

Function Approximation with Numerical Redundancy

Astrid Herremans
joint work with Daan Huybrechs

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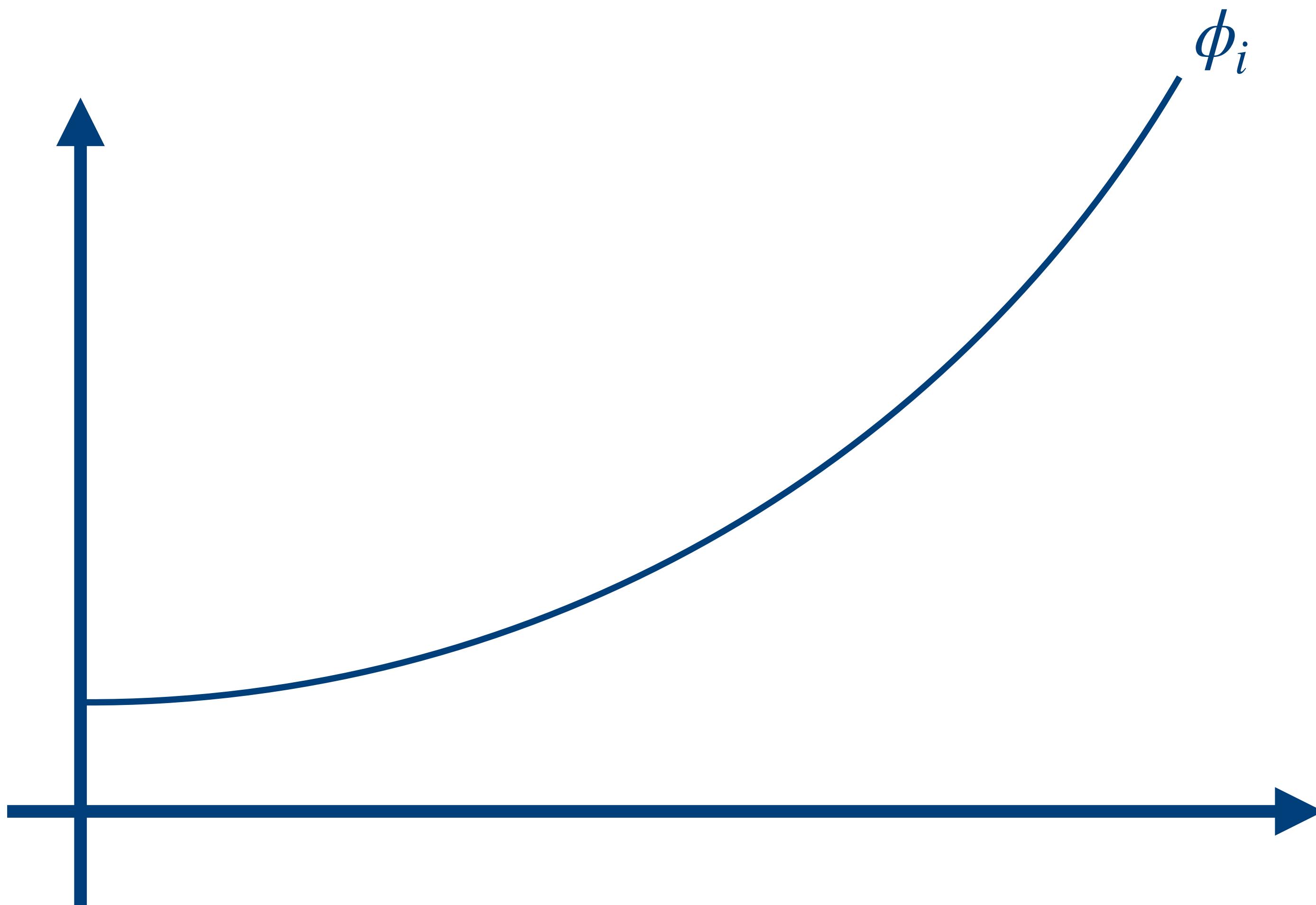
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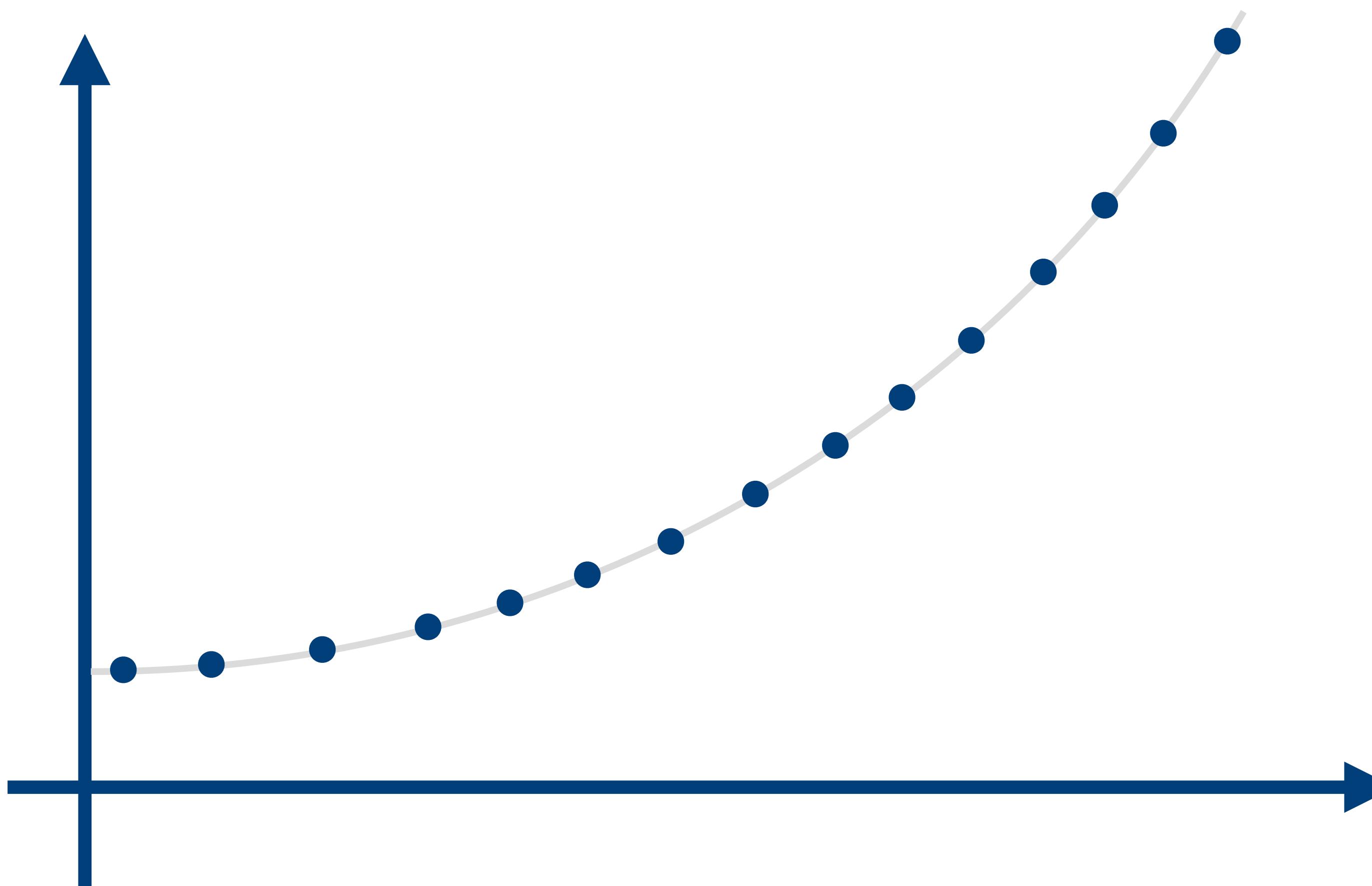
