Cminus: A Language for Practice Implementation

1 Purpose

This document describes the Cminus programming language. It is intended to provide enough detail to allow implementation of a parser (context-free analyzer). This document contains information that is irrelevant to the some portions of each assignment in this course. Read the programming descriptions carefully to ensure that you understand the scope of each assignment.

2 Introduction

Cminus is a programming language designed for practice implementation. Cminus is an extremely simplified version of C in which one may perform simple integer calculations. Cminus is intended to be simple enough to implement in a single semester by any student willing to put in some effort. Each feature included in the language was added specifically to illustrate some problem that arises in the design and implementation of a very simple compiler.

Cminus supports two basic data types: *integer* and *float*. These types may be aggregated into a one dimensional array. The language is intended to be *strongly typed*; that is, the type of each expression should be determinable at compile time. Since there is no *boolean* data type, integers are used as logical values. Control structures in Cminus are limited. It has an *if* statement, a *while* statement and a *compound* statement. Procedure abstraction is also limited. One cannot declare a procedure inside another procedure.

3 Lexical Properties of Cminus

In this section, any sequence of characters denoted inside of a pair of "'s indicates the actual string literal found inside of the "'s.

- 1. In Cminus, blanks are significant.
- 2. In Cminus, all keywords are reserved; that is, the programmer cannot use a Cminus keyword as the name of a variable. The valid keywords are:

```
"else"
(ELSE)
\langle EXIT \rangle
                       "exit"
⟨FLOAT⟩
                       "float"
\langle IF \rangle
                      "if"
                      "int"
\langle INT \rangle
\langle READ \rangle
                       "read"
(RETURN)
                       "return"
(WHILE)
                       "while"
                       "write"
(WRITE)
```

(Note that Cminus is *case sensitive*, that is, the variable X differs from x. Thus, if is a keyword, but IF can be a variable name.)

3. The following special characters have meanings in a Cminus program.

```
\langle AND \rangle
                                 "&&"
⟨ASSIGN⟩
                                 "="
\langle CM \rangle
(DIVIDE)
                                "/"
\langle EQ \rangle
                                 "=="
\langle GE \rangle
                                 ">="
                                ">"
\langle GT \rangle
                                " { "
\langle LBR \rangle
                                " ["
\langle LBK \rangle
\langle LE \rangle
                                "<="
\langle LP \rangle
                              "("
\langle LT \rangle
                        → "<"</p>
                                 "-"
⟨MINUS⟩
                                "!="
\langle NE \rangle
\langle NOT \rangle
                                ...i...
                                "||"
\langle OR \rangle
(PLUS)
                                "+"
                               "}"
\langle RBR \rangle
\langle RBK \rangle
                               "]"
                                ")"
\langle RP \rangle
                                 ";"
\langle SC \rangle
\langle DQ \rangle
                                "\""
                                 "*"
\langle TIMES \rangle
```

- 4. Comments are delimited by the characters /* and */. A /* begins a comment; it is valid in no other context. A */ ends a comment; it cannot appear inside a comment. Comments may appear before or after any other token.
- 5. Identifiers are written with upper and lowercase letters and are defined as follows:

The implementor may restrict the length of identifiers so long as identifiers of at least 31 characters are legal.

6. Constants are defined as follows:

Special string constants are acceptable in write statements:

$$\langle STRING \rangle \rightarrow \langle DQ \rangle (\sim (\langle DQ \rangle))^* \langle DQ \rangle$$

4 Cminus Syntax

This section gives a syntactic description of Cminus. The sections following the grammar provide implementation notes on the various parts of the grammar.

The grammar, as stated, defines the language. It may require some massaging before implementation with any particular parser generator system.

4.1 BNF

The following grammar describes the context-free syntax of Cminus:

 $\langle Program \rangle \rightarrow \langle DeclList \rangle \langle Procedures \rangle$

| (Procedures)

 $\langle Procedure \rangle \rightarrow \langle Procedure Decl \rangle \langle Procedures \rangle$

(ProcedureDecl)

 $\langle Procedure Decl \rangle \rightarrow \langle Procedure Head \rangle \langle Procedure Body \rangle$

 $\langle ProcedureHead \rangle \rightarrow \langle FunctionDecl \rangle \langle DeclList \rangle$

| \langle FunctionDecl \rangle

 $\langle FunctionDecl \rangle \hspace{1cm} \rightarrow \hspace{1cm} \langle Type \rangle \hspace{1cm} \langle IDENTIFIER \rangle \hspace{1cm} \langle LP \rangle \hspace{1cm} \langle RP \rangle \hspace{1cm} \langle LBR \rangle$

 $\begin{array}{ccc} \langle \operatorname{ProcedureBody} \rangle & \to & \langle \operatorname{StatementList} \rangle & \langle \operatorname{RBR} \rangle \\ \langle \operatorname{DeclList} \rangle & \to & \langle \operatorname{Type} \rangle & \langle \operatorname{IdentifierList} \rangle & \langle \operatorname{SC} \rangle \\ \end{array}$

 $\langle DeclList \rangle \langle Type \rangle \langle IdentifierList \rangle \langle SC \rangle$

