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CPPCON2014

# EFFICIENCY & PERFORMANCE ALGORITHMS & DATA STRUCTURES

SO WHAT?  
WHO CARES?

“Software is getting slower more rapidly  
than hardware becomes faster.”

-NIKLAUS WIRTH, A PLEA FOR LEAN SOFTWARE

HE ALSO PREDICTED  
SOME ASPECTS OF  
THIS TALK...

found, systematic,  
and parametrizing  
treatment of basic  
and dynamic data  
structures, sorting,  
recursing algorithms,  
language structures,  
and compiling

PRENTICE-HALL  
SERIES IN  
AUTOMATIC  
COMPUTATION

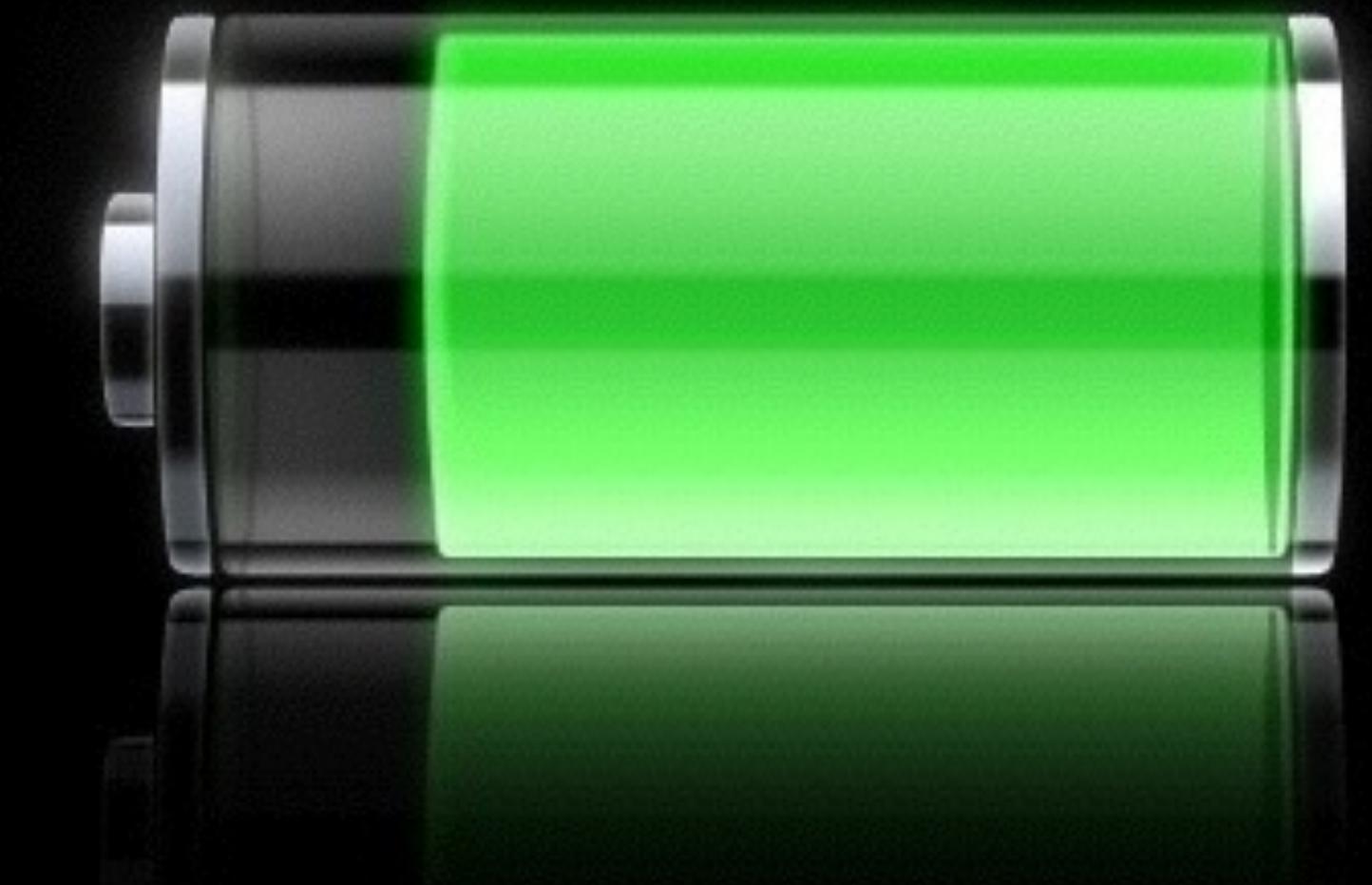
NIKLAUS WIRTH

Algorithms +  
Data  
Structures =  
Programs

EVEN MORE RELEVANT TODAY



SMALL, FAST — CHOOSE ONE







COMPUTE/WATT DOMINATES

RAW SPEED ISN'T EVERYTHING...

AFTER ALL, JAVA IS FASTER!



**TROLL**

Java is faster!

In terms of throughput in a long running system...

For specific applications...

On a well tuned JVM...

With near-optimal GC strategy selected...

SO WHY ARE WE USING C++?  
JAVA CON!!!

NOT SO FAST... (BA-DUM)

C++ DOESN'T GIVE YOU PERFORMANCE,  
IT GIVES YOU CONTROL OVER PERFORMANCE

# HIGH LEVEL THESIS:

- Efficiency through Algorithms
- Performance through Data Structures

# EFFICIENCY: HOW MUCH WORK IS REQUIRED BY A TASK

- Improving efficiency involves *doing less work*
- An efficient program is one which does the minimum (that we're aware of) amount of work to accomplish a given task

# PERFORMANCE: HOW QUICKLY A PROGRAM DOES ITS WORK

- Improving performance involves doing work faster
- But there is no such thing as “performant”, not even a word in the English language
- There is essentially no point at which a program cannot do work any faster... until you hit Bremermann’s limit...

