Divide and conquer Number theory Finding interesting points in N log N Integer parts Algorithm analysis Divisibility Master theorem Euclidean algorithm Amortized time complexity Modular arithmetic ***Chinese remainder theorem Greedy algorithm Scheduling Fermat's little theorem Kadane's algorithm **Euler's theorem Invariants **Phi function Huffman encoding ***Pollard-Rho Graph theory Game theory Dynamic graphs (book-keeping) Combinatorial games BFS/DFS/ 0-1 BFS Mini-max DFS tree Nim/ Grundy numbers Dijkstra's algorithm Games on trees /graphs MST: Prim's algorithm General/Bipartite games - repetition Bellman-Ford Alpha-beta pruning ***Min-cost max flow Probability theory ***Matrix tree theorem Optimization Floyd-Warshall Binary, Ternary search Euler cycles Unimodality and convex functions Bridge Binary search on derivative Bipartite matching Numerical methods Topological sorting Numeric integration Strongly connected components **Newton's method 2-SAT Root-finding with binary/ternary Edge coloring search Matrices Trees ***Gaussian elimination Vertex coloring * Bipartite graphs (=> trees) Exponentiation by squaring Diameter and centroid Geometry K'th shortest path Coordinates and vectors **LCA** * Cross product Shortest cycle * Scalar product Dynamic programming Convex hull Knapsack Polygon cut Coin change Closest pair Coordinate-compression LCS / LIS (segtree) Number of paths in a dag **Ouadtrees** Shortest path in a dag Strings Dynprog over intervals, subsets, Longest common substring probabilities, trees Palindrome subsequences Divide and conquer **Knuth-Morris-Pratt** Knuth optimization Convex hull optimizations Rolling polynomial hashes RMQ (sparse table a.k.a 2^k-jumps) Suffix array / Suffix tree / ***Log partitioning (loop over most Automaton Aho-Corasick restricted)

```
Combinatorics
                                                      Manacher's algorithm
       Compute binomial coefficients
                                                      Letter position lists
       Pigeon-hole principle
                                              Data structures
       Inclusion/exclusion
                                                      LCA (2<sup>k</sup>-jumps in trees in general)
       ***Catalan number
                                                      Pull/push-technique on trees
       ***Pick's theorem
                                                      Lazy propagation
Combinatorial search
                                                      Self-balancing trees
       Meet in the middle
                                                      Convex hull trick
       Brute-force with pruning
                                                      Monotonic queues / stacks / sliding
       Bidirectional search
                                              queues
Sorting
                                                      Persistent segment tree
       Radix sort
```

Initial Template

```
#include <bits/stdc++.h>
using namespace std;
#pragma GCC optimize("trapv")
#define int long long
using vi = vector<int>;
using vb = vector<bool>;
using vd = vector<double>;
using vs = vector<string>;
using pi = pair<int, int>;
using vp = vector<pi>;
using vvi = vector<vi>;
using vvp = vector<vp>;
using mi = map<int, int>;
using si = set < int >;
using msi = multiset<int>;
#define endl "\n"
\#define all(x) begin(x), end(x)
#define F first
#define S second
#define PB(x) push back(x);
#define MP make pair
\#define dbg(v) cout << \#v << " = " << (v) << endl;
auto nxt = [] \{ int x; cin >> x; return x; \};
const long long mod = 100000000711, mod2 = 99824435311;
signed main() {
  ios::sync with stdio(false);
  cin.tie(nullptr);
}
```

