

# Introduction to OpenMP

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DATA SCIENCE &  
ARTIFICIAL INTELLIGENCE



SCIENTIFIC &  
DATA-INTENSIVE COMPUTING

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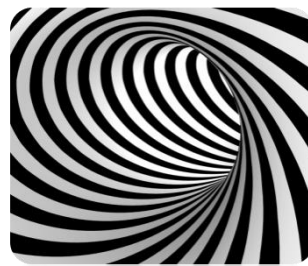
# OpenMP Outline



Introduction  
&  
Concept



Parallel  
Regions



Parallel  
Loops



NUMA  
AWARENESS

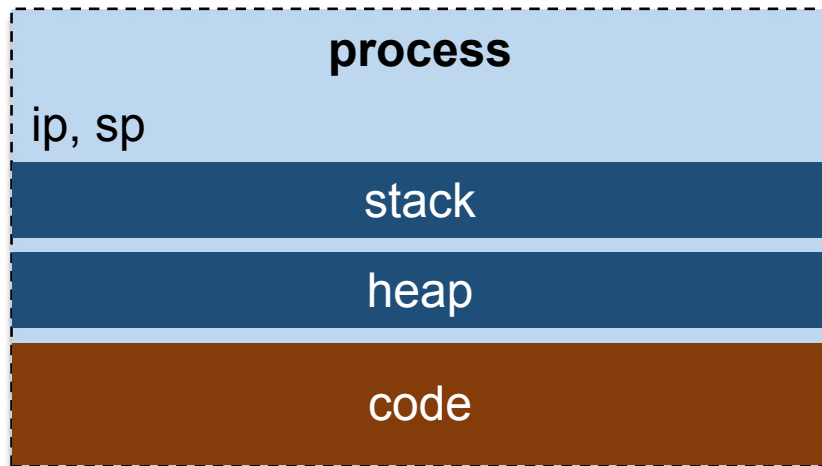


A **process** is an independent sequence of instructions *and* the ensemble of resources needed for their execution.

A program needs much more than just its binary code (i.e. the list of ops to be executed): it needs to access to a protected memory space and to access system resources (e.g. files and network).

A “process” is then a program that has been allocated with the necessary resources by the operating system.

There may be different **instances** of the same program as different, independent processes



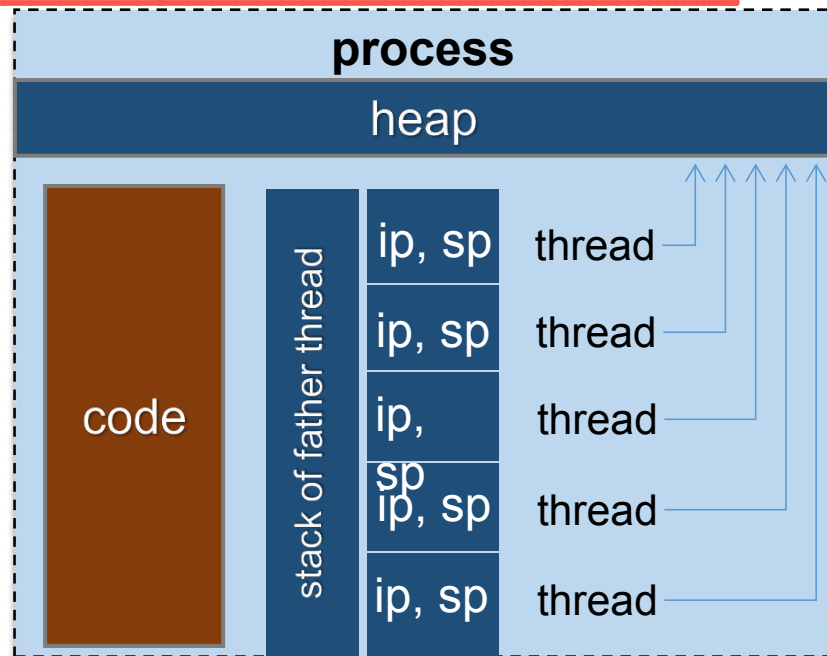


A **thread** is an independent instance of code execution *within* a process. There may be from one to many threads within the same process.

Each thread shares the same code, memory address space and resources than its father process.

While each thread has its own stack, ip and sp, the heap will be shared among threads, which then operate in *shared-memory*. threads also share the stack of the father thread.

