

Advanced Programming

INFO135

Lecture 4: **Sorting algorithms** in Python (part2)

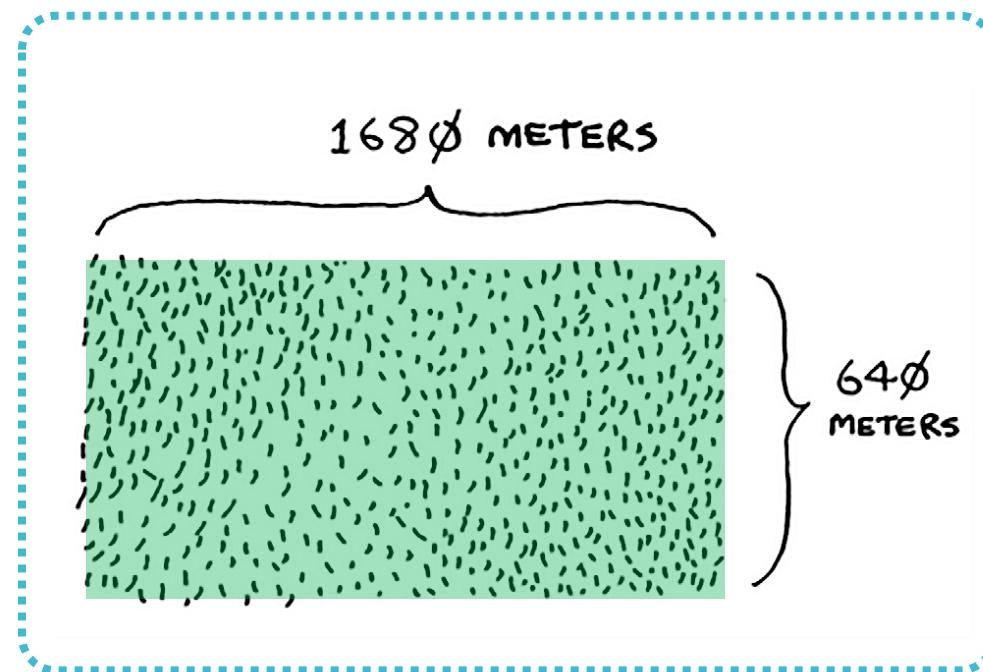
Mehdi Elahi

University of Bergen (UiB)

Divide & Conquer

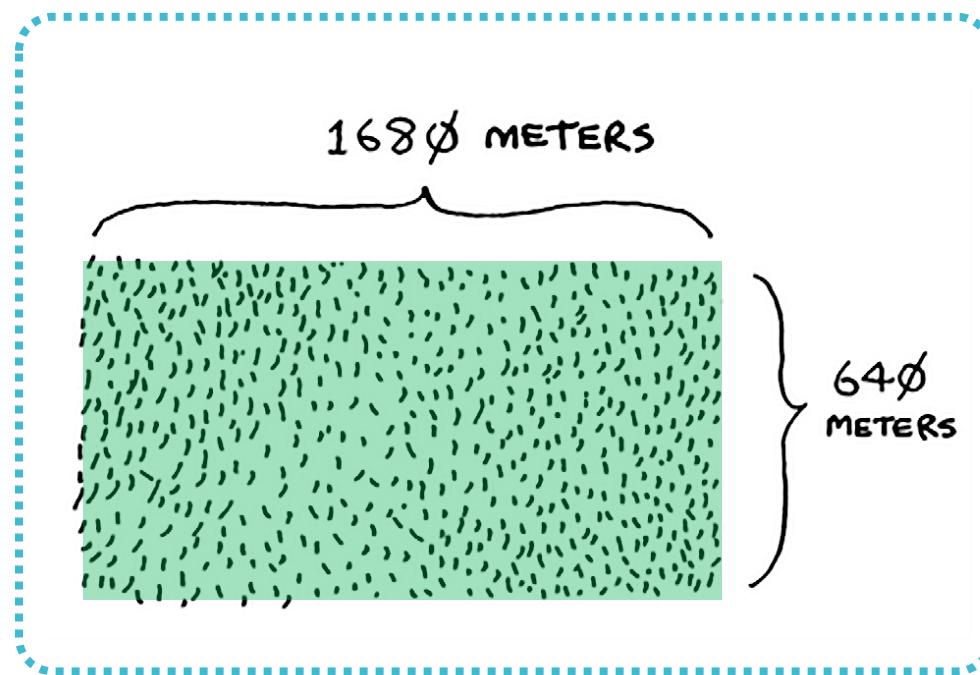
❖ Divide & Conquer (D&C) is a strategy to solve problems.

- ❖ Example: Suppose you have a plot of land.
- ❖ The size of your land is **1680 x 640 (m)**



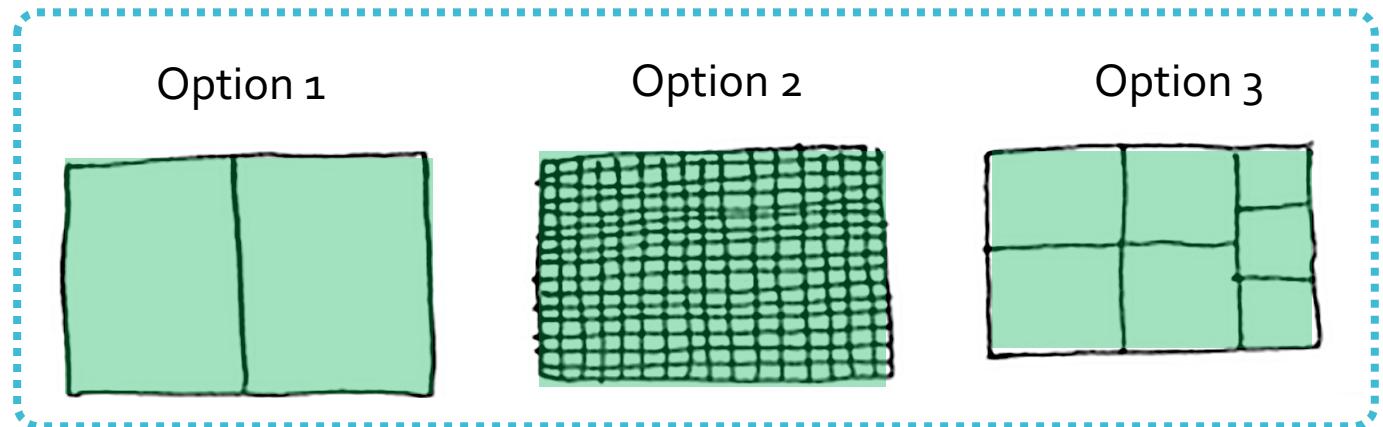
Divide & Conquer

- ❖ You want to divide this farm evenly into **square plots**.
- ❖ Plots should be **as big as possible**.



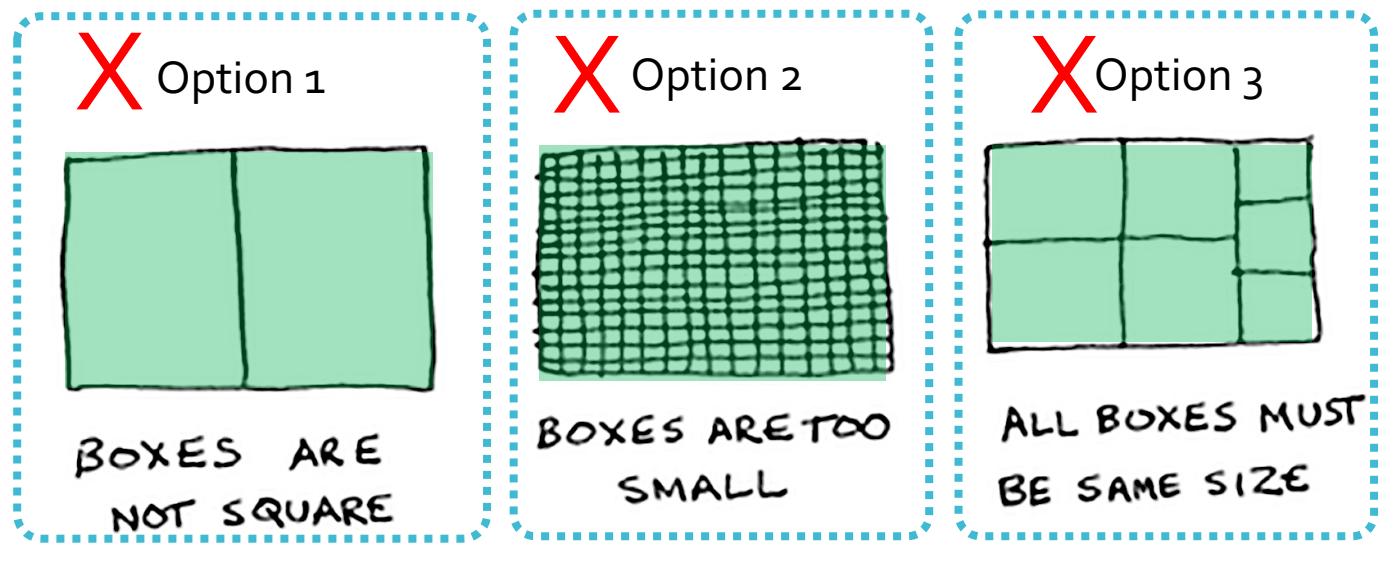
Divide & Conquer

- ❖ We can explore some of the **splitting** methods?
- ❖ These can be options:



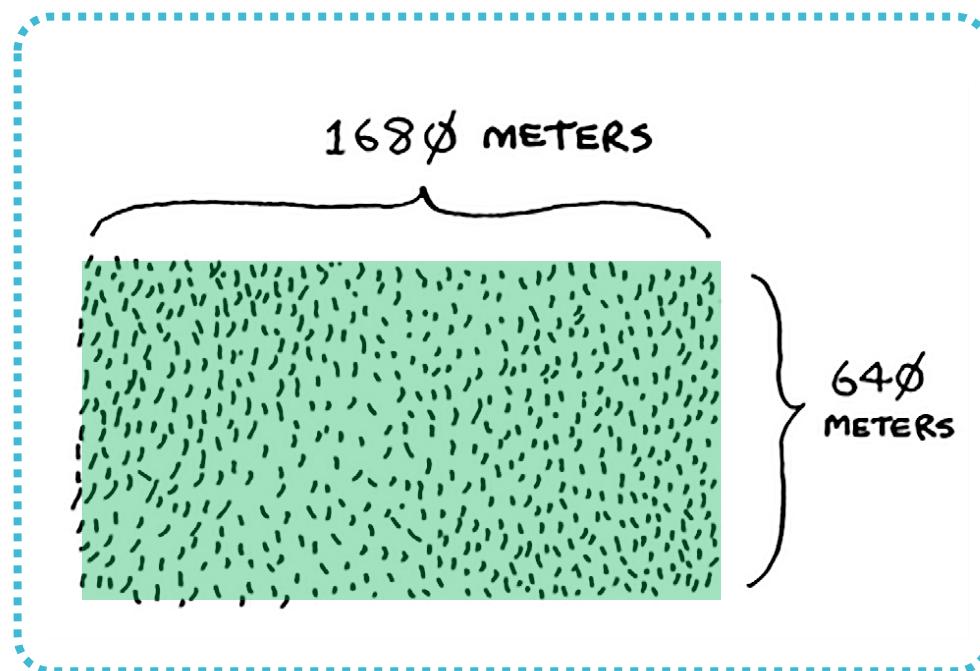
Divide & Conquer

- ❖ Well ... there are issues with each of these options.
- ❖ So **none of them** will work.



Divide & Conquer

❖ **Question revisited:** How do you find the **largest square size?**



Divide & Conquer

❖ Answer:

You can use **Divide & Conquer (D&C)** strategy!

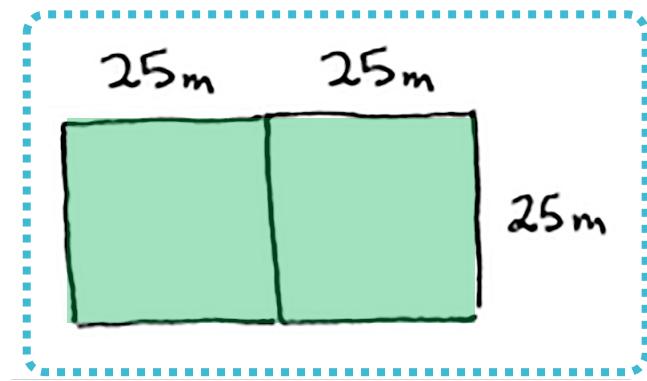
❖ **Divide & Conquer** strategy based algorithms are recursive and solve problem in two steps:

- **step 1:** figure out the base case (the *simplest* case).
- **step 2:** divide the problem until it becomes base case.

Divide & Conquer

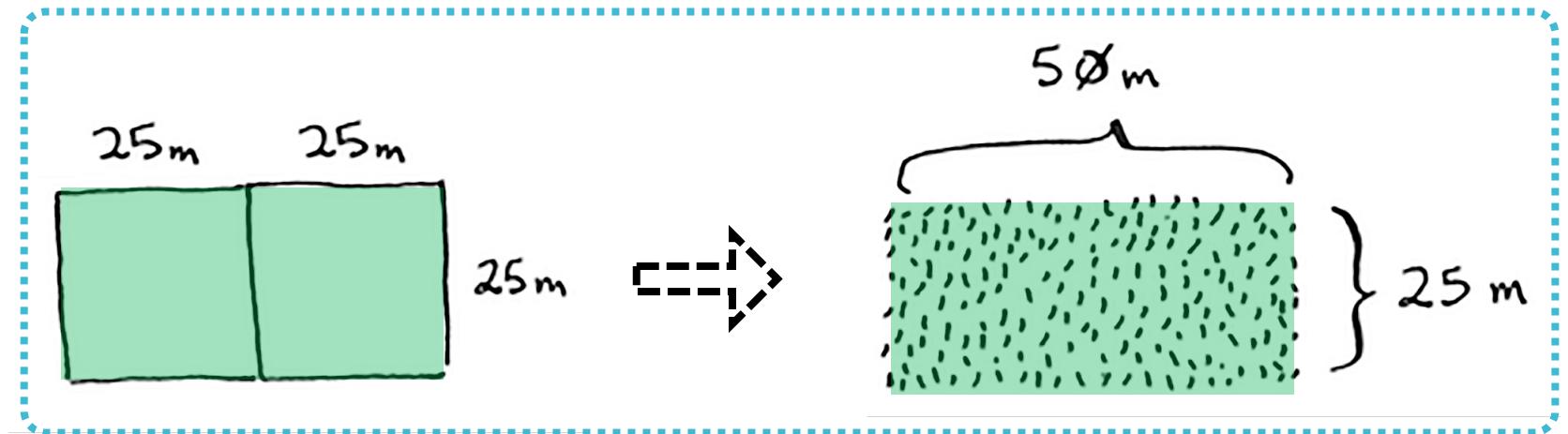
❖ Let's Divide & Conquer to find a solution to this problem:

❖ **step 1:** figure out the base case. Easiest case would be when one side was a multiple of another side, right?



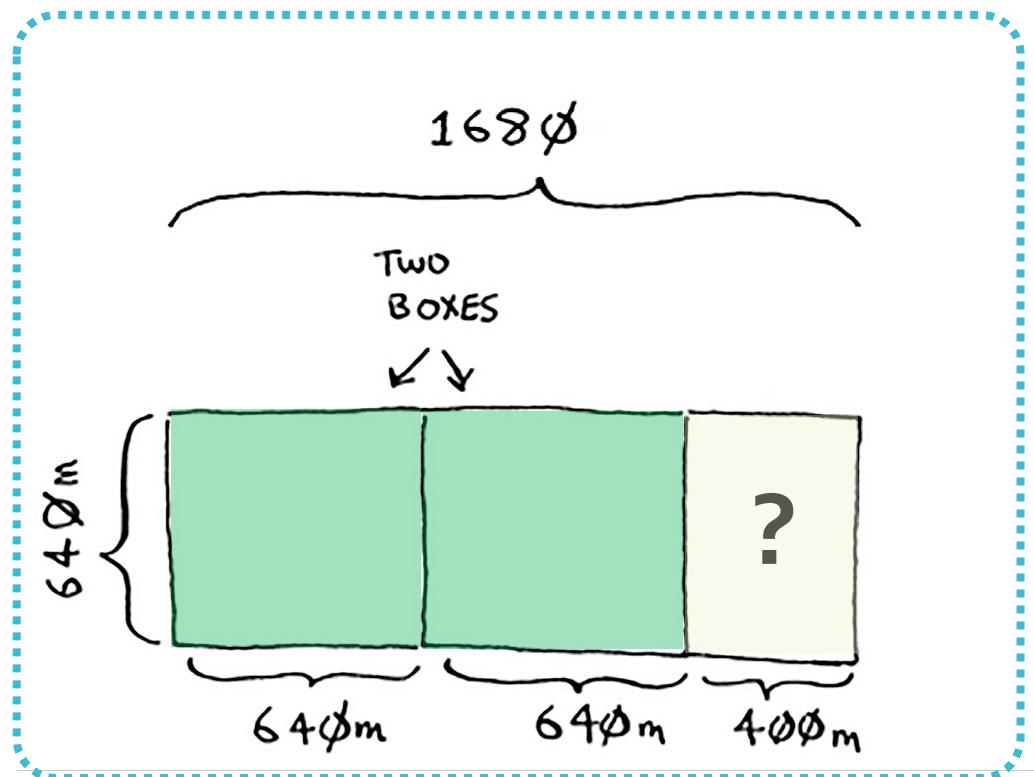
Divide & Conquer

- Suppose one side is **25 (m)** and the other side is **50 (m)**.
- Then, the largest box you can use is **$25 \text{ m} \times 25 \text{ m}$** .