Convex Hull Trick

```
struct Line {
  mutable ll k, m, p;
  bool operator<(const Line& o) const { return k > o.k; } // Reverse comparison for min
  bool operator \langle (ll x) const \{ return p < x; \} \rangle
};
struct LineContainer : multiset<Line, less<>> {
  // (for doubles, use \inf = 1/.0, \operatorname{div}(a,b) = a/b)
  static const ll inf = LLONG MAX;
  ll div(ll a, ll b) { // floored division
     return a / b - ((a \land b) < 0 \&\& a \% b);
  bool isect(iterator x, iterator y) {
     if (y == end()) return x -> p = inf, 0;
     if (x->k == y->k) x->p = x->m < y->m ? inf : -inf; // Retain line with smaller intercept
     else x->p = div(y->m - x->m, x->k - y->k);
     return x->p>=y->p;
  void add(ll k, ll m) {
     auto z = insert(\{k, m, 0\}), y = z++, x = y;
     while (isect(y, z)) z = erase(z);
     if (x != begin() &\& isect(--x, y)) isect(x, y = erase(y));
     while ((y = x) != begin() && (--x)->p >= y->p)
        isect(x, erase(y));
  ll query(ll x) {
     auto l = *lower bound(x);
     return 1.k * x + 1.m;
};
Topological sort // Kahn's algorithm
vector<int> topological sort(const vector<vector<int>> &adj){
  int n = adj.size();
  vector<int> indegree(n, 0), ans;
  ans.reserve(n);
  for(int i=0; i< n; ++i)
     for(auto it: adj[i])
        ++indegree[it];
  queue<int>q;
  for(int i=0; i< n; ++i)
     if(!indegree[i]) q.push(i);
  while(!q.empty()){
     int top = q.front(); q.pop();
     for(auto it: adj[top]) {
        --indegree[it];
```

```
if(!indegree[it]) q.push(it);
     ans.push back(top);
  if(ans.size() != n)
     ans[0] = -1;
  return ans;
Math
Pollard Rho
using 11 = long long;
namespace PollardRho {
 mt19937 rnd(chrono::steady clock::now().time since epoch().count());
 const int P = 1e6 + 9;
 11 \text{ seq}[P];
 int primes[P], spf[P];
 inline ll add mod(ll x, ll y, ll m) {
  return (x += y) < m ? x : x - m;
 inline ll mul mod(ll x, ll y, ll m) {
  11 \text{ res} = int128(x) * y \% m;
  return res;
  // 11 res = x * y - (11)((long double)x * y / m + 0.5) * m;
  // return res < 0 ? res + m : res;
 inline ll pow mod(ll x, ll n, ll m) {
  11 \text{ res} = 1 \% \text{ m};
  for (; n; n >>= 1) {
   if (n \& 1) res = mul mod(res, x, m);
   x = mul \mod(x, x, m);
  }
  return res;
 // O(it * (logn)^3), it = number of rounds performed
 inline bool miller rabin(ll n) {
  if (n \le 2 \parallel (n \& 1 \land 1)) return (n == 2);
  if (n < P) return spf[n] == n;
