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COLLEGE OF SOFTWARE 软件学院
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Chapter 3

Distributed Memory Programming with MPI

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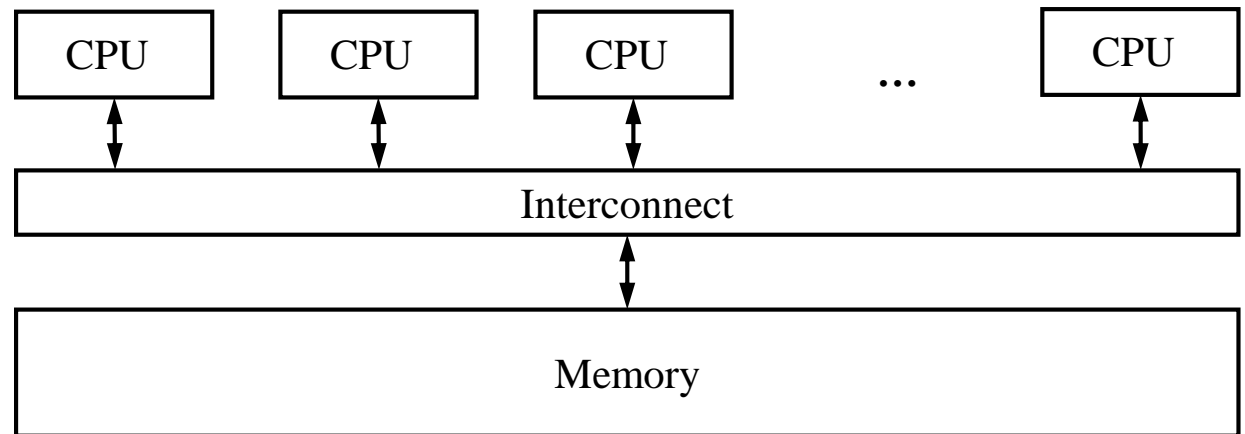
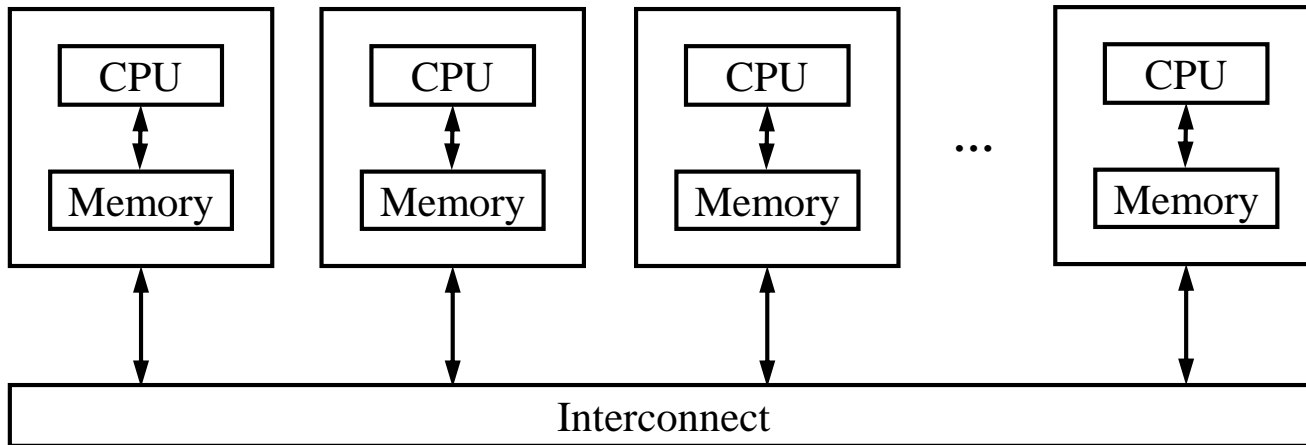
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A distributed memory system



A shared memory system



1. WRITING OUR FIRST MPI PROGRAM

```
#include <stdio.h>
int main(void) {
    printf("hello, world\n");

    return 0;
}
```

(a classic)



Identifying MPI processes

- Common practice to identify processes by nonnegative integer ranks.
- p processes are numbered $0, 1, 2, \dots, p-1$



Our first MPI program

```
1  #include <stdio.h>
2  #include <string.h>    /* For strlen          */
3  #include <mpi.h>       /* For MPI functions, etc */
4
5  const int MAX_STRING = 100;
6
7  int main(void) {
8      char    greeting[MAX_STRING];
9      int     comm_sz;    /* 处理器个数 */
10     int     my_rank;    /* 当前进程 (My process) 序号 */
11
12     MPI_Init(NULL, NULL);
13     MPI_Comm_size(MPI_COMM_WORLD, &comm_sz);
14     MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);
15
16     if (my_rank != 0) {
17         sprintf(greeting, "Greetings from process %d of %d!",
18                 my_rank, comm_sz);
19         MPI_Send(greeting, strlen(greeting)+1, MPI_CHAR, 0, 0,
20                 MPI_COMM_WORLD);
21     } else {
22         printf("Greetings from process %d of %d!\n", my_rank, comm_sz);
23         for (int q = 1; q < comm_sz; q++) {
24             MPI_Recv(greeting, MAX_STRING, MPI_CHAR, q,
25                     0, MPI_COMM_WORLD, MPI_STATUS_IGNORE);
26             printf("%s\n", greeting);
27         }
28     }
29
30     MPI_Finalize();
31     return 0;
32 } /* main */
```



mpi_hello.c

MPI程序基本结构

```
#include <mpi.h>

void main(int argc, char *argv[]) {
    int comm_sz, my_rank, ierr;
    ierr = MPI_Init(&argc, &argv);
    MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);
    MPI_Comm_size(MPI_COMM_WORLD, &comm_sz);
    /* Do some works */
    ierr = MPI_Finalize();
}
```

MPI include file

变量定义

MPI 环境初始化

执行程序
进程间通信

退出MPI环境



Compilation

wrapper script to compile

source file

```
mpicc -g -Wall -o mpi_hello mpi_hello.c
```

produce debugging information

*create this executable file name
(as opposed to default a.out)*

turns on all warnings



Execution

```
mpirexec -n <number of processes> <executable>
```

```
mpirexec -n 1 ./mpi_hello
```

run with 1 process

```
mpirexec -n 4 ./mpi_hello
```

run with 4 processes



Execution

```
mpiexec -n 1 ./mpi_hello
```



Greetings from process 0 of 1 !

```
mpiexec -n 4 ./mpi_hello
```



Greetings from process 0 of 4 !

Greetings from process 1 of 4 !

Greetings from process 2 of 4 !

Greetings from process 3 of 4 !



MPI Programs

- Written in C.
 - Has main.
 - Uses `stdio.h`, `string.h`, etc.
- Need to add **mpi.h** header file.
- Identifiers defined by MPI start with “MPI_”.
- First letter following underscore is uppercase.
 - For function names and MPI-defined types.
 - Helps to avoid confusion.



MPI Components

■ MPI_Init

- Tells MPI to do all the necessary setup.

```
int MPI_Init(  
    int*      argc_p    /* in/out */,  
    char***   argv_p    /* in/out */);
```

■ MPI_Finalize

- Tells MPI we're done, so clean up anything allocated for this program.

```
int MPI_Finalize(void);
```



Basic Outline

```
...  
#include <mpi.h>  
...  
int main(int argc, char* argv[]) {  
    ...  
    /* No MPI calls before this */  
    MPI_Init(&argc, &argv);  
    ...  
    MPI_Finalize();  
    /* No MPI calls after this */  
    ...  
    return 0;  
}
```



Communicators

- A collection of processes that can send messages to each other.
- MPI_Init defines a communicator that consists of all the processes created when the program is started.
- Called **MPI_COMM_WORLD**.



Communicators



```
int MPI_Comm_size(  
    MPI_Comm comm      /* in */,  
    int* comm_sz_p     /* out */);
```

number of processes in the communicator

```
int MPI_Comm_rank(  
    MPI_Comm comm      /* in */,  
    int* my_rank_p     /* out */);
```

*my rank
(the process making this call)*



SPMD

- **Single-Program Multiple-Data**
- We compile one program.
- **Process 0 does something different.**
 - Receives messages and prints them while the other processes do the work.
- The **if-else** construct makes our program SPMD.

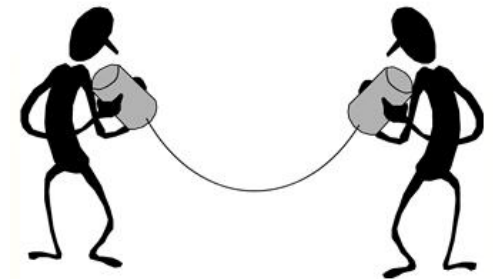


Communication

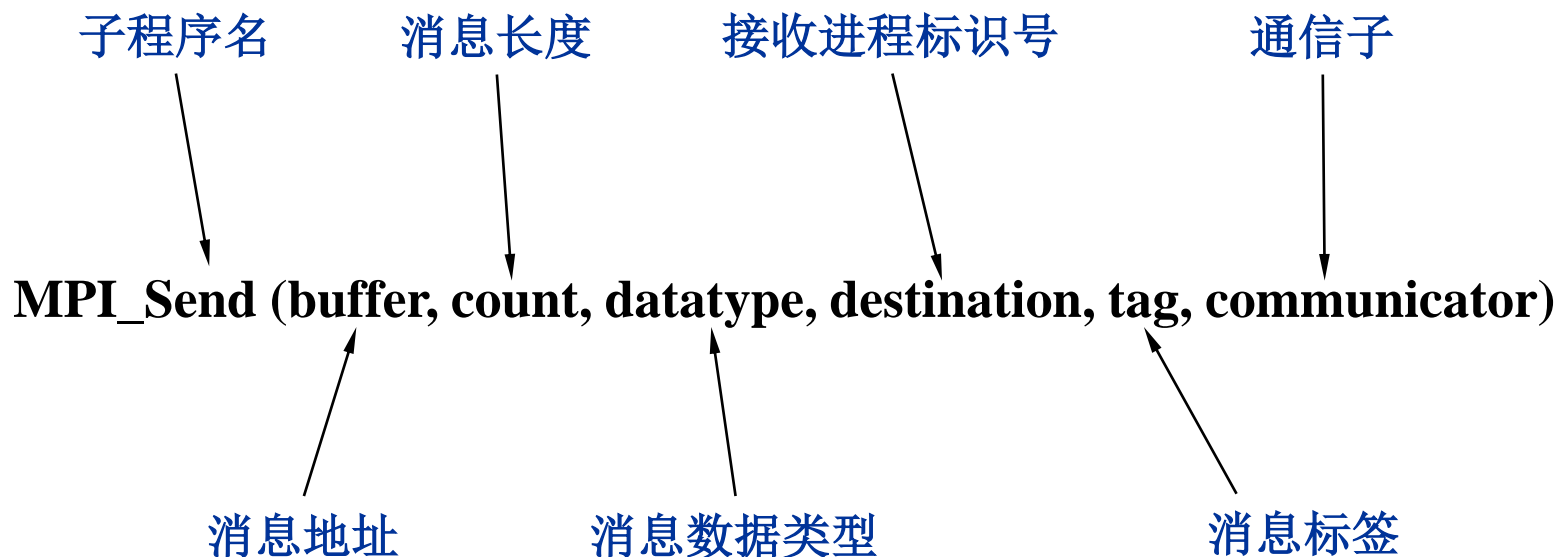
```
int MPI_Send(  
    void*      msg_buf_p      /* in */,  
    int        msg_size       /* in */,  
    MPI_Datatype msg_type      /* in */,  
    int        dest           /* in */,  
    int        tag            /* in */,  
    MPI_Comm   communicator   /* in */);
```

消息缓冲

消息信封



Message



快递详情单

025800001

寄件人姓名		始发地		收件人姓名		目的地	
单位名称				单位名称			
寄件人详细地址				收件人详细地址			
联系电话 (非常重要)		手机		联系电话 (非常重要)		手机	
<input type="checkbox"/> 文件 <input type="checkbox"/> 物品		如系物品, 请如实填写内件名称及数量。 如需保价, 请如实申报保价金额并交纳保价费。		付款方式 <input type="checkbox"/> 现金 <input type="checkbox"/> 到付 <input type="checkbox"/> 协议结算			
重量	千克	体积	长 × 宽 × 高 = 厘米	<input type="checkbox"/> 保价 保价金额	万 仟 佰 拾 元 (大写)		
内件品名及数量		特别注意: 请阅读背面快递服务协议, 贵重物品请保价, 未保价物品的理赔限额为运费的5倍。		代收货款 ¥	万 仟 佰 拾 元 (大写)		
寄件人签名:		寄件人员签章:		运费 ¥	加急费 ¥	包装费 ¥	保价费 ¥
证件号		年 月 日 时		收件人签名:	证件号	年 月 日 时	
025800001		填写本详单前, 请务必阅读背面快递服务协议! 使用本单表示您理解并接受协议内容。		我们不是无所不能, 但一定竭尽所能!			

备注:



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Data types

MPI datatype	C datatype
MPI_CHAR	signed char
MPI_SHORT	signed short int
MPI_INT	signed int
MPI_LONG	signed long int
MPI_LONG_LONG	signed long long int
MPI_UNSIGNED_CHAR	unsigned char
MPI_UNSIGNED_SHORT	unsigned short int
MPI_UNSIGNED	unsigned int
MPI_UNSIGNED_LONG	unsigned long int
MPI_FLOAT	float
MPI_DOUBLE	double
MPI_LONG_DOUBLE	long double
MPI_BYTE	
MPI_PACKED	

为什么要使用消息标签(Tag)?

