

CS5413 - Data Structures and Algorithm Analysis III - Exam #3

(Total 100 points)

(In class, Closed book/notes)

(Justify every step of your answers)

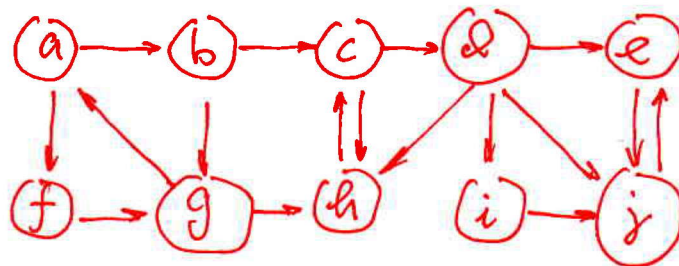
Name (Last name, First name):

Key

1. (30 points) Given the following algorithms and graph as described in the adjacency matrix, answer the following questions.

	a	b	c	d	e	f	g	h	i	j
a	0	1	0	0	0	1	0	0	0	0
b	0	0	1	0	0	0	1	0	0	0
c	0	0	0	1	0	0	0	1	0	0
d	0	0	0	0	1	0	0	1	1	1
e	0	0	0	0	0	0	0	0	0	1
f	0	0	0	0	0	0	1	0	0	0
g	1	0	0	0	0	0	0	1	0	0
h	0	0	1	0	0	0	0	0	0	0
i	0	0	0	0	0	0	0	0	0	1
j	0	0	0	0	1	0	0	0	0	0

From adjacency matrix, we have directed graph as following:



(a) Show the execution of the following algorithm. Mark each vertex with its discover time (d) and finish time (f), (d/f).

DFS(G)

```

1 for each vertex  $u \in V[G]$ 
2   do  $color[u] \leftarrow WHITE$ 
3    $\pi[u] \leftarrow NIL$ 
4 time  $\leftarrow 0$ 
5 for each vertex  $u \in V[G]$ 
6   do if  $color[u] = WHITE$ 
7     then DFS-VISIT(u)
  
```

DFS-VISIT(u)

```

1  $color[u] \leftarrow GRAY$ 
2 time  $\leftarrow time + 1$ 
3  $d[u] \leftarrow time$ 
4 for each  $v \in Adj[u]$ 
5   do if  $color[v] = WHITE$ 
6     then  $\pi[v] \leftarrow u$ 
7     DFS-VISIT(v)
8  $color[u] \leftarrow BLACK$ 
9  $f[u] \leftarrow time \leftarrow time + 1$ 
  
```

(a)

Execution step:

Assumption: The nodes ~~is~~ has been sorted by node's name in $V[G]$

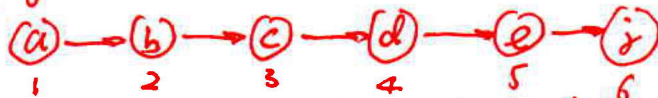
actions:

0;

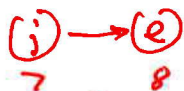
written out all nodes and ~~sort by node's name~~, time = 0

1

grayed out nodes and set times (d)



black out nodes and set finish time (f)



2:

grayed out nodes and set times: (h) ~~is~~ blacked out (h)

3:

grayed out: (i), blacked out: (i)

4:

black out (d)

5:

black out (c)

6:

grayed out (g), blacked out (g)

7:

black out (b)

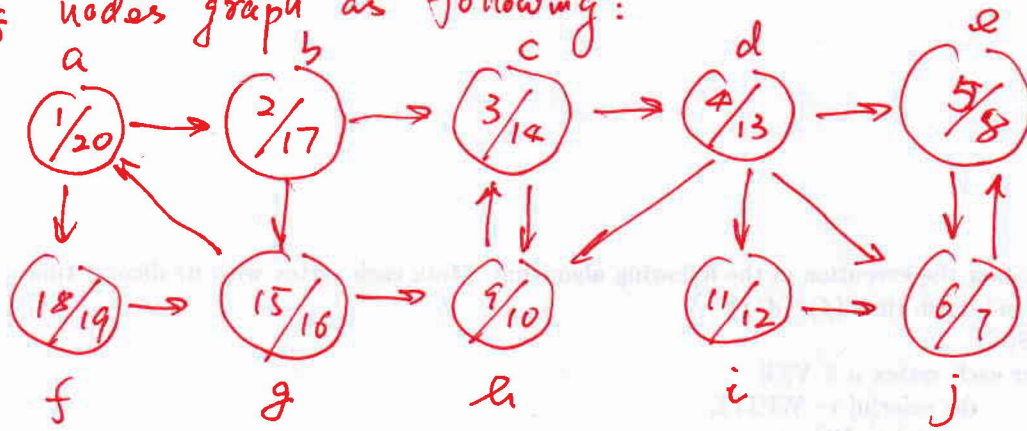
8:

grayed out (f), blacked out (f)

9:

black out (a)

d/s nodes graph as following:



if a node has two children, it is not a leaf node. The nodes are labeled a through j.