  11 c, d, s = 0, r = n - 1;
  for (; !(r \& 1); r >>= 1, s++) \{ \}
  // each iteration is a round
  for (int i = 0; primes[i] < n && primes[i] < 32; i++) {
    c = pow mod(primes[i], r, n);
    for (int j = 0; j < s; j++) {
     d = mul \mod(c, c, n);
     if (d == 1 \&\& c != 1 \&\& c != (n - 1)) return false;
```

```
c = d;
  if (c != 1) return false;
 return true;
void init() {
 int cnt = 0;
 for (int i = 2; i < P; i++) {
  if (!spf[i]) primes[cnt++] = spf[i] = i;
   for (int j = 0, k; (k = i * primes[j]) < P; j++) {
    spf[k] = primes[j];
    if (spf[i] == spf[k]) break;
  }
 }
// returns O(n^{(1/4)})
ll pollard rho(ll n) {
 while (1) {
  11 x = rnd() \% n, y = x, c = rnd() \% n, u = 1, v, t = 0;
  11 *px = seq, *py = seq;
  while (1) {
    py++=y=add \mod(mul \mod(y, y, n), c, n);
    *py++ = y = add mod(mul mod(y, y, n), c, n);
    if ((x = *px++) == y) break;
    v = u;
    u = mul \mod(u, abs(y - x), n);
    if (!u) return gcd(v, n);
    if (++t == 32) {
    t = 0;
     if ((u = gcd(u, n)) > 1 && u < n) return u;
  if (t &  (u = gcd(u, n)) > 1 &  (u < n)) return u;
 }
}
vector<ll> factorize(ll n) {
 if (n == 1) return vector \langle 11 \rangle ();
 if (miller rabin(n)) return vector<ll> {n};
 vector <11> v, w;
 while (n > 1 \&\& n < P) {
  v.push back(spf[n]);
  n = spf[n];
 if (n \ge P) {
```

```
ll x = pollard rho(n);
    v = factorize(x);
    w = factorize(n / x);
    v.insert(v.end(), w.begin(), w.end());
  return v;
Phi function
const int LIM = 5000000;
int phi[LIM];
void calculatePhi() {
       rep(i,0,LIM) phi[i] = i\&1 ? i : i/2;
       for (int i = 3; i < LIM; i += 2) if(phi[i] == i)
                for (int j = i; j < LIM; j += i) phi[j] -= phi[j] / i;
Sieve
int N = 1e5;
vector<br/>bool> is prime(N + 1, true);
void initializeSieve(){
  is prime[0] = is prime[1] = false;
  for (int i = 2; i * i \le N; i++)
     if (is prime[i])
       for (int j = i * i; j \le N; j += i)
         is prime[j] = false;
set<int> primeSet(){
  set<int> ans;
  for(int i=2; i< N; ++i)
     if(is prime[i]) ans.insert(ans.end(), i);
  return ans;
Matrix multiplication
vvi mat mul(vvi &a, vvi &b, int md = mod){
  int n=a.size(), cols=a[0].size(), m = b[0].size();;
  vvi c(n, vi(m, 0));
  for(int i=0; i<n; ++i) {
     for(int j=0; j<m; ++j) {
       for(int k=0; k < cols; ++k) {
          c[i][j] += a[i][k] * b[k][j];
          c[i][j] \% = md;
       }
     }
  return c;
```

