CSE 127: Introduction to Security

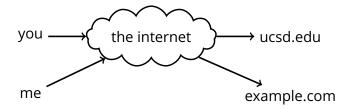
Lecture 10: Intro to Networking

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Winter 2023

The Internet



Original Idea:

- Network is dumb
- Simple, robust service
- Shift complexity to endpoints
- Acts like postal system (packet-based) rather than traditional phone system (circuit-based)

Need protocol to actually communicate

A protocol is an agreement on how to communicate.

Includes syntax and semantics.

- **Syntax:** How communication is specified and structured.
 - Format, order messages are sent and received.

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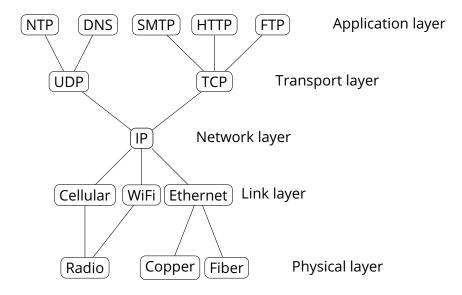
- Syntax: How communication is specified and structured.
 - Format, order messages are sent and received.
- Semantics: What a communication means
 - Actions taken when transmitting, receiving, or timer expires.
- Example: RFC 2616 (HTTP/1.1)
 - Section 5: Syntax of HTTP Requests
 - Section 9.3: Semantics of GET Requests

Protocols are layered

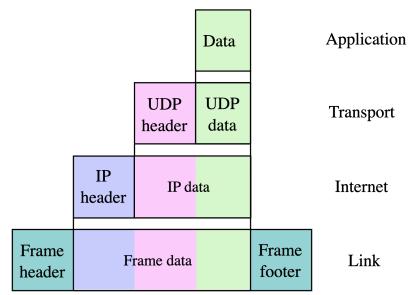
- Networks use a stack of layers
- Lower layers provide services to layers above
 - Don't care what higher layers do
- Higher layers use services of layers below
 - Don't care how lower layers implement services
- Layers define abstraction boundaries
 - At a given layer, all layers above and below are opaque

Basic Internet Archictecture "Hourglass"

Narrow waist = interoperability



Packet encapsulation at each layer



Source: Wikipedia

Link layer: Connecting hosts to local network

Most common link layer protocol: **Ethernet**



- Messages organized into frames
- Every node has a globally unique 6-byte MAC address

Source: Wikipedia

Link layer: Connecting hosts to local network

- Originally a broadcast protocol: every node on network received every packet
- Now switched: switch learns the physical port for each MAC address and sends packets to correct port if known
- WiFi similar to Ethernet, but nodes can move

IP: Internet Protocol

- Connectionless delivery model
- "Best effort" = no guarantees about delivery
- No attempt to recover from failure
- Packets might be lost, delivered out of order, delivered multiple times
- Packets might be fragmented
- Provides hierarchical addressing scheme

IP: Internet Protocol

- IPv4
 - 32-bit host addresses
 - Written as 4 bytes in decimal,
 - e.g. 192.168.1.1
- IPv6
 - 128-bit host addresses
 - Written as 16 bytes in hex
 - :: implies zero bytes
 - e.g. 2620:0:e00:b::53 = 2620:0:e00:b:0:0:0:53

```
0
1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1 2 3 4 5 6 7 8 9 0 1
     |Type of Service|
   IHL
               Total Length
|Flags|
   Identification
                Fragment Offset
Time to Live |
      Protocol |
               Header Checksum
Source Address
Destination Address
Options
                   Padding
```

Example Internet Datagrarm Header

Note that each tick mark represents one bit position.

ARP: Address Resolution Protocol

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- ARP request: source MAC, dest MAC, "Who has IP address N?"
- ARP reply: source MAC, dest MAC, "IP address N is at MAC address M."

