

//Important==Floyd's cycle approach

On Abstract data type (ADT) we only talk about features/operation (mathematical/logical)

possible use of array as dynamic list (my code school)

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## # Array and Linked List

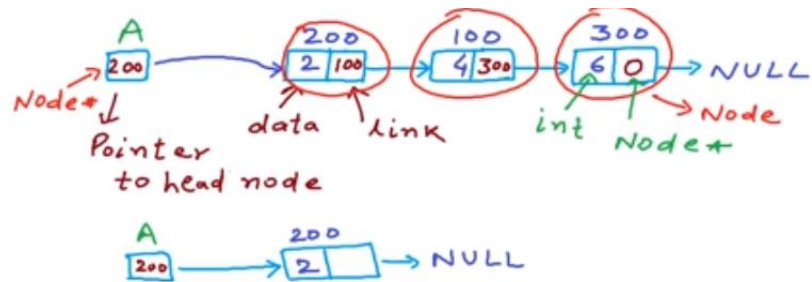
Array	Linked List
<ul style="list-style-type: none"> <li>Both store data of same type.</li> <li>Collection of elements of similar data type</li> <li>Elements are stored in contiguous memory location</li> </ul>	<ul style="list-style-type: none"> <li>Ordered collection of element of same type, connected to each other using pointers.</li> <li>Stored anywhere in the memory (address of new element is stored in the prev. node of linked list, hence forming a link b/w the two nodes)</li> </ul>
<p>Cost of accessing an element</p> <p><math>O(1)</math> - constant time</p> <p>base address = 200</p> <p>Address of <math>A[i] = 200 + 4(i)</math></p>	<p><math>O(n)</math></p> <p>head → [12] → [4] → [8] → Null</p> <p>↑ data ↑ link</p>
<p>Memory Usage</p> <p>fixed size</p> <p>used unused</p> <p>memory may not be available as one large block</p>	<p>No unused memory</p> <p>extra memory for pointer variable</p> <p>memory may available as multiple small blocks</p>
<p>Cost of Inserting/deleting</p> <ul style="list-style-type: none"> <li>at beginning - <math>O(n)</math></li> <li>at end - <math>O(1)</math></li> <li>if array is full, copying to new array - <math>O(n)</math></li> </ul>	<p><math>O(1)</math></p> <p><math>O(n)</math> (traversing whole array and then inserting at end, can be <math>O(1)</math> using stack)</p> <p><math>O(n)</math></p>
<p>average</p> <p>at <math>i^{th}</math> position - <math>O(n)</math></p>	<p><math>O(n)</math></p>
<p>case of use</p>	<p>X</p>

## LINKED LIST BASIC(CREATING A LINK LIST):

```

Struct Node
{
    int data;
    Node* link;
}
Node* A;
A = NULL;
Node* temp =
    (Node*)malloc(sizeof(Node))
(*temp).data = 2;
(*temp).link = NULL;
A = temp;

```

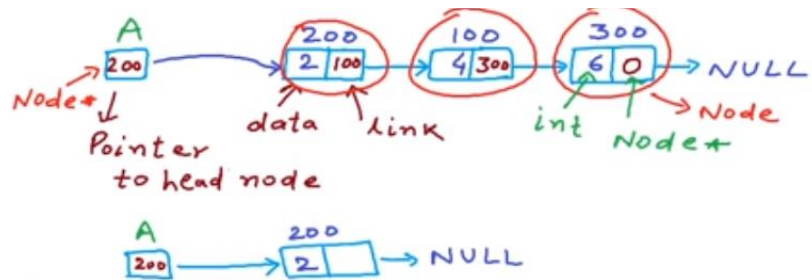


## NOW IN CPP WE DO IT LIKE THIS

```

Struct Node
{
    int data;
    Node* link;
}
Node* A;
A = NULL;
Node* temp = new Node();
temp->data = 2;
temp->link = NULL;
A = temp;

```




## ADDING ELEMENTS AT THE END:

Node\* A; ✓  
 A = NULL; ✓  
 Node\* temp = new Node();  
 temp->data = 2;  
 temp->link = NULL;  
 A = temp;  
 temp = new Node();  
 temp->data = 4;  
 temp->link = NULL;

Traversal

Node\* temp1 = A;  
 while (temp1->link != NULL)  
 {  
   temp1 = temp1->link;  
 }  
 temp1->link = temp;



```

graph LR
    A[A] --> Node1[200 | 2 | 100]
    Node1 --> Node2[100 | 4 | NULL]
  
```

### Reversing a Linked list:

```

// Reverse a nexted list
#include<stdio.h>
#include<stdlib.h>
struct Node {
    int data;
    struct Node* next;
};
struct Node* Reverse(struct Node* head) {
    struct Node *current,*prev,*next;
    current = head;
    prev = NULL;
    while(current != NULL)
    {
        next = current->next;
        current->next = prev;
        prev = current;
        current = next;
    }
    head = prev;
    return head;
}

int main()
{
    struct Node* head =NULL; // local variable
    head = Insert(head,2); // Insert: struct Node* Insert(struct Node* head,int data)
    head = Insert(head,4);
    head = Insert(head,6);
    head = Insert(head,8);
    Print(head);
    head = Reverse(head);
    Print(head);
}
  
```