                                                      ld sector = sector_area(theta, radius);
                                                                                                                                             if (v < 0) return -1; // clockwise</pre>
                                                                      ld tri_area = radius * radius * cos(theta) * sin(theta);
                                                                                                                                             if (v > 0) return +1; // counter-clockwise
ld lerp(ld t0, ld t1, ld x0, ld x1, ld t) {
                                                                      return sector - tri area:
   ld slope = (x1 - x0) / (t1 - t0);
                                                                                                                                         }:
   return x0 + slope * (t - t0);
                                                                   //dist = distance from center
                                                                                                                                         function<bool(vec2, vec2, vec2)> collinear =
                                                                   ld chord_area_dist(ld dist, ld radius) {
                                                                                                                                              [&orientation](vec2 a, vec2 b, vec2 c) -> bool {
vec2 mul(vec2 a, ld s) {
                                                                      ld theta = acos(dist / radius);
                                                                                                                                             return orientation(a, b, c) == 0;
   a.x *= s;
                                                                      return chord_area(theta, radius);
   a.y *= s;
   return a:
                                                                                                                                         function<bool(vec2, vec2, vec2, bool)> cw =
    [&orientation](vec2 a, vec2 b, vec2 c, bool
                                                                   //length of chord
                                                                  ld chord_area_length(ld length, ld radius) {
                                                                                                                                              include_collinear) -> bool {
vec2 normalize(vec2 a){
                                                                                                                                             int o = orientation(a, b, c);
                                                                      ld theta = asin((length / 2.0) / radius):
   ld len = length(a);
                                                                      return chord_area(theta, radius);
   vec2 ret;
                                                                                                                                             return o < 0 || (include_collinear && o == 0);</pre>
   ret.x = a.x / len;
   ret.y = a.y / len;
                                                                   //given a point inside and outside a circle, find the point
                                                                                                                                          vec2 p0 = *min_element(a.begin(), a.end(), [](vec2 a,
   return ret;
                                                                       along the line that intersects the circle.
                                                                   vec2 find_circle_intersect(vec2 in, vec2 out, vec2 c_center,
                                                                                                                                             return make_pair(a.y, a.x) < make_pair(b.y, b.x);</pre>
//angle from the +x axis in range (-pi, pi)
                                                                       ld c_radius) {
                                                                       //just binary search :D
ld polar_angle(vec2 a) {
                                                                                                                                          sort(a.begin(), a.end(), [&p0, &orientation](const vec2&
   return atan2(a.y, a.x);
                                                                       //i think we can reduce this to some sort of quadratic.
                                                                                                                                              a, const vec2& b)
                                                                      ld low = 0:
                                                                                                                                             int o = orientation(p0, a, b);
                                                                      1d high = 1:
                                                                                                                                             if (o == 0)
//project a onto b
                                                                      ld len = length(sub(in, out));
                                                                                                                                                 return (p0.x-a.x)*(p0.x-a.x) +
vec2 project(vec2 a, vec2 b) {
                                                                      vec2 norm = normalize(sub(out, in));
                                                                                                                                                      (p0.y-a.y)*(p0.y-a.y)
                                                                      while(abs(high - low) > epsilon) {
   b = normalize(b);
                                                                                                                                                     <(p0.x-b.x)*(p0.x-b.x) + (p0.y-b.y)*(p0.y-b.y);
   ld proj_mag = dot(a, b);
                                                                          ld mid = (high + low)^{-1}/2.0;
                                                                                                                                             return o < 0:
   return mul(b, proj_mag);
                                                                          vec2 mid_pt = add(in, mul(norm, len * mid));
                                                                                                                                         }):
                                                                          ld mid_dist = length(sub(mid_pt, c_center));
                                                                                                                                         if (include_collinear) {
                                                                          if(mid_dist < c_radius) {</pre>
                                                                                                                                             int i = (int)a.size()-1;
                                                                              low = mid:
vec2 rotateCCW(vec2 a, ld theta) {
                                                                                                                                             while (i \ge 0 \&\& collinear(p0, a[i], a.back())) i--;
   vec2 ret(0, 0);
                                                                                                                                             reverse(a.begin()+i+1, a.end());
   ret.x = a.x * cos(theta) - a.v * sin(theta);
                                                                          else {
   ret.v = a.x * sin(theta) + a.v * cos(theta):
                                                                             high = mid;
   return ret:
                                                                                                                                         vector<vec2> st;
                                                                                                                                         for (int i = 0; i < (int)a.size(); i++) {</pre>
                                                                      return add(in. mul(norm. len * low));
                                                                                                                                             while (st.size() > 1 && !cw(st[st.size()-2],
//returns the coefficients s and t, where p1 + v1 * s = p2 +
                                                                                                                                                  st.back(), a[i], include_collinear))
                                                                                                                                                st.pop_back();
vector<ld> lineLineIntersect(vec2 p1, vec2 v1, vec2 p2, vec2
                                                                   //returns the area of the polygon.