Divide & Conquer

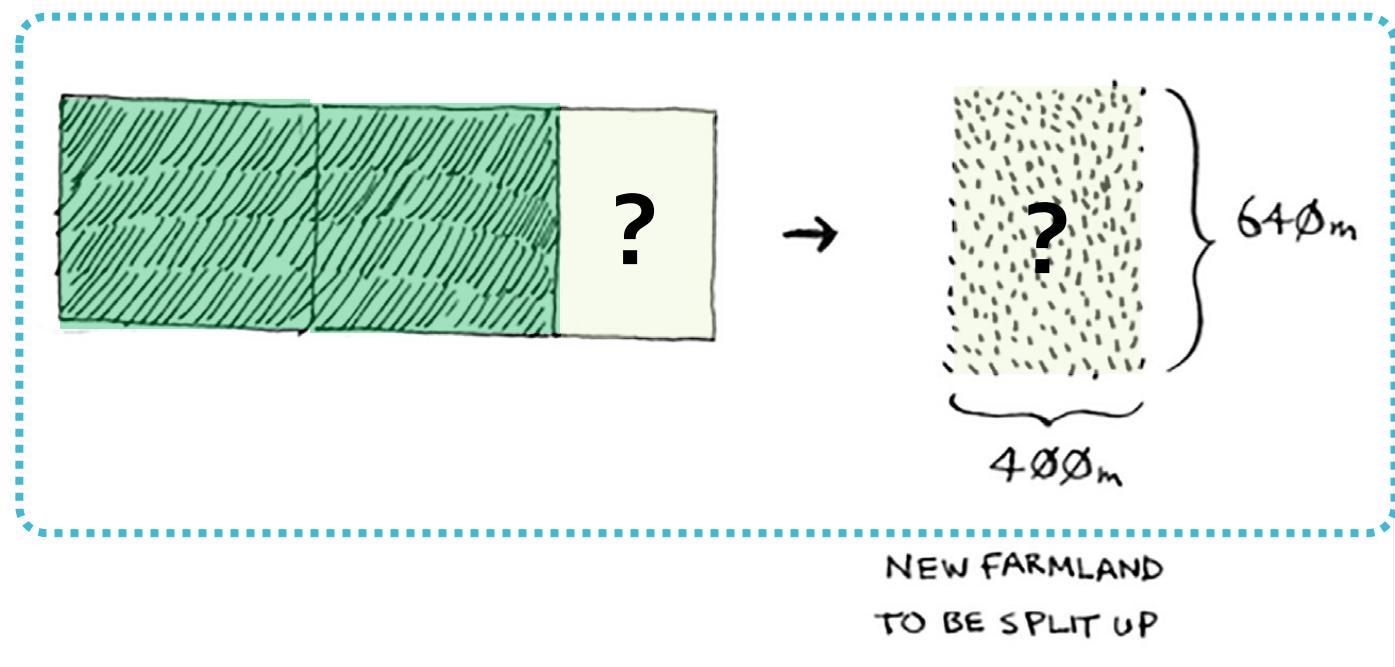
- ❖ **second step:** with every recursive call, you have to reduce your problem to a smaller problem.
- ❖ But How?
- ❖ Let's mark out **two biggest boxes** (640×640) to use.



Divide & Conquer

❖ But there's some land left to be divided.

❖ Why not to apply same algorithm here?



Divide & Conquer

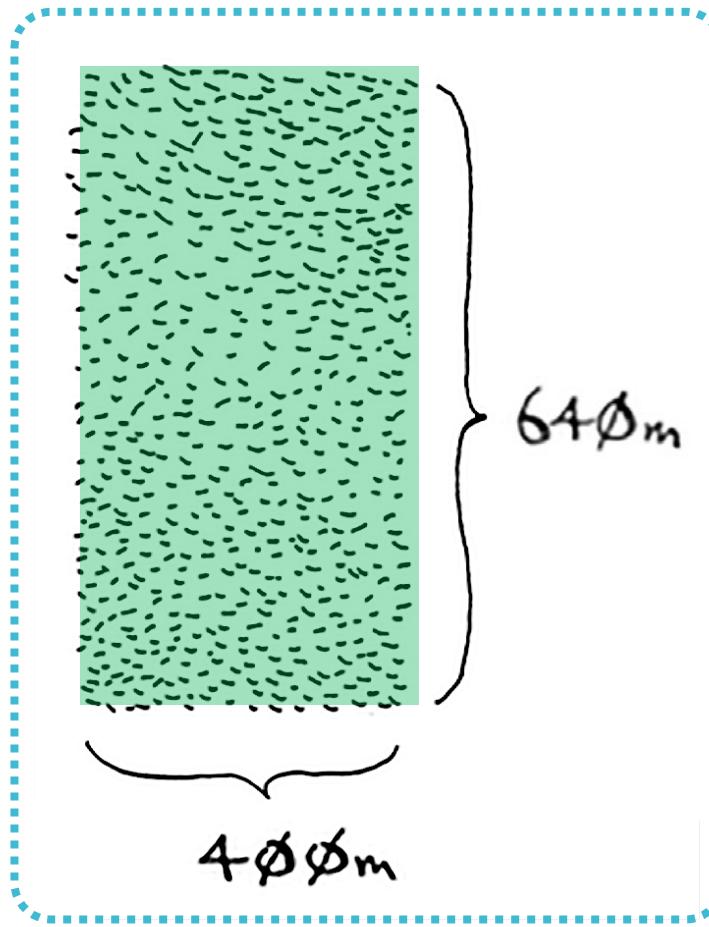
- ❖ What happened?
- ❖ You started out with a **1680×640 (m)** farm that needed to be split up.
- ❖ Now you need to split a smaller segment, **640×400 (m)**.

Divide & Conquer

- ❖ You just reduced the problem:
 - ❖ **from** a 1680×640 (m) farm
 - ❖ **to** a 640×400 (m) farm!
- ❖ You can continue and get the problem smaller and smaller.

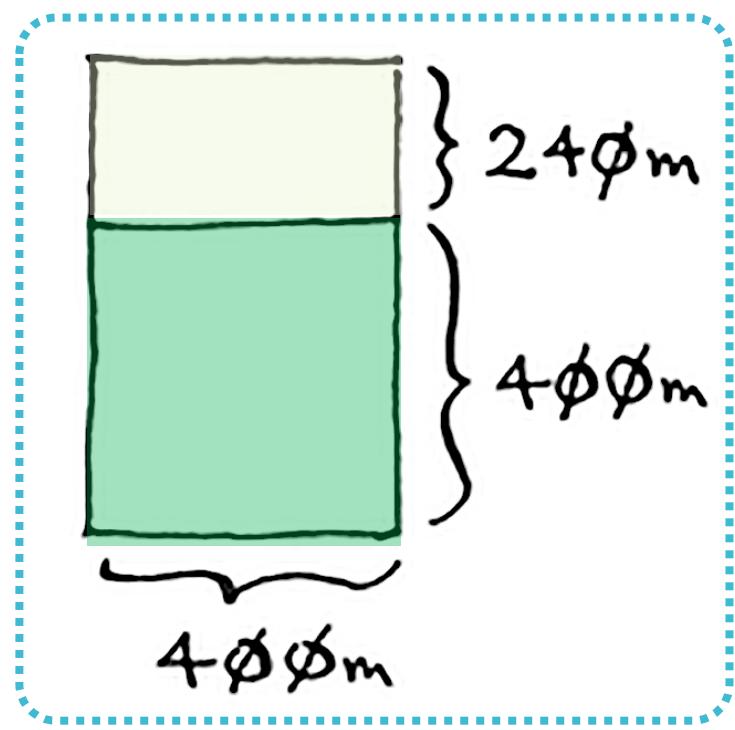
Quiz

- ❖ Smaller problem: a **land** with a **size 640×400 (m)**.
- ❖ Need to be divides it evenly into **square plots**.
- ❖ What will be the **biggest plot**?



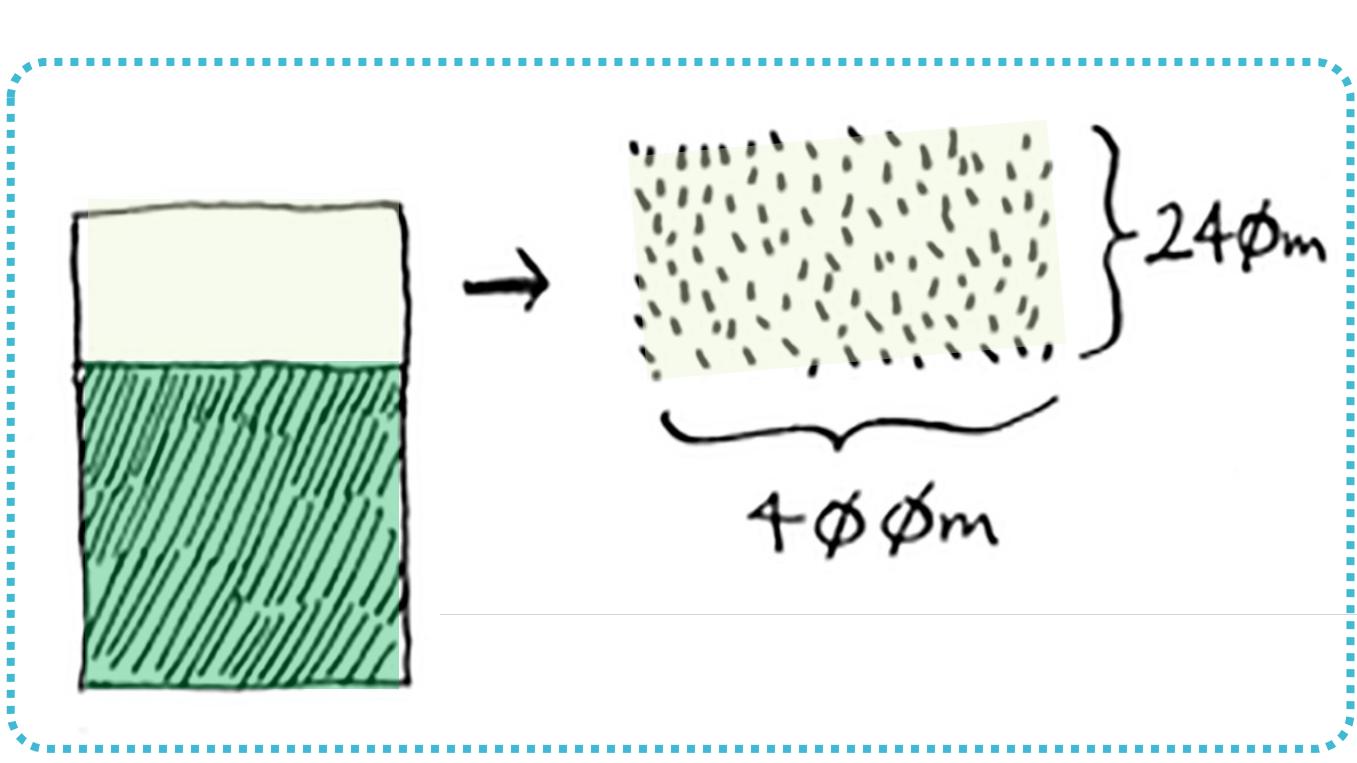
Answer

❖ Lets start with a plot of 400×400 (m).



Answer

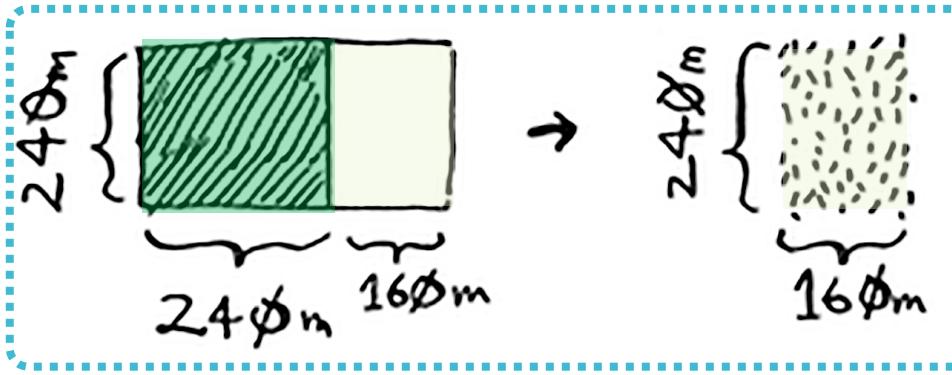
❖ And that leaves you with a smaller plot of 400×240 (m)



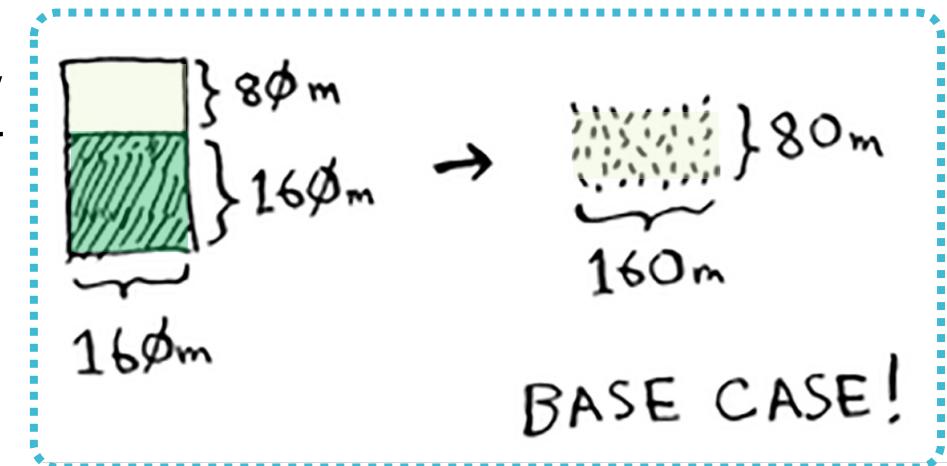
Answer

- ❖ And again, you can draw another box to get a smaller segment, 80×160 (m).

- ❖ This can be your **Base Case**.

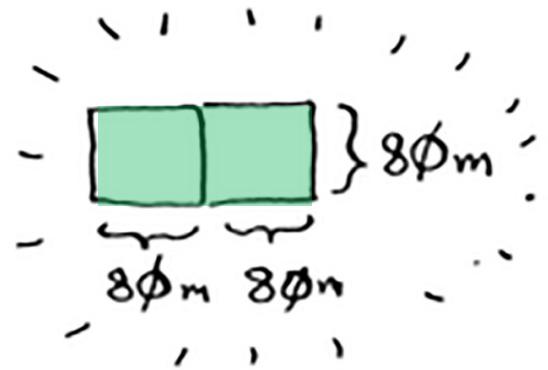


❖ You can draw a box on that to get a smaller segment, 240×160 (m).



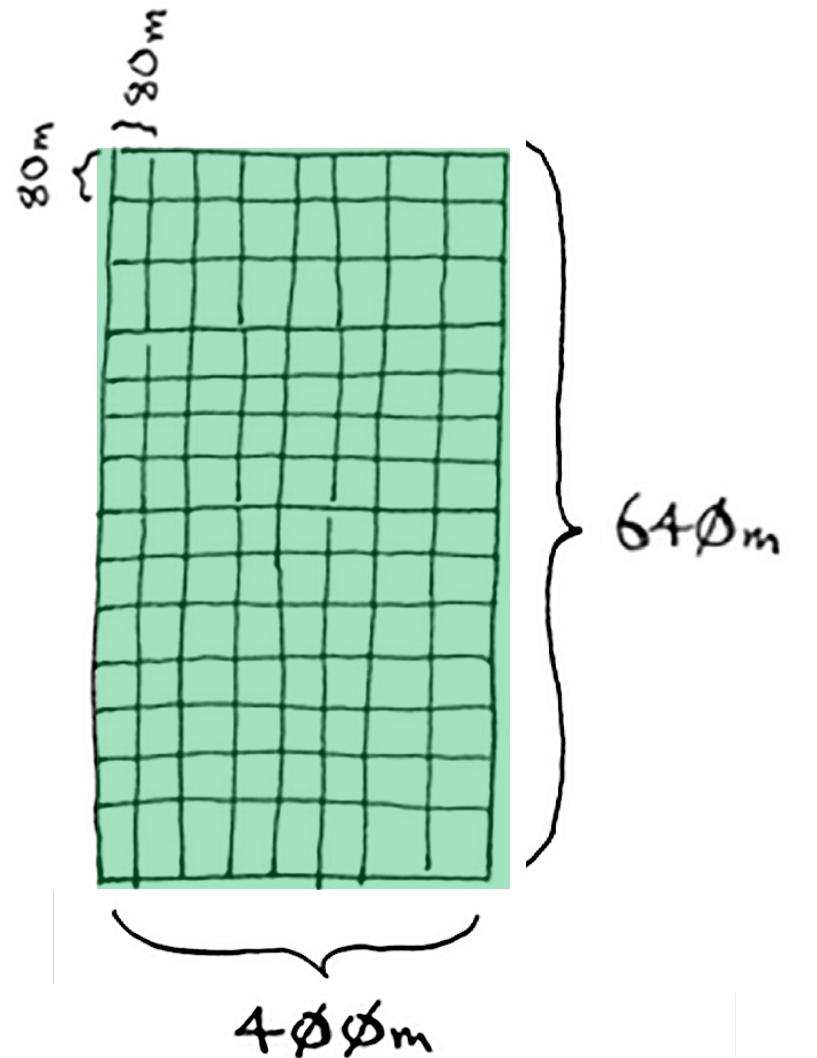
Answer

- ❖ Why Base Case?
- ❖ Because **80 is a factor of 160** and you can perfectly split up this segment using two boxes.
- ❖ And nothing is left over!



Answer

- ❖ At the end, for the original land, the biggest plot size you can use is 80×80 m.



Application

❖ Now we will apply this approach in a sorting algorithm.

Quick Sort

- ❖ **Quick Sort** is a sorting algorithm developed based on **Divide & Conquer**.
- ❖ It is faster than other sorts (e.g., selection sort).
- ❖ Developed by British computer scientist **Tony Hoare** in 1959.



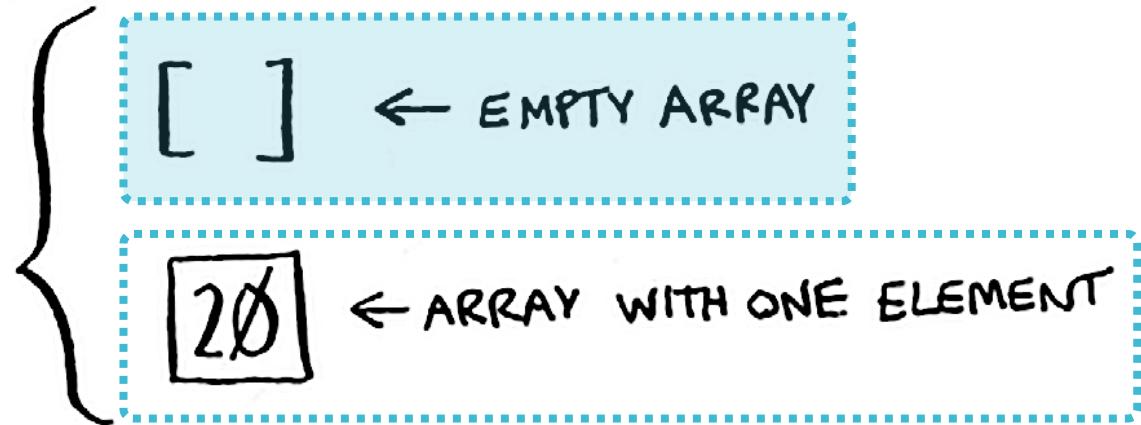
Quick Sort

- ❖ Suppose you would like to apply **Quicksort** on an array.
- ❖ What might be the **simplest** array to sort?