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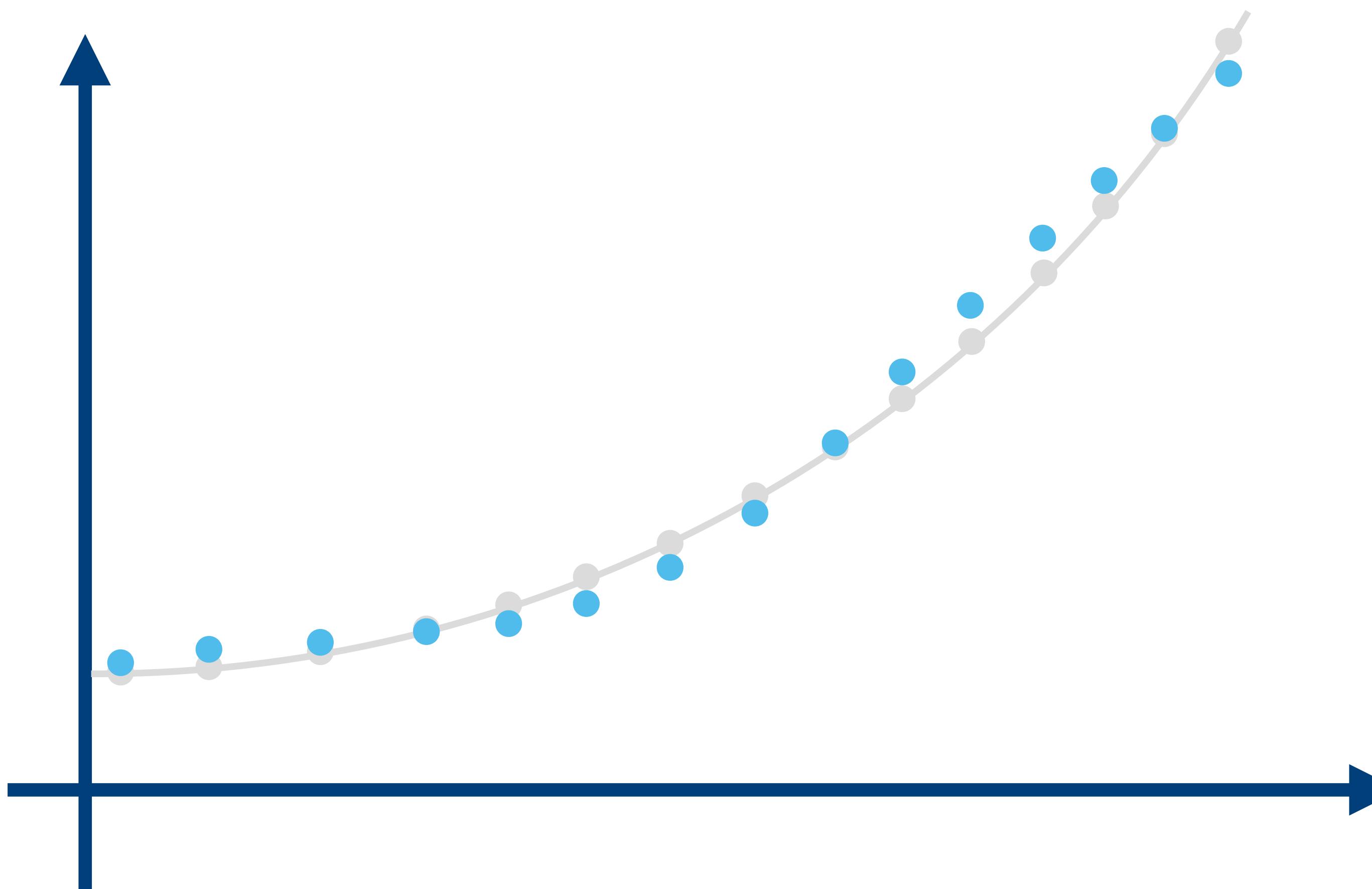
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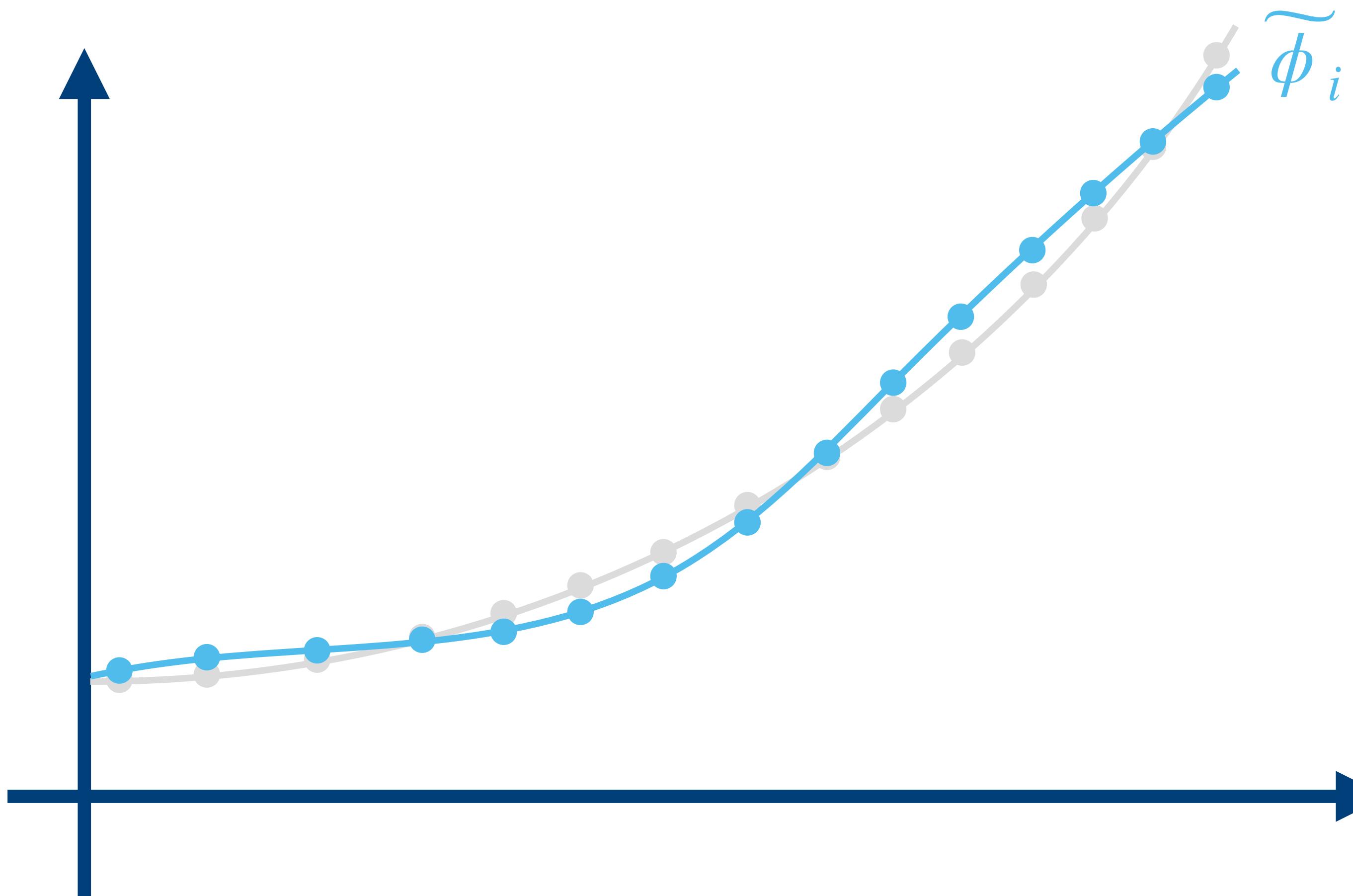
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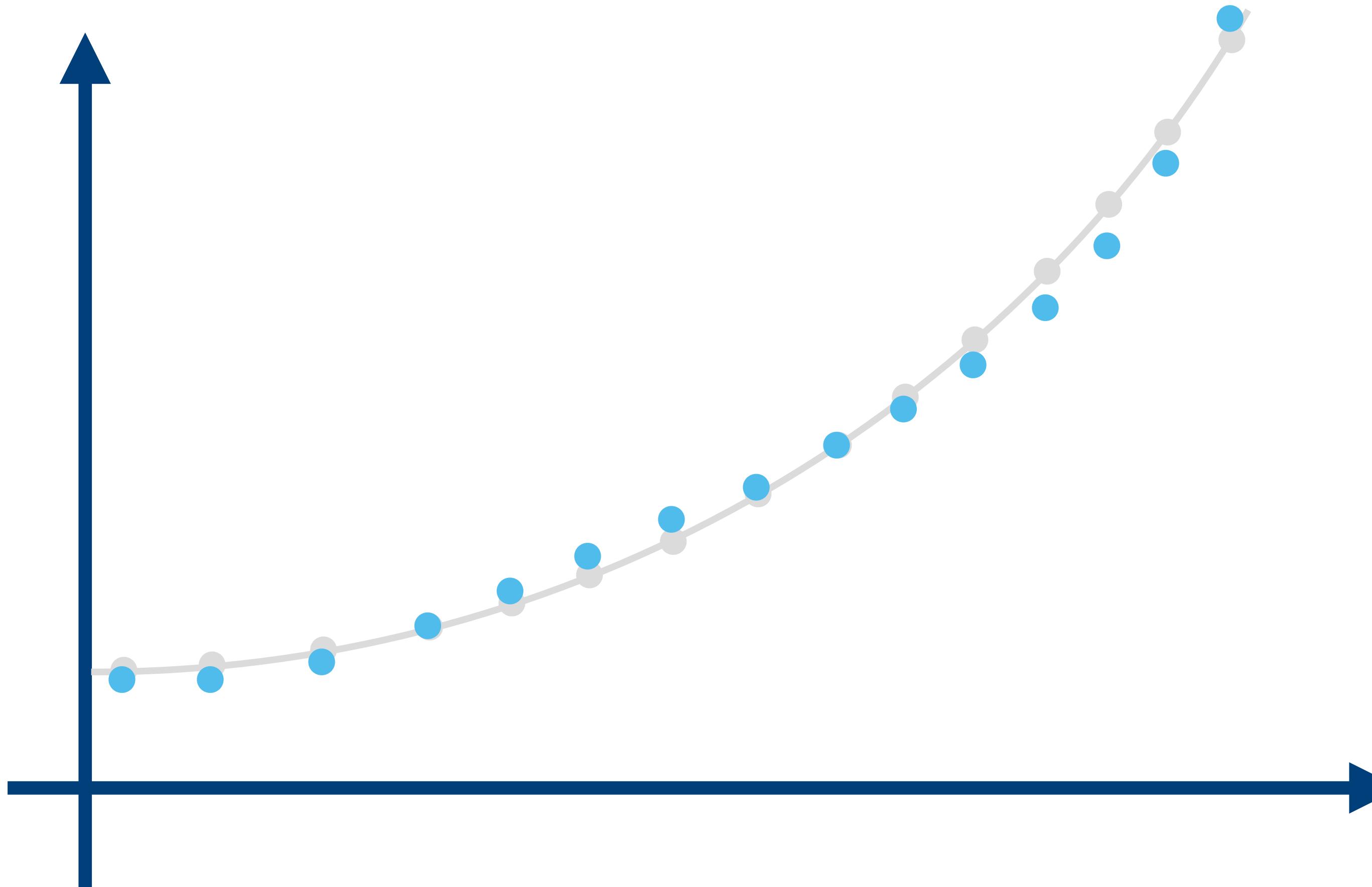
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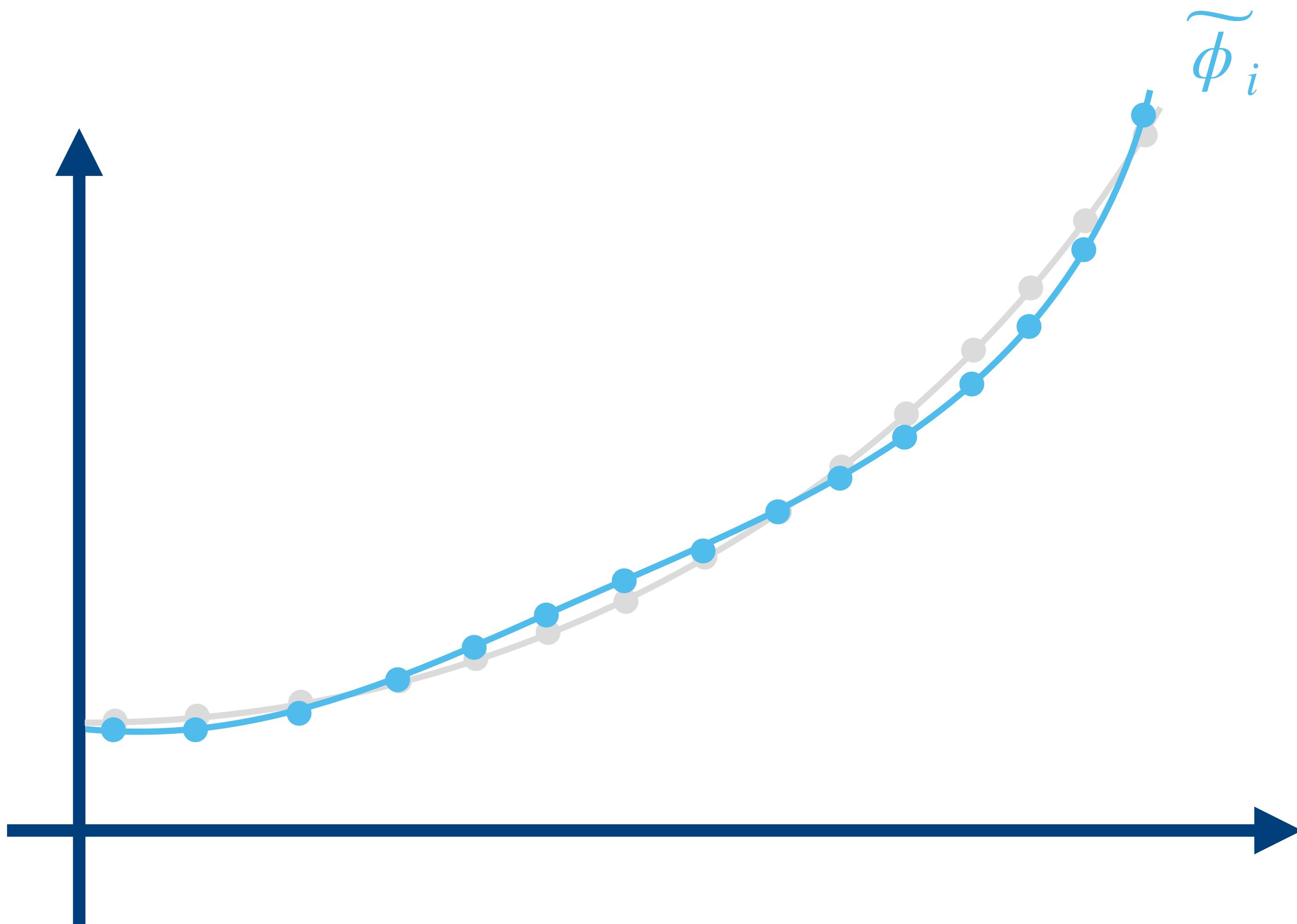