为了说明为什么要用标签，我们先来看右面一段没有使用标签的代码。

这段代码打算传送A的前32个字节进入X，传送B的前16个字节进入Y。但是，如果消息B尽管后发送但先到达进程Q，就会被第一个recv()接收在X中。

使用标签可以避免这个错误。

未使用标签

Process P:

```
send(A, 32, Q)
send(B, 16, Q)
```

Process Q:

```
recv(X, 32, P)
recv(Y, 16, P)
```

使用了标签

Process P:

```
send(A, 32, Q, tag1)
send(B, 16, Q, tag2)
```

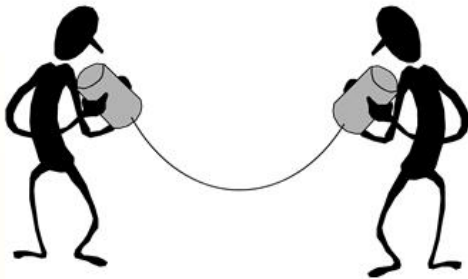
Process Q:

```
recv (X, 32, P, tag1)
recv (Y, 16, P, tag2)
```



Communication

```
int MPI_Recv(  
    void*          msg_buf_p    /* out */,  
    int            buf_size     /* in  */,  
    MPI_Datatype    buf_type     /* in  */,  
    int            source        /* in  */,  
    int            tag           /* in  */,  
    MPI_Comm        communicator /* in  */,  
    MPI_Status*     Status_p     /* out */);
```



Message matching

```
MPI_Send(send_buf_p, send_buf_sz, send_type, dest, send_tag,  
send_comm);
```

MPI_Send

src = q



MPI_Recv

dest = r

```
MPI_Recv(recv_buf_p, recv_buf_sz, recv_type, src, recv_tag,  
recv_comm, &status);
```



Message matching “wildcard” arguments

- It's possible that the receiving process doesn't know the order in which the messages will be sent, and one process can receive multiple messages with different tags from another process.

```
for (i = 1; i < comm_sz; i++) {  
    MPI_Recv(result, result_sz, result_type, MPI_ANY_SOURCE,  
             MPI_ANY_TAG, comm, MPI_STATUS_IGNORE);  
    Process_result(result);  
}
```

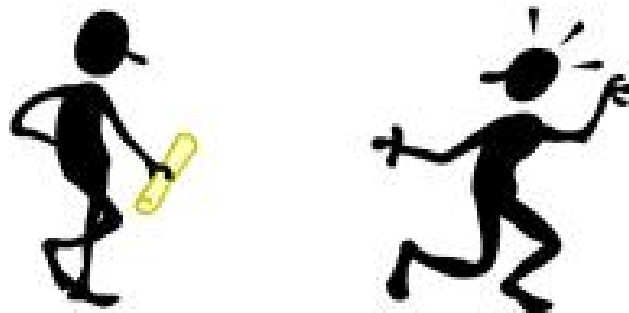
recv_tag (points to MPI_ANY_TAG)

q (points to MPI_ANY_SOURCE)



Receiving messages

- A receiver can get a message without knowing:
 - the amount of data in the message,
 - the sender of the message,
 - or the tag of the message.



status_p argument

```
MPI_Recv(recv_buf_p, recv_buf_sz, recv_type, src, recv_tag,  
recv_comm, &status);
```

MPI_Status*



```
MPI_Status* status;
```

```
status.MPI_SOURCE
```

```
status.MPI_TAG
```

```
status.MPI_ERROR
```

MPI_SOURCE

MPI_TAG

MPI_ERROR



How much data am I receiving?

```
int MPI_Get_count(  
    MPI_Status*   Status_p   /* in */,  
    MPI_Datatype  type       /* in */,  
    int*          count_p    /* out */);
```



Issues with send and receive

- Exact behavior is determined by the MPI implementation.
- MPI_Send may behave differently with regard to buffer size, cutoffs and blocking.
- MPI_Recv always blocks until a matching message is received.
- Know your implementation; don't make assumptions!



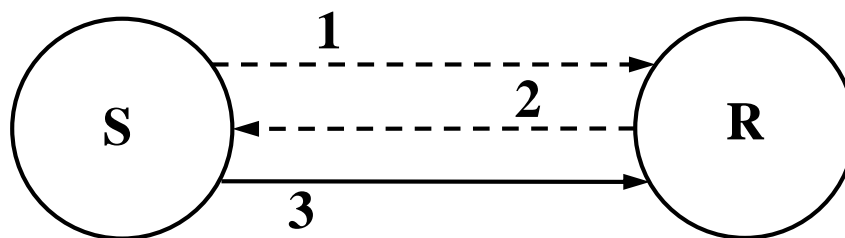
MPI中的四种通信模式 (communication mode)

① 同步的(synchronous)

发送进程直到相应的接收过程已经启动才返回，因此接收端要有存放到达消息的应用缓冲。

注意：在MPI中可以有非阻塞的同步发送，它的返回不意味着消息已经被发出！它的实现不需要在接收端有附加的缓冲，但需要在发送端有一个系统缓冲。为了消除额外的消息拷贝，应使用阻塞的同步发送。

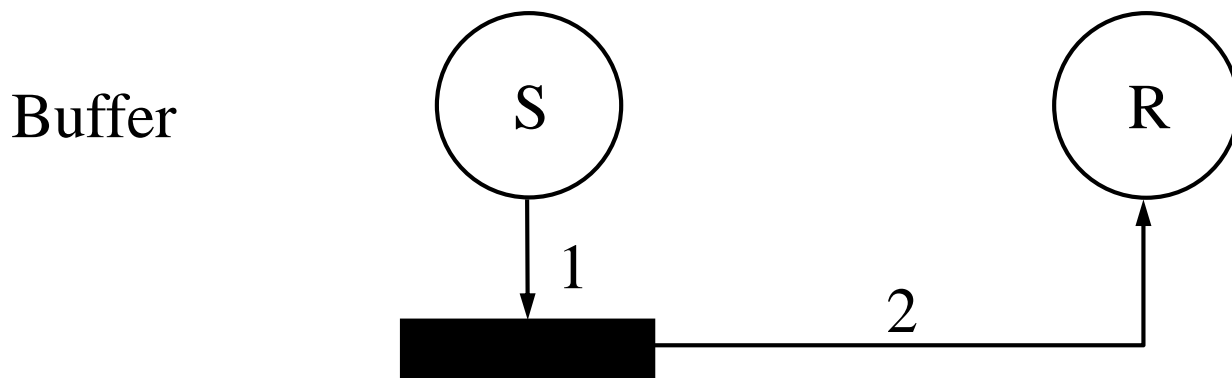
Synchronous



MPI中的四种通信模式

② 缓冲的(buffered)

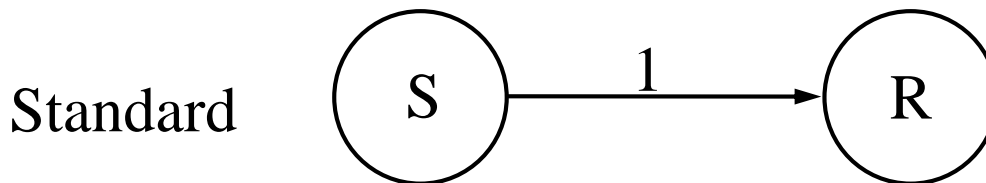
缓冲的发送假定能得到一定大小的缓冲空间，它必须事先由用户程序通过调用子例程MPI_Buffer_attach(buffer, size)来定义，由它来分配大小为size的用户缓冲。这个缓冲可以用MPI_Buffer_detach(*buffer, *size)来实现。



MPI中的四种通信模式

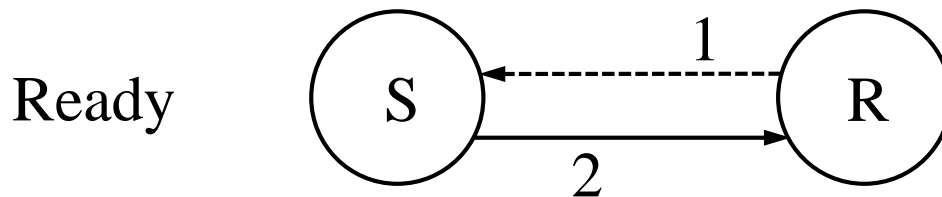
③ 标准的(standard)

发送可以是同步的或缓冲的，取决于实现。



④ 就绪的(ready)

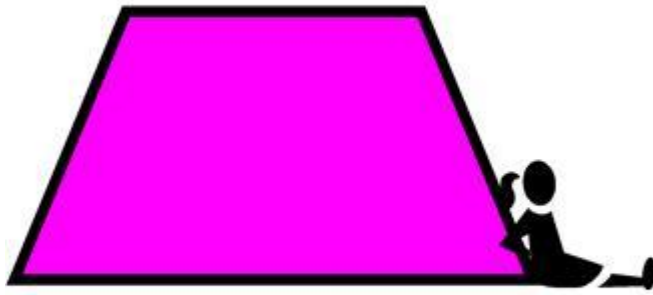
在肯定相应的接收已经开始才进行发送。它不像在同步模式中那样需要等待。这就允许在相同的情况下实际使用一个更有效的通信协议。



MPI点对点通信中不同的发送和接收操作

通信模式	MPI Primitive	Blocking	Non-Blocking
标准通信	Standard Send	MPI_Send	MPI_Isend
同步通信	Synchronous Send	MPI_Ssend	MPI_Issend
缓冲通信	Buffered Send	MPI_Bsend	MPI_Ibsend
就绪通信	Ready Send	MPI_Rsend	MPI_Irsend
	Receive	MPI_Recv	MPI_Irecv
	Completion Check	MPI_Wait	MPI_Test

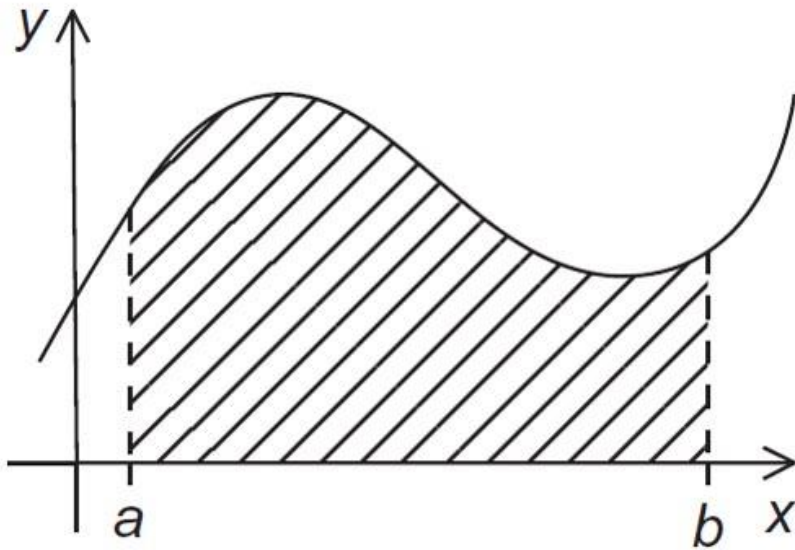




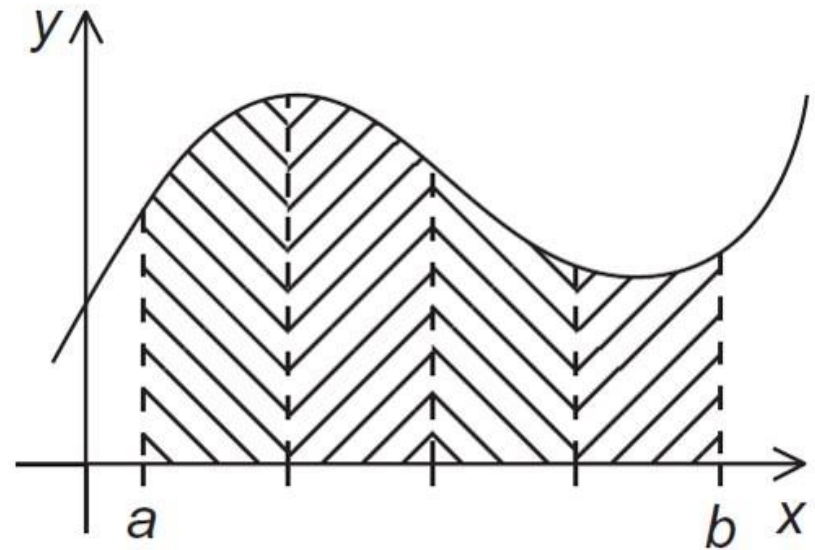
2. TRAPEZOIDAL RULE IN MPI



The Trapezoidal Rule



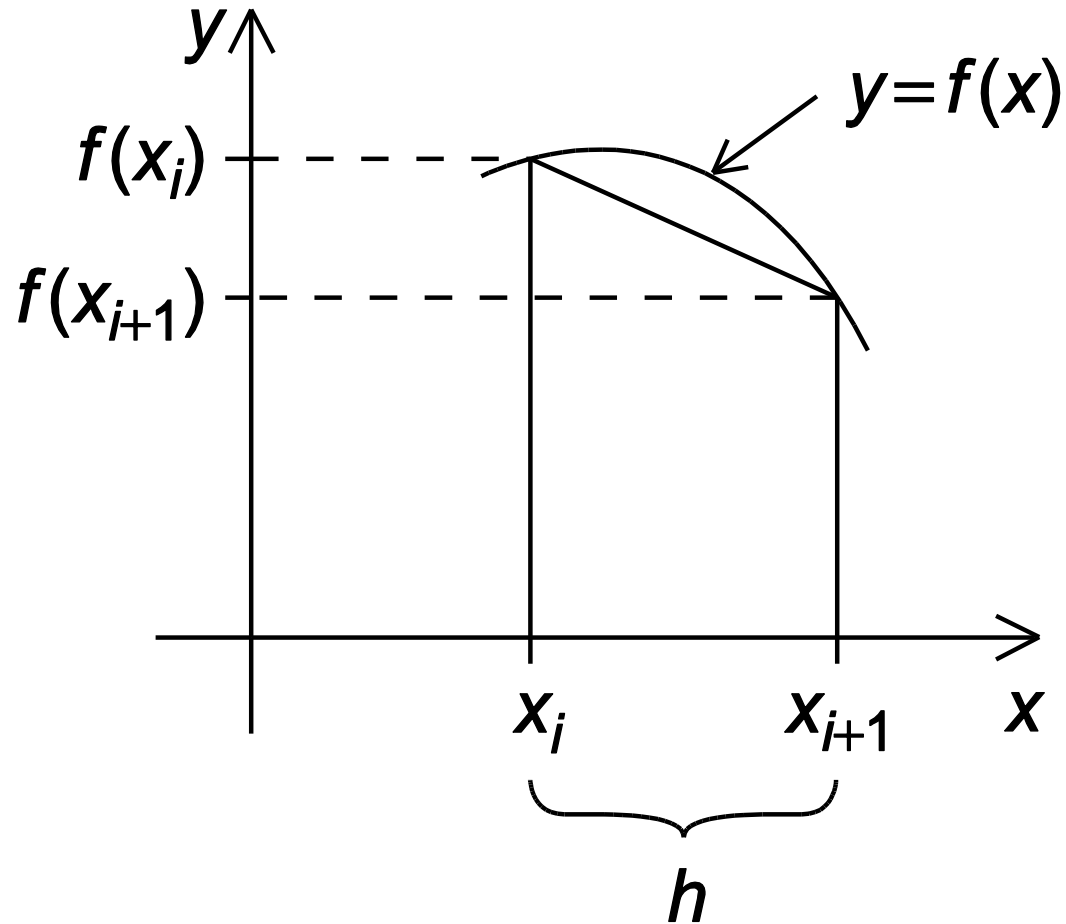
(a)



(b)