Quick Sort

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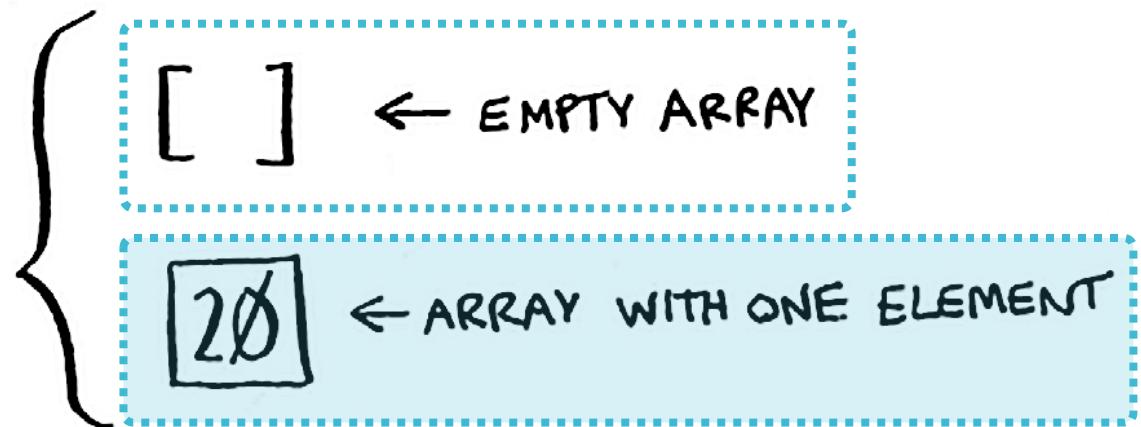
NO NEED
TO SORT
ARRAYS
LIKE THIS



Quick Sort

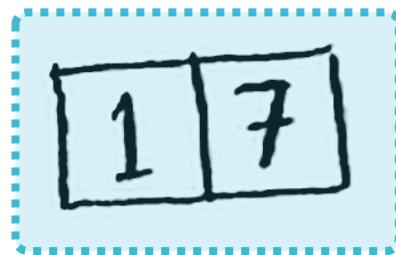
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Quick Sort

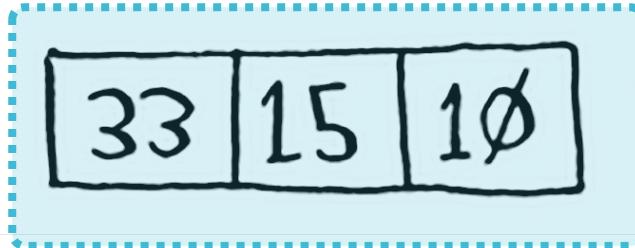
❖ What about an array with **two elements**?



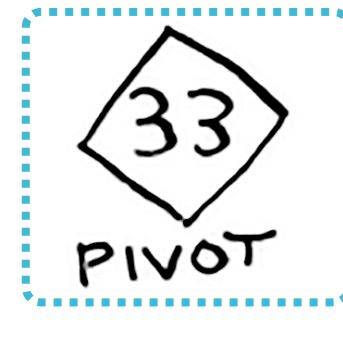
CHECK IF FIRST ELEMENT IS SMALLER THAN THE SECOND.
IF IT ISN'T, SWAP THEM.

Quick Sort

- ❖ And an array of **three elements**?

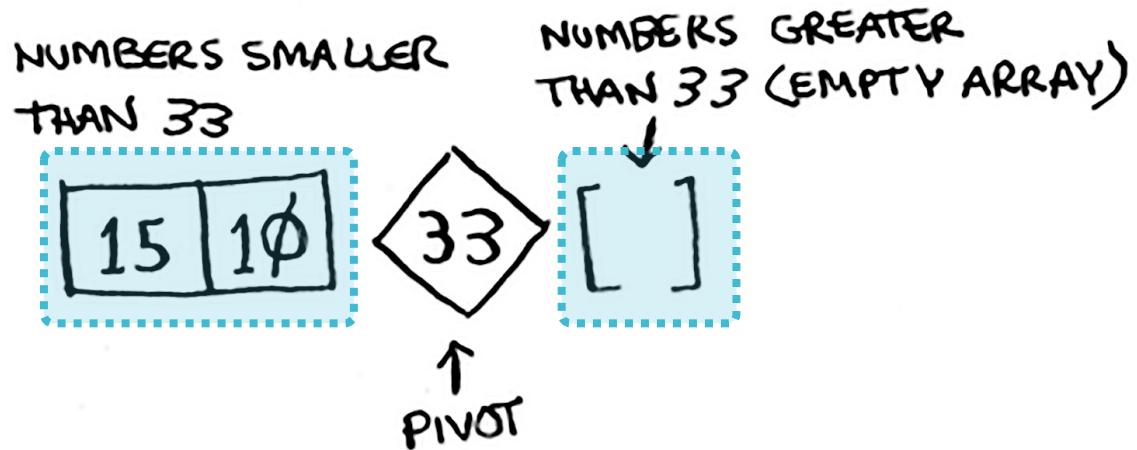


- ❖ First, **pick an element** in the array (called **Pivot**). Let's say the first element is the pivot.



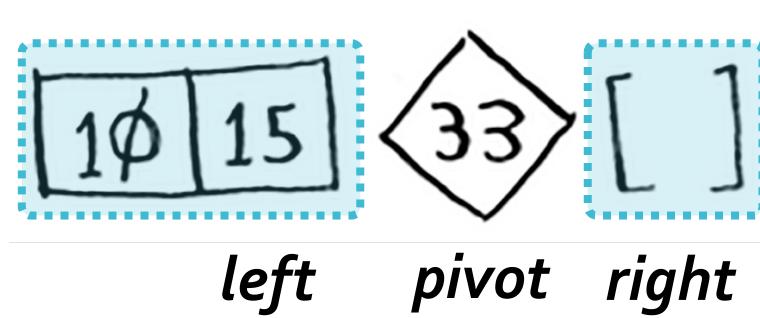
Quick Sort

- ❖ Now find the elements **smaller than the pivot** and the elements **larger than the pivot**.



Quick Sort

- ❖ This is called **partitioning**. The result is a:
 - ❖ **sub-array “left”**: elements less than pivot
 - ❖ **pivot**
 - ❖ **sub-array “right”**: elements greater than pivot



- ❖ Sub-arrays **may not be sorted**.

Quick Sort

- ❖ Idea: What if the sub-arrays are sorted?
- ❖ then you can just combine: `left` + pivot + `right`

- ❖ This will give you the sorted array:

$$[10, 15] + [33] + [] = [10, 15, 33]$$

Quick Sort

❖ Quick sort algorithm:

- 1) Pick a **pivot**.
- 2) Partition the array into two sub-arrays:
 - a) *left*: elements **less than pivot**
 - b) *right*: elements **greater than pivot**.
- 3) Call quicksort **recursively** on two sub-arrays.

Quick Sort

❖ Quick sort algorithm:

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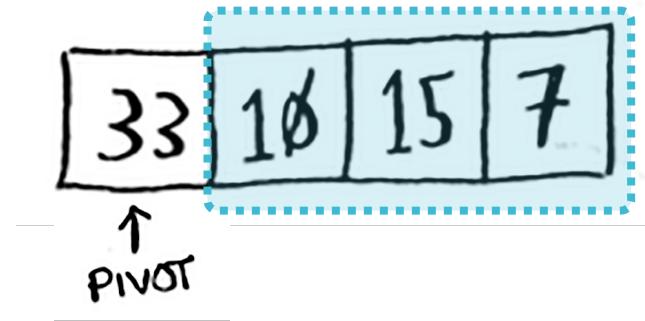
Quick Sort

❖ Quick sort algorithm:

- 1) Pick a **pivot**.
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Quick Sort

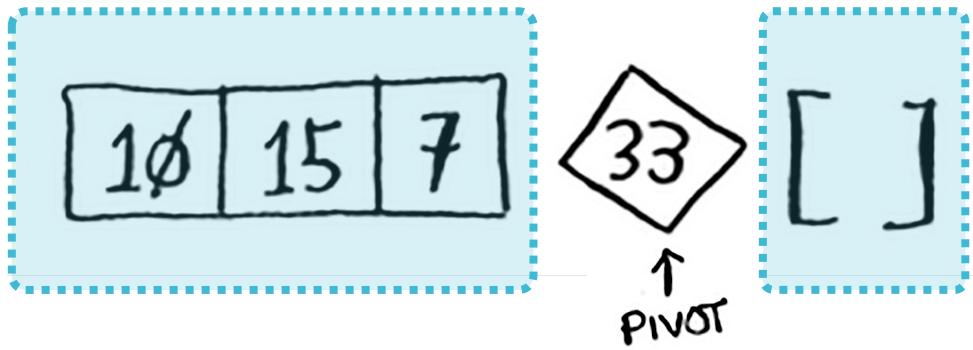
- ❖ We did for array of **three elements**.
- ❖ Suppose you have an array of **four elements**.



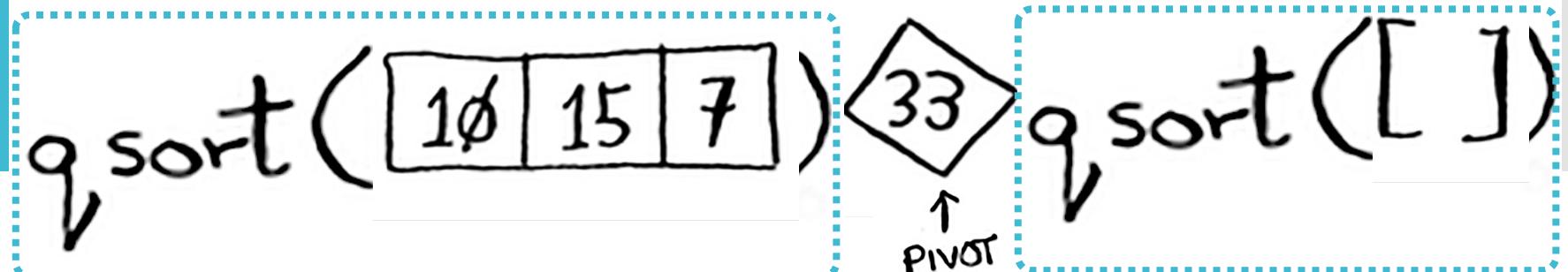
- ❖ First, pick a **pivot** (for example 33).

Quick Sort

- ❖ The array on the left has three elements.

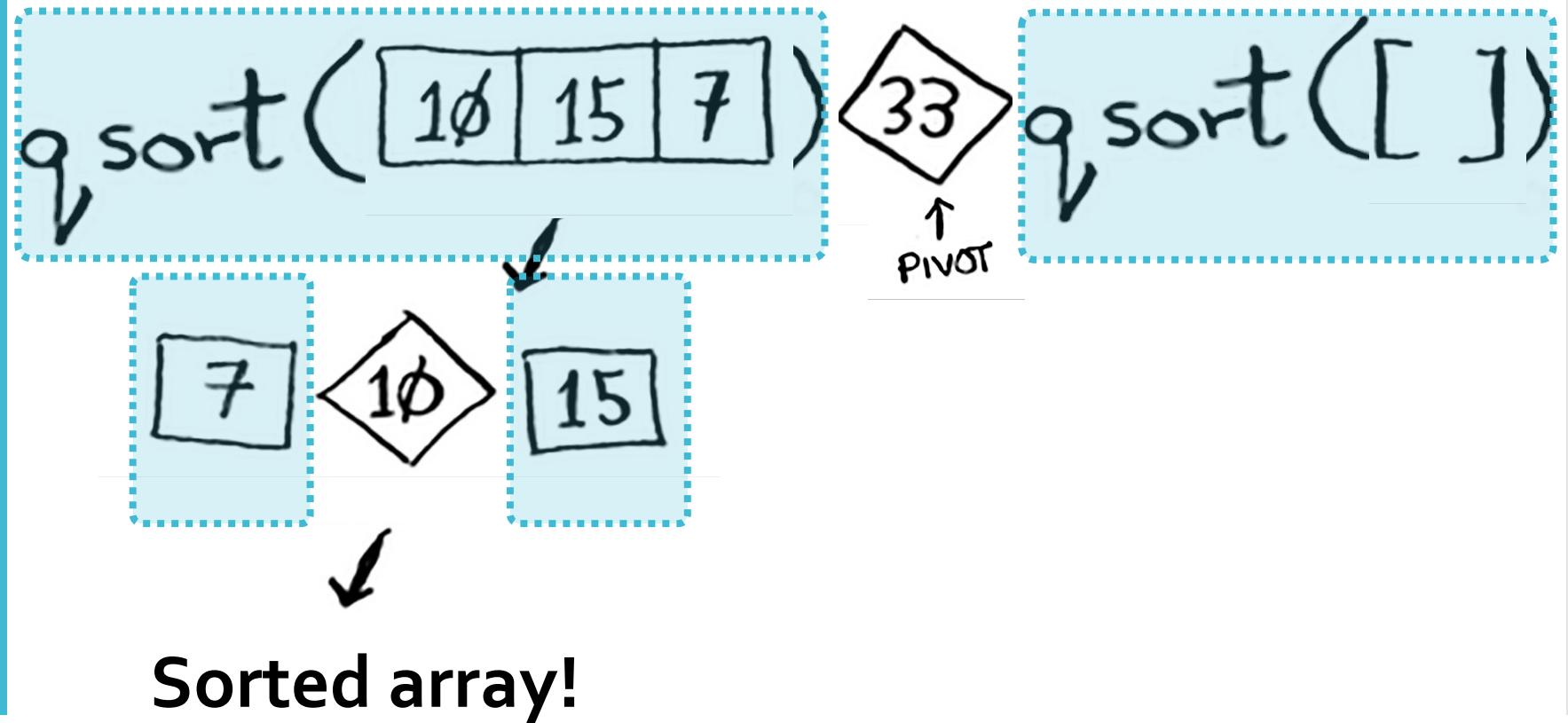


- ❖ So you can call quicksort on it **recursively**.



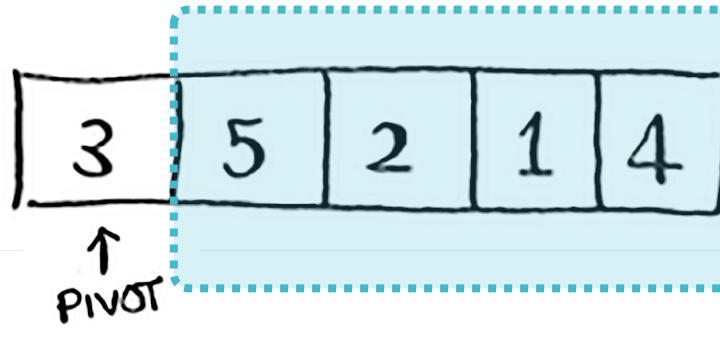
Quick Sort

❖ You already done quicksort on array of size three.



Quick Sort

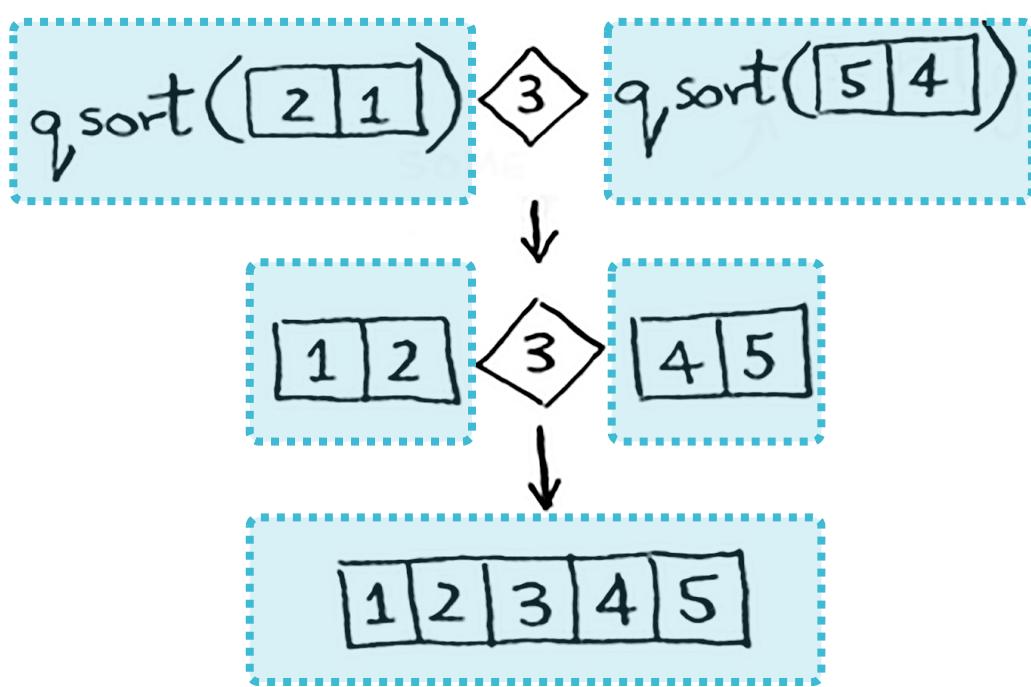
- ❖ What about **larger arrays**?
- ❖ Suppose you have an array of **five elements**.



- ❖ First, pick a pivot (for example 3).

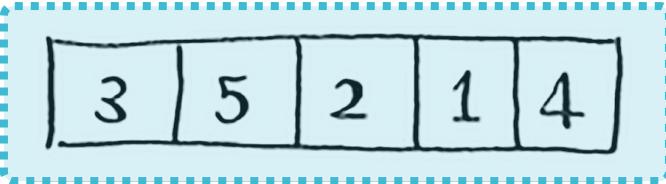
Quick Sort

- ❖ After partitioning, you have the sub-arrays.
- ❖ Then, you call quicksort (**qsort**) on sub-arrays.
- ❖ Finally, you combine the sub-arrays and pivot.

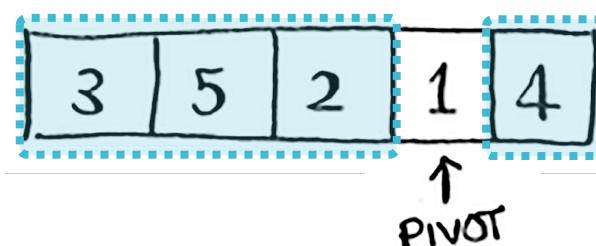
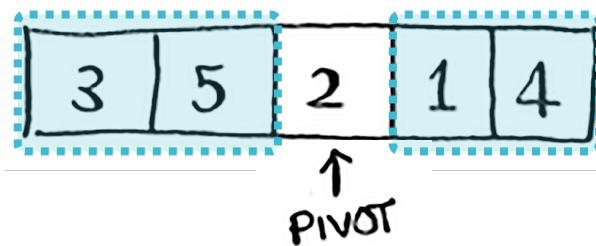
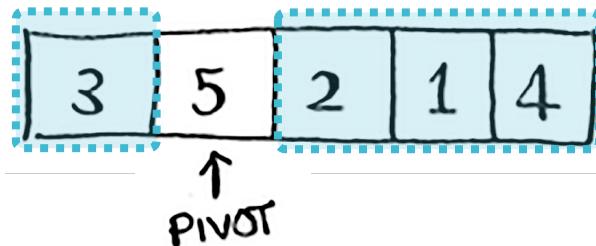
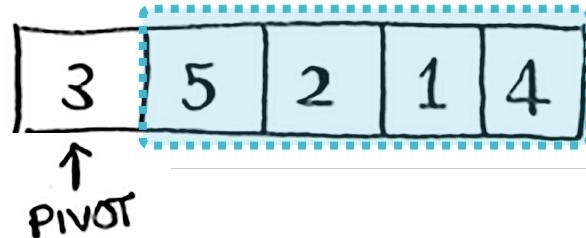


Quiz

- ❖ Suppose you have an array of five elements.

