

Exam 26 October 2012, 8:00–13:00, Sparta:A–B

## EDAN55 Advanced Algorithms

*preliminary alpha draft*

The exam consists of 4 large questions; each consisting of a number of smaller subquestions.

1. The exam is “open book,” so you can bring whatever material you want, including textbooks, a dictionary, and your own course notes.
2. You can bring an electronic calculator, even though I don’t see how that is useful.
3. We try to minimise the dependencies among subquestions. In particular, you can solve them in any order and are free to *use* the result of subquestion  $x$  to answer subquestion  $y$ , even if you didn’t answer  $x$ .
4. Scoring: Answering “I don’t know” (and nothing else) scores  $\frac{1}{4}$  of a subquestion’s points. An empty or wrong answer scores 0 points.
5. You can answer in Swedish or English.

Some tips:

1. Shorter is better.
2. An example is better than a failed attempt at explaining something in general.
3. Drawings, pseudocode, and formulas are good. “Wall of text” is bad.
4. Admit ignorance.
5. Be tidy.

*Good luck!*

### Question 1, Approximation

The **Maximum Independent Set** problem is defined as follows. Given an undirected graph  $G = (V, E)$ , find a **subset of vertices** such that **no two vertices are adjacent**. The **size of the independent set** is maximised. Formally, a **subset of vertices** is an independent set if **no two vertices are adjacent**. Define the **Maximum Independent Set** problem as follows:

We want to find a **subset of vertices** of maximum **size**.

►1a (1 pt.) Find a maximum **independent set** the graph in fig. 1.<sup>1</sup>

Consider the following simple algorithm for **Maximum Independent Set**. Let  $V = \{v_1, \dots, v_n\}$  and for **each vertex  $v_i$  in  $V$**

1. Set  **$S = \emptyset$**
2. For **each vertex  $v_i$  in  $V$**   
**if  $v_i$  is not adjacent to any vertex in  $S$**   
**then add  $v_i$  to  $S$**

►1b (1 pt.) Run the algorithm on the graph in fig. 1 and give the result.<sup>2</sup>

►1c (1 pt.) Give an example where the algorithm finds a **subset of vertices**

►1d (3 pts.) Show that the algorithm is guaranteed to **find a maximum independent set**.<sup>4</sup>

►1e (2 pts.) Prove that there cannot be an algorithm for **Maximum Independent Set** that approximates the optimum solution within **factor  $c$**  unless  **$P = NP$** . You can freely use that  **$P \neq NP$**  is **conjectured**.

►1f (1 pt.) Modify the algorithm and analysis for the **Maximum Independent Set** problem, where **vertices have weights** and **the goal is to find a maximum weight independent set**.

►1g (1 pt.) Modify the algorithm and analysis for the **Maximum Independent Set** problem, where **vertices have weights** and **the goal is to find a maximum weight independent set**.

<sup>1</sup> Your answer is a drawing showing **the vertices of the maximum independent set** and an integer **representing its size**.

<sup>2</sup> Your answer is the **size of the maximum independent set**.

<sup>3</sup> Your answer is a concrete graph, an optimum solution to that instance and the solution found by the algorithm.

<sup>4</sup> Your answer is a short proof. It includes a lower bound on the solution found by the algorithm and an upper bound on the optimum solution.

<sup>5</sup> Your answer is a short proof.

<sup>6</sup> Your answer contains an algorithm and a brief analysis of its approximation guarantee; the most important part is to state the resulting approximation factor.

<sup>7</sup> Your answer contains an algorithm and a brief analysis of its approximation guarantee; the most important part is to state the resulting approximation factor.





### Question 4, Randomized Algorithms

We consider the [redacted]

Name: [redacted]

Input: A simple, undirected graph  $G = (V, E)$  with  $n$  vertices and  $m$  edges.

Output: A [redacted]  
[redacted] (or both).

Consider the following randomized algorithm for this problem:

1. For every [redacted] pick [redacted] at random from [redacted]
2. Let [redacted]  
[redacted]  
[redacted]

where  $N(v)$  denotes the neighbours of  $v$ .

- 4a (1 pt.) Run the algorithm on the graph in fig. ?? [redacted]  
[redacted]
- 4b (1 pt.) [redacted] Find the probability that [redacted]<sup>11</sup>
- 4c (1 pt.) Consider running the algorithm on [redacted]  
[redacted] Are the events [redacted] and [redacted] independent? Why or why not?<sup>12</sup>
- 4d (1 pt.) What is the probability that [redacted]
- 4e (2 pt.) Assume that  $G$  is  $d$ -regular (i.e., every vertex as degree exactly  $d$ ). Find [redacted] expected size of [redacted]  
[redacted]
- 4f (1 pt.) Assume  $G$  is the  $n$ -cycle. (That is,  $E = \{\{i, i+1\}: 1 \leq i < n\} \cup \{n, 1\}$ .) Find the expected [redacted]  
[redacted]
- 4g (1 pt.) Assume  $G$  is the  $n$ -star. (That is,  $E = \{\{1, i\}: 2 \leq i \leq n\}$ .) What is the probability that [redacted]  
[redacted]<sup>16</sup>

<sup>10</sup> Your answer is a drawing s [redacted]

<sup>11</sup> Your answer is an expression and an argument for it.

<sup>12</sup> Your answer is the word "yes" or the word "no", followed by an argument.

<sup>13</sup> This is not meant to be a trick question. But if you think your answer is weird, it's probably correct.

<sup>14</sup> Your answer is an expression and an argument for it.

<sup>15</sup> Your answer is an expression and an argument for it.

<sup>16</sup> Your answer is an expression and an argument for it.

►48 (1 pt.) Assume  $G$  is the  $n$ -star. (That is,  $E = \{\{1, i\} : 2 \leq i \leq n\}$ .) Calculate the

probability of

<sup>17</sup>

<sup>17</sup> Your answer is an expression and an argument for it.