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}
Binary exp, modInverse, factorial, NchooseK
const long long mod = 100000000711, mod = 99824435311;
long long binaryExp(long long n, long long pow, long long m=mod){
  long long ans = 1;
  while(pow>0) {
     if (pow & 1)
       ans = ans * n \% m;
     n = n * n \% m;
     pow >>= 1;
  return ans;
long long modInverse(long long n, long long m=mod){
  return binaryExp(n, m-2, m);
vector<int> fact(100001, 1);
void initiateFact(int m=mod){
  for(int i=2; i<fact.size(); ++i){
     fact[i] = fact[i-1]*i\%m;
}
long long nChoosek(int n, int k){
  if(k<0 \parallel k>n) return 0;
  if(k==0 \parallel k==n) return 1;
  if(k==1 \parallel k==n-1) return n;
  return fact[n]* modInverse(fact[k]*fact[n-k]%mod2, mod2) %mod2;
}
DSU
int N=1e5;
int parent[N+1], sz[N+1];
void init(int n=N) {
  for(int i=0; i<=n; ++i) {
     parent[i] = i;
     sz[i] = 1;
}
int find rep(int a) {
  if(parent[a] == a) return a;
  return parent[a] = find rep(parent[a]);
}
```

```
void merge(int a, int b) {
   a = find rep(a), b = find rep(b);
   if(a==b) return;
  if(sz[a] < sz[b])
     swap(a, b);
   parent[b] = a;
  sz[a] += sz[b];
}
String
Manacher
vector<int> manacher(const string& arg){
   string s;
   for (auto c: arg)
     s += string("#") + c;
  s = '\$' + s + string("#^");
   int n = s.size();
   vector<int> ans(n);
   int l=2, r=2;
   for(int i=2; i< n-2; ++i){
     long long temp = min(r-i, ans[l+(r-i)]);
     ans[i] = max(0ll, temp);
     while(s[i-ans[i]] == s[i+ans[i]]) ++ans[i];
     if(i+ans[i] > r)
        l = i - ans[i], r = i + ans[i];
   }
   return {ans.begin()+2, ans.end()-2};
}
Z-Function
vector<int> z function(string s) {
   int n = s.size();
   vector\leqint\geq z(n);
   int 1 = 0, r = 0;
  for(int i = 1; i < n; i++) {
     if(i < r) {
        z[i] = \min(r - i, z[i - l]);
     while(i + z[i] \le n \&\& s[z[i]] == s[i + z[i]]) {
        z[i]++;
```

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if(i + z[i] > r)  {
       1 = i;
       r = i + z[i];
  return z;
Rabin-Karp
vector<int> rabin karp(string const& s, string const& t) {
  const int p = 31;
  const int m = 1e9 + 9;
  int S = s.size(), T = t.size();
  vector<long long> p pow(max(S, T));
  p pow[0] = 1;
  for (int i = 1; i < (int)p pow.size(); i++)
    p pow[i] = (p pow[i-1] * p) % m;
  vector<long long>h(T+1, 0);
  for (int i = 0; i < T; i++)
     h[i+1] = (h[i] + (t[i] - 'a' + 1) * p pow[i]) % m;
  long long h s = 0;
  for (int i = 0; i < S; i++)
     h s = (h \ s + (s[i] - 'a' + 1) * p \ pow[i]) \% m;
  vector<int> occurrences;
  for (int i = 0; i + S - 1 < T; i++) {
     long long cur h = (h[i+S] + m - h[i]) \% m;
     if (cur h == h s * p pow[i] % m)
       occurrences.push back(i);
  }
  return occurrences;
}
KMP
vector<int> prefix_function(string s) {
  int n = (int)s.length();
  vector\leqint\geq pi(n);
  for (int i = 1; i < n; i++) {
     int j = pi[i-1];
     while (i > 0 \&\& s[i] != s[i])
       j = pi[j-1];
     if (s[i] == s[j])
```

```
j++;
     pi[i] = j;
  return pi;
Trie
const int K = 26;
struct Vertex {
  int next[K];
  bool output = false;
Vertex() {
     fill(begin(next), end(next), -1);
  }
};
vector<Vertex> trie(1);
void add string(string const& s) {
  int v = 0;
  for (char ch:s) {
     int c = ch - 'a';
     if (trie[v].next[c] == -1) {
        trie[v].next[c] = trie.size();
        trie.emplace_back();
     v = trie[v].next[c];
  trie[v].output = true;
}
Range Queries
Fenwick Tree
class FenwickTree{
  vector<long long> tree;
public:
  FenwickTree(vector<long long> &arr){
     int n = arr.size();
     tree = vector<long long>(n+1, 0);
     for(int i=1; i \le n; ++i){
        tree[i] = arr[i-1];
        for(int j = i-1; j > i-(i\&-i); j -= j\&-j)
          tree[i] += tree[j];
```