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(b) Show the execution of the following algorithm.

TOPOLOGICAL-SORT(G)

- 1 call DFS(G) to compute finishing times $f[v]$ for each vertex v
- 2 as each vertex is finished, insert it onto the front of a linked list
- 3 **return** the linked list of vertices

(c) Given directed graph $G = (V, E)$, a **strongly connected component (SCC)** of G is a maximal set of vertices $C \subseteq V$ such that for all $u, v \in C$, both $u \rightsquigarrow^b v$ and $v \rightsquigarrow^b u$. Find the SCCs in the graph.

- (b) \therefore The new node is always added before the head of the linkedlist.
 \therefore The earliest node would be the ~~first~~ ^{end} node.
 The last node would be the beginning node.
 \therefore The node insertion order is by the finish time.
 \therefore The order of nodes is like:

Finish time 20 19 17 16 14 13 12 10 8 7
 Linklist nodes: $\boxed{a} \rightarrow \boxed{f} \rightarrow \boxed{b} \rightarrow \boxed{g} \rightarrow \boxed{c} \rightarrow \boxed{d} \rightarrow \boxed{i} \rightarrow \boxed{h} \rightarrow \boxed{e} \rightarrow \boxed{j}$

(c) From the directed graph, we can see:

$\therefore (a \rightarrow b, b \rightarrow a), (a \rightarrow g, g \rightarrow a), (b \rightarrow g, g \rightarrow b), (a \rightarrow f, f \rightarrow a)$
 $(f \rightarrow g, g \rightarrow f), (b \rightarrow f, f \rightarrow b)$

$\therefore SCC_1 = \{a, b, f, g\}$

$\therefore (c \rightarrow d, d \rightarrow e), (c \rightarrow h, h \rightarrow c), (d \rightarrow h, h \rightarrow d)$

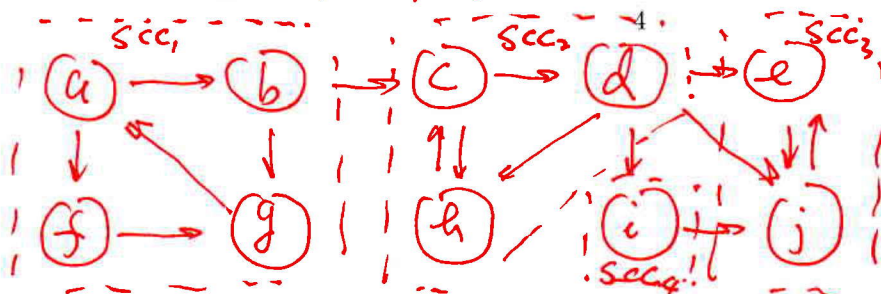
$\therefore SCC_2 = \{c, d, h\}$

$\therefore (e \rightarrow j, j \rightarrow e)$

$\therefore SCC_3 = \{e, j\}$

\therefore

$\therefore SCC_4 = \{i\}$. The SCCs in the graph as following:



2. (25 points) Given the following graph as described in the adjacency matrix,

	a	b	c	d	e	f	g	h	i
a	0	4	0	0	0	0	0	8	0
b	4	0	8	0	0	0	0	11	0
c	0	8	0	7	0	4	0	0	2
d	0	0	7	0	9	14	0	0	0
e	0	0	0	9	0	10	0	0	0
f	0	0	4	14	10	0	2	0	0
g	0	0	0	0	0	2	0	1	6
h	8	11	0	0	0	0	1	0	7
i	0	0	2	0	0	0	6	7	0

Show the execution of the following algorithm.

$$\text{MST-PRIM}(G, w, r)$$
1 **for** each $u \in V[G]$ 2 **do** key[u] $\leftarrow \infty$
$$3 \quad \pi[u] \leftarrow \text{NIE}$$
$$4 \text{ key}[r] \leftarrow 0$$
$$5 \quad Q \leftarrow V[G]$$