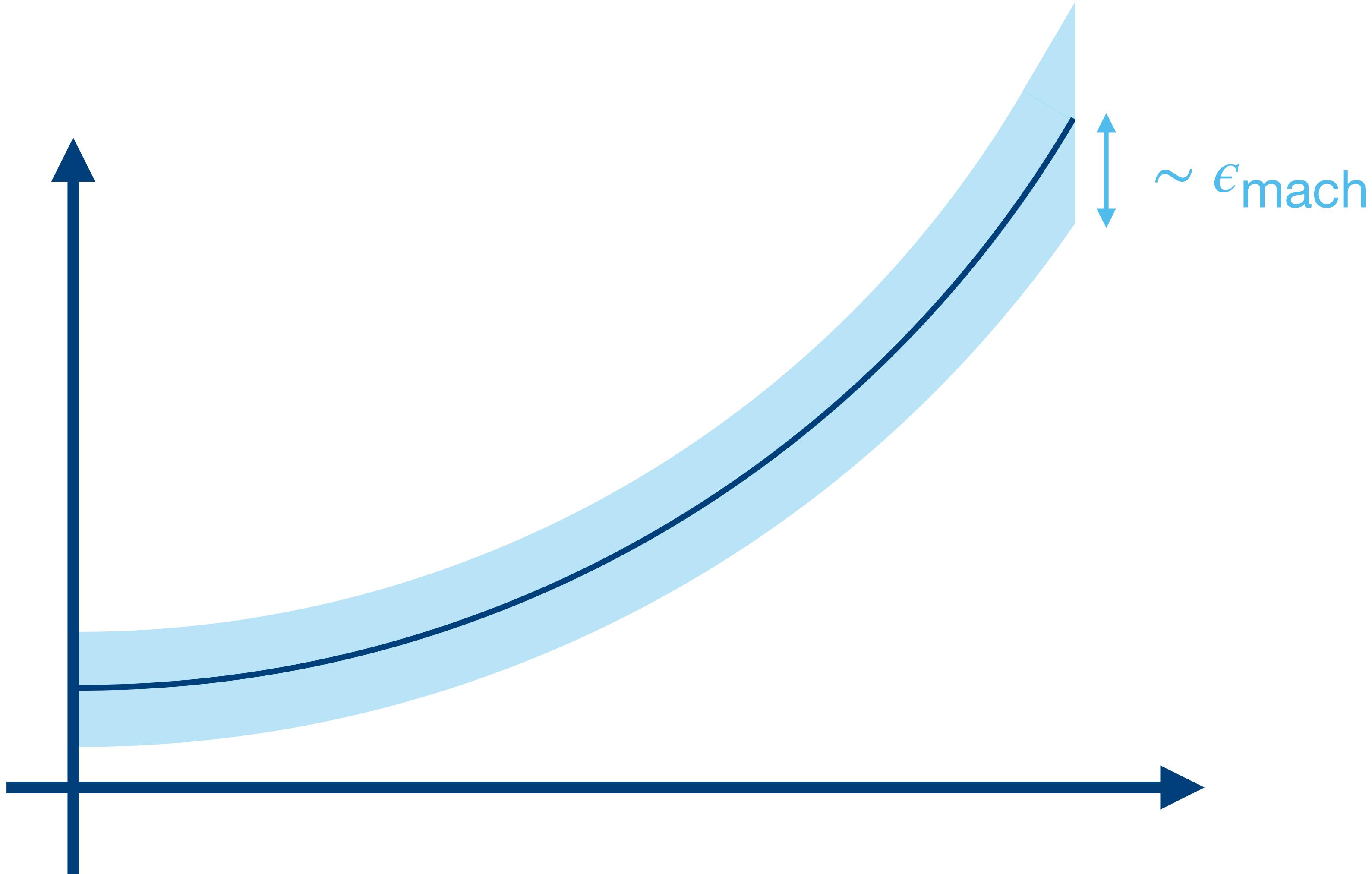


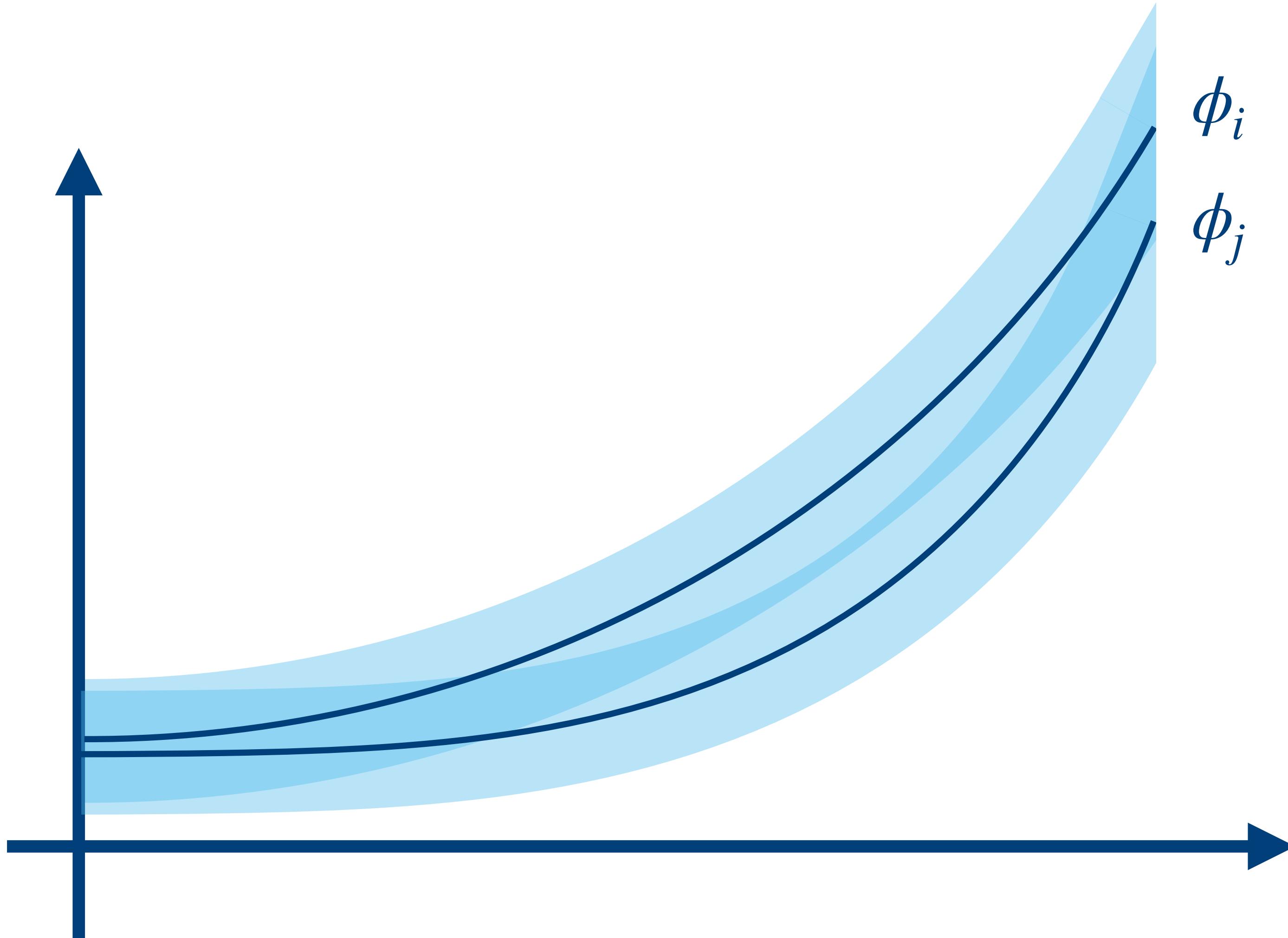


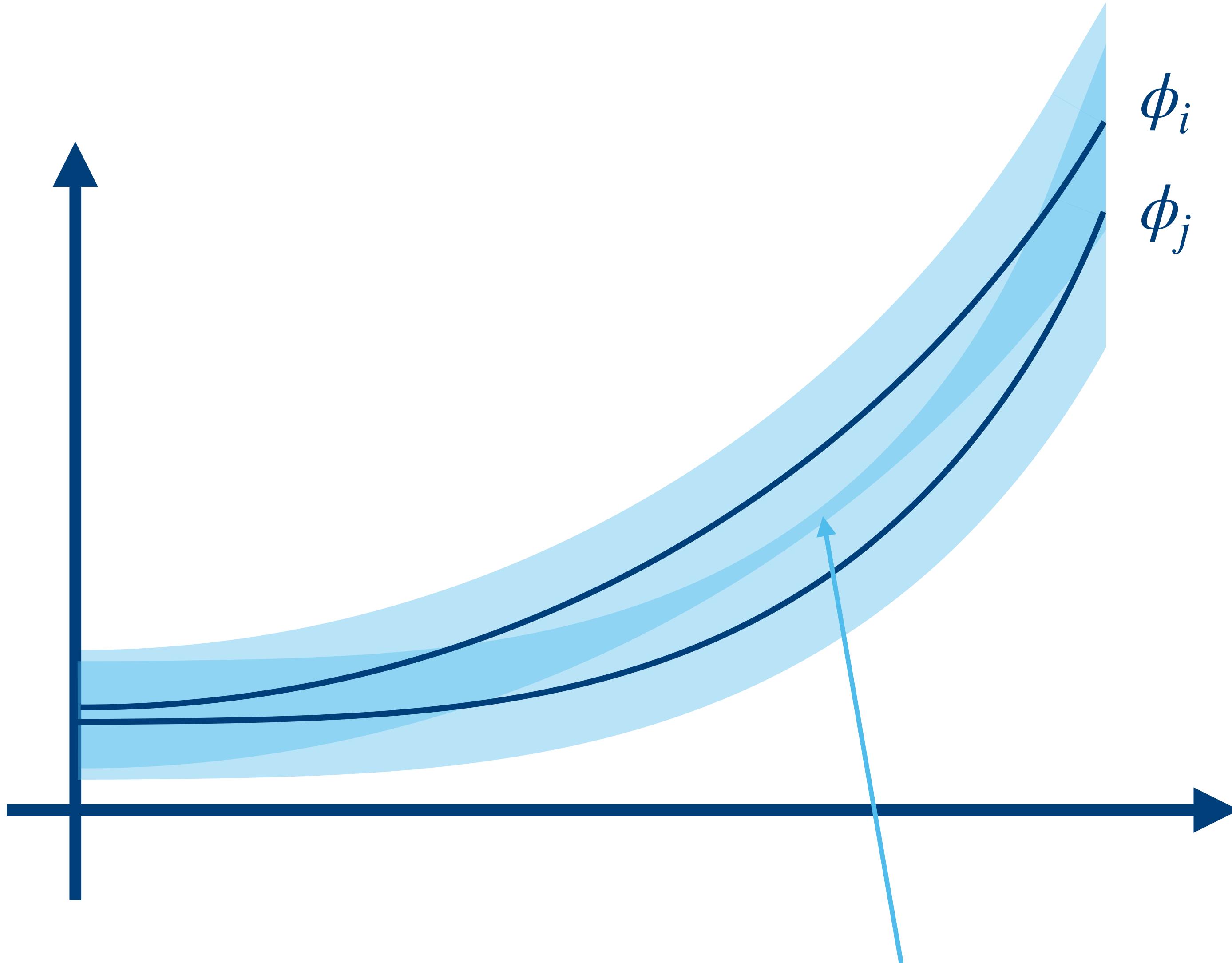












ϕ_i and ϕ_j are indistinguishable from a numerical point of view

Numerically redundant sets

span a lower dimensional space when analysed [numerically](#) rather than analytically

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This is equivalent to: the singular values of the synthesis operator

$$\mathcal{T}_n : \mathbb{C}^n \rightarrow H, \quad c \mapsto \sum_{i=1}^n c_i \phi_i$$