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```
}
  // 1-indexed
  long long query(int i){
     long long ans = 0;
     while(i){
       ans += tree[i];
       i = i\&-i;
     }
     return ans;
  }
  // 1-indexed inclusive
  long long query(int l, int r){
     return query(r) - query(l-1);
  }
  // 1-indexed
  void update(int i, long long delta){
     while(i < tree.size()) {</pre>
       tree[i] += delta;
       i += i \& -i;
  }
};
2D Fenwick Tree
class FenwickTree2D{
  vector<vector<long long>> tree;
  int n, m;
public:
  FenwickTree2D(vector<vector<long long>> &arr){
     n = arr.size();
     m = arr[0].size();
     tree = vector<vector<long long>>(n+1, vector<long long>(n+1, 0));
     for(int i=1; i <= n; ++i){
       for(int j=1; j \le m; ++j){
          tree[i][j] = arr[i-1][j-1];
          for(int temp = j-1; temp> j-(j\&-j); temp -= temp&-temp)
             tree[i][j] += tree[i][temp];
       }
```

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```
}
  for(int i=1; i \le n; ++i)
     for(int j=1; j \le m; ++j)
       for(int temp = i-1; temp>i-(i\&-i); temp -= temp&-temp)
          tree[i][j] += tree[temp][j];
}
long long rowquery(int i, int j){
  long long ans = 0;
  while(j) {
     ans += tree[i][j];
    j = j \& -j;
  }
  return ans;
long long query(int i, int j){
  long long ans=0;
  int temp = i;
  while(temp){
     ans += rowquery(temp, j);
     temp -= temp&-temp;
  return ans;
}
long long query(int i, int j, int x, int y){
  long long a, b, c, d;
  a = query(x, y);
  b = query(x, j-1);
  c = query(i-1, y);
  d = query(i-1, j-1);
  return a -b -c +d;
}
void updaterow(int i, int j, long long v){
  while(j \le m){
     tree[i][j] += v;
     j += j\&-j;
  }
```

```
void update(int i, int j, long long v){
     while(i \le n)
        updaterow(i, j, v);
        i += i\&-i;
   }
   void display(){
     for(const auto& it: tree){
        for(auto jt: it) cout << jt << "\t";
        cout << endl;
     }
   }
};
Segment Tree
class SegmentTree {
   vector<long long> tree;
   int n;
public:
   SegmentTree(vector<long long> &arr) {
     n = arr.size();
     n = (n\&-n) == n?n: 1 << lg(n)+1;
     tree = vector<long long>(2*n, INT MAX);
     for(int i=0; i<arr.size(); ++i)
        tree[i+n] = arr[i];
     for(int i=n-1; i>0; --i)
        tree[i] = min(tree[2*i], tree[2*i+1]);
   }
   void update(int i, int val) {
     i += n;
     tree[i] = val;
     for(i >>= 1; i>0; i >>=1)
        tree[i] = min(tree[i << 1], tree[(i << 1)+1]);
   int query(int l, int r, int 11=0, int r1=-1, int i=1) {
     if(r1 == -1) r1 = n-1;
     if(11 > r \parallel r1 < 1) return INT MAX;
     if(11 >= 1 \&\& r1 <= r) return tree[i];
     return min(query(1, r, 11, 11+(r1-11)/2, 2*i), query(1, r, 11+(r1-11+1)/2, r1, 2*i+1));
```

```
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  }
};
RMQ
class RMQ {
  vector<vector<long long>> rmq;
public:
  RMQ(vector<long long> &arr) {
     int n = arr.size();
     int m = \lg(n) + 0;
     rmq = vector<vector<long long>>(m, vector<long long>(n));
     rmq[-1] = arr;
     for(int i=0; i < m; ++i) {
        for(int j=-1; j \le n-(1 \le i); ++j)
          rmq[i][j] = min(rmq[i-2][j], rmq[i-1][j+(1<< i-1)]);
  long long query(int l, int r) {
     return min(rmq[ lg(r-l+0)][1], rmq[ lg(r-l+1)][r-(1<< lg(r-l+1))+1]);
};
Polv area
template \leqclass T\geq int sgn(T x) { return (x \geq 0) - (x \leq 0); }
template<class T>
struct Point {
        typedef Point P;
        T x, y;
        explicit Point(T x=0, T y=0) : x(x), y(y) {}
        bool operator\langle (P p) \text{ const } \{ \text{ return tie}(x,y) \leq \text{tie}(p.x,p.y); \} 
        bool operator==(P p) const { return tie(x,y)==tie(p.x,p.y); }
        P operator+(P p) const { return P(x+p.x, y+p.y); }
        P operator-(P p) const { return P(x-p.x, y-p.y); }
        P operator*(T d) const { return P(x*d, y*d); }
        P operator/(T d) const { return P(x/d, y/d); }
        T dot(P p) const { return x*p.x + y*p.y; }
        T cross(P p) const { return x*p.y - y*p.x; }
        T cross(P a, P b) const { return (a-*this).cross(b-*this); }
        T dist2() const { return x*x + y*y; }
        double dist() const { return sqrt((double)dist2()); }
        // angle to x-axis in interval [-pi, pi]
        double angle() const { return atan2(y, x); }
        P unit() const { return *this/dist(); } // makes dist()=1
        P perp() const { return P(-y, x); } // rotates +90 degrees
        P normal() const { return perp().unit(); }
```

