

Advanced Programming

INFO135

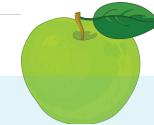
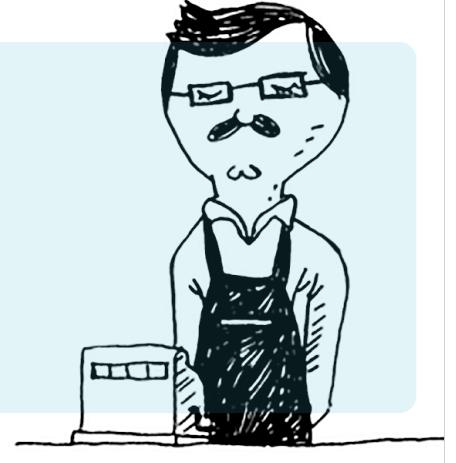
Lecture 6: Hashing

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Hashing

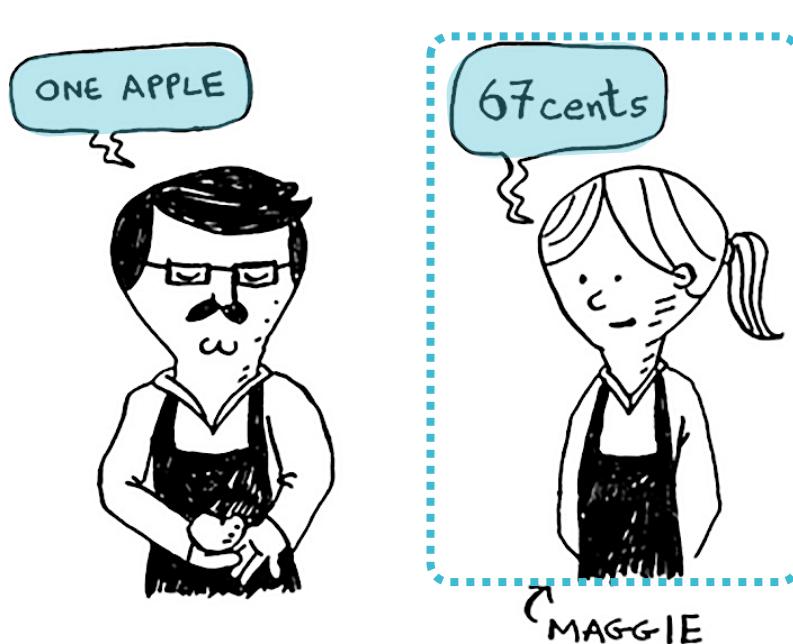
- ❖ Suppose you work at a **grocery** store.
- ❖ When a customer buys a product (e.g. apple), you **look up price** in pricebook.



- a) What if pricebook is **un-alphabetized**:
 - it would take a long time to look through every single line for apple: **O(n)** time.
- b) What if the book is **alphabetized**:
 - it would be quick since you could run **Binary Search** to find the price: **O(log n)** time.

Hashing

- ❖ Binary Search is fast, but **looking up in a book** is a challenge, even in a sorted book.
- ❖ What if your assistant “Maggie” can **memorize all names and prices!**
- ❖ Then, no need to look up anything, and you can **find the prices instantly.**



Hashing

- ❖ This is wonderful:
- ❖ **Maggie** can give you the price in **O(1)** time.
- ❖ No matter which item or **how big** the pricebook is.
- ❖ Maggie will be even **faster than Binary Search**.

# OF ITEMS IN THE BOOK	SIMPLE SEARCH	BINARY SEARCH	MAGGIE
100	$O(n)$	$O(\log n)$	$O(1)$
1000	10 sec	1 sec	INSTANT
10000	1.6 min	1 sec	INSTANT
100000	16.6 min	2 sec	INSTANT

Hashing

- ❖ The pricebook is actually an **Array**.
- ❖ Each item in the Array contains two data:

- Name
- Price

(EGGS, 2.49)	(MILK, 1.49)	(PEAR, 0.79)
--------------	--------------	--------------

- ❖ If Array is sorted by name, you can run **Binary Search** to find the price.
- ❖ But Maggie can find the price in $O(1)$ time.
- ❖ This is basically how **Hashing** works.

Hashing

❖ Some definitions related to Hashing:

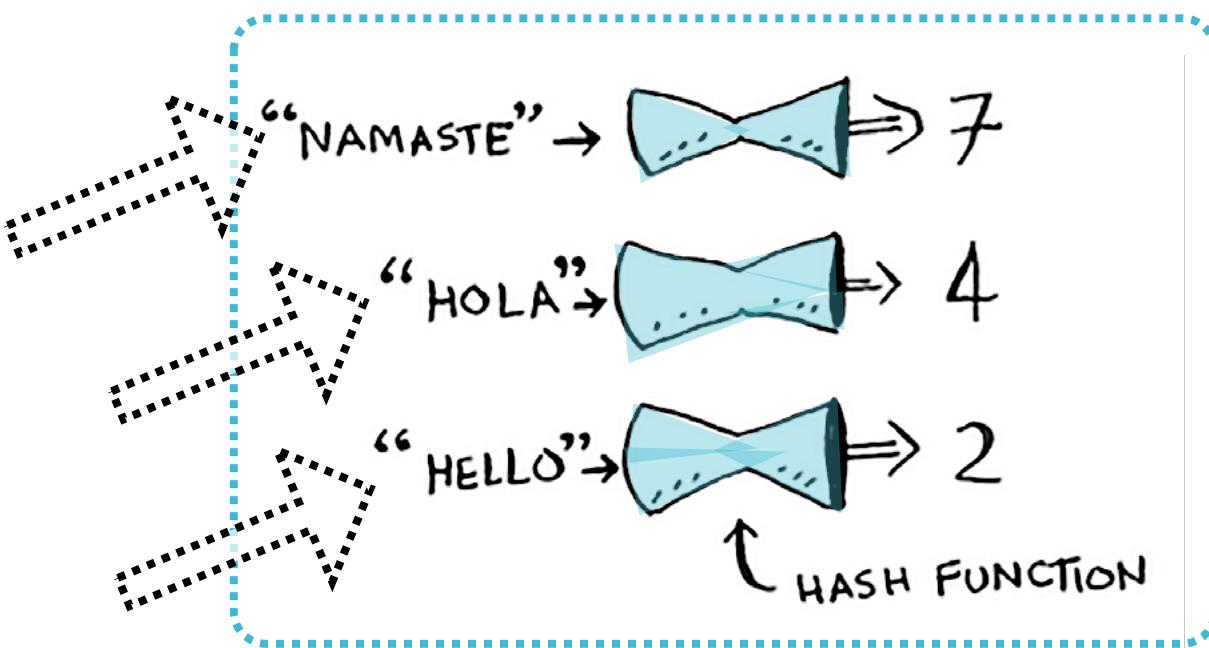
❖ **Hash Table:** a collection of items that are stored in a form which is **easy to find** them later.

❖ **Each position** of the hash table are called **a Slot** and can hold an item and is named by an integer value.

❖ **Hash Function:** is a function that maps an item to the slot where that item belongs in the **Hash Table**.

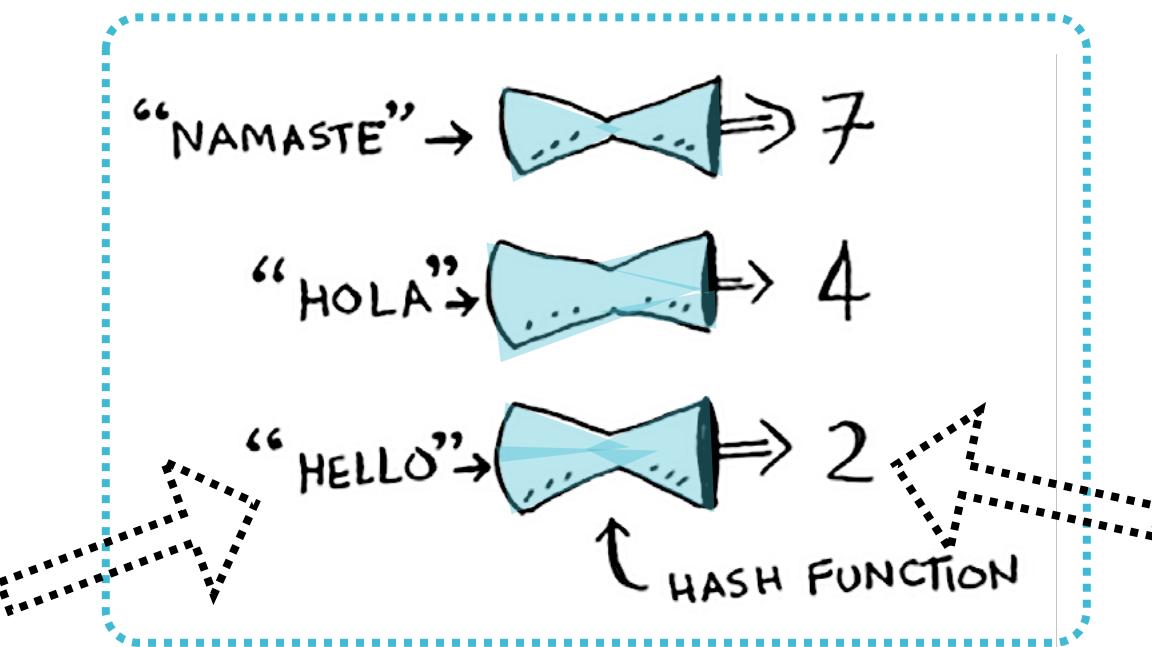
Hash Function

- ❖ Hash function takes an item in the collection and returns an integer in the range of slots.
- ❖ If there are m slots, Hash function returns an integer between 0 and $m-1$.



Hash Function

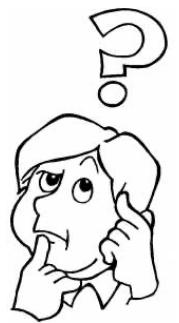
- ❖ Hash function should be **consistent**.
- ❖ For example, if you put “Hello”, then it returns “2”.
- ❖ All the time you put “Hello” it must return “2”.



Quiz

❖ Which of these functions are **consistent**?

