All variables are ints.

```
PRECONDITION: x > 0
z = 0;
y = x;
while (y \% 10 == 0) {
 y = y / 10 // integer division
 z = z + 1
POSTCONDITION: { x == y * 10^z \&\& (y \% 10 != 0) }
Example: x = 1100; x = 11*10^2
LI: \{x == y * 10^z\}
Base case:
y == x \& \& z == 0
{x == y * 10^0} == {y == x * 1} == {y == x}
Iteration k:
assume: x == y_k * 10^{z_k}
y_{k+1} = \frac{y_k}{10}
z_{k+1} = z_k + 1
\{y_{k+1} * 10^{z_{k+1}} == \frac{y_k}{10} * 10^{z_k+1}\} == \{y_k * 10^{z_k} == x\}
At exit:
```

# Termination:

One choice for D = the number of trailing 0's in y

 $\{!(y\%10 == 0) \& \& x == y*10^z\} \rightarrow \{y\%10 \neq 0 \& \& x == y*10^z\}$ 

```
D = String.valueOf(y).length() -
String.valueOf(y).replaceAll("0*$","").length();
```

```
Alternative: (y % 10 == 0 ? 1:0) * int(log10(y)) int(log(y)) is one less than the number of digits in y – fails for y=0 i.e. only works for precondition for x > 0.
```

D = number of trailing zeros in y, didn't equal 0 or we would have exited loop

D\_new = number of trailing zeros in y\_new

 $y_new = y_old/10$ 

D\_new has 1 less trailing 0.

D\_new < D\_old

D = (y % 10 == 0 ? 1 : 0) \* int(log10(y)) = int(log10(y)) // D is not zero or we would exit

D\_new = 
$$(y_new \% 10 == 0 ? 1 : 0) * int(log10(y_new))$$
  
=  $(y/10 \% 10 == 0 ? 1 : 0) * int(log10(y/10))$   
=  $1 * (int(log10(y)) - int(log10(10)))$   
= D - 1 if first term was 1, otherwise D\_new = 0 and D decreases either way.

At Exit:

$$D = 0 =>$$
 no trailing zeros  $=>$  y % 10 != 0

Using the other decrement function

$$D = 0 \Rightarrow y \% 10 = 0$$

gcd is the greatest common divisor of two positive integers, i.e. the largest integer number that evenly divides both numbers.

```
PRECONDITION: \{ x1 > 0 \& x2 > 0 \}
y1 = x1
y2 = x2
while ( y1 != y2 ) {
  if ( y1 > y2 ) {
    y1 = y1 - y2
   élse {
     y2 = y2 - y1
}
POSTCONDITION: { y1 == gcd(x1, x2) }
Some gcd facts:
gcd(x,x)=x
gcd(x,y) = gcd(x-y, y)
Proof:
x=ad, y=bd
x-y = ad - bd = (a-b)d => d is a divisor of x-y, as well as x and y
At exit, we want y1==gcd(x,y) & y1==y2 (exit condition)
since gcd(y1,y1) == y1 and at exit y2==y1, a good guess might be
LI: gcd(y1, y2) = gcd(x1, x2)
Let's see if it works
```

## Initial step:

$$y1 = x1, y2 = x2;$$
  
 $gcd(y1, y2) = gcd(x1, x2)$ 

#### Iteration k+1:

assume: 
$$gcd(y1_k, y2_k) = gcd(x1, x2)$$
  
 $y1_k < y2_k$   
 $y1_{k+1} = y1_k - y2_k$   
 $gcd(y1_{k+1}, y2_{k+1}) = gcd(y1_k - y2_k, y2_k) = gcd(y1_k, y2_k) = gcd(x1, x2)$ 

Similar proof for y2 > y1. If y1 = = y2, we exit loop.

### At Exit:

#### Termination:

At each iteration, we choose max(y1, y2). At the end, y1 == y2 == gcd(x1,x2).

A reasonable choice for D might be:

 $D = \max(y1,y2) - \gcd(y1,y2)$ 

The minimum is 0 and it should decrease at each iteration.

Minimum occurs when y1 == y2

$$\begin{split} &D_k = \max(y1_k, y2_k) - \gcd(y1_k, y2_k) \\ &D_{k+1} = \max(y1_{k+1}, y2_{k+1}) - \gcd(y1_{k+1}, y2_{k+1}) \\ &y1_k > y2_k \\ &D_{k+1} = \max(y1_k - y2_k, y2_k) - \gcd(y1_k - y2_k, y2_k) < \max(y1_k, y2_k) - \gcd(y1_k, y2_k) \\ &\therefore D_{k+1} < D_k \end{split}$$

// reduce larger y

Similar proof for y2>y1.

At exit: D == 0 max(y1, y2) - gcd(y1, y2) ==0 max(y1, y2) == gcd(y1, y2) => y1 == y2