satisfy $\sigma_{\min} \leq \epsilon_{mach} \sigma_{\max}$

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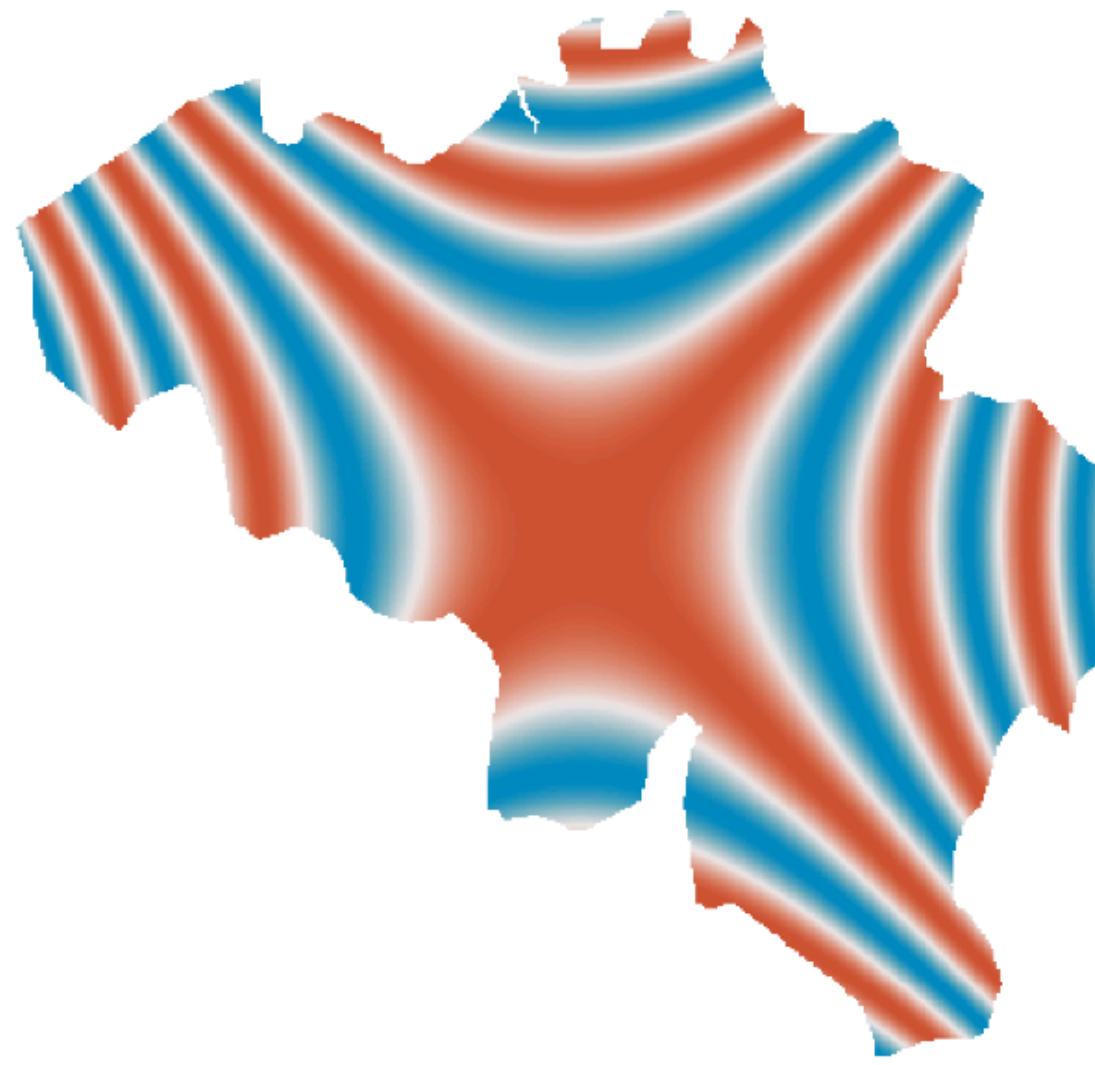
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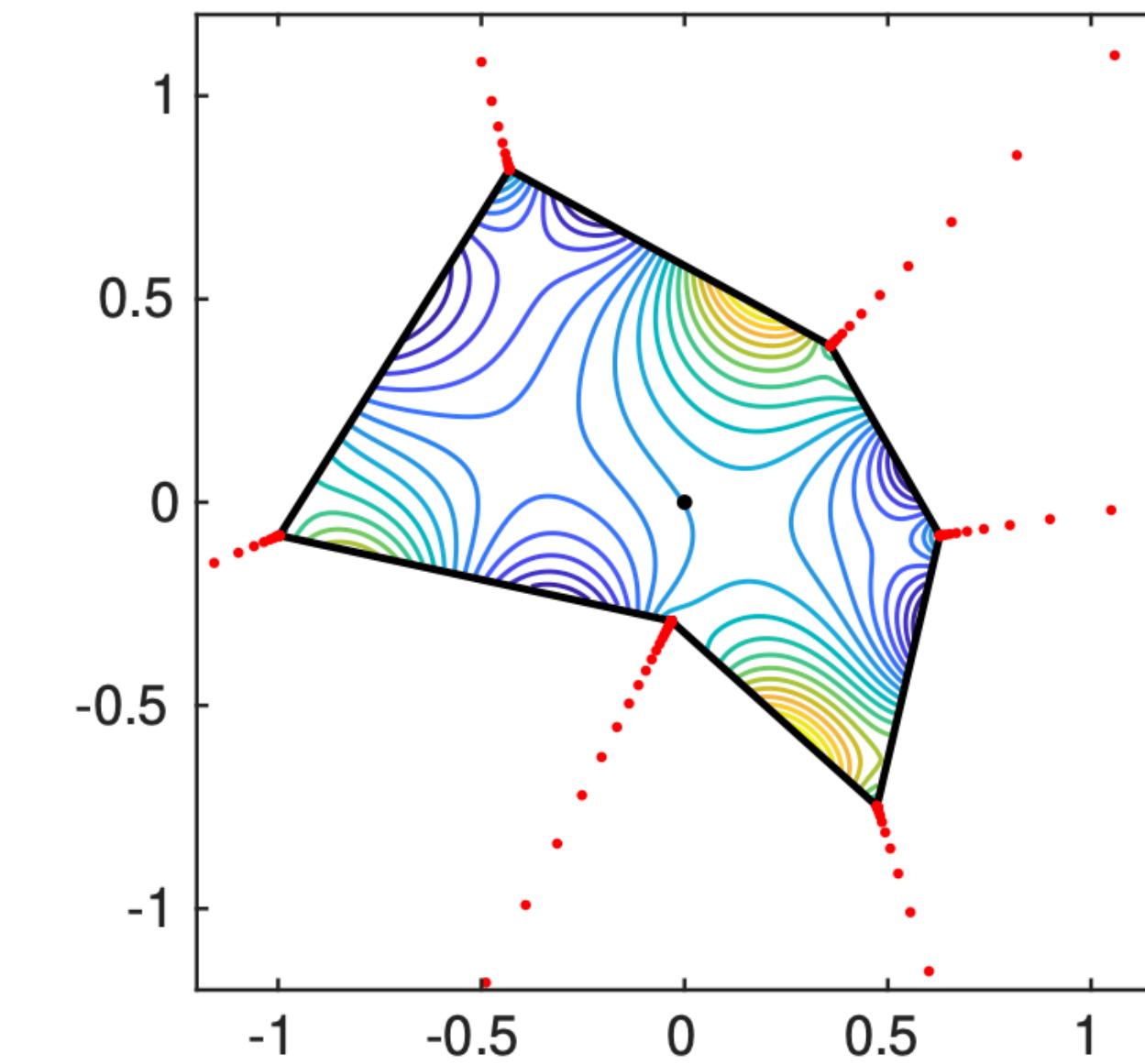
You recognise this if: you have plenty of data on f yet the system of equations to compute coefficients c is ill-conditioned anyway

Numerically redundant sets

offer a lot of flexibility



approximate on irregular domains

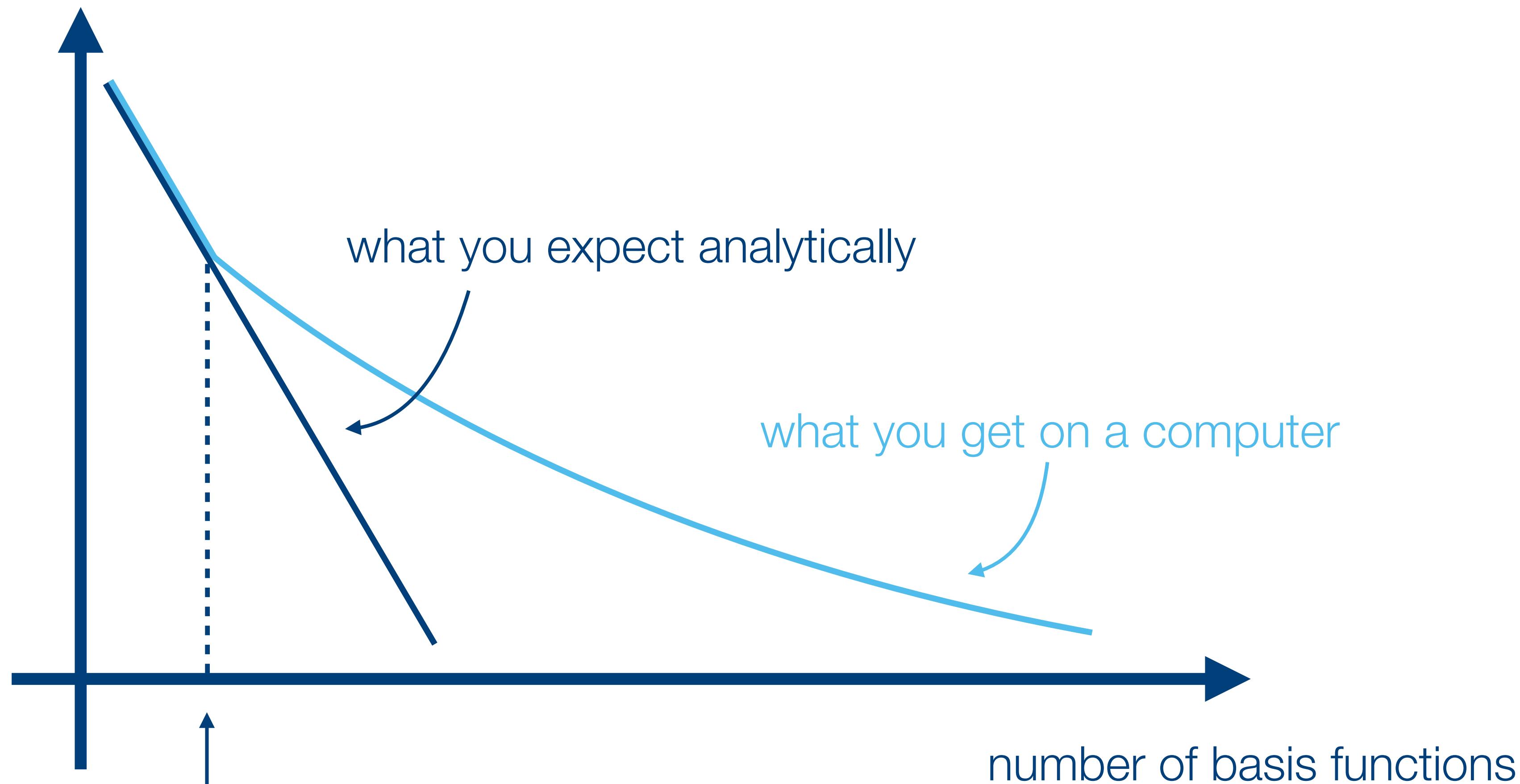


incorporate knowledge on f
by combining / weighting bases

- ▶ **The bad news** - slower convergence
- ▶ **The ugly news** - regularization
- ▶ **The good news** - less data

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approximation error



the basis becomes
numerically redundant

number of basis functions

Achievable accuracy

On a computer, the basis functions ϕ_i are perturbed to $\tilde{\phi}_i$

$$\text{best approximation error in } \text{span}(\phi_i) = \inf_{c \in \mathbb{C}^n} (\|f - \mathcal{T}_n c\|)$$

$$\text{best approximation error in } \text{span}(\tilde{\phi}_i) \leq \inf_{c \in \mathbb{C}^n} (\|f - \mathcal{T}_n c\| + \epsilon_{\text{mach}} \|\mathcal{T}_n\| \|c\|_2)$$

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- ▶ the numerical accuracy depends on the norm of the coefficients $\|c\|_2$

Achievable accuracy

On a computer, the basis functions ϕ_i are perturbed to $\tilde{\phi}_i$

best approximation error in $\text{span}(\phi_i)$

$$\inf_{c \in \mathbb{C}^n} (\|f - \mathcal{T}_n c\|)$$

best approximation error in $\text{span}(\tilde{\phi}_i)$

$$\leq \inf_{c \in \mathbb{C}^n} (\|f - \mathcal{T}_n c\| + \epsilon_{\text{mach}} \|\mathcal{T}_n\| \|c\|_2)$$

- ▶ the numerical accuracy depends on the norm of the coefficients $\|c\|_2$
- ▶ the difference is only significant if \mathcal{T}_n has small singular values