                                                                                                                                             st.push_back(a[i]);
    v2) {
                                                                   //winding direction doesn't matter
   if(cross(v1, v2) == 0){
                                                                   //polygon can be self intersecting i think...
       return {};
                                                                   ld polygon_area(vector<vec2>& poly) {
                                                                                                                                         //make sure there are no duplicate vertices
                                                                                                                                         vector<vec2> ans(0);
   ld s = cross(sub(p2, p1), v2) / cross(v1, v2);
                                                                      for(int i = 0; i < poly.size(); i++){</pre>
                                                                                                                                         for(int i = 0; i < st.size(); i++){</pre>
   ld t = cross(sub(p1, p2), v1) / cross(v2, v1);
                                                                          vec2 v0 = poly[i];
                                                                                                                                             vec2 v0 = st[i];
vec2 v1 = st[(i + 1) % st.size()];
   return {s, t};
                                                                          vec2 v1 = poly[(i + 1) \% poly.size()];
                                                                          area += cross(v0, v1);
                                                                                                                                             if(v0.x == v1.x \&\& v0.y == v1.y) {
                                                                                                                                                 continue:
ld tri area(vec2 t1, vec2 t2, vec2 t3) {
                                                                      return abs(area / 2.0);
   vec2 v1 = sub(t1, t2);
```

```
ans.push_back(st[i]);
   return ans;
//checks if the area of the triangle is the same as the three
     triangle areas formed by drawing lines from pt to the
//i don't think triangle winding order matters
bool point_inside_triangle(vec2 pt, vec2 t0, vec2 t1, vec2
   1d a1 = abs(cross(sub(t1, t0), sub(t2, t0)));
   1d a2 = abs(cross(sub(t0, pt), sub(t1, pt))) +
        abs(cross(sub(t1, pt), sub(t2, pt))) +
        abs(cross(sub(t2, pt), sub(t0, pt)));
   return abs(a1 - a2) < epsilon;
//\text{runs in }O(n * \log(n)) \text{ time.}
//has to do O(n * log(n)) preprocessing, but after
    preprocessing can answer queries online in O(log(n))
vector<br/>bool> points_inside_convex_hull(vector<vec2>& pts,
    vector<vec2>& hull) {
   vector<bool> ans(pts.size(), false);
   //edge case
   if(hull.size() <= 2){</pre>
       return ans;
   //find point of hull that has minimum x coordinate
   //if multiple elements have same x, then minimum y.
