

Data Structure

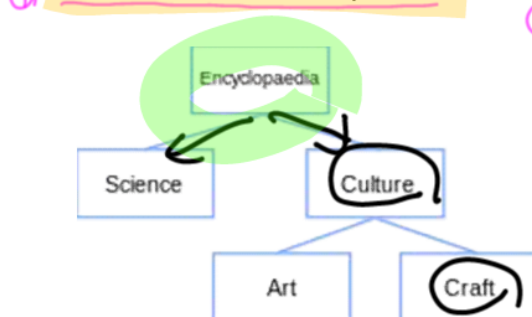
Tree

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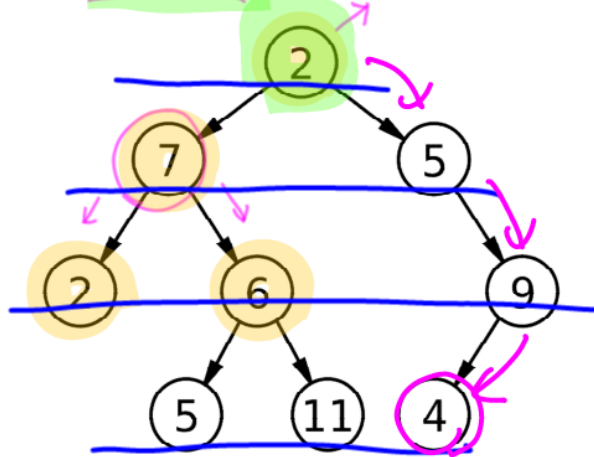
Tree

- ✓ A tree data structure can be defined recursively (locally) as a collection of nodes (starting at a root node), where each node is a data structure consisting of a value, together with a list of references to nodes (the "children"), with the constraints that no reference is duplicated, and none points to the root.



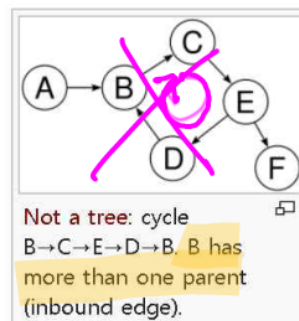
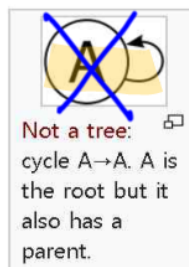
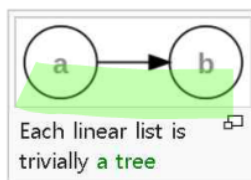
```
struct Node
{
    char name[20];
    Node* left;
    Node* right;
};
```

- ✓ A simple **unordered tree**; in this diagram, the node labeled 7 has two **children**, labeled 2 and 6, and one **parent**, labeled 2. The **root node**, at the top, has no parent.



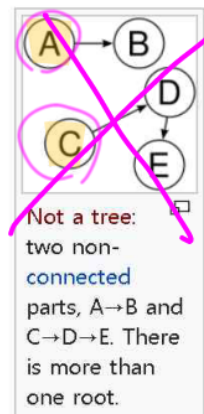
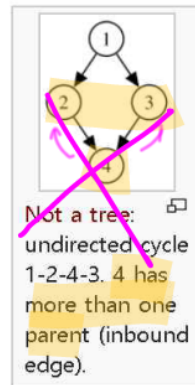
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- ✓ A tree is a data structure made up of **nodes or vertices** and **edges without having any cycle**.
- ✓ The tree with no nodes is called the **null or empty** tree.



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- ✓ A tree that is not empty consists of a **root node** and potentially many **levels** of additional nodes that form a **hierarchy**.



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Terms

- ✓ **Root**
 - The top node in a tree.
- ✓ **Child**
 - A node directly connected to another node when moving away from the root.
- ✓ **Parent**
 - The converse notion of a child.
- ✓ **Siblings**
 - A group of nodes with the same parent.
- ✓ **Descendant**
 - A node reachable by repeated proceeding from parent to child. Also known as subchild.
- ✓ **Ancestor**
 - A node reachable by repeated proceeding from child to parent.