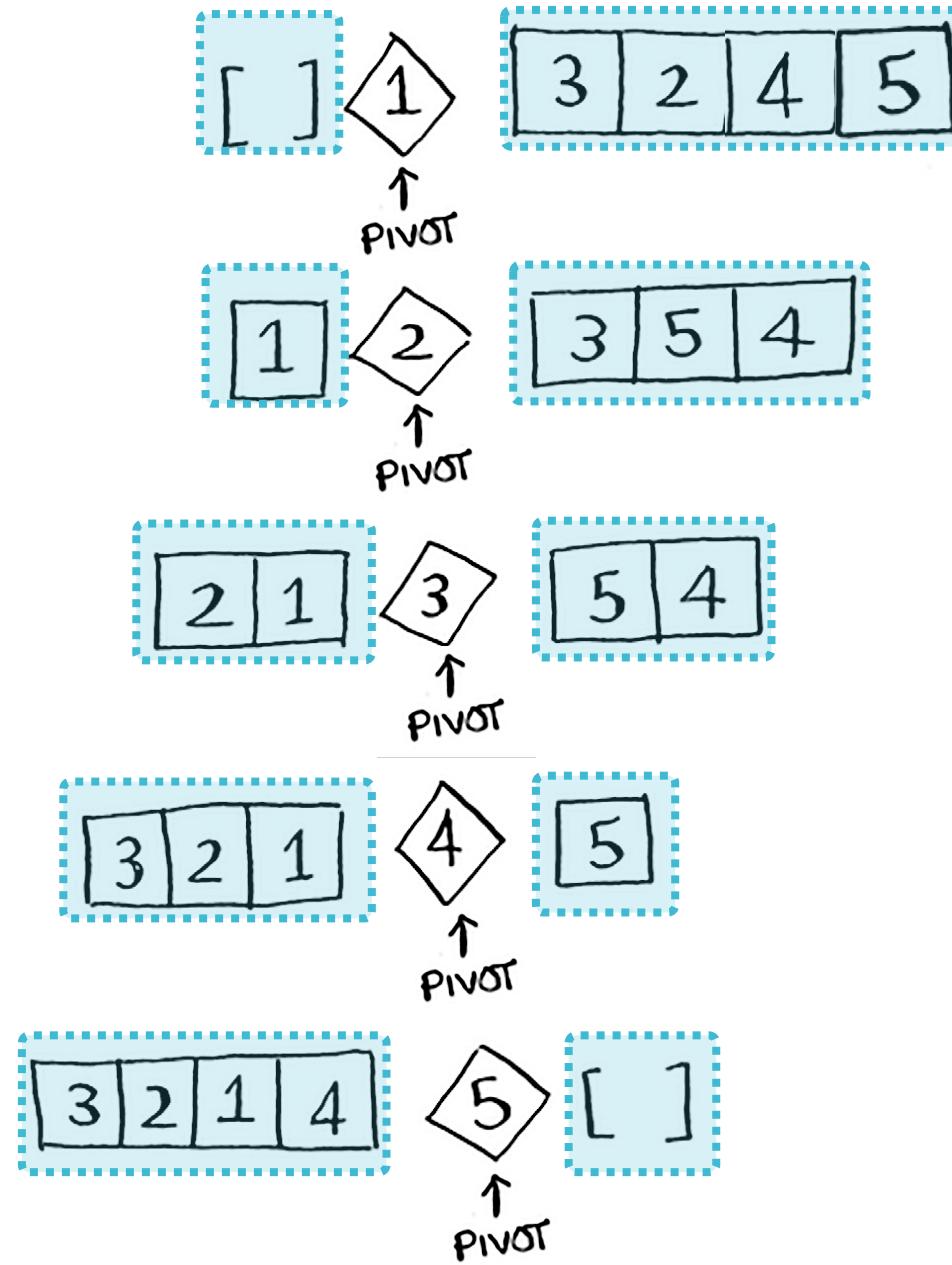


- ❖ Write down the ways you can partition this array, depending on different **pivots**?



Answer

- ❖ You can partition the array, in different ways, depending on different pivots.

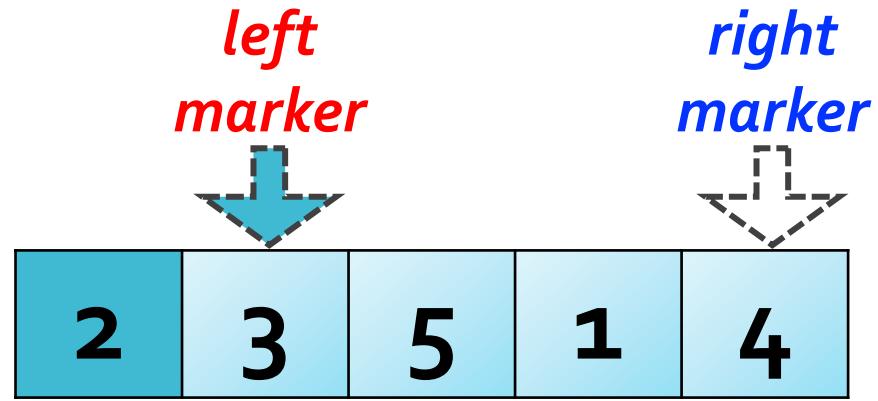


Quick Sort

2	3	5	1	4
---	---	---	---	---

*Lets Quick sort this!
We choose 2 as Pivot*

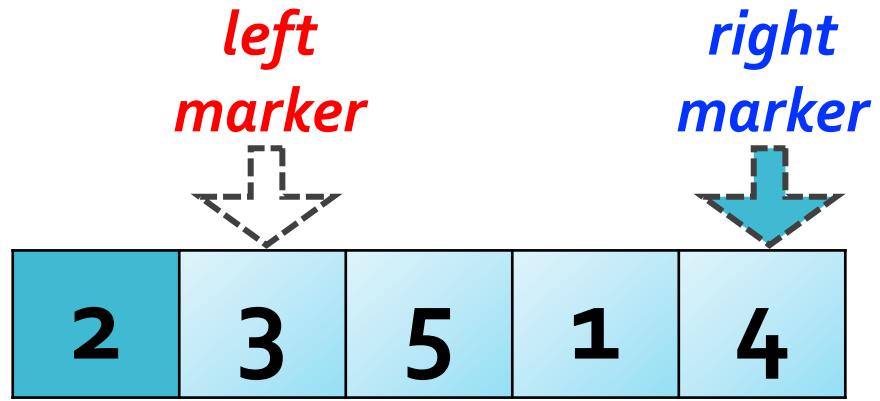
Quick Sort



left marker:

*3 is larger than 2 (pivot):
stop!*

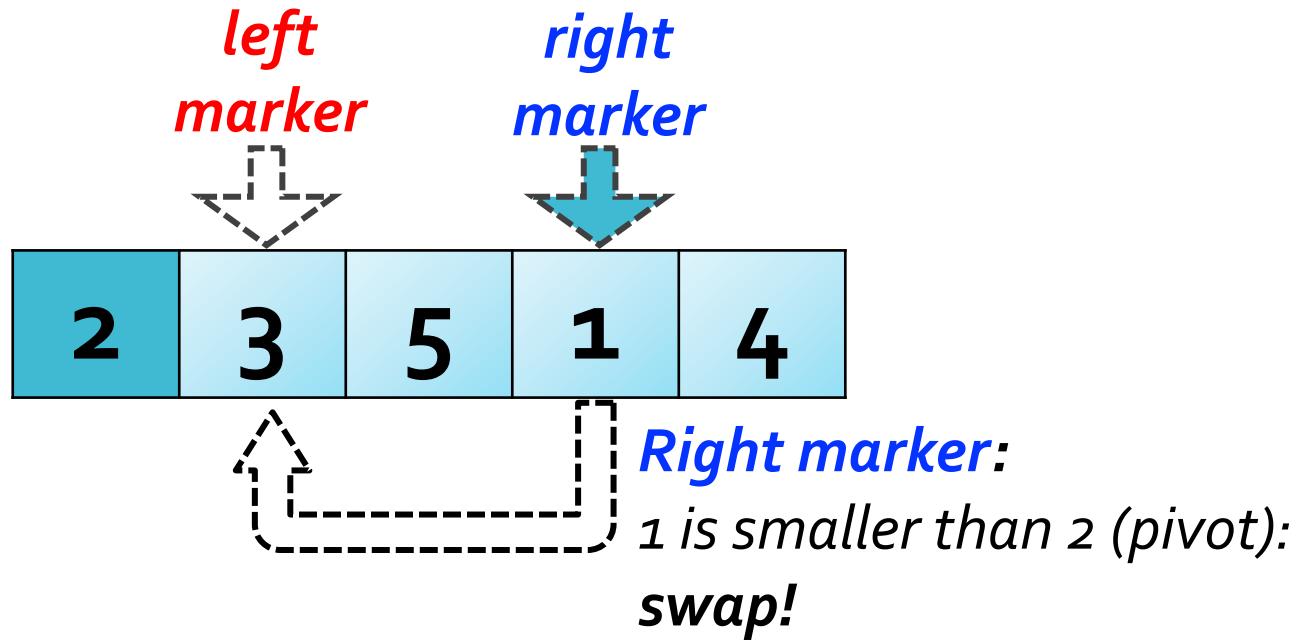
Quick Sort



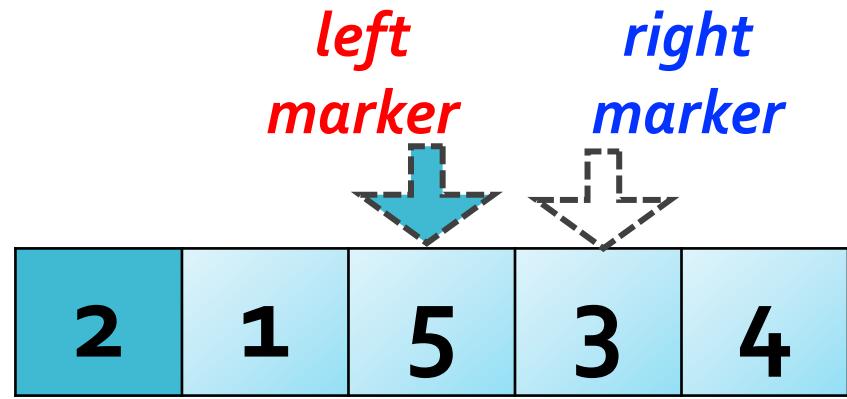
Right marker:

*4 is not smaller than 2 (pivot):
move!*

Quick Sort



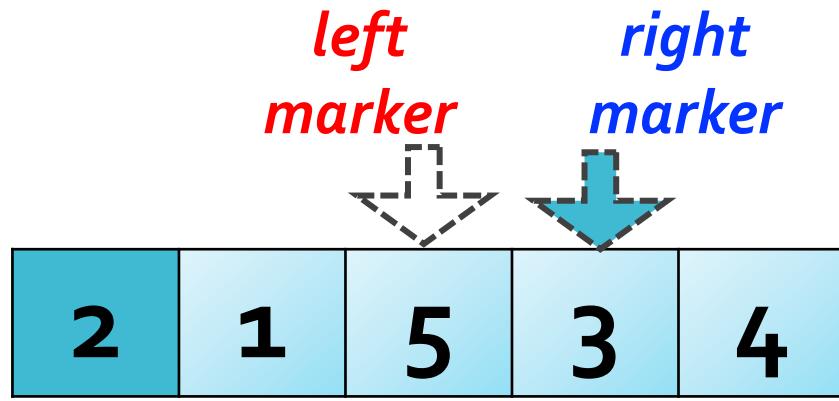
Quick Sort



left marker:

*5 is greater than 2 (pivot):
stop!*

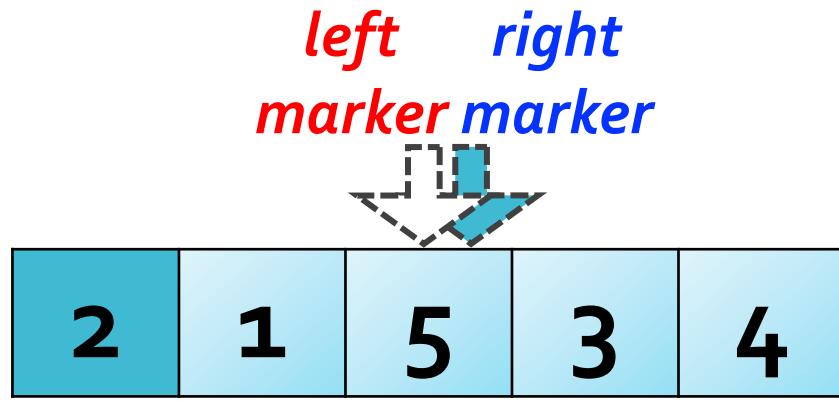
Quick Sort



right marker:

*3 is not smaller than 2 (pivot):
move!*

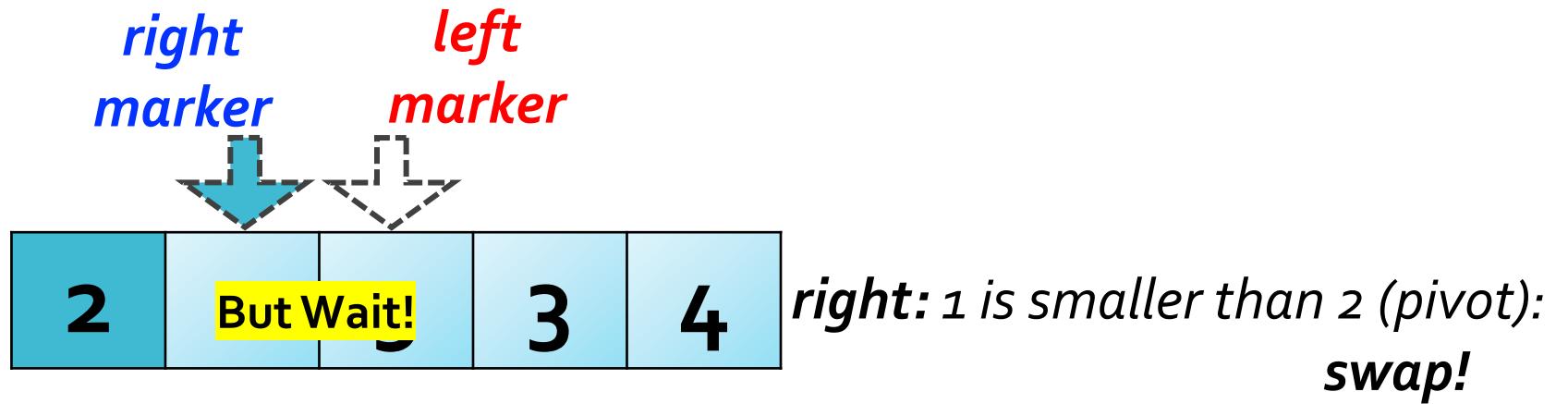
Quick Sort



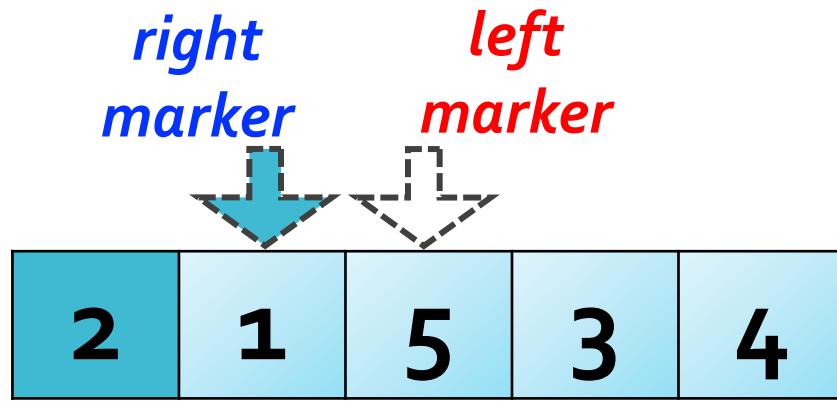
right marker:

*5 is not smaller than 2 (pivot):
move!*

Quick Sort

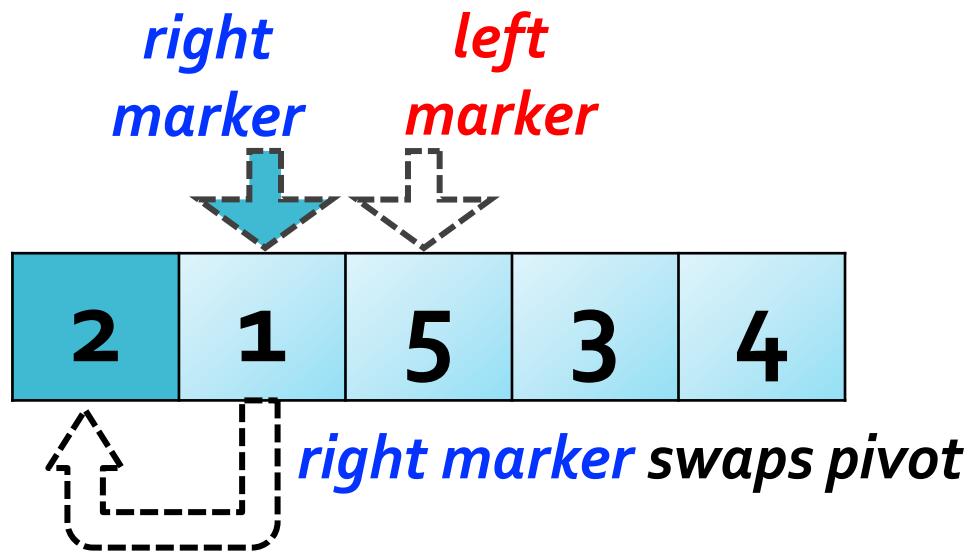


Quick Sort



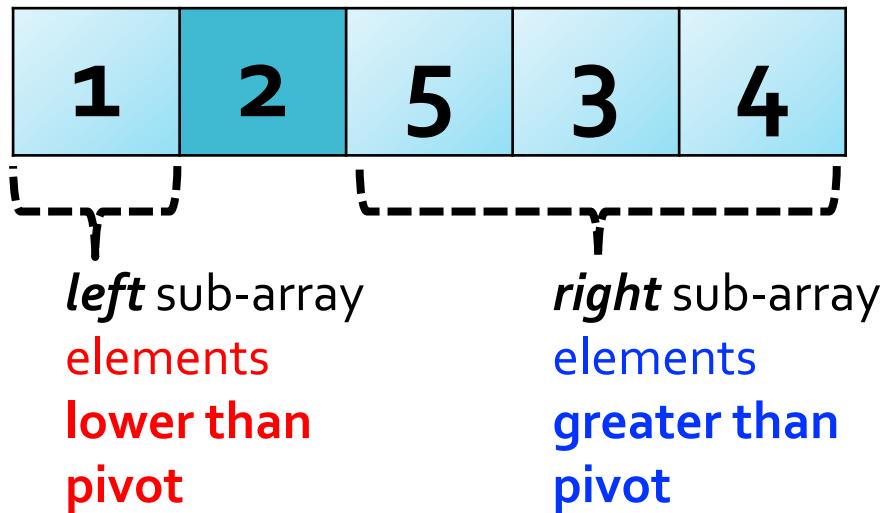
right marker crossed *red marker*
Then, **loop ends!**

Quick Sort



Quick Sort

Now you call Quick sort on the *left* and *right* sub-arrays.
At the end, the array gets sorted.



