Module 2 Fundamentals of Algorithm Analysis

1. Recursion vs Iteration

Concepts:

Recursion: A process where a function calls itself directly or indirectly.

Example: Calculating factorial using recursion.

Analogy: Think of opening nested Russian dolls. Each doll represents a function call until you reach the smallest doll (base case).

Iteration: Uses repetition structures like loops to execute a set of instructions repeatedly.

Example: Calculating factorial using iteration.

Analogy: Imagine climbing stairs one step at a time until you reach the top (end condition).

C++ Examples:

```
// Recursive Factorial
int factorialRecursive(int n) {
   if (n == 0) return 1; // Base case
   return n * factorialRecursive(n - 1);
}

// Iterative Factorial
int factorialIterative(int n) {
   int result = 1;
   for (int i = 1; i <= n; ++i) {
     result *= i;
   }
   return result;
}</pre>
```

Types of Recursion

```
Direct Recursion:
```

```
Function calls itself directly
     void directRecursive() {
  // Some Code
  directRecursive();
  // Some Code
Indirect Recursion:
Function fun calls another function fun2, which in turn calls fun.
void fun2(); // Forward declaration
void fun() {
  // Some Code
  fun2();
  // Some Code
}
void fun2() {
  // Some Code
  fun();
  // Some Code
Tail Recursion:
The recursive call is the last operation in the function
void tailRecursive(int n) {
  if (n == 0) return;
  std::cout << n << " ";
  tailRecursive(n - 1); // Last operation
}
```

Non-Tail Recursion:

There are operations after the recursive call.

```
void nonTailRecursive(int n) {
  if (n == 0) return;
  nonTailRecursive(n - 1);
  std::cout << n << " "; // Operation after recursive call
}</pre>
```

Analysis Framework

Efficiency Metrics:

Time Complexity: Indicates how fast an algorithm runs. Space Complexity: Amount of memory required by the algorithm.

Example:

Consider finding the sum of all elements in an array.

```
// Time Complexity: O(n)
// Space Complexity: O(1)
int sumArray(int arr[], int n) {
   int sum = 0;
   for (int i = 0; i < n; ++i) {
      sum += arr[i];
   }
   return sum;
}</pre>
```

Analogy: Think of a library cataloging books. The time it takes depends on the number of books (input size), and the space needed might include a small notepad to keep track..

Asymptotic Notations

Big-O Notation (O):

Represents the upper bound of an algorithm's running time.

Example: If an algorithm has a complexity of O(n²), it means the running time grows quadratically with input size.

Big-Omega Notation (Ω):

Represents the lower bound of an algorithm's running time.

Example: $\Omega(n)$ implies that the algorithm will take at least linear time.

Theta Notation (Θ):

Represents both upper and lower bounds.

Example: $\Theta(n \log n)$ indicates that the algorithm's running time is tightly bound by n log n.

```
// O(1) - Constant Time
int getFirstElement(int arr[]) {
  return arr[0];
}
// O(n) - Linear Time
int findElement(int arr[], int n, int key) {
  for (int i = 0; i < n; ++i) {
     if (arr[i] == key) return i;
  return -1;
}
// O(n^2) - Quadratic Time
void printPairs(int arr[], int n) {
  for (int i = 0; i < n; ++i) {
     for (int j = i + 1; j < n; ++j) {
        std::cout << "(" << arr[i] << ", " << arr[j] << ")" << std::endl;
     }
  }
```

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Best, Worst, and Average Case Analysis Sequential Search Example:

```
int sequentialSearch(int arr[], int n, int key) {
   for (int i = 0; i < n; ++i) {
      if (arr[i] == key) return i;
   }
   return -1;
}</pre>
```

Best Case: O(1) – The key is found at the first position.

Worst Case: O(n) – The key is either at the last position or not present.

Average Case: O(n) – Assuming uniform distribution of the key's position.

Analogy: Searching for a book in a disorganized shelf:

Best Case: The book is right at the front.

Worst Case: The book is at the back or missing.

Average Case: On average, you might need to check halfway through the shelf.