One trapezoid



The Trapezoidal Rule

$$\text{Area of one trapezoid} = \frac{h}{2} [f(x_i) + f(x_{i+1})]$$

$$h = \frac{b - a}{n}$$

$$x_0 = a, \quad x_1 = a + h, \quad x_2 = a + 2h, \dots, \quad x_{n-1} = a + (n-1)h, \quad x_n = b$$

$$\text{Sum of trapezoid areas} = h[f(x_0)/2 + f(x_1) + f(x_2) + \dots + f(x_{n-1}) + f(x_n)/2]$$



Pseudo-code for a serial program

```
/* Input: a, b, n */  
h = (b - a) / n;  
approx = (f(a) + f(b)) / 2.0;  
for (i = 1; i <= n - 1; i++) {  
    x_i = a + i * h;  
    approx += f(x_i);  
}  
approx = h * approx;
```



Parallelizing the Trapezoidal Rule

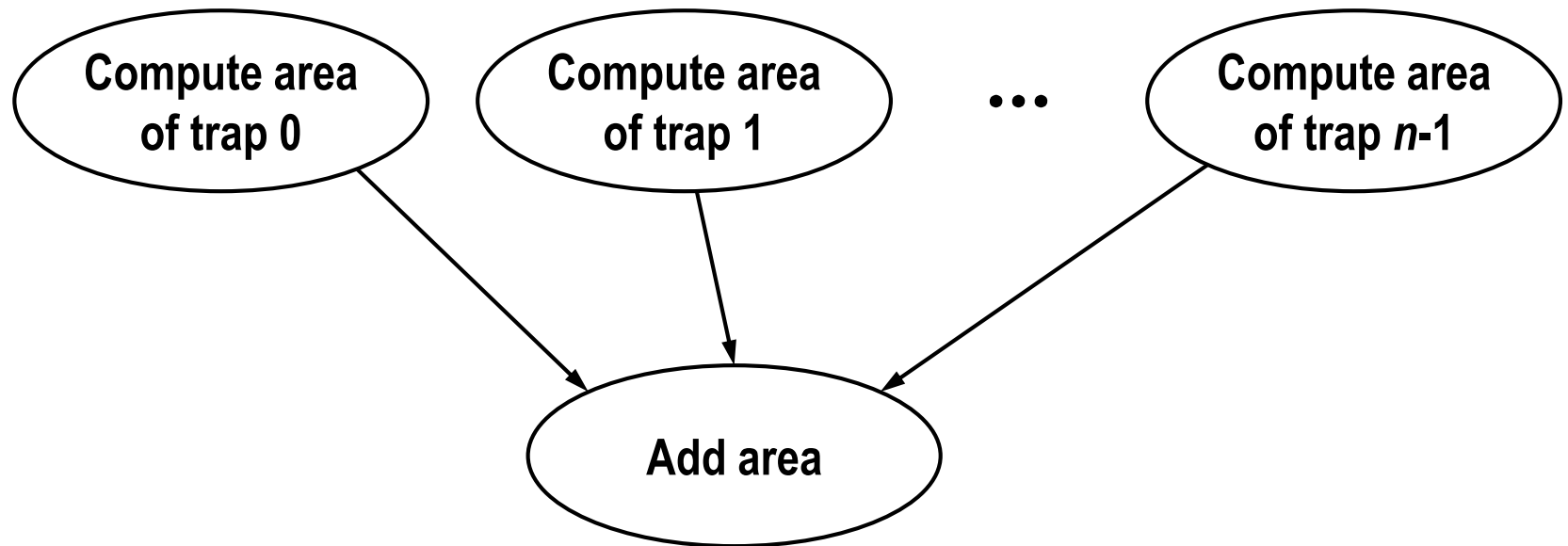
- 1. Partition problem solution into tasks.**
- 2. Identify communication channels between tasks.**
- 3. Aggregate tasks into composite tasks.**
- 4. Map composite tasks to cores.**



Parallel pseudo-code

```
1  Get a, b, n;
2  h = (b - a) / n;
3  local_n = n / comm_sz;
4  local_a = a + my_rank * local_n * h;
5  local_b = local_a + local_n * h;
6  local_integral = Trap(local_a, local_b, local_n, h);
7  if (my_rank != 0)
8      Send local_integral to process 0;
9  else /* my_rank == 0 */
10     total_integral = local_integral;
11     for (proc = 1; proc < comm_sz; proc++) {
12         Receive local_integral from proc;
13         total_integral += local_integral;
14     }
15 }
16 if (my_rank == 0)
17     print result;
```

Tasks and communications for Trapezoidal Rule



First version (1)

```
1  int main(void) {
2      int my_rank, comm_sz, n = 1024, local_n;
3      double a = 0.0, b = 3.0, h, local_a, local_b;
4      double local_int, total_int;
5      int source;
6
7      MPI_Init(NULL, NULL);
8      MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);
9      MPI_Comm_size(MPI_COMM_WORLD, &comm_sz);
10
11     h = (b - a)/n;          /* h is the same for all processes */
12     local_n = n/comm_sz; /* So is the number of trapezoids */
13
14     local_a = a + my_rank * local_n * h;
15     local_b = local_a + local_n * h;
16     local_int = Trap(local_a, local_b, local_n, h);
17
18     if (my_rank != 0) {
19         MPI_Send(&local_int, 1, MPI_DOUBLE, 0, 0,
20                 MPI_COMM_WORLD);
```



First version (2)

```
21 } else {
22     total_int = local_int;
23     for (source = 1; source < comm_sz; source++) {
24         MPI_Recv(&local_int, 1, MPI_DOUBLE, source, 0,
25                 MPI_COMM_WORLD, MPI_STATUS_IGNORE);
26         total_int += local_int;
27     }
28 }
29
30 if (my_rank == 0) {
31     printf("With n = %d trapezoids, our estimate\n", n);
32     printf("of the integral from %f to %f = %.15e\n",
33           a, b, total_int);
34 }
35 MPI_Finalize();
36 return 0;
37 } /* main */
```



First version (3)

```
1 double Trap(  
2     double left_endpt  /* in */,  
3     double right_endpt /* in */,  
4     int     trap_count /* in */,  
5     double base_len    /* in */) {  
6     double estimate, x;  
7     int i;  
8  
9     estimate = (f(left_endpt) + f(right_endpt))/2.0;  
10    for (i = 1; i <= trap_count - 1; i++) {  
11        x = left_endpt + i*base_len;  
12        estimate += f(x);  
13    }  
14    estimate = estimate*base_len;  
15  
16    return estimate;  
17 } /* Trap */
```



Dealing with I/O

```
#include <stdio.h>
#include <mpi.h>

int main(void) {
    int my_rank, comm_sz;

    MPI_Init(NULL, NULL);
    MPI_Comm_size(MPI_COMM_WORLD, &comm_sz);
    MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);

    printf("Proc %d of %d > Does anyone have a toothpick?\n",
           my_rank, comm_sz);

    MPI_Finalize();
    return 0;
} /* main */
```

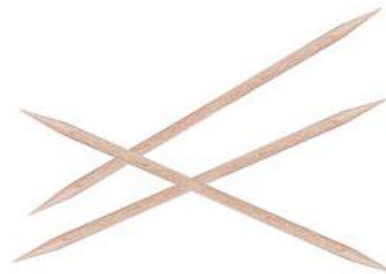
*Each process just
prints a message.*



Running with 6 processes

```
Proc 0 of 6 > Does anyone have a toothpick?  
Proc 1 of 6 > Does anyone have a toothpick?  
Proc 2 of 6 > Does anyone have a toothpick?  
Proc 4 of 6 > Does anyone have a toothpick?  
Proc 3 of 6 > Does anyone have a toothpick?  
Proc 5 of 6 > Does anyone have a toothpick?
```

unpredictable output



Input

- Most MPI implementations only allow process 0 in `MPI_COMM_WORLD` access to **stdin**.
- Process 0 must read the data (`scanf`) and send to the other processes.

```
...  
MPI_Comm_rank(MPI_COMM_WORLD, &my_rank);  
MPI_Comm_size(MPI_COMM_WORLD, &comm_sz);  
Get_input(my_rank, comm_sz, &a, &b, &n);  
  
H = (b - a) / n  
...
```

教材有误



Function for reading user input

```
void Get_input(  
    int      my_rank    /* in */,  
    int      comm_sz    /* in */,  
    double*  a_p        /* out */,  
    double*  b_p        /* out */,  
    int*     n_p        /* out */) {  
    int dest;  
  
    if (my_rank == 0) {  
        printf("Enter a, b, and n\n");  
        scanf("%lf %lf %d", a_p, b_p, n_p);  
        for (dest = 1; dest < comm_sz; dest++) {  
            MPI_Send(a_p, 1, MPI_DOUBLE, dest, 0, MPI_COMM_WORLD);  
            MPI_Send(b_p, 1, MPI_DOUBLE, dest, 0, MPI_COMM_WORLD);  
            MPI_Send(n_p, 1, MPI_INT, dest, 0, MPI_COMM_WORLD);  
        }  
    } else { /* my_rank != 0 */  
        MPI_Recv(a_p, 1, MPI_DOUBLE, 0, 0, MPI_COMM_WORLD,  
                 MPI_STATUS_IGNORE);  
        MPI_Recv(b_p, 1, MPI_DOUBLE, 0, 0, MPI_COMM_WORLD,  
                 MPI_STATUS_IGNORE);  
        MPI_Recv(n_p, 1, MPI_INT, 0, 0, MPI_COMM_WORLD,  
                 MPI_STATUS_IGNORE);  
    }  
} /* Get_input */
```

3. COLLECTIVE COMMUNICATION



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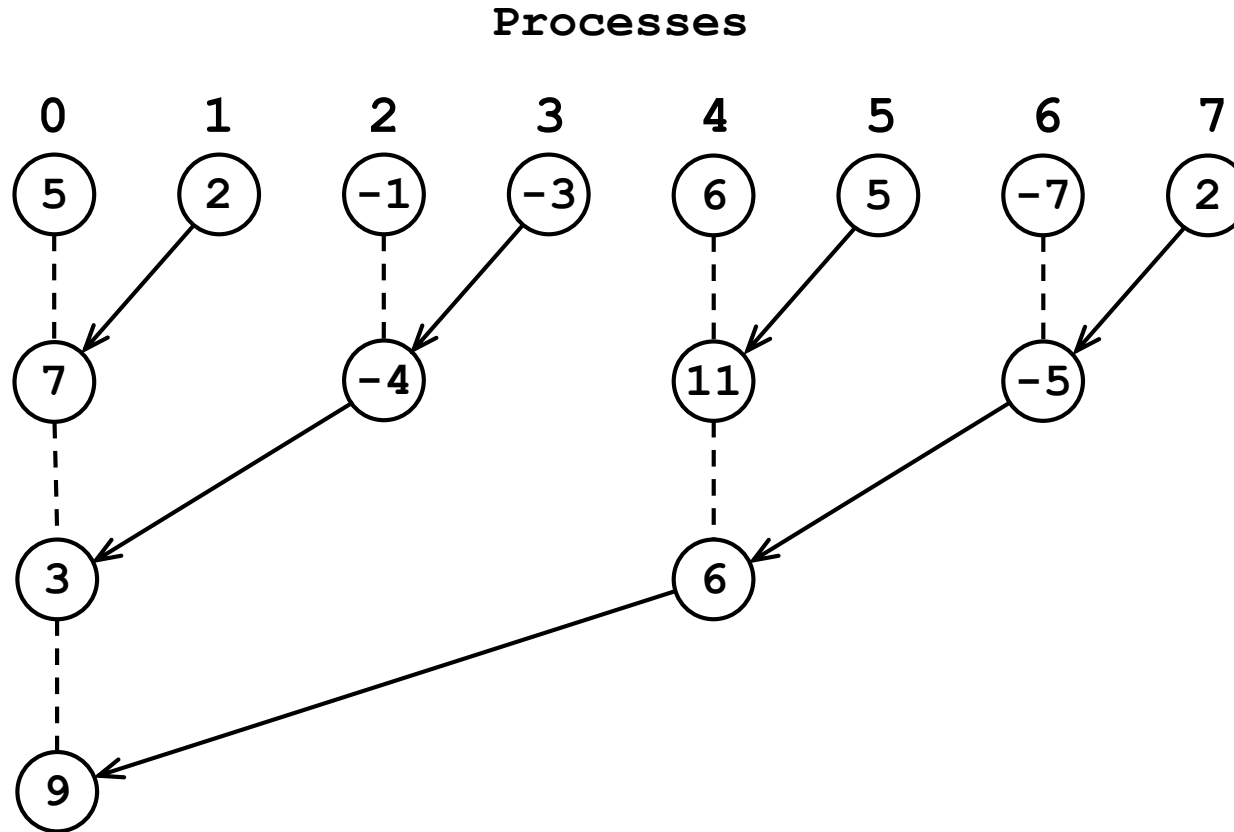
Optimize the trapezoidal rule program

The “global sum” after each process has computed its part of the integral.

```
if (my_rank != 0) {
    MPI_Send(&local_int, 1, MPI_DOUBLE, 0, 0,
             MPI_COMM_WORLD);
} else {
    total_int = local_int;
    for (source = 1; source < comm_sz; source++) {
        MPI_Recv(&local_int, 1, MPI_DOUBLE, source, 0,
                 MPI_COMM_WORLD, MPI_STATUS_IGNORE);
        total_int += local_int;
    }
}
```