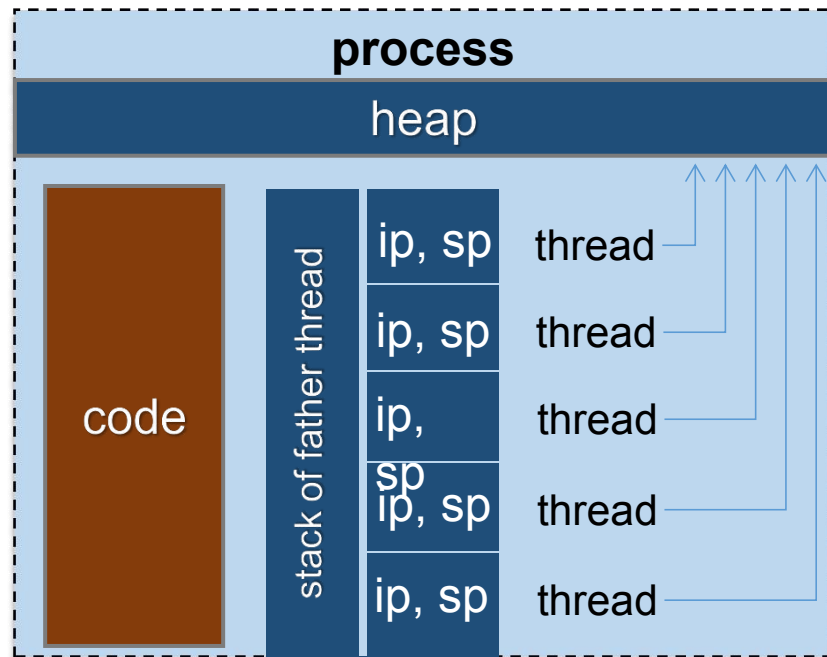
In geneal spawning threads inside a process is much less costly than creating processes.



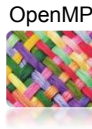


A thread can run either on the same computational units of its father process or on a different one.

A computational unit nowadays amounts to a **core**, either inside the same CPU (socket) on which the father process runs, or inside a sibling socket in the same NUMA region.





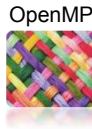


OpenMP is a standard API to enable shared-memory parallel programming:  
**Open** specifications for **MultiProcessing**

It allows to write multi-threaded programs with a standard behaviour through the usage of a set of compiler directives to be inserted in the source code:

- Pragmas '#' in C/C++
- Specially formatted comments in Fortran

Both fine- and coarse-grain parallelism are possible, from loop-level to explicit assignment to threads.



i.e

## Shared-Memory

vs

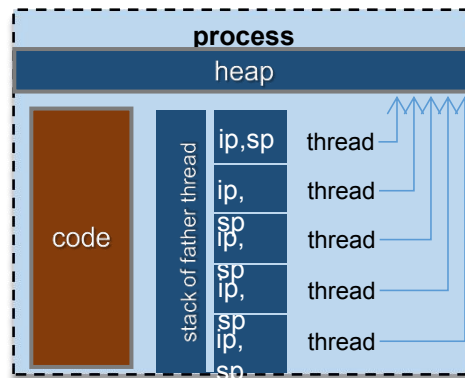
## Distributed-Memory



## Shared-Memory

(e.g. OpenMP)

A unique process that spawns a number of threads. There is a unique memory space that is accessible by all the threads



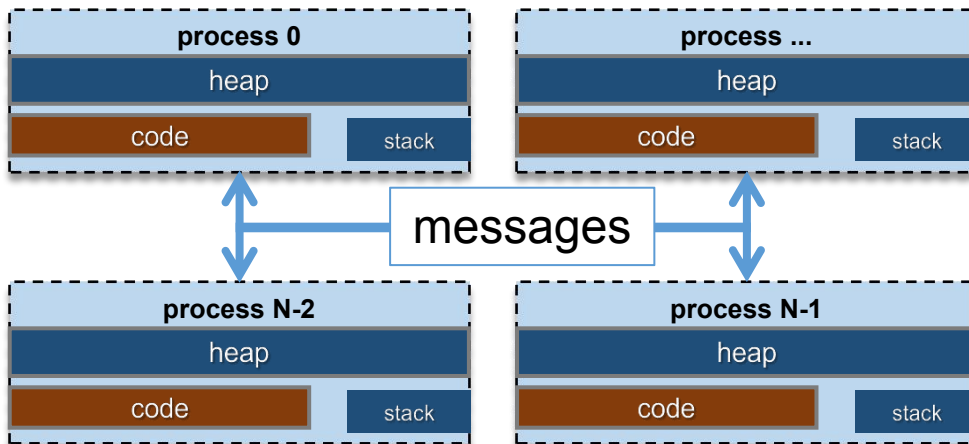
## Distributed-Memory

(e.g. MPI)

$N$  processes are created, each with its own copy of the code and its own memory space.

