

# Asymptotic Analysis + Dynamical Systems

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# 1. Basics of Asymptotic Analysis

## 1.1 Basic Notions and Definitions

One of the central theorems of the asymptotic analysis is the Taylor's theorem which is stated as below.

**Theorem 1.1 — Taylor's approximation theorem with reminder.** Let  $f : (a, b) \rightarrow \mathbb{R}$  be  $n + 1$  times differentiable. fix  $x_0 \in (a, b)$ . Then for  $x \in (a, b)$  we have

$$f(x) = f(x_0) + f'(x_0)(x-x_0) + \frac{f''(x_0)}{2!}(x-x_0)^2 + \cdots + \frac{f^{(n)}(x_0)}{n!}(x-x_0)^n + \frac{f^{(n+1)}(\xi)}{(n+1)!}(x-x_0)^{n+1}$$

for some  $\xi$  between  $x$  and  $x_0$ . This form of reminder is called the Lagrange form.

■ **Remark** There are also other possible forms of the reminder such as the Peano form, Cauchy form, Schlömilch form, and etc. For more details consult the Wikipedia page of the Taylor's approximation theorem.

■ **Example 1.1** Let  $f(x) = \sin(x)$ . Then from the Taylor's theorem with remainder we can write

$$\sin(x) = x - \frac{\sin(\xi)}{2}x^2$$

where  $\xi \in (0, x)$ . Note that in the expression above, the right hand side is not an approximate of  $\sin(x)$ , but it is actually  $\sin(x)$ . Note that in the reminder, we have  $\xi$  that belongs to  $(0, x)$ , and for different values of  $x$  we might find a different  $\xi$  that works. **So don't confuse the theorem above with the Taylor's series which is in fact an infinite sum** where for  $\sin(x)$  we have

$$\sin(x) = x - x^3/3! + x^5/5! - x^7/7! + \mathcal{O}(x^9)$$

■


**Definition 1.1 — Big-O notation.** Let  $f$  and  $g$  be functions from  $\mathbb{R}$  to  $\mathbb{R}$  “where  $f(x), g(x) \rightarrow 0$ , or  $f(x), g(x) \rightarrow \infty$  as  $x \rightarrow x_0$ ”<sup>a</sup>. We write

$$f = \mathcal{O}(g) \quad \text{as } x \rightarrow x_0$$

and read “ $f$  is big-Oh of  $g$ , if  $\exists k, \delta > 0$  such that

$$|x - x_0| < \delta \implies |f(x)| < k|g(x)|$$

<sup>a</sup>see the Be Careful Here section below

**Be Careful Here**  I searched all over the internet, and the definition above without the expression inside the quotation mark happens to be the standard one. But I think there is a problem with this definition, as it leads to some contradictions. For instance, define the following functions

$$f(x) = 100x^2 + 1, \quad g(x) = \frac{x}{100} + 1.$$

Then, following the definition above, I can both show that  $f = \mathcal{O}(g)$  and  $g = \mathcal{O}(f)$ , while we know that the function  $f$  goes to 1 faster than  $g$ . To show  $f = \mathcal{O}(g)$ , let  $\delta = 5$ . Then  $|f(x)| < 10^4|g(x)|$  for all  $|x| < 5$ . And to show  $g = \mathcal{O}(f)$ , let  $\delta = 5$ . Then we will have  $|g(x)| < 2|f(x)|$  for all  $|x| < 5$ . This issue will be fixed, and the definition will make the correct sense, if we assume  $f(x) \rightarrow 0$  and  $g(x) \rightarrow 0$  as  $x \rightarrow x_0$ .


**Definition 1.2 — Little-o definition.** Let  $f$  and  $g$  be functions from  $\mathbb{R}$  to  $\mathbb{R}$  “where  $f(x), g(x) \rightarrow 0$ , or  $f(x), g(x) \rightarrow \infty$  as  $x \rightarrow x_0$ ”<sup>a</sup>. We write

$$f = o(g) \quad \text{as } x \rightarrow x_0$$

and read “ $f$  is little-Oh of  $g$ , if  $\forall \delta > 0, \exists k > 0$  such that

$$0 < |x - x_0| < \delta \implies |f(x)| < k|g(x)|.$$

<sup>a</sup>see the Be Careful Here section above

**Observation**  **1.1.1 — Intuitive meaning of the big-O and little-o notation.** First, note that we use the big-O notation to compare two functions of the same kind of behaviour, i.e. both goes to 0 or both goes to  $\infty$ . Then, when we write  $f = \mathcal{O}(g)$  as  $x \rightarrow x_0$ , then it means that the function  $f$  will eventually be smaller than an scaled version of  $g$  as we get closer to  $x_0$ . However, if  $f = o(g)$ , then  $f(x)$  will get smaller than any scaled version of  $g(x)$  as  $x \rightarrow x_0$ .