A tree-structured global sum

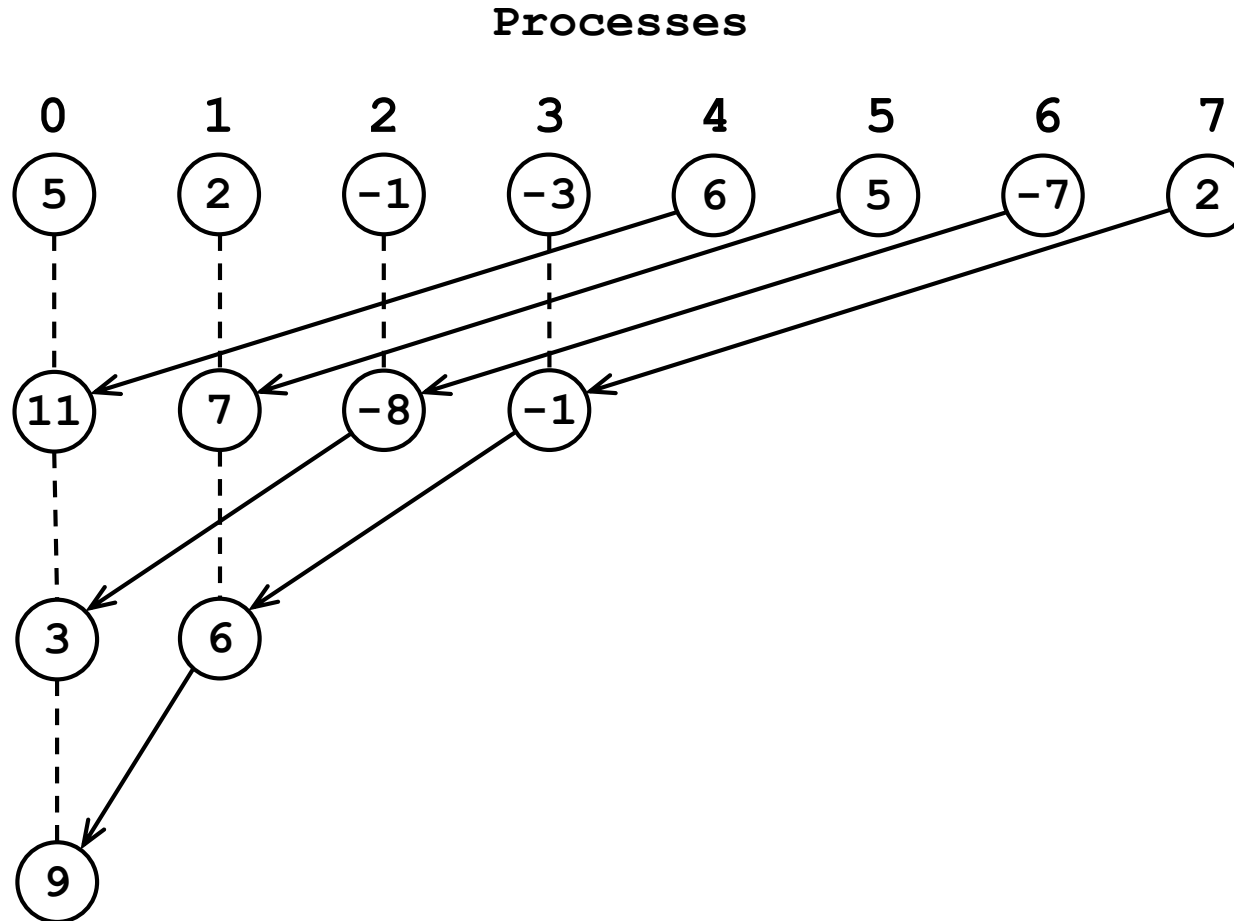


Tree-structured communication

- 1. (a) Process 1 sends to 0, 3 sends to 2, 5 sends to 4, and 7 sends to 6.
(b) Processes 0, 2, 4, and 6 add in the received values.**
- 2. (a) Processes 2 and 6 send their new values to processes 0 and 4, respectively.
(b) Processes 0 and 4 add the received values into their new values.**
- 3. (a) Process 4 sends its newest value to process 0.
(b) Process 0 adds the received value to its newest value.**



An alternative tree-structured global sum



(1) MPI_Reduce

```
int MPI_Reduce(  
    void*      input_data_p    /* in */,  
    void*      output_data_p   /* out */,  
    int        count           /* in */,  
    MPI_Datatype datatype      /* in */,  
    MPI_Op      operator       /* in */,  
    int        dest_process    /* in */,  
    MPI_Comm    comm          /* in */);
```

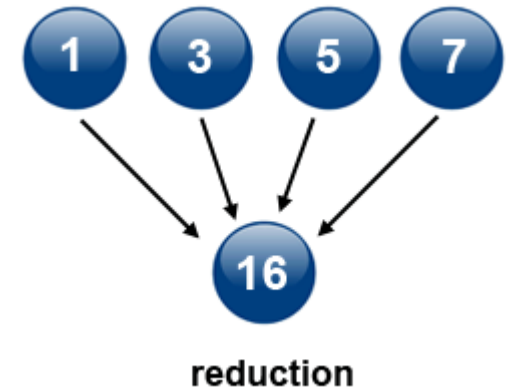
```
MPI_Reduce(&local_int, &total_int, 1, MPI_DOUBLE, MPI_SUM,  
          0, MPI_COMM_WORLD);
```

```
double local_x[N], sum[N];  
...  
MPI_Reduce(local_x, sum, N, MPI_DOUBLE, MPI_SUM, 0,  
          MPI_COMM_WORLD);
```



Predefined reduction operators in MPI

Operation Value	Meaning
MPI_MAX	Maximum
MPI_MIN	Minimum
MPI_SUM	Sum
MPI_PROD	Product
MPI_LAND	Logical and
MPI_BAND	Bitwise and
MPI_LOR	Logical or
MPI_BOR	Bitwise or
MPI_LXOR	Logical exclusive or
MPI_BXOR	Bitwise exclusive or
MPI_MAXLOC	Maximum and location of maximum
MPI_MINLOC	Minimum and location of minimum



Collective vs. Point-to-Point Communications

- All the processes in the communicator must call the same collective function.
- For example, a program that attempts to match a call to **MPI_Reduce** on one process with a call to **MPI_Recv** on another process is erroneous, and, in all likelihood, the program will hang or crash.



Collective vs. Point-to-Point Communications

- The arguments passed by each process to an MPI collective communication must be “compatible.”
- For example, if one process passes in 0 as the **dest_process** and another passes in 1, then the outcome of a call to **MPI_Reduce** is erroneous, and, once again, the program is likely to hang or crash.



Collective vs. Point-to-Point Communications

- The **output_data_p** argument is only used on **dest_process**.
- However, all of the processes still need to pass in an actual argument corresponding to **output_data_p**, even if it's just **NULL**.



Collective vs. Point-to-Point Communications

- **Point-to-point communications are matched on the basis of tags and communicators.**
- **Collective communications don't use tags. They're matched solely on the basis of the communicator and the order in which they're called.**



Example (1)

Time	Process 0	Process 1	Process 2
0	a = 1; c = 2	a = 1; c = 2	a = 1; c = 2
1	MPI_Reduce (&a, &b, ...)	MPI_Reduce (&c, &d, ...)	MPI_Reduce (&a, &b, ...)
2	MPI_Reduce (&c, &d, ...)	MPI_Reduce (&a, &b, ...)	MPI_Reduce (&c, &d, ...)

Multiple calls to MPI_Reduce



Example (2)

- Suppose that each process calls **MPI_Reduce** with operator **MPI_SUM**, and destination process 0.
- At first glance, it might seem that after the two calls to **MPI_Reduce**, the value of b will be 3, and the value of d will be 6.



Example (3)

- However, the names of the memory locations are irrelevant to the matching of the calls to **MPI_Reduce**.
- The order of the calls will determine the matching so the value stored in b will be $1+2+1 = 4$, and the value stored in d will be $2+1+2 = 5$.



Another question

```
MPI_Reduce(&x, &x, 1, MPI_DOUBLE, MPI_SUM, 0, comm);
```

- How about using the same buffer as input and output?

—The result is unpredictable.

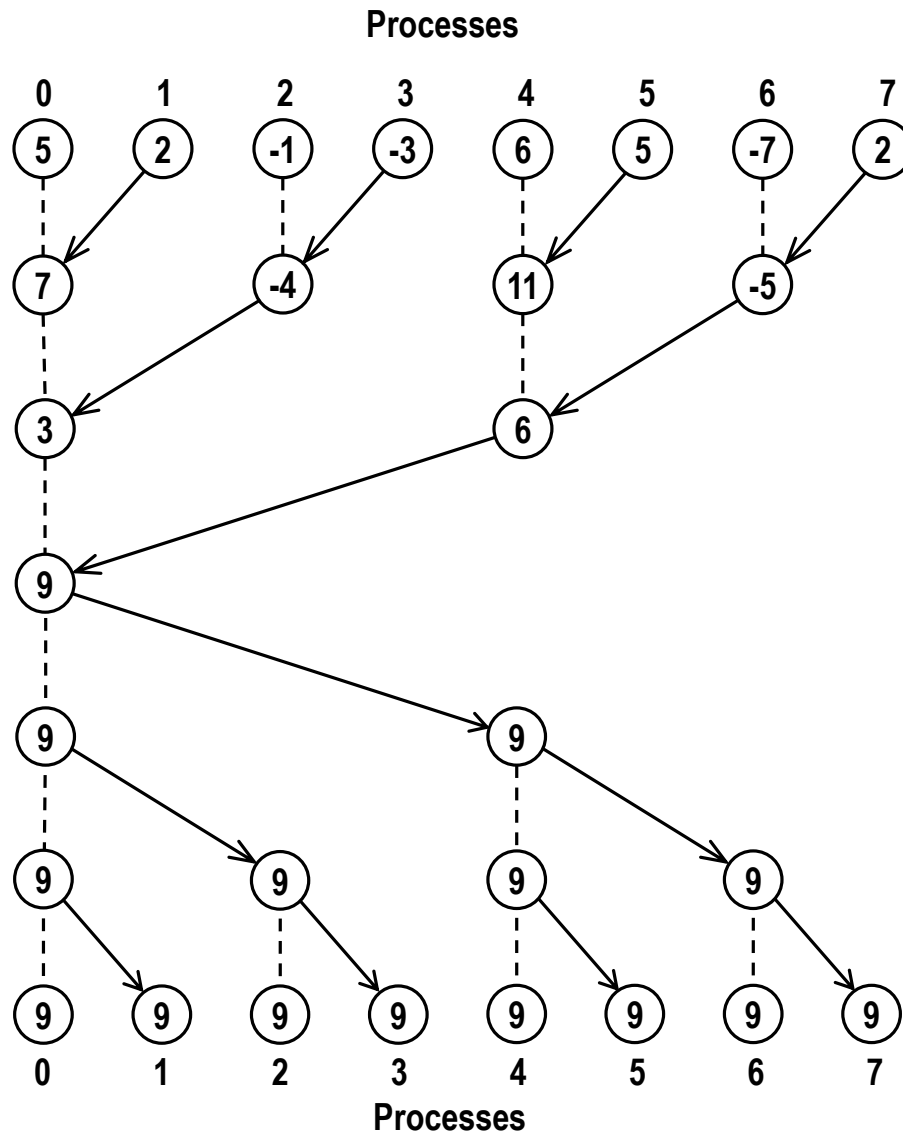


(2) MPI_Allreduce

- Useful in a situation in which all of the processes need the result of a global sum in order to complete some larger computation.

```
int MPI_Allreduce(  
    void*      input_data_p    /* in */,  
    void*      output_data_p   /* out */,  
    int        count           /* in */,  
    MPI_Datatype datatype      /* in */,  
    MPI_Op      operator        /* in */,  
    MPI_Comm    comm           /* in */);
```



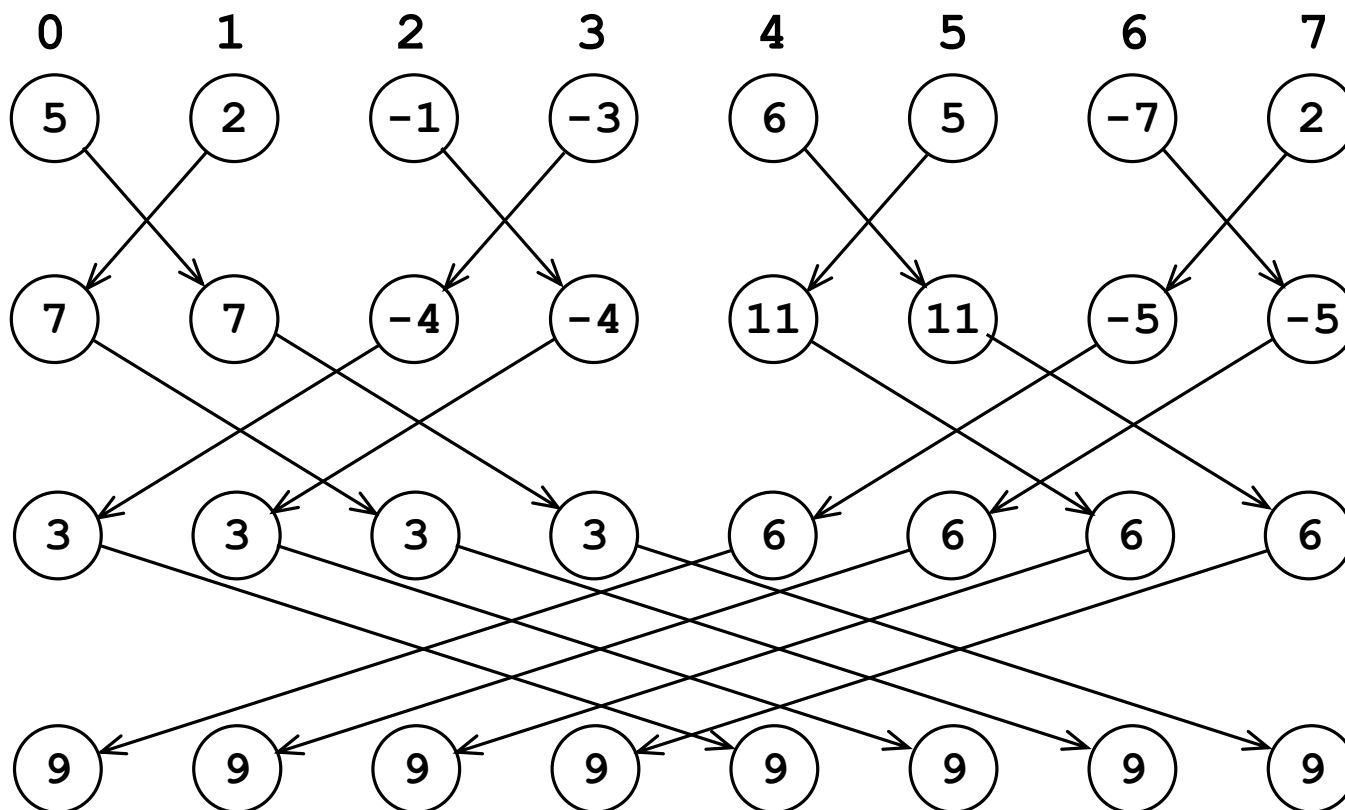


A global sum followed by distribution of the result.





Processes



A butterfly-structured global sum.