 $\langle IdentifierList \rangle \rightarrow \langle VarDecl \rangle$

 $\langle IdentifierList \rangle \langle CM \rangle \langle VarDecl \rangle$

 $\langle VarDecl \rangle \rightarrow \langle IDENTIFIER \rangle$

| $\langle IDENTIFIER \rangle \langle LBK \rangle \langle INTCON \rangle \langle RBK \rangle$

 $\langle \text{Type} \rangle \longrightarrow \langle \text{INT} \rangle$

 $\langle FLOAT \rangle$

 $\langle StatementList \rangle \rightarrow \langle Statement \rangle$

| \langle StatementList \rangle \langle Stmt \rangle

 $\langle Statement \rangle \rightarrow \langle Assignment \rangle$

⟨IfStatement⟩
⟨WhileStatement⟩
⟨IOStatement⟩
⟨ReturnStatement⟩
⟨ExitStatement⟩
⟨CompoundStatement⟩

 $\langle Assignment \rangle \rightarrow \langle Variable \rangle \langle ASSIGN \rangle \langle Expr \rangle \langle SC \rangle$

 $\langle \text{IfStatement} \rangle \hspace{1cm} \rightarrow \hspace{1cm} \langle \text{IF} \rangle \hspace{1cm} \langle \text{TestAndThen} \rangle \hspace{1cm} \langle \text{ELSE} \rangle \hspace{1cm} \langle \text{CompoundStatement} \rangle$

 $\langle IF \rangle \langle TestAndThen \rangle$

 $\langle \text{TestAndThen} \rangle \rightarrow \langle \text{Test} \rangle \langle \text{CompoundStatement} \rangle$

 $\langle \text{Test} \rangle \qquad \rightarrow \quad (\langle \text{Expr} \rangle)$

 $\langle \text{WhileStatement} \rangle \hspace{1cm} \rightarrow \hspace{1cm} \langle \text{WhileToken} \rangle \hspace{1cm} \langle \text{WhileExpr} \rangle \hspace{1cm} \langle \text{Statement} \rangle$

 $\begin{array}{lll} \langle \mathrm{WhileToken} \rangle & & \rightarrow & \langle \mathrm{WHILE} \rangle \\ \langle \mathrm{WhileExpr} \rangle & & \rightarrow & (\ \langle \mathrm{Expr} \rangle) \end{array}$

 $\langle IOStatement \rangle \rightarrow \langle READ \rangle \langle LP \rangle \langle Variable \rangle \langle RP \rangle \langle SC \rangle$

 $| \quad \langle WRITE \rangle \ \langle LP \rangle \ \langle Expr \rangle \ \langle RP \rangle \ \langle SC \rangle$

 $\langle WRITE \rangle \langle LP \rangle \langle StringConstant \rangle \langle RP \rangle \langle SC \rangle$

 $\langle ReturnStatement \rangle \rightarrow \langle RETURN \rangle \langle Expr \rangle \langle SC \rangle$

 $\langle \text{ExitStatement} \rangle \longrightarrow \langle \text{EXIT} \rangle \langle \text{SC} \rangle$

 $\langle CompoundStatement \rangle \quad \rightarrow \quad \langle LBR \rangle \; \langle StatementList \rangle \; \langle RBR \rangle$

```
\langle Expr \rangle
                                            \langle \text{Expr} \rangle \langle \text{AND} \rangle \langle \text{SimpleExpr} \rangle
                                            \langle \text{Expr} \rangle \langle \text{OR} \rangle \langle \text{SimpleExpr} \rangle
                                            \langle SimpleExpr \rangle
                                            \langle NOT \rangle \langle SimpleExpr \rangle
(SimpleExpr)
                                            \langle \text{SimpleExpr} \rangle \langle \text{EQ} \rangle \langle \text{AddExpr} \rangle
                                            \langle SimpleExpr \rangle \langle NE \rangle \langle AddExpr \rangle
                                            \langle SimpleExpr \rangle \langle LE \rangle \langle AddExpr \rangle
                                            \langle SimpleExpr \rangle \langle LT \rangle \langle AddExpr \rangle
                                            (SimpleExpr) (GE) (AddExpr)
                                            \langle SimpleExpr \rangle \langle GT \rangle \langle AddExpr \rangle
                                            \langle AddExpr \rangle
(AddExpr)
                                            ⟨AddExpr⟩ ⟨PLUS⟩ ⟨MulExpr⟩
                                            ⟨AddExpr⟩ ⟨MINUS⟩ ⟨MulExpr⟩
                                            (MulExpr)
                                            \langle \text{MulExpr} \rangle \langle \text{TIMES} \rangle \langle \text{Factor} \rangle
\langle MulExpr \rangle
                                            ⟨MulExpr⟩ ⟨DIVIDE⟩ ⟨Factor⟩
(Factor)
                                            (Variable)
                                            (Constant)
                                            \langle IDENTIFIER \rangle \langle LP \rangle \langle RP \rangle
                                            \langle LP \rangle \langle Expr \rangle \langle RP \rangle
                                            ⟨IDENTIFIER⟩
(Variable)
                                            (IDENTIFIER) (LBK) (Expr) (RBK)
\langle StringConstant \rangle
                                   \rightarrow
                                            \langle STRING \rangle
(Constant)
                                            ⟨INTCON⟩
                                            ⟨FLOATCON⟩
```