WAT



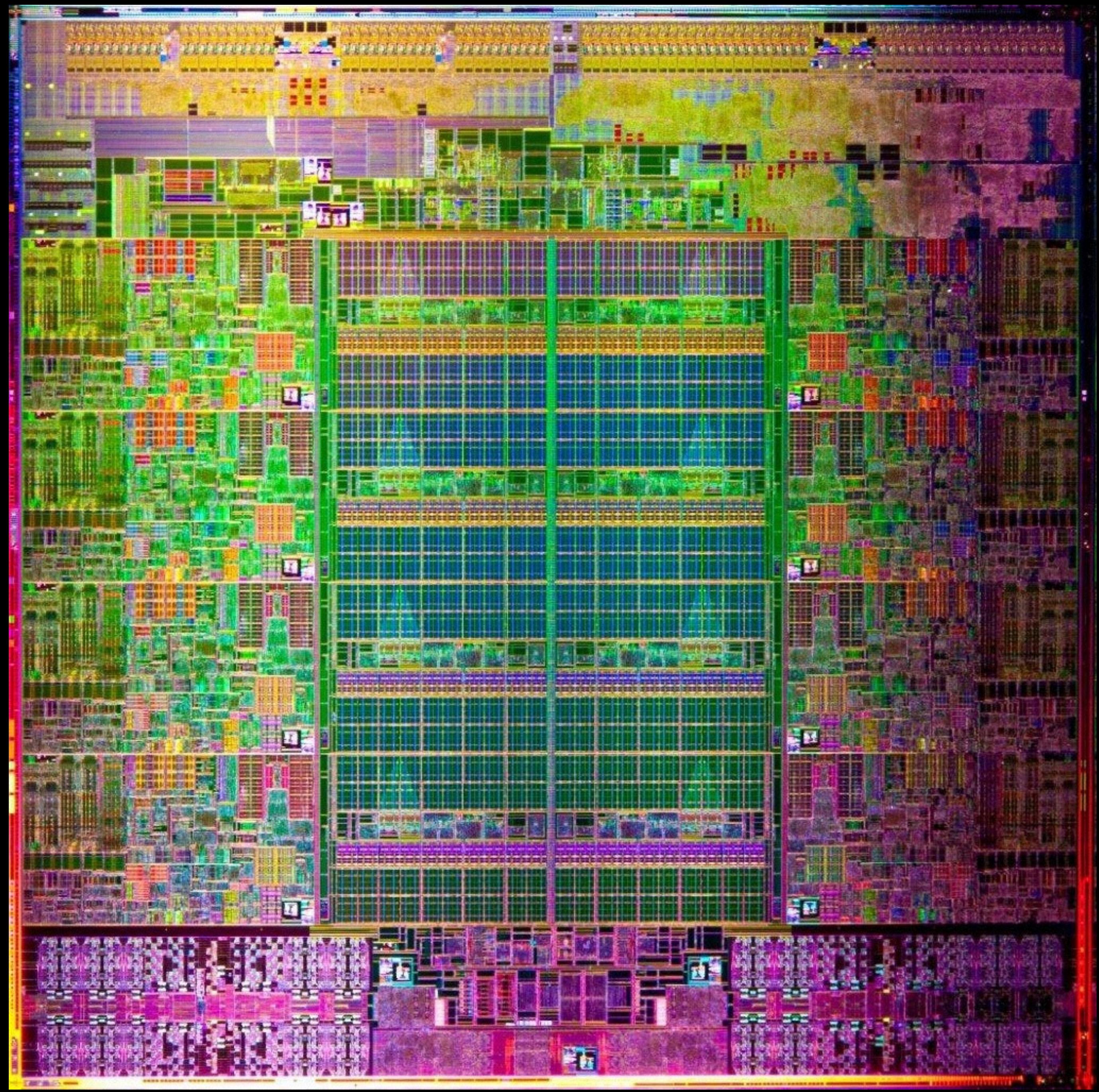
# PERFORMANCE: HOW QUICKLY A PROGRAM DOES ITS WORK

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# PERFORMANCE: HOW QUICKLY A PROGRAM DOES ITS WORK

- What does it mean to improve the performance of software?
- The software is going to run on a specific, real machine
- There is some theoretical limit on how quickly it can do work

PERFORMANCE:  
LIGHTING UP ALL OF THE TRANSISTORS



# ALL COMES BACK TO WATTS

- Every circuit not used on a processor is wasting power
- Don't reduce this to the absurd — it clearly doesn't make sense to use more parts of the CPU without improving performance!

EFFICIENCY → ALGORITHMS

# ALGORITHMS

- Complexity theory & analysis
- Common across higher level languages, etc.
- Very well understood by most (I hope?)
- Improving algorithmic efficiency requires finding a different way of solving the problem

# SUB-STRING SEARCHING

- Initially, you might have a basic  $O(n^2)$  algorithm
- Next, we have Knuth-Morris-Pratt (a table to skip)
- Finally, we have Boyer-Moore (use the end of the needle)

# DO LESS WORK BY NOT WASTING EFFORT

```
std::vector<X> f(int n) {  
    std::vector<X> result;  
    for (int i = 0; i < n; ++i)  
        result.push_back(X(...));  
    return result;  
}
```

# DO LESS WORK BY NOT WASTING EFFORT

```
std::vector<X> f(int n) {  
    std::vector<X> result;  
    result.reserve(n);  
    for (int i = 0; i < n; ++i)  
        result.push_back(X(...));  
    return result;  
}
```

# DO LESS WORK BY NOT WASTING EFFORT

```
X *getX(std::string key,  
         std::unordered_map<std::string,  
                           std::unique_ptr<X>> &cache) {  
    if (cache[key])  
        return cache[key].get();  
  
    cache[key] = std::make_unique<X>(...);  
    return cache[key].get();  
}
```