Routing: BGP (Border Gateway Protocol)

- Internet organized into ASes (Autonomous Systems) with peer, provider, or customer relationships between them
- Rough tree shape, with a small number of backbone ASes in a clique at the root

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- Internet organized into ASes (Autonomous Systems) with peer, provider, or customer relationships between them
- Rough tree shape, with a small number of backbone ASes in a clique at the root
- BGP allows routers to exchange information about their routing tables
- Routers maintain global table of routes
- Each router announces what it can route to its neighbors
- Routes propagate through network

TCP (Transmission Control Protocol)

- Want abstraction of a stream of bytes delivered reliably and in-order between applications on different hosts
- TCP provides:
 - Reliable in-order byte stream
 - Connection-oriented protocol
 - Explicit setup/teardown
 - End hosts (processes) have multiple concurrent long-lived dialogs
 - Congestion control: adapt to network path capacity, receiver's ability to receive packets

TCP Header Format

Ports

- Each application is identified by a port number
- TCP connection established between port A on host address M to port B on host address N. Ports are 16 bits, 1–65535
- Some destination ports are used for particular applications by convention

```
80 HTTP (web)
```

443 HTTPS (web)

25 SMTP (mail)

67 DHCP (host configuration)

22 SSH (secure shell)

23 telnet

TCP Sequence Numbers

- Bytes in application data stream numbered with 32-bit sequence number
- Data sent in segments: sequences of contiguous bytes sent in a single IP datagram
- Sequence number indicates where data belongs in byte sequence
- Sequence number in packet header is the sequence number of the first byte in the payload

TCP Sequence Numbers and Acknowledgement

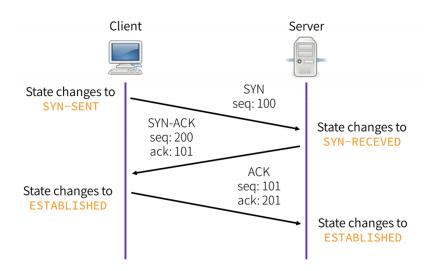
- Two logical data streams in a TCP connection: one in each direction
- Receiver acknowledges received data: acknowledgement number is sequence number of next expected byte of stream in opposite direction
- ACK flag set to acknowledge data
- Sender retransmits lost data
- Congestion control: sender adapts retransmission according to timeouts

TCP 3-Way Handshake

Starting a TCP connection

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FIN/RST: Closing TCP connections

• FIN initiates a clean close of a TCP connection, waits for ACK from receiver

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- FIN initiates a clean close of a TCP connection, waits for ACK from receiver
- If a host receives a TCP packet with RST flag, it tears down the connection
- Designed to handle spurious TCP packets from previous connections

UDP (User Datagram Protocol)

- UDP offers no service quality guarantee
- Essentially a transport layer protocol that is a wrapper around IP
- Adds ports to let applications demultiplex traffic
- Useful for applications that only need best-effort guarantee
- e.g. DNS, NTP

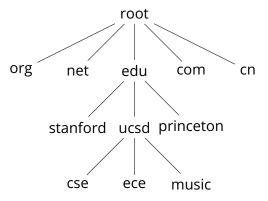
J. Postel ISI 28 August 1980

	User Da	atagram	Protocol	
0	7 8	15 16	23 24	31
+	+	+	+	+
	Source	1	Destination	1
	Port	1	Port	- 1
+	+	+	+	+
		1		- 1
	Length	1	Checksum	- 1
+	+	+	+	+
1				
data octets				
+				

User Datagram Header Format

DNS (Domain Name Service)

- Handle mapping between host names (e.g. ucsd.edu) and IP addresses (e.g. 132.239.180.101)
- DNS is a delegatable, hierarchical name space



DNS Records

```
$ dig cseweb.ucsd.edu
: <<>> DiG 9.10.6 <<>> cseweb.ucsd.edu
;; global options: +cmd
;; Got answer:
:: ->>HEADER<<- opcode: QUERY, status: NOERROR, id: 3727
;; flags: qr rd ra; QUERY: 1, ANSWER: 2, AUTHORITY: 0, ADDITIONAL: 1
;; OPT PSEUDOSECTION:
; EDNS: version: 0, flags:; udp: 4096
;; QUESTION SECTION:
:cseweb.ucsd.edu. TN A
;; ANSWER SECTION:
cseweb.ucsd.edu. 3140 IN CNAME roweb.eng.ucsd.edu.
roweb.eng.ucsd.edu. 2855 IN A 132.239.8.30
;; Query time: 57 msec
:: SERVER: 192.168.1.254#53(192.168.1.254)
:: WHEN: Sun Nov 03 20:49:08 PST 2019
;; MSG SIZE rcvd: 84
```