4.2 Section Notes

4.2.1 Declarations

Cminus has two standard types: int and float. Integers and floats occupy a single machine "word". This standard types may be composed into a structured array type. An identifier may represent one of four types of objects:

- 1. an integer variable
- 2. an integer array
- 3. a floating point variable
- 4. a floating point array

Identifiers are declared to be variables or arrays in a variable declaration. Only singly dimensioned arrays are permitted in Cminus. Indexing begins at 0 as in C and Java.

```
\it Example:
```

```
int x,y;
int a[15];
float vector[100];
```

4.2.2 Function Declarations

The semantics of function definition are simple. A function returns the value of the expression specified in the first return statement that it executes.

Example:

```
int max () {
   int a,b;
    read(a); read(b);
   if (a < b) {
      return b;
   }
   else {
      return a;
   }
}</pre>
```

4.2.3 Assignment Statement

The assignment statement requires that the *left hand side* (the $\langle \text{Variable} \rangle$ non-terminal) and *right hand side* (the $\langle \text{Expr} \rangle$ non-terminal) evaluate to have the same type. If they have different types, a context-sensitive error has occurred.

4.2.4 If Statement

The grammar for the **if-else** construct is written to eliminate the dangling else ambiguity. This is done by forcing a (CompoundStatement) in each of the then- and else-clauses. To evaluate an if statement, the expression is evaluated. For an integer value, Cminus defines 0 as *false*; any other value is equivalent to *true*.

Examples:

```
if (c == d) { d = a; }
if (b == 0) { b = 2*a; } else { b = b/2; }
```

4.2.5 While Statement

The while statement provides a simple mechanism for iteration. Cminus's while statement behaves like the while statement in many other languages; it executes the statement in the loop's body until the controlling expression becomes false. The controlling expression will be treated as a boolean value encoded into an int expression.

4.2.6 Input-Output Statements

Cminus provides two primitives for input and output. The read and write statements are intended to provide direct access to primitives implemented in the target abstract machine. Note that the write statement outputs a newline by default.

Examples:

```
read (x);
write (x+y);
write (''error'');
```

4.2.7 Return Statements

Cminus allows functions to return values. The type of the return value needs to be the same type as the type of the function, or it must be converted to that type.

```
Example:
    int f() {
        return 7;
}
```

4.2.8 Exit Statement

An exit statement in Cminus exits the program completely.

4.2.9 Expressions

Cminus expressions compute simple values of type int or float. For both integer and floating point numbers, arithmetic and comparison are defined.

Coercion: If an expression contains operands of only one type, evaluation is straight forward. When an operand contains mixed types, the situation is more complex. If an *Addop* or *Mulop* has an int operand and a float operand, the int operand should be converted to a float before the operation is performed.

Relational operators always produce an integer. Comparisons between integers and floats produce integer results. To perform the comparison, the integer is converted to a float. For the numbers, comparison is based on both sign and magnitude.

Note: in an assignment, the value of a numeric expression gets converted to match the type of the variable that appears on its left hand side.

Operator Precedence Operator precedences in Cminus are specified in the table below. Multiplication and division have the highest priority, && and || have the lowest.

Operator	Precedence
*, /	5
+, -	4
<, <=, =, >=, >, !=	3
!	2
&&,	1

Relational operators always produce an integer.

Booleans: Because Cminus has no booleans, relational expressions are defined to yield integer results. Thus, a relational expression of the form a == b is considered to be an arithmetic expression whose value is 0 if the relation does not hold and not 0 otherwise. Hence, both the if-else and while statements test integer values; the expression is considered *false* if it evaluates to 0 and to *true* if it evaluates to anything else. Consider the following example which tests for either of two conditions being true:

```
{
    read (a); read (b); read (c); read (d);
    if ((a == b) + (c < d)) { write (''error''); }
}</pre>
```

Note that relational expressions should be enclosed in parentheses because they have very low precedence. In the above example, the special operator || could have been used. In Cminus the operator || takes two integer operands. || produces the result 0 if both operands evaluate to 0; otherwise, it produces 1. The operator && evaluates to 1 if both operands are nonzero; otherwise it evaluates to 0. The unary logical operator! evaluates to 1 if its argument is zero and to 0 otherwise.

Function Invocation: Cminus uses parentheses to indicate invocation and square brackets to indicate subscripting of an array. This simplifies the grammar — many languages use parentheses for both purposes. Cminus has no parameters. Global variables are required to pass values to a function.

5 An Example Program

The following program represents a simple example program written in Cminus. This program successively reads pairs of integers from the input file and prints out their product.

```
int main() {
  int x, y;
    read(x); read(y);
  while ((x != 0) || (y != 0)) {
      write (x*y);
      read (x); read (y);
    }
    exit;
}
```