Advanced Programming INFO135

Lecture 9: Greedy Algorithms

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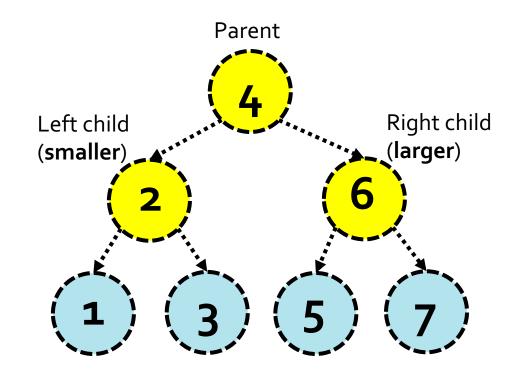
In the previous lecture we introduced **binary tree** structure.

- A special case of a binary tree that follows a specific **property** is called **Binary Search Tree**.
- Data stored in Binary Search Tree is more efficient to find (i.e., to search) than in an ordinary binary trees.

- **Binary Search Tree (BST)** is a binary tree where:
 - *values that are less than parent node are in **left** subtree
 - *values that are greater than parent are in **right** subtree
- This is the property of Binary Search Tree. Value stored in a node is also called Key.
- Duplicates are **not allowed** in binary search tree (similar to Python Set).

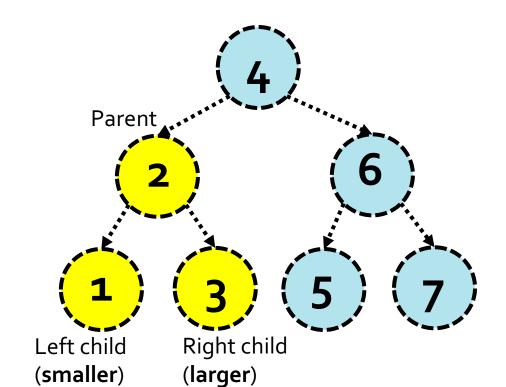
Example:

- check the highlighted nodes of this binary tree:
 - value of parent (4) is larger than left child (2)
 - value of parent (4) is smaller than right child (6)

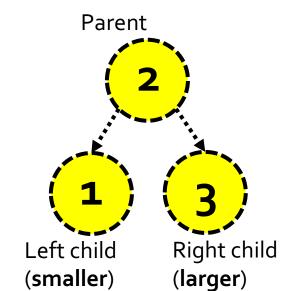


Example:

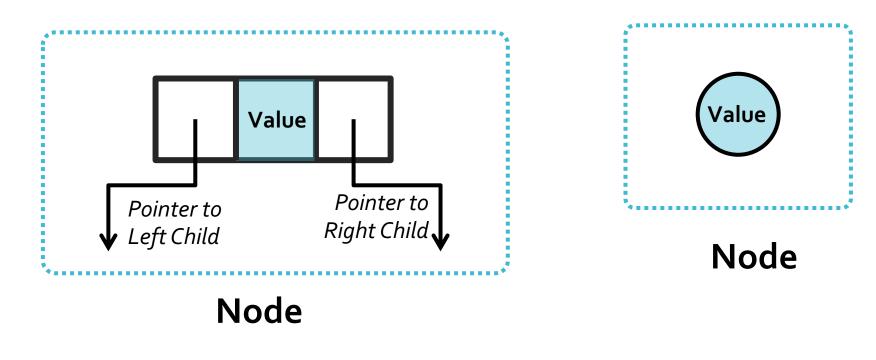
- check the highlighted nodes of this binary tree:
 - value of parent (2) is larger than left child (1)
 - value of parent (2) is smaller than right child (3)



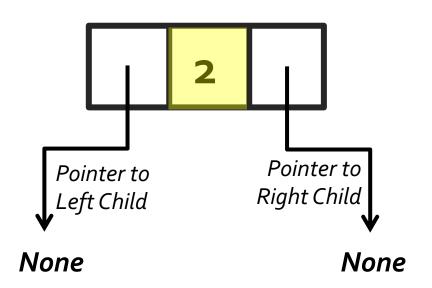
Lets check how this **Binary Search Tree** can be formed.

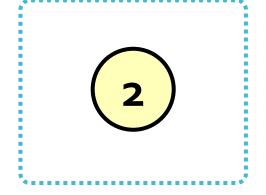


*We define a Node structure, that has a value and a "pointer" to the left child & a "pointer" to right child.



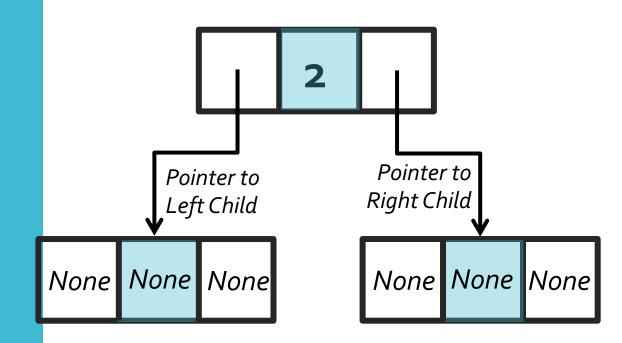
First node with value 2 is created, but it does not have any (left or right) child node.





No child can be represented with an **empty node** where values are *None*.