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Terms

- ✓ Leaf
- ✓ External node (not common)
 - A node with no children.
- ✓ Branch node
- ✓ Internal node
 - A node with at least one child.
- ✓ Degree
 - For a given node, its number of children. A leaf is necessarily degree zero.
- ✓ Edge
 - The connection between one node and another.
- ✓ Path
 - A sequence of nodes and edges connecting a node with a descendant.

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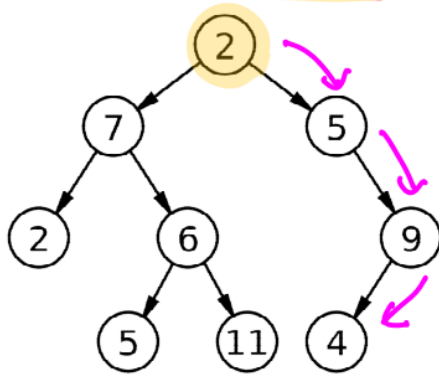
Terms

- ✓ Level
 - The level of a node is defined as: 1 + the number of edges between the node and the root.
- ✓ Depth
 - The depth of a node is defined as: the number of edges between the node and the root.
- ✓ Height of node
 - The height of a node is the number of edges on the longest path between that node and a leaf.
- ✓ Height of tree
 - The height of a tree is the height of its root node.
- ✓ Forest
 - A forest is a set of $n \geq 0$ disjoint trees.

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Binary Tree

- ✓ A **binary tree** is a **tree data structure** in which each node has at most two **children**, which are referred to as the **left child** and the **right child**.



```

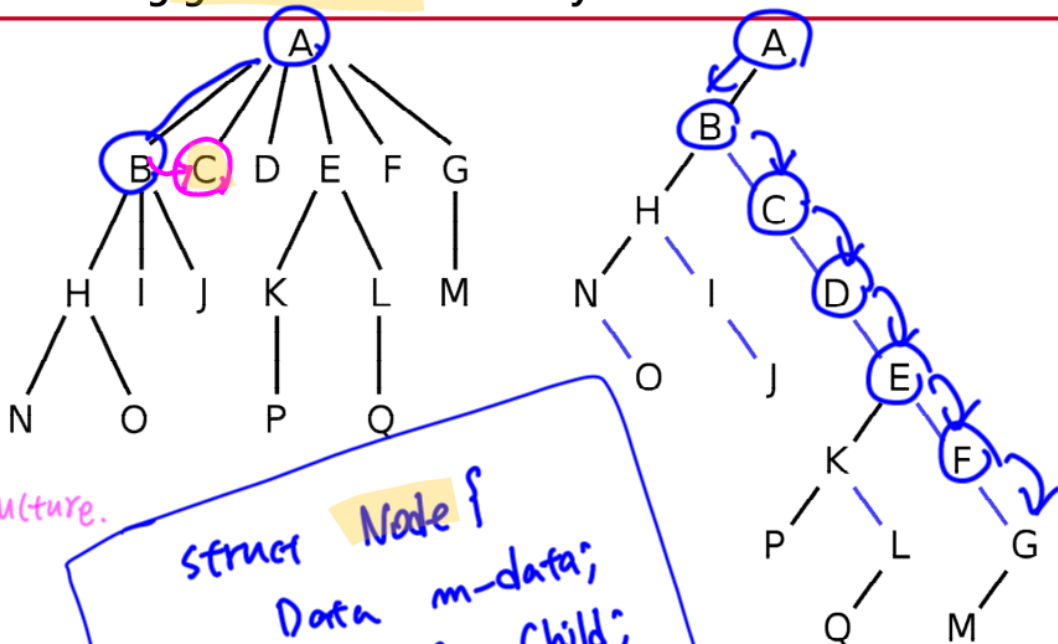
struct Node {
    Node * left;
    Node * right;
    m_data;
}
    
```

- A labeled binary tree of **size 9** and **height 3**, with a **root node** whose value is 2.