- a) $f(x) = 1$
- b) $f(x) = \text{rand}()$
- c) $f(x) = \text{next_empty_slot_in_memory}()$
- d) $f(x) = \text{len}(x)$





❖ Which of these functions are **consistent**?

a) $f(x) = 1$ **is consistent:**

it returns always “1”

b) $f(x) = \text{rand}()$ **is not consistent:**

it returns a random number every time

c) $f(x) = \text{next_empty_slot_in_memory}()$ **not consistent:**

it returns the index of an empty slot in memory

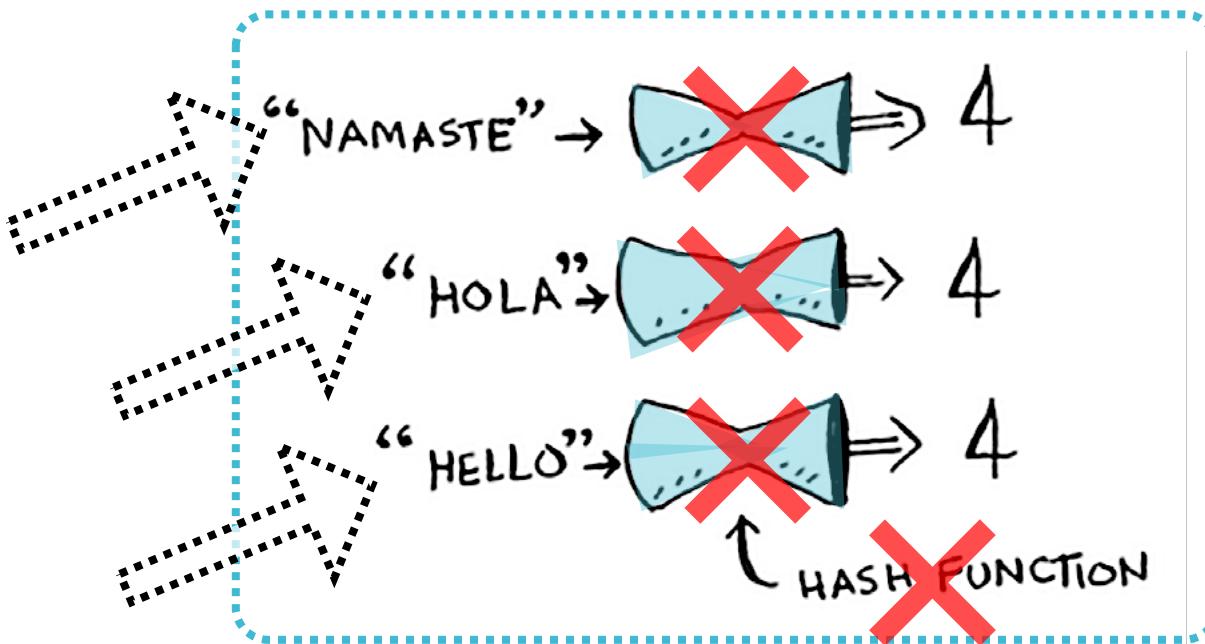
d) $f(x) = \text{len}(x)$ **is consistent:**

it returns the length (of the list or string)

Answer

Hash Function

- ❖ But in addition to consistency, Hash function must map different items to different slots.
- ❖ Example: a hash function is **not good** if:
 - ❖ it returns “4” when input is “Hello”
 - ❖ it returns “4” when input is “Hola”
 - ❖ it returns “4” when input is “Namaste”.



Hash Function

- ❖ How can we choose a good **Hash Function**?
- ❖ And how does a Hash Function **works**?

- ❖ There are **different methods** for that:
 - Division
 - Truncation
 - Folding
 - Hashing Strings

Hash Function

- ❖ **Division:** divide the key by the size of the Hash table and take the remainder:

$$h(\text{key}) = \text{key} \% m$$

- ❖ **Example:** If the key is **4873152**, and the size of the hash table is **100**, then the hash value is **52**

$$4873152 \% 100 = 52$$

Hash Function

- ❖ A more sophisticated function is called **Multiply-Add-and-Divide (MAD)**:

$$h(\text{key}) = [(a * \text{key} + b) \% p] \% m$$

- ❖ In this formula:
 - ❖ m is the size of the table
 - ❖ p is a **prime number** typically larger than m
 - ❖ a & b are random integers within interval $[0, p]$ & $a > 0$

Hash Function

- ❖ MAD function has the advantage of **eliminating the repeated patterns** in the hash codes.
- ❖ It can get closer to a **good hash function**. It is actually used for **compression**!
- ❖ The probability of **collisions** is nearly **1/N**

Quiz

- ❖ Write a has function that uses **Multiply-Add-and-Divide (MAD)** to generate hash values of the keys, with the following setting:
 - $m = 12000$
 - $p = 109345121$

- ❖ Then **call the function 3 times** and print the hash values of 2022.

- ❖ Did you **expect** the output?

Answer

❖ Example implementation of the function

```
from random import randrange
```

```
def hash_function(key):  
    m = 12000  
    p = 109345121  
    a = 1 + randrange(p - 1)  
    b = randrange(p)  
    return (a * key + b) % p % m
```

```
print(hash_function(2022))  
print(hash_function(2022))  
print(hash_function(2022))
```

Example [Output]: 7762

10372

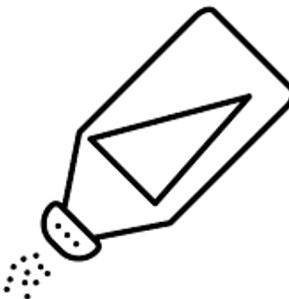
4421

Answer

Did you **expect** the output?

In real-world scenarios the hash values are “salted” with an **unpredictable** random value.

Although they **remain constant** within individual processes, they are not predictable between repeated invocations of functions.



Example [Output]: 7762

10372

4421

Hash Function

❖ **Truncation:** For large integers, some columns in the key are ignored and not used in the computation of the Hash table index

❖ **Example:** If the key is **4873152**, hash value is **812**.



Hash Function

❖ **Folding:** split a key into multiple parts & combine them into a single integer value, by adding them.

❖ **Example:** If the key is **4873152** the hash value **831**

The diagram illustrates the folding process. On the left, the key **4873152** is shown in a row of boxes. The first two digits, **48**, are highlighted with a dotted blue border. The next three digits, **731**, are highlighted with a dotted blue border. The last two digits, **52**, are highlighted with a dotted blue border. An arrow points from the sum **48 + 731 + 52 =** to the right, where the result is shown as the hash value **831**, which is also enclosed in a row of boxes with a dotted blue border.

$$48 + 731 + 52 = 831$$

Hash Function

❖ **Hashing Strings:** sum the ASCII values of the individual characters.

ASCII printable characters

32	space	64	@	96	.
33	!	65	A	97	a
34	"	66	B	98	b
35	#	67	C	99	c
36	\$	68	D	100	d
37	%	69	E	101	e
38	&	70	F	102	f
39	'	71	G	103	g
40	(72	H	104	h
41)	73	I	105	i
42	*	74	J	106	j
43	+	75	K	107	k
44	,	76	L	108	l
45	-	77	M	109	m
46	.	78	N	110	n
47	/	79	O	111	o
48	0	80	P	112	p
49	1	81	Q	113	q
50	2	82	R	114	r
51	3	83	S	115	s
52	4	84	T	116	t
53	5	85	U	117	u
54	6	86	V	118	v
55	7	87	W	119	w
56	8	88	X	120	x
57	9	89	Y	121	y
58	:	90	Z	122	z
59	;	91	[123	{
60	<	92	\	124	
61	=	93]	125	}
62	>	94	^	126	~
63	?	95	_		

Hash Function

❖ Hashing Strings: sum ASCII values of individual characters.

❖ Example: If the key is “Hello”
the hash value is 500

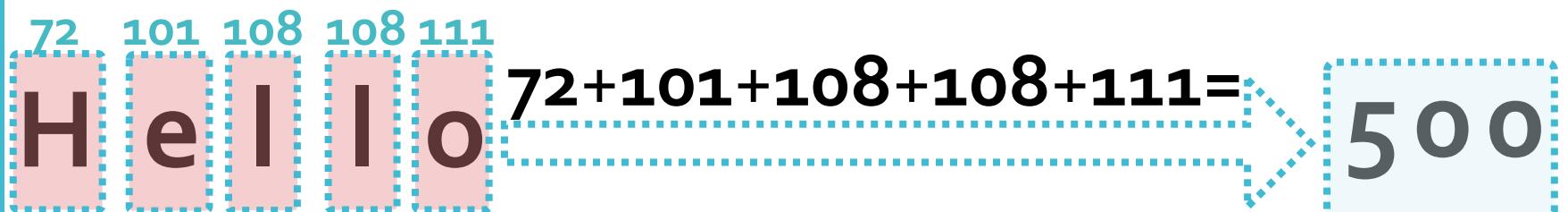
❖ Why?

ASCII printable characters

32	space	64	@	96	`
33	!	65	A	97	a
34	"	66	B	98	b
35	#	67	C	99	c
36	\$	68	D	100	d
37	%	69	E	101	e
38	&	70	F	102	f
39	'	71	G	103	g
40	(72	H	104	h
41)	73	I	105	i
42	*	74	J	106	j
43	+	75	K	107	k
44	,	76	L	108	l
45	-	77	M	109	m
46	.	78	N	110	n
47	/	79	O	111	o
48	0	80	P	112	p
49	1	81	Q	113	q
50	2	82	R	114	r
51	3	83	S	115	s
52	4	84	T	116	t
53	5	85	U	117	u
54	6	86	V	118	v
55	7	87	W	119	w
56	8	88	X	120	x
57	9	89	Y	121	y
58	:	90	Z	122	z
59	;	91	[123	{
60	<	92	\	124	
61	=	93]	125	}
62	>	94	^	126	~
63	?	95	_		

Hash Function

❖ Example: If the key is “Hello” the hash value is **500**



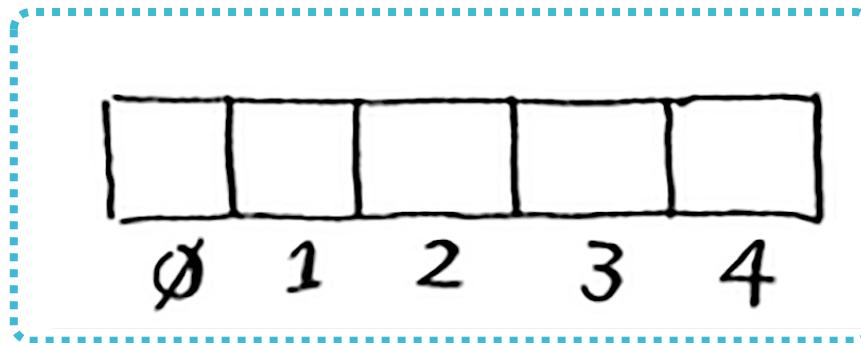
Hash Table

- ❖ If you put a hash function & an array together, you can build a new Data Structure called **Hash Table**.
- ❖ Array and List can map straight to memory, but **Hash Table is smarter.**
- ❖ Hash Table uses Hash Function to intelligently figure out **where in the memory** to store the items.

Hash Table

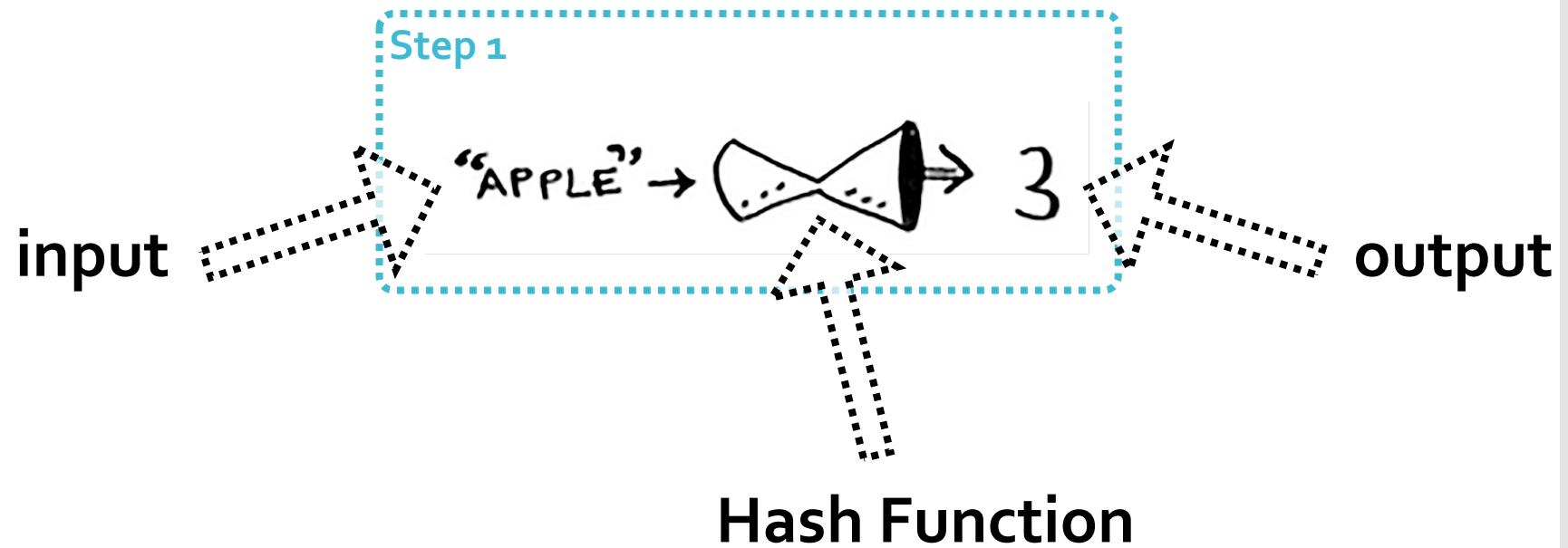
- ❖ Lets assume we chose a good Hash function.
- ❖ How can we make a **Hash Table** for Maggie's task?