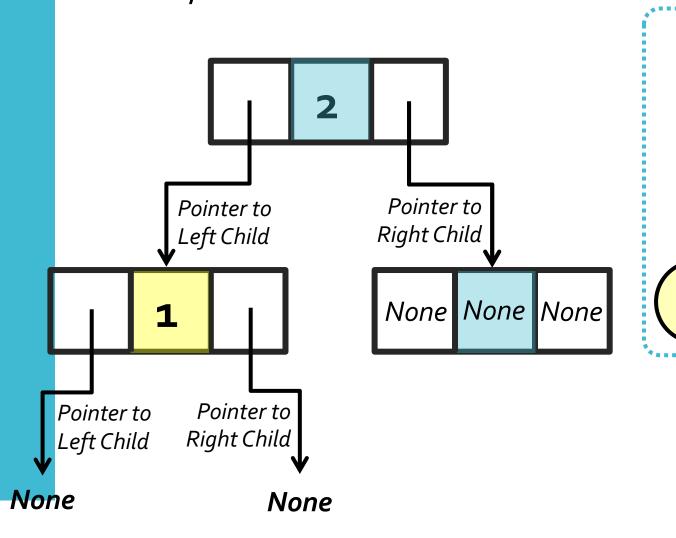


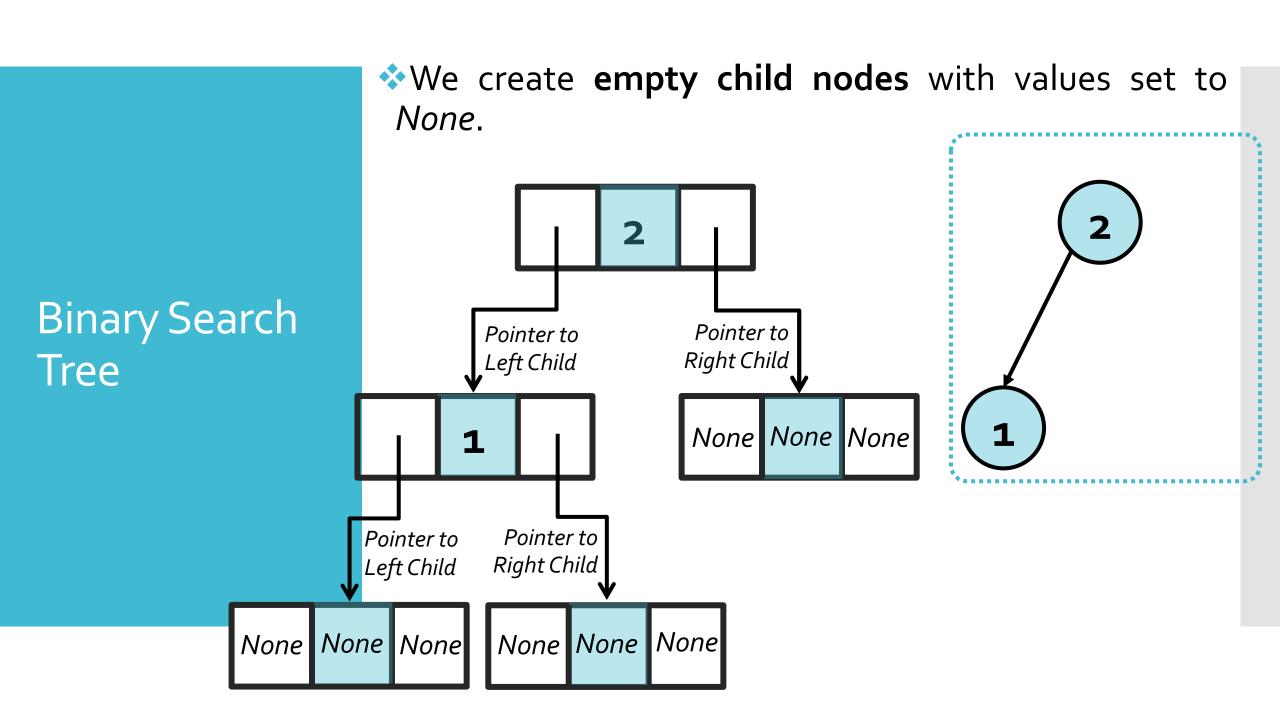




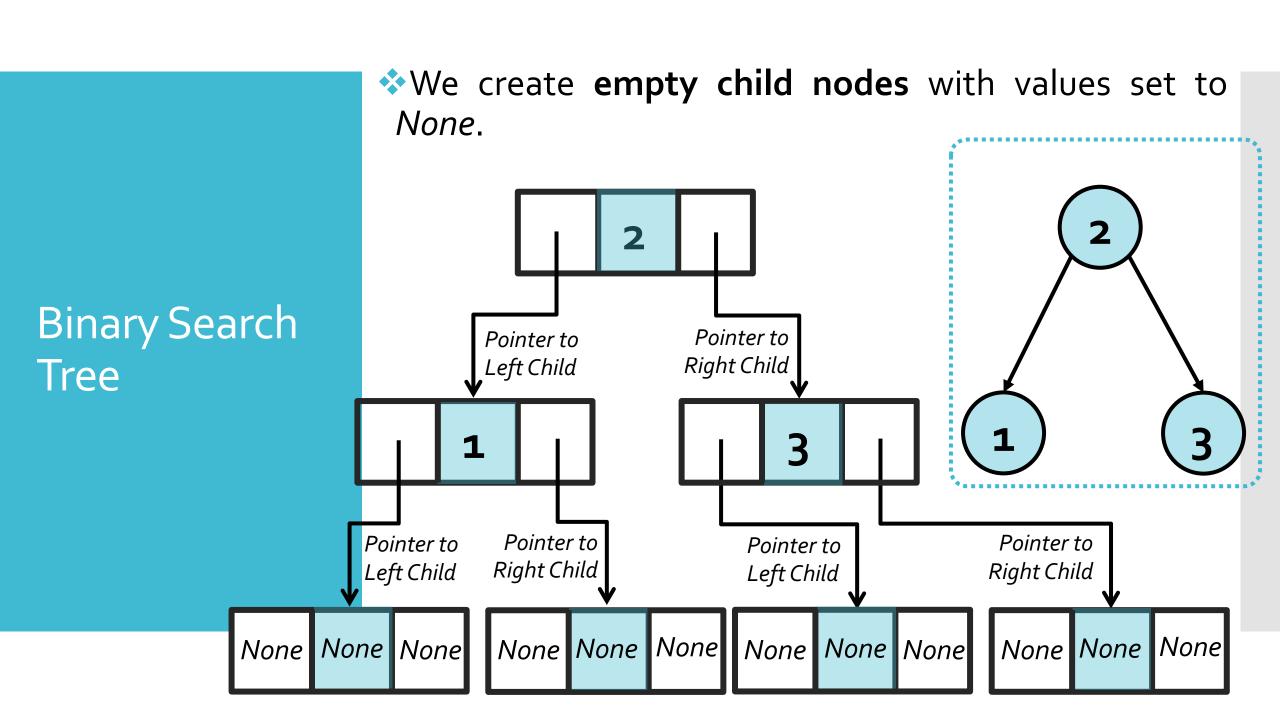
Next node with value 1 is created (to insert into the tree). This node has a value smaller than the first node and so, it is added as the left child.

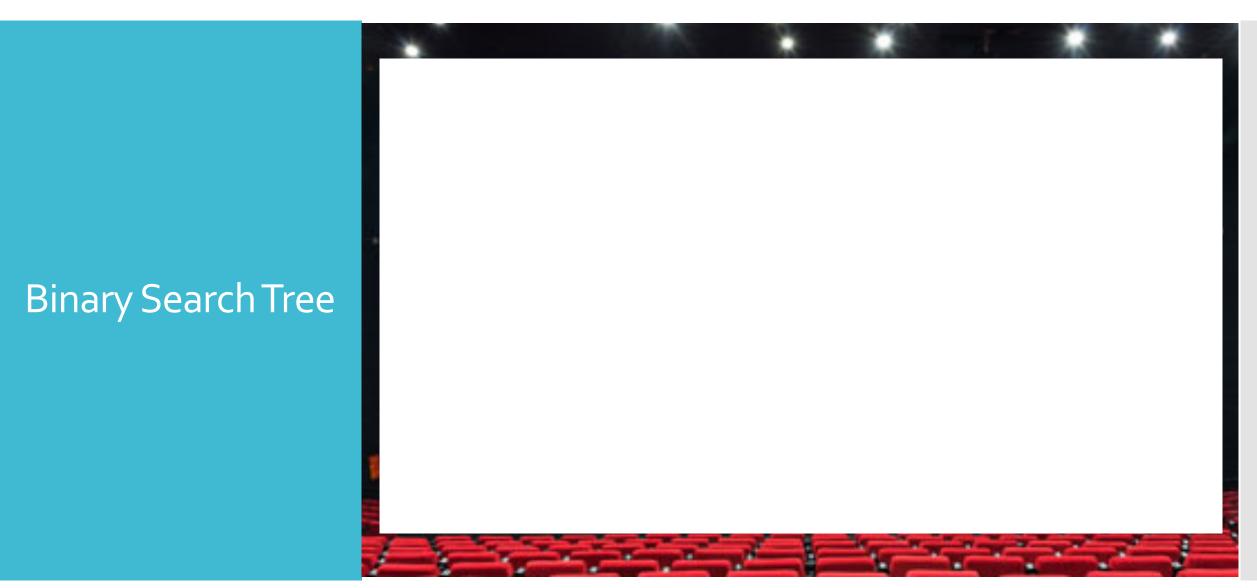
Binary Search Tree





The next is value 3 & we create a node for it (to insert it into tree). This node has a value larger than root node (2) and so, it is added as the right child. Binary Search Pointer to Pointer to Left Child Right Child Tree Pointer to Pointer to Pointer to Pointer to Right Child Right Child Left Child Left Child None None None None None None None None