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Encoding **general trees** as binary trees



* Data Structure.

```

struct Node {
    Data m_data;
    Node * firstChild;
    Node * nextSibling;
    Node * parent;
}
    
```

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Common operations

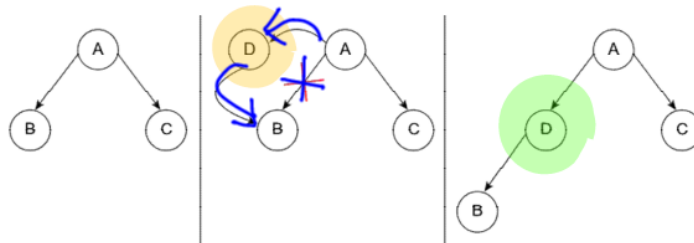
✓ Insertion

- Nodes can be inserted into binary trees in between two other nodes or added after a leaf node.
- In binary trees, a node that is inserted is specified as to which child it is

✓ Leaf nodes

- To add a new node after leaf node A, A assigns the new node as one of its children and the new node assigns node A as its parent.

✓ Internal nodes

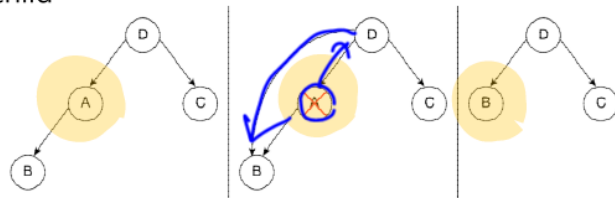


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Deletion

✓ Node with zero or one children

- Suppose that the node to delete is node A.
- If A has no children, deletion is accomplished by setting the child of A's parent to null.
- If A has one child, set the parent of A's child to A's parent and set the child of A's parent to A's child



✓ Node with two children

- In a binary tree, a node with two children cannot be deleted unambiguously.
- However, in certain binary trees (including binary search trees) these nodes *can* be deleted, though with a rearrangement of the tree structure

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next Heapify

Traversal: Depth-first search

- ✓ These searches are referred to as *depth-first search* (DFS), as the search tree is *deepened* as much as possible on each child before going to the next sibling.
- (L) Recursively traverse its *left subtree*. This step is finished at the node N again.
- (R) Recursively traverse its *right subtree*. This step is finished at the node N again.
- (N) Process N itself.