Convergence guarantees

Assume $\{\phi_i\}_{i=1}^n \subset \{\phi_i\}_{i=1}^\infty$ and $f \in \overline{\text{span}}(\{\phi_i\}_{i=1}^\infty)$, then $f = \sum_{i=1}^\infty a_i \phi_i$ and

orthonormal basis

a is unique

$$\|a\|_2 = \|f\|$$

Riesz basis

a is unique

$$A\|a\|_2^2 \leq \|f\|^2 \leq B\|a\|_2^2$$

(overcomplete) frame

a is not unique

$$\exists a : A\|a\|_2^2 \leq \|f\| \leq B\|a\|_2^2$$

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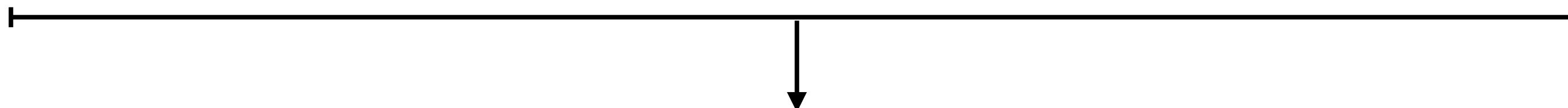
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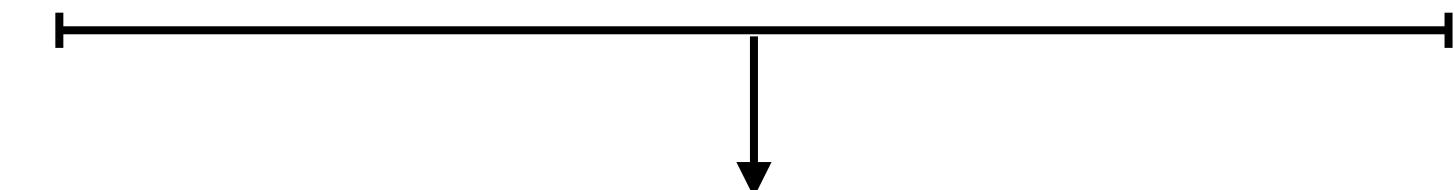
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subsequence is again an
orthonormal / Riesz basis



subsequence is
numerically redundant as $n \rightarrow \infty$

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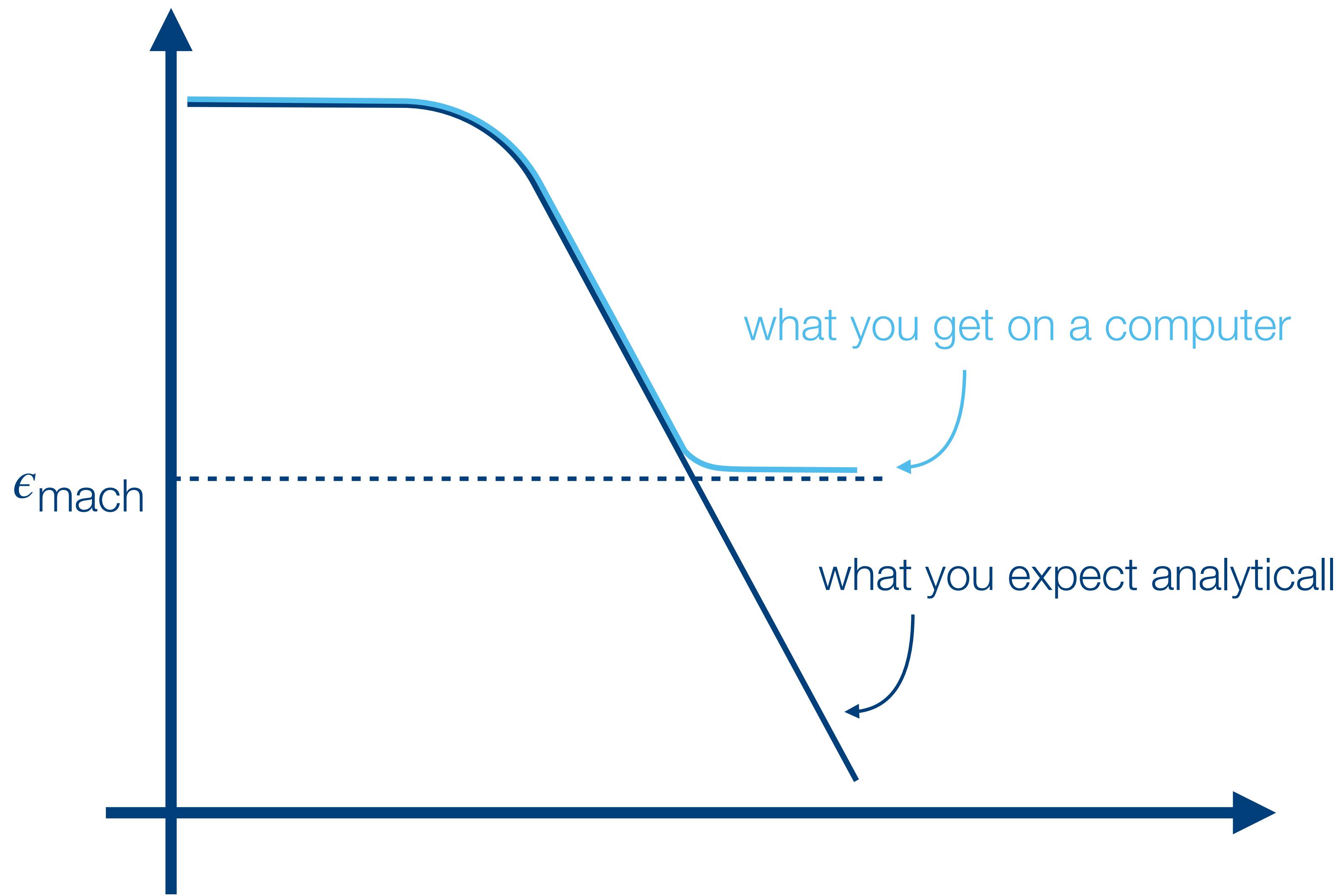
$$\exists a : A\|a\|_2^2 \leq \|f\|^2 \leq B\|a\|_2^2$$

The existence of bounded coefficients $\{a_i\}_{i=1}^\infty$ guarantees convergence to ϵ_{mach}

$$\lim_{n \rightarrow \infty} \left(\inf_{c \in \mathbb{C}^n} (\|f - \mathcal{T}_n c\| + \epsilon_{\text{mach}} \|\mathcal{T}_n\| \|c\|_2) \right) \leq \epsilon_{\text{mach}} \sqrt{\frac{B}{A}} \|f\|$$

- ▶ **The bad news** - slower convergence
- ▶ **The ugly news** - regularization
- ▶ **The good news** - less data

singular values of the
system of equations



Backward stability

We look for coefficients c that minimize

$$\|Ac - b\|_2 \quad \text{with } (A)_{i,j} = l_i(\phi_j) \text{ and } (b)_i = l_i(f)$$

where $\{l_i\}_{i=1}^m$ are linear sampling functionals

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Numerical algorithms guarantee to compute

$$\hat{c} = \arg \min_c \| (A + \Delta A)c - (b + \Delta b) \|_2 \quad \text{where } \|\Delta \cdot\|_2 \lesssim \epsilon_{\text{mach}} \|\cdot\|_2$$

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such that

$$\|A\hat{c} - b\|_2 \lesssim \inf_c \|Ac - b\|_2 + \epsilon_{\text{mach}} (\|A\|_2 (\|\hat{c}\|_2 + \|c\|_2) + \|b\|_2)$$

Backward stability does not suffice

We look for coefficients c that minimize

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where $\{l_i\}_{i=1}^m$ are linear sampling functionals

Numerical algorithms

$$\hat{c} = \arg \min_x$$

For numerically redundant sets, A is

heavily ill-conditioned and $\|\hat{c}\|_2$ can be huge!

such that

$$\|A\hat{c} - b\|_2 \lesssim \inf_c \|Ac - b\|_2 + \epsilon_{\text{mach}} (\|A\|_2 (\|\hat{c}\|_2 + \|c\|_2) + \|b\|_2)$$



ℓ^2 -regularization

If we penalize the norm of the coefficients

$$\min_c \|Ac - b\|_2^2 + \epsilon^2 \|c\|_2^2 \quad \text{where } \epsilon \sim \epsilon_{\text{mach}} \|A\|_2$$

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Remember: the numerically achievable accuracy equals

$$\inf_c (\|\mathcal{T}_n c - f\| + \epsilon_{\text{mach}} \|\mathcal{T}_n\| \|c\|_2)$$

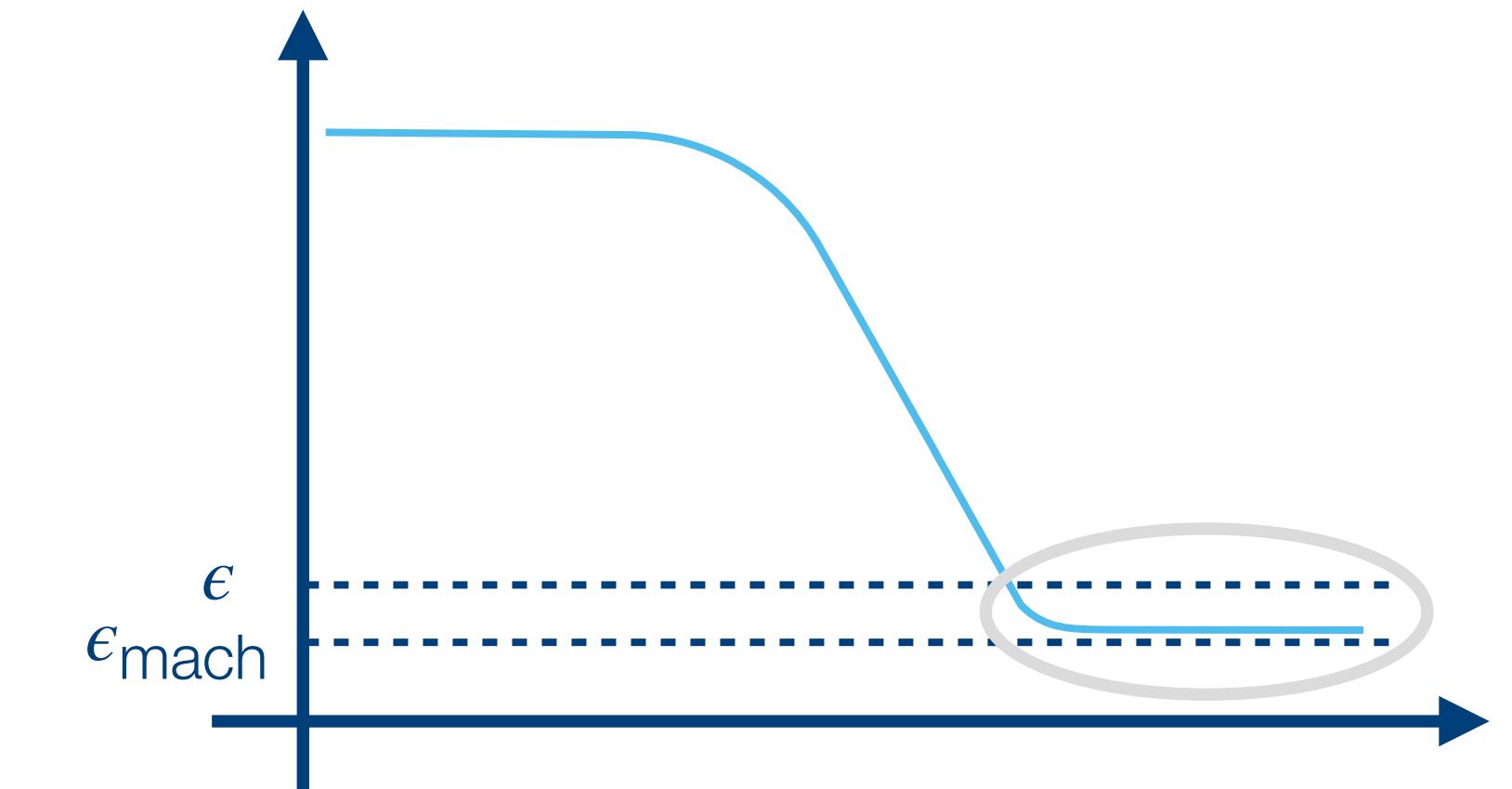
Common strategies

ℓ^2 -regularization

- ▶ Tikhonov regularization
- ▶ truncated singular value decomposition (TSVD)