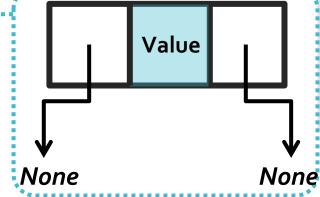




ref : youtu.be/7vw2iIdqHlM

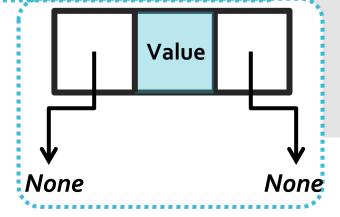
Implementing Binary Search Tree class (part 1)

```
class BinarySearchTree:
    def init (self, value=None):
        self.value = value
        if self.value:
            self.left child = BinarySearchTree()
            self.right child = BinarySearchTree()
        else:
            self.left child = None
            self.right child = None
```



Implementing Binary Search Tree class (part 1)

```
class BinarySearchTree:
   def init (self, value=None):
        self.value = value
        if self.value:
            self.left_child = BinarySearchTree()
            self.right_child = BinarySearchTree()
        else:
            self.left child = None
            self.right child = None
```



constructor

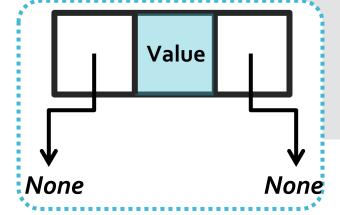
Implementing Binary Search Tree class (part 1)

```
class BinarySearchTree:
```

```
def __init__(self, value=None):
    self.value = value

    if self.value:
        self.left_child = BinarySearchTree()
        self.right_child = BinarySearchTree()

    else:
        self.left_child = None
        self.right_child = None
```

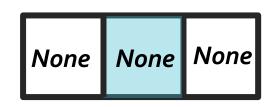


Implementing Binary Search Tree class (part 1)

class BinarySearchTree:

```
def __init__(self, value=None
    self.value = value
    if self.value:
        self.left_child = Binar *SearchTree()
        self.right_child = I** arySearchTree()

    else:
        self.left_child = None
        self.right_child = None
```



Implementing Binary Search Tree class (part 1)

```
class BinarySearchTree:
    def init (self, value=None):
        self.value = value
        if self.value:
            self.left child = BinarySearchTree()
            self.right child = BinarySearchTree()
        else:
            self.left child = None
            self.right child = None
    def is empty(self):
                                         checks if the
        return self.value is None
                                       node is empty
```

Implementing Binary Search Tree class (part 2)

```
def insert(self, value):
    if self.is empty():
        self.value = value
        self.left child = BinarySearchTree()
        self.right child = BinarySearchTree()
    elif value < self.value:</pre>
        self.left child.insert(value)
    elif value > self.value:
        self.right child.insert(value)
```

Implementing Binary Search Tree class (part 2)

```
if node is empty,
def insert(self, value):
                                    it adds a new
    if self.is empty():
                                            node
        self.value = value
        self.left child = BinarySearchTree()
        self.right child = BinarySearchTree()
    elif value < self.value:</pre>
        self.left child.insert(value)
    elif value > self.value:
        self.right child.insert(value)
```

Implementing Binary Search Tree class (part 2)

```
if new value is less
def insert(self, value):
                                 than current node,
    if self.is empty():
                                  insert to left side
        self.value = value
        self.left_child = Bina ySearchTree()
        self.right child = ** marySearchTree()
    elif value < self.value:</pre>
        self.left child.insert(value)
    elif value > self.value:
        self.right child.insert(value)
```

Implementing Binary Search Tree class (part 2)

```
def insert(self, value):
    if self.is empty():
        self.value = value
        self.left child = BinarySearchTree()
        self.right child = BinarySearchTree()
    elif value < self.value:</pre>
        self.left child.insert(value)
    elif value > self.value:
        self.right_child.insert(value)
```

if new value is larger than current node insert to right side

Implementing Binary Search Tree class (part 3)

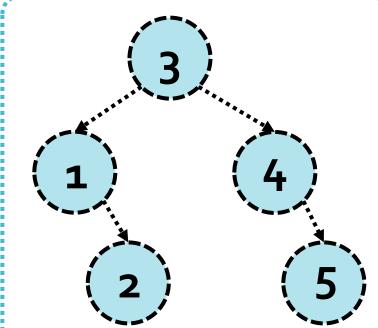
```
def in order(self):
    if self.is empty():
        return []
    else:
        return self.left child.in order() +
                [self.value] + \
                self.right child.in order()
def print_tree(self):
                                        recursively
                                  traverses the tree
    print(self.in ord
```

Implementing Binary Search Tree class (part 3)

Testing the implementation.

```
my tree = BinarySearchTree()
my tree.insert(3)
my tree.insert(1)
my tree.insert(4)
my tree.insert(2)
my tree.insert(5)
my tree.print tree()
[Output:]
```

[1, 2, 3, 4, 5]



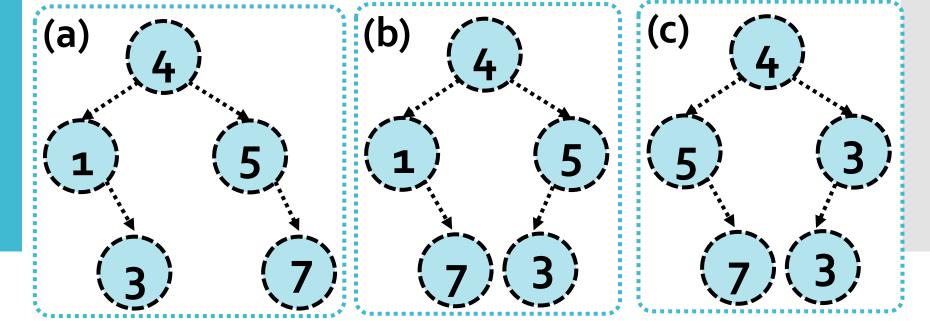
Which Tree does the following code creates?



```
Quiz
```

```
my_tree = BinarySearchTree()

my_tree.insert(4)
my_tree.insert(1)
my_tree.insert(5)
my_tree.insert(3)
my_tree.insert(7)
```