# DO LESS WORK BY NOT WASTING EFFORT

```
X *getX(std::string key,
          std::unordered_map<std::string,
                                std::unique_ptr<X>> &cache) {
    std::unique_ptr<X> &entry = cache[key];
    if (entry)
        return entry.get();

    entry = std::make_unique<X>(...);
    return entry.get();
}
```

ALWAYS DO LESS WORK

DESIGN APIs TO HELP

PERFORMANCE → DATA STRUCTURES

DISCONTIGUOUS DATA STRUCTURES ARE  
THE ROOT OF ALL (PERFORMANCE) EVIL

JUST SAY NO TO LINKED LISTS

# MODERN CPUS ARE TOO FAST

- 1,000,000,000 cycles per second
- 12+ cores per socket
- 3+ execution ports per core
- 36,000,000,000 instructions per second
- All of the time spent waiting for data (ok, 50%)

WE NEED FASTER MEMORY

# CPUS HAVE A HIERARCHICAL CACHE SYSTEM

One cycle on a 3 GHz processor	1	ns			
L1 cache reference	0.5	ns			
Branch mispredict	5	ns			
L2 cache reference	7	ns	14x L1 cache		
Mutex lock/unlock	25	ns			
Main memory reference	100	ns	20x L2, 200x L1		
Compress 1K bytes with Snappy	3,000	ns			
Send 1K bytes over 1 Gbps network	10,000	ns	0.01 ms		
Read 4K randomly from SSD*	150,000	ns	0.15 ms		
Read 1 MB sequentially from memory	250,000	ns	0.25 ms		
Round trip within same datacenter	500,000	ns	0.5 ms		
Read 1 MB sequentially from SSD*	1,000,000	ns	1 ms	4X memory	
Disk seek	10,000,000	ns	10 ms	20x datacenter RT	
Read 1 MB sequentially from disk	20,000,000	ns	20 ms	80x memory, 20X SSD	
Send packet CA->Netherlands->CA	150,000,000	ns	150 ms		

# STD::LIST

- Doubly-linked list
- Each node separately allocated
- All traversal operations chase pointers to totally new memory
- In most cases, every step is a cache miss
- Only use this when you rarely traverse the list, but very frequently update the list

JUST USE STD::VECTOR

# STACKS? QUEUES? MAPS?

- Just use `std::vector`. Really.
- It is already a perfectly good stack.
- If the queue can have a total upper bound and/or is short-lived, consider using a vector with an index into the front
- Grow the vector forever, chase the tail with the index

USING STD::MAP IS AN EXERCISE  
IN SLOWING DOWN CODE

# STD::MAP

- It's just a linked list, oriented as a binary tree
- All the downsides of linked lists
- Insertion and removal are **also** partial traversals
- Even with hinting, every rebalancing is a traversal
- But you can use std::unordered\_map, right?

NO.



# STD::UNORDERED\_MAP

- Essentially required to be implemented with buckets of key-value pairs for each hash entry.
- These buckets are... you guessed it... linked lists.
- You essentially always have some pointer chasing.

# A GOOD HASH TABLE DESIGN

- No buckets! Use open addressing into a table of the key-value pairs.
- Table stored as contiguous range of memory
- Use local probing on collisions to find an open slot in the same cache line (usually)
- Keep both key and values **small**

EFFICIENCY → ALGORITHMS  
PERFORMANCE → DATA STRUCTURES

EXCEPT THAT THIS IS ALL A LIE

# DATA STRUCTURES AND ALGORITHMS

- They're tightly coupled, see Wirth's book
  - You have to keep both factors in mind to balance between them
- Algorithms can also influence the data access pattern regardless of the data structure used
- Worse is better: bubble sort and cuckoo hashing

# CONCLUSION

- Both efficiency and performance matter, today more than ever
- C++ helps you control them
- Attend to your algorithms!
- Develop pervasive habits of using APIs in a way that avoids wasted work
- Use contiguous, dense, cache-oriented data structures
- Have fun writing crazy fast C++ code





THANK YOU!

QUESTIONS?