(3) Broadcast

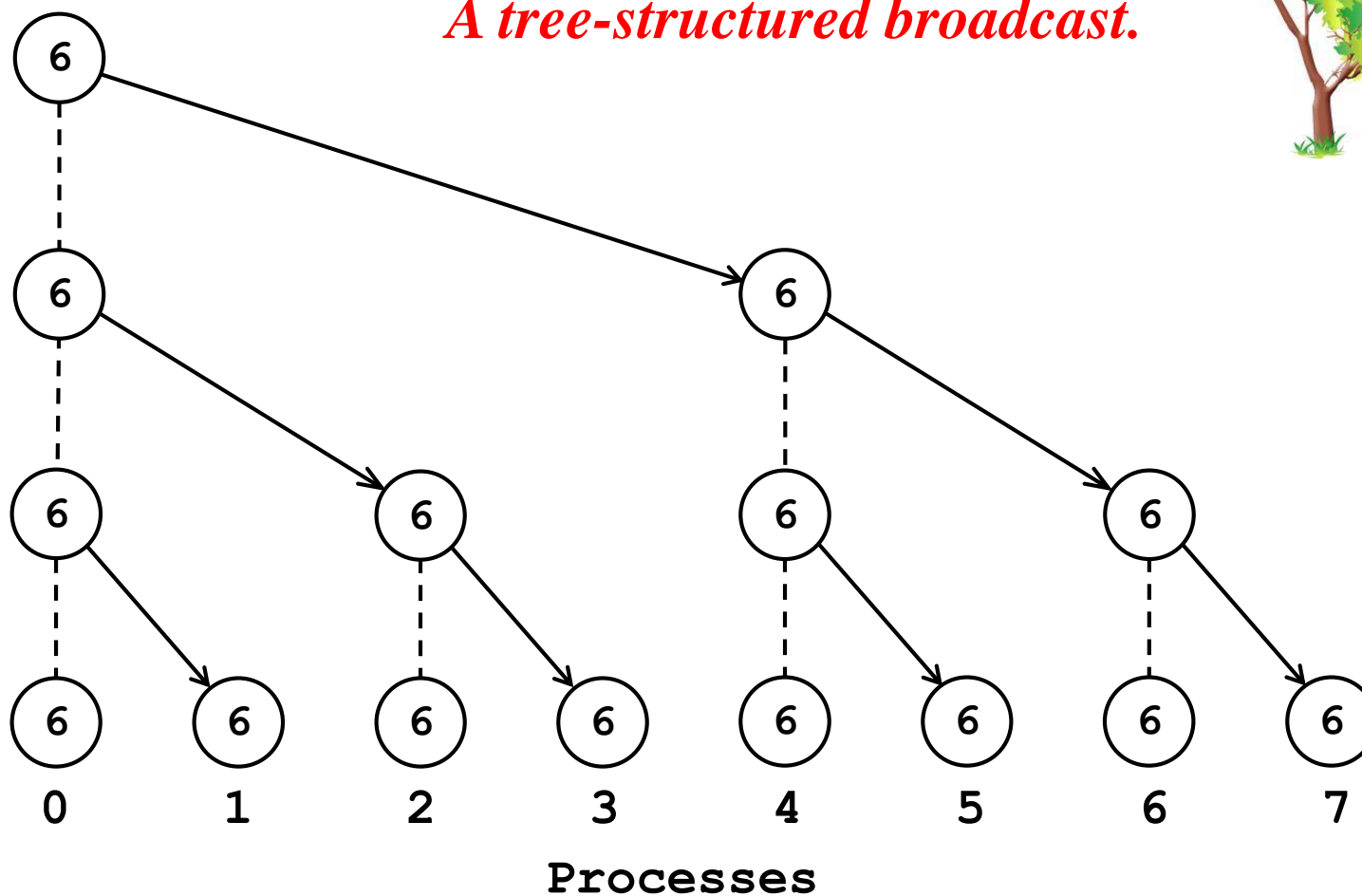
- **Data belonging to a single process is sent to all of the processes in the communicator.**

```
int MPI_Bcast (
    void*          data_p          /* in/out */,
    int            count           /* in      */,
    MPI_Datatype    datatype       /* in      */,
    int            source_proc     /* in      */,
    MPI_Comm        comm           /* in      */);
```





A tree-structured broadcast.



A version of Get_input that uses MPI_Bcast

```
void Get_input(  
    int      my_rank    /* in */,  
    int      comm_sz    /* in */,  
    double*  a_p        /* out */,  
    double*  b_p        /* out */,  
    int*     n_p        /* out */) {  
  
    if (my_rank == 0) {  
        printf("Enter a, b, and n\n");  
        scanf("%lf %lf %d", a_p, b_p, n_p);  
    }  
  
    MPI_Bcast(a_p, 1, MPI_DOUBLE, 0, MPI_COMM_WORLD);  
    MPI_Bcast(b_p, 1, MPI_DOUBLE, 0, MPI_COMM_WORLD);  
    MPI_Bcast(n_p, 1, MPI_INT, 0, MPI_COMM_WORLD);  
} /* Get_input */
```



Data distributions

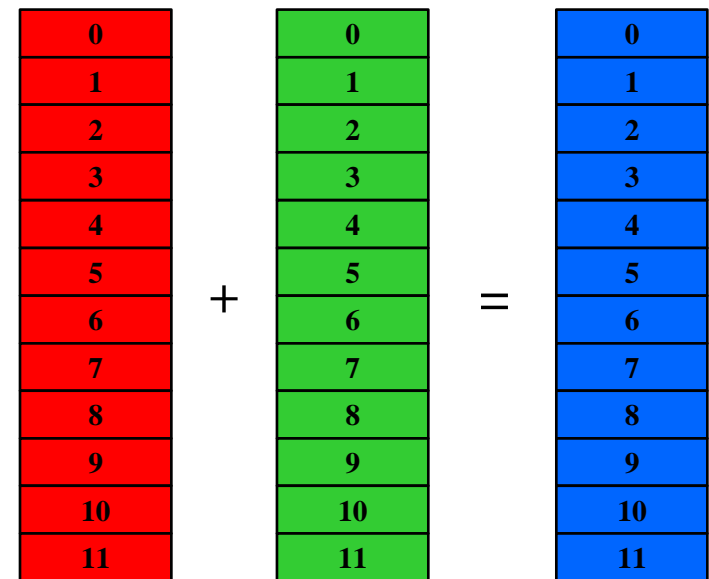
$$\begin{aligned}x + y &= (x_0, x_1, \dots, x_{n-1}) + (y_0, y_1, \dots, y_{n-1}) \\&= (x_0 + y_0, x_1 + y_1, \dots, x_{n-1} + y_{n-1}) \\&= (z_0, z_1, \dots, z_{n-1}) \\&= z\end{aligned}$$

Compute a vector sum.



Serial implementation of vector addition

```
void Vector_sum(double x[], double y[], double z[], int n) {  
    int i;  
  
    for (i = 0; i < n; i++)  
        z[i] = x[i] + y[i];  
} /* Vector_sum */
```



Different partitions of a 12-component vector among 3 processes

Process	Components											
	Block				Cyclic				Block-cyclic Blocksize = 2			
0	0	1	2	3	0	3	6	9	0	1	6	7
1	4	5	6	7	1	4	7	10	2	3	8	9
2	8	9	10	11	2	5	8	11	4	5	10	11



Partitioning options

- **Block partitioning**
 - Assign blocks of consecutive components to each process.
- **Cyclic partitioning**
 - Assign components in a round robin fashion.
- **Block-cyclic partitioning**
 - Use a cyclic distribution of blocks of components.



Parallel implementation of vector addition

```
void Parallel_vector_sum(  
    double  local_x[]    /* in  */,  
    double  local_y[]    /* in  */,  
    double  local_z[]    /* out */,  
    int      local_n     /* in  */) {  
    int local_i;  
  
    for (local_i = 0; local_i < local_n; local_i++)  
        local_z[local_i] = local_x[local_i] + local_y[local_i];  
} /* Parallel_vector_sum */
```



(4) Scatter

- **MPI_Scatter** can be used in a function that reads in an entire vector on process 0 but only sends the needed components to each of the other processes.

```
int MPI_Scatter(  
    void*          send_buf_p      /* in */,  
    int            send_count      /* in */,  
    MPI_Datatype   send_type       /* in */,  
    void*          recv_buf_p      /* out */,  
    int            recv_count      /* in */,  
    MPI_Datatype   recv_type       /* in */,  
    int            src_proc        /* in */,  
    MPI_Comm       comm           /* in */);
```



Reading and distributing a vector

```
void Read_vector(  
    double    local_a[]    /* out */,  
    int       local_n      /* in  */,  
    int       n            /* in  */,  
    char      vec_name[]   /* in  */,  
    int       my_rank      /* in  */,  
    MPI_Comm  comm         /* in  */) {  
    double* a = NULL;  
    int i;  
  
    if (my_rank == 0) {  
        a = malloc(n*sizeof(double));  
        printf("Enter the vector %s\n", vec_name);  
        for (i = 0; i < n; i++)  
            scanf("%lf", &a[i]);  
        MPI_Scatter(a, local_n, MPI_DOUBLE, local_a, local_n, MPI_DOUBLE,  
                    0, comm);  
        free(a);  
    } else {  
        MPI_Scatter(a, local_n, MPI_DOUBLE, local_a, local_n, MPI_DOUBLE,  
                    0, comm);  
    }  
} /* Read_vector */
```



(5) Gather

- Collect all of the components of the vector onto process 0, and then process 0 can process all of the components.

```
int MPI_Gather(  
    void*          send_buf_p    /* in */,  
    int            send_count    /* in */,  
    MPI_Datatype   send_type     /* in */,  
    void*          recv_buf_p    /* out */,  
    int            recv_count    /* in */,  
    MPI_Datatype   recv_type     /* in */,  
    int            dest_proc     /* in */,  
    MPI_Comm       comm         /* in */);
```



Print a distributed vector (1)

```
int Print_vector(  
    double        local_b[]      /* in */,  
    int           local_n        /* in */,  
    int           n              /* in */,  
    char          title[]        /* in */,  
    int           my_rank        /* in */,  
    MPI_Comm      comm          /* in */) {  
  
    double* b = NULL;  
    int i;
```



Print a distributed vector (2)

```
if (my_rank == 0) {
    b = malloc(n*sizeof(double));
    MPI_Gather(local_b, local_n, MPI_DOUBLE, b, local_n,
               MPI_DOUBLE, 0, comm);
    printf("%s\n", title);
    for (i = 0; i < n; i++)
        printf("%f ", b[i]);
    printf("\n");
    free(b);
} else {
    MPI_Gather(local_b, local_n, MPI_DOUBLE, b, local_n,
               MPI_DOUBLE, 0, comm);
}
} /* Print_vector */
```



(6) Allgather

- Concatenates the contents of each process' **send_buf_p** and stores this in each process' **recv_buf_p**.
- As usual, **recv_count** is the amount of data being received from each process.

```
int MPI_Allgather(  
    void*      send_buf_p      /* in */,  
    int        send_count      /* in */,  
    MPI_Datatype send_type      /* in */,  
    void*      recv_buf_p      /* out */,  
    int        recv_count      /* in */,  
    MPI_Datatype recv_type      /* in */,  
    MPI_Comm    comm           /* in */);
```



Matrix-vector multiplication

$A = (a_{ij})$ is an $m \times n$ matrix

x is a vector with n components

$y = Ax$ is a vector with m components

$$y_i = a_{i,0}x_0 + a_{i,1}x_1 + a_{i,2}x_2 + \cdots + a_{i,n-1}x_{n-1}$$

i -th component of y

*Dot product of the i th
row of A with x .*



Matrix-vector multiplication

$a_{0,0}$	$a_{0,1}$	\cdots	$a_{0,n-1}$
$a_{1,0}$	$a_{1,1}$	\cdots	$a_{1,n-1}$
\vdots	\vdots		\vdots
$a_{i,0}$	$a_{i,1}$	\cdots	$a_{i,n-1}$
\vdots	\vdots		\vdots
$a_{m-1,0}$	$a_{m-1,1}$	\cdots	$a_{m-1,n-1}$

 \bullet

x_0
x_1
\vdots
x_{n-1}

 $=$

y_0
y_1
\vdots
$y_i = a_{i,0}x_0 + a_{i,1}x_1 + \cdots + a_{i,n-1}x_{n-1}$
\vdots
y_{m-1}



Multiply a matrix by a vector

```
/* For each row of A */  
for (i = 0; i < m; i++) {  
    /* Form dot product of ith row with x */  
    y[i] = 0.0;  
    for (j = 0; j < n; j++)  
        y[i] += A[i][j] * x[j];  
}
```

Serial pseudo-code



C style arrays

$$\begin{pmatrix} 0 & 1 & 2 & 3 \\ 4 & 5 & 6 & 7 \\ 8 & 9 & 10 & 11 \end{pmatrix}$$

stored as

0 1 2 3 4 5 6 7 8 9 10 11



Serial matrix-vector multiplication

```
void Mat_vec_mult(  
    double  A[]    /* in  */,  
    double  x[]    /* in  */,  
    double  y[]    /* out */,  
    int      m      /* in  */,  
    int      n      /* in  */) {  
    int i, j;  
  
    for (i = 0; i < m; i++) {  
        y[i] = 0.0;  
        for (j = 0; j < n; j++)  
            y[i] += A[i*n+j] * x[j];  
    }  
} /* Mat_vec_mult */
```



An MPI matrix-vector multiplication function (1)

```
void Mat_vec_mult(  
    double    local_A[]    /* in */,  
    double    local_x[]    /* in */,  
    double    local_y[]    /* out */,  
    int        local_m      /* in */,  
    int        n            /* in */,  
    int        local_n      /* in */,  
    MPI_Comm   comm        /* in */) {  
    double* x;  
    int local_i, j;  
    int local_ok = 1;
```