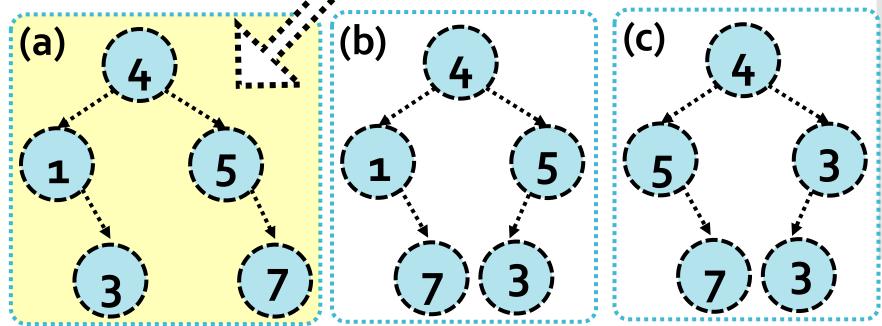
Which Tree does the following code creates?



```
Answer
```

```
my_tree = BinarySearchTree()

my_tree.insert(4)
my_tree.insert(1)
my_tree.insert(5)
my_tree.insert(3)
my_tree.insert(7)
```



*We need a **find(value)** method that searches the tree [-may be- recursively?] in case whether or not a value is found.

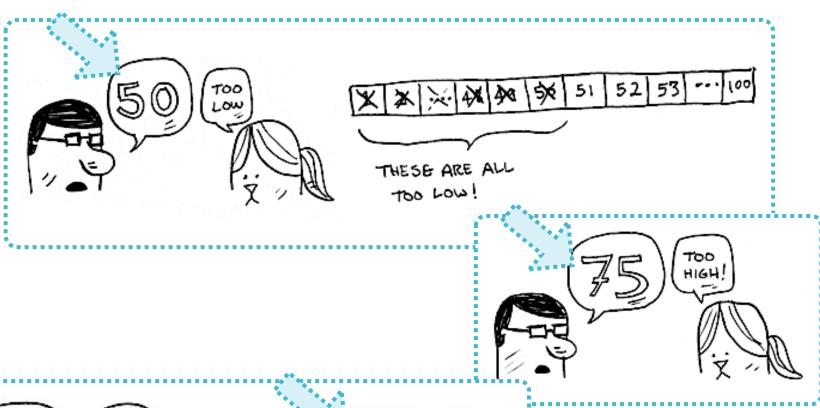
- Then we may ended up with these conditions:
 - we may **find** a node with matching value.
 - we may reach a non-matching **leaf node**.

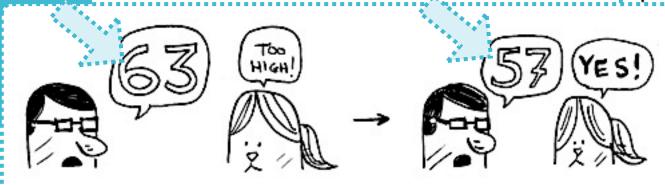
- \diamond Suppose we are searching for value k, then we:
 - a) check the value of **current node (v)**
 - b) move to the **left** child if k < v
 - c) move to the **right** child if k > v

Does it sound familiar?

Do you remember the game?

Binary Search Tree





Implementing Binary Search Tree class (part 4)

Binary Search Tree

```
def find(self, value):
    if self.is empty():
        return False
    elif value == self.value:
        return True
    elif self.value > value:
        return self.left child.find(value)
    elif self.value < value:</pre>
        return self.right child.find(value)
```

Implementing Binary Search Tree class (part 4)

```
def find(self, value):
    if self.is_empty(): if current node is empty
        return False
    elif value == self.value:
        return True
    elif self.value > value:
        return self.left child.find(value)
    elif self.value < value:</pre>
        return self.right child.find(value)
```

Implementing Binary Search Tree class (part 4)

Binary Search Tree

```
def find(self, value):
                                     if current value
    if self.is empty():
                                      is what we are
        return False
                                      searching for
    elif value == self.value:
        return True
    elif self.value > value:
        return self.left child.find(value)
    elif self.value < value:</pre>
        return self.right child.find(value)
```

Implementing Binary Search Tree class (part 4)

```
def find(self, value):
    if self.is empty():
        return False
    elif value == self.value:
                                    if current value
                                is larger than
        return True
    elif self.value > value: ...
        return self.left child.find(value)
    elif self.value < value:</pre>
        return self.right child.find(value)
```

Implementing Binary Search Tree class (part 4)

```
def find(self, value):
    if self.is empty():
        return False
    elif value == self.value:
        return True
                                     if current value
    elif self.value > value:
                                      is smaller than
        return self.left_child,fin, search value
    elif self.value < value:</pre>
        return self.right child.find(value)
```

Testing the implementation.

```
my tree = BinarySearchTree()
my tree.insert(3)
my tree.insert(1)
my tree.insert(4)
my tree.insert(2)
my tree.insert(5)
print('Found 4?')
print(my_tree.find(4))
[Output:]
  Found 4?
  True
```

- •We can extend the Binary Search Tree by implementing more methods.
- May be checking if we reached a leaf note.
- *We can also implement a **copy method** for copy (for instance coping a child node).
- And also a **delete** method.

```
def is leaf(self):
    return self.left child.is empty() and \
           self.right_child.is_empty()
def make empty(self):
    self.value = None
    self.left child = None
    self.right child = None
def copy_child(self, child):
    if child == 'left':
        self.value = self.left child.value
        self.right child = self.left child.right child
        self.left child = self.left child.left child
    elif child == 'right':
        self.value = self.right child.value
        self.left child = self.right child.left child
        self.right child = self.right_child.right_child
```

Implementing Binary Search Tree class (part 5)

```
def is leaf(self):
    return self.left child.is empty() and \
           self.right child.is empty()
def make empty(self):
    self.value = None
                                         But why?
    self.left child = None
                                         We will need it later!
    self.right child = None
def copy_child(self, child):
    if child == 'left':
        self.value = self.left child.value
        self.right child = self.left child.right child
        self.left child = self.left child.left child
    elif child == 'right':
        self.value = self.right child.value
        self.left child = self.right child.left child
```

self.right child = self.right child.right child

```
def is_leaf(self):
    return self.left_child.is_empty() and \
        self.right_child.is_empty()

def make_empty(self):
    self.value = None
    self.left_child = None
    self.right_child = None
```

```
def copy_child(self, child):
    if child == 'left':
        self.value = self.left_child.value
        self.right_child = self.left_child.right_child
        self.left_child = self.left_child.left_child
    elif child == 'right':
        self.value = self.right_child.value
        self.left_child = self.right_child.left_child
        self.right_child.left_child
```

```
def delete(self, value):
    if self.is_empty():
        print('Binary tree is empty.')
    elif value < self.value:</pre>
        self.left child.delete(value)
    elif value > self.value:
        self.right child.delete(value)
    elif value == self.value:
        if self.is_leaf():
            self.make empty()
        elif self.left child.is empty():
            self.copy child('right')
        else:
            self.value = self.left child.delete max()
```

```
def delete(self, value):
    if self.is empty():
        print('Binary tree is empty.')
    elif value < self.value:</pre>
        self.left child.delete(value)
    elif value > self.value:
        self.right child.delete(value)
    elif value == self.value:
        if self.is_leaf():
            self.make empty()
        elif self.left child.is empty():
            self.copy child('right')
        else:
            self.value = self.left child.delete max()
```

```
def delete(self, value):
    if self.is empty():
        print('Binary tree is empty.')
    elif value < self.value:</pre>
        self.left child.delete(value)
    elif value > self.value:
        self.right child.delete(value)
    elif value == self.value:
        if self.is_leaf():
            self.make empty()
        elif self.left child.is empty():
            self.copy child('right')
        else:
            self.value = self.left child.delete max()
```

```
def delete(self, value):
    if self.is empty():
        print('Binary tree is empty.')
    elif value < self.value:</pre>
        self.left child.delete(value)
    elif value > self.value:
        self.right child.delete(value)
    elif value == self.value:
        if self.is leaf():
            self.make empty()
        elif self.left child.is empty():
            self.copy child('right')
        else:
            self.value = self.left child.delete max()
```