① DFS
② BFS

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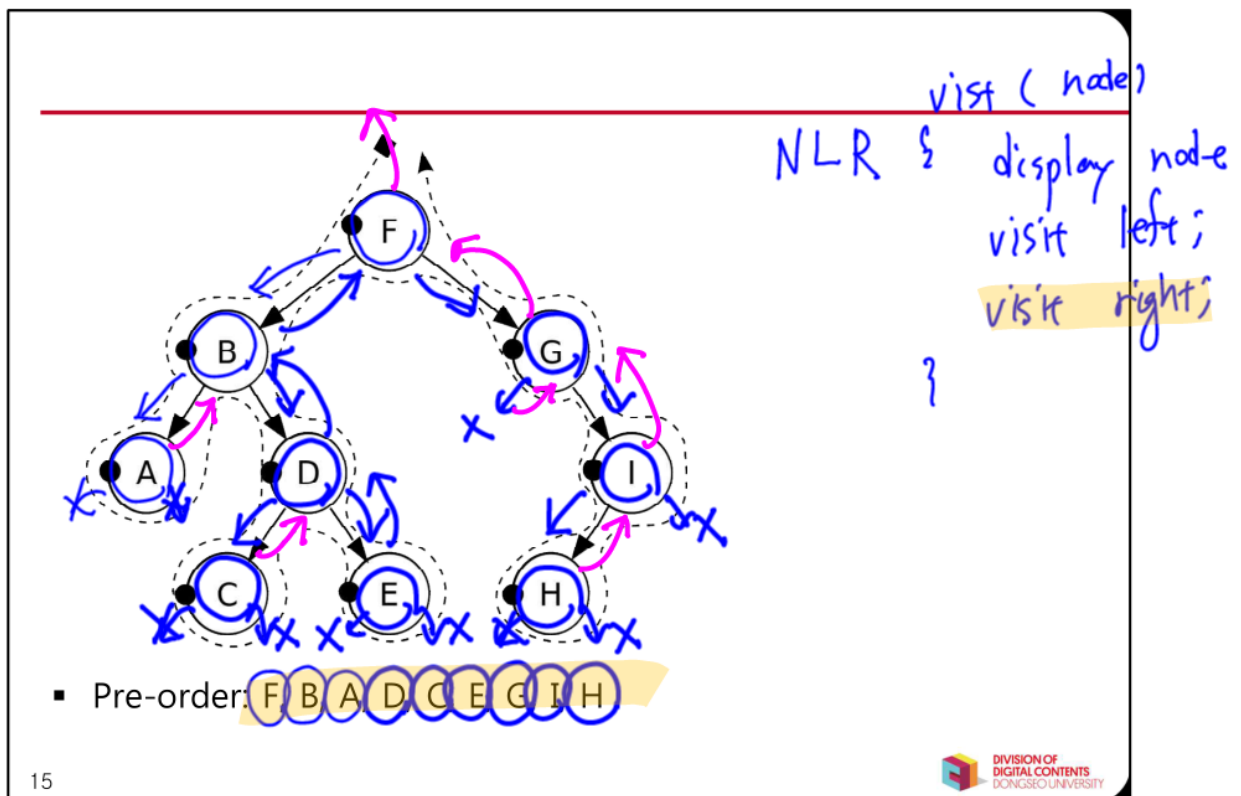


Pre-order(NLR)

1. Check if the *current node is empty or null*.
2. *Display* the data part of the root (or current node).
3. Traverse the *left subtree* by recursively calling the pre-order function.
4. Traverse the *right subtree* by recursively calling the pre-order function.

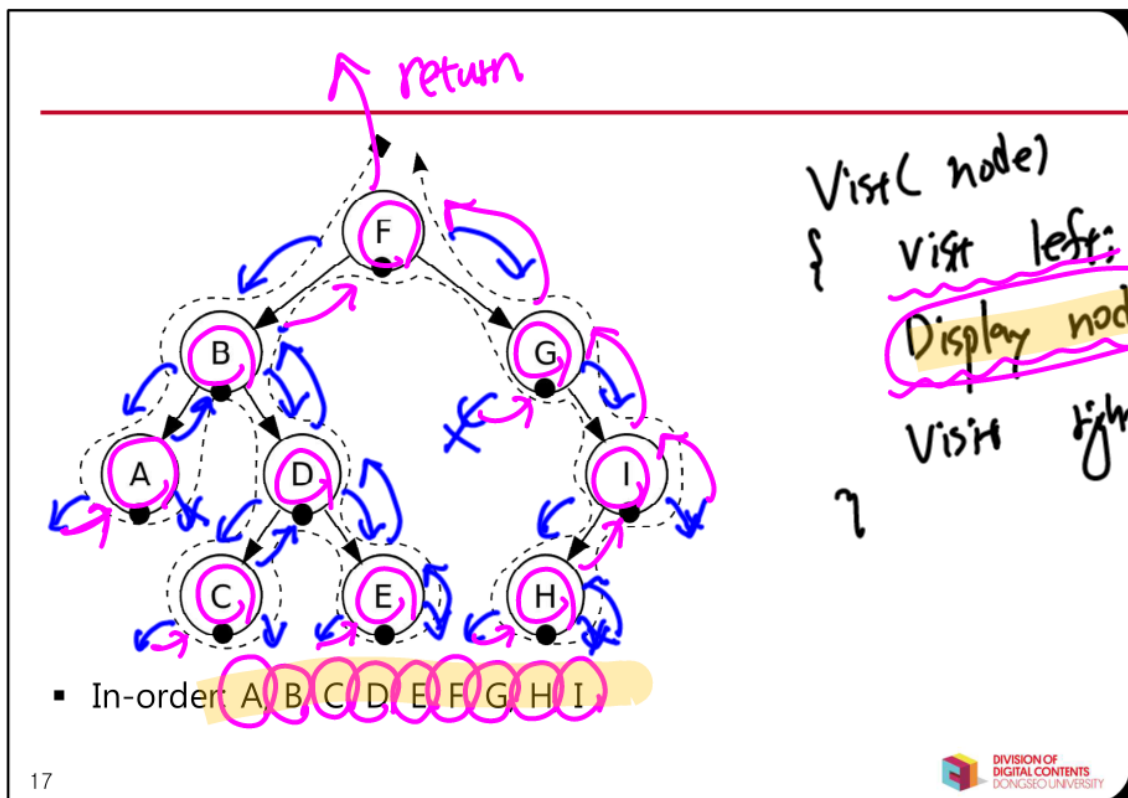
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In-order(LNR)