DNS Details

- 13 main DNS root servers
- DNS responses are cached for quicker responses
- DNS authorities queried progressively according to domain name hierarchy

```
$ dig cseweb.ucsd.edu +trace
: <<>> DiG 9.10.6 <<>> cseweb.ucsd.edu +trace
;; global options: +cmd
. 105604 IN NS d.root-servers.net.
. 105604 IN NS h.root-servers.net.
. 105604 IN NS c.root-servers.net.
. 105604 IN NS i.root-servers.net.
. 105604 IN NS l.root-servers.net.
. 105604 IN NS i.root-servers.net.
. 105604 IN RRSIG NS 8 0 518400 20191115050000 20191102040000 22545 . Z14B+vD/MKz0X1UBwu04kzw0Naihg1Af1K7i5Jvd9N
;; Received 525 bytes from 192.168.1.254#53(192.168.1.254) in 44 ms
edu 172800 IN NS b edu-servers net
edu 172800 IN NS f edu-servers net
edu. 172800 IN NS i.edu-servers.net.
```

9DHS4EP5G85PF9NUFK06HEK0048QGK77.edu, 86400 IN NSEC3 1 1 0 - 9V5L4LUB1VNJ9EQQLIHEQCBREACL2500 NS SOA RRSIG DNSKI 9DHS4EP5G85PF9NUFK06HEK0048QGK77.edu, 86400 IN RRSIG NSEC3 8 2 86400 20191111043435 20191104032435 47252 edu, M5 3FTB9RSLR00JU0PDNLJJE2T31U25M4MG.edu. 86400 IN NSEC3 1 1 0 - 4586U2HHMPSEA0HJD6R9INNA38P0F8KL NS DS RRSIG 3FTB9RSLR0QJU0PDNLJJE2I31U25M4MG.edu. 86400 IN RRSIG NSEC3 8 2 86400 20191111041950 20191104030950 47252 edu. BK

```
edu. 86400 IN RRSIG DS 8 1 86400 20191116170000 20191103160000 22545 . BsoO9WI4UphacN5rLOB4f3bCzVPptbmTCKHwcMgb6
;; Received 1174 bytes from 192.58.128.30#53(j.root-servers.net) in 20 ms
```

edu. 172800 IN NS c.edu-servers.net. edu. 172800 IN NS e.edu-servers.net. edu 172800 IN NS d edu-servers net

ucsd.edu. 172800 IN NS ns-auth2.ucsd.edu.

ucsd.edu. 172800 IN NS ns-auth3.ucsd.edu.

cseweb.ucsd.edu. 3600 IN CNAME roweb.eng.ucsd.edu. roweb.eng.ucsd.edu. 3600 IN A 132.239.8.30 ;; Received 84 bytes from 132.239.252.186#53(ns-auth3.ucsd.edu) in 14 ms

:: Received 671 bytes from 192.41.162.30#53(1.edu-servers.net) in 9 ms

edu. 86400 IN DS 28065 8 2 4172496CDE85534E51129040355BD04B1FCFEBAE996DFDDE652006F6 F8B2CE76

You connect your laptop to a cafe wifi network and type ucsd.edu into your browser's URL bar. What happens?

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 - Broadcasts DHCPDISCOVER to 255.255.255.255 with its MAC address
 - DHCP server responds with config: lease on host IP address, gateway IP address, DNS server information

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 - Every connection outside the local network will be encapsulated in a link-layer frame with the local router's MAC address as the desination.
 - Your laptop encapsulates each IP packet in a WiFi Ethernet frame addressed to the local router.
 - The local router decapsulates these Ethernet frames and re-encodes them to forward them on its fiber connection to its upstream ISP, or to another part of the network.
 - Each hop re-encodes the link layer for its own network.

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 - Each response tells the laptop what authority to query, until it learns the final IP address (75.2.44.127) for ucsd.edu
 - This address is cached, along with the authorities for the hierarchy in the hostname.

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 - The local router has a routing table that contains IP prefixes that it matches against the IP address that tells it what address to forward the packets to.
 - The packet passes through a series of Autonomous Systems (ASes).
 - From cafe network (ATT), go through sbcglobal.net → att.net → level3.net → cenic.net → ucsd.edu.

- 5. Your laptop sends a HTTP GET request inside the TCP connection.
- 6. Based on the HTTP response, the laptop performs a new DNS lookup, TCP handshake, and HTTP GET requests for every resource in the HTML as it renders.