! standard routines such as Matlab's backslash
regularize under the hood

singular values of the
system of equations

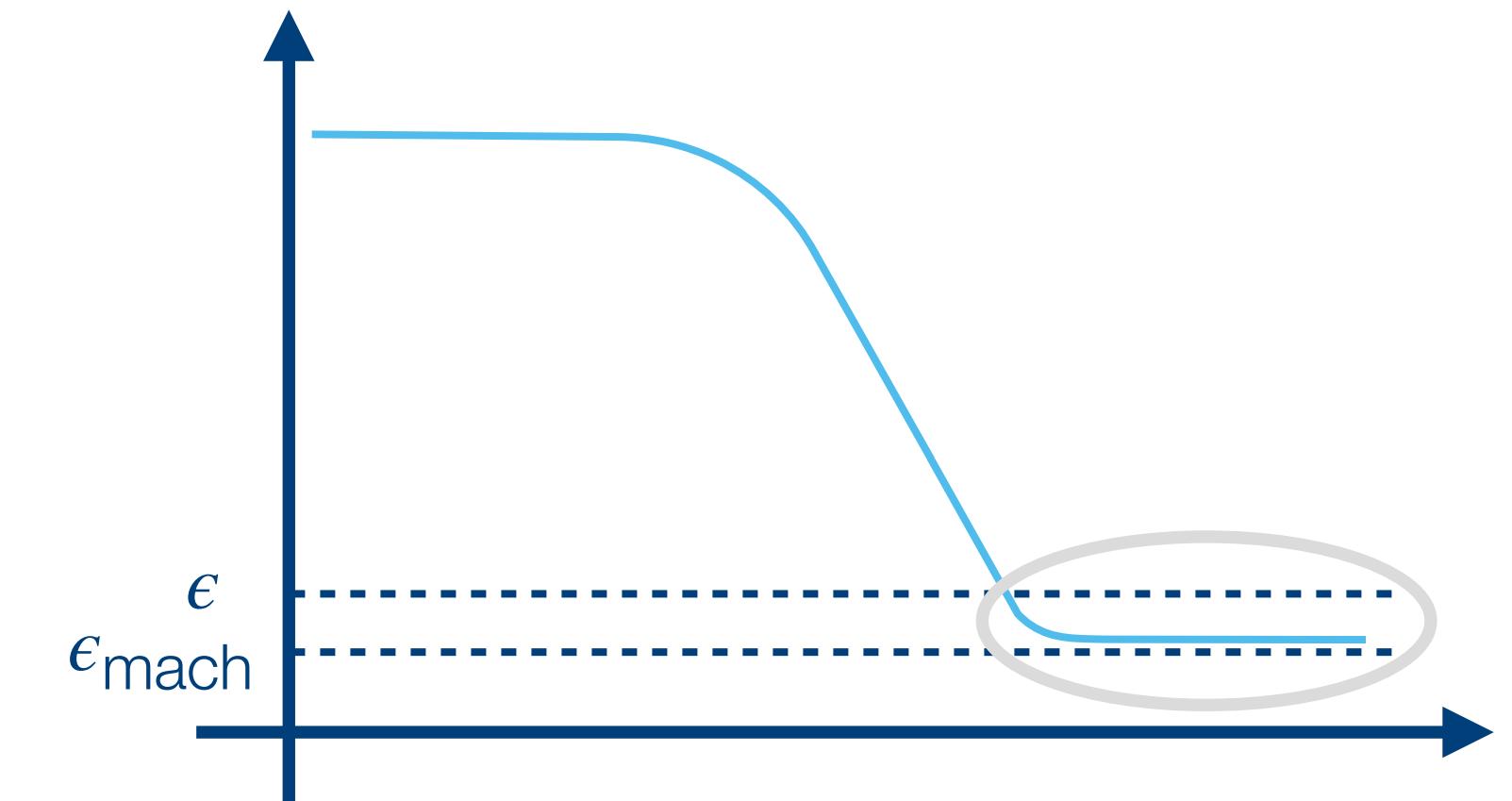


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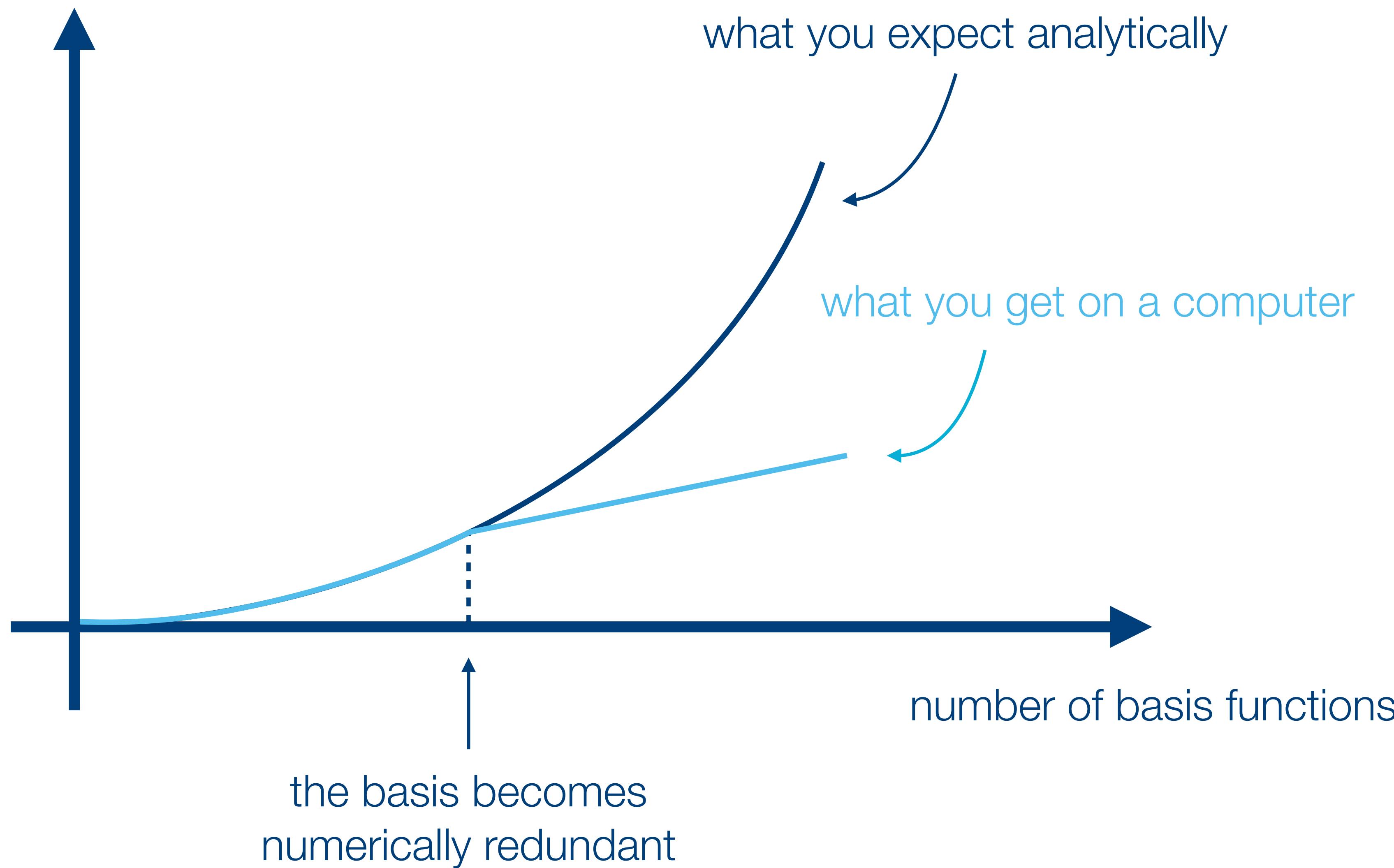
Numerical orthogonalization on a dense grid $\{t_j\}_{j=1}^m$

$$(T + \Delta T) = QR$$

where $T = \begin{bmatrix} \phi_1(t_1) & \dots & \phi_n(t_1) \\ \vdots & & \vdots \\ \phi_1(t_m) & \dots & \phi_n(t_m) \end{bmatrix}$ and $\|\Delta T\| \lesssim \epsilon_{mach} \|T\|_2$

- ▶ **The bad news** - slower convergence
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required number
of samples



Error analysis

We discretize using the sampling operator $\mathcal{M}_m : f \mapsto \{l_j(f)\}_{j=1}^m$ defining $\|\cdot\|_m = \|\mathcal{M}_m \cdot\|_2$

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Analytical behaviour

$$c = \arg \min_x \|\mathcal{T}_n x - f\|_m^2$$

then we obtain

$$\|\mathcal{T}_n c - f\| \leq \left(1 + \frac{\|\mathcal{M}_m\|}{\sqrt{A_{n,m}}} \right) \inf_x \|\mathcal{T}_n x - f\|$$

where

$$A_{n,m} \|v\|^2 \leq \|v\|_m^2, \quad \forall v \in \text{span}(\{\phi_i\}_i)$$

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$$c = \arg \min_x \|\mathcal{T}_n x - f\|_m^2 + \epsilon^2 \|x\|_2^2$$

then if

$$\|\mathcal{T}_n c - f\| \leq \left(1 + \frac{\|\mathcal{M}_m\|}{\sqrt{A_{n,m}^\epsilon}} \right) \inf_x \|\mathcal{T}_n x - f\| + \epsilon \|x\|_2$$

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Discretization condition

Analytical behaviour

$$A_{n,m} \|v\|^2 \leq \|v\|_m^2, \quad \forall v \in \text{span}(\{\phi_i\}_i)$$

$$\Leftrightarrow A_{n,m} G_n \preceq G_{n,m}$$

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$$(G_n)_{i,j} = \langle \phi_i, \phi_j \rangle \text{ and } (G_{n,m})_{i,j} = \langle \mathcal{M}_m \phi_i, \mathcal{M}_m \phi_j \rangle_2$$

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- ▶ Independent of the spanning set $\{\phi_i\}_i$: we can use an ONB for the analysis s.t. $G_n = I$

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- Dependent on the spanning set $\{\phi_i\}_i$

How do we find sampling functionals
that satisfy these norm inequalities?

Christoffel sampling

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$A_{n,m}$ close to 1 w.h.p. when using

$$m \geq Cn \log(n)$$

pointwise random samples with probability depending on

$$k_n(x) = \sum_{i=1}^n |u_i(x)|^2$$

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the inverse Christoffel function / continuous analogue of leverage scores

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$A_{n,m}^\epsilon$ close to 1 w.h.p. when using $n^\epsilon = \sum_{i=1}^n \frac{\sigma_i^2}{\sigma_i^2 + \epsilon^2}$

$$m \geq Cn^\epsilon \log(n^\epsilon)$$

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$$A_{n,m} \|v\|^2 \leq \|v\|_m^2, \quad \forall v \in \text{span}(\{\phi_i\}_i)$$

$$\Leftrightarrow A_{n,m} I \leq G_{n,m}^{ONB}$$

$A_{n,m}$ close to 1 w.h.p. when using

$$m \geq Cn \log(n)$$

pointwise random samples with probability depending on

$$k_n(x) = \sum_{i=1}^n |u_i(x)|^2$$

Numerical behaviour

$$A_{n,m}^\epsilon \|\mathcal{T}_n x\|^2 \leq \|\mathcal{T}_n x\|_m^2 + \epsilon^2 \|x\|_2^2, \quad \forall x \in \mathbb{C}^n$$

$$\Leftrightarrow A_{n,m}^\epsilon G_n \leq G_{n,m} + \epsilon^2 I$$

$A_{n,m}^\epsilon$ close to 1 w.h.p. when using

$$m \geq Cn^\epsilon \log(n^\epsilon)$$

pointwise random samples with probability depending on

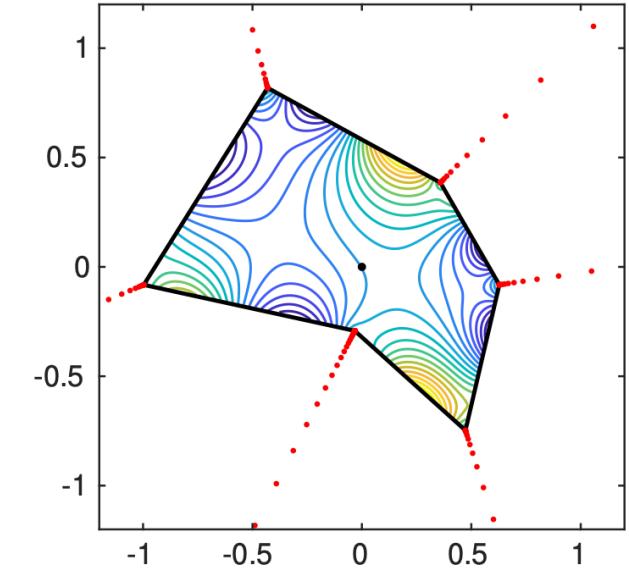
$$k_n^\epsilon(x) = \sum_{i=1}^n \frac{\sigma_i^2}{\sigma_i^2 + \epsilon^2} |u_i(x)|^2$$