- ❖ We start with an empty array:



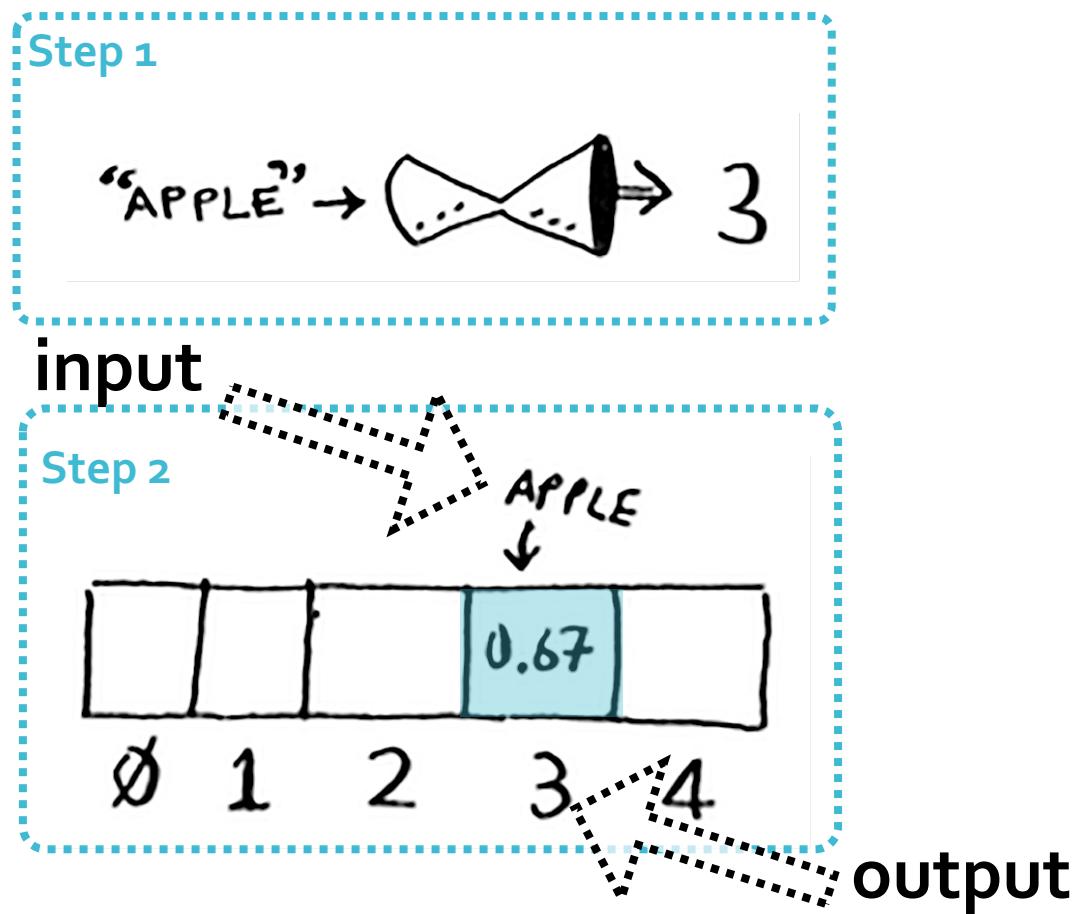
Hash Table

- ❖ We store all the prices in this array.
- ❖ Lets add **apple**, the Hash function **outputs “3”**.



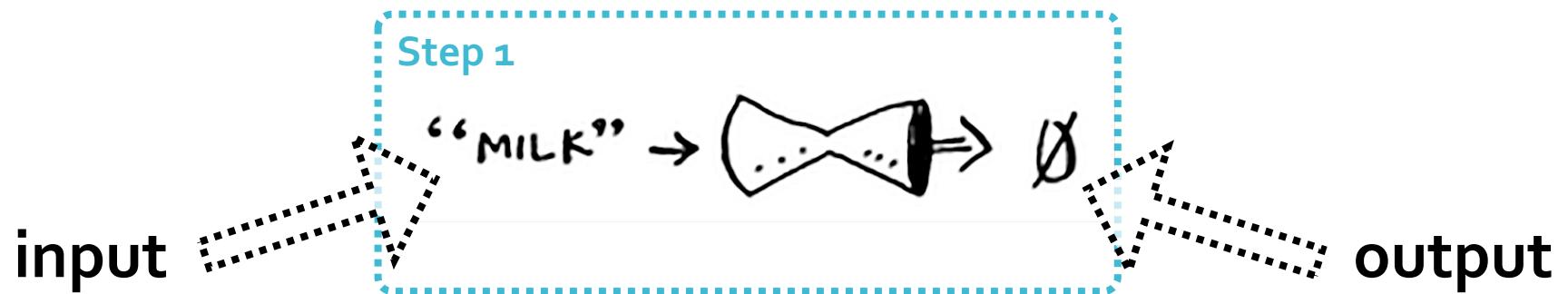
Hash Table

- ❖ We add the **price** of apple at **index 3** in the array.



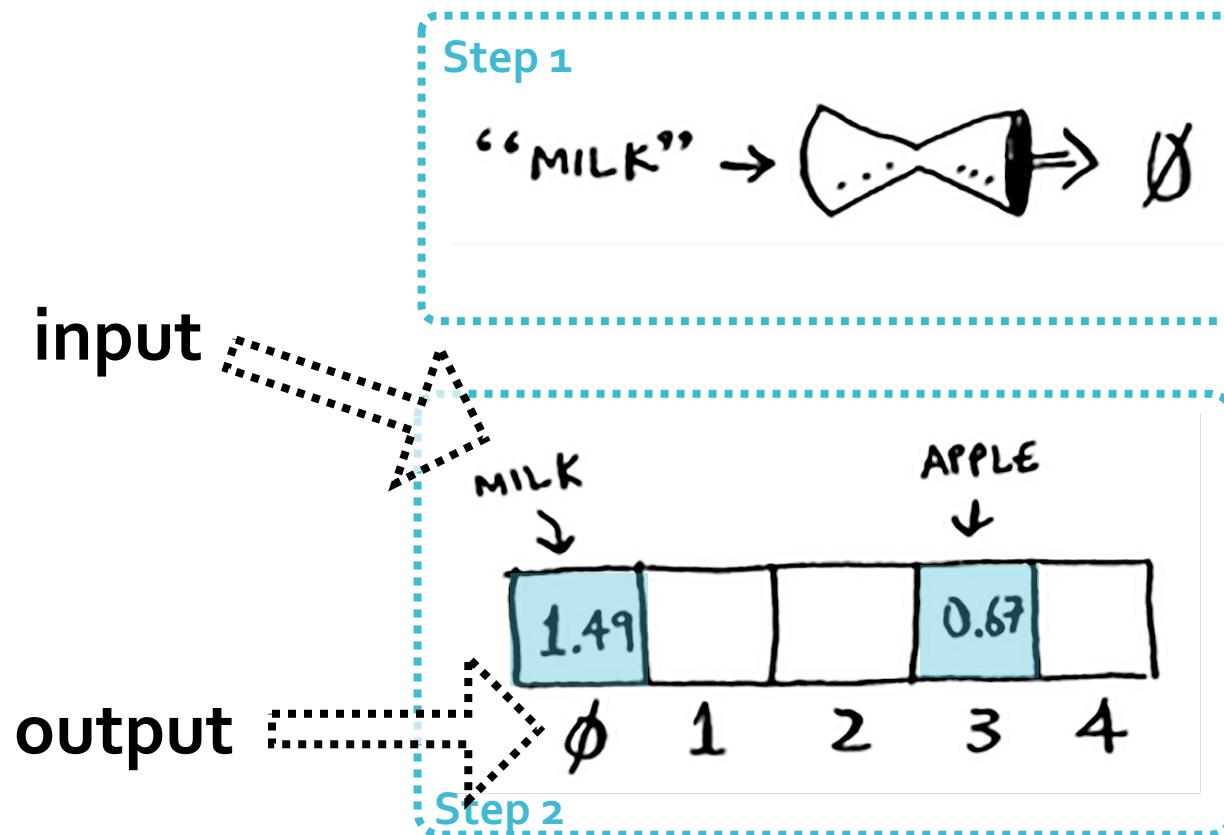
Hash Table

- ❖ Lets add **milk**, the Hash function **outputs “o”**.



Hash Table

❖ Lets add **milk**, the Hash function **outputs “o”**.



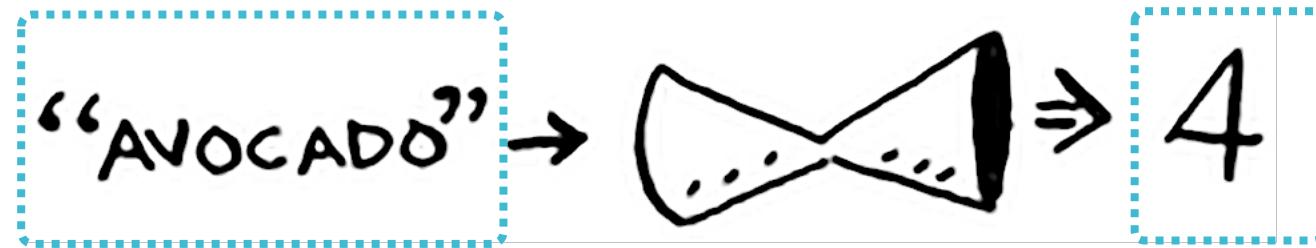
HashTable

- ❖ Keep going, till the **whole array** will be full of prices.

1.49	0.79	2.49	0.67	1.49
------	------	------	------	------

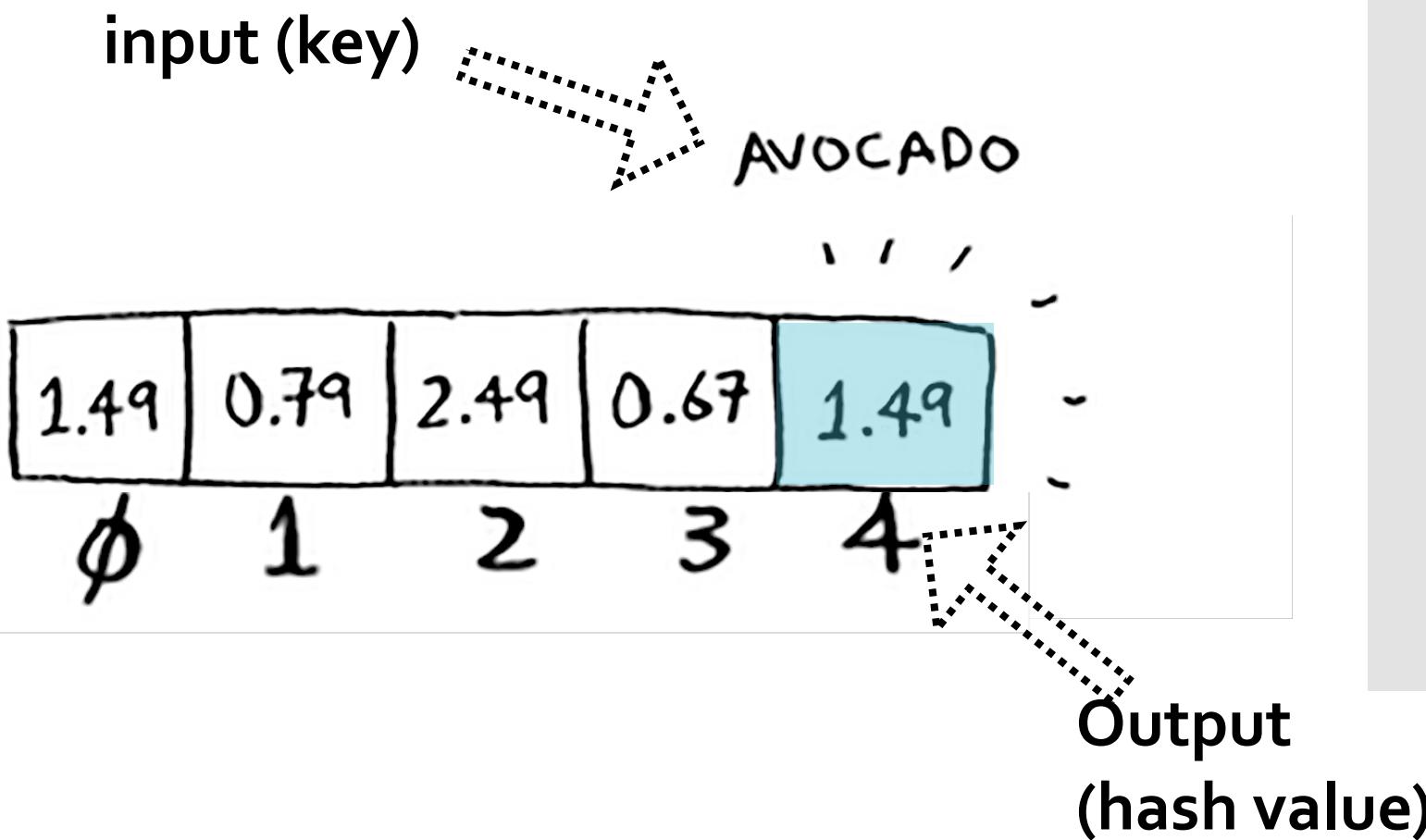
Hash Table

- ❖ Now you want to know the price of an **Avocado**:
- ❖ **No need to wait for a Search algorithm** to do all the operations to find it.
- ❖ Just put “Avocado” into the **Hash Function!**



Hash Table

- ❖ It will instantly tell that the **price** is stored at **index 4**.



Hash Table

- ❖ Hash Table is one of the **most useful data structure**.
- ❖ It is really fast!
- ❖ It has other names **Dictionary, Hash Maps, or Associative Array**.

Hash Table

- ❖ The **efficiency of hashing** depends in large part on the selection of a good hash function.
- ❖ **Perfect Hash Function:** given a collection of items, a hash function that maps each item into a unique slot.
- ❖ But this is **rarely achieved**.
- ❖ Instead, we can choose a **good Hash Function** that will **distribute the keys** across the range of Hash Table indices as evenly as possible.

Hash Table

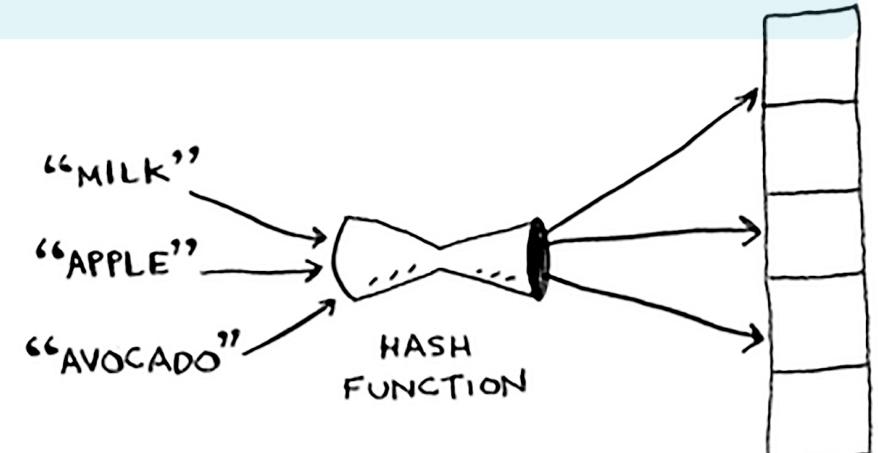
- ❖ **Guidelines for choosing a Hash Function:**
 - ❖ Computation should be simple to produce **quick results.**
 - ❖ Table size should be a **prime number** to produce a better distribution and fewer collisions.

Hash Function

- ❖ But we care about **performance**?
- ❖ Well.... we need to first understand what **Collisions** are.

Hashing

- ❖ Before, we assumed a hash function **always maps different keys to different slots in the array.**
- ❖ In reality, this is almost **impossible**.



- ❖ Lets check an **example**:

Hashing

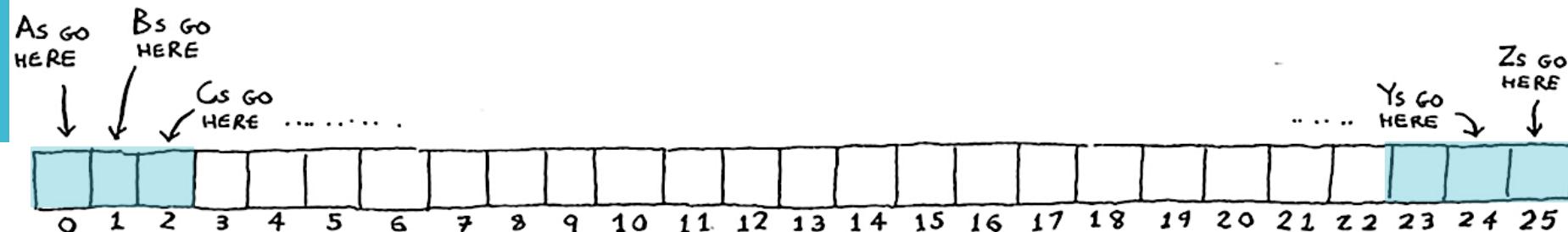
- ❖ Suppose you have an array with **26 slots**.
- ❖ Similar to what we said previously, it stores the prices of products.

❖ "A" -> 0

❖ "B" -> 1

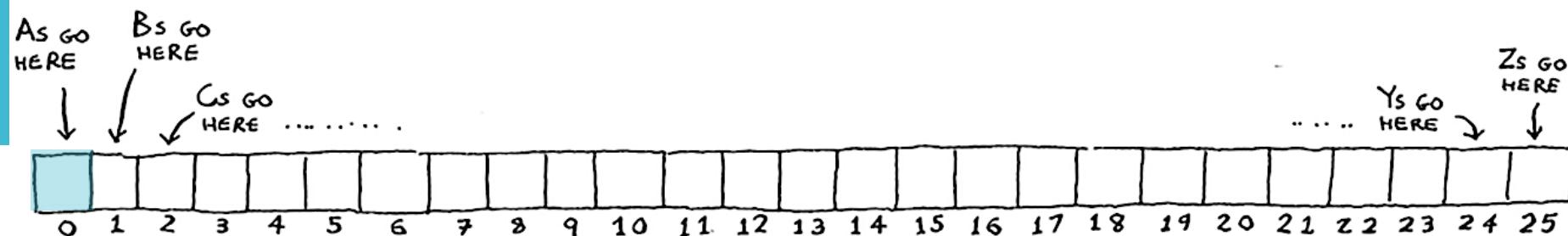
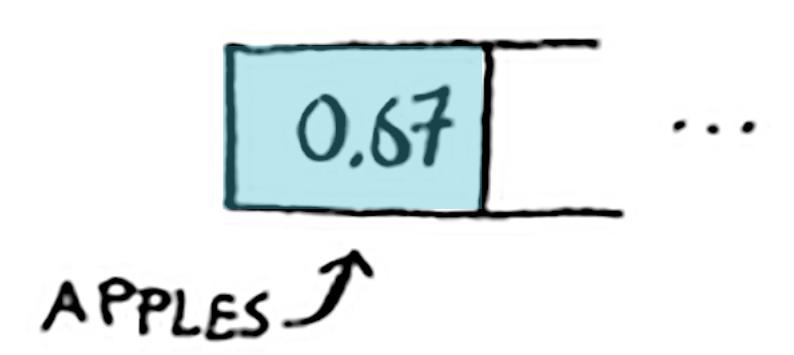
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❖ "Z" -> 25



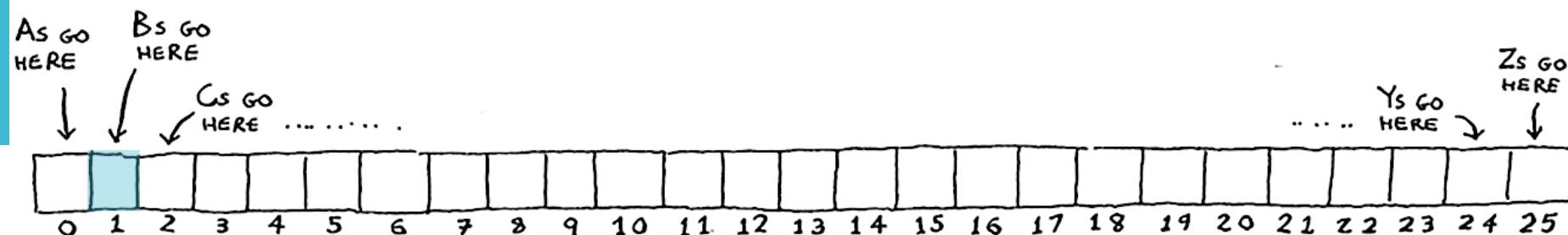
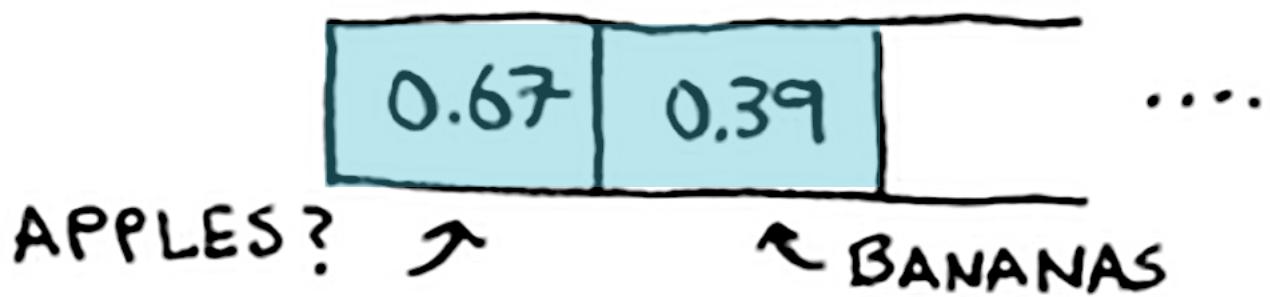
Hashing

❖ Now, you want to put the price of apples in your hash. You get assigned the first slot.



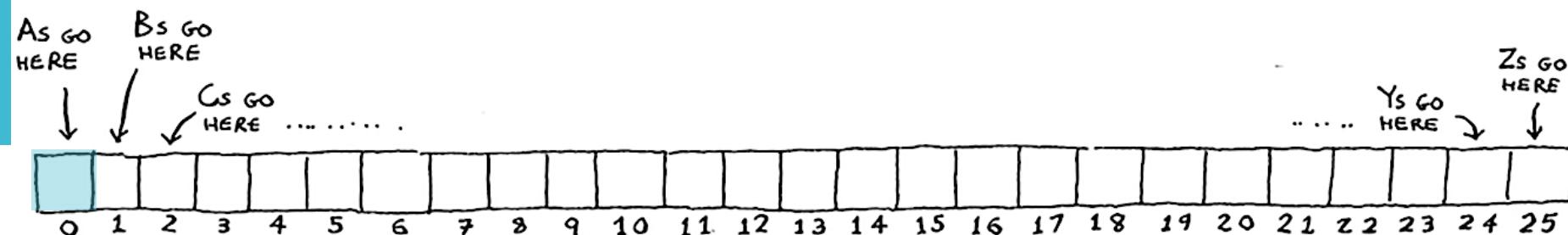
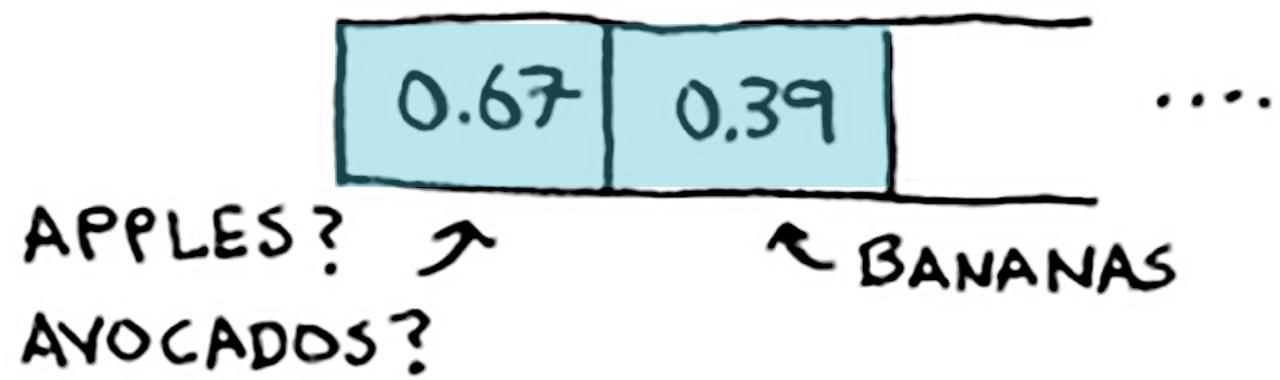
Hashing

- ❖ Now you want to put the price of banana.
- ❖ You get assigned the second slot again!



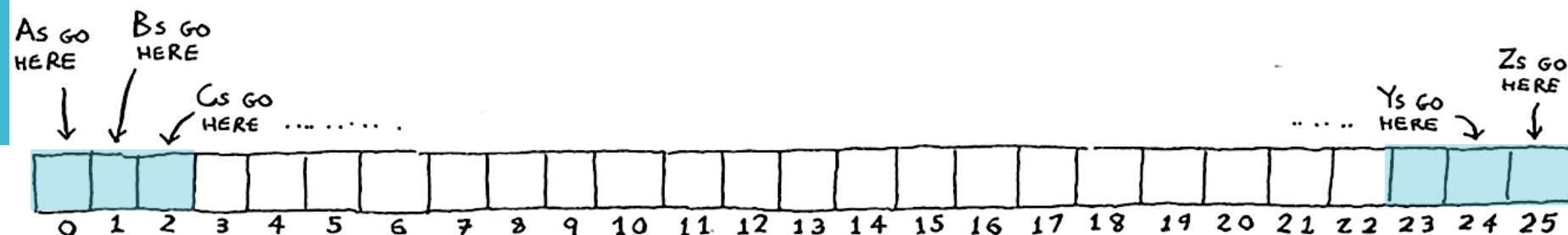
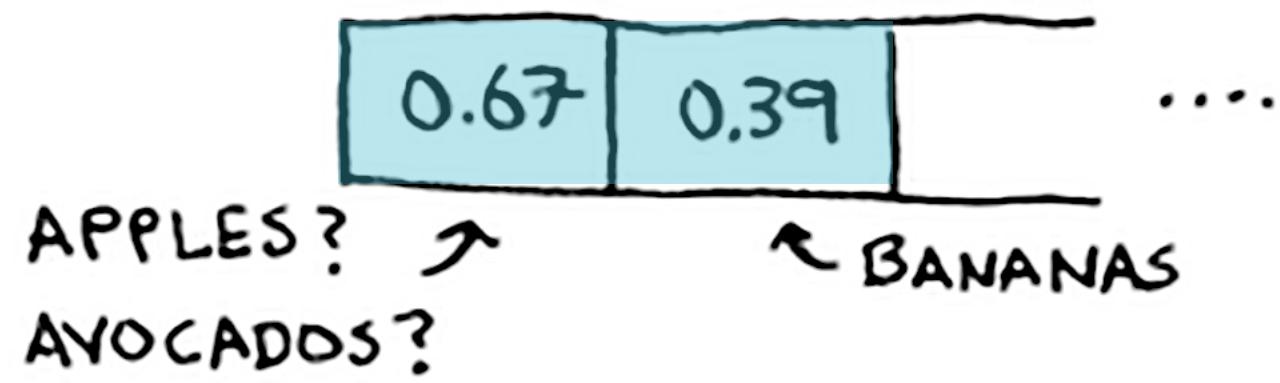
Hashing

- ❖ But if you want to put the **price of avocados**.
- ❖ You get assigned the **first slot** again!



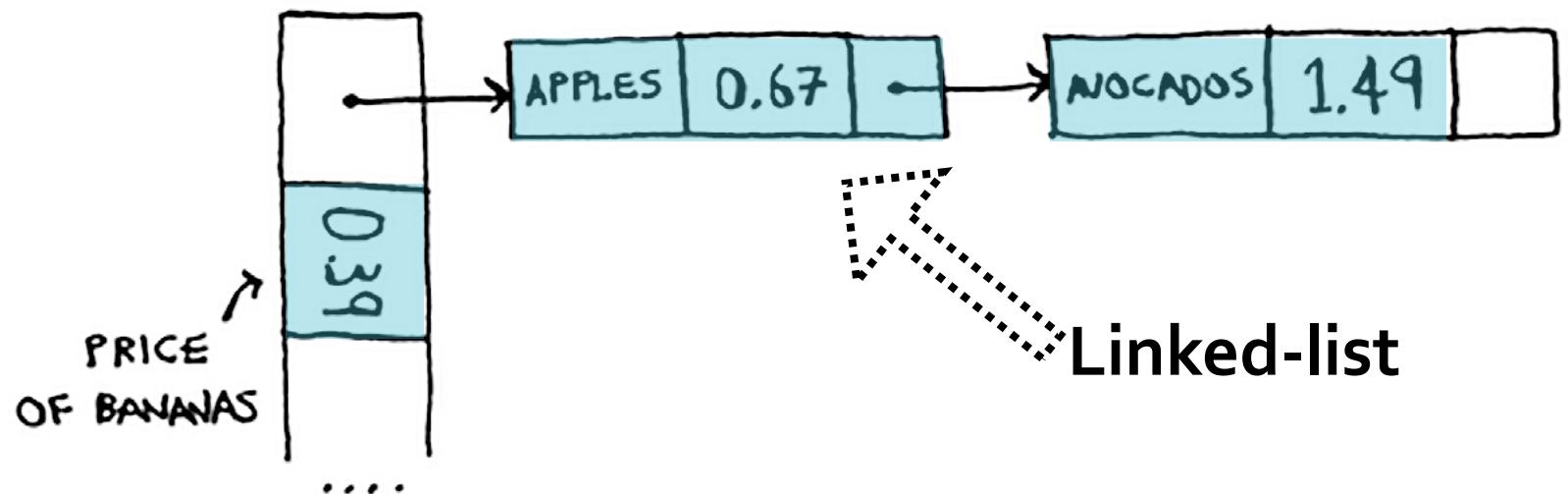
Collision

- ❖ But apples have that slot already!
- ❖ This is called a collision: two keys have been assigned the same slot.



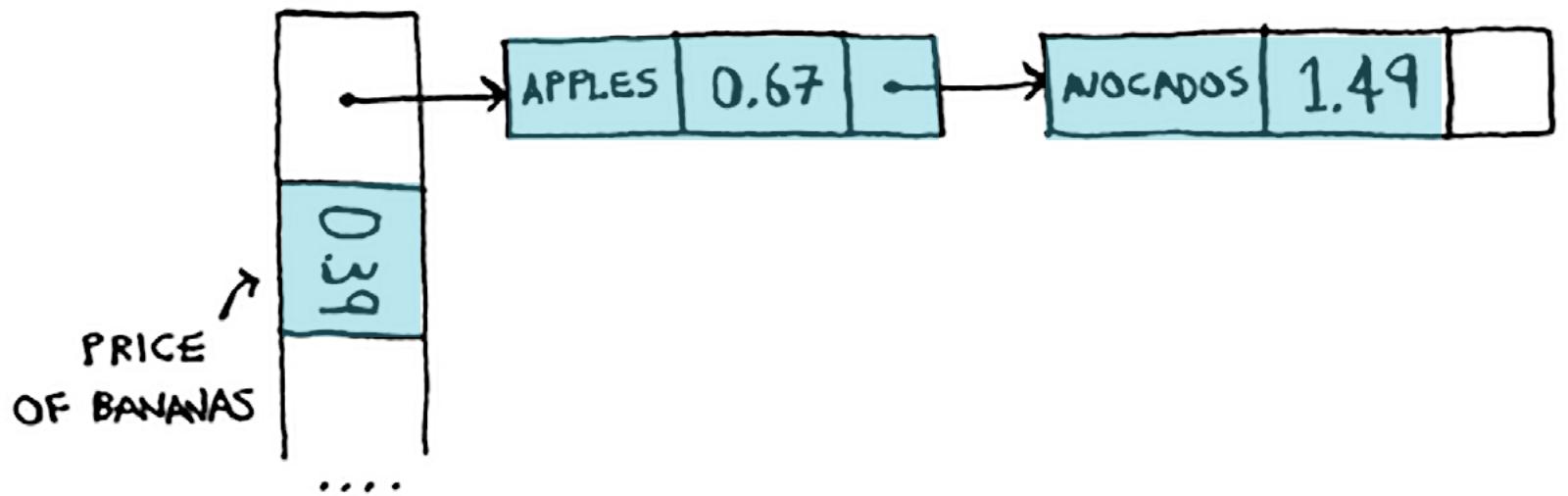
Collision

- ❖ There are different ways to deal with **collisions**.
- ❖ **The simplest one is this:** if multiple keys map to the same slot, start a **Linked-list** at that slot.



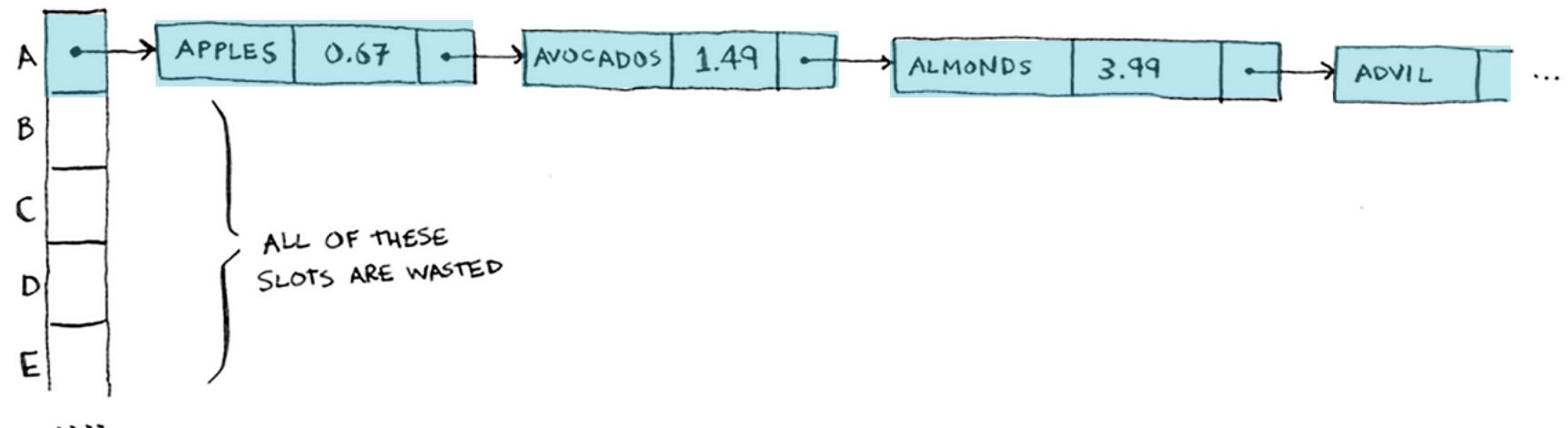
Collision

- ❖ If you need to know:
 - ❖ the price of **bananas**, it's still **quick**.
 - ❖ the price of **apples** or **avocados**, it's a **slower**.



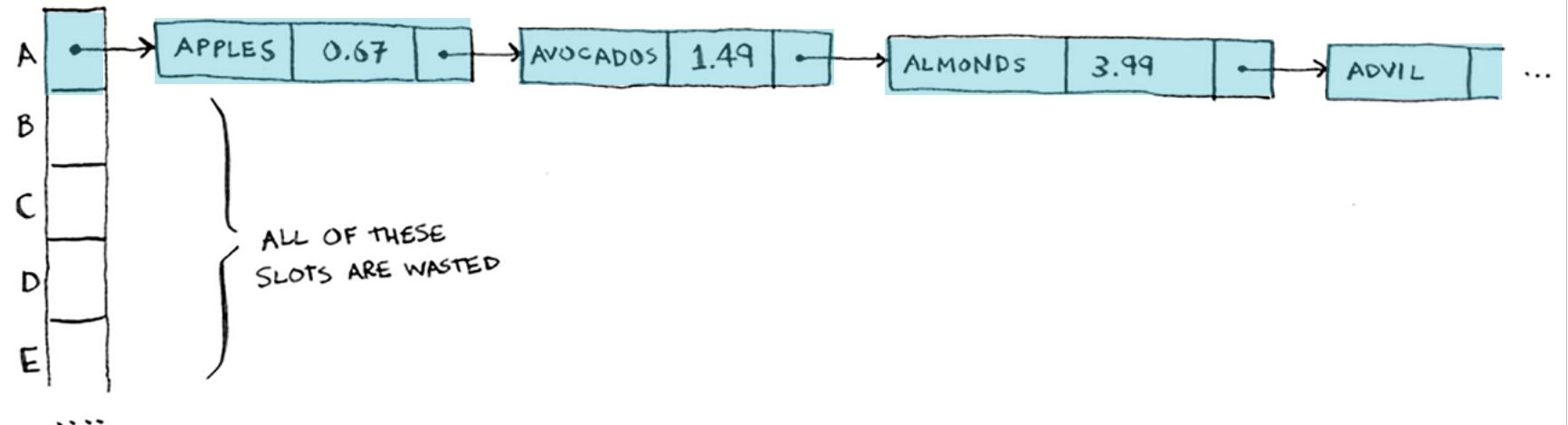
Collision

- ❖ If the Linked-list is small, no big deal.
- ❖ But suppose you work at a grocery store where you only sell **a lot of products that starts with the letter A**.



Collision

- ❖ Then Hash Table is totally empty **except for 1 slot!**
- ❖ And that slot has a **giant Linked-list!**
- ❖ Every item in this Hash table is in the Linked-list.
- ❖ It's going to **slow down** your hash table.



Collision

❖ Three lessons here:

- Such long **Linked-lists slows down** a lot a hash table.
- Hash function should **map keys evenly** all over hashes.
- **Good Hash function** should give a **very few collisions**.

Collision

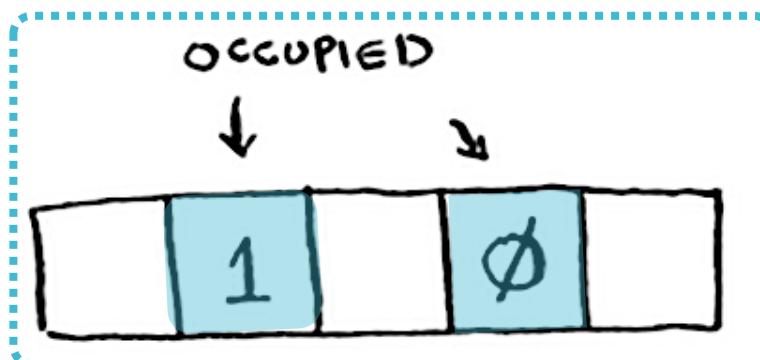
- ❖ **Open Addressing:** process of resolving collisions by looking into the Hash table and trying to find another open slot.
- ❖ **Example:** start at the original hash value position and then move in a sequential manner through the slots until finding an empty slot that is empty.
- ❖ Visiting each slot one at a time, with open addressing technique, is called **Linear Probing**.

Performance

- ❖ Load factor of a Hash Table is an important factor.
- ❖ It is calculated by the **number of items** divided by the total **number of slots** in Hash Table.

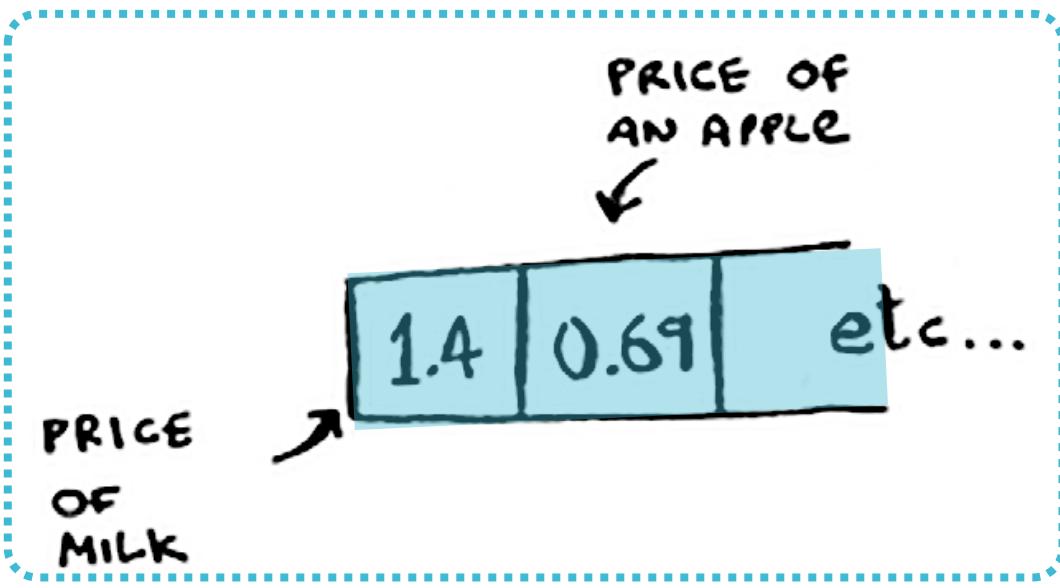
$$\frac{\text{NUMBER OF ITEMS IN HASH TABLE}}{\text{TOTAL NUMBER OF SLOTS}}$$

- ❖ Example: this Hash table has a load factor of **2/5**



Performance

- ❖ Suppose you need to store the **price of 100 items** in your hash table.
- ❖ Your hash table has **100 slots**, and hence in the best case, **each item will get its own slot**.
- ❖ This hash table has a **load factor of 1**.

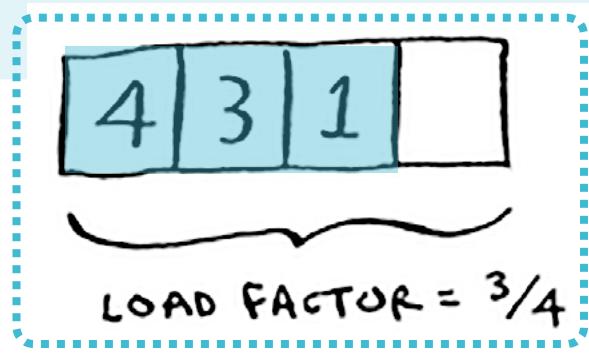


Performance

- ❖ What if your hash table has only 50 slots? Then it has a **load factor of 2**.
- ❖ Having a **load factor greater than 1** means you have more items than slots in your array.
- ❖ **Resizing:** when the **load factor grows**, more slots to have to be added to the hash table.

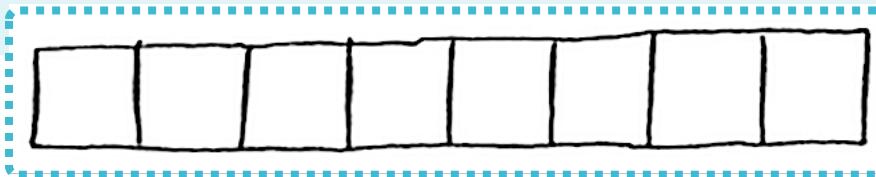
Performance

❖ Example: suppose you have this hash table that is getting pretty full.

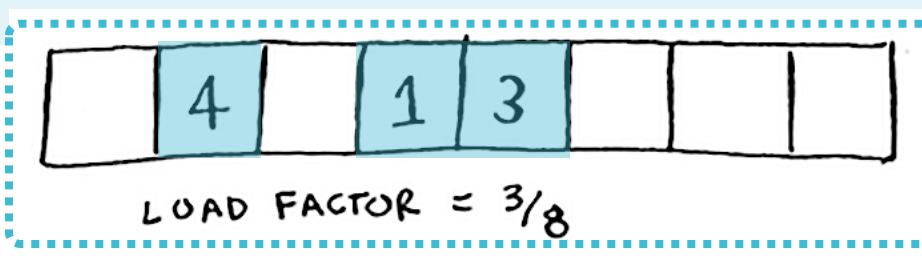


❖ You need to **resize** this hash table:

- 1) create a new array that's bigger (e.g., twice the size)

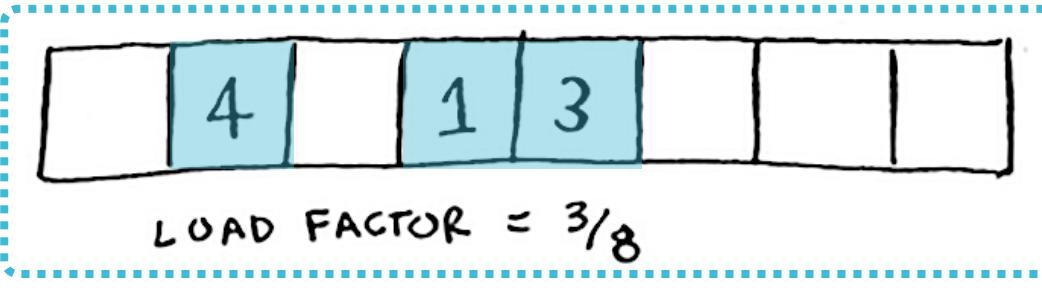


- 2) then, **re-insert** all items into this new Hash table using the hash function:



Performance

- ❖ This new table has a load factor of $3/8$: Much better!



- ❖ Lower load factor means fewer collisions the Hash table will perform better.
- ❖ Rule of thumb: resize when load factor is greater than 0.7

Performance

- ❖ Resizing is expensive and takes a lot of time.
- ❖ Resizing should not occur too often.
- ❖ But still, by averaged, Hash tables take $O(1)$ even with resizing.

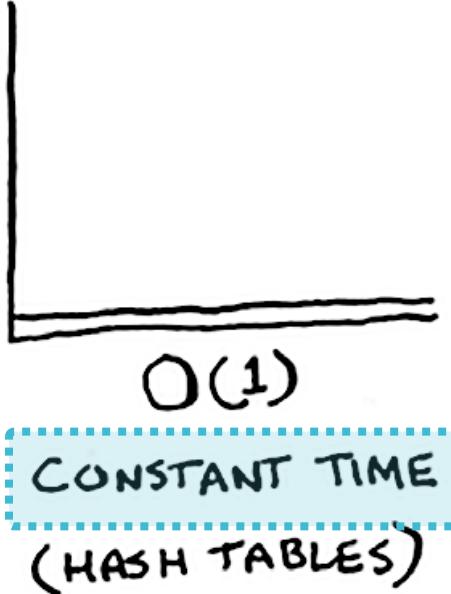
Performance

❖ So what is the conclusion?

- hash tables are still **really fast**: average case is $O(1)$ (constant time).
- $O(1)$ means the **time** taken is **the same**, no matter **how big** is the hash table.

Performance

❖ Hash Tables



Performance

❖ Hash Tables vs Binary Search



$O(1)$
CONSTANT TIME
(HASH TABLES)



$O(\log n)$
LOG TIME
(BINARY SEARCH)

Performance

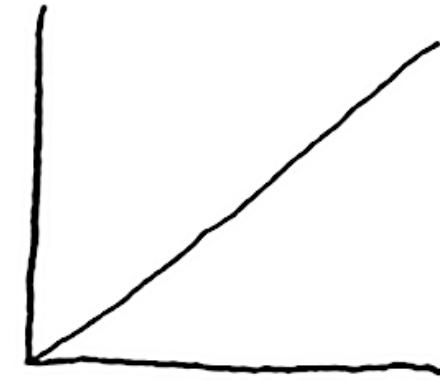
❖ Hash Tables vs Binary Search vs Simple Search



$O(1)$
CONSTANT TIME
(HASH TABLES)



$O(\log n)$
LOG TIME
(BINARY SEARCH)



$O(n)$
LINEAR TIME
(SIMPLE SEARCH)

Performance

❖ Hash Tables vs Binary Search vs Simple Search

	HASH TABLES (AVERAGE) ARRAYS	LINKED LISTS
SEARCH	$O(1)$	$O(1)$
INSERT	$O(1)$	$O(n)$
DELETE	$O(1)$	$O(n)$

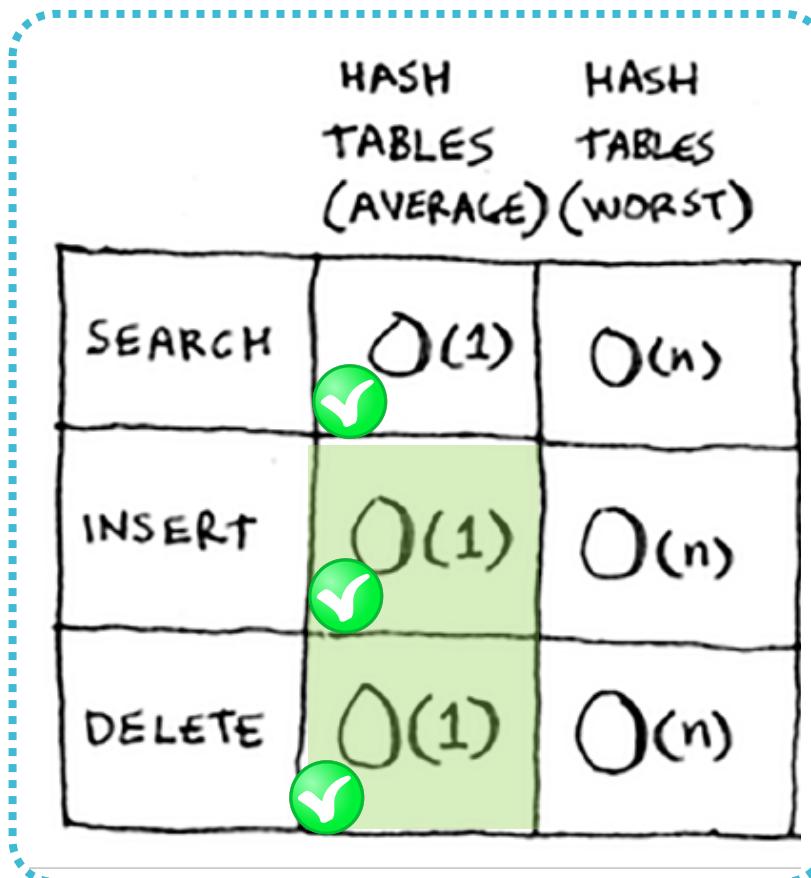
A hand-drawn table comparing the performance of Hash Tables, Average Arrays, and Linked Lists. The table has columns for SEARCH, INSERT, and DELETE operations. Hash Tables and Average Arrays are grouped under "HASH TABLES (AVERAGE) ARRAYS" and are highlighted in green. Linked Lists are under "LINKED LISTS" and are white. Checkmarks are placed in the first three cells of each row.

	HASH TABLES (AVERAGE) ARRAYS	LINKED LISTS
SEARCH	$O(1)$	$O(1)$
INSERT	$O(1)$	$O(n)$
DELETE	$O(1)$	$O(n)$

Performance

- ❖ Hash tables (in average) are $O(1)$:
 - ❖ as fast as arrays at searching (getting a value at an index).
 - ❖ as fast as linked lists at insert and delete.
- ❖ Hash tables (in worst case) are $O(n)$ which is slow at all of those.

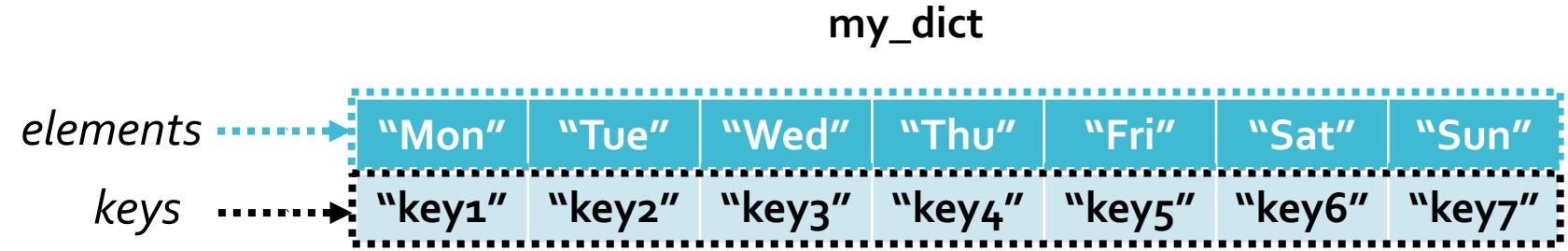
		HASH TABLES (AVERAGE)	HASH TABLES (WORST)
SEARCH	$O(1)$	$O(n)$	
INSERT	$O(1)$	$O(n)$	
DELETE	$O(1)$	$O(n)$	



Hashing

- ❖ Python and other languages have **built-in hash tables** and no need to implement the hash tables.

Remember?



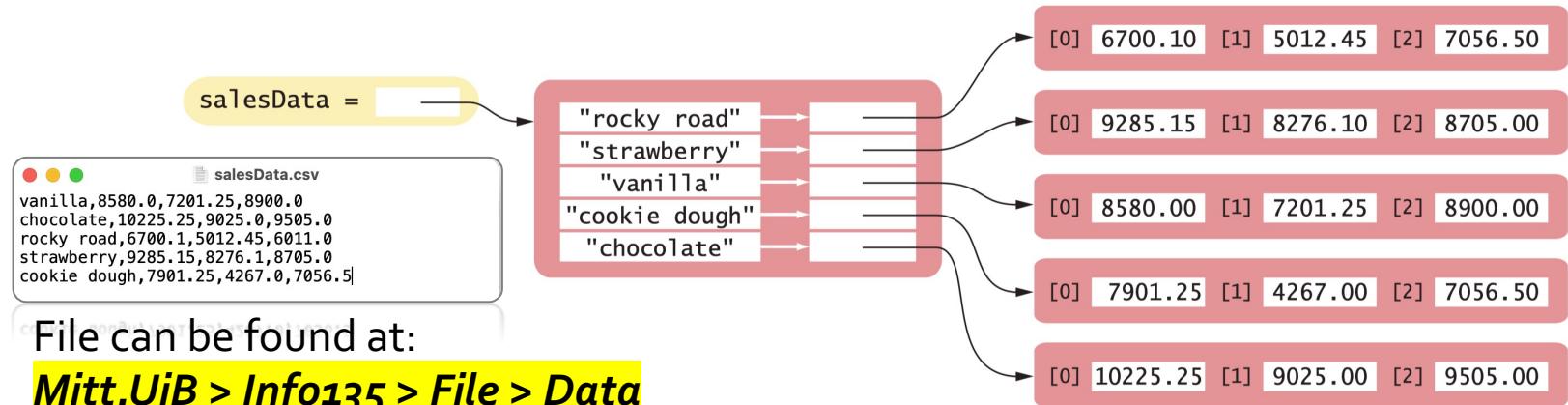
- ❖ Data collection that is **unordered**, **mutable**, and **indexed**
- ❖ Can have **duplicates**
- ❖ Represented with **curly brackets** and specifying the **key & value** pairs {key:value}

- ❖ Example:

```
my_dict =  
{"key1": "Mon", "key2": "Tue", "key3": "Wed",  
 "key4": "Thu", "key5": "Fri", "key6": "Sat",  
 "key7": "Sun"}
```

Quiz

- ❖ Write a Python program **build_dict_from_data()** that receives the name of file as input parameter and goes through and reads it.
- ❖ Use the program to read a file called **salesData.csv**, that contains a collection of sales for an ice-cream shop.
- ❖ The program should create and print a Python dictionary where the **keys** are ice-cream **flavors** and the **values** are **lists** of numbers representing the sales of the shop within different months.



Quiz

Reminder for opening files:

❖ Syntax:

```
with open("filename" , "r") as file:  
    data = file.read()
```

❖ Example:

```
with open("Hello.txt") as file:  
    data = file.read()
```

Answer

- ❖ Reading the file and building a dictionary for **every row of the file** (key:values)

Part 1

```
def build_dict_from_date(filename):
    salesData = {}

    with open(filename, "r") as infile:
        for line in infile:
            fields = line.split(",")
            flavor = fields[0]
            salesData[flavor] = make_list(fields)

    return salesData
```

Answer

❖ Generating list of numbers (monthly sales of the shop).

```
def make_list(fields):  
    storeSales = []  
    for i in range(1, len(fields)):  
        sales = float(fields[i])  
        storeSales.append(sales)  
  
    return storeSales
```

Part 2

```
salesData = build_dict_from_date("salesData.csv")  
  
for key, value in salesData.items():  
    print(key, ' : ', value)
```

Answer

❖ Generating list of numbers (monthly sales of the shop).

```
def make_list(fields):  
    storeSales = []  
    for i in range(1, len(fields)):  
        sales = float(fields[i])  
        storeSales.append(sales)  
  
    return storeSales
```

```
salesData = build_dict_from_date("salesData.csv")  
  
for key, value in salesData.items():  
    print(key, ' : ', value)
```

Part 3

[Output]: vanilla : [8580.0, 7201.25, 8900.0]
chocolate : [10225.25, 9025.0, 9505.0]
rocky road : [6700.1, 5012.45, 6011.0]
strawberry : [9285.15, 8276.1, 8705.0]
cookie dough : [7901.25, 4267.0, 7056.5]

Hashing

- ❖ Python has also a built in Hash function:
hash(*object*)

- ❖ It return hash value of the object. You can also **override** it!

- ❖ Numbers that equal value (even with different types) have the same hash value.

- ❖ **Example of usage:** to quickly compare dictionary keys during a dictionary lookup.

Hashing

❖ Hashing

```
class Car:  
    def __init__(self, model, brand):  
        self.model = model  
        self.brand = brand  
  
    def __eq__(self, other):  
        return self.model == other.model \  
               and self.brand == other.brand  
  
    def __hash__(self):  
        return self.model \  
               * ord(self.brand[0]) \  
               * ord(self.brand[1]) \  
               * ord(self.brand[2]) \  
               * ord(self.brand[3])
```

Hashing

❖ Hashing

```
my_car1 = Car(2020, 'Benz')
my_car2 = Car(2019, 'Audi')
print("My Hash value for car 1:", hash(my_car1))
print("My Hash value for car 2:", hash(my_car2))
```

[Output:]

My Hash value for car 1: 180704594400

My Hash value for car 2: 161222197500

Is this a good hash function? why?

Hashing

❖ Hashing

```
my_car1 = Car(2020, 'Benz')
my_car2 = Car(2019, 'Audi')
print("My Hash value for car 1:", hash(my_car1))
print("My Hash value for car 2:", hash(my_car2))
```

[Output:]

My Hash value for car 1: 180704594400

My Hash value for car 2: 161222197500

Is this a good hash function? why?

No. There will be collisions when the cars have **the same brand and model**. This has not resolved.

Quiz

- 1) Write a class **Employee** where you override Hash function. Define **your own** Hash function which works based on the **id_number** of an employee.

```
class Employee:  
    def __init__(self, id_number):  
        [REDACTED]  
  
    def __eq__(self, other):  
        [REDACTED]  
  
    def __hash__(self):  
        [REDACTED]
```

Quiz

- 2) In Python only **immutable objects can be hashed**.
But why?

❖ Hash function based on **id_number** of employees.

```
class Employee:  
    def __init__(self, id_number):  
        self.id_number = id_number  
  
    def __eq__(self, other):  
        return self.id_number == other.id_number  
  
    def __hash__(self):  
        return hash(self.id_number)
```

```
employee1 = Employee(30122000)  
employee2 = Employee(29121999)  
print("Hash value for employee 1:", hash(employee1))  
print("Hash value for employee 2:", hash(employee2))
```

[Output:]

```
Hash value for employee 1: 30122000  
Hash value for employee 2: 29121999
```

Answer

Answer

- ❖ If a mutable objects are also allowed to be hashed, then we must **update Hash Table** every time the the objects are changed.
- ❖ This may mean that the object must be moved to a totally different location and this can be an **expensive operation**.

Hashing

❖ Popular Hash functions:

- **Message Digest algorithm 5 (MD5)**
- **First Published:** April 1992
- **By Computer Scientist Ronald Rivest**



- ## ❖ MD5 hash is created by taking a **string of any length** and encoding it into a **128-bit fingerprint**.



79054025
255fb1a2
6e4bc422
aef54eb4

=

Not recommended anymore!

- ## ❖ Demo: www.md5hashgenerator.com

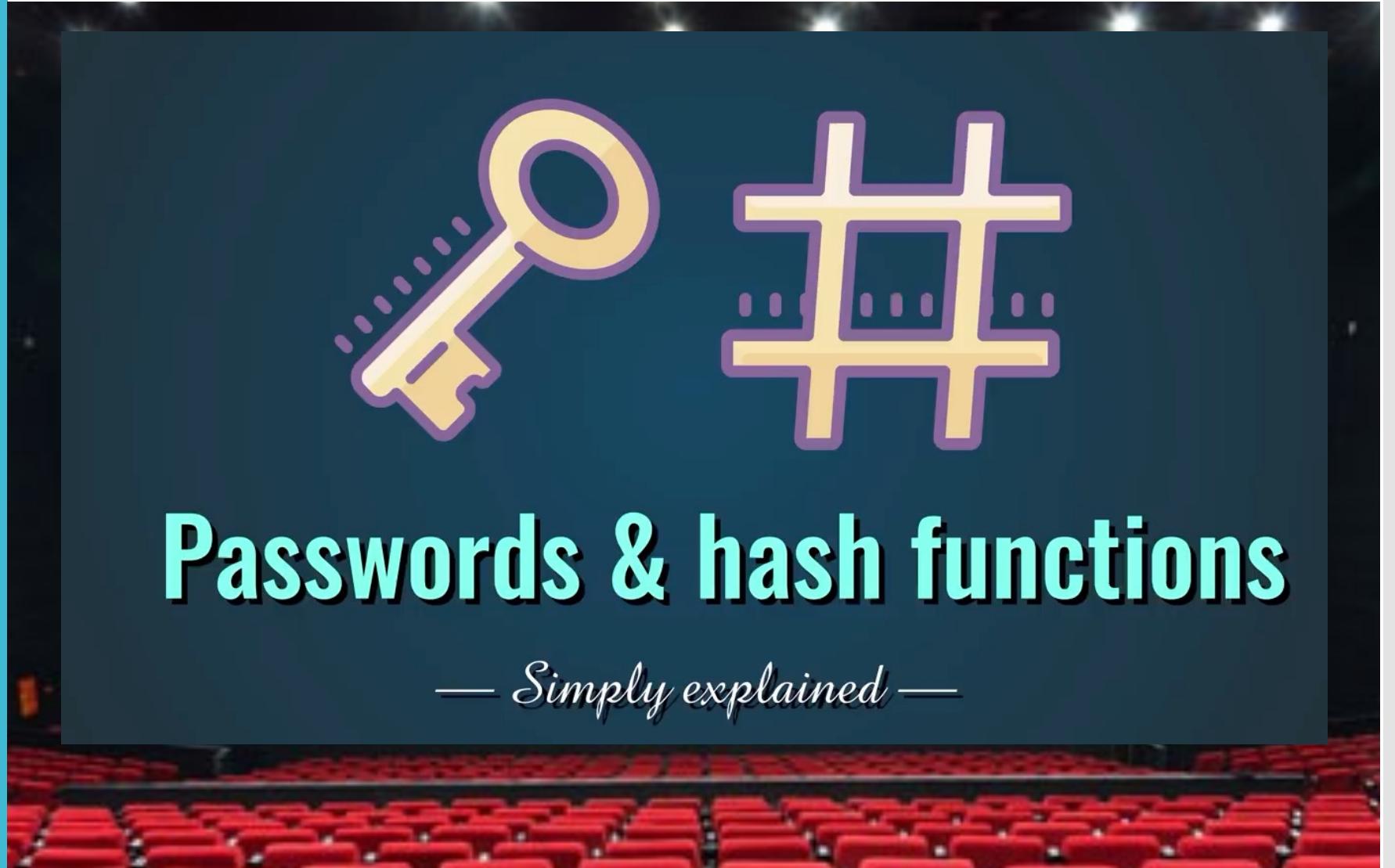
Hashing

- ❖ A popular Hash function:
 - **Secure Hash Algorithms (SHA)**
 - **First Published:** April 1993 (SHA-0)
 - **By:** National Institute of Standards and Tech
- ❖ The message generated by SHA ranges from **160 to 512 bits.**
- ❖ **Demo:** xorbin.com/tools/sha256-hash-calculator

Hashing

❖ One of the popular applications of these algorithms is **password hashing**.

Hash
Application



ref: youtu.be/cczlpieu42M

Hashing

❖ Lets try hashing with this algorithms!

❖ This is a function that uses **SHA hashing**

```
import hashlib as hl
```

```
def hash_it_by_sha1(user_str):  
    hash_str = hl.sha1(user_str.encode()).hexdigest()  
    return hash_str
```

```
print('1) Hash by sha:', hash_it_by_sha1('Hello Bergen'))  
print('2) Hash by sha:', hash_it_by_sha1('Hello Bergen!'))  
print('3) Hash by sha:', hash_it_by_sha1('Hello Bergen!!'))
```

Hashing

[Output:]

- 1) Hash by sha: a527128b69070db7c29dbe7694c57db1898c1a2e
- 2) Hash by sha: 054a328b3889660a1e000a5c6938f6c97a02ba47
- 3) Hash by sha: a875245dfceb3728af7f2fc7fefe490039a8527

Hashing

- ❖ SHA object has the following methods:
 - ❖ **digest()** returns the digest of the strings passed to the update method so far (in byte format where it may contain non-ASCII characters).
 - ❖ **hexdigest()** is similar to **digest()** except this returns hexadecimal digits.
 - ❖ **update(arg)** updates the object with the string *arg*.
 - ❖ **copy()** return a copy (“clone”) of the object. This can be used to efficiently compute the digests of strings that share a common initial substring.

❖ Lets write a Password class

```
import hashlib as hl

class Password:

    def __init__(self):
        user_pass = input('Please enter your password: ')
        self.hash_pass = self.hash_it(user_pass)
```

Hashing

❖ Lets write a Password class

```
import hashlib as hl
```

```
def hash_it(self, user_pass):  
    self.hash_pass = hl.sha1(user_pass.encode()).hexdigest()  
    return self.hash_pass
```

Hashing

❖ Lets write a Password class

```
import hashlib as hl
```

Hashing

```
def print_it(self):  
    print('Your hashed password is:', self.hash_pass)  
    print('Size of hash password is:', len(self.hash_pass))
```

Hashing

```
import hashlib as hl

class Password:

    def __init__(self):
        user_pass = input('Please enter your password: ')
        self.hash_pass = self.hash_it(user_pass)

    def hash_it(self, user_pass):
        self.hash_pass = hl.sha1(user_pass.encode()).hexdigest()
        return self.hash_pass

    def print_it(self):
        print('Your hashed password is:', self.hash_pass)
        print('Size of hash password is:', len(self.hash_pass))

my_password = Password()
my_password.print_it()
```

[Output:]

```
Please enter your password: bergen2020
Your hashed password is: e17b5c9a9735e58c0ba6b8974c7721428c7f23fb
Length of hashed password is: 40
```

❖ In Password class, write a new method called **log_in()** that checks and prints if a **password**, entered by user, is **correct or not!**

```
class Password:  
  
    def __init__(self):...  
  
    def hash_it(self, user_pass):...  
  
    def print_it(self):...  
  
    def log_in(self):
```

Quiz



Quiz

```
class Password:  
    |  
    |     def __init__(self):...  
    |  
    |     def hash_it(self, user_pass):...  
    |  
    |     def print_it(self):...  
    |  
    |     def log_in(self):
```

```
my_password = Password()  
my_password.log_in()
```

[Output:]

Please enter your password: *bergen2020*

To login, please re-enter your password: *bergen2020*

Password is correct.



Answer

```
class Password:

    def __init__(self):...

    def hash_it(self, user_pass):...

    def print_it(self):...

    def log_in(self):
        entered_pass = input('To login,re-enter password: ')
        if self.hash_pass == self.hash_it(entered_pass):
            print('Password is correct.')
        else:
            print('Sorry, password not correct.')

my_password = Password()
my_password.log_in()
```

[Output:]

Please enter your password: bergen2020

To login, please re-enter your password: bergen2020

Password is correct.



Next Lesson

❖ **Graphs**