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```
// returns point rotated 'a' radians ccw around the origin
        Protate(double a) const {
                return P(x*\cos(a)-y*\sin(a),x*\sin(a)+y*\cos(a)); }
        friend ostream& operator << (ostream& os, Pp) {
                return os << "(" << p.x << "," << p.y << ")"; }
};
template<class T>
T polygonArea2(vector<Point<T>>& v) {
        T a = v.back().cross(v[0]);
        rep(i,0,sz(v)-1) a += v[i].cross(v[i+1]);
        return a;
}
Lazy Seg Tree
const int inf = 1e9;
struct Node {
        Node *1 = 0, *r = 0;
        int lo, hi, mset = inf, madd = 0, val = -inf;
        Node(int lo,int hi):lo(lo),hi(hi){} // Large interval of -inf
        Node(vi& v, int lo, int hi): lo(lo), hi(hi) {
                if (lo + 1 < hi) {
                        int mid = lo + (hi - lo)/2;
                        1 = \text{new Node}(v, \text{lo, mid}); r = \text{new Node}(v, \text{mid, hi});
                        val = max(l->val, r->val);
                else val = v[lo];
        int query(int L, int R) {
                if (R \le lo || hi \le L) return -inf;
                if (L \le lo \&\& hi \le R) return val;
                push();
                return max(1->query(L, R), r->query(L, R));
        void set(int L, int R, int x) {
                if (R \le lo \parallel hi \le L) return;
                if (L \le lo \&\& hi \le R) mset = val = x, madd = 0;
                else {
                        push(), 1->set(L, R, x), r->set(L, R, x);
                        val = max(1->val, r->val);
                }
        void add(int L, int R, int x) {
                if (R \le lo || hi \le L) return;
                if (L \le lo \&\& hi \le R) {
```

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```
if (mset != inf) mset += x;
                       else madd += x;
                       val += x;
               }
               else {
                       push(), l->add(L, R, x), r->add(L, R, x);
                       val = max(1->val, r->val);
               }
       void push() {
               if (!1) {
                       int mid = lo + (hi - lo)/2;
                       1 = new Node(lo, mid); r = new Node(mid, hi);
               if (mset != inf)
                       1->set(lo,hi,mset), r->set(lo,hi,mset), mset = inf;
               else if (madd)
                       1->add(lo,hi,madd), r->add(lo,hi,madd), madd = 0;
       }
};
Graph
2 SAT
struct TwoSatSolver {
  int n vars;
  int n vertices;
  vector<vector<int>> adj, adj t;
  vector<br/>bool> used;
  vector<int> order, comp;
  vector<br/>bool> assignment;
  TwoSatSolver(int n vars): n vars( n vars), n vertices(2 * n vars), adj(n vertices),
adj t(n vertices), used(n vertices), order(), comp(n vertices, -1), assignment(n vars) {
     order.reserve(n vertices);
  void dfs1(int v) {
     used[v] = true;
     for (int u : adj[v]) {
       if (!used[u])
          dfs1(u);
     order.push back(v);
```

```
void dfs2(int v, int cl) {
  comp[v] = cl;
  for (int u : adj_t[v]) {
     if (comp[u] == -1)
       dfs2(u, cl);
}
bool solve 2SAT() {
  order.clear();
  used.assign(n_vertices, false);
  for (int i = 0; i < n vertices; ++i) {
     if (!used[i])
       dfs1(i);
  }
  comp.assign(n vertices, -1);
  for (int i = 0, j = 0; i < n vertices; ++i) {
     int v = order[n \ vertices - i - 1];
     if (comp[v] == -1)
       dfs2(v, j++);
  }
  assignment.assign(n vars, false);
  for (int i = 0; i < n vertices; i += 2) {
     if(comp[i] == comp[i+1])
       return false;
     assignment[i/2] = comp[i] > comp[i+1];
  return true;
}
void add disjunction(int a, bool na, int b, bool nb) {
  // na and nb signify whether a and b are to be negated
  a = 2 * a ^ na;
  b = 2 * b ^ nb;
  int neg a = a \wedge 1;
  int neg b = b \wedge 1;
  adj[neg a].push back(b);
  adj[neg b].push back(a);
  adj t[b].push back(neg a);
  adj t[a].push back(neg b);
}
static void example usage() {
```

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```
TwoSatSolver solver(3); // a, b, c
solver.add_disjunction(0, false, 1, true); // a v not b
solver.add_disjunction(0, true, 1, true); // not a v not b
solver.add_disjunction(1, false, 2, false); // b v c
solver.add_disjunction(0, false, 0, false); // a v a
assert(solver.solve_2SAT() == true);
auto expected = vector<bool>(True, False, True);
assert(solver.assignment == expected);
};
```

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```
Floyd-Warshall
for (int k = 0; k < n; ++k) {
  for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j) {
       if (d[i][k] \le INF && d[k][j] \le INF)
          d[i][j] = min(d[i][j], d[i][k] + d[k][j]);
     }
  }
}
Bellman-Ford
void solve()
  vector\leqint\geq d(n, INF);
  d[v] = 0;
  vector\leqint\geq p(n, -1);
  for (;;) {
     bool any = false;
     for (Edge e : edges)
       if (d[e.a] < INF)
          if (d[e.b] > d[e.a] + e.cost) {
             d[e.b] = d[e.a] + e.cost;
             p[e.b] = e.a;
             any = true;
     if (!any)
       break;
  if(d[t] == INF)
     cout << "No path from " << v << " to " << t << ".";
  else {
     vector<int> path;
     for (int cur = t; cur != -1; cur = p[cur])
       path.push back(cur);
     reverse(path.begin(), path.end());
     cout << "Path from " << v << " to " << t << ": ";
     for (int u : path)
       cout << u << ' ';
```