continuous analogue of ridge leverage scores

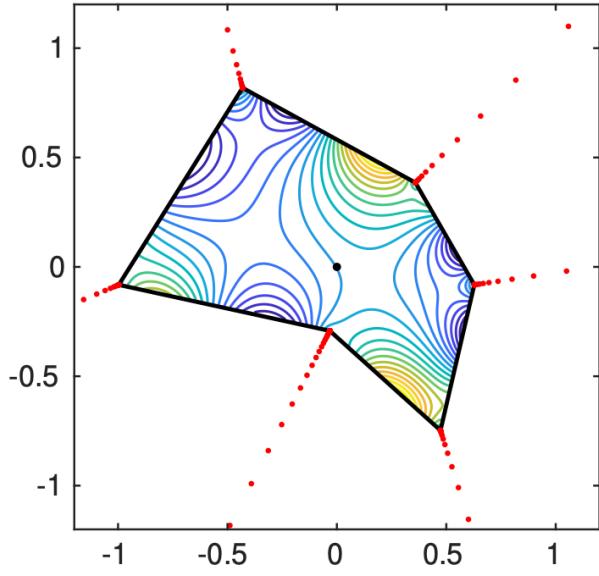
Deterministic sampling

Approximation of $f(x) = J_{1/2}(x + 1) + \frac{1}{x^2 + 1}$ on $[-1, 1]$ using the basis

$$\{p_i(x)\}_{i=1}^{40} \cup \{w(x) p_i(x)\}_{i=1}^{40} \quad \text{with } w(x) = \sqrt{x + 1}$$

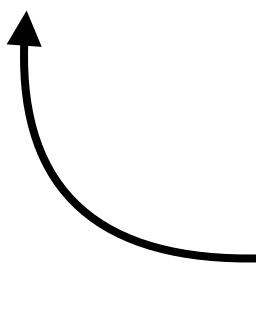


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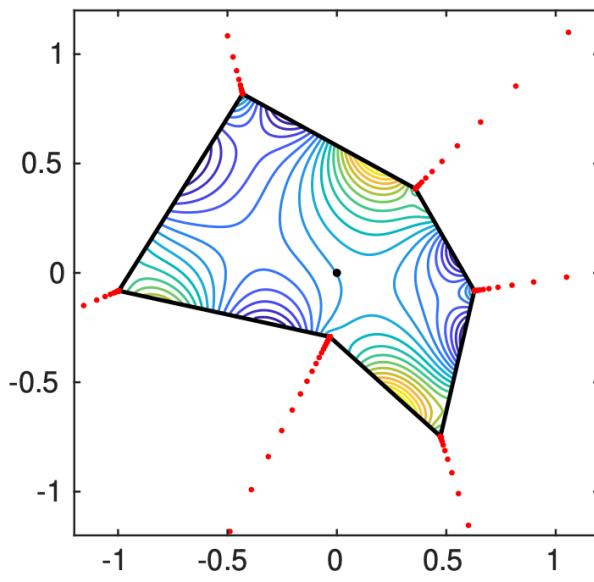
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 this is a subsequence of an overcomplete frame!

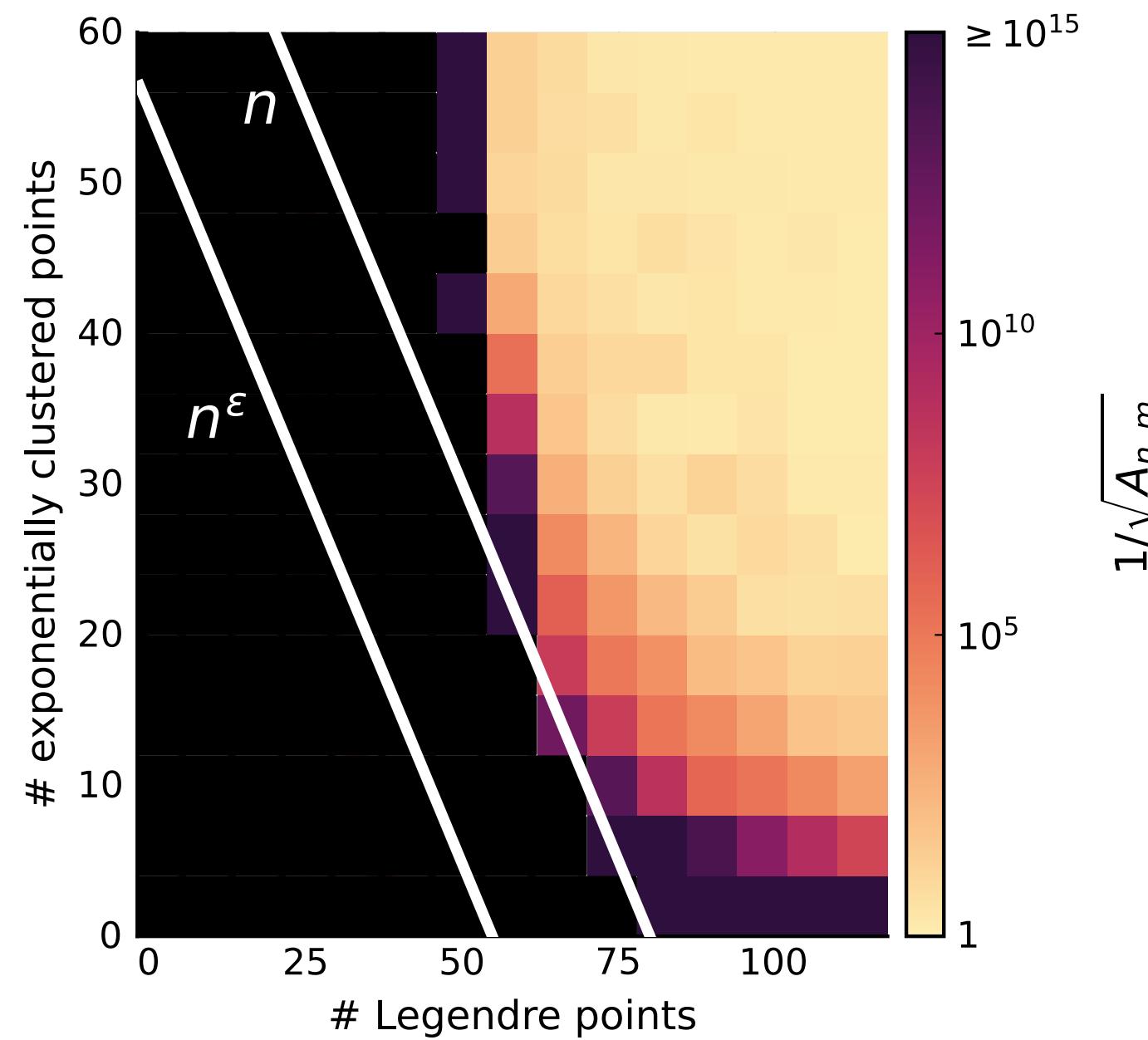
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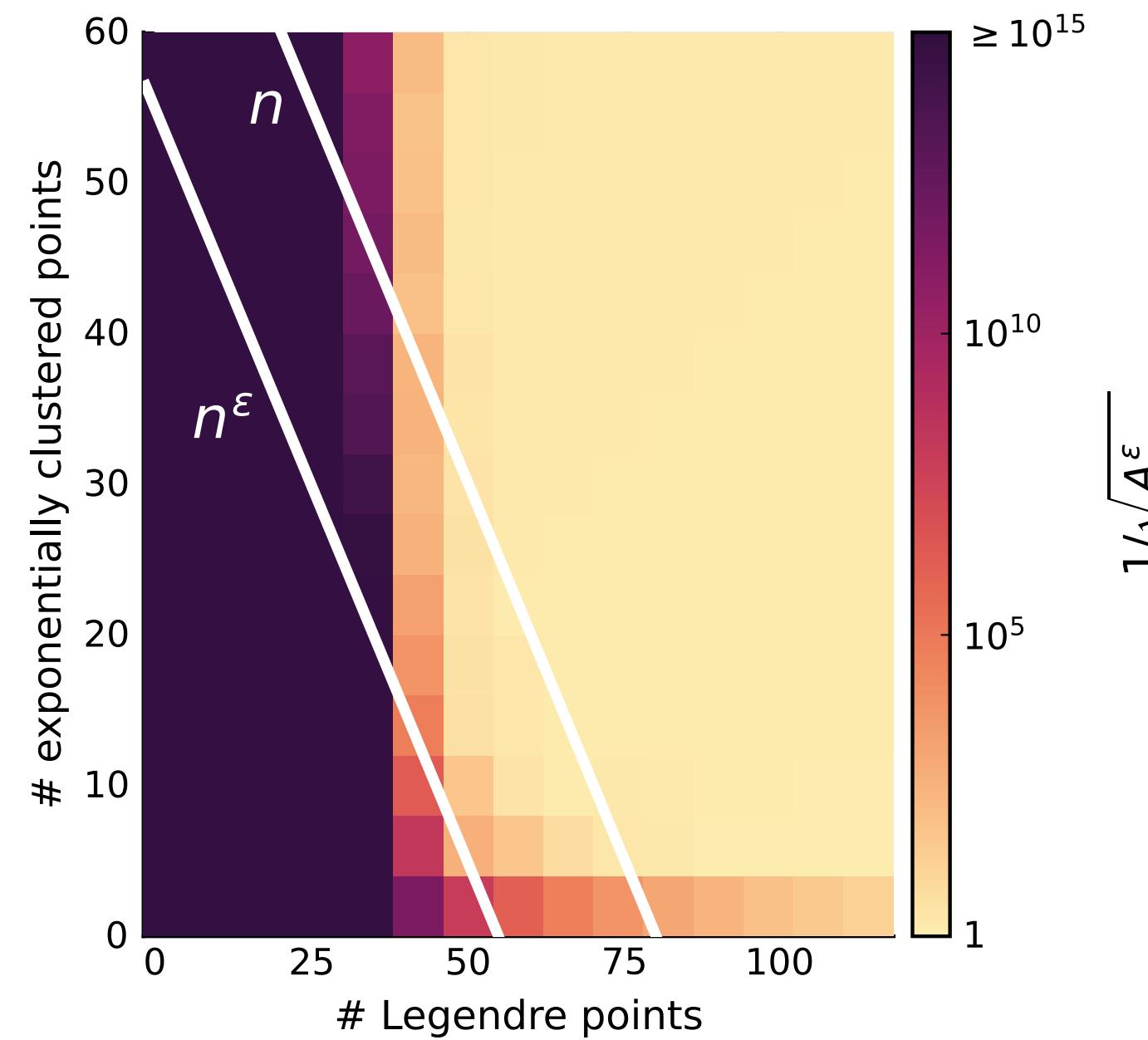
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analytical analysis