1. Check if the current node is empty or null.
 2. Traverse the left subtree by recursively calling the in-order function.
 3. Display the data part of the root (or current node).
 4. Traverse the right subtree by recursively calling the in-order function.
- ✓ In a binary search tree, in-order traversal retrieves data in sorted order



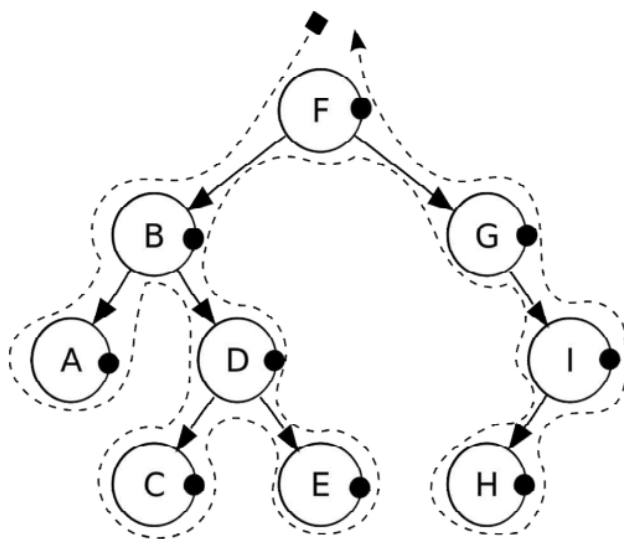
Post-order(LRN)

1. Check if the current node is empty or null.
2. Traverse the left subtree by recursively calling the post-order function.
3. Traverse the right subtree by recursively calling the post-order function.
4. Display the data part of the root (or current node).

visit (node)
{ visit left;
visit right;
Display node;
}

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* Do it!

- Post-order: A, C, E, D, B, H, I, G, F

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Generic tree

1. Perform pre-order operation.
2. For each i from 1 to the number of children do:
 1. Visit i -th, if present.
 2. Perform in-order operation.
3. Perform post-order operation.

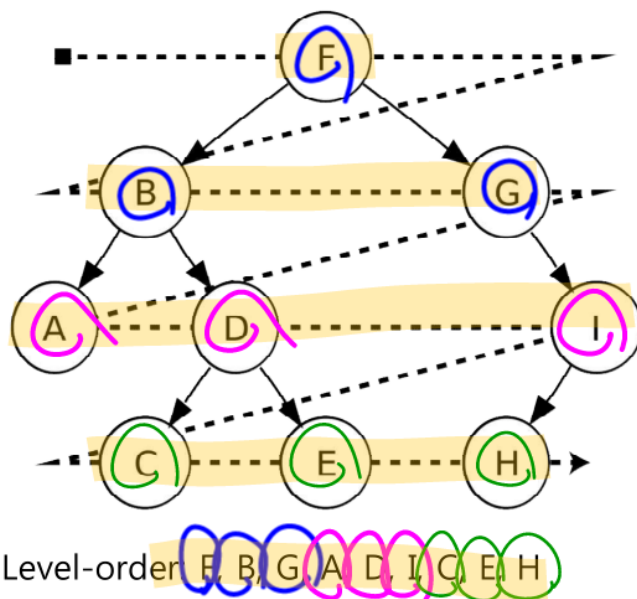
Uniq.