   int pivot_ind = 0;
   for(int i = 1; i < hull.size(); i++){</pre>
       if(hull[i].x < hull[pivot_ind].x || (hull[i].x ==</pre>
            hull[pivot_ind].x && hull[i].y <
            hull[pivot_ind].y)) {
           pivot_ind = i;
   //sort all the remaining elements according to polar
        angle to the pivot
   vector<vec2> h_pts(0);
   vec2 pivot = hull[pivot_ind];
   for(int i = 0; i < hull.size(); i++){</pre>
       if(i != pivot_ind) {
           h_pts.push_back(hull[i]);
   }
   sort(h_pts.begin(), h_pts.end(), [&pivot](vec2& a, vec2&
        b) -> bool {
       return polar_angle(sub(a, pivot)) < polar_angle(sub(b,</pre>
            pivot));
   //for each point we want to check, compute it's polar
        angle, then binary search for the sector that should
         contain it
   for(int i = 0; i < pts.size(); i++){</pre>
       vec2 pt = pts[i];
       ld pt_ang = polar_angle(sub(pt, pivot));
int low = 0;
       int high = h_pts.size() - 2;
       int tri_ind = low;
       while(low <= high) {</pre>
           int mid = low + (high - low) / 2;
           if(polar_angle(sub(h_pts[mid], pivot)) <= pt_ang) {</pre>
               tri_ind = max(tri_ind, mid);
               low = mid + 1;
           else {
               high = mid - 1;
       ans[i] = point_inside_triangle(pt, pivot,
            h_pts[tri_ind], h_pts[tri_ind + 1]);
   return ans;
```

## 6 Misc

# 6.1 Hungarian Algorithm

Given J jobs and W workers ( $J \le W$ ), computes the minimum cost to assign each prefix of jobs to distinct workers. Input: a matrix of dimensions JxW such that  $\mathfrak{C}[\mathfrak{j}][\mathfrak{w}]$  is the cost to assign J-th job to W-th worker (possibly negative). Returns a vector of length J, with the J-th entry equaling the minimum cost to assign the first J+1 jobs to distinct workers

```
template <class T>
bool ckmin(T &a, const T &b) { return b < a ? a = b, 1 : 0: }
template <class T>
vector<T> hungarian(const vector<vector<T>> &C) {
   const int J = (int)size(C), W = (int)size(C[0]);
   // job[w] = job assigned to w-th worker, or -1 if no job
        assigneď
   // note: a W-th worker was added for convenience
   vector<int> job(W + 1, -1);
   vector<T> ys(J), yt(W + 1); // potentials
   // -yt[W] will equal the sum of all deltas
   vector<T> answers:
   const T inf = numeric_limits<T>::max();
   for (int j_cur = 0; j_cur < J; ++j_cur) { // assign</pre>
        j_cur-th job
       int w_cur = W;
       job[w_cur] = j_cur;
       // min reduced cost over edges from Z to worker w
       vector<T> min_to(W + 1, inf);
       vector<int> prv(W + 1, -1); // previous worker on
            alternating path
       vector<bool> in_Z(W + 1); // whether worker is in Z
       while (job[w_cur] != -1) { // runs at most j_cur + 1
           in_Z[w_cur] = true;
          const int j = job[w_cur];
T delta = inf;
           int w_next;
           for (int w = 0; w < W; ++w) {
              if (!in_Z[w]) {
                  if (ckmin(min_to[w], C[j][w] - ys[j] -
                       yt[w]))
                      prv[w] = w_cur;
                  if (ckmin(delta, min_to[w])) w_next = w;
           // delta will always be non-negative.
           // except possibly during the first time this loop
           // if any entries of C[j_cur] are negative
           for (int w = 0; w \le W; ++w) {
              if (in_Z[w]) ys[job[w]] += delta, yt[w] -=
               else min_to[w] -= delta;
           w_cur = w_next;
       // update assignments along alternating path
       for (int w; w_cur !=-1; w_cur = w) job[w_cur] = job[w
            = prv[w_cur]];
       answers.push_back(-yt[W]);
   return answers;
```

# 6.2 Fast Unordered Set / Map

```
#include <ext/pb_ds/assoc_container.hpp>
#include <ext/pb_ds/tree_policy.hpp>
using namespace __gnu_pbds;
struct chash {
    static uint64_t splitmix64(uint64_t x) {
        // http://xorshift.di.unimi.it/splitmix64.c
        x += 0x9e3779b97f4a7c15;
       x = (x \hat{ } (x >> 30)) * 0xbf58476d1ce4e5b9;

x = (x \hat{ } (x >> 27)) * 0x94d049bb133111eb;
        return x \hat{ } (x >> 31);
    size t operator()(uint64 t x) const {
        static const uint64_t FIXED_RANDOM =
             chrono::steady_clock::now().time_since_epoch().count()
        return splitmix64(x + FIXED_RANDOM);
};
template<typename T> using pb_set = gp_hash_table<T,</pre>
     null_type, chash>; // unordered_set but faster
template<typename T, typename U> using pb_map =
     gp_hash_table<T, U, chash>; // unordered_map but faster
```

## 6.3 Ordered Set / Multiset

#### 6.4 Vector Print