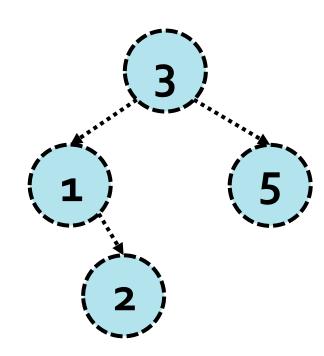
```
def delete max(self):
    if self.right child.is empty():
        max val = self.value
        if self.left child.is empty():
            self.make empty()
        else:
            self.copy child('left')
        return max val
    else:
        return self.right child.delete max()
```

Testing the implementation.

```
my_tree.delete(4)
my_tree.print_tree()

print('Found 4?')
print(my_tree.find(4))
```

```
[Output:] [1, 2, 3, 5]
Found 4?
False
```

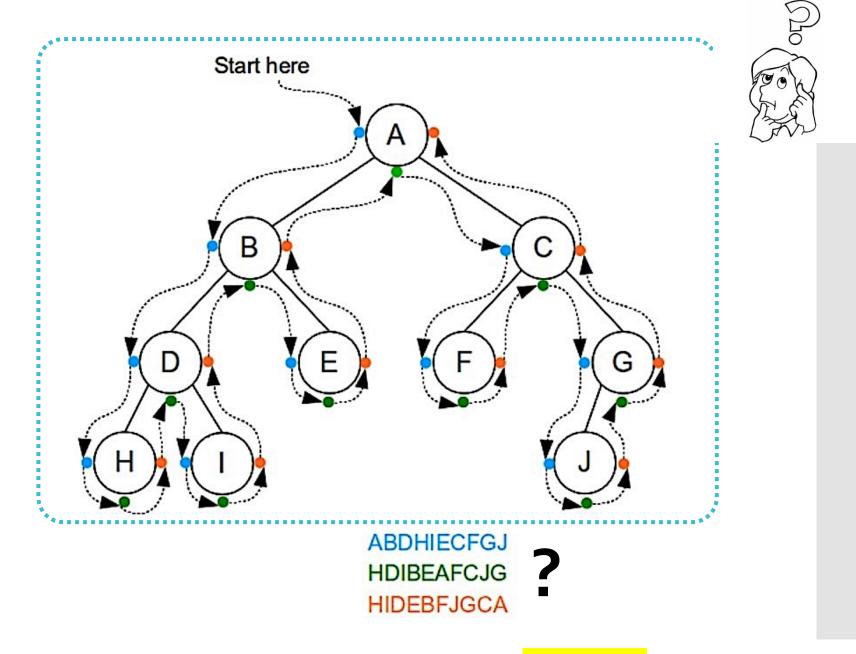


This implementation has one traverse method, i.e., inorder!

❖Inorder recursively do a (inorder) traversal on the left subtree, then visit the root node, and finally do a recursive (inorder) traversal of the right subtree.

But we can have **more** ways of traverse!

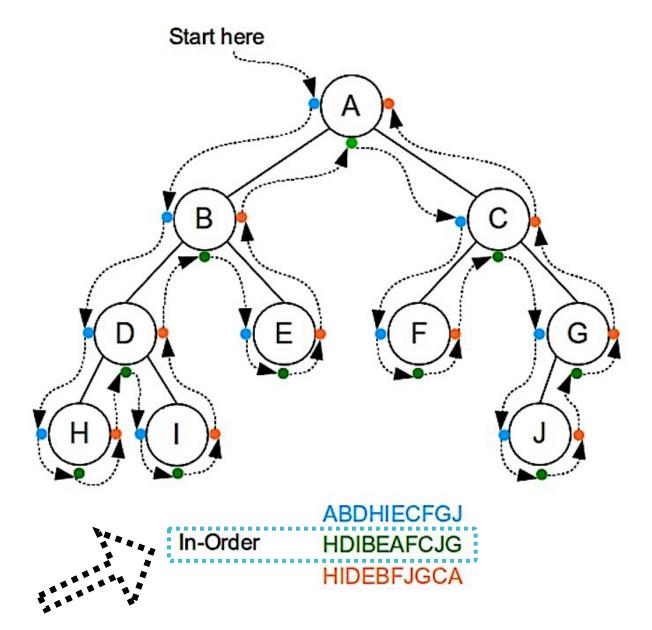
Quiz



Question: Which Color indicates In-order traverse?

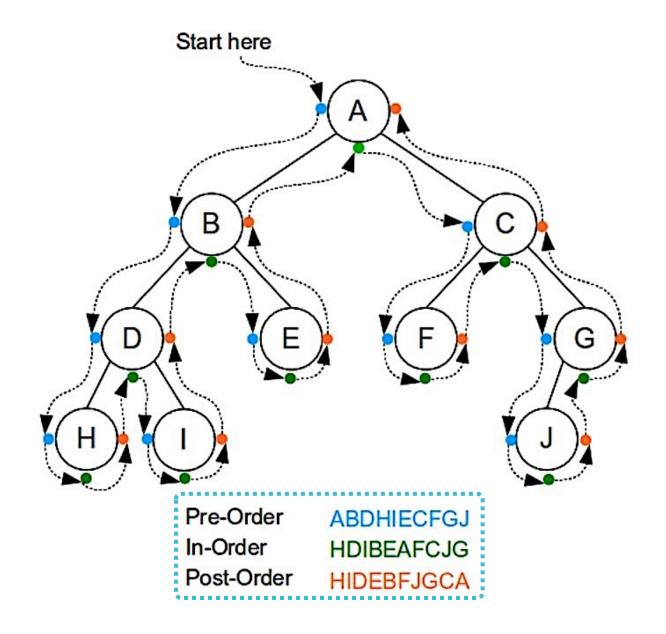


Answer



But what about the **other colors**?

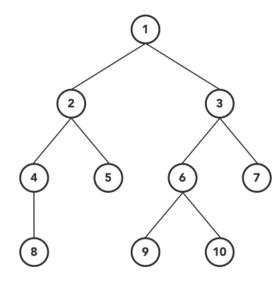
Answer



- ❖ Preorder: visits the root node first, then recursively do a preorder traversal of the left subtree, followed by a recursive preorder traversal of the right subtree.
- ❖ Postorder: recursively do a postorder traversal of the left subtree and the right subtree followed by a visit to the root node.
- Now lets see them **traversing** an example tree!

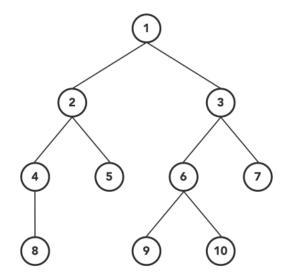
Preorder:

- ❖visit root node,
- ❖go to left-subtree,
- go to right-subtree



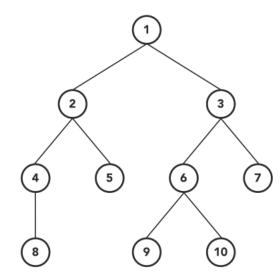
❖Inorder:

- go to left-subtree,
- ❖visit root node,
- go to right-subtree



Postorder:

- go to left-subtree,
- go to right-subtree,
- ❖visit root node



```
class BinarySearchTree:
    def pre_order(self):
        if self.is_empty():
            return []
        else:
            return [self.value] + \
                   self.left child.pre order() + \
                   self.right child.pre order()....
```

Implementation of Preorder and Postorder

Implementation of Preorder and Postorder

Implementation of Preorder and Postorder

Testing our implement

```
my_tree = BinarySearchTree()
my_tree.insert(3)
my_tree.insert(1)
my_tree.insert(4)
my_tree.insert(2)
my_tree.insert(5)

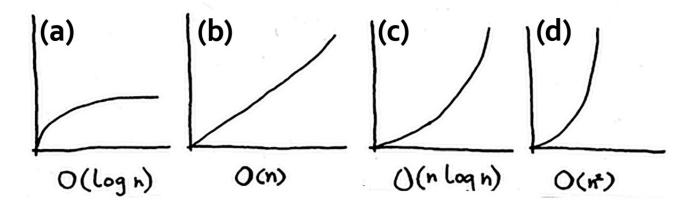
print("Pre-order traversal:", my_tree.pre_order())
print("In-order traversal:", my_tree.in_order())
print("Post-order traversal:", my_tree.post_order())
```

[Output:]

```
Pre-order traversal: [3, 1, 2, 4, 5]
In-order traversal: [1, 2, 3, 4, 5]
Post-order traversal: [2, 1, 5, 4, 3]
```

Quiz

Which of the images represent the Big O notation for Preorder, Inorder, Postorder?

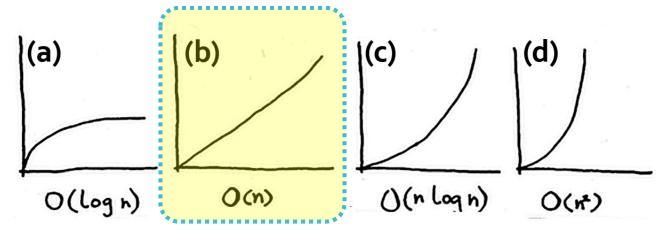


- 2) Which of the the tree traverse could be better:
 - for making a copy of a tree
 - for **deleting a tree** (from leaf to root)



Answer

1) Which of the images represent the Big O notation for **Preorder, Inorder, Postorder**?



> Because we traverse each node **only once**.



Which of the tree traverse could be better:

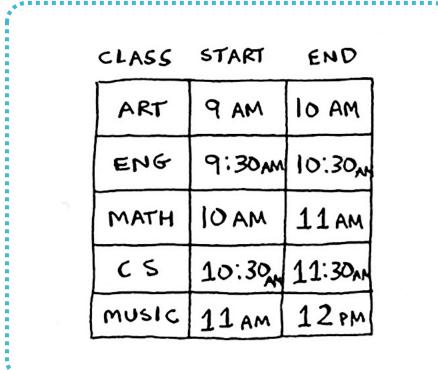
- for making a copy of a tree --> Preorder
- •for **deleting a tree** (from leaf to root) --> Postorder

Answer

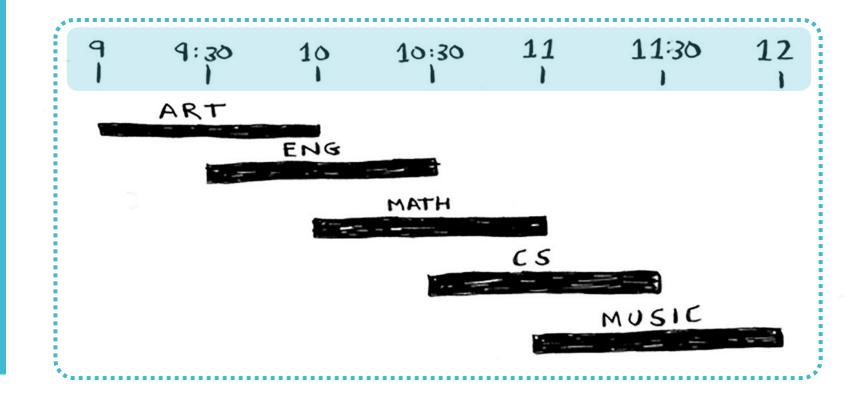


- Greedy Algorithms follow an effective strategy to solve problems.
- Greedy Algorithms: make choices that seem to be the best solution at a moment (i.e., locally-optimal) hoping that it will lead to globally-optimal solution.
- ❖If there is a (objective) function that needs to be optimized (such as cost to be minimized):
 - Greedy algorithm makes choices step-by-step to ensure that the objective function is optimized.

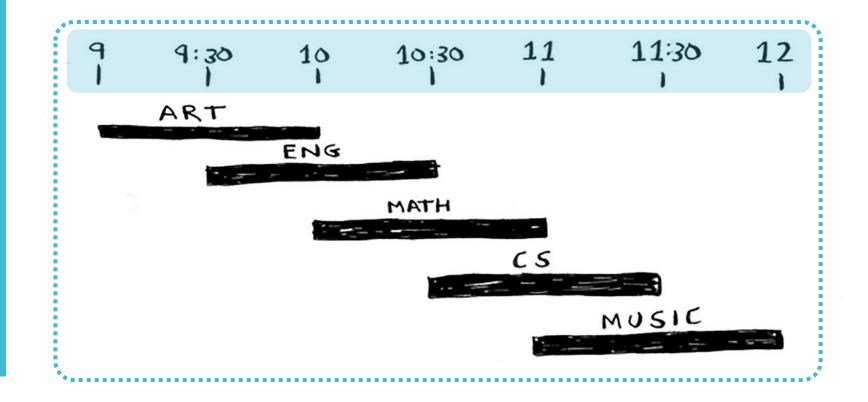
- Suppose you have a classroom and want to hold as many classes here as possible.
- This is the list of possible classes.



❖ You can't hold all of these classes in there, because some of them overlap.



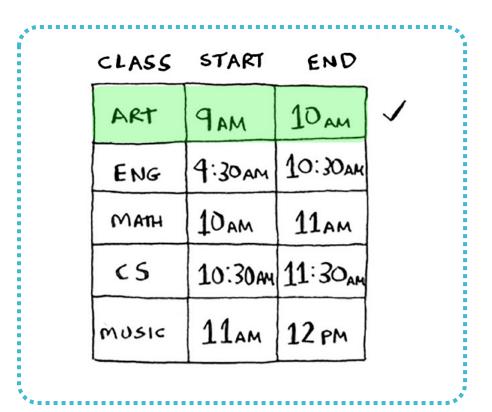
How do you pick what set of classes to hold, so that you get the biggest set of classes possible?



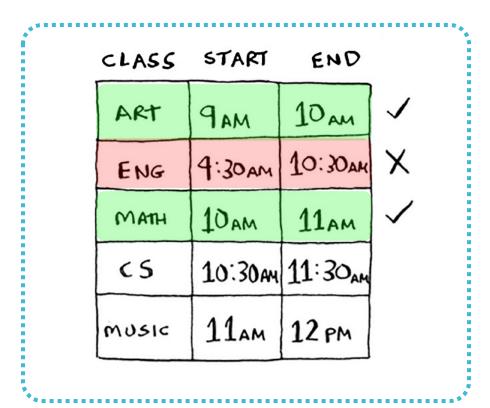
Greedy Algorithms

- The algorithm is the following:
 - 1. First, pick the class that **ends the soonest**. This is the first class you'll hold.

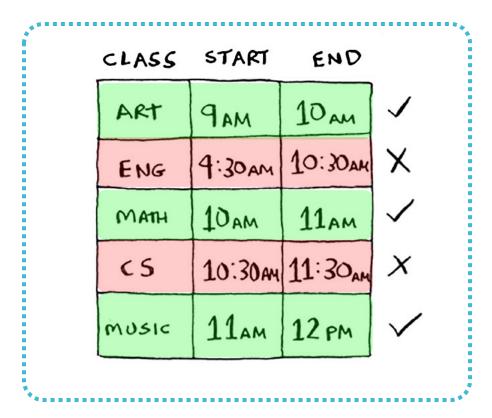
Art ends the soonest (at 10:00 am) and we pick it.



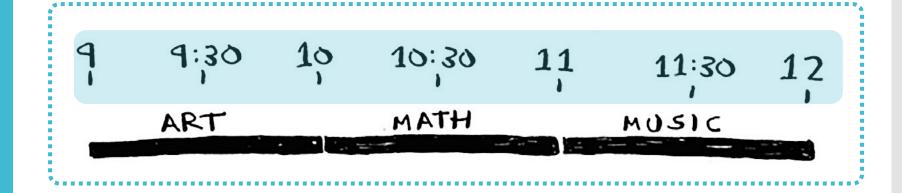
- The algorithm is the following:
 - 2. Then, pick a class that starts after the first class and ends the soonest. This is the second class you'll hold. English is out because it conflicts with Art, but Math works.



- The algorithm is the following:
 - 3. Finally, CS conflicts with Math, but Music works.



So these are the **three classes** you'll hold in this classroom.



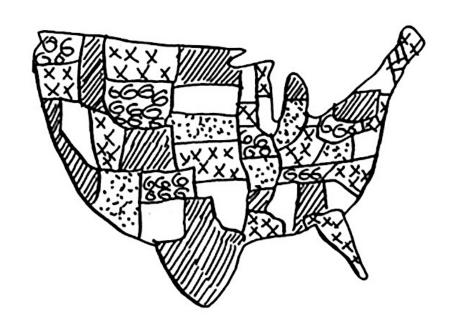
- This was a Greedy Algorithm, since:
 - *at each step, it picks an **optimal move**.
 - in this case, it picks a class that **ends the** soonest.

- In technical terms:
 - at each step, it picks a **locally optimal** solution.
 - at the end, it reaches **globally optimal** solution.

Suppose:

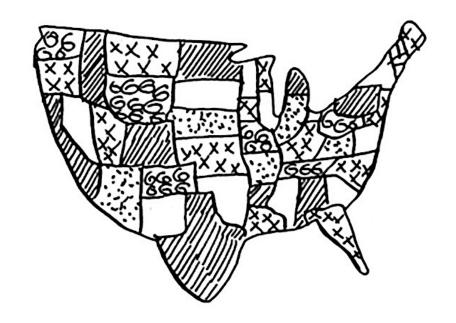
- ❖ You're starting a radio show in USA.
- You want to reach listeners in all states.
- But a station in a new state will cost you.
- ❖You try to **minimize the cost** by choosing stations you will play on.



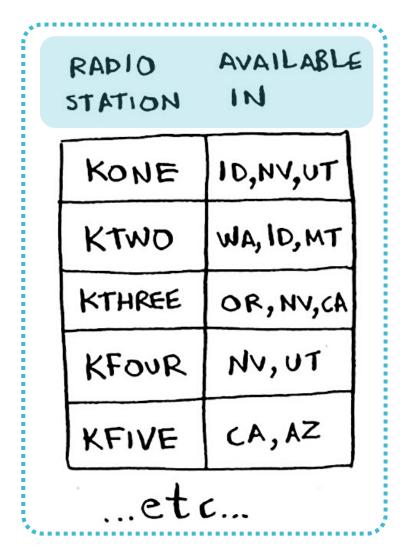


- Each station covers a region, and there's overlap. How to figure out the smallest set of stations you can play on to cover all states?
- Lets see how you can do it.

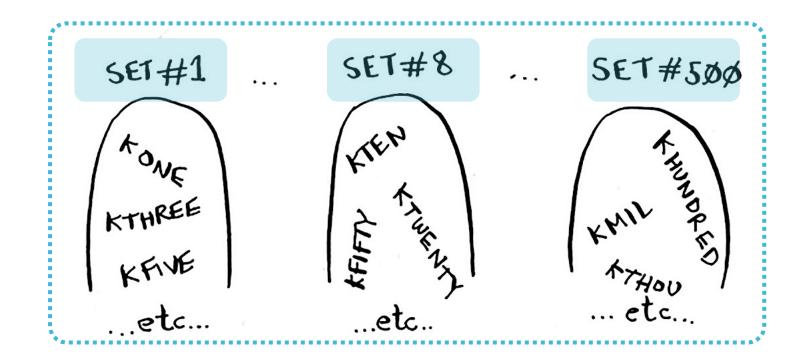




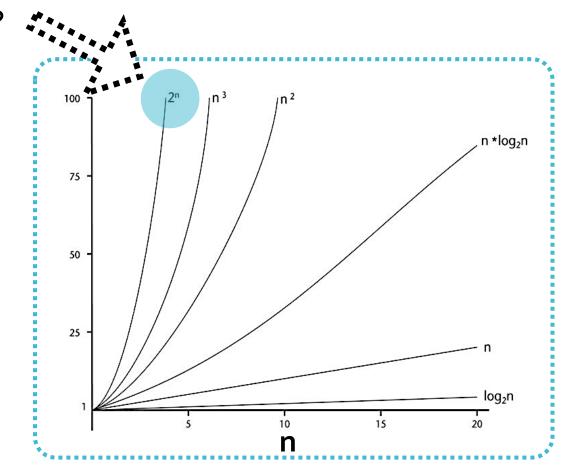
 List every possible subset of stations. This is called the power set.