Bubble sort vs Quick Sort



ref: Youtube

Quick Sort

❖ Implementation of the Quick sort

```
def quick_sort(arr):

    if len(arr) < 2:
        return arr

    pivot = arr[0]
    left = [i for i in arr[1:] if i <= pivot]
    right = [i for i in arr[1:] if i > pivot]

    return quick_sort(left) + [pivot] + \
           quick_sort(right)

music_count = [156, 141, 35, 94, 88, 61, 111]
music_sorted = quick_sort(music_count)

print(music_sorted)
```

[Output]: [35, 61, 88, 94, 111, 141, 156]

Quick Sort

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[Output]: [35, 61, 88, 94, 111, 141, 156]
```

list comprehension

- ❖ always returns a **result list**.
- ❖ concise way to create lists.

- ❖ the following:

```
[expression for item in list if condition]
```

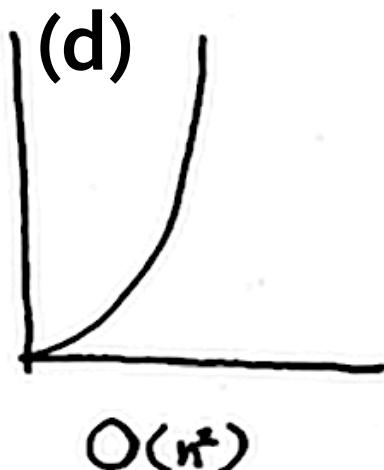
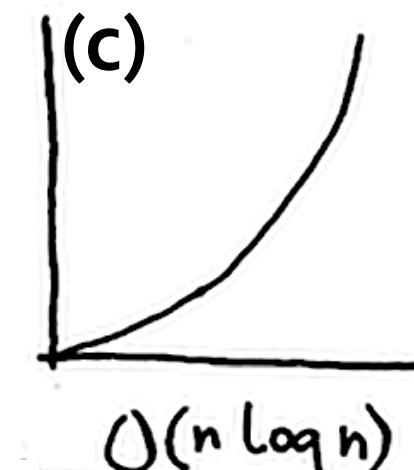
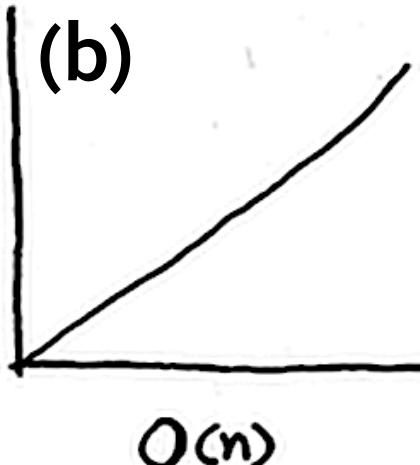
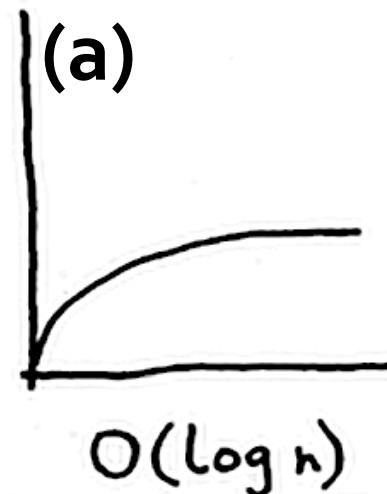
- ❖ is equivalent to