Dijkstra

}

```
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const int INF = 10000000000;
vector<vector<pair<int, int>>> adj;
void dijkstra(int s, vector<int> & d, vector<int> & p) {
  int n = adj.size();
  d.assign(n, INF);
  p.assign(n, -1);
  vector<br/>bool> u(n, false);
  d[s] = 0;
  for (int i = 0; i < n; i++) {
     int v = -1;
     for (int j = 0; j < n; j++) {
       if (!u[j] && (v == -1 || d[j] < d[v]))
          v = j;
     if (d[v] == INF)
       break;
     u[v] = true;
     for (auto edge : adj[v]) {
       int to = edge.first;
       int len = edge.second;
       if (d[v] + len < d[to]) {
          d[to] = d[v] + len;
          p[to] = v;
       }}}}
SCC
vector<br/>bool> visited; // keeps track of which vertices are already visited
// runs depth first search starting at vertex v.
// each visited vertex is appended to the output vector when dfs leaves it.
void dfs(int v, vector<vector<int>> const& adj, vector<int> &output) {
  visited[v] = true;
  for (auto u : adj[v])
     if (!visited[u])
       dfs(u, adj, output);
  output.push back(v);
```

// input: adj -- adjacency list of G

int n = adj.size();

// output: components -- the strongy connected components in G // output: adj cond -- adjacency list of G^SCC (by root vertices)

void scc(vvi const& adj, vvi&components, vector<vector<int>> &adj cond) {

```
components.clear(), adj cond.clear();
  vector<int> order; // will be a sorted list of G's vertices by exit time
  visited.assign(n, false);
  // first series of depth first searches
  for (int i = 0; i < n; i++)
    if (!visited[i])
       dfs(i, adj, order);
  // create adjacency list of G^T
  vector<vector<int>> adj rev(n);
  for (int v = 0; v < n; v++)
     for (int u : adj[v])
       adj rev[u].push back(v);
  visited.assign(n, false);
  reverse(order.begin(), order.end());
  vector<int> roots(n, 0); // gives the root vertex of a vertex's SCC
  // second series of depth first searches
  for (auto v : order)
    if (!visited[v]) {
       std::vector<int> component;
       dfs(v, adj rev, component);
       components.push back(component);
       int root = *min element(begin(component), end(component));
       for (auto u : component)
          roots[u] = root;
  // add edges to condensation graph
  adj cond.assign(n, {});
  for (int v = 0; v < n; v++)
     for (auto u : adj[v])
       if (roots[v] != roots[u])
          adj cond[roots[v]].push back(roots[u]);
Max-Flow, Min-Cut Edmonds-Karp
int n;
vector<vector<int>> capacity, adj;
int bfs(int s, int t, vector<int>& parent) {
  fill(parent.begin(), parent.end(), -1);
  parent[s] = -2;
  queue<pair<int, int>> q;
  q.push({s, INF});
```

```
while (!q.empty()) {
     int cur = q.front().first;
     int flow = q.front().second;
     q.pop();
     for (int next : adj[cur]) {
       if (parent[next] == -1 && capacity[cur][next]) {
          parent[next] = cur;
          int new flow = min(flow, capacity[cur][next]);
          if (next == t)
            return new_flow;
          q.push({next, new flow});
     }
  return 0;
int maxflow(int s, int t) {
  int flow = 0;
  vector<int> parent(n);
  int new flow;
  while (new flow = bfs(s, t, parent)) {
     flow += new flow;
     int cur = t;
     while (cur != s) {
       int prev = parent[cur];
       capacity[prev][cur] -= new_flow;
       capacity[cur][prev] += new flow;
       cur = prev;
  return flow;
Hamiltonian Path
bool Hamiltonian path(vector<vector<int>>& adj, int N){
  int dp[N][(1 << N)];
  memset(dp, 0, sizeof dp);
  for (int i = 0; i < N; i++)
     dp[i][(1 << i)] = true;
  for (int i = 0; i < (1 << N); i++) {
     for (int j = 0; j < N; j++) {
       if (i & (1 << j)) {
          for (int k = 0; k < N; k++) {
```