A process **can not** access the memory space of another process.

The processes communicate through *messages*.







## Shared-Memory

(e.g. OpenMP)

A unique process  
number of threads  
memory space that  
the threads

Actually MPI 3.0 introduced special tools to

- (1) allow shared-like accesses among tasks that run on cores that share the memory;
- (2) allow direct memory access to the memory of other MPI tasks in general, which is called Remote Memory Access

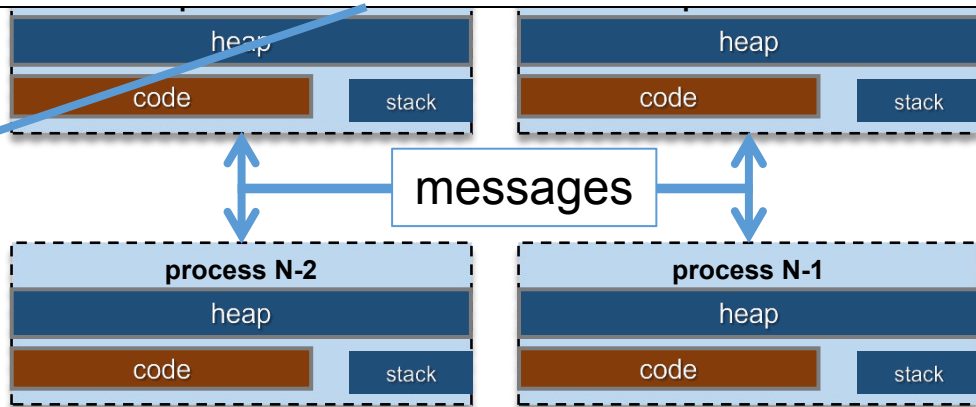
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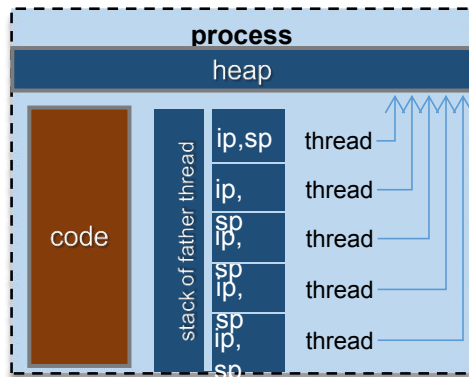




## Shared-Memory

(e.g. OpenMP)

A unique process that spawns a number of threads. There is a unique memory space that is accessible by all the threads

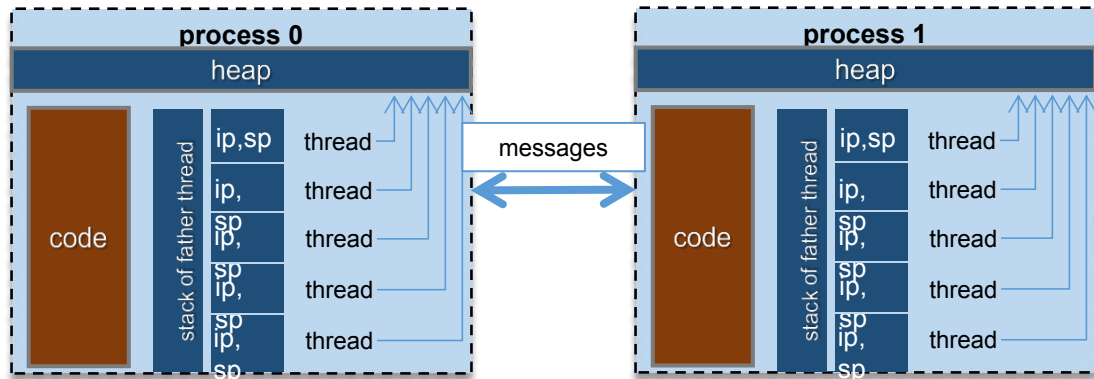


## Distributed-Memory

(e.g. MPI) + **Shared-Memory**

$N$  processes are created, each with its own copy of the code and its own memory space. Each process may spawn a number of threads as in shared-memory.

