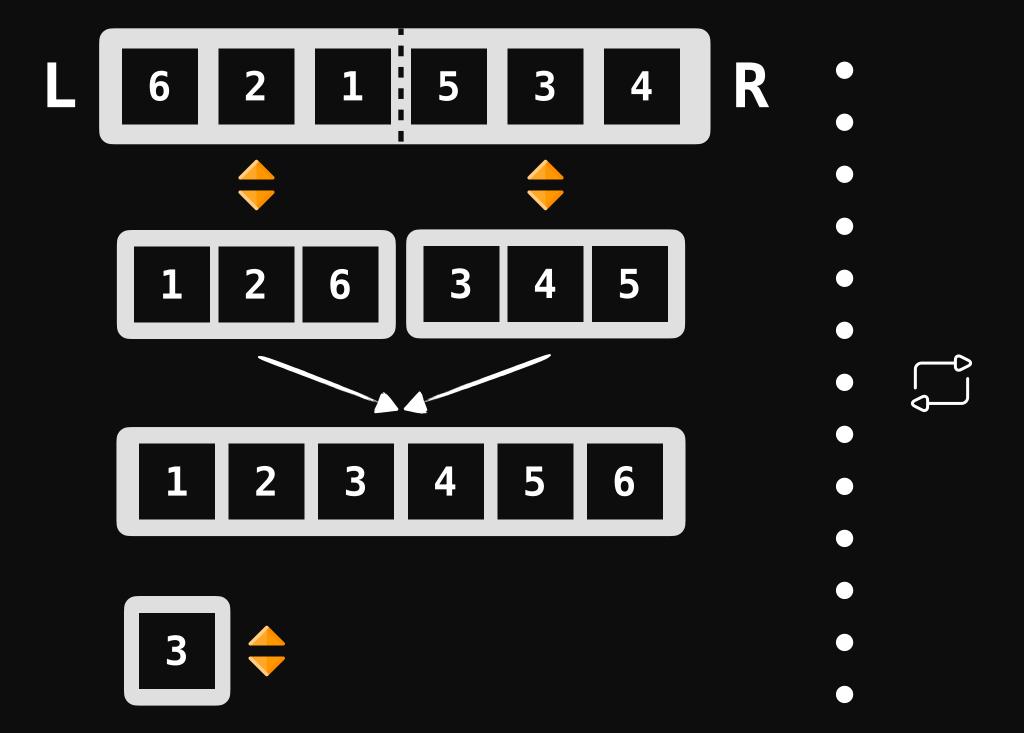
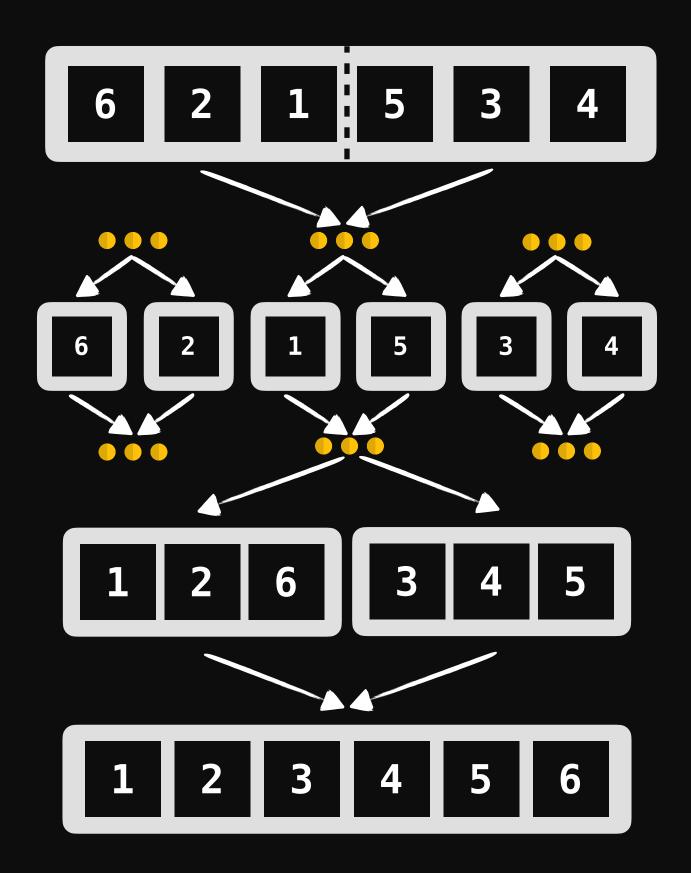
- Find the *middle* of the incoming array/slice
- Sort the *LEFT* side
- Sort the RIGHT side
- Merge the LEFT & RIGHT side
- Array/Slice of 1 element is considered sorted
- Repeat the process using divide & conquer