6 while $Q \neq \emptyset$

```

7   do  $u \leftarrow \text{EXTRACT-MIN}(Q)$ 

```

8 **for** each $v \in \text{Adj}[u]$

```

9      do if  $v \in Q$  and  $w(u,v) < \text{key}[v]$ 

```

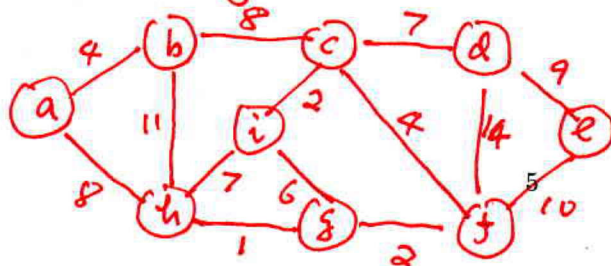
10 **then** $\pi[v] \leftarrow u$

```

11         key[v] ← w(u,v)

```

From the adjacency matrix, we have the nodes graph with weights as following:



Execution steps

Actions

#0.

choose a node 'a' as starting node, add a to V_A

#1.

find a light edge crossing cut $(V_A, V - V_A) = 4$.
add b to V_A .

#2.

the light edge $(V_A, V - V_A) = 8$. ^{here, I choose c and} add c to V_A

#3.

the light edge $(V_A, V - V_A) = 2$. add i to V_A

#4.

the light edge $(V_A, V - V_A) = 4$. add f to V_A

#5.

the light edge $(V_A, V - V_A) = 2$, add g to V_A

#6.

the light edge $(V_A, V - V_A) = 1$, add h to V_A

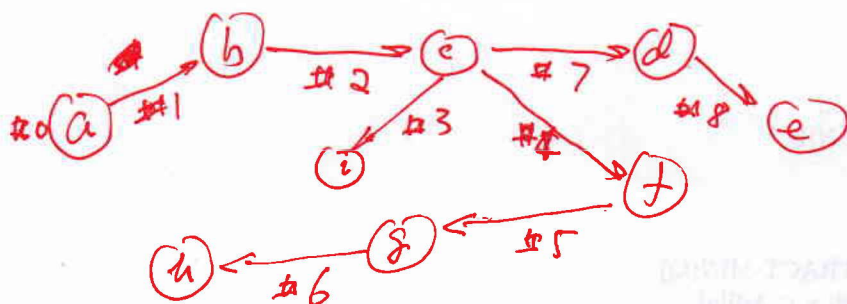
#7.

the light edge $(V_A, V - V_A) = 7$. add d to V_A

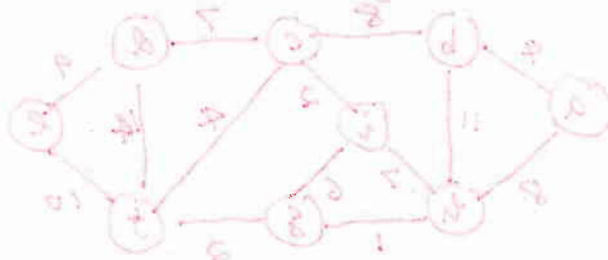
#8.

the light edge $(V_A, V - V_A) = 9$. add e to V_A

The execution path as following:



From the adjacency matrix, we have the undirected graph with weights as following:



3. (25 points) Given the following graph as described in the adjacency matrix, show the execution of the following algorithm.

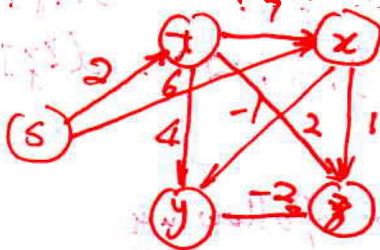
	s	t	x	y	z
s	0	2	6	0	0
t	0	0	7	4	2
x	0	0	0	-1	1
y	0	0	0	0	-2
z	0	0	0	0	0

a-shortest-path(V, E, w, s)

- 1 topologically sort the vertices in V
- 2 initialize-single-source(V, s)
- 3 for each vertex u , taken in topologically sorted order
- 4 do for each vertex $v \in \text{Adj}[u]$
- 5 do RELAX(u, v, w)

directed

From the adjacency matrix, we have the nodes graph as following:



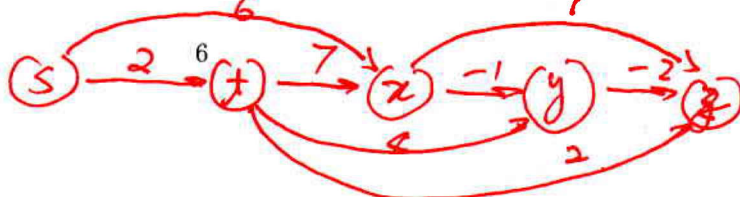
Execution steps:

Actions:

0

use the way of question (c) to do topological sorting.
 \therefore (z) would be the first reached the finish time.
 and (y) would be 2nd, and (x) would be 3rd, and
 soon, the last one is (s).

\therefore We have the DAG as:



1.

init-single-source (V, s)

we have $d[s]$ as:

$d[s] = 0, d[t] = \infty, d[x] = \infty, d[y] = \infty, d[z] = \infty$

2.

RELAX($s, t, 2$) when s is u node, t is v node. 2 is weight

we have $d[t] = 2$

3.

~~RELAX($t, x, 7$)~~ and RELAX($s, x, 6$)

we have $d[x] = 6$

~~# 4.~~

~~RELAX($t, y, 4$) and RELAX($x, y, 4$)~~

~~we have $d[y] = 6$~~

4

RELAX($x, x, 7$), $d[x]$ don't change

5

RELAX($t, y, 4$), $d[y] = 2 + 4 = 6$

6

RELAX($t, z, 2$), $d[z] = 2 + 2 = 4$

7

RELAX($x, y, -1$), $d[y] = 6 - 1 = 5$

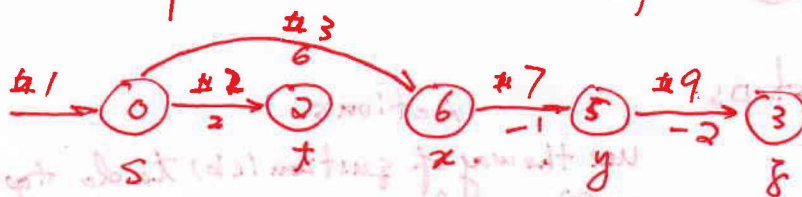
8

RELAX($x, z, 1$), $d[z]$ no change

9

RELAX($y, z, -2$), $d[z] = 5 - 2 = 3$

Effective Execution path with $d[s]$ as following



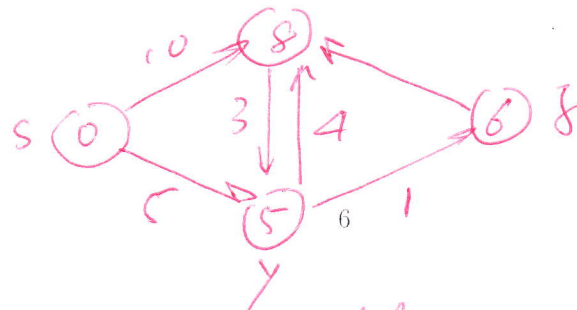
3. (20 points) Given the following graph as described in the adjacency matrix, show the execution of the following algorithm.

	s	x	y	z
s	0	10	5	0
x	0	0	3	0
y	0	4	0	1
z	0	2	0	0

```

DIJKSTRA(G,w,s)
1 INITIALIZE-SINGLE-SOURCE(G,s)
2 S ← ∅
3 Q ← V[G]
4 while Q ≠ ∅
5     do u ← EXTRACT-MIN(Q)
6       S ← S ∪ {u}
7       for each vertex v ∈ Adj[u]
8         do RELAX(u,v,w)
    
```

Init: $d[s] = 0$
 #1. $s \rightarrow y$, $d[y] = 5$. #2. $s \rightarrow x$, $d[x] = 10$
 #3. $y \rightarrow x$ $\because 5 + 4 = 9 < 10 \therefore d[x] = 9$
 #4. $y \rightarrow z$, $d[z] = 6$. #5. $x \rightarrow y$ $\because 9 + 3 = 12 > 5 \therefore d[y] = 5$
 #6. $x \rightarrow z$, $\because 9 + 2 = 11 > 6 \therefore d[z] = 6$
 #7. $z \rightarrow x$ $\because 6 + 2 = 8 < 9 \therefore d[x] = 8$



Order of adding to S: s, y, z, x