```
for item in list:  
    if condition:  
        expression
```

Quiz

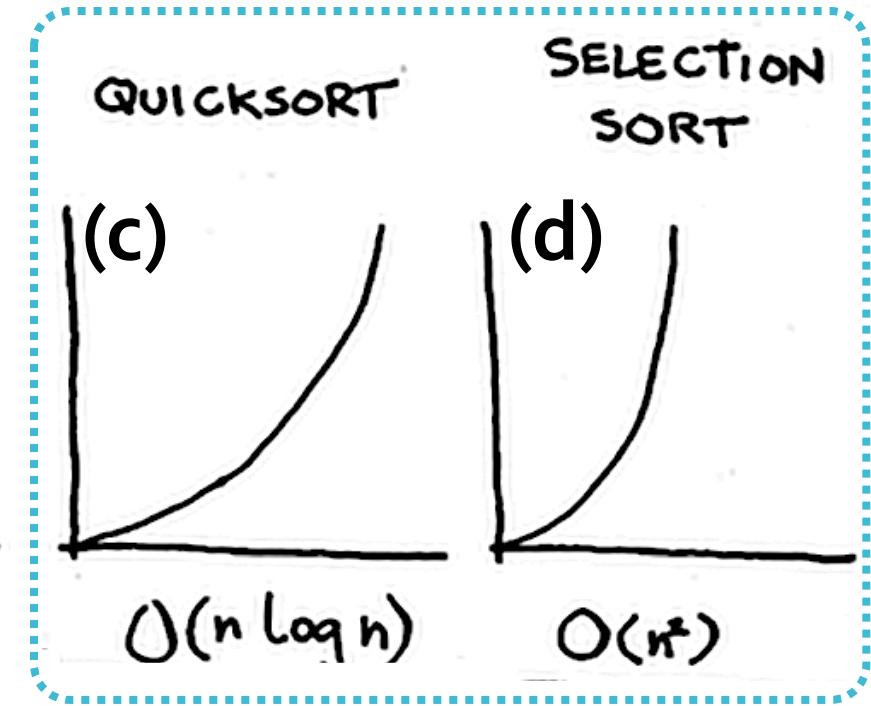
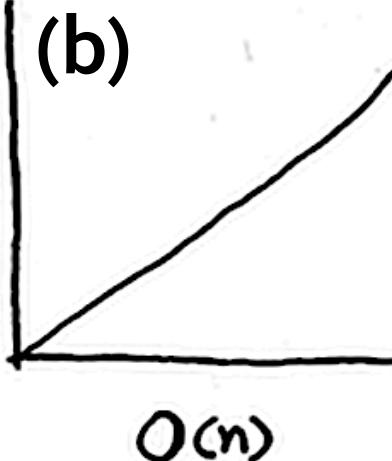
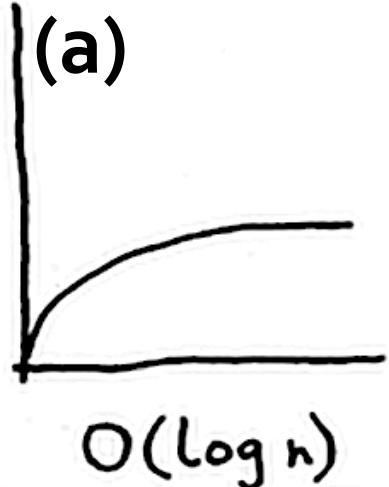
❖ Which of the following figures could represent the Big O notation of for:

- 1) Quick sort
- 2) Selection sort



Answer

❖ Which of the following figures could represent the Big O notation of for:

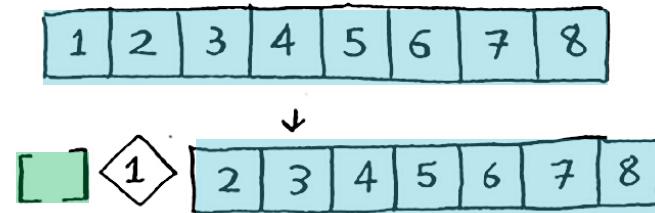


- ❖ Quick sort is good, but its performance heavily **depends on the chosen pivot.**
- ❖ Lets assume a list **already sorted:**

1	2	3	4	5	6	7	8
---	---	---	---	---	---	---	---

Quick Sort

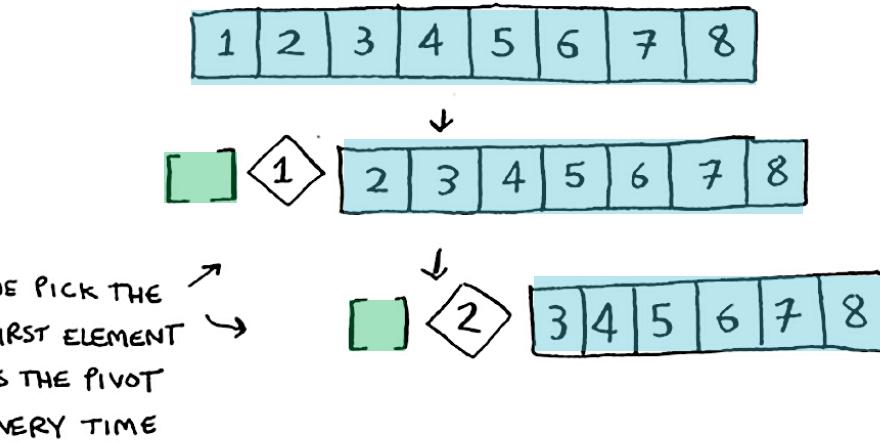
- ❖ Performance of quicksort heavily **depends on** the chosen **pivot**.



Quick Sort

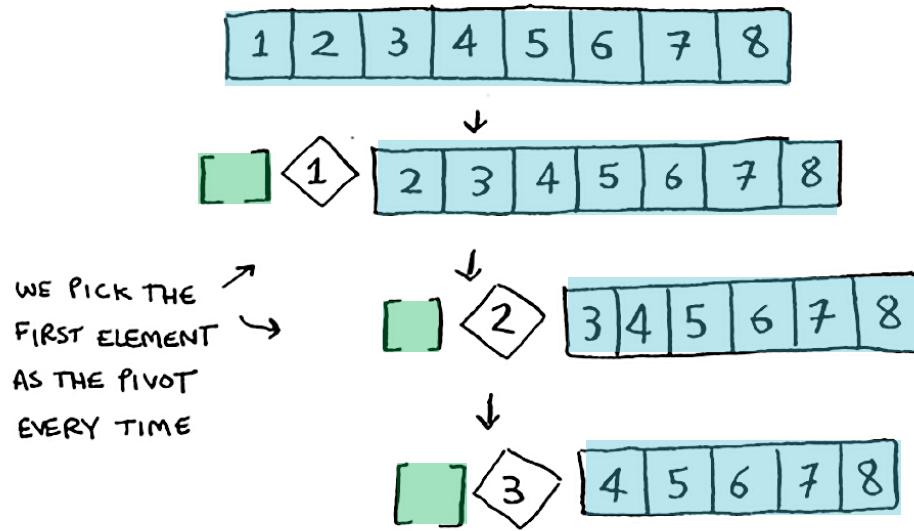
Quick Sort

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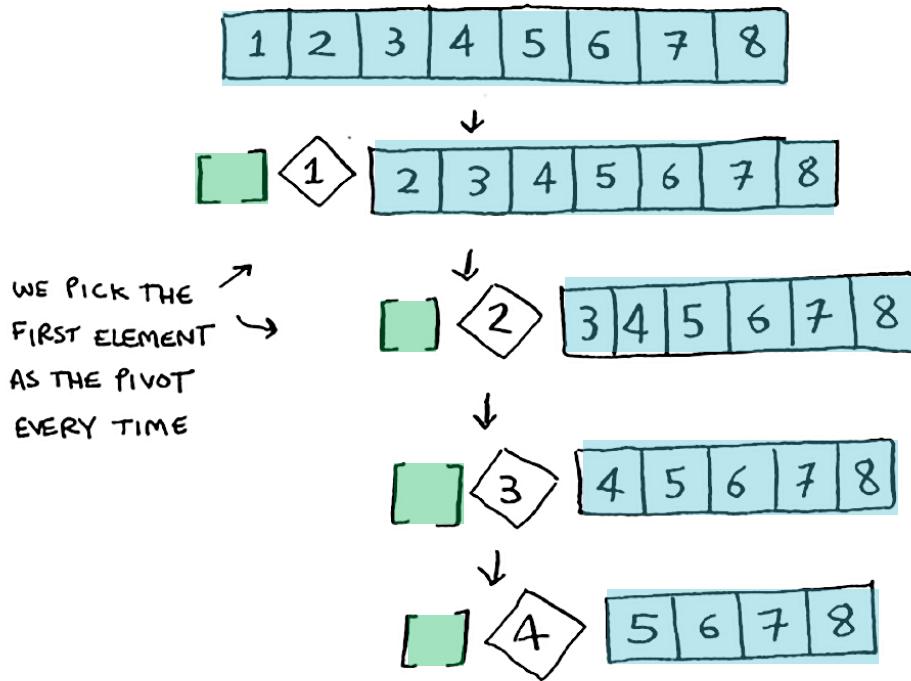
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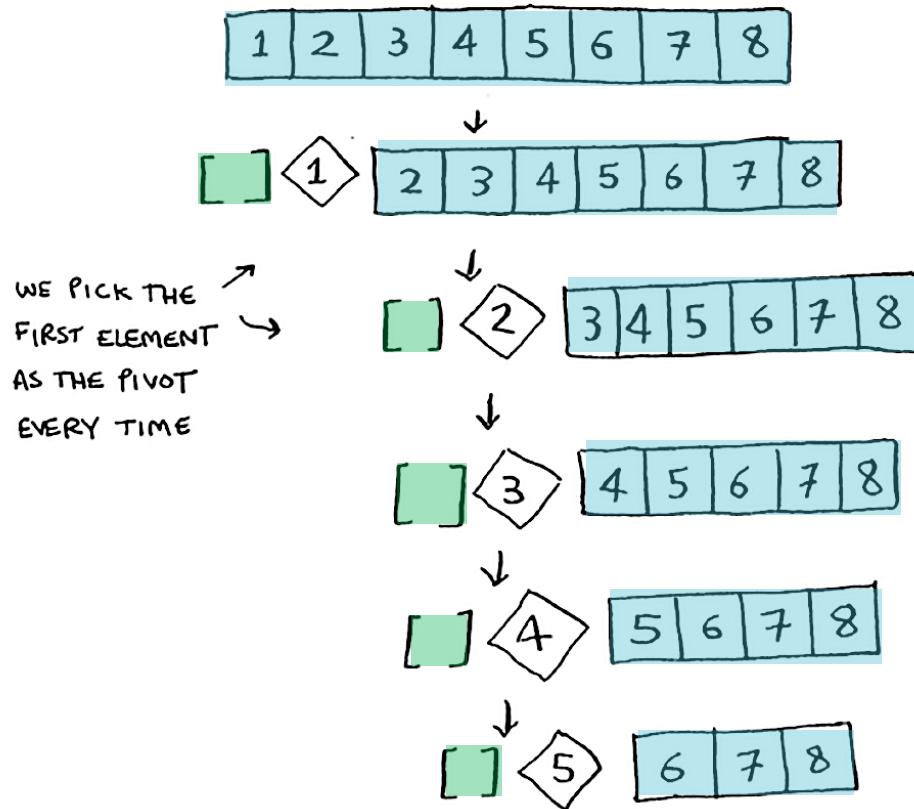
Quick Sort

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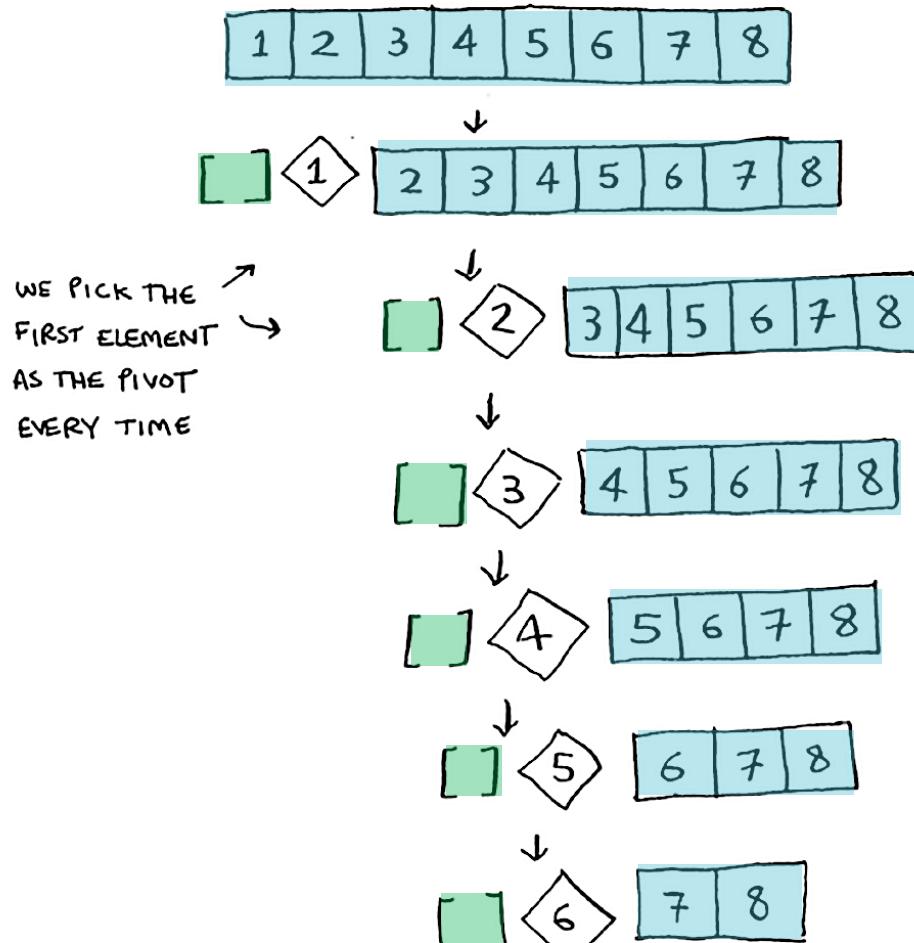
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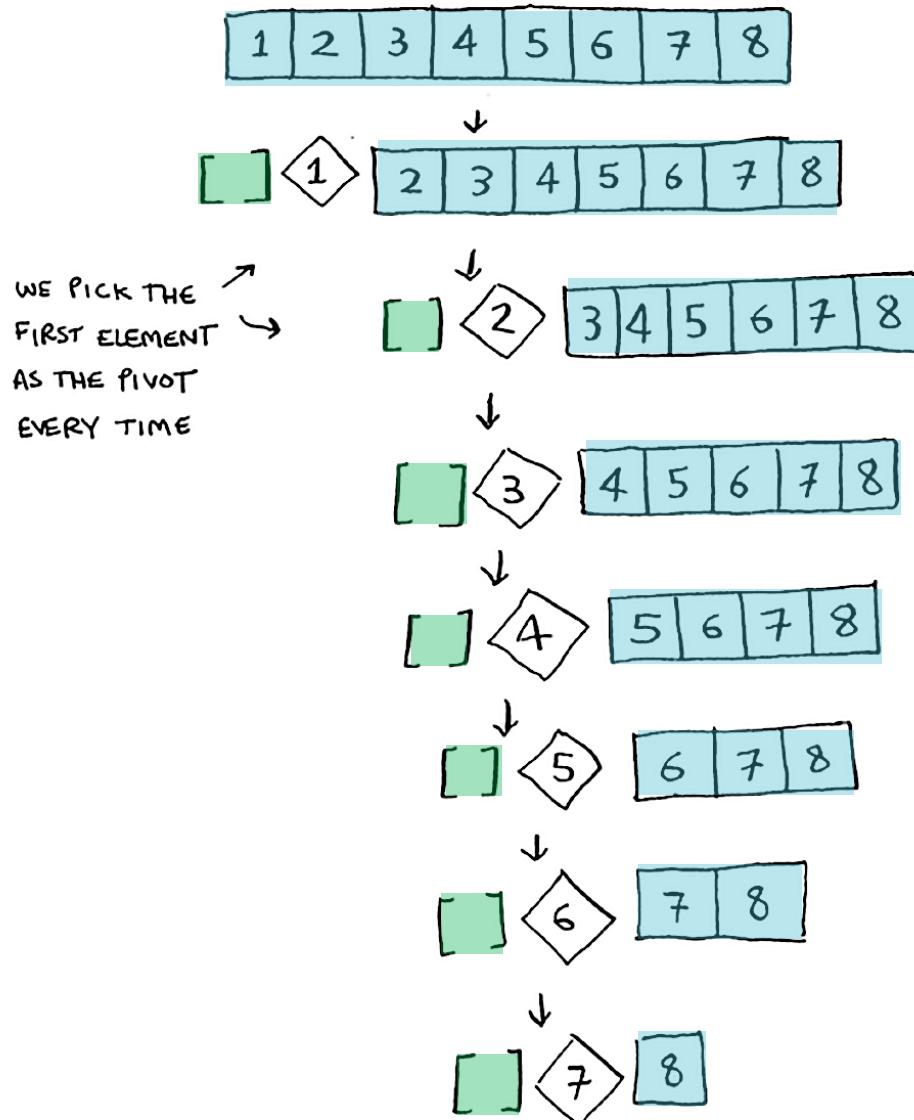
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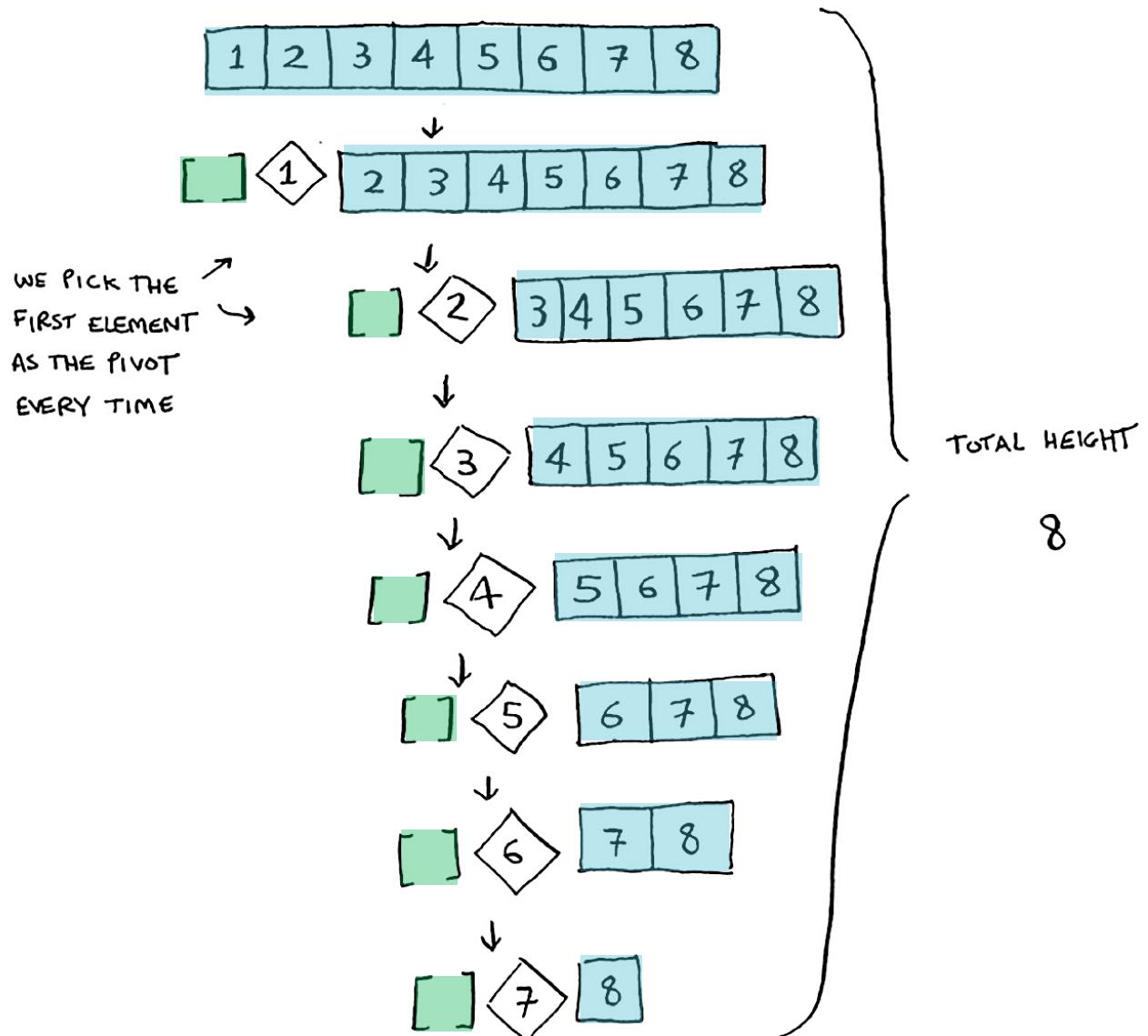
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Quick Sort

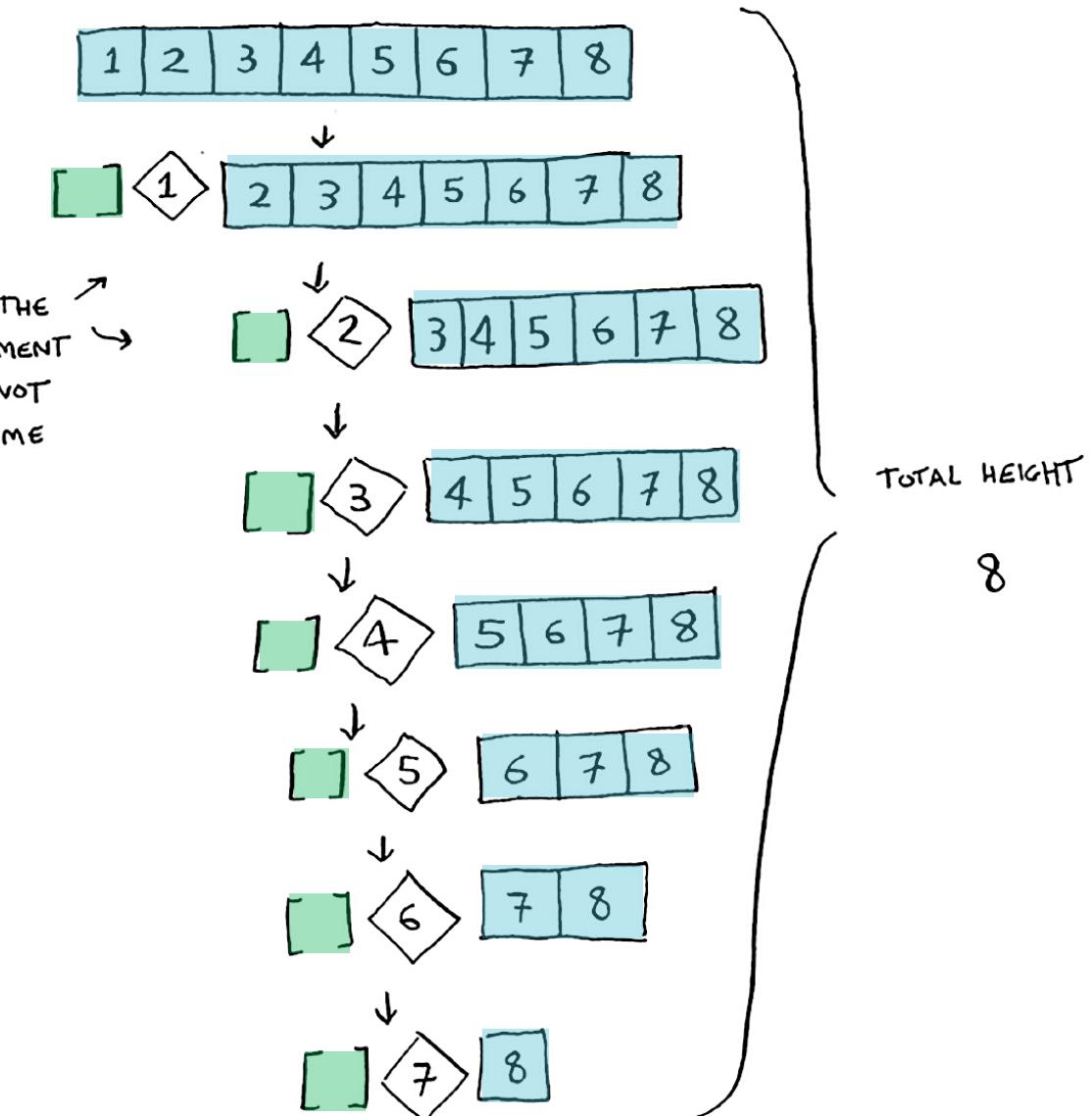
❖ Performance of quicksort heavily **depends on the chosen pivot.**



Quick Sort

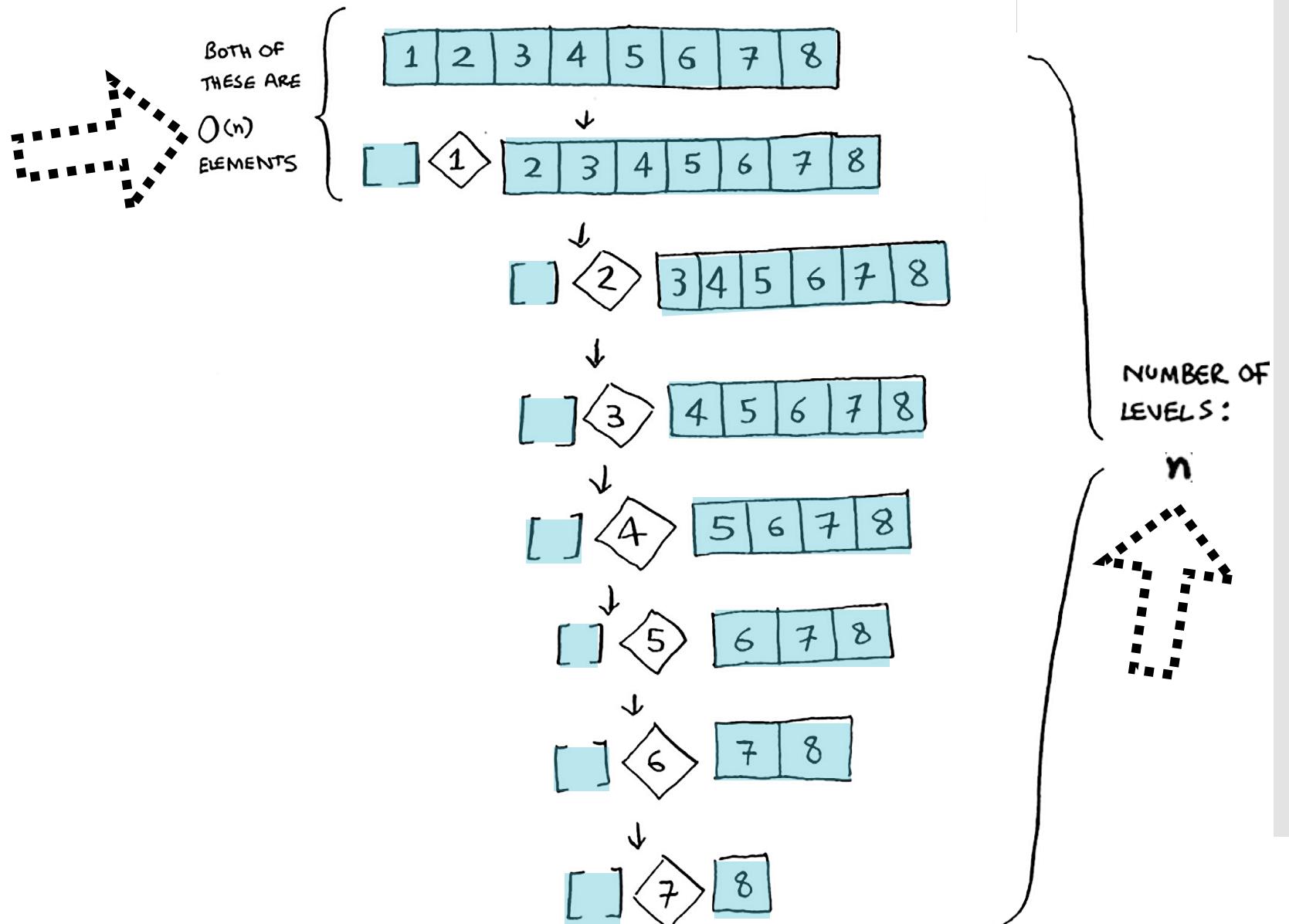
Notice that the
left sub-arrays is
always empty!

- ❖ Performance of quicksort heavily **depends on the chosen pivot.**



❖ Performance of quicksort heavily depends on the chosen pivot.

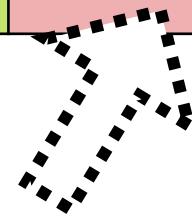
Quick Sort



Quick Sort

- ❖ Performance of quicksort heavily depends on the chosen **pivot**.

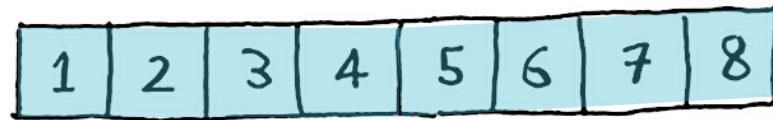
	Best case	Average case	Worse case
Quick Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$



Quick Sort

- ❖ Even if you partition the array differently, you're still touching **O(n) elements every time.**

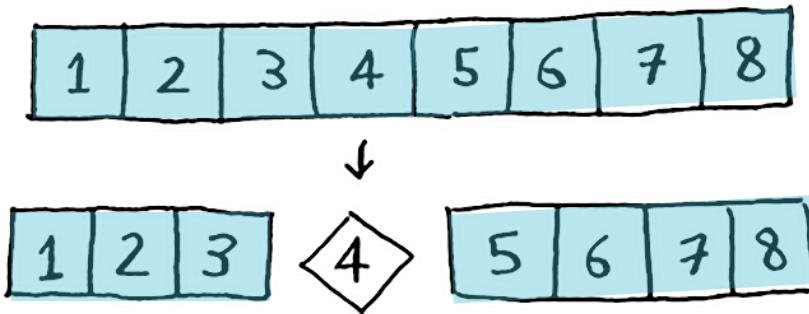
8 ELEMENTS
i.e. $O(n)$ ELEMENTS {



Quick Sort

8 ELEMENTS
i.e. $O(n)$ ELEMENTS {

8 ELEMENTS
i.e. $O(n)$ ELEMENTS {



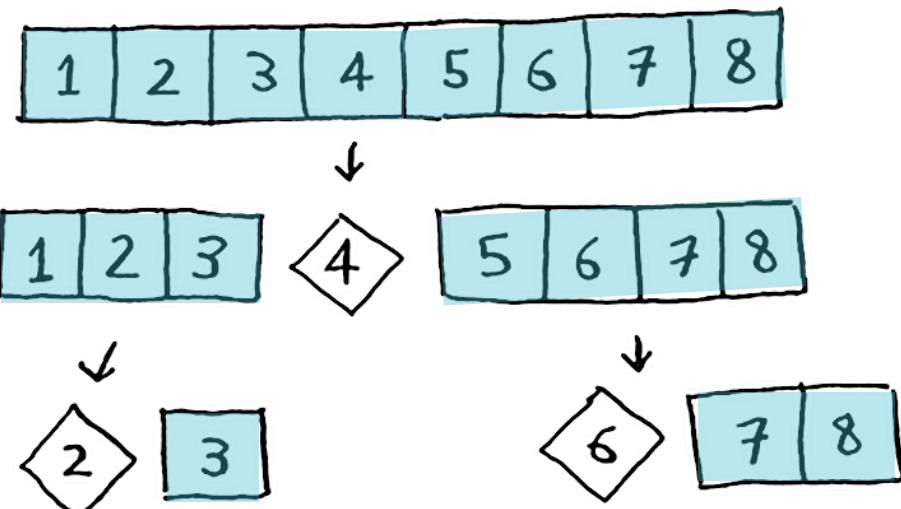
- ❖ Even if you partition the array differently, you're still touching **$O(n)$ elements every time.**