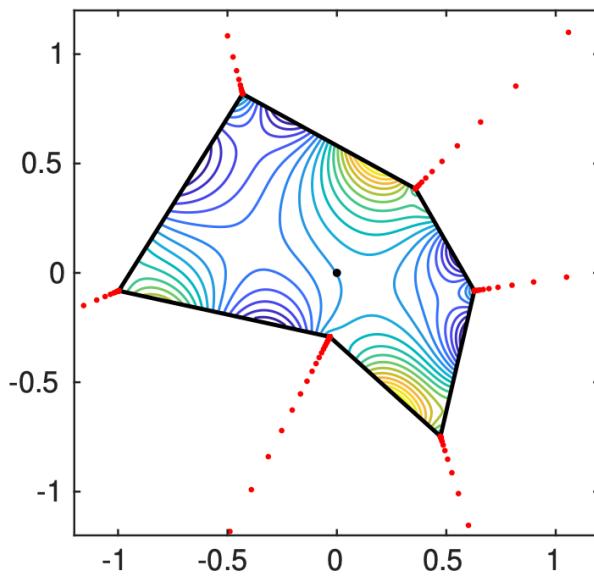
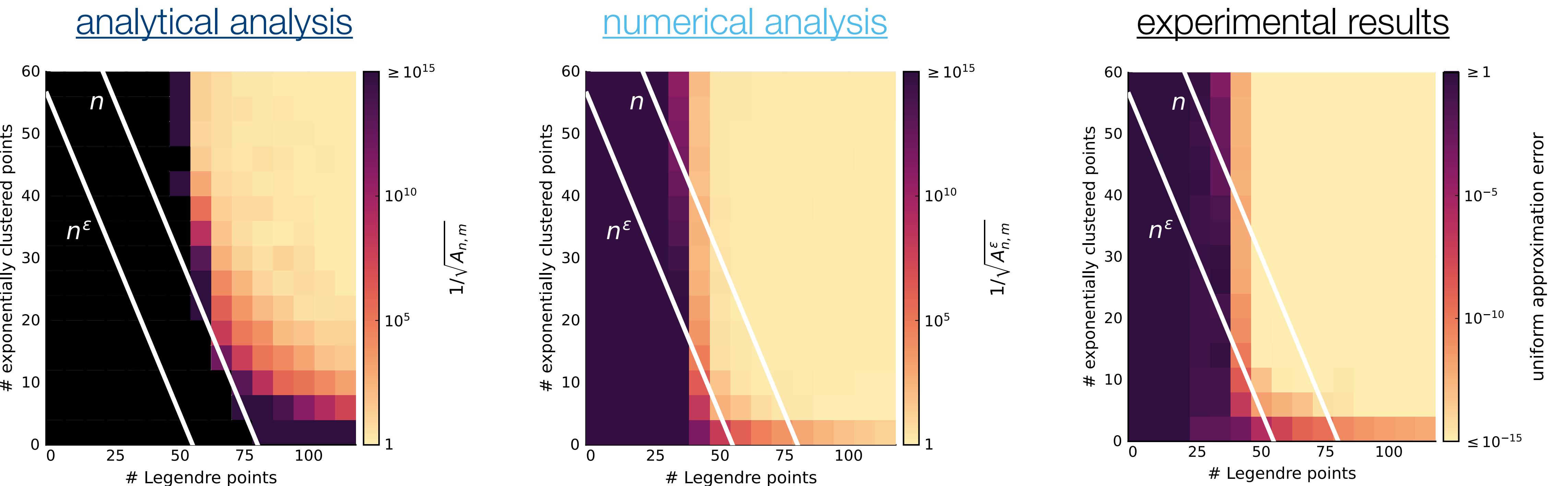
numerical analysis



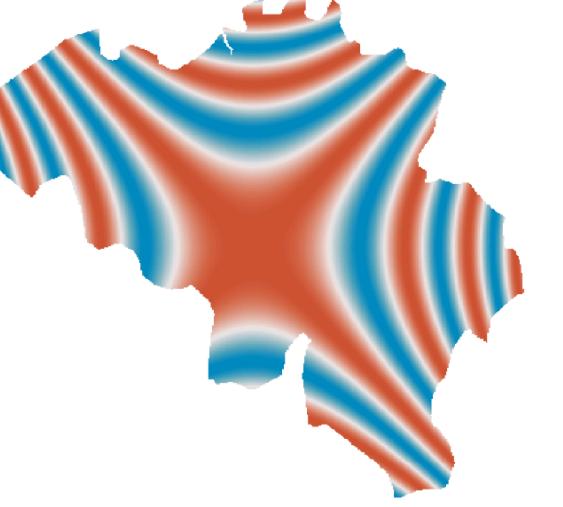
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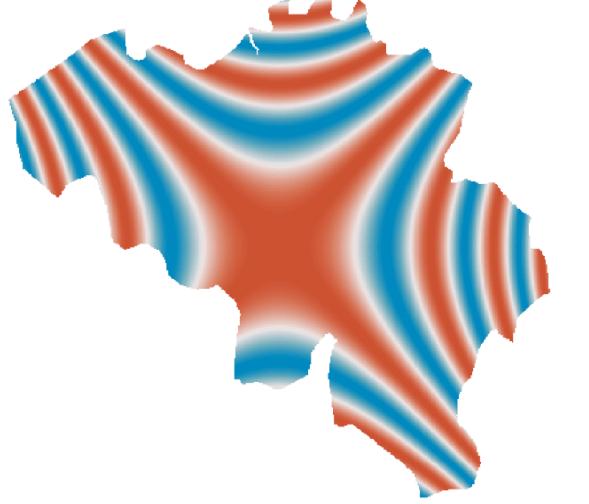
Random sampling



Approximation on $[-0.3, 0.3]$ using a Fourier extension on $[-0.5, 0.5]$

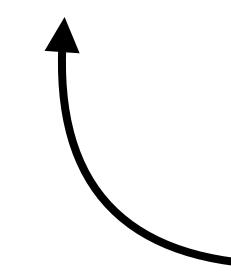
$$\phi_k = \exp(2\pi i x k), \quad - (n-1)/2 \leq k \leq (n-1)/2$$

Random sampling

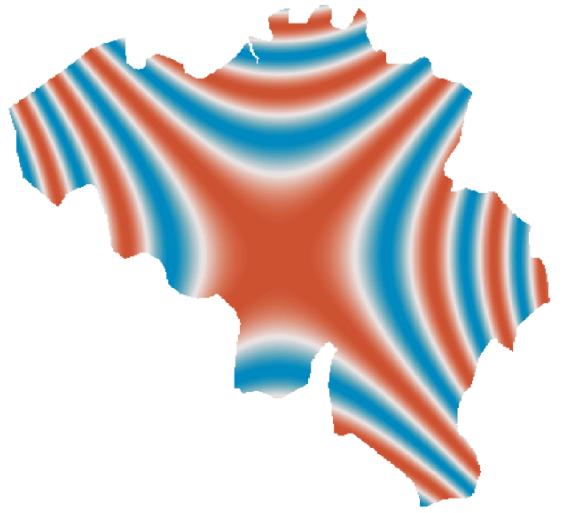


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When using uniformly random samples, the required number of samples equals

$$m \geq C \|k_n\|_\infty \log(n)$$

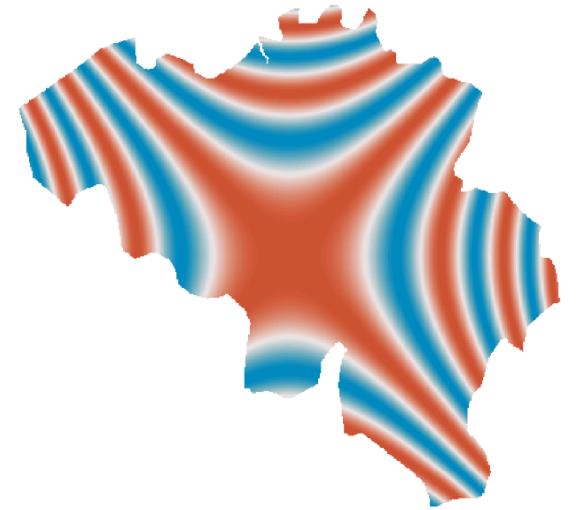
(analytically)

vs

$$m \geq C \|k_n^\epsilon\|_\infty \log(n^\epsilon)$$

(numerically)

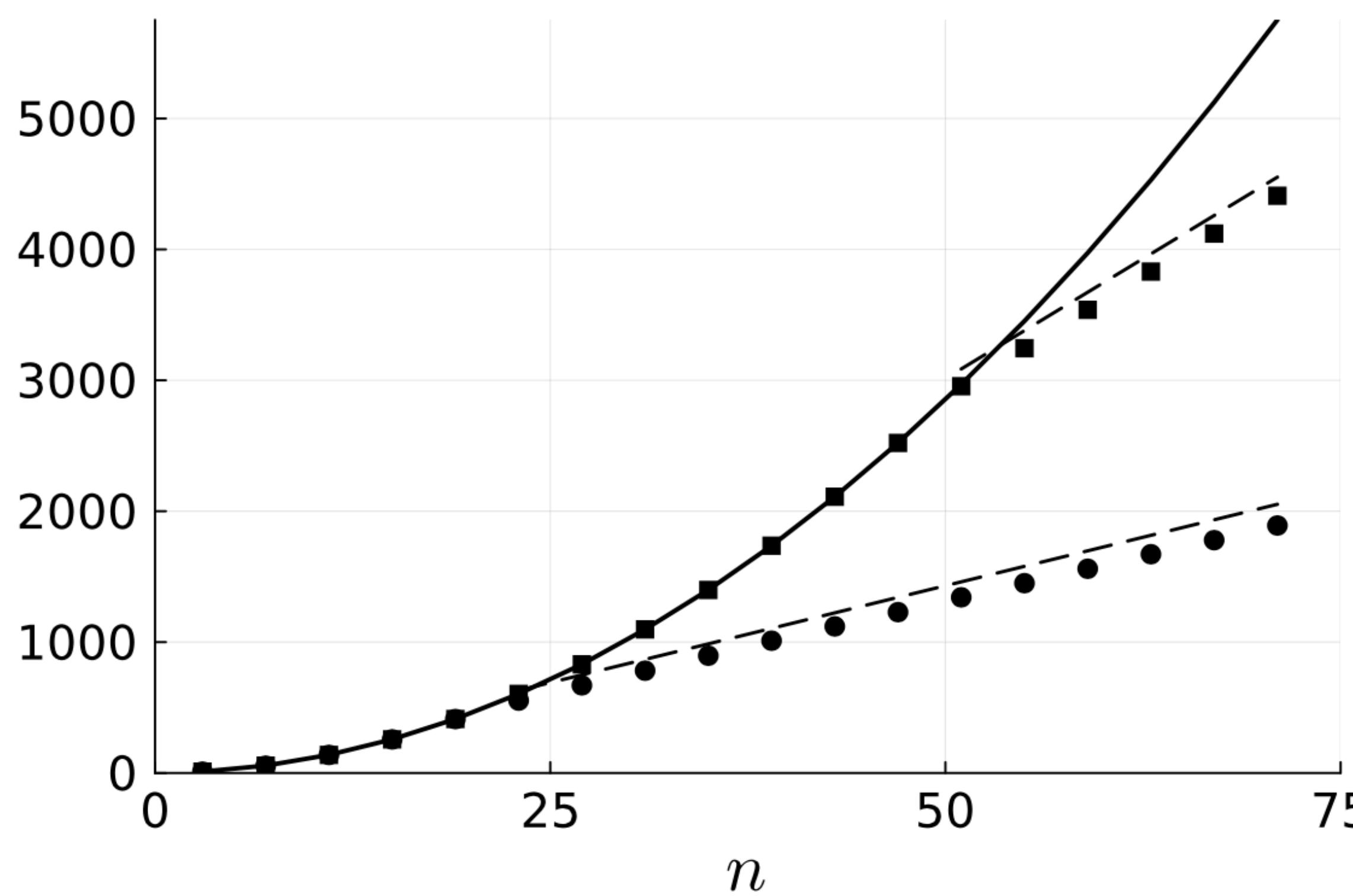
Random sampling



$m \geq C\|k_n\|_\infty \log(n)$
(analytically)

vs

$m \geq C\|k_n^\epsilon\|_\infty \log(n^\epsilon)$
(numerically)



$$\|k_n\|_\infty = \mathcal{O}(n^2)$$

$$\|k_n^\epsilon\|_\infty = \mathcal{O}(n) \text{ (double precision)}$$

$$\|k_n^\epsilon\|_\infty = \mathcal{O}(n) \text{ (single precision)}$$

Conclusions

- ▶ **The bad news** - slower convergence
- ▶ **The ugly news** - regularization
- ▶ **The good news** - less data

Conclusions

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Not immediately
clear who wins...

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+ more in the paper (on arXiv)

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Sampling Theory for Function Approximation
with Numerical Redundancy

Astrid Herremans* and Daan Huybrechs*

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- + I'm working on efficient Christoffel samplers

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