Using the notion of limit superior we can define the the order notation in an alternative way.



**Definition 1.3** Let  $f, g : (a, b) \rightarrow \mathbb{R}$  and  $x_0 \in (a, b)$ , where  $f(x), g(x) \rightarrow \infty$  or  $f(x), g(x) \rightarrow 0$  as  $x \rightarrow x_0$ . Then we write  $f = \mathcal{O}(g)$  if

$$\limsup_{x \rightarrow x_0} \left| \frac{f(x)}{g(x)} \right| = L < \infty,$$

and  $f = o(g)$  if

$$\limsup_{x \rightarrow x_0} \left| \frac{f(x)}{g(x)} \right| = 0$$

Note that in both cases we assume  $g(x)$  is not zero on an open neighborhood of  $x_0$ .

*Proof.* Here, we prove that these two definitions are equivalent. For the first part we want to show that  $\limsup_{x \rightarrow x_0} \left| \frac{f(x)}{g(x)} \right| = L < \infty$  is equivalent to  $f = \mathcal{O}(g)$  as  $x \rightarrow x_0$ . To show this, assume  $\limsup_{x \rightarrow x_0} \left| \frac{f(x)}{g(x)} \right| = L < \infty$ . Then from the definition of  $\limsup$ , we know that  $\forall \epsilon > 0$  there exists  $\delta > 0$  such that  $|x - x_0| < \delta$  then  $\left| \frac{f(x)}{g(x)} \right| < L + \epsilon$ . Let  $\epsilon = 1$ , and define  $\hat{L} = L + 1$ . Then we know  $\exists \delta > 0$  such that for all  $|x - x_0| < \delta$  we have  $\left| \frac{f(x)}{g(x)} \right| < \hat{L}$ . This implies  $|f(x)| < \hat{L}|g(x)|$ . To show the converse, assume  $f = \mathcal{O}(g)$  as  $x \rightarrow x_0$ . Then there exists  $k, \delta > 0$  such that  $|x - x_0| < \delta$  implies  $|f(x)| < k|g(x)|$ . Then we have  $\left| \frac{f(x)}{g(x)} \right| < k$  for all  $|x - x_0| < \delta$ . From the definition of the  $\limsup$  we can conclude that  $\limsup_{x \rightarrow x_0} \left| \frac{f(x)}{g(x)} \right| < L$ .

For the second case, we want to show  $\limsup_{x \rightarrow x_0} \left| \frac{f(x)}{g(x)} \right| = 0$  if and only if  $f = o(g)$  as  $x \rightarrow x_0$ . For the forward direction, assume  $\limsup_{x \rightarrow x_0} \left| \frac{f(x)}{g(x)} \right| = 0$ . Then  $\forall \epsilon > 0$  there exists  $\delta > 0$  such that  $|x - x_0| < \delta$  implies  $\left| \frac{f(x)}{g(x)} \right| < \epsilon$ . This is exactly the definition of the little-O, thus we can conclude  $f = o(g)$ . For the converse, assume  $f = o(g)$  as  $x \rightarrow x_0$ . Then  $\forall k > 0$  there exists  $\delta > 0$  such that  $|x - x_0| < \delta$  implies  $|f(x)| < k|g(x)|$ . Since in the definition we have assumed that  $g(x) \neq 0$  in a neighborhood of  $x_0$ , then we can write  $\left| \frac{f(x)}{g(x)} \right| < k$ . Since for every  $k > 0$  we can find a  $\delta$  that holds the inequality, then from the definition of  $\limsup$  we can write

$$\limsup_{x \rightarrow x_0} \left| \frac{f(x)}{g(x)} \right| = 0,$$

which completes the proof. □

**Corollary 1.1** Let  $f, g : (a, b) \rightarrow \mathbb{R}$ , and fix  $x_0 \in (a, b)$ . Assume  $g(x) \neq 0$  in an open neighborhood of  $x_0$ . Also, assume  $f(x), g(x) \rightarrow 0$  or  $f(x), g(x) \rightarrow \infty$  as  $x \rightarrow x_0$ . Then

we have

$$\begin{aligned} \lim_{x \rightarrow x_0} \left| \frac{f(x)}{g(x)} \right| = L < \infty &\implies f = \mathcal{O}(g) \text{ as } x \rightarrow x_0, \\ \lim_{x \rightarrow x_0} \left| \frac{f(x)}{g(x)} \right| = 0 &\implies f = o(g) \text{ as } x \rightarrow x_0. \end{aligned}$$

*Proof.* The proof follows immediