

Computer Networks (01)

An Introduction

lijt@fudan.edu.cn

Some slides from Jim Kurose

Introduction

Our goal:

- ❑ get “feel” and terminology
- ❑ more depth, detail *later* in course
- ❑ approach:
 - ❖ use Internet as example

Overview:

- ❑ what's the Internet
- ❑ what's a protocol?
- ❑ network edge
- ❑ network core
- ❑ access net, physical media
- ❑ Internet/ISP structure
- ❑ performance: loss, delay
- ❑ protocol layers, service models
- ❑ network modeling

Course Non-goals

❑ Learn how to configure a Router

- ◆ That requires a class all by itself
- ◆ Huawei/Cisco teaches those classes
- ◆ Our perspective will be broader

❑ Become "Internet Experts"

- ◆ The Internet will be our frequent *motivating example*
- ◆ *Our perspective will be broader*
- ◆ *Focus on Internet, but understand important common issues and challenges*

Why *not* an "Internet class"?

□ Is there anything other than the Internet?

Give two examples ?

Course Goals

- ❑ Think “the network way”
 - ◆ Distributed coordination is hard
- ❑ Learn how computer networks work
 - ◆ Problems, approaches, protocols, software
- ❑ Learn how to write network applications
- ❑ Hands-on understanding of network internals
 - ◆ Course Projects

What Is a Network?

- ❑ Collection of nodes and links that connect them
- ❑ This is vague. Why? Consider different networks:
 - ❖ Internet
 - ❖ Social networks
 - ❖ Telephone
 - ❖ Others
 - - transport nets, sensor nets, cellular nets, ...

What is a Computer Network

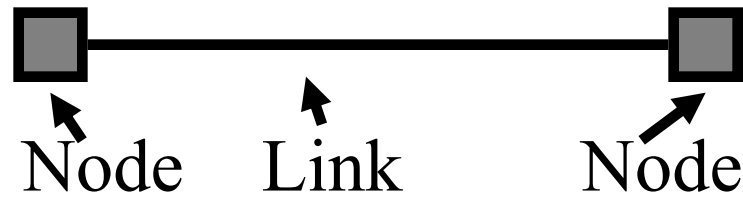
- ❑ Interconnected;
- ❑ Autonomous;
- ❑ Collection.

Abstraction:

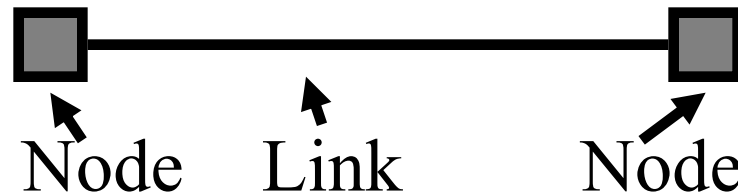
End system + Link + Switching node

Host + Link + Router (for Internet)

How to Draw a Computer Network

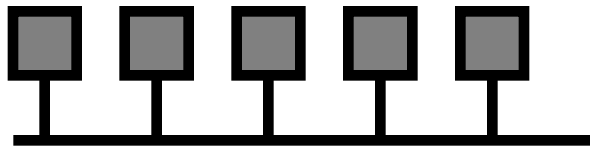


Building block: The Links

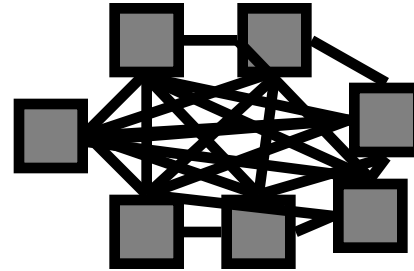


- ❑ Electrical questions
 - ❖ Voltage, frequency, ...
 - ❖ Wired or wireless?
- ❑ Link-layer issues: How to send data?
 - ❖ When to talk - can everyone talk at once?
 - ❖ What to say - low-level format?
 - ❖ Stay tuned for Chapter 5
- ❑ Okay... what about more nodes?

□ ... But what if we want more hosts?



One wire

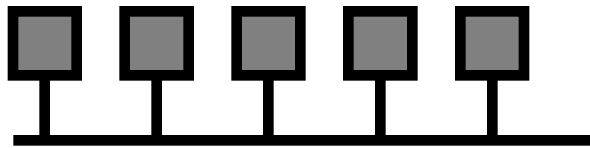


Wires for everybody!

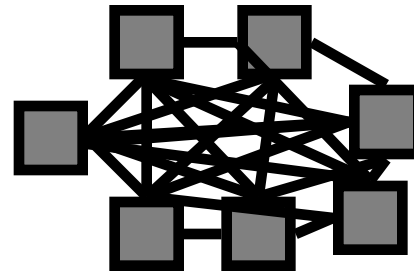
□ Scalability?!

Scalability: how to solve ?

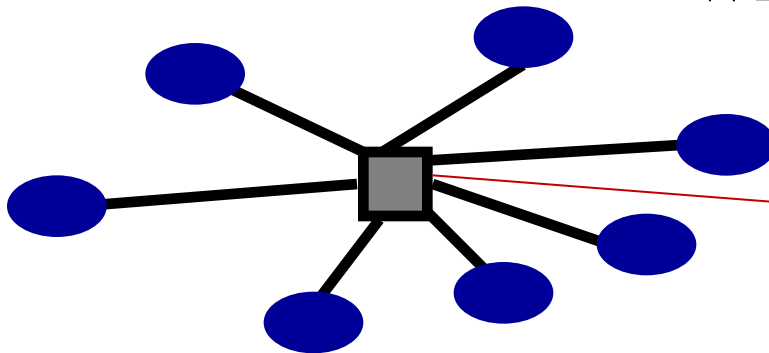
□ ... But what if we want more hosts?



One wire



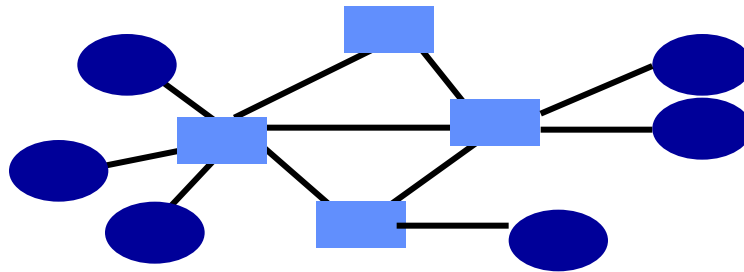
Wires for everybody!



Switching node

Multiplexing!

- Need to share network resources



- How? Switched network
 - ❖ Party "A" gets resources sometimes
 - ❖ Party "B" gets them sometimes
- Interior nodes act as "Routers" or "Switches"
- What mechanisms can share resources?

Review:

□ Layering

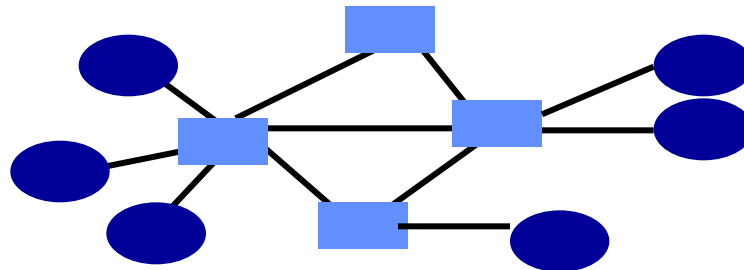
- ❖ Abstraction
- ❖ OSI vs. TCP/IP

□ Control plane and Data Plane

- ❖ Design considerations

□ Basic Model for computer networks

End system + Link + Switching node



Networks Juggle Many Goals

❑ Efficiency - resource use; cost

❑ The "ilities":

- ❖ Evolvability
- ❖ Managability
- ❖ Security (securability, if you must)
- ❖ Ease of:
 - Creation
 - Deployment
 - Management
 - *Creating useful applications*
- ❖ Scalability

A Protocol for Packet Network Interconnection. Vinton Cerf and Bob Kahn
The Design Philosophy of the DARPA Internet Protocols. David D. Clark

Roadmap

1.1 What *is* the Internet?

1.2 Network edge

1.3 Network core

1.4 Network access and physical media

1.5 Internet structure and ISPs

1.6 Delay & loss in packet-switched networks

1.7 Protocol layers, service models

1.8 History

What's the Internet: “nuts and bolts” view



□ millions of connected computing devices:
hosts

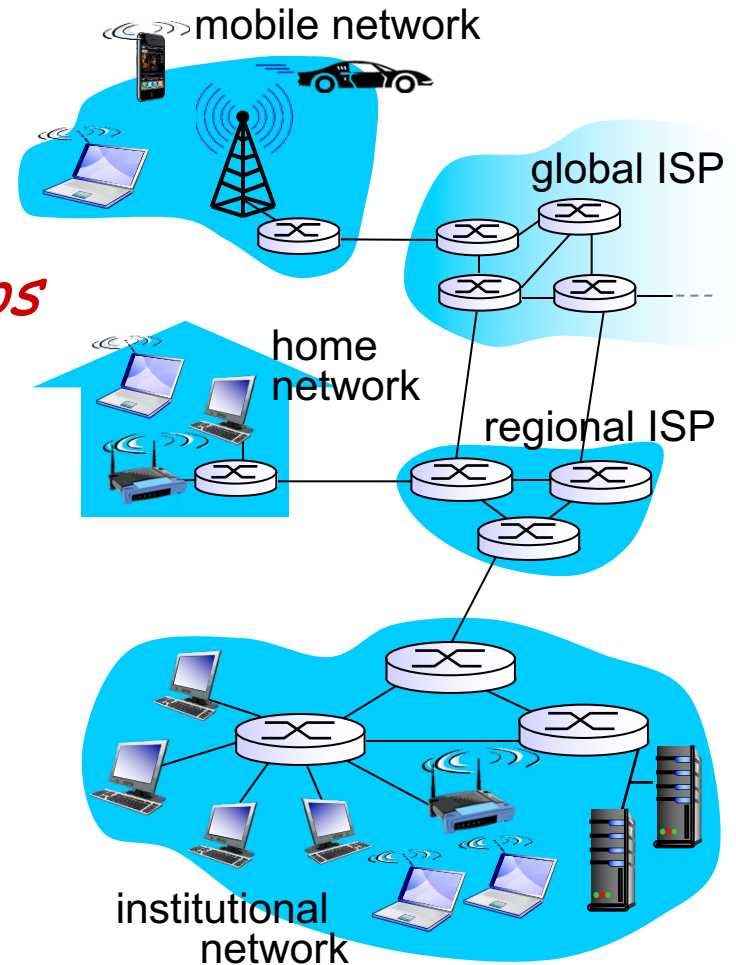
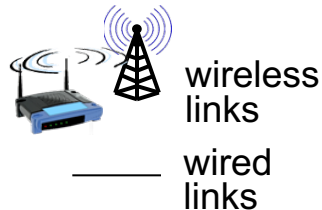
❖ running *network apps*

❖ *communication links*

■ fiber, copper, radio, satellite

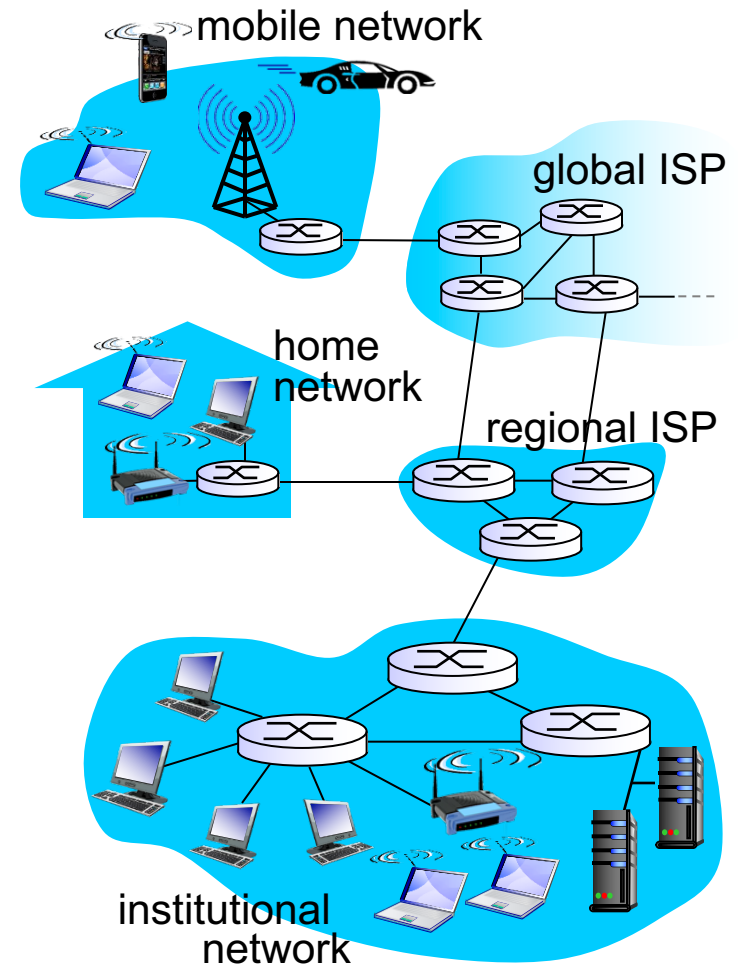
■ transmission rate: *bandwidth*

❖ *routers*: forward packets (chunks of data)



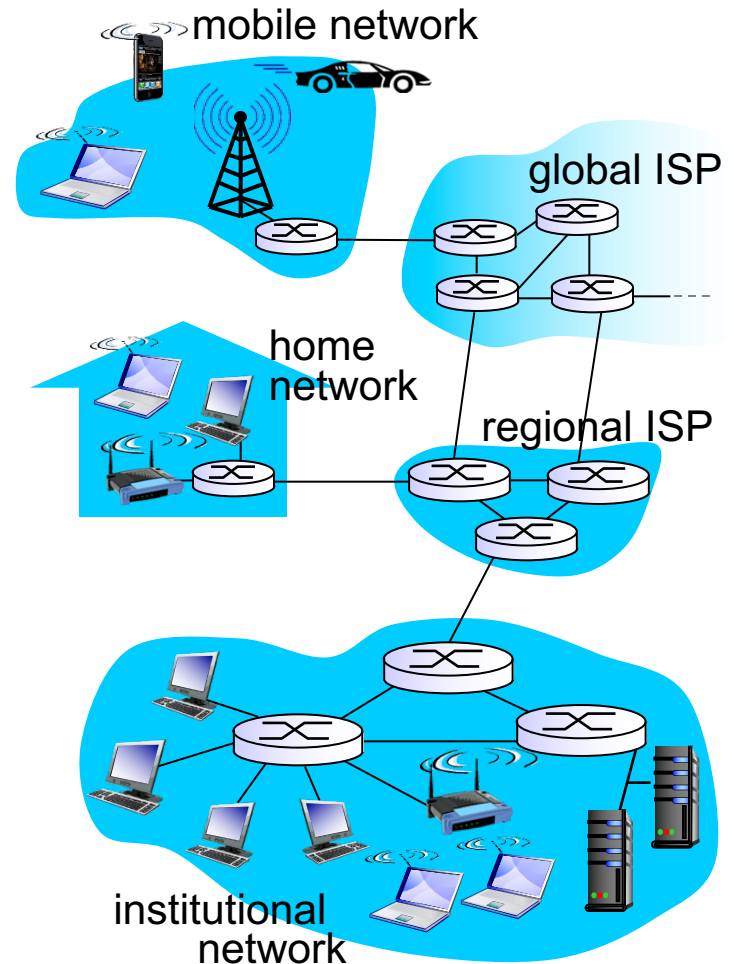
What's the Internet: “nuts and bolts” view

- ❑ *protocols* control sending, receiving of msgs
 - ❖ e.g., TCP, IP, HTTP, Skype, Ethernet
- ❑ *Internet: “network of networks”*
 - ❖ loosely hierarchical
 - ❖ public Internet versus private intranet
- ❑ Internet standards
 - ❖ RFC: Request for comments
 - ❖ IETF: Internet Engineering Task Force



What's the Internet: a service view

- ❑ *communication infrastructure* enables distributed applications:
 - ❖ Web, VoIP, email, games, e-commerce, file sharing
- ❑ *communication services provided to apps:*
 - ❖ reliable data delivery from source to destination
 - ❖ “best effort” (unreliable) data delivery

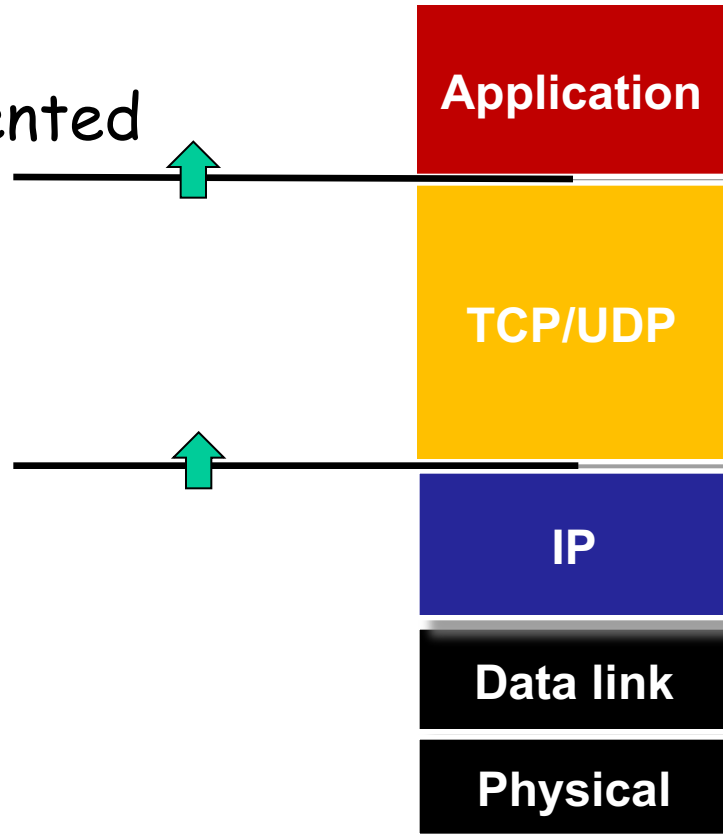


What's the Internet: a service view

Provided Service:

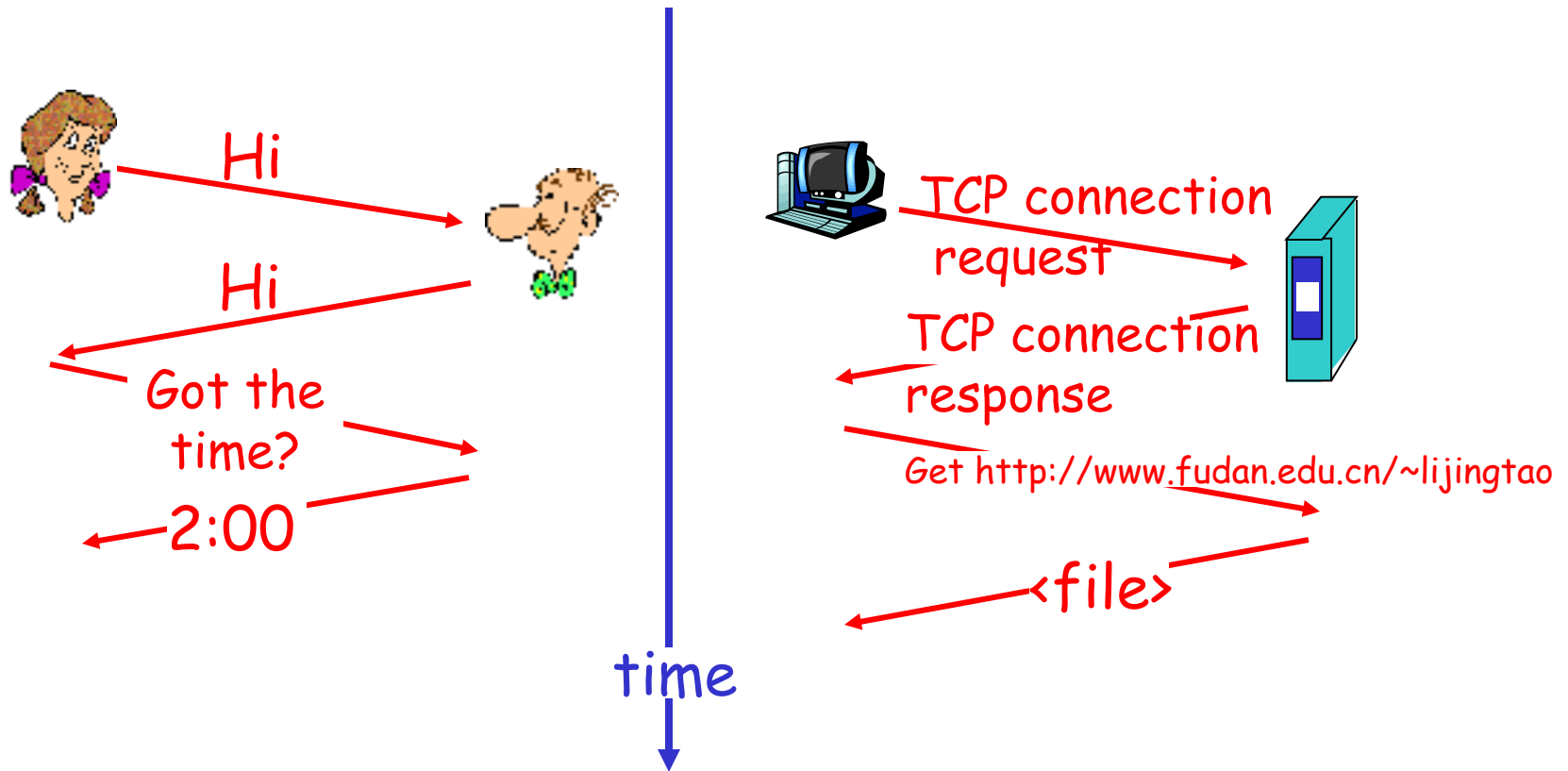
- TCP: connection-oriented
- UDP: connectionless

Provided Service: ?



What's a protocol?

a human protocol and a computer network protocol:



Q: Other human protocols?

What's a protocol?

human protocols:

- ❑ "what's the time?"
- ❑ "I have a question"
- ❑ introductions

... specific msgs sent

... specific actions taken
when msgs received,
or other events

network protocols:

- ❑ machines rather than humans
- ❑ all communication activity in Internet governed by protocols

*protocols define format,
order of msgs sent and
received among network
entities, and actions
taken on msg
transmission, receipt*

Introduction: roadmap

1.1 What *is* the Internet?

1.2 Network edge

1.3 Network core

1.4 Network access and physical media

1.5 Internet structure and ISPs

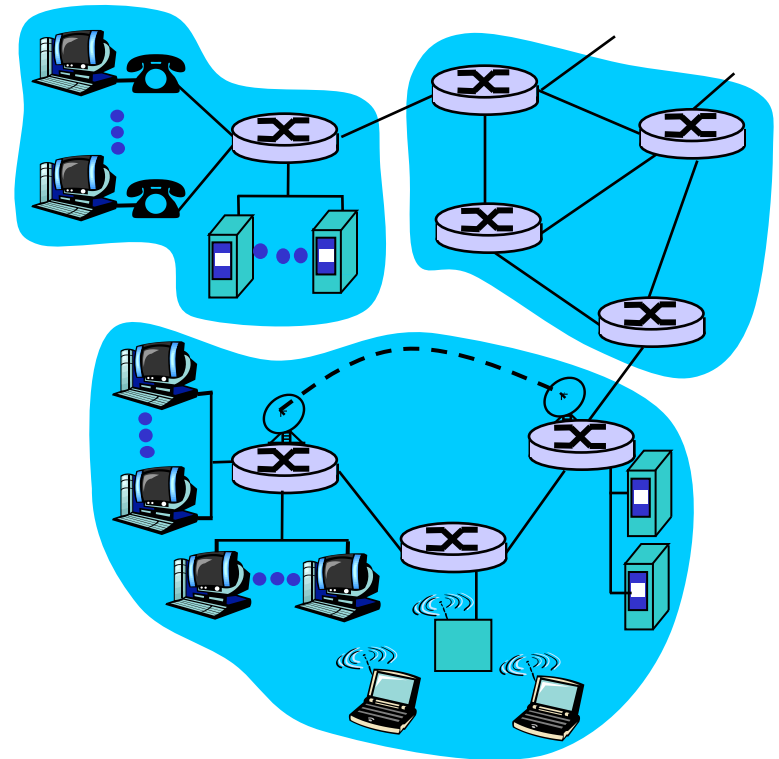
1.6 Delay & loss in packet-switched networks

1.7 Protocol layers, service models

1.8 History

A closer look at network structure:

- ❑ **network edge:**
applications and hosts
- ❑ **network core:**
 - ❖ routers
 - ❖ network of routers
- ❑ **access networks, physical media:**
 - ❖ communication links
 - ❖ routers



The network edge:

❑ end systems (hosts):

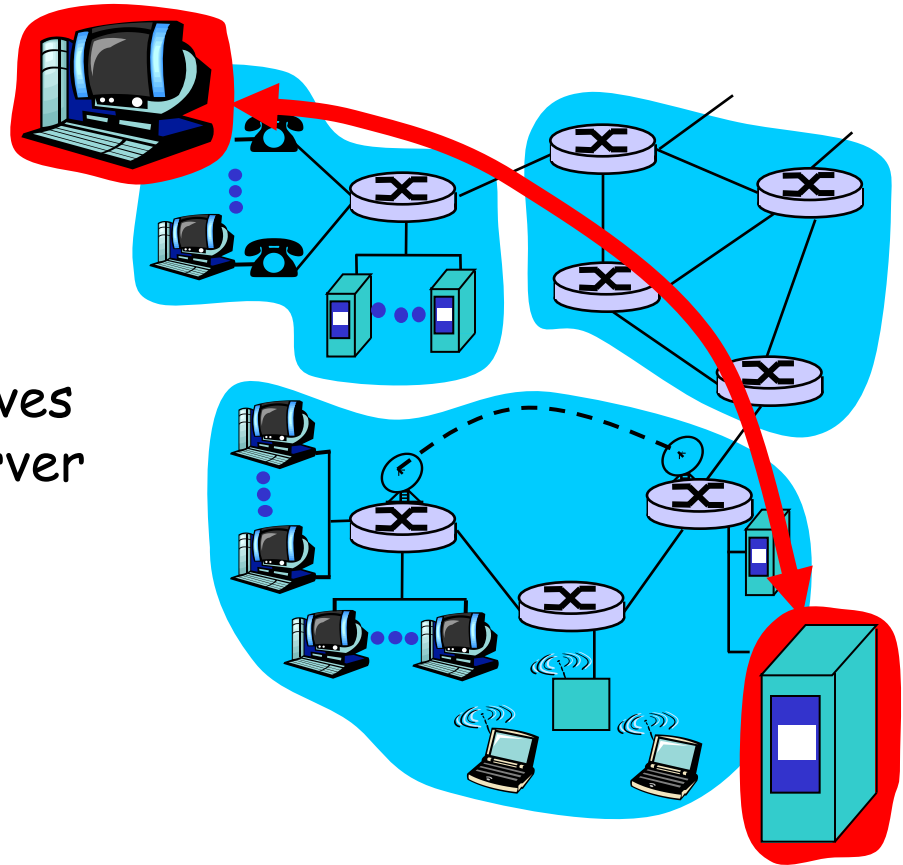
- ❖ run application programs
- ❖ e.g. Web, email
- ❖ at “edge of network”

❑ client/server model

- ❖ client host requests, receives service from always-on server
- ❖ e.g. Web browser/server; email client/server

❑ peer-peer model:

- ❖ minimal (or no) use of dedicated servers
- ❖ e.g. Skype, BitTorrent, KaZaA



Network edge: connection-oriented service

- Goal: data transfer
between end systems
- ❑ *handshaking*: setup
(prepare for) data
transfer ahead of time
 - ❖ Hello, hello back human
protocol
 - ❖ *set up "state"* in two
communicating hosts
 - ❑ TCP - Transmission
Control Protocol
 - ❖ Internet's connection-
oriented service

TCP service [RFC 793]

- ❑ *reliable, in-order* byte-
stream data transfer
 - ❖ loss: acknowledgements
and retransmissions
- ❑ *flow control*:
 - ❖ sender won't overwhelm
receiver
- ❑ *congestion control*:
 - ❖ senders "slow down sending
rate" when network
congested

Network edge: connectionless service

Goal: data transfer
between end systems

- ❖ same as before!

- ❑ **UDP** - User Datagram Protocol [RFC 768]:

- ❖ connectionless
- ❖ unreliable data transfer
- ❖ no flow control
- ❖ no congestion control

App's using TCP:

- ❑ HTTP (Web), FTP (file transfer), Telnet (remote login), SMTP (email)

App's using UDP:

- ❑ streaming media, QQ, teleconferencing, DNS, Internet telephony

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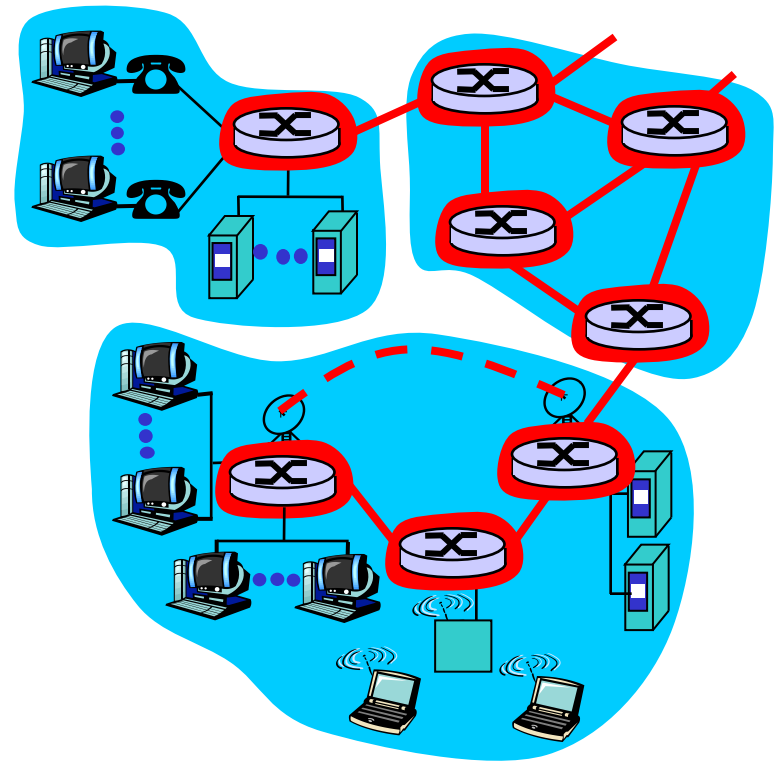
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The Network Core

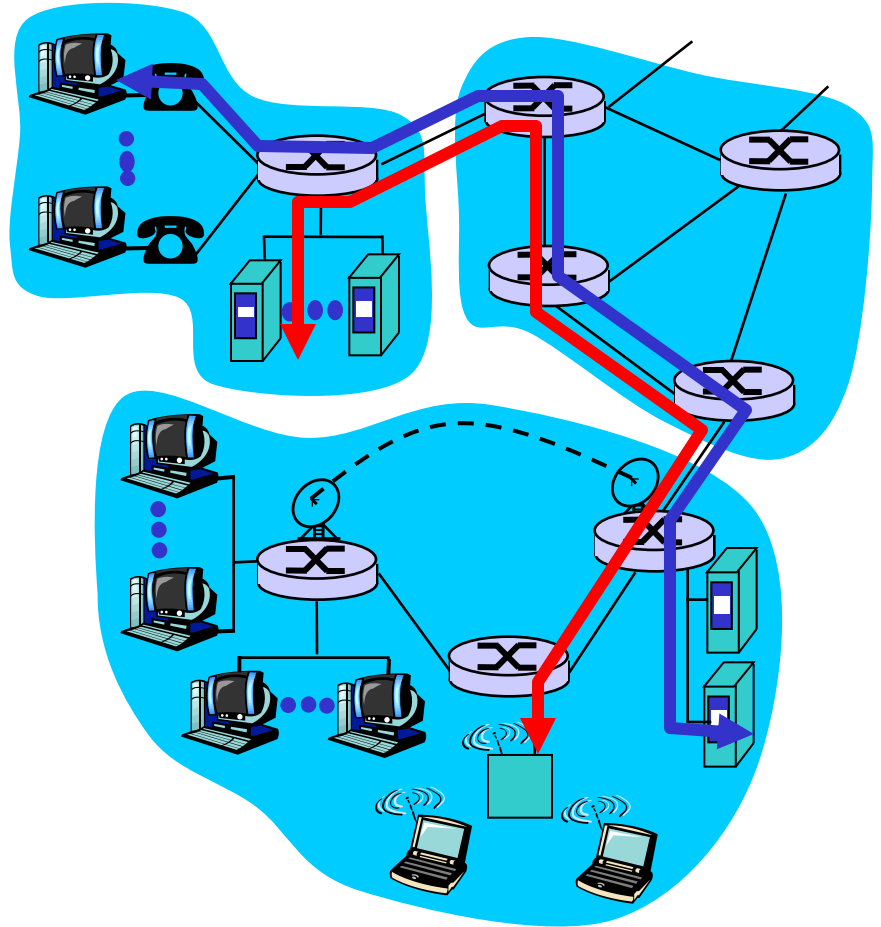
- ❑ mesh of interconnected routers
- ❑ the fundamental question: how is data transferred through net?
 - ❖ circuit switching: dedicated circuit per call: telephone net
 - ❖ packet-switching: data sent thru net in discrete "chunks"



Network Core: Circuit Switching

End-end resources
reserved for "call"

- ❑ link bandwidth, switch capacity
- ❑ dedicated resources: no sharing
- ❑ circuit-like (guaranteed) performance
- ❑ call setup required



Network Core: Circuit Switching

network resources
(e.g., bandwidth)
divided into "pieces"

- ❑ pieces allocated to calls
- ❑ resource piece *idle* if not used by owning call
(*no sharing*)

- ❑ dividing link bandwidth into "pieces"
 - ❖ frequency division
 - ❖ time division

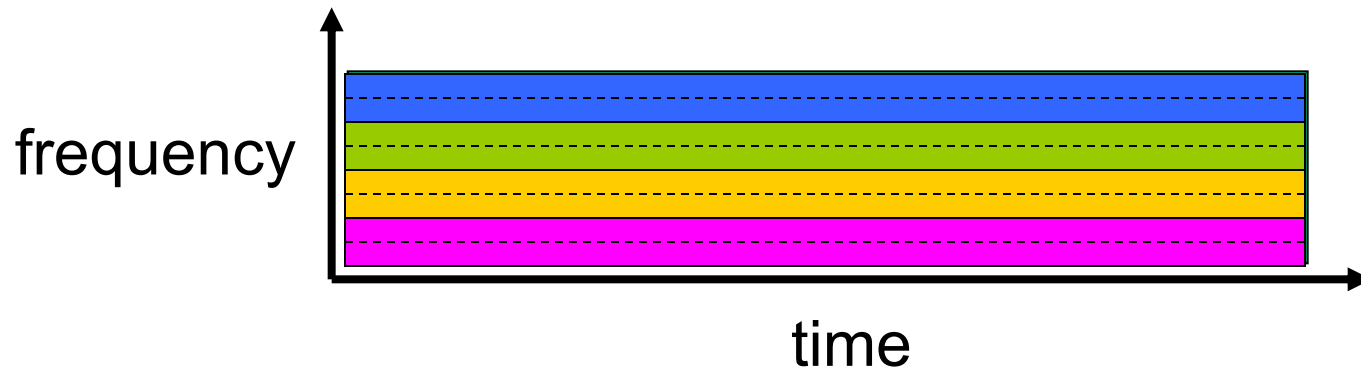
Circuit Switching: FDM and TDM

Example:

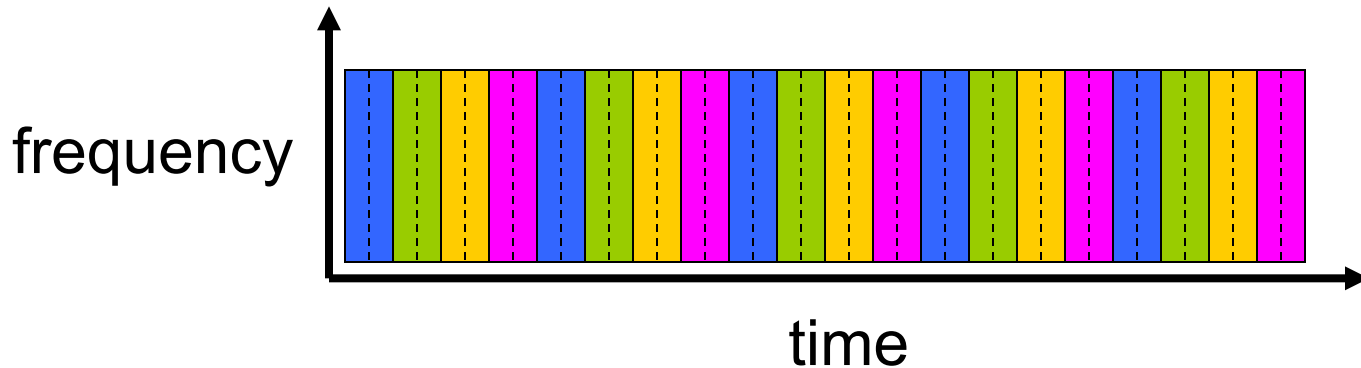
4 users



FDM



TDM



More Examples

GSM (a single cell)

- 4 Carriers --FDM
- 8 time slots for each carrier --TDM

	0	1	2	3	4	5	6	7
TRX 1	CCH*	SDCCH/8	TCH	TCH	TCH	TCH	TCH	TCH
TRX 2	TCH	TCH	TCH	TCH	TCH	TCH	TCH	TCH
TRX 3	TCH	TCH	TCH	TCH	TCH	TCH	TCH	TCH
TRX 4	TCH	TCH	TCH	TCH	TCH	TCH	TCH	TCH

T1/DS1 (Digital Signal Designator)

- ❑ In the T1 system, voice or other analog signals are sampled 8,000 times a second and each sample is digitized into an 8-bit word.
- ❑ With 24 channels being digitized at the same time, a 192-bit/slot frame (24 channels each with an 8-bit word) is thus being transmitted 8,000 times a second.
- ❑ Each frame is separated from the next by a single bit, making a 193-bit block. The 192 bit frame multiplied by 8,000 and the additional 8,000 framing bits make up the T1's 1.544 Mbps data rate.

T1 Frame Format



- Each DS0 called a **time slot**
- $8000 \text{ frames/sec} * 8 \text{ bits/slot} = 64 \text{ Kbps}$
- $24 * 8 + 1 = 193 \text{ bits/frame}$
- $8000 \text{ frames/sec} * 193 \text{ bits/frame} = 1.544 \text{ Mbps}$
- 8000 Framing bits sent per second

Transmission Rates

Digital Signal Designator	Data Rate	DS0 Multiple	T-Carrier
DS0	64 Kbps	1	-
DS1	1.544 Mbps	24	T1
DS3	44.736 Mbps	672	T3
OC3	155.52 Mbps		Optical Carrier
OC12	622.08 Mbps		
OC192	10 Gbps		
OC256	13.271 Gbps		
OC768	40 Gbps		

Numerical example

- How long does it take to send a file of 640,000 bits from host A to host B over a circuit-switched network?
 - ❖ All links are 1.536 Mbps
 - ❖ Each link uses TDM with 24 slots/sec
 - ❖ 500 msec to establish end-to-end circuit

Let's work it out!

Numerical example

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Let's work it out! (10.5sec)

Network Core: Packet Switching

each end-end data stream
divided into *packets*

- ❑ user A, B packets *share* network resources
- ❑ each packet uses full link bandwidth
- ❑ resources used *as needed*

Bandwidth division into "pieces"

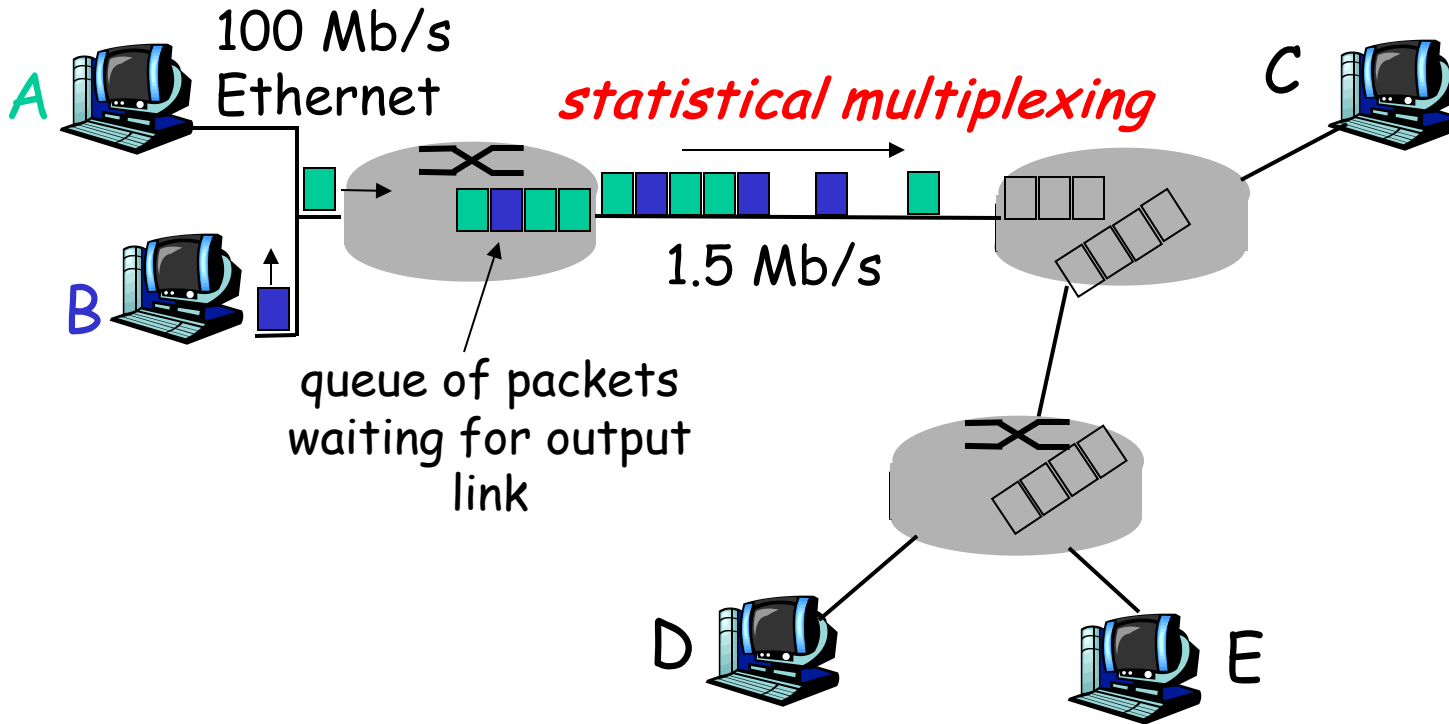
Dedicated allocation

Resource reservation

resource contention:

- ❑ aggregate resource demand can exceed amount available
- ❑ congestion: packets queue, wait for link use
- ❑ store and forward: packets move one hop at a time
 - ❖ Node receives complete packet before forwarding

Packet Switching: Statistical Multiplexing



Sequence of A & B packets does not have fixed pattern, shared on demand ➡ *statistical multiplexing*.

TDM: each host gets same slot in revolving TDM frame.

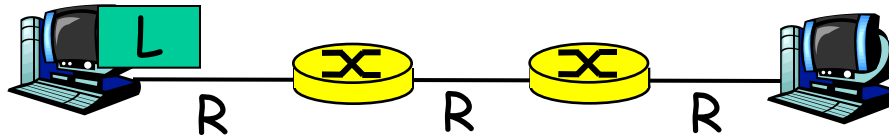
Packet switching versus circuit switching

Is packet switching a "winner?"

- ❑ Great for bursty data
 - ❖ resource sharing
 - ❖ simpler, no call setup
- ❑ **Excessive congestion:** packet delay and loss
 - ❖ protocols needed for reliable data transfer, congestion control
- ❑ **Q: How to provide circuit-like behavior?**
 - ❖ bandwidth guarantees needed for audio/video apps
 - ❖ still an unsolved problem (known as "QoS")

Q: human analogies of reserved resources (circuit switching) versus on-demand allocation (packet-switching)?

Packet-switching: store-and-forward



- Takes L/R seconds to transmit (push out) packet of L bits on to link at R bps
- Entire packet must arrive at router before it can be transmitted on next link: *store and forward*
- $\text{delay} = 3L/R$ (assuming zero propagation delay)

Example:

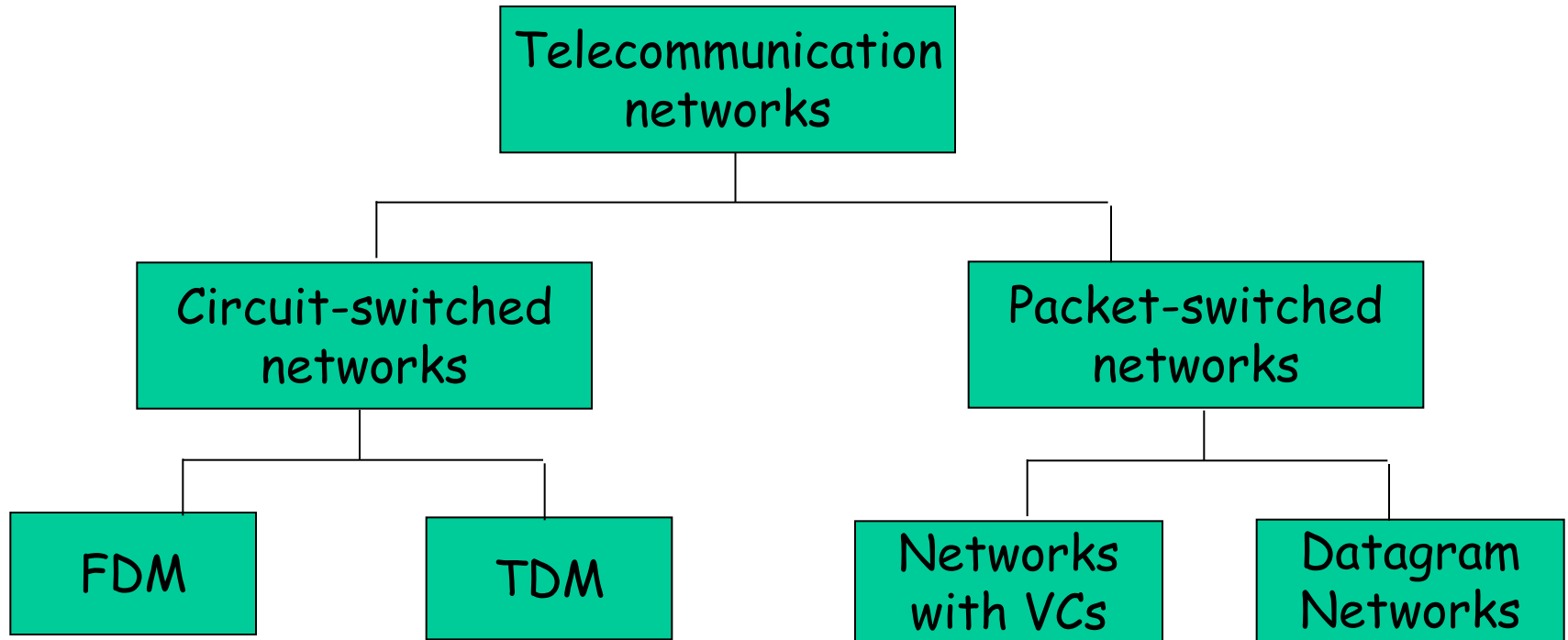
- $L = 7.5$ Mbits
- $R = 1.5$ Mbps
- $\text{delay} = 15$ sec

} more on delay shortly ...

Packet-switched networks: forwarding

- *Goal:* move packets through routers from source to destination
 - ❖ we'll study several path selection (i.e. routing) algorithms (chapter 4)
- **datagram network:**
 - ❖ *destination address* in packet determines next hop
 - ❖ routes may change during session
 - ❖ analogy: driving, asking directions
- **virtual circuit network:**
 - ❖ each packet carries tag (virtual circuit ID), tag determines next hop
 - ❖ fixed path determined at *call setup time*, remains fixed thru call
 - ❖ *routers maintain per-call state*

Network Taxonomy



- Datagram network is not either connection-oriented or connectionless.
- Internet provides both connection-oriented (TCP) and connectionless services (UDP) to apps.

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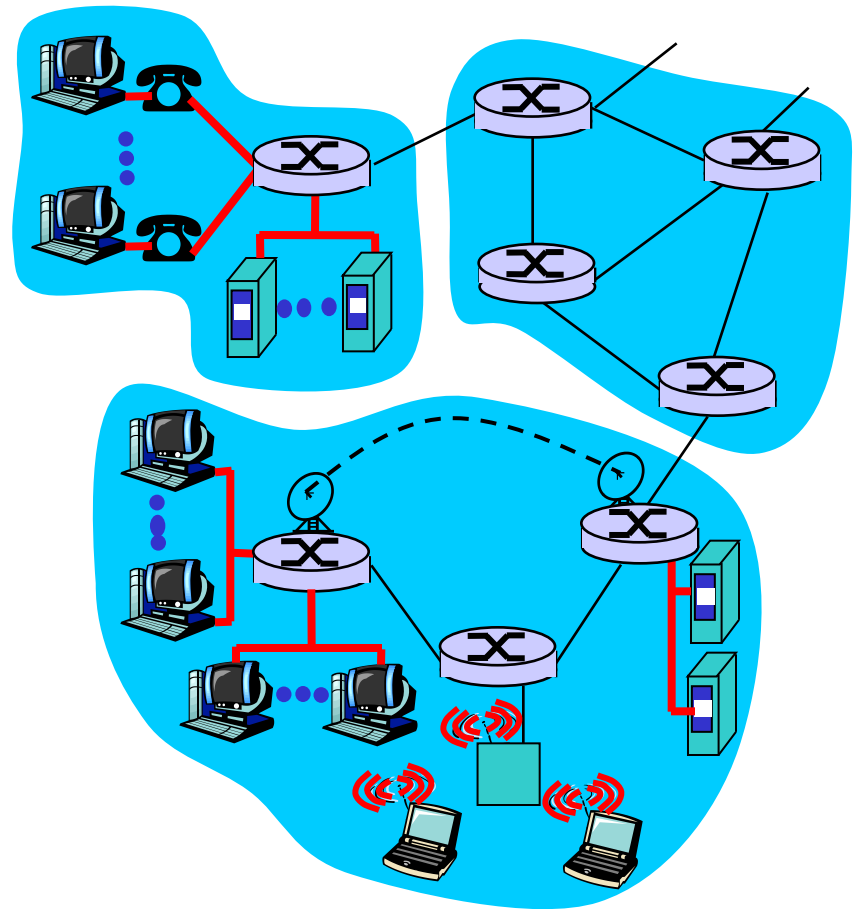
Access networks and physical media

Q: How to connect end systems to edge router?

- ❑ residential access nets
- ❑ institutional access networks (school, company)
- ❑ mobile access networks

Keep in mind:

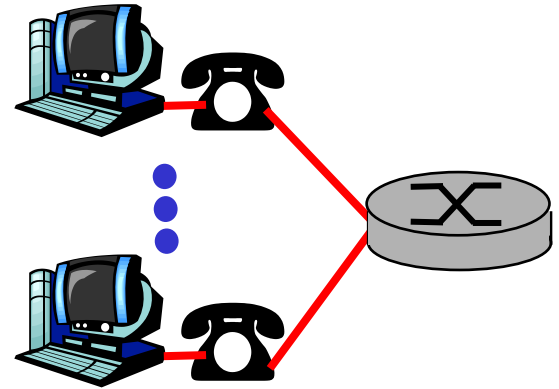
- ❑ bandwidth (bits per second) of access network?
- ❑ shared or dedicated?



Residential access: point to point access

❑ Dialup via modem

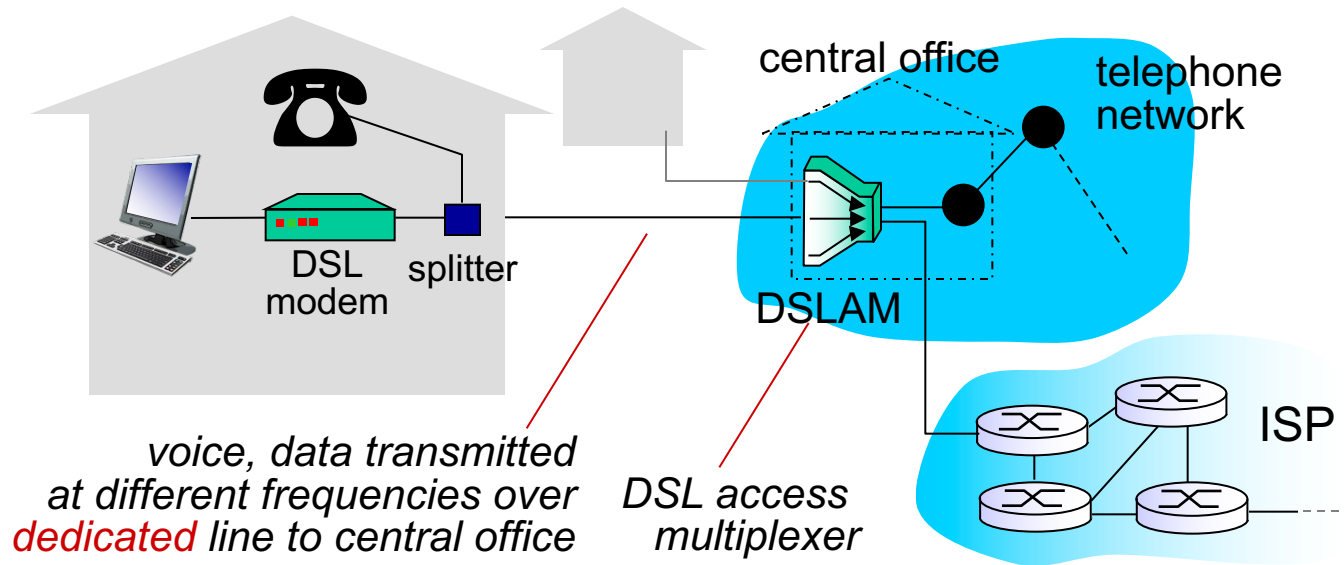
- ❖ up to 56Kbps direct access to router (often less)
- ❖ Can't surf and phone at same time: can't be "always on"



❑ ADSL: asymmetric digital subscriber line

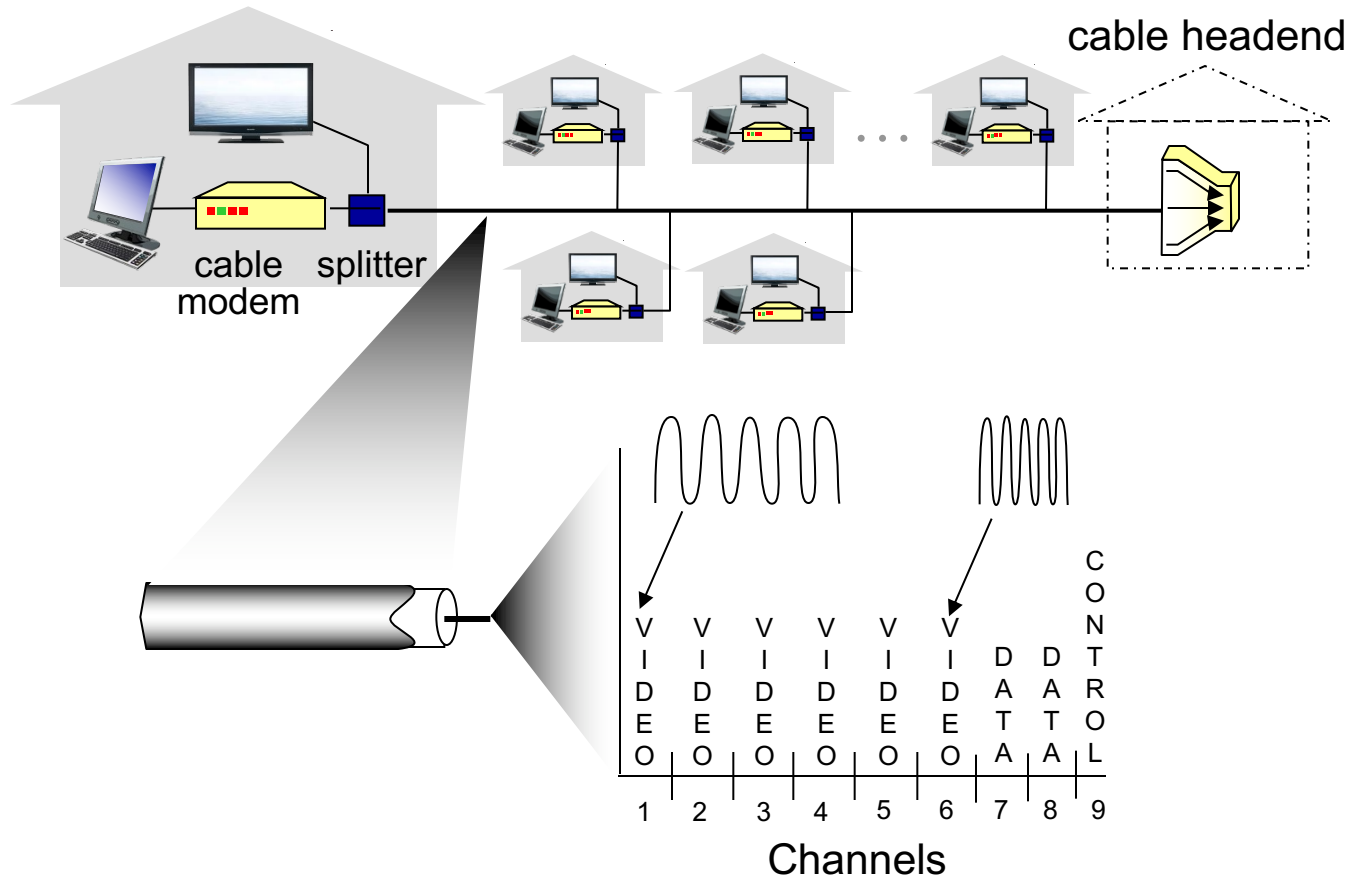
- ❖ FDM: 50 kHz - 1 MHz for downstream
4 kHz - 50 kHz for upstream
0 kHz - 4 kHz for ordinary telephone

Access net: digital subscriber line (DSL)



- ❖ use existing telephone line to central office DSLAM
 - data over DSL phone line goes to Internet
 - voice over DSL phone line goes to telephone net
- ❖ < 2.5 Mbps upstream transmission rate (typically < 1 Mbps)
- ❖ < 24 Mbps downstream transmission rate (typically < 10 Mbps)

Access net: cable network



frequency division multiplexing: different channels transmitted in different frequency bands

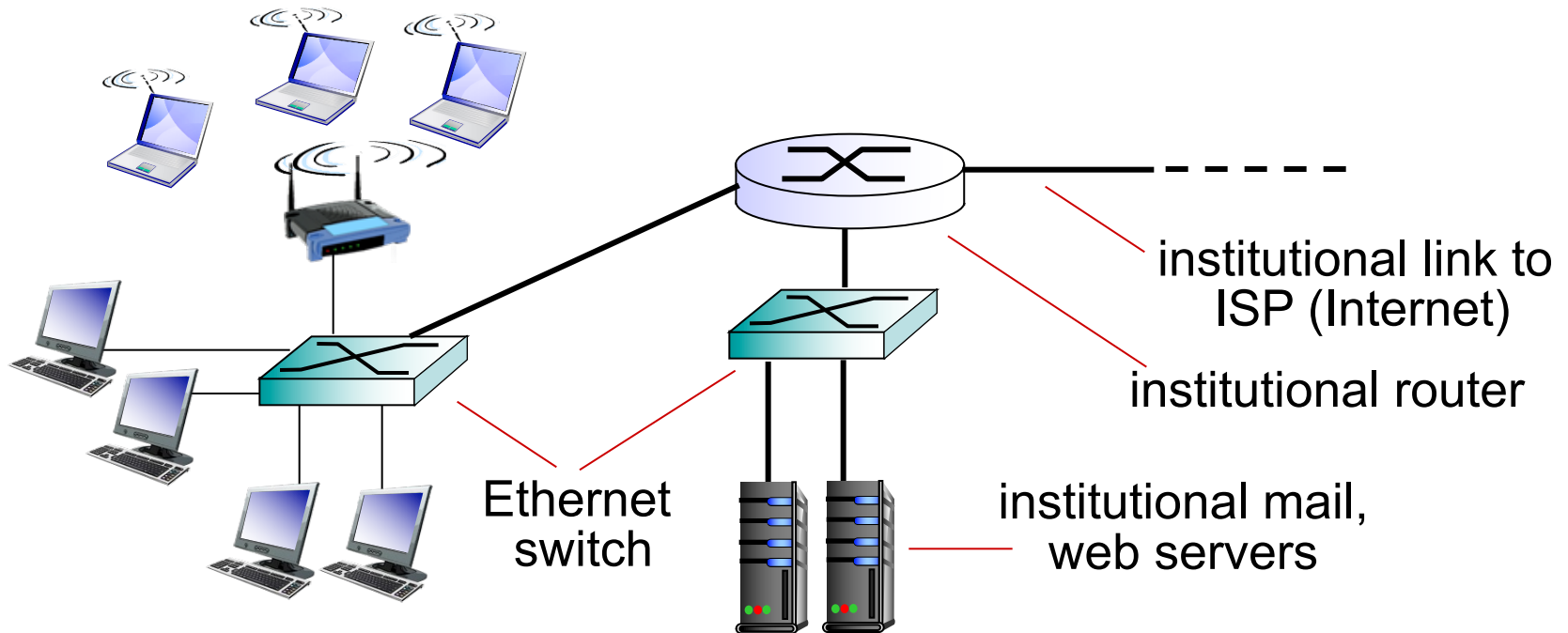
Residential access: cable modems

- ❑ HFC: hybrid fiber coax
 - ❖ asymmetric: up to 30Mbps downstream, 2 Mbps upstream
- ❑ network of cable and fiber attaches homes to ISP router
 - ❖ homes share access to router
- ❑ deployment: available via cable TV companies

Enterprise access nets: local area networks

- ❑ company/univ **local area network** (LAN) connects end system to edge router
- ❑ **Ethernet:**
 - ❖ shared or dedicated link connects end system and router
 - ❖ 10 Mbps, 100Mbps, Gigabit Ethernet

Enterprise access networks



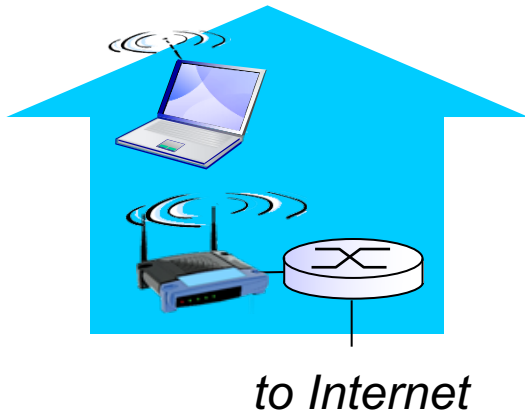
- ❖ 10 Mbps, 100Mbps, 1Gbps, 10Gbps transmission rates
- ❖ today, end systems typically connect into Ethernet switch

Wireless access networks

- ❑ shared *wireless* access network connects end system to router
 - ❖ via base station aka “access point”

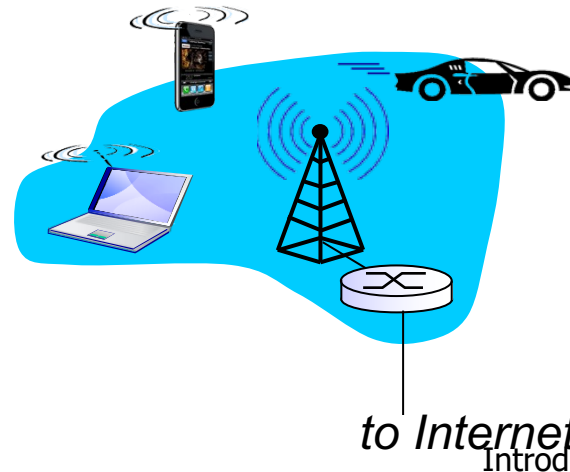
wireless LANs:

- within building (50 m)
- 802.11b/g/n (WiFi): 11, 54, 540Mbps transmission rate



wide-area wireless access

- provided by telco (cellular) operator, 10's Km
- between 1 and X00 Mbps
- 3G, 4G, LTE, WiMax



Home networks

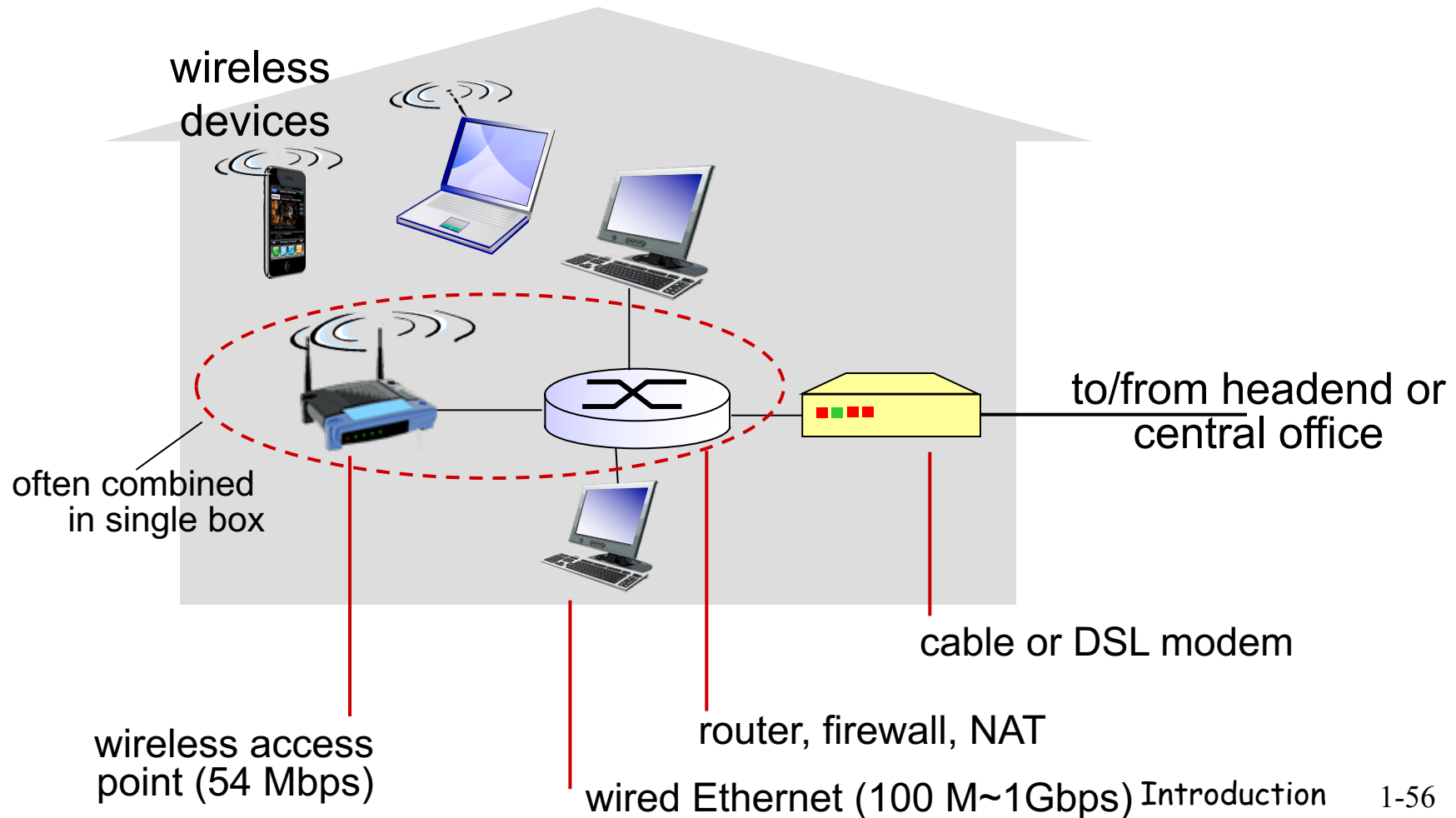
Typical home network components:

□ ADSL, cable modem, or FTTH (eg. EPON)

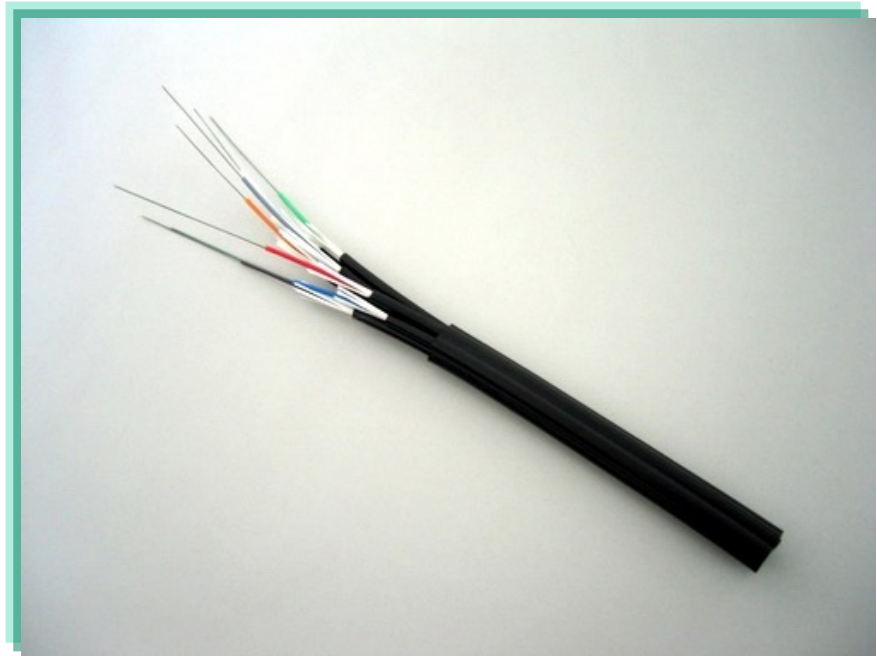
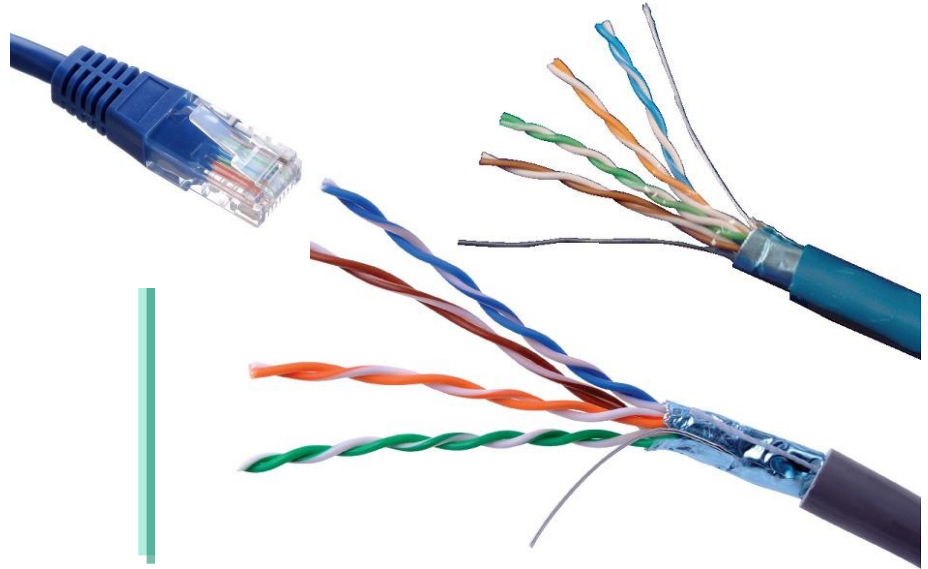
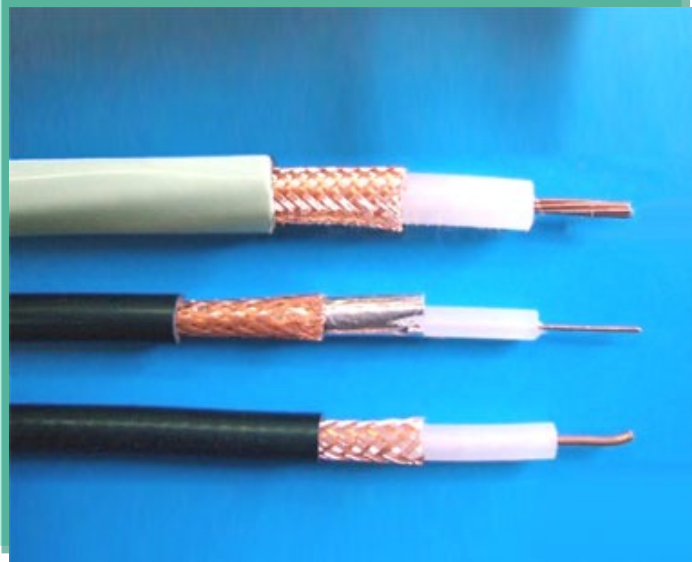
□ router/firewall/NAT

□ Ethernet

□ wireless access point



Physical Media



Physical Media

- ❑ **Bit:** propagates between transmitter/rcvr pairs
- ❑ **physical link:** what lies between transmitter & receiver
- ❑ **guided media:**
 - ❖ signals propagate in solid media: copper, fiber, coax
- ❑ **unguided media:**
 - ❖ signals propagate freely, e.g., radio

Twisted Pair (TP)

- ❑ two insulated copper wires
 - ❖ Category 3: traditional phone wires, 10 Mbps Ethernet
 - ❖ Category 5: 100Mbps Ethernet 1 Gpbs Ethernet
 - ❖ Category 6: 10Gbps



Physical Media: coax, fiber

Coaxial cable:

- ❑ two concentric copper conductors
- ❑ bidirectional
- ❑ baseband:
 - ❖ single channel on cable
 - ❖ legacy Ethernet
- ❑ broadband:
 - ❖ multiple channels on cable
 - ❖ HFC



Fiber optic cable:

- ❑ glass fiber carrying light pulses, each pulse a bit
- ❑ high-speed operation:
 - ❖ high-speed point-to-point transmission (e.g., 10's-100's Gps)
- ❑ low error rate: repeaters spaced far apart ; immune to electromagnetic noise



Physical media: radio

- ❑ signal carried in electromagnetic spectrum
- ❑ no physical "wire"
- ❑ bidirectional
- ❑ propagation environment effects:
 - ❖ reflection
 - ❖ obstruction by objects
 - ❖ interference

Radio link types:

- ❑ **terrestrial microwave**
 - ❖ e.g. up to 45 Mbps channels
- ❑ **LAN** (e.g., Wifi)
 - ❖ 11Mbps, 54 Mbps, 540Mbps...
- ❑ **wide-area** (e.g., cellular)
 - ❖ e.g. 3G: ~2Mbps; 4G: 1xxMbps
- ❑ **satellite**
 - ❖ Kbps to 45Mbps channel (or multiple smaller channels)
 - ❖ 270 msec end-end delay
 - ❖ geosynchronous versus low altitude