A process *can not* access the memory space of another process (nor any of its threads can). The processes communicate through *messages*.





# What is OpenMP



1997

OpenMP for Fortran, 1.0

1998

OpenMP for C/C++, 1.0

2000

OpenMP for Fortran, 2.0

2002

OpenMP for C/C++, 2.0

2005

OpenMP 2.5

2008

OpenMP 3.0

2013

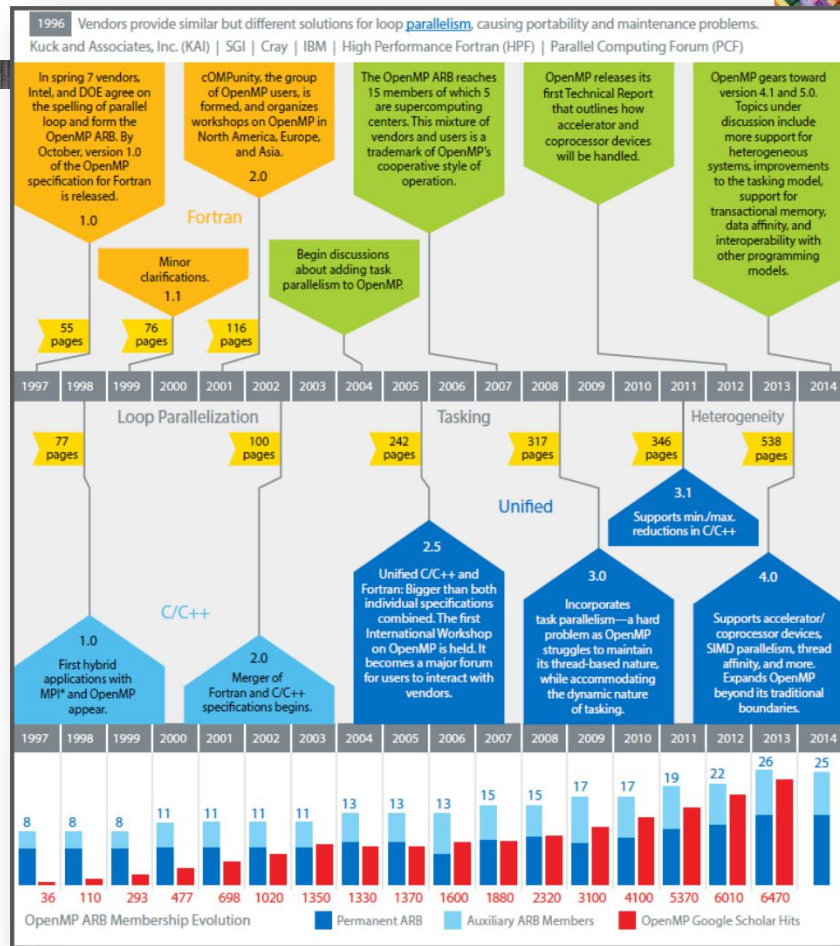
OpenMP 4.0

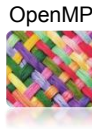
2015

OpenMP 4.5

2016

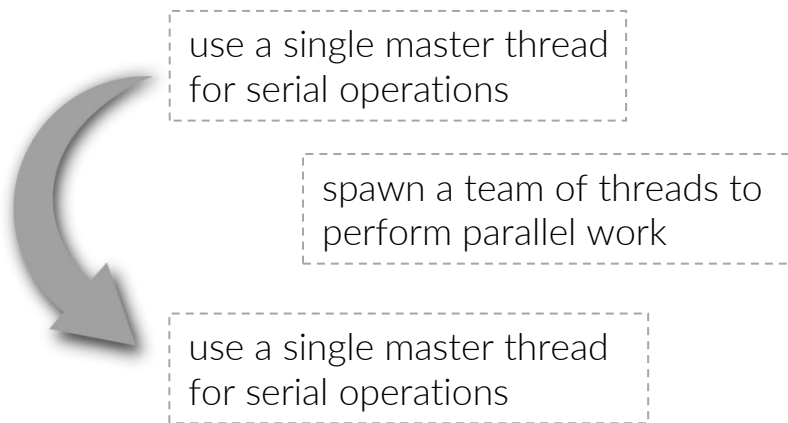
OpenMP 5.0

regular loops  
parallelizationfrom <http://openmp.org>Irregular parallelism  
task-based parallelismaccelerators / tasks+ /  
atomics / affinity /  
SIMD / user reductionoff-loading to  
accelerators;  
more on tasks and  
workshares, ...