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Breadth-first search

- ✓ Trees can also be traversed in level-order, where we visit every node on a level before going to a lower level.
- ✓ This search is referred to as breadth-first search (BFS), as the search tree is broadened as much as possible on each depth before going to the next depth.

Queue.



Implementation

✓ Pre-order

```
preorder(node)
  if (node = null)
    return
  visit(node)
  preorder(node.left)
  preorder(node.right)
```

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in-order

```
inorder(node)
  if (node = null)
    return
  inorder(node.left)
  visit(node)
  inorder(node.right)
```

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post-order

```
postorder(node)
  if (node = null)
    return

  postorder(node.left)

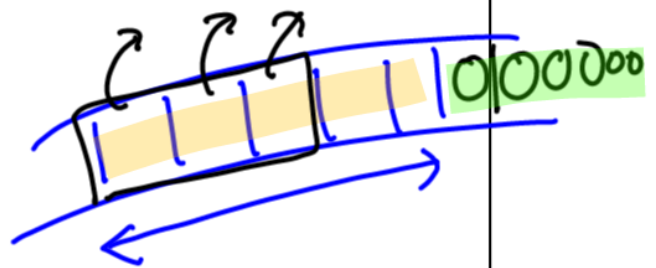
  postorder(node.right)
  visit(node)
```

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level-order

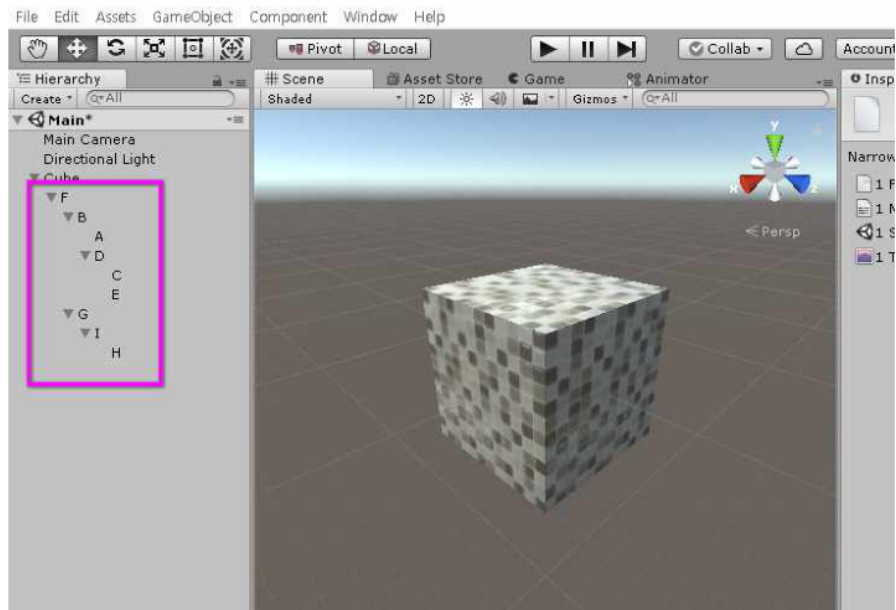
```
levelorder(root)
  q ← empty queue
  q.enqueue(root)
  while (not q.isEmpty())
    node ← q.dequeue()
    visit(node)
    if (node.left ≠ null)
      q.enqueue(node.left)
    if (node.right ≠ null)
      q.enqueue(node.right)
```

BFS



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Practice in Unity



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Pre-order and Post-order

```
void TreeTraverse( Transform go )
{
    foreach( Transform child in go )
    {
        Debug.Log( child.name ); // pre-order
        TreeTraverse( child );
        //Debug.Log( child.name ); // post-order
    }
}

// Use this for initialization
void Start () {
    TreeTraverse( transform );
}
```