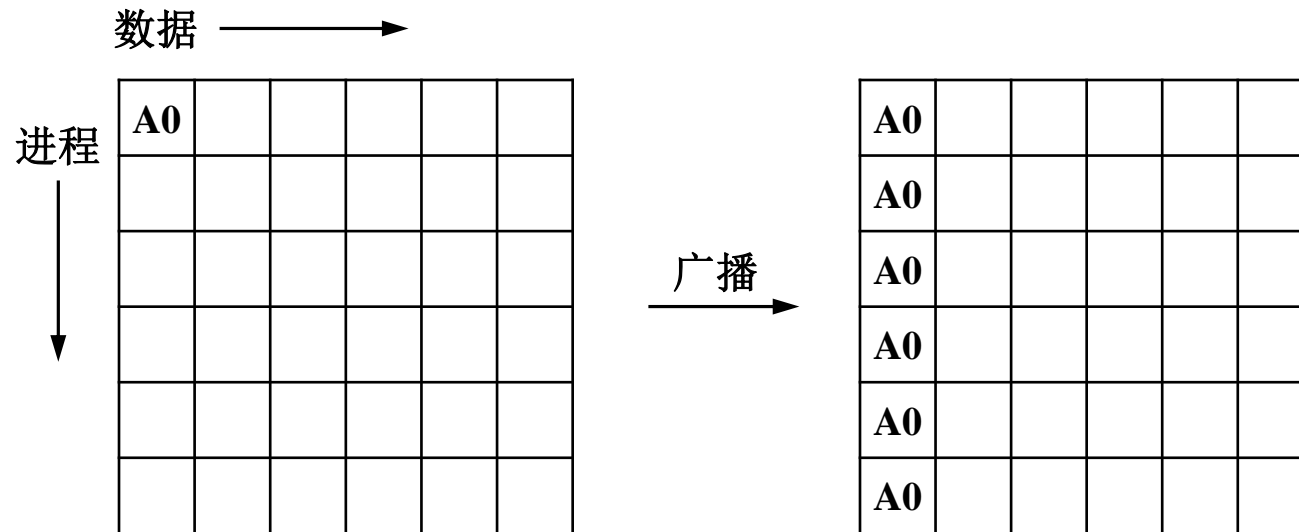
An MPI matrix-vector multiplication function (2)

```
x = malloc(n*sizeof(double));
MPI_Allgather(local_x, local_n, MPI_DOUBLE,
              x, local_n, MPI_DOUBLE, comm);

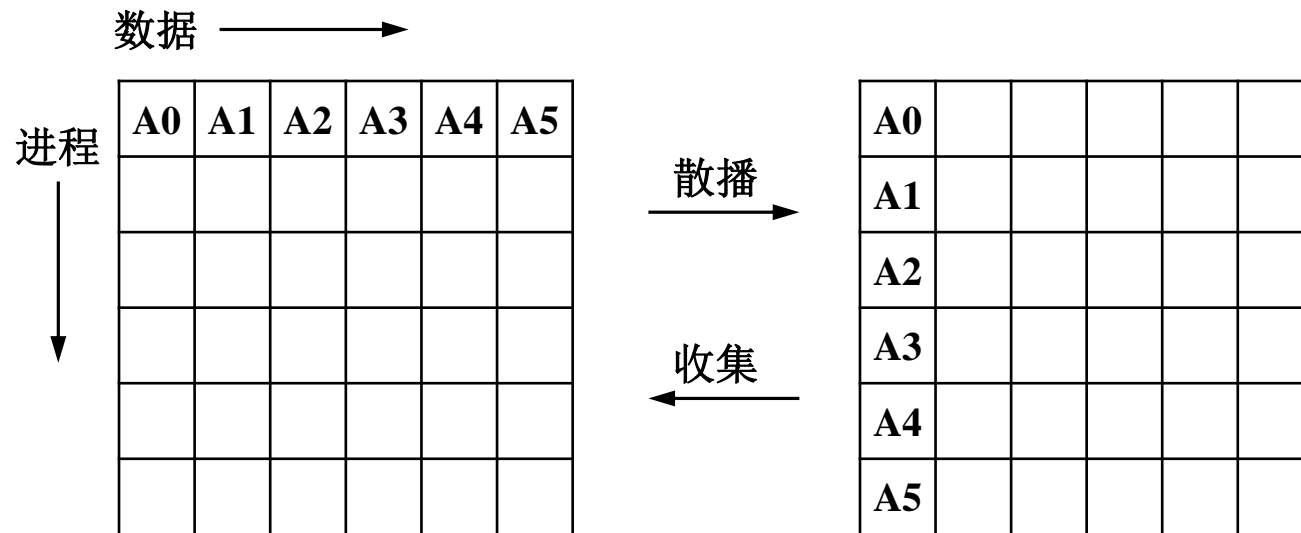
for (local_i = 0; local_i < local_m; local_i++) {
    local_y[local_i] = 0.0;
    for (j = 0; j < n; j++)
        local_y[local_i] += local_A[local_i*n+j] * x[j];
}
free(x);
} /* Mat_vec_mult */
```



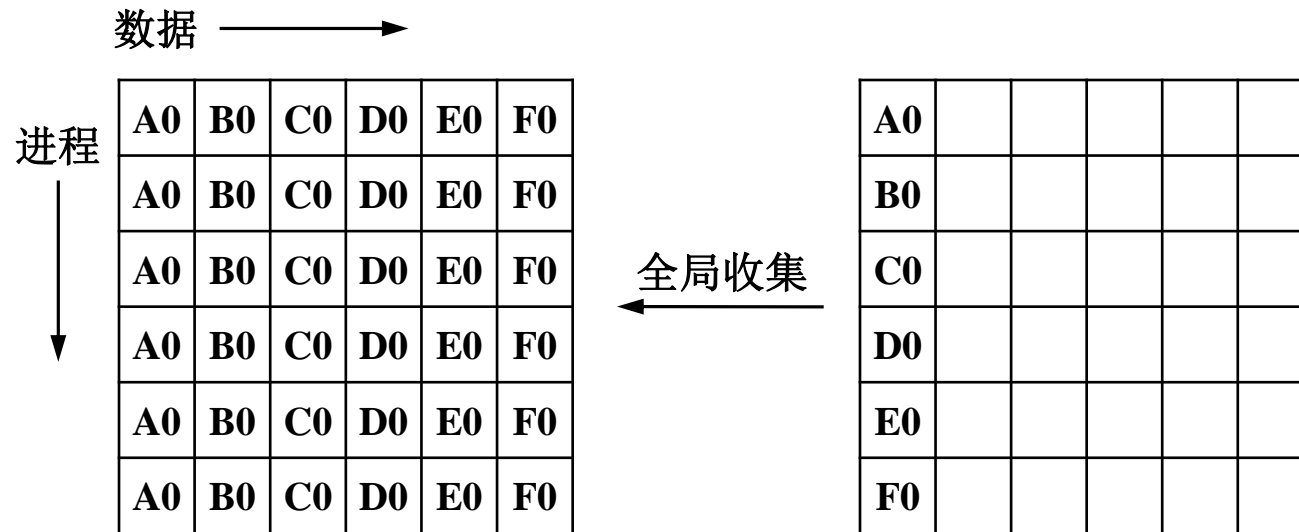
广播 (MPI_Bcast)



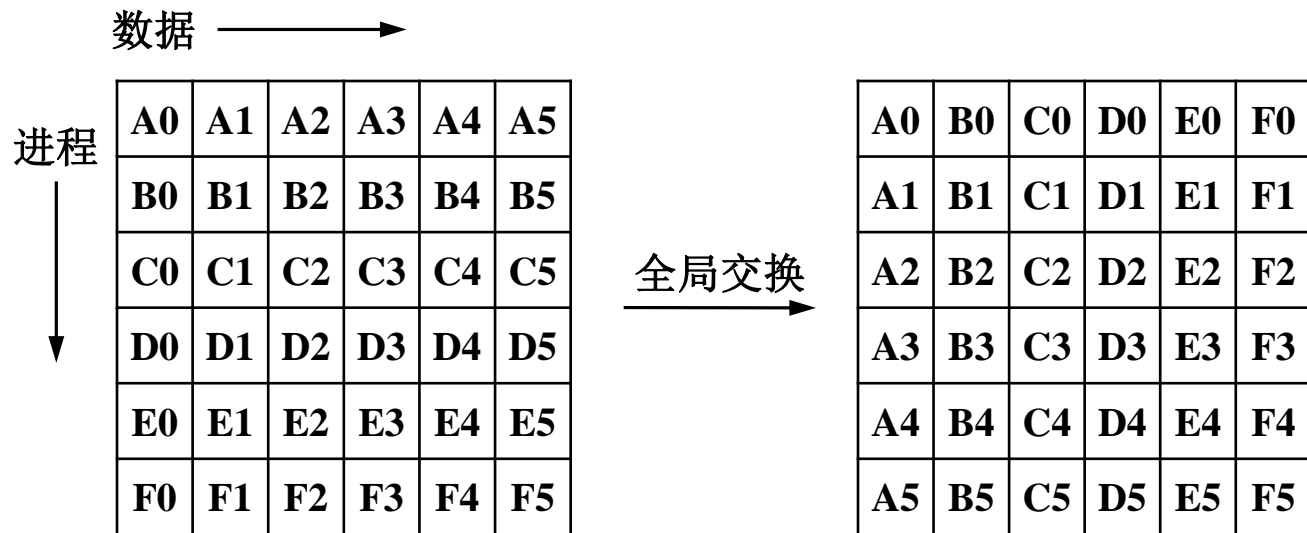
收集和散播 (MPI_Gather & MPI_Scatter)



全局收集 (MPI_Allgather)



全局交换 (MPI_Alltoall)





4. MPI DERIVED DATATYPES



Derived datatypes



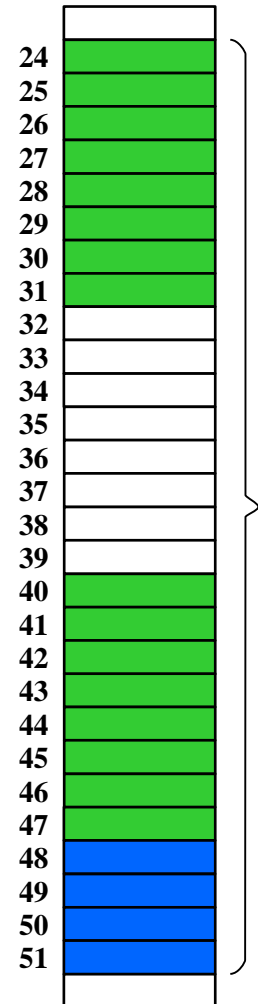
- Used to represent any collection of data items in memory by storing both the types of the items and their relative locations in memory.
- The idea is that if a function that sends data knows this information about a collection of data items, it can collect the items from memory before they are sent.
- Similarly, a function that receives data can distribute the items into their correct destinations in memory when they're received.



Derived datatypes

- Formally, consists of a sequence of basic MPI data types together with a displacement for each of the data types.
- Trapezoidal Rule example:

Variable	Address
a	24
b	40
n	48



$\{(\text{MPI_DOUBLE}, 0), (\text{MPI_DOUBLE}, 16), (\text{MPI_INT}, 24)\}$



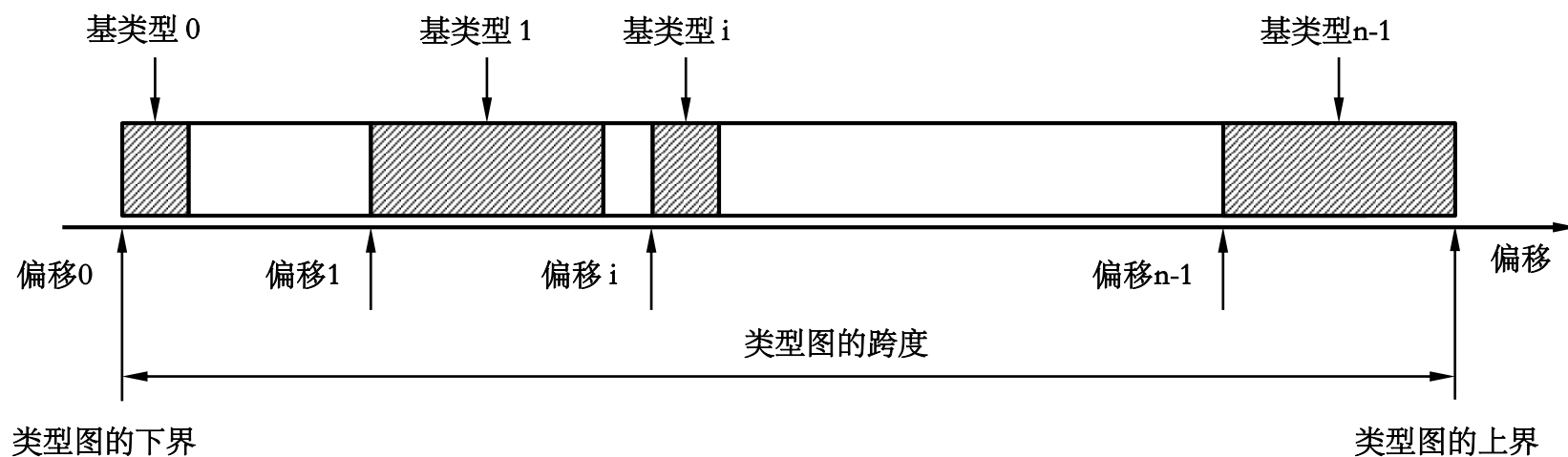
MPI_Type_create_struct

- Builds a derived datatype that consists of individual elements that have different basic types.

```
int MPI_Type_create_struct (
    int          count          /* in */,
    int          array_of_blocklengths[] /* in */,
    MPI_Aint      array_of_displacements[] /* in */,
    MPI_Datatype  array_of_types[] /* in */,
    MPI_Datatype* new_type_p     /* out */);
```



派生数据类型的类型图示意结构



MPI_Get_address

- Returns the address of the memory location referenced by **location_p**.
- The special type **MPI_Aint** is an integer type that is big enough to store an address on the system.

```
int MPI_Get_address(  
    void*          location_p    /* in */,  
    MPI_Aint*      address_p     /* out */);
```



MPI_Type_commit

- **Allows the MPI implementation to optimize its internal representation of the datatype for use in communication functions.**

```
int MPI_Type_commit(MPI_Datatype* new_mpi_t_p /* in/out */);
```



MPI_Type_free

- When we're finished with our new type, this frees any additional storage used.

```
int MPI_Type_free(MPI_Datatype* old_mpi_t_p /* in/out */);
```



Get_input function with a derived datatype (1)

```
void Build_mpi_type(  
    double*      a_p      /* in */,  
    double*      b_p      /* in */,  
    int*         n_p      /* in */,  
    MPI_Datatype* input_mpi_t_p /* out */) {  
  
    int array_of_blocklengths[3] = {1, 1, 1};  
    MPI_Datatype array_of_types[3] = {MPI_DOUBLE,  
        MPI_DOUBLE, MPI_INT};  
    MPI_Aint a_addr, b_addr, n_addr;  
    MPI_Aint array_of_displacements[3] = {0};
```



Get_input function with a derived datatype (2)

```
MPI_Get_address(a_p, &a_addr);
MPI_Get_address(b_p, &b_addr);
MPI_Get_address(n_p, &n_addr);
array_of_displacements[1] = b_addr - a_addr;
array_of_displacements[2] = n_addr - a_addr;
MPI_Type_create_struct(3, array_of_blocklengths,
                      array_of_displacements, array_of_types,
                      input_mpi_t_p);
MPI_Type_commit(input_mpi_t_p);    /*派生出的新数据类型，必须
                                   先要经过MPI系统的确认后才能使用。*/
} /* Build_mpi_type */
```



Get_input function with a derived datatype (3)

```
void Get_input(int my_rank, int comm_sz, double* a_p,  
              double* b_p, int* n_p) {  
    MPI_Datatype input_mpi_t;  
  
    Build_mpi_type(a_p, b_p, n_p, &input_mpi_t);  
  
    if (my_rank == 0) {  
        printf("Enter a, b, and n\n");  
        scanf("%lf %lf %d", a_p, b_p, n_p);  
    }  
  
    MPI_Bcast(a_p, 1, input_mpi_t, 0, MPI_COMM_WORLD);  
  
    MPI_Type_free(&input_mpi_t);  
} /* Get_input */
```





5. PERFORMANCE EVALUATION




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Elapsed parallel time

- Returns the number of seconds that have elapsed since some time in the past.

```
double MPI_Wtime(void);  
  
double start, finish;  
. . .  
start = MPI_Wtime();  
/* Code to be timed */  
. . .  
finish = MPI_Wtime();  
printf("Proc %d > Elapsed time = %e seconds\n"  
       my_rank, finish-start);
```



Elapsed serial time

- In this case, you don't need to link in the MPI libraries.
- Returns time in microseconds elapsed from some point in the past.

```
#include "timer.h"  
.  
.  
double now;  
.  
.  
GET_TIME(now);
```



Elapsed serial time

```
#include "timer.h"
. . .
double start, finish;
. . .
GET_TIME(start);
/* Code to be timed */
. . .
GET_TIME(finish);
printf("Elapsed time = %e seconds\n", finish - start);
```



MPI_Barrier

- Ensures that no process will return from calling it until every process in the communicator has started calling it.

```
int MPI_Barrier(MPI_Comm comm /* in */);
```



MPI_Barrier

```
double local_start, local_finish, local_elapsed, elapsed;
. . .
MPI_Barrier(comm);
local_start = MPI_Wtime();
/* Code to be timed */
. . .

local_finish = MPI_Wtime();
local_elapsed = local_finish - local_start;
MPI_Reduce(&local_elapsed, &elapsed, 1, MPI_DOUBLE,
          MPI_MAX, 0, comm);

if (my_rank == 0)
    printf("Elapsed time = %e seconds\n", elapsed);
```



Run-times of serial and parallel matrix-vector multiplication

comm_sz	Order of Matrix				
	1024	2048	4096	8192	16,384
1	4.1	16.0	64.0	270	1100
2	2.3	8.5	33.0	140	560
4	2.0	5.1	18.0	70	280
8	1.7	3.3	9.8	36	140
16	1.7	2.6	5.9	19	71

(milliseconds)



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Speedup

$$S(n, p) = \frac{T_{\text{serial}}(n)}{T_{\text{parallel}}(n, p)}$$



Efficiency

$$E(n, p) = \frac{S(n, p)}{p} = \frac{T_{\text{serial}}(n)}{p \times T_{\text{parallel}}(n, p)}$$



Speedups of Parallel Matrix-Vector Multiplication

comm_sz	Order of Matrix				
	1024	2048	4096	8192	16,384
1	1.0	1.0	1.0	1.0	1.0
2	1.8	1.9	1.9	1.9	2.0
4	2.1	3.1	3.6	3.9	3.9
8	2.4	4.8	6.5	7.5	7.9
16	2.4	6.2	10.8	14.2	15.5



Efficiencies of Parallel Matrix-Vector Multiplication

comm_sz	Order of Matrix				
	1024	2048	4096	8192	16,384
1	1.00	1.00	1.00	1.00	1.00
2	0.89	0.94	0.97	0.96	0.98
4	0.51	0.78	0.89	0.96	0.98
8	0.30	0.61	0.82	0.94	0.98
16	0.15	0.39	0.68	0.89	0.97



Scalability

- A program is **scalable** if the problem size can be increased at a rate so that the efficiency doesn't decrease as the number of processes increase.



Scalability

- Programs that can maintain a constant efficiency without increasing the problem size are sometimes said to be **strongly scalable**.
- Programs that can maintain a constant efficiency if the problem size increases at the same rate as the number of processes are sometimes said to be **weakly scalable**.



6. A PARALLEL SORTING ALGORITHM



Sorting

- **n keys and $p = \text{comm_sz}$ processes.**
- **n/p keys assigned to each process.**
- **No restrictions on which keys are assigned to which processes.**
- **When the algorithm terminates:**
 - **The keys assigned to each process should be sorted in (say) increasing order.**
 - **If $0 \leq q < r < p$, then each key assigned to process q should be less than or equal to every key assigned to process r .**



Serial bubble sort

```
void Bubble_sort(  
    int  a[]      /* in/out */,  
    int  n        /* in      */) {  
    int list_length, i, temp;  
  
    for (list_length = n; list_length >= 2; list_length--)  
        for (i = 0; i < list_length-1; i++)  
            if (a[i] > a[i+1]) {  
                temp = a[i];  
                a[i] = a[i+1];  
                a[i+1] = temp;  
            }  
} /* Bubble sort */
```



Odd-even transposition sort

- A sequence of phases.
- Even phases, compare swaps:

$(a[0], a[1]), (a[2], a[3]), (a[4], a[5]), \dots$

- Odd phases, compare swaps:

$(a[1], a[2]), (a[3], a[4]), (a[5], a[6]), \dots$



Example

Start: 5, 9, 4, 3

**Even phase: compare-swap (5,9) and (4,3)
getting the list 5, 9, 3, 4**

**Odd phase: compare-swap (9,3)
getting the list 5, 3, 9, 4**

**Even phase: compare-swap (5,3) and (9,4)
getting the list 3, 5, 4, 9**

**Odd phase: compare-swap (5,4)
getting the list 3, 4, 5, 9**

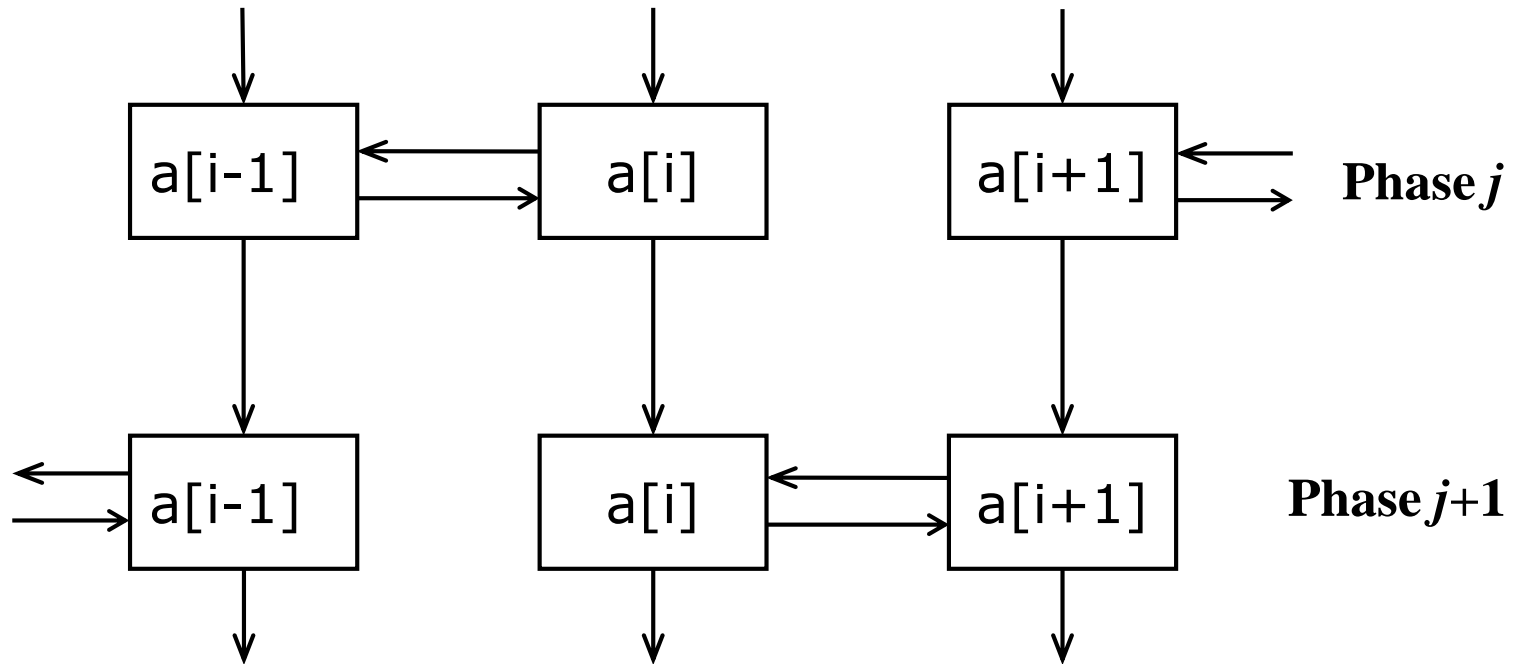


Serial odd-even transposition sort

```
void Odd_even_sort(  
    int    a[]    /* in/out */,  
    int    n      /* in      */) {  
    int phase, i, temp;  
  
    for (phase = 0; phase < n; phase++)  
        if (phase % 2 == 0) { /* Even phase */  
            for (i = 1; i < n; i += 2)  
                if (a[i-1] > a[i]) {  
                    temp = a[i]; a[i] = a[i-1]; a[i-1] = temp;  
                }  
        } else { /* Odd phase */  
            for (i = 1; i < n-1; i += 2)  
                if (a[i] > a[i+1]) {  
                    temp = a[i]; a[i] = a[i+1]; a[i+1] = temp;  
                }  
        }  
    } /* Odd_even_sort */
```



Communications among tasks in odd-even sort



Tasks determining $a[i]$ are labeled with $a[i]$.



Parallel odd-even transposition sort

15, 11, 9, 16 , 3, 14, 8, 7, 4, 6, 12, 10, 5, 2, 13, 1

Time	Process			
	0	1	2	3
Start	15, 11, 9, 16	3, 14, 8, 7	4, 6, 12, 10	5, 2, 13, 1
After Local Sort	9, 11, 15, 16	3, 7, 8, 14	4, 6, 10, 12	1, 2, 5, 13
After Phase 0	3, 7, 8, 9	11, 14, 15, 16	1, 2, 4, 5	6, 10, 12, 13
After Phase 1	3, 7, 8, 9	1, 2, 4, 5	11, 14, 15, 16	6, 10, 12, 13
After Phase 2	1, 2, 3, 4	5, 7, 8, 9	6, 10, 11, 12	13, 14, 15, 16
After Phase 3	1, 2, 3, 4	5, 6, 7, 8	9, 10, 11, 12	13, 14, 15, 16



Pseudo-code

```
Sort local keys;
for (phase = 0; phase < comm_sz; phase++) {
    partner = Compute_partner(phase, my_rank);
    if (I'm not idle) {
        Send my keys to partner;
        Receive keys from partner;
        if (my_rank < partner)
            Keep smaller keys;
        else
            Keep larger keys;
    }
}
```

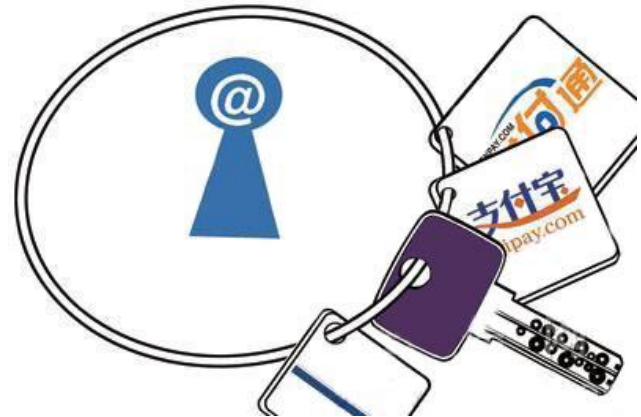


Compute_partner

```
if (phase % 2 == 0)                /* Even phase */
    if (my_rank % 2 != 0)           /* Odd rank */
        partner = my_rank - 1;
    else                             /* Even rank */
        partner = my_rank + 1;
else                                /* Odd phase */
    if (my_rank % 2 != 0)           /* Odd rank */
        partner = my_rank + 1;
    else                             /* Even rank */
        partner = my_rank - 1;
if (partner == -1 || partner == comm_sz)
    partner = MPI_PROC_NULL;
```



SAFETY IN MPI PROGRAMS



Safety in MPI programs

- **The MPI standard allows MPI_Send to behave in two different ways:**
 - it can simply copy the message into an MPI managed buffer and return,
 - or it can block until the matching call to MPI_Recv starts.



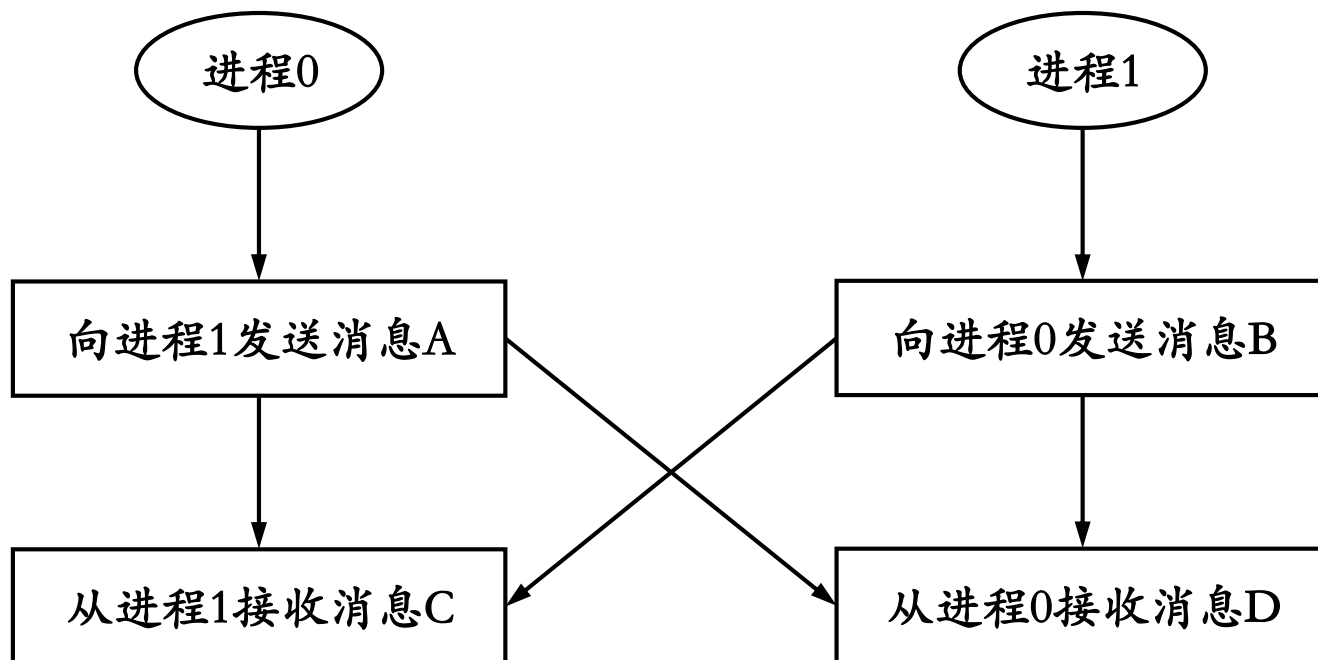
Safety in MPI programs

- Many implementations of MPI set a threshold at which the system switches from buffering to blocking.
- Relatively small messages will be buffered by MPI_Send.
- Larger messages, will cause it to block.