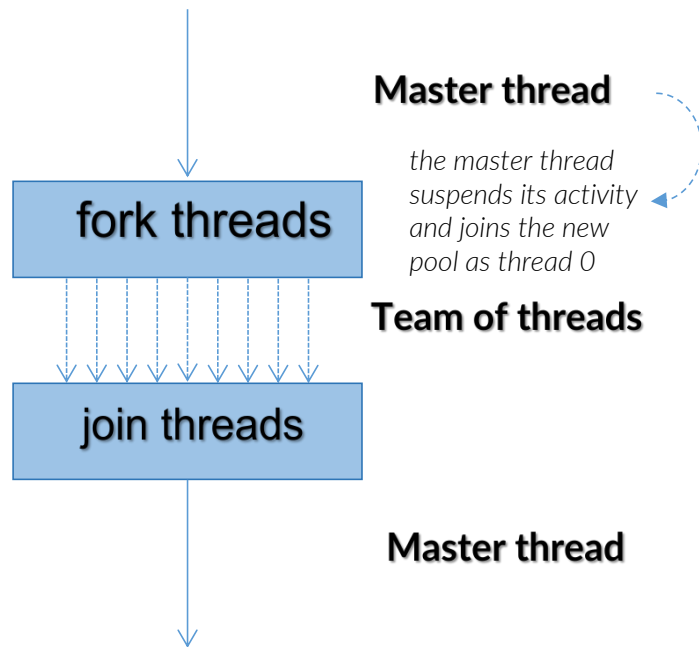


## Advantages of a directive-based approach

- **Abstraction**  
Subtleties of `pthread` and hardware-specific aspects are hidden. You can focus on data and workflow much more easily.
- **Efficiency.**  
The learning curve to achieve reasonable results is much shallower. The code's design is easier, the result/effort ratio is favourable with respect to `pthread`.
- **Incremental approach**  
No need to re-write your whole code. You start concentrating on some sections only, following a the suggestions from profiling.
- **Portability.**  
The compiler will take care of this for you. You still have to develop a design able to adapt to different topologies.
- **One source**  
Through conditional compilation, serial and parallel versions can easily coexist.



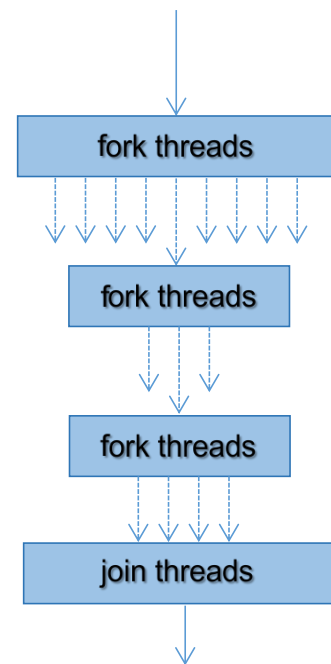
This is called “fork-join” model: a thread meets, at some point in its existence, a *directive* that activates the creation of a pool of children threads.



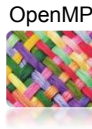


- Threads access and modify shared memory regions
  - explicit or implicit synchronization protect against race conditions
  - there is no concept like explicit “message-passing”
  - loop-carried dependencies hamper any parallel speedup
  - shared-variable attributes are vital to reduce or avoid race conditions or the need for synchronization
- Each thread performs its part of parallel work in a separate space and stack that are not visible to other threads or outside the parallel region
- Nested parallelism is explicitly permitted
- The number of threads can be dynamically changed before a parallel region

## NESTED PARALLELISM







An OpenMP directive is a specially-formatted pragma for C/C++ and comment for FORTRAN codes.

Most of the directives apply to *structured code block*, i.e. a block with a single input and a single output points and no branch within it.