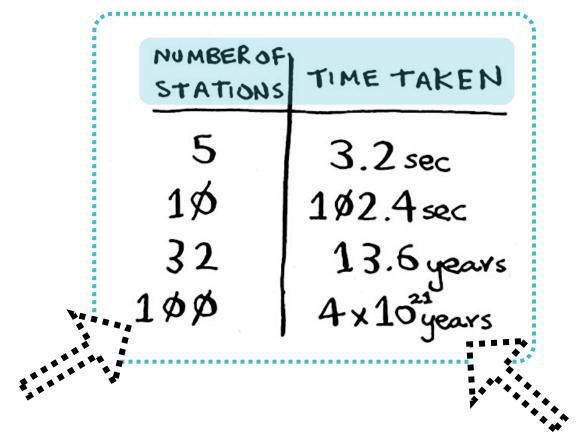
2) Then, pick the set with the **smallest number** of stations that **covers all states**.



- Can you say how many subsets we may have?
- ❖There are 2ⁿ possible subsets (=2ⁿ stations)!
- It will take **O(2ⁿ) time**. This really big.
- ❖But how big?



- **⋄**O(2ⁿ) time
- If you can calculate 10 subsets per second.



- Greedy algorithms can be the rescue, again:
 - 1) pick the station that covers **the most** states, that haven't been covered yet.
 - 2) repeat until all the states are covered.
- This algorithm takes O(n²) time, where n is the number of radio stations.
- It is an approximation algorithm.
- If computing the exact solution takes too much time, such approximation algorithm can help.

- Approximation algorithms are judged by
 - how fast they are?
 - how **close** they are to the optimal solution?
- Greedy algorithms are good choices because:
 - they are **simple** to come up with.
 - they usually run fast.

Lets implement our Greedy Algorithm to find the states, step-by-step.

- First, we build a **python dictionary** of stations:
 - each **key** is a **station** name
 - each values is the set of states that station covers
- **Example:** "k1" station covers:

that is Idaho (id), Nevada (nv), and Utah (ut).

```
states_needed = {"mt", "wa", "or", "id", "nv", "ut", "ca", "az"}

stations = dict()
stations["k1"] = {"id", "nv", "ut"}
stations["k2"] = {"wa", "id", "mt"}
stations["k3"] = {"or", "nv", "ca"}
stations["k4"] = {"nv", "ut"}
stations["k5"] = {"ca", "az"}
```

final_stations()

is a Python Set that keeps the final set of stations.

```
def find_states(states_needed, stations):
    final_stations = set()
    while states_needed:
        best_station = None
        states_covered = set()
```

while loop

- will go through all states and pick the station that covers the most uncovered states.
- we call this best_station

```
def find_states(states needed, stations):
     final stations = set()
     while states needed:
       • best station = None
       $\ states_covered = set()
best station
   is initially set to None
```

and haven't yet been covered.

is the set of all states that this station covers

states_covered

for loop

will go through all station to check which one is the **best** station.

```
for station, states_for_station in stations.items():
    covered = states_needed & states_for_station

if len(covered) > len(states_covered):
    best_station = station
    states_covered = covered
```

covered

- it is the intersection of two sets:
- states_needed set
- \$\states_for_station set

```
for station, states_for_station in stations.items()
covered = states_needed & states_for_station

if len(covered) > len(states_covered):
    best_station = station
    states_covered = covered
```

if clause checks whether covered station has more
states than the current best_station

```
for station, states_for_station in stations.items()
    covered = states_needed & states_for_station

if len(covered) > len(states_covered):
    best_station = station
    states_covered = covered
```

if so, then:
 covered station is the new best station.

```
for station, states_for_station in stations.items()
    covered = states_needed & states_for_station

if len(covered) > len(states_covered):
    best_station = station
    states_covered = covered
```

states_needed

is updated by removing the states that aren't needed anymore.

```
states_needed -= states_covered
final_stations.add(best_station)
print('Final stations are:', final_stations)
```

when **for** loop is over:

- we add best_station to the final list of stations, called final stations
- and print the final_stations

```
print('Final stations are:', final_stations)
```

Full implementation of the Greedy Algorithm that can find the states to cover.

Testing our implementation.

```
stations = dict()
stations["k1"] = {"id", "nv", "ut"}
 stations["k2"] = {"wa", "id", "mt"}
 stations["k3"] = {"or", "nv", "ca"}
  stations["k4"] = {"nv", "ut"}
 tations["k5"] = {"ca", "az"}
  find states(states needed, stations);
[Output:]
Final stations are: {'k1', 'k3', 'k5',
```

Greedy Algorithm vs Exact Algorithm

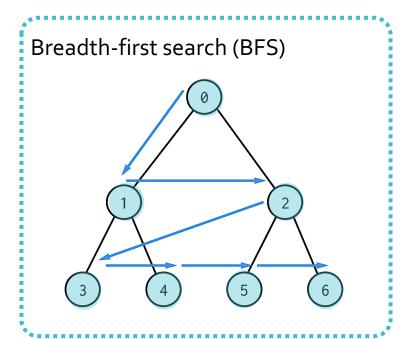
	O(n!)	$O(n^2)$
NUMBER OF	EXACT	GREEDY
STATIONS	ALGORITHM	ALGORITHM
5	3.2 sec	2.3 sec
1 Ø	1Ø2.4 sec	1Ø sec
32	13.6 yrs	102.4 sec
1ØØ	4 x 1021 yrs	16.67 min

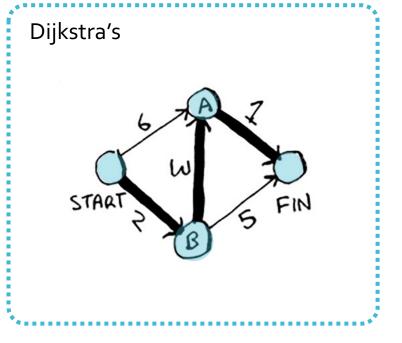


Which of the following algorithms we discussed before, are Greedy Algorithms.

- a) Breadth-first search (BFS)
- b) Dijkstra's algorithm

Quiz

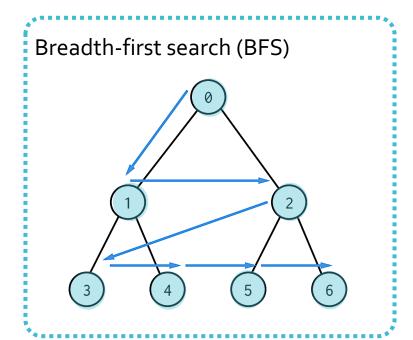


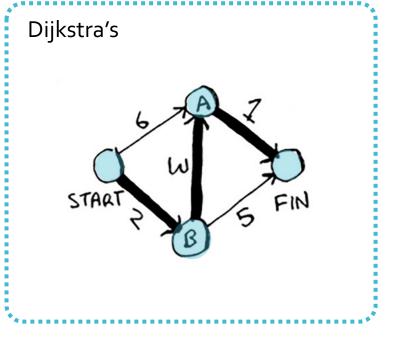


Which of the following algorithms we discussed before, are Greedy Algorithms.

Answer: Both of them.

Answer





Next Lesson

Dynamic Programming