

## Dynamic Programming

Dynamic Programming is an optimization over plain recursion.

It is a programming technique used to solve the problems comprises -

- (a) Overlapping sub-problems.
- (b) Optimal sub-structure.

**Overlapping sub-problems** → Jb same sub-problem is getting repeat again and again.

**Optimal sub-structure** → Jb ek badi problem ka optimal solution depend kr rha hai same type ki chotti problem ke optimal solution pr.

\* DP is like agr ek problem solve kr di ek baar, then fir se solve nhi krenge kyoki uske corresponding answer store kr rkha hai.

### \* Approaching a DP problem -

- Top-Down Approach (Recursion + memoization)
- Bottom-Up Approach (Tabulation method)
- Pattern finding for space optimisation.



Which is better  $\rightarrow$  Top-down or Bottom-up?

Top-down recursive approach hai.

Bottom-up iterative approach hai.

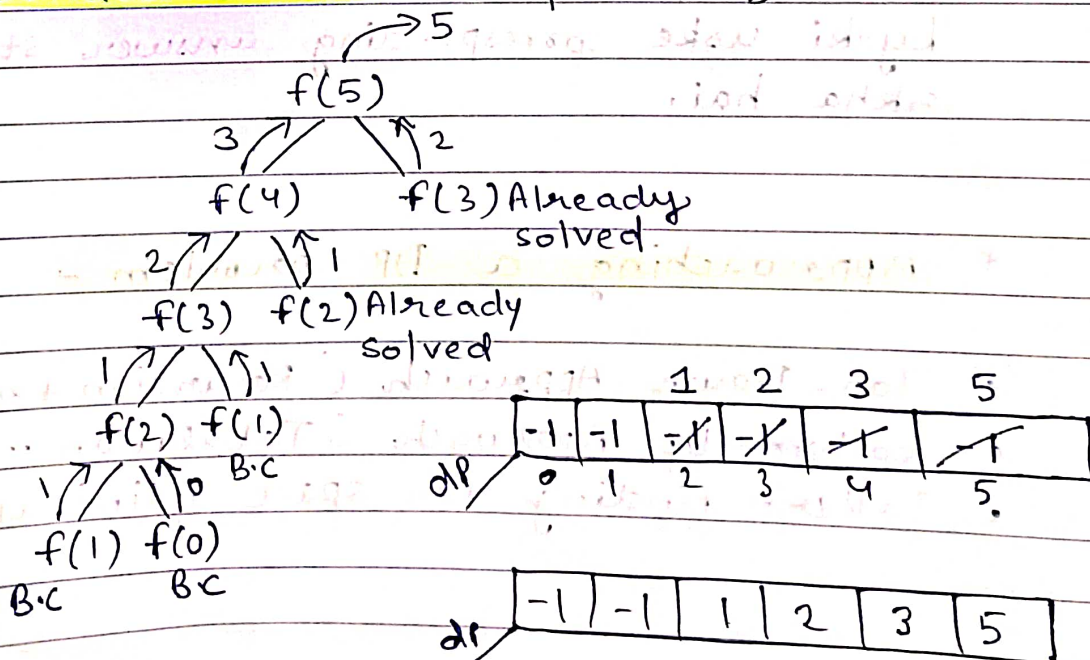
Top-down me function calls ka overhead aata hai whereas Bottom-up me nhi.

Possibility hai koi solution bottom-up approach se chale top-down se nhi kyoki bottom-up me function calls ka overhead nhi aata.

### \* Top-Down Approach (Recursion + memoization)

- Step 1 - Create a DP array.
- Step 2 - Store/return answer in DP array.
- Step 3 - If answer already exists in DP array, then return it.

### \* Fibonacci number (Top-Down)





Code -

```
class Solution {
public:
    // top-down approach
    int solve(int n, vector<int>&dp) {
        if (n == 0 || n == 1) return n;
        // step 3: if ans already exists, return it
        if (dp[n] != -1) return dp[n];
        // step 2: store/return using DP array
        dp[n] = solve(n-1, dp) + solve(n-2, dp);
        return dp[n];
    }

    int fib(int n) {
        // step 1: create a DP array
        vector<int> dp(n+1, -1);
        int ans = solve(n, dp);
        return ans;
    }
};
```

T.C  $\rightarrow O(n)$ . The function is called each time for  $(0 \rightarrow n)$  numbers.

S.C  $\rightarrow O(n) \rightarrow$  due to DP array  
 $O(n) \rightarrow$  due to recursion stack  
Overall,  $O(n)$ .

## \* Bottom-up Approach (Tabulation method)

- Step 1- Create DP array.  
Step 2- Analyse base case and fill DP array accordingly. This will maintain an initial state in DP array.  
Step 3- Fill remaining DP array using logic/formula of recursive relation.

Fibonacci series - 0, 1, 1, 2, 3, 5, 8, 13, 21, ---

0	1		1	2	3	5
0	1		2	3	4	5

Fill initial array using base case      Fill remaining array using recursive relation

$$dp[2] = dp[1] + dp[0]$$

$$dp[3] = dp[2] + dp[1]$$

$$dp[4] = dp[3] + dp[2]$$

$$dp[5] = dp[4] + dp[3]$$



Code -

```
class Solution {
public:
    // bottom-up approach

    int fib(int n) {
        // create DP array
        vector<int> dp(n+1, -1);

        // analyse base case to maintain initial state
        // of DP array
        dp[0] = 0; if (n == 0) return 0;
        dp[1] = 1;

        // array size = n+1
        // index 0 and 1 fill ho chuke hai
        // remaining array to fill, index (2 → n)

        // fill remaining DP array using loop
        for (int i = 2; i <= n; i++) {
            dp[i] = dp[i-1] + dp[i-2];
        }

        // return ans
        return dp[n];
    }
};
```

T.C →  $O(n)$

S.C →  $O(n)$  → due to DP array

**Note** - Space optimisation krte time shifting krna bhul jate hai, vo yaad rkhnna.

## Fibonacci number (Space optimisation)

Fibonacci series - 0, 1, 1, 2, 3, 5, 8, 13, 21 - - -

$$f(n) = f(n-1) + f(n-2)$$

The  $n^{\text{th}}$  fibonacci number is the sum of  $(n-1)^{\text{th}}$  and  $(n-2)^{\text{th}}$  fibonacci number. Using this pattern, we can optimise the code space-wise.

**Code -**

```
class Solution {  
public:
```

```
    int fib(int n) {
```

```
        int prev = 0;
```

```
        if (n == 0) return 0;
```

```
        int curr = 1;
```

```
        if (n == 1) return 1;
```

```
        int ans = 0;
```

```
        for (int i = 2; i <= n; i++) {
```

```
            ans = prev + curr;
```

```
            prev = curr;
```

```
            curr = ans;
```

```
        }
```

```
        return ans;
```

```
    }  
};
```



\* 1-D, 2-D and 3-D DP  $\rightarrow$

1-D = Recursive function has only one varying argument.

2-D = Recursive function has two varying arguments.

3-D = Recursive function has three varying or non-constant arguments.

Note - (1) 99% cases mein memoization wale code mein jo argument change ho rha hota hai in 1-D dp. vhi return kr rhe hote hai.

(2) Agr memoization wala code TLE de rha h to 99% cases mein argument ko by reference pass krne se chl jata hai code.

Note - hum top-down and bottom-up approach me generally opposite direction me ja rhe hote hai.

For instance, if ~~kisi~~ top-down approach se kisi code mein  $0 \rightarrow n$  move kr rhe hai to bottom-up mein  $n \rightarrow 0$  move kr rhe honge.