Quick Sort

8 ELEMENTS
i.e. $O(n)$ ELEMENTS {

8 ELEMENTS
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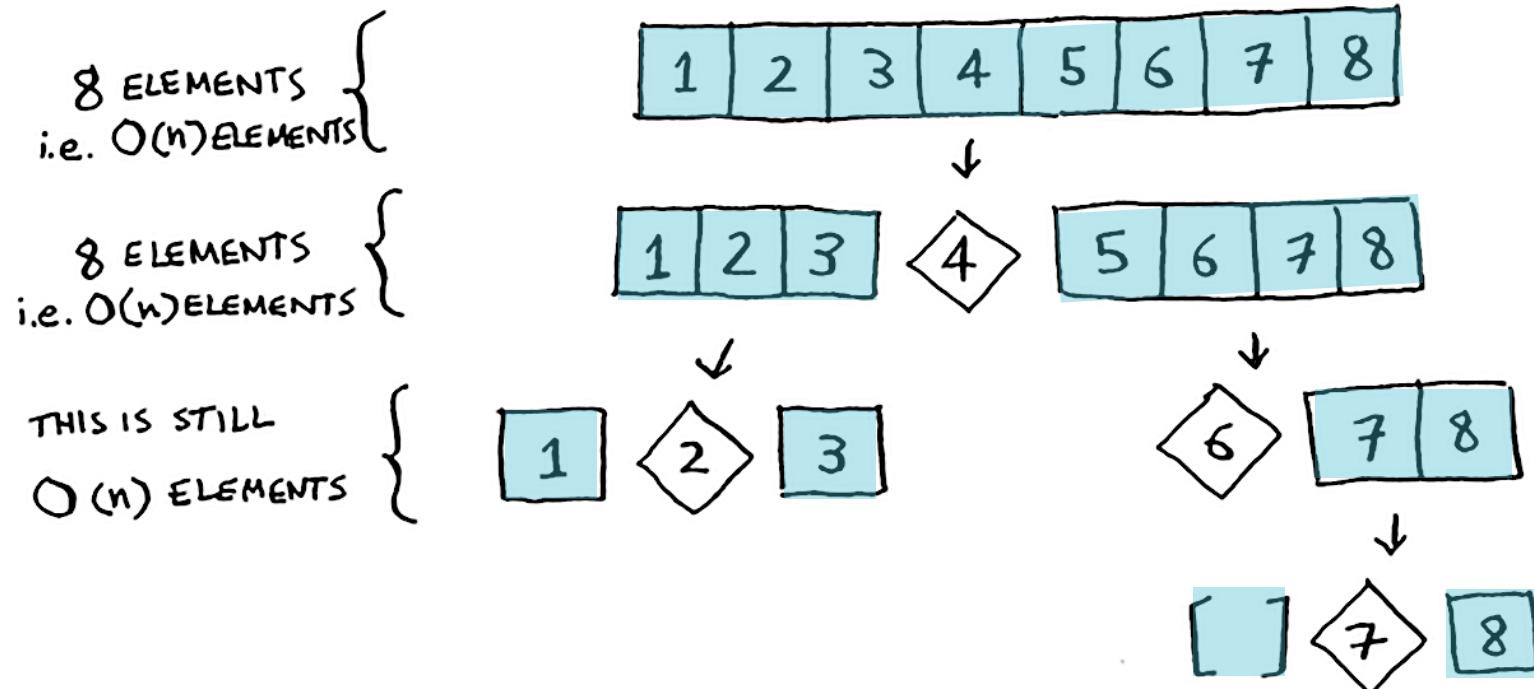
THIS IS STILL
 $O(n)$ ELEMENTS {



- ❖ Even if you partition the array differently, you're still touching **$O(n)$ elements every time.**

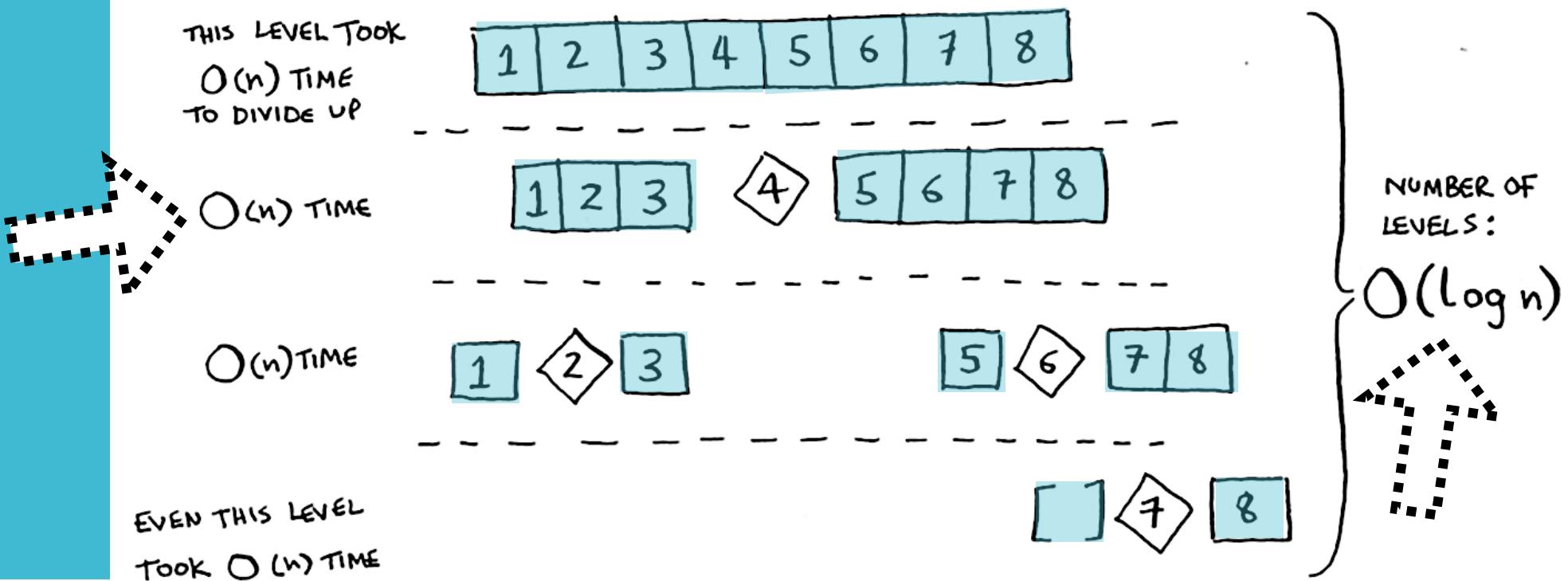
Quick Sort

- ❖ Even if you partition the array differently, you're still touching **O(n)** elements every time.



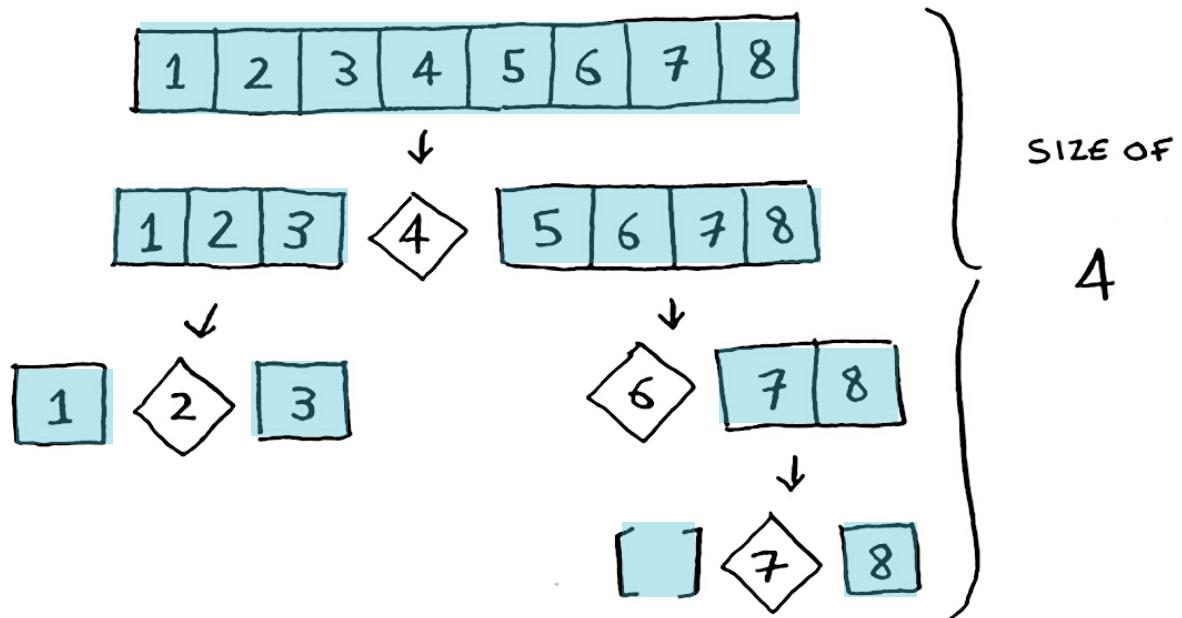
Quick Sort

- ❖ Even if you partition the array differently, you're still touching $O(n)$ elements every time.



Quick Sort

- ❖ It's much **quicker** now!
- ❖ Because the **array is divided into half** every time and no need to make many calls.
- ❖ Base case is hit **sooner**.



Quick Sort

- ❖ 1st example is the **worst-case scenario**:
 - $O(n*n) = O(n^2)$
 - all elements are checked and this operation takes **$O(n)$ time**, repeated “n” times.

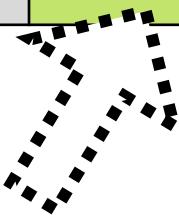
- ❖ 2nd example is the **best-case scenario**:
 - $O(n \log n)$

- ❖ All elements are checked and this operation takes **$O(n)$ time**. And this is repeated “ $\log n$ ” times.

Quick Sort

- ❖ Performance of quicksort heavily depends on the chosen **pivot**.

	Best case	Average case	Worse case
Quick Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$



Quick Sort

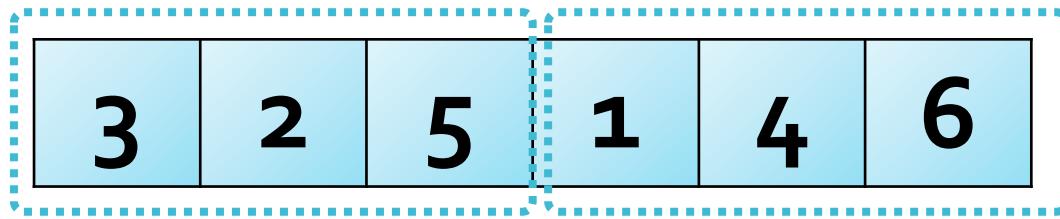
	Best case	Average case	Worse case
Bubble Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Insertion Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
Quick Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$

Merge Sort

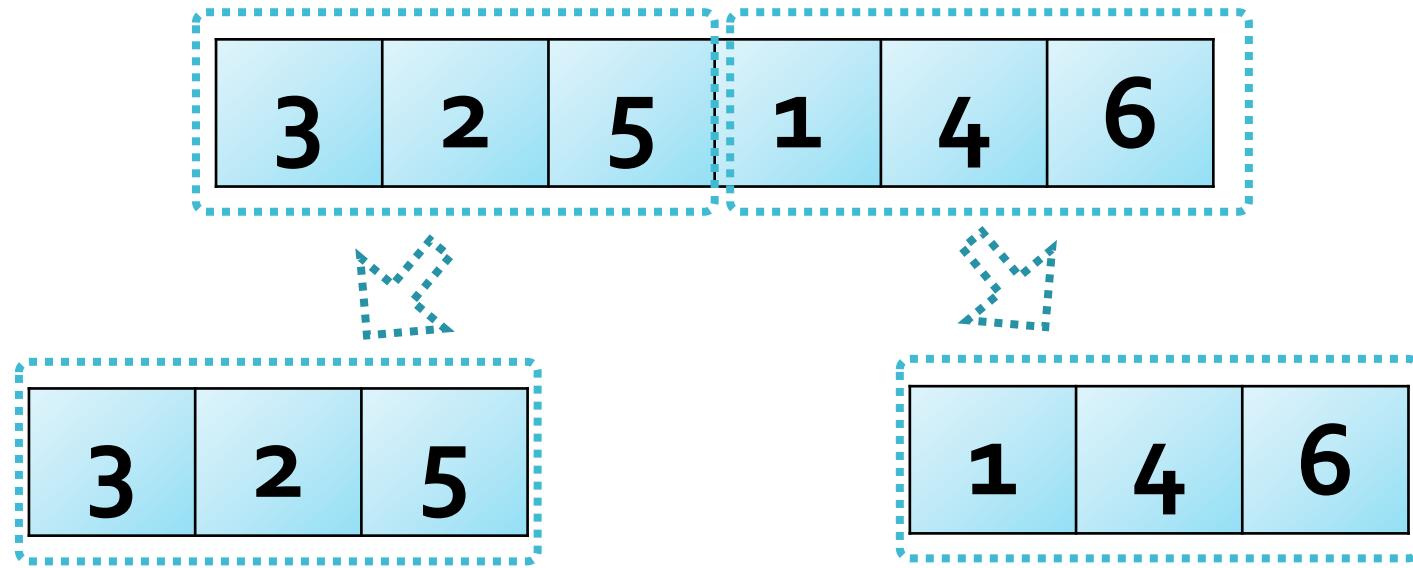
- ❖ **Merge Sort** is a sorting algorithm developed based on uses **Divide & Conquer**.
- ❖ It is faster than other sorts (e.g., selection sort).
- ❖ Developed by Hungarian-American computer scientist **John von Neumann** in 1945.



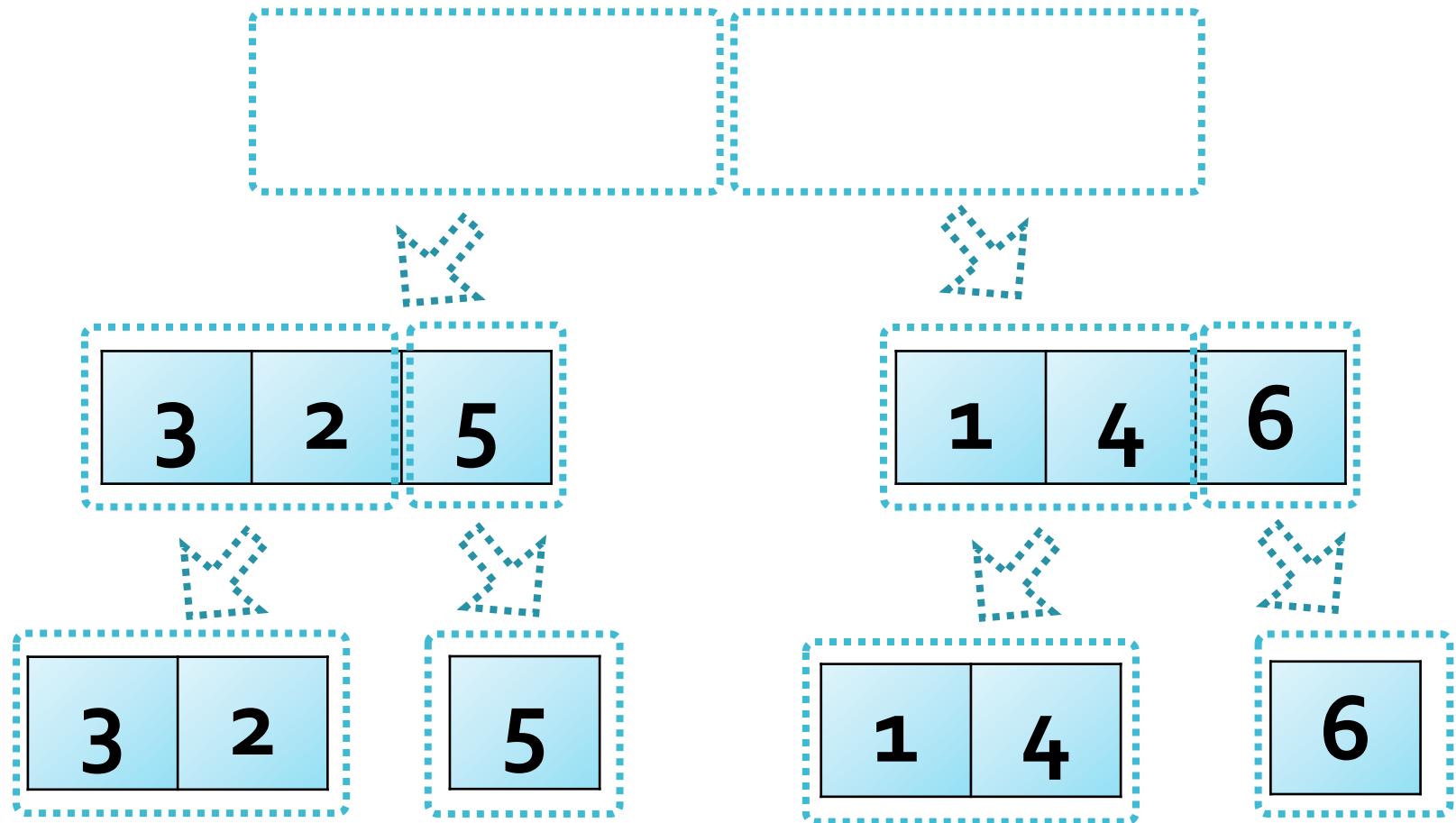
Merge sort



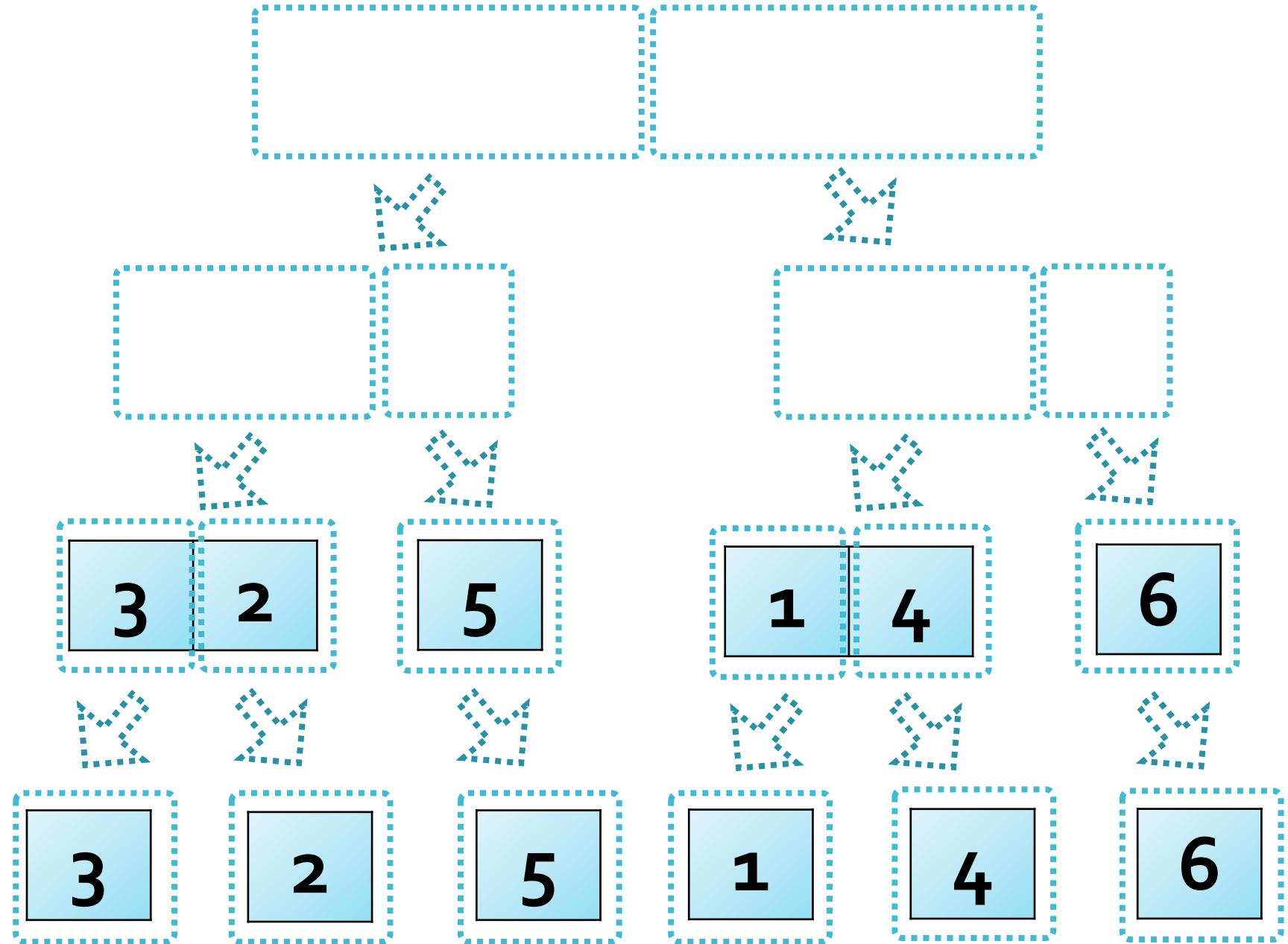
Merge sort



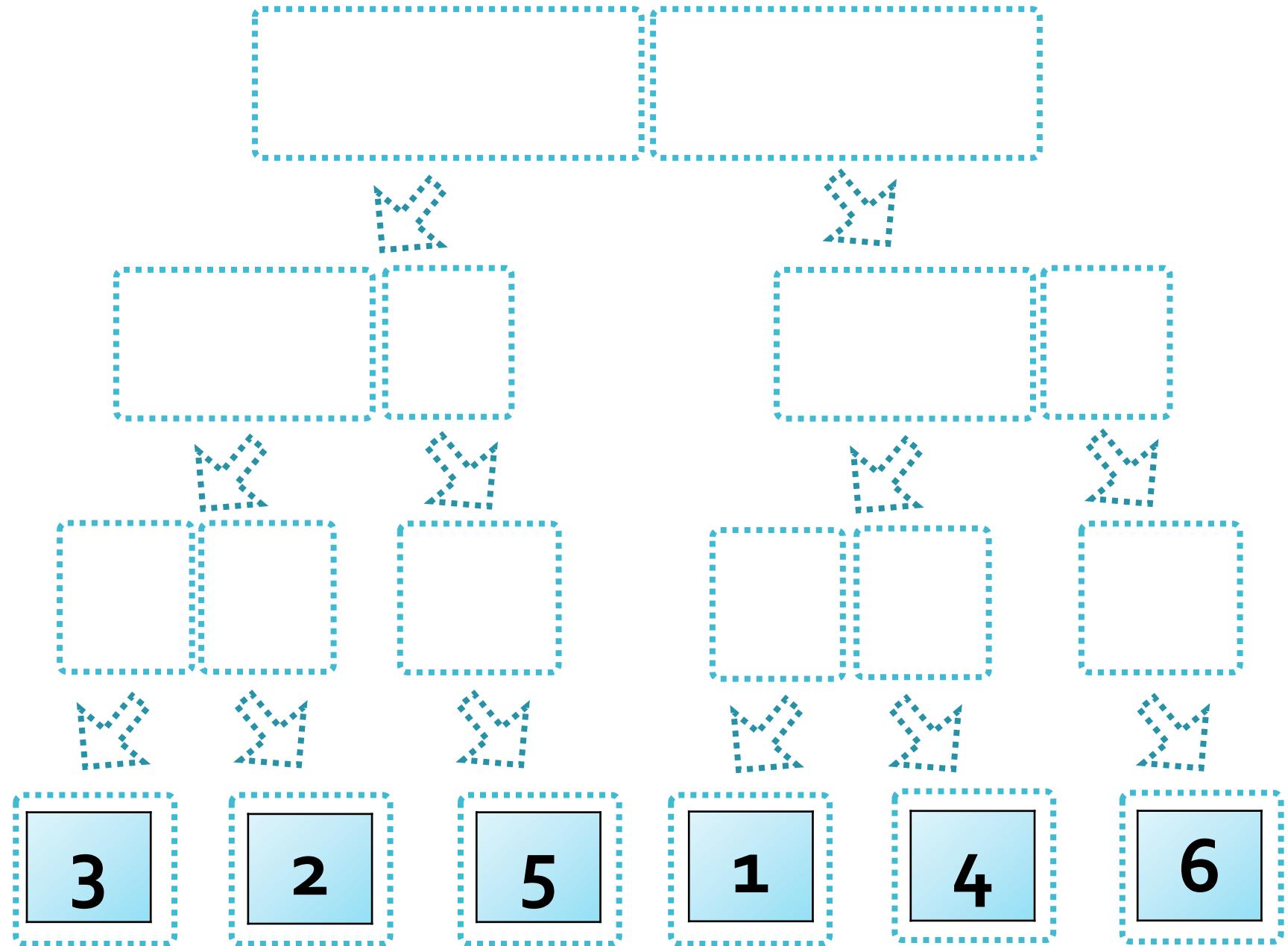
Merge sort



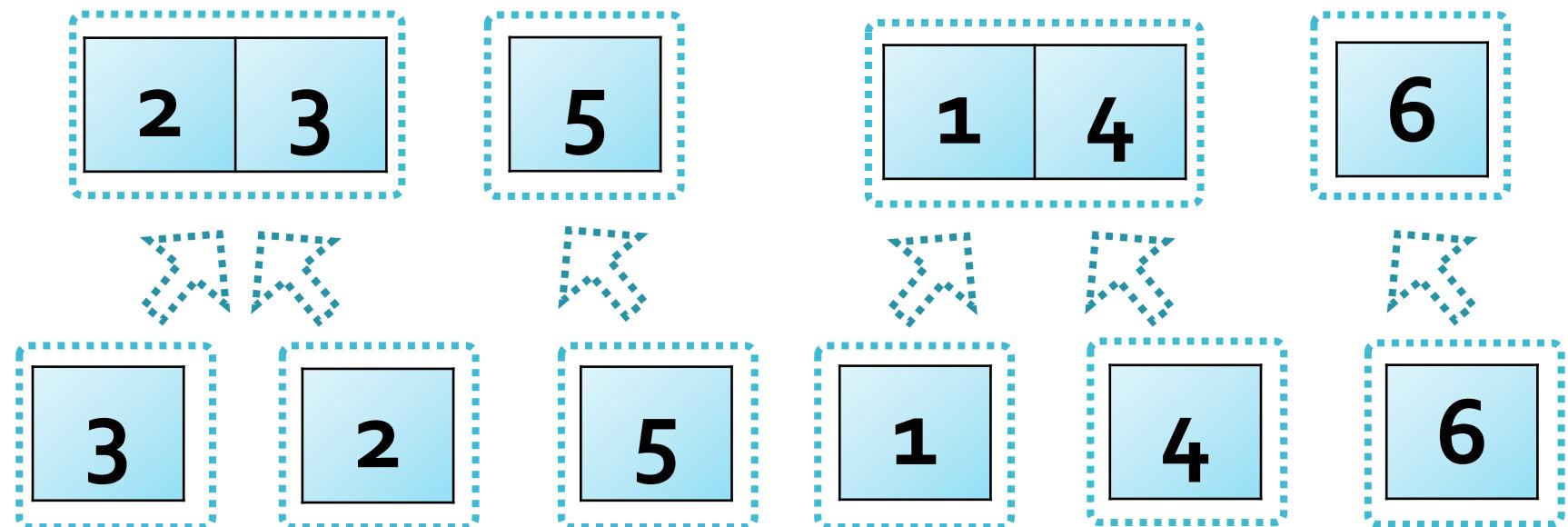
Merge sort



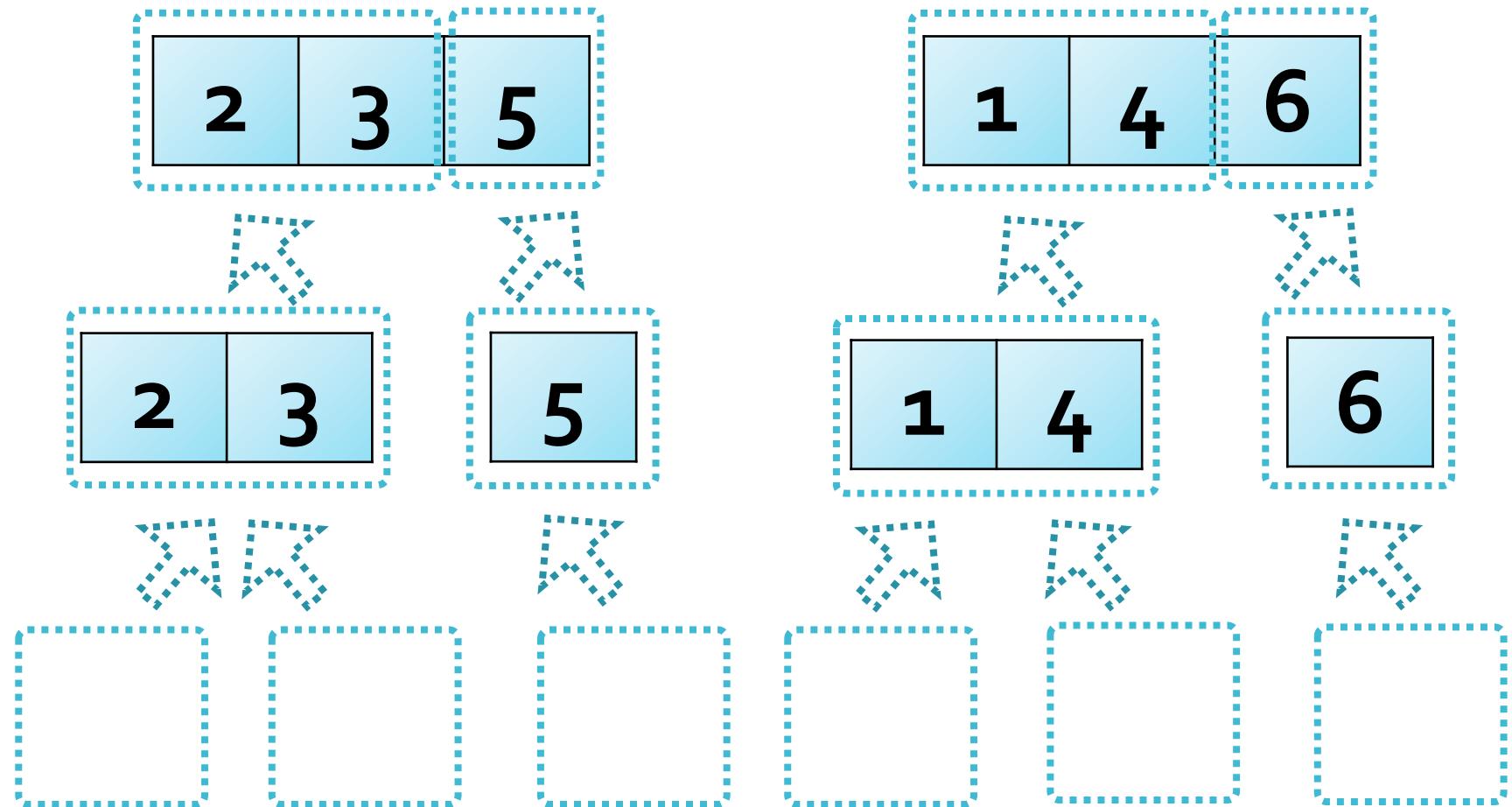
Merge sort



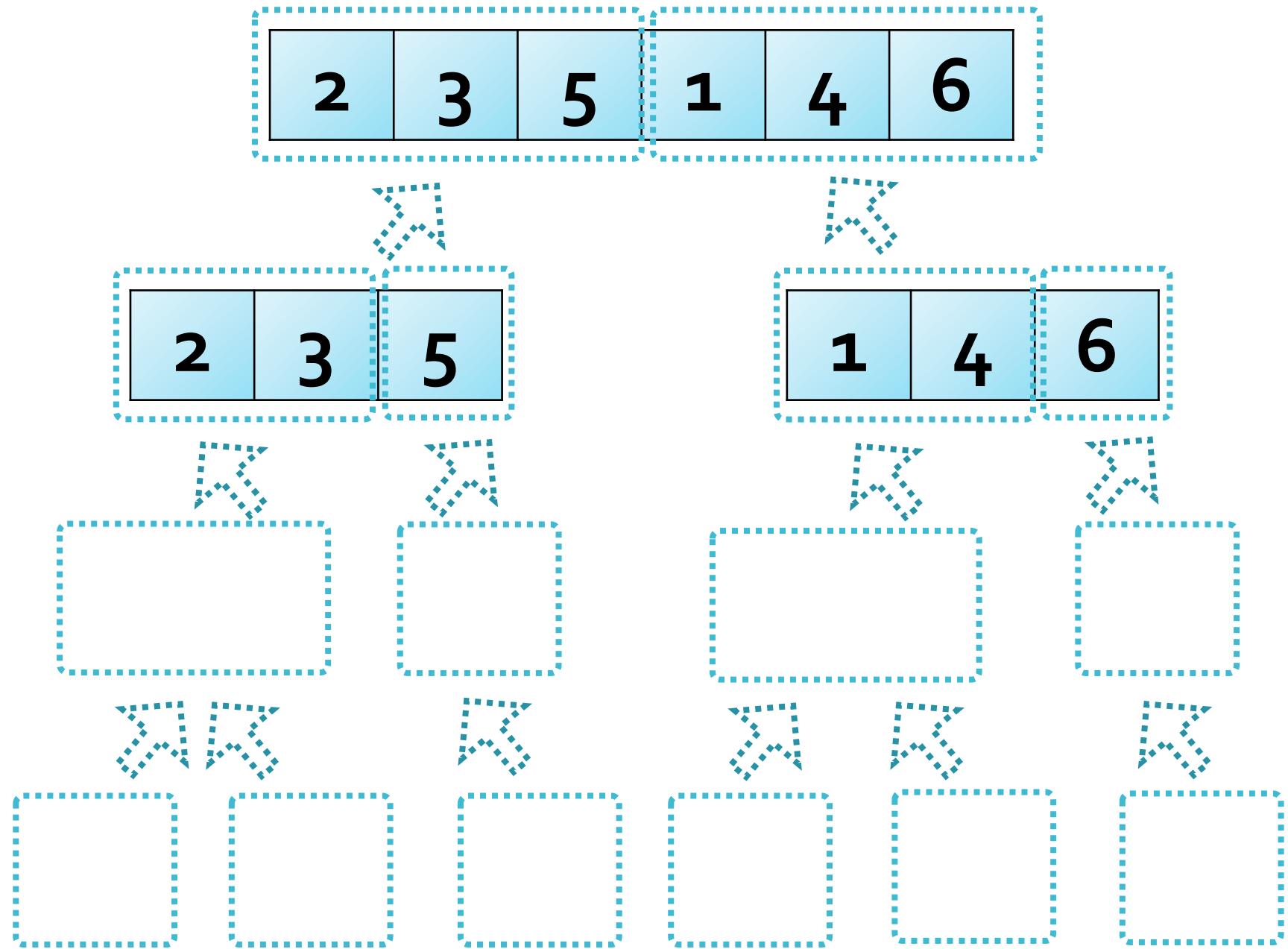
Merge sort



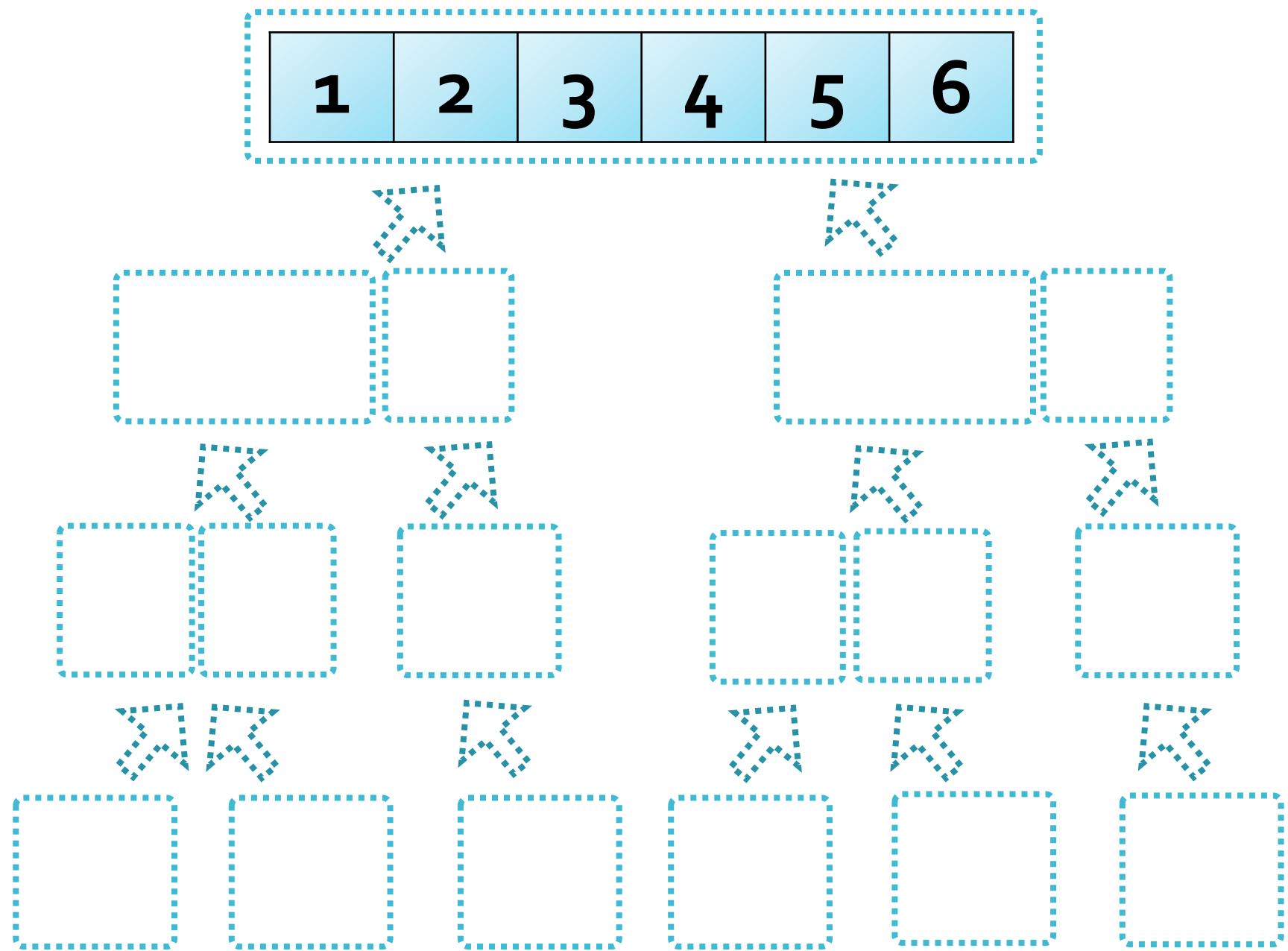
Merge sort



Merge sort



Merge sort



Merge Sort (part 1)

❖ Implementation of the **Merge sort**:

Part 1