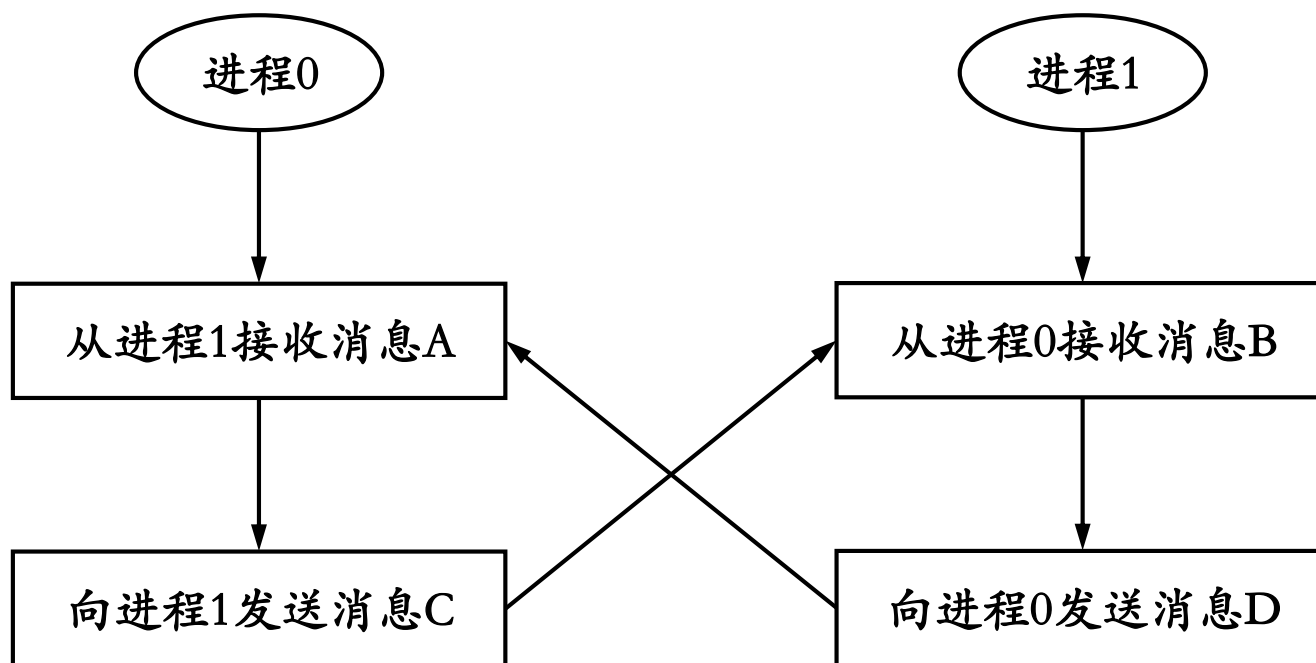


避免死锁deadlock

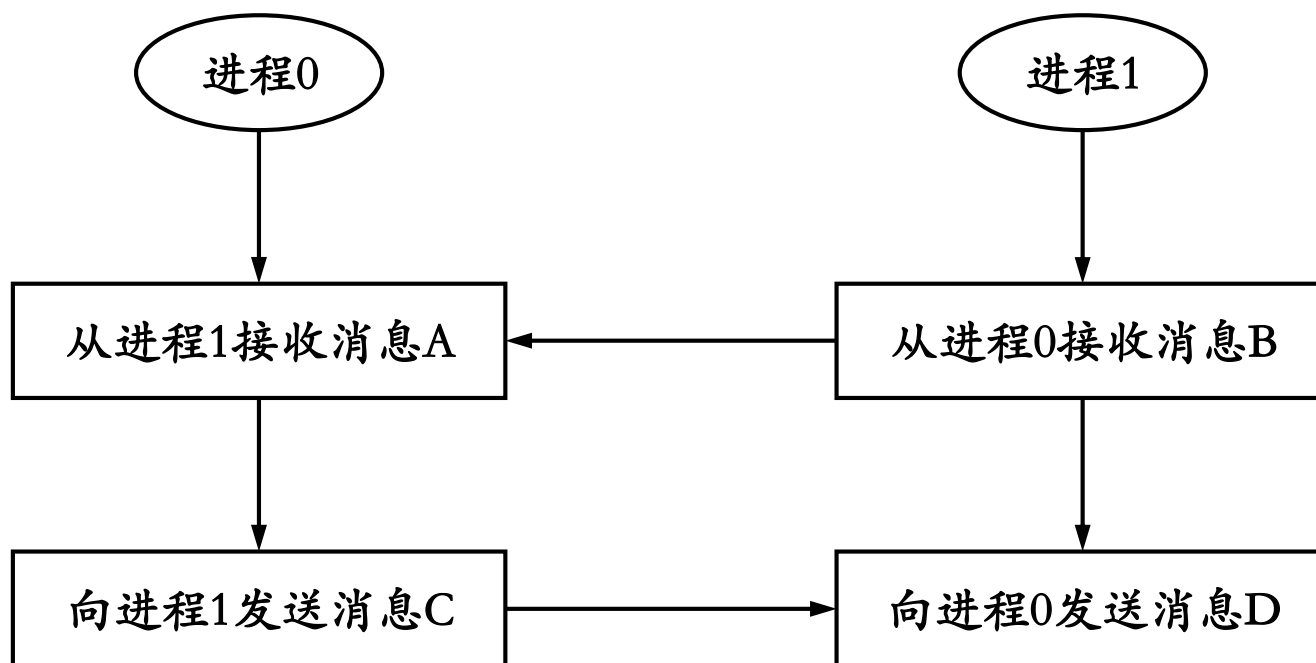
- 发送和接收是成对出现的，忽略这个原则 很可能产生死锁



不安全的通信调用次序

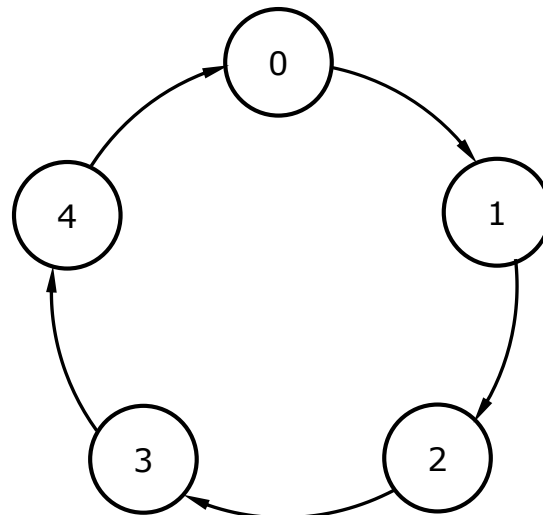


安全的通信调用次序



Safety in MPI programs

- If the MPI_Send executed by each process blocks, no process will be able to start executing a call to MPI_Recv, and the program will hang or **deadlock**.
- Each process is blocked waiting for an event that will never happen.



Safety in MPI programs

- A program that relies on MPI provided buffering is said to be **unsafe**.
- Such a program may run without problems for various sets of input, but it may hang or crash with other sets.



MPI_Ssend

- An alternative to MPI_Send defined by the MPI standard.
- The extra “s” stands for **s**ynchronous and MPI_Ssend is guaranteed to block until the matching receive starts.

```
int MPI_Ssend(  
    void*          msg_buf_p      /* in */,  
    int            msg_size       /* in */,  
    MPI_Datatype    msg_type      /* in */,  
    int            dest           /* in */,  
    int            tag            /* in */,  
    MPI_Comm        communicator /* in */);
```



Restructuring communication

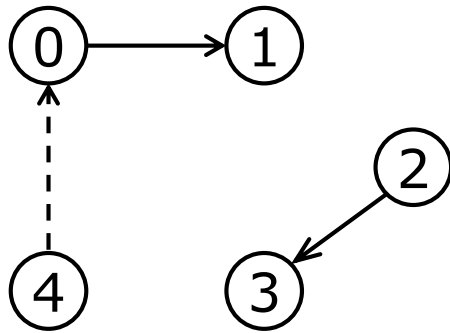
```
MPI_Send(msg, size, MPI_INT, (my_rank+1)%comm_sz, 0, comm);  
MPI_Recv(new_msg, size, MPI_INT, (my_rank+comm_sz-1)%comm_sz,  
0, comm, MPI_STATUS_IGNORE).
```



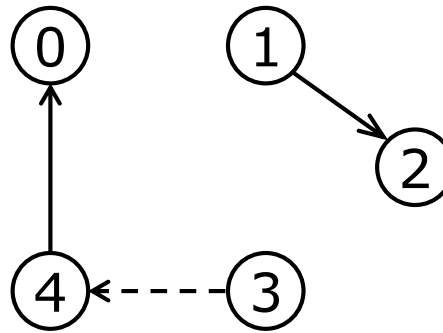
```
if (my_rank % 2 == 0) {  
    MPI_Send(msg, size, MPI_INT, (my_rank+1)%comm_sz, 0, comm);  
    MPI_Recv(new_msg, size, MPI_INT, (my_rank+comm_sz-1)%comm_sz,  
0, comm, MPI_STATUS_IGNORE).  
} else {  
    MPI_Recv(new_msg, size, MPI_INT, (my_rank+comm_sz-1)%comm_sz,  
0, comm, MPI_STATUS_IGNORE).  
    MPI_Send(msg, size, MPI_INT, (my_rank+1)%comm_sz, 0, comm);  
}
```



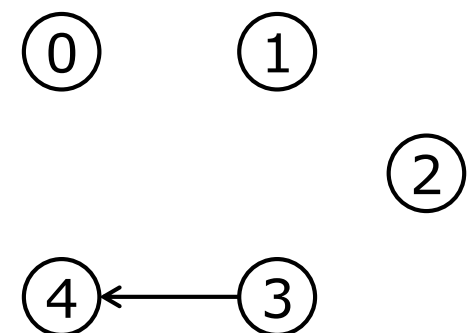
Safe communication with five processes



Time 0
偶发奇收



Time 1
奇发偶收



Time 2
偶发奇收



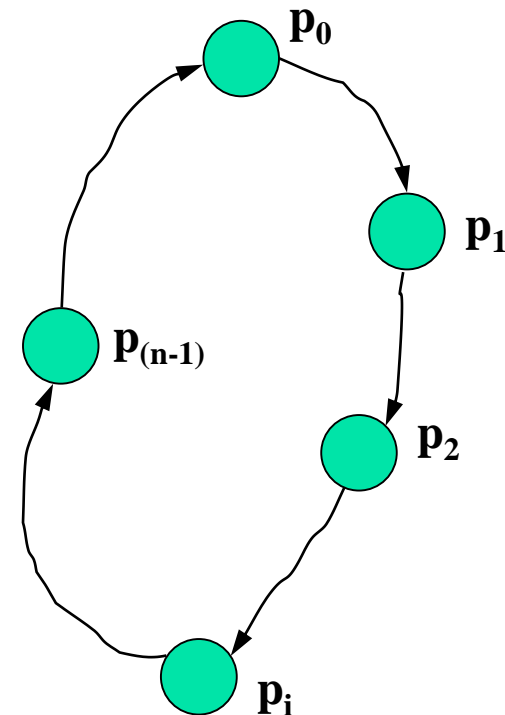
MPI_Sendrecv

- An alternative to scheduling the communications ourselves.
- Carries out a blocking send and a receive in a single call.
- The dest and the source can be the same or different.
- Especially useful because MPI schedules the communications so that the program won't hang or crash.



MPI_Sendrecv

```
int MPI_Sendrecv(  
    void*      send_buf_p      /* in */,  
    int        send_buf_size   /* in */,  
    MPI_Datatype send_buf_type /* in */,  
    int        dest            /* in */,  
    int        send_tag        /* in */,  
    void*      recv_buf_p      /* out */,  
    int        recv_buf_size   /* in */,  
    MPI_Datatype recv_buf_type /* in */,  
    int        source          /* in */,  
    int        recv_tag        /* in */,  
    MPI_Comm   communicator    /* in */,  
    MPI_Status* status_p       /* in */);
```



MPI_Sendrecv用法示意

```
...  
int a, b;  
...  
MPI_Status status;  
int dest = (rank + 1)%p;  
int source = (rank + p -1)%p;    /* p为进程个数 */  
MPI_Sendrecv( &a, 1, MPI_INT, dest, 99, &b, 1, MPI_INT,  
              source, 99, MPI_COMM_WORLD, &status);
```

该函数被每一进程执行一次。



Parallel odd-even transposition sort

```
void Merge_low(  
    int my_keys[],          /* in/out      */ My_keys[ ]和recv_keys[ ]  
    int recv_keys[],        /* in         */ 都是已排序的列表!  
    int temp_keys[],        /* scratch    */  
    int local_n             /* = n/p, in */){  
    int m_i, r_i, t_i;  
  
    m_i = r_i = t_i = 0;  
    while (t_i < local_n) {  
        if (my_keys[m_i] <= recv_keys[r_i]){  
            temp_keys[t_i] = my_keys[m_i];  
            t_i++; m_i++;  
        } else {  
            temp_keys[t_i] = recv_keys[r_i];  
            t_i++; r_i++;  
        }  
    }  
    for (m_i = 0; m_i < local_n; m_i++)  
        my_keys[m_i] = temp_keys[m_i];  
} /* Merge_low */
```



Run-times of parallel odd-even sort

Processes	Number of Keys (in thousands)				
	200	400	800	1600	3200
1	88	190	390	830	1800
2	43	91	190	410	860
4	22	46	96	200	430
8	12	24	51	110	220
16	7.5	14	29	60	130

(times are in milliseconds)



Concluding Remarks (1)

- **MPI or the Message-Passing Interface is a library of functions that can be called from C, C++, or Fortran programs.**
- **A communicator is a collection of processes that can send messages to each other.**
- **Many parallel programs use the single-program multiple data or SPMD approach.**



Concluding Remarks (2)

- **Most serial programs are deterministic: if we run the same program with the same input we'll get the same output.**
- **Parallel programs often don't possess this property.**
- **Collective communications involve all the processes in a communicator.**



Concluding Remarks (3)

- When we time parallel programs, we're usually interested in elapsed time or “wall clock time”.
- Speedup is the ratio of the serial run-time to the parallel run-time.
- Efficiency is the speedup divided by the number of parallel processes.



Concluding Remarks (4)

- If it's possible to increase the problem size (n) so that the efficiency doesn't decrease as p is increased, a parallel program is said to be scalable.
- An MPI program is unsafe if its correct behavior depends on the fact that MPI_Send is buffering its input.



MPI常用例程一览表

1	MPI_Init	
2	MPI_Finalize	
3	MPI_Comm_rank	
4	MPI_Comm_size	
5	MPI_Wtime	
6	MPI_Send	MPI_Ssend, MPI_Isend
7	MPI_Recv	MPI_Irecv
8	MPI_Sendrecv	MPI_Sendrecv_replace
9	MPI_Status	
10	MPI_Barrier	
11	MPI_Bcast	
12	MPI_Gather	MPI_Allgather
13	MPI_Scatter	
14	MPI_Reduce	MPI_Allreduce