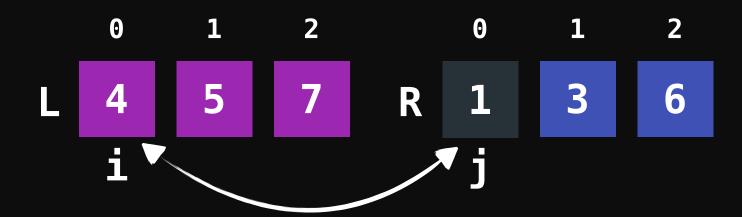


- Divide a problem into smaller problems of the same type
- Repeat the process till the problem cannot be divided
- Collect & merge indivisible problems (solutions) (Conquer)



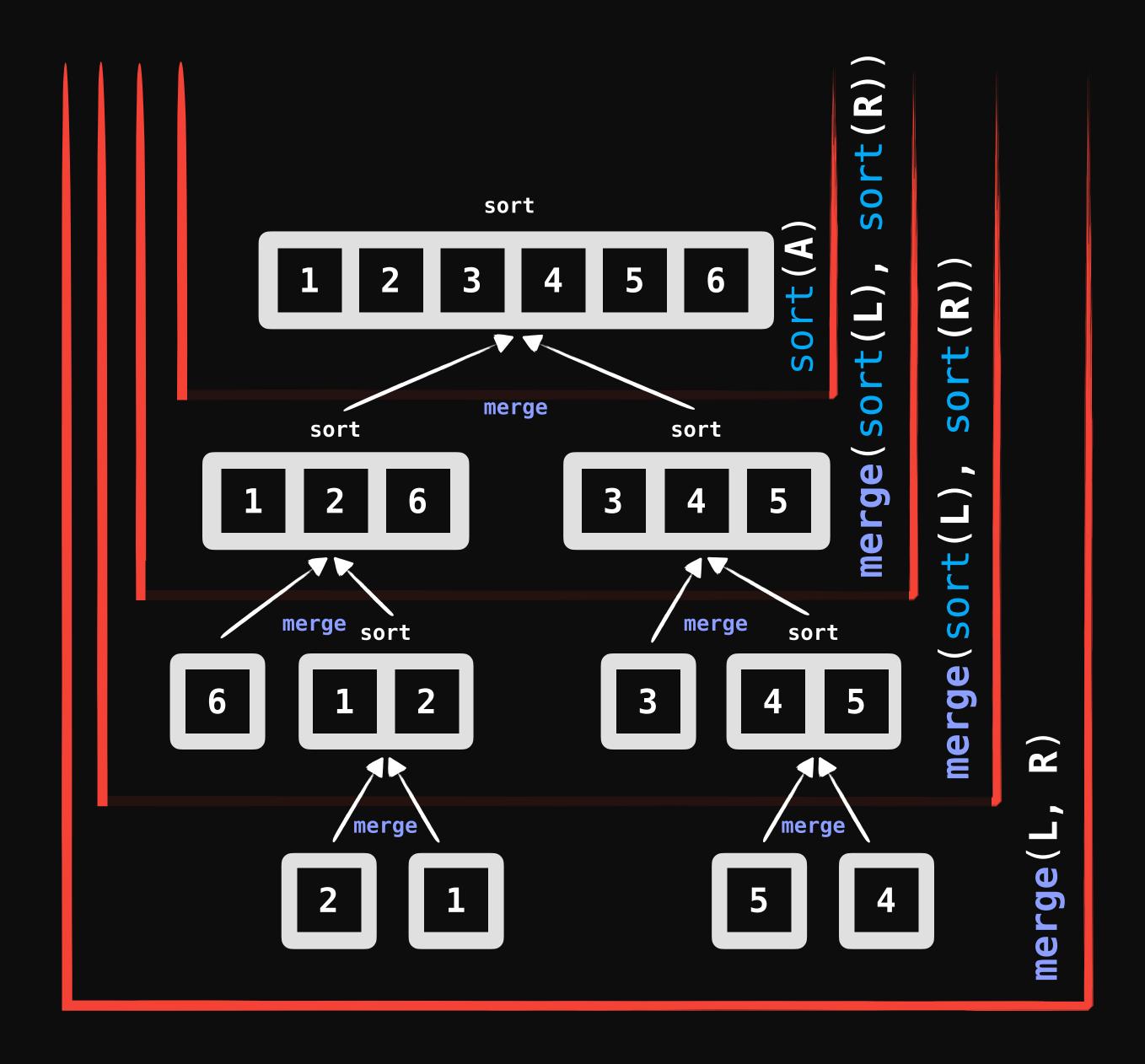
## THE ALGORITHM

- We need a **merge** & a **sort** function
- Split the slice in 2 sides
- Call **sort** recursively on the **Left** side
- Call **sort** recursively on the **Right** side
- Call merge recursively for Left & Right side
- Return the sorted merged slice





```
A := []int{6,2,1,3,5,4})
sort([]int{6,2,1,3,5,4})
sort([]int{6,2,1})
sort([]int{2,1})
merge([]int{2}, []int{1})
merge([]int{6}, []int{1,2})
sort([]int{3,5,4})
sort([]int{5,4})
merge([]int{5}, []int{4})
merge([]int{3}, []int{4,5})
merge([]int{1,2,6}, []int{3,4,5})
```



```
function merge(L, R) {
  let A[0..L.length+R.length]
  i, j, k = 0, 0, 0
 while i < L.length and j < R.length {</pre>
    if L[i] < R[j] {</pre>
      A[k] = L[i]; i++
    } else {
      A[k] = R[j]; j++
 while i < L.length {</pre>
    A[k] = L[i]; i++; k++
 while j < R.length {</pre>
    A[k] = L[j]; j++; k++
```

```
function sort(A) {
  if A.length > 1 {
   n1 = A.length/2
    n2 = A.length-n1
    let L[0..n1] and R[0..n2]
    for i=0 to n1
                       m = A \cdot length/2
      L[i] = A[i]
                       L = A[:m]
    for j=0 to n2
                       R = A[m:]
      R[j] = A[n1+j]
    A = merge(sort(R), sort(L))
 return A
```

## **IMPROVEMENTS**

- Merge sort uses temporary memory (L, R slices) (problem for big data sets)
- Merge sort works nicely on big data sets & is stable & predictable
- Insertion sort is faster for small data sets
- Insertion sort does not use temporary memory

## INSERTION MERGE SORT

- Use **merge sort** mostly
- Check data length before splitting the array/slice
- Use insertion sort if the data length is small enough

## **FACTS**

- Merge Sort is Efficient & Very Fast for big data sets
- Merge Sort Time Complexity: 0(n log(n))
- Merge Sort: Space Complexity: O(n)
- Insertion Sort is Efficient for small data
- Insertion Sort Best: 0(n)
- Insertion Sort Average/Worst: 0(n\*n)
- Insertion Sort: Space Complexity: 0(1)