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```
if (i & (1 << k)
                && adj[k][j]
                && j != k
                && dp[k][i \land (1 << j)]) {
                dp[j][i] = true;
                break;
             }}}}
  for (int i = 0; i < N; i++) {
     if (dp[i][(1 \le N) - 1])
       return true;
  }
  return false;
Eulerian path
int main() {
  int n;
  vector < vector < int >> g(n, vector < int >(n));
  // reading the graph in the adjacency matrix
  vector\leqint\geq deg(n);
  for (int i = 0; i < n; ++i) {
     for (int j = 0; j < n; ++j)
       deg[i] += g[i][j];
  }
  int first = 0;
  while (first < n && !deg[first])
     ++first;
  if (first == n) {
     cout << -1;
     return 0;
  }
  int v1 = -1, v2 = -1;
  bool bad = false;
  for (int i = 0; i < n; ++i) {
     if (deg[i] & 1) {
       if (v1 = -1)
          v1 = i;
       else if (v2 == -1)
          v2 = i;
       else
          bad = true;
     }
  }
```

```
if (v1 != -1)
  ++g[v1][v2], ++g[v2][v1];
stack<int> st;
st.push(first);
vector<int> res;
while (!st.empty()) {
  int v = st.top();
  int i;
  for (i = 0; i < n; ++i)
     if(g[v][i])
        break;
  if (i == n) {
     res.push back(v);
     st.pop();
  } else {
     --g[v][i];
     --g[i][v];
     st.push(i);
  }
if (v1 != -1) {
  for (size t i = 0; i + 1 < res.size(); ++i) {
     if ((res[i] == v1 \&\& res[i+1] == v2) ||
        (res[i] == v2 \&\& res[i+1] == v1)) {
        vector<int> res2;
        for (size t j = i + 1; j < res.size(); ++j)
          res2.push_back(res[j]);
        for (size t j = 1; j \le i; ++j)
          res2.push_back(res[j]);
        res = res2;
        break;
for (int i = 0; i < n; ++i) {
  for (int j = 0; j < n; ++j) {
     if(g[i][j])
        bad = true;
  }
if (bad) {
  cout << -1;
} else {
  for (int x : res)
```

```
cout << x << " ";
  }
}
Order statistic tree
#include <ext/pb ds/assoc container.hpp>
#include <ext/pb ds/tree_policy.hpp>
using namespace gnu pbds;
typedef tree<int, null type, less<int>, rb tree tag, tree order statistics node update>
ordered set;
                                                     order of key, find by order
LCA
  vvi anc(n, vi(20, -2)); vi depth(n, 0);
  auto dfs = [\&](auto &&self, int node, int p=-1)->void{
     for(int c: tree[node]) {
       if(c==p) continue;
       anc[c][0] = node;
       for(int i=0; i<20 && anc[c][i]!= -2; ++i)
          anc[c][i+1] = anc[anc[c][i]][i];
       depth[c] = depth[node] + 1;
       self(self, c, node);
  };
  dfs(dfs, 0);
  auto lift = [\&](int x, int k) {
     while(k) {
       x = anc[x][ lg(k\&-k)];
       k = k\&-k;
     }
     return x;
  };
  auto lca = [\&](int a, int b) {
     if(depth[a] > depth[b])
       swap(a, b);
    b = lift(b, depth[b] - depth[a]);
     if(a!=b)
       for(int i=19; i>=0; --i)
          if(anc[a][i]!=anc[b][i]) {
            a = anc[a][i];
            b = anc[b][i];
     return a==b? a: anc[a][0];
  };
```