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C# generic: Queue

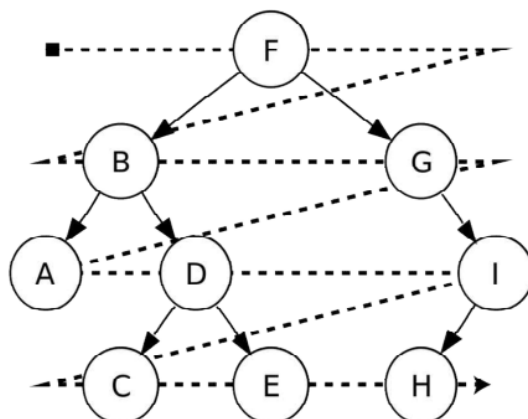
```
Queue<string> numbers = new Queue<string>();  
    numbers.Enqueue("one");  
    numbers.Enqueue("two");  
    numbers.Enqueue("three");  
    numbers.Enqueue("four");  
    numbers.Enqueue("five");  
  
    foreach( string number in numbers )  
    {  
        Console.WriteLine(number);  
    }
```

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Practice: Implement LevelOrder

✓ @see "TreeTraverse_20181203_jintaeks.unitypackage"



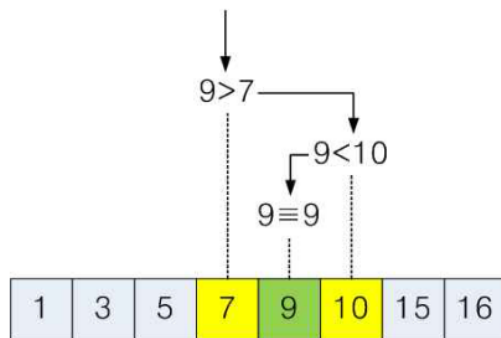
▪ Level-order: F, B, G, A, D, I, C, E, H

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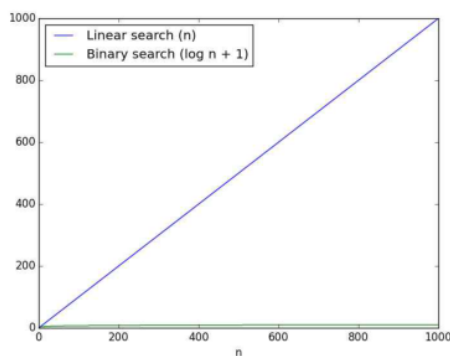
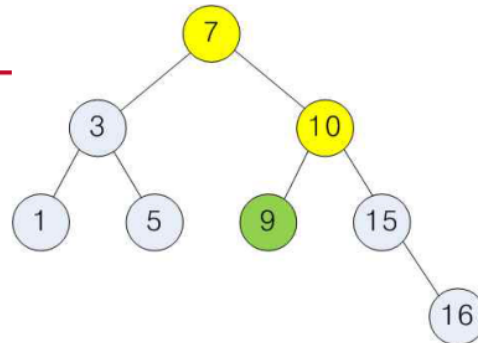
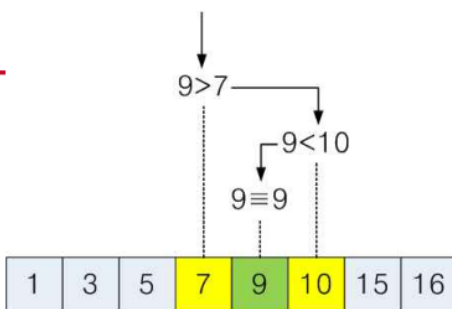


Binary search

- ✓ In [computer science](#), **binary search**, also known as **half-interval search** or **logarithmic search**,^[1] is a [search algorithm](#) that finds the position of a target value within a [sorted array](#).



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: Number of iterations for binary and linear search for n from 1 to 1000.

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QnA

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