```
def merge_sort(arr):  
  
    if len(arr) > 1:  
        mid = len(arr) // 2  
        left = arr[:mid]  
        right = arr[mid:]  
        merge_sort(left)  
        merge_sort(right)
```

Merge Sort (part 2)

Part 2

```
i = j = k = 0
while i < len(left) and j < len(right):

    if left[i] < right[j]:
        arr[k] = left[i]
        i += 1
    else:
        arr[k] = right[j]
        j += 1
    k += 1

while i < len(left):
    arr[k] = left[i]
    i += 1
    k += 1

while j < len(right):
    arr[k] = right[j]
    j += 1
    k += 1
```

Merge Sort (part 3)

```
i = j = k = 0
while i < len(left) and j < len(right):

    if left[i] < right[j]:
        arr[k] = left[i]
        i += 1
    else:
        arr[k] = right[j]
        j += 1
    k += 1

while i < len(left):
    arr[k] = left[i]
    i += 1
    k += 1

while j < len(right):
    arr[k] = right[j]
    j += 1
    k += 1
```

Merge sort

	Best case	Average case	Worse case
Bubble Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Insertion Sort	$O(n)$	$O(n^2)$	$O(n^2)$
Selection Sort	$O(n^2)$	$O(n^2)$	$O(n^2)$
Quick Sort	$O(n \log n)$	$O(n \log n)$	$O(n^2)$
Merge Sort	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$

Merge sort vs Quick Sort



ref: Youtube

Quiz

❖ Given the following list:

[21, 1, 26, 45, 29, 28, 2, 9, 16, 49, 39, 27, 43, 34, 46, 40]

❖ which of the following is the list after **three recursive calls** to Merge sort?

- 1) [16, 49, 39, 27, 43, 34, 46, 40]
- 2) [21, 1]
- 3) [21, 1, 26, 45]
- 4) [21]

❖ which of the following is the list after **three recursive calls** to Merge sort?

[21, 1, 26, 45, 29, 28, 2, 9, 16, 49, 39, 27, 43, 34, 46, 40]

Answer

- 1) [16, 49, 39, 27, 43, 34, 46, 40]
- 2) [21, 1]
- 3) [21, 1, 26, 45]
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❖ which of the following is the list after **three recursive calls** to Merge sort?

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Answer

- 1) [16, 49, 39, 27, 43, 34, 46, 40]
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❖ which of the following is the list after **three recursive calls** to Merge sort?

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Answer

- 1) [16, 49, 39, 27, 43, 34, 46, 40]
- 2)  [21, 1]
- 3) [21, 1, 26, 45]
- 4) [21]

Next Lesson

- Recursion