The directives allows to

- create team of threads for parallel execution
- manage the sharing of workload among threads
- specify which memory regions (i.e. variables) are shared and which are private to each threads
- drive the update of shared memory regions
- synchronize threads and determine atomic/exclusive operations

## DECLARE PARALLEL REGION

```
!$OMP PARALLEL
```

```
...
```

```
!$OMP END PARALLEL
```

```
#pragma omp parallel  
{  
    ...  
}
```



As we have seen in the previous slide, the lexical scope of structured blocks defines the *static* extent of an OpenMP parallel region.

What happens if a function is called from a parallel region?

```
#pragma omp parallel
{
    double *array;
    int N;
    ...
    sum = foo(array, N);
    ...
}
```

} static  
extent



As we have seen in the previous slide, the lexical scope of structured blocks defines the *static* extent of an OpenMP parallel region.

Every function call from within a parallel region determines the creation of a *dynamic* extent to which the same directives apply.

**The dynamic extent includes the original static extent and all the instructions and further calls along the call tree.**

```
#pragma omp parallel
{
    double *array;
    int N;
    ...
    sum = foo(array, N);
    ...
}
```

static  
extent

```
double foo( double *A, int N )
{
    double mysum = 0;

    for ( int ii = 0; ii < N; ii++ )
        mysum += A[ii];
    return sum;
}
```

dynamic  
extent



As we have seen in the previous slide, the lexical scope of structured blocks defines the *static* extent of an OpenMP parallel region.

Every function call from within a parallel region determines the creation of a *dynamic* extent to which the same directives apply.

**The dynamic extent includes the original static extent and all the instructions and further calls along the call tree.**

The functions called in the dynamic extent can contain additional OpenMP directives.

```
#pragma omp parallel
{
    double *array;
    int N;
    ...
    sum = foo(array, N);
    ...
}
```

static  
extent

```
double foo( double *A, int N )
{
    double mysum = 0;
    #pragma parallel for reduction(+:sum)
    for ( int ii = 0; ii < N; ii++ )
        mysum += A[ii];
    return sum;
}
```

dynamic  
extent

“orphan”  
directive



As we have seen in the previous slide, the lexical scope of structured blocks defines the *static* extent of an OpenMP parallel region.

Every function call from within a parallel region determines the creation of a *dynamic* extent to which the same directives apply.

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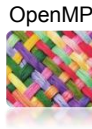
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    double *array;
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    sum = foo(array, N);
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}
```

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{
    double mysum = 0;
    #pragma parallel for reduction(+:sum)
    for ( int ii = 0; ii < N; ii++ )
        mysum += A[ii];
    return sum;
}
```

static  
extent

These will  
be thread-  
specific

dynamic  
extent



OpenMP is made of 3 components:

## 1. **Compiler directives**

give indication to the compiler about how to manage threads internals

## 2. **Run-time libraries** linked by the compiler

## 3. **Environment variables**

set by the user, determine the behaviour of the omp library; for instance, the number of threads to be spawned or the requirements about the thread-cores-memory affinity





directives for  
parallel regions

`#pragma omp parallel directive`

work-sharing  
constructs

`#pragma omp for`  
`#pragma omp section`  
`#pragma omp task`

data attributes

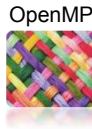
`shared, private, firstprivate, lastprivate,`  
`copyin, copyprivate, threadprivate`  
clauses

Synchronization

`critical, atomic, barrier, ordered, flush,`  
`nowait`

run-time

Lot of run-time functions and utils  
`omp_set_num_threads(),`  
`omp_get_thread_num(),`  
`omp_get_num_threads(), ...`



By default, when the compiler is instructed to activate the processing of OpenMP directives,

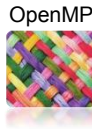
```
gcc -fopenmp ...  
icc -fopenmp ...  
pgcc -mp ...
```

it defines a macro that let you to conditionally compile sections of the code:

```
#ifdef _OPENMP  
my_thread_id = omp_some_function();  
#endif  
...
```



omp\_versions.c : with this example, you see how to determine what version of OpenMP standard is supported by your compiler



What the `_OPENMP` is useful for?

To write code that works as well also without OpenMP.

That helps you in assessing the correctness and portability of your code (mostly if you are writing an hybrid code, for instance MPI+OpenMP).

that's all, have fun



“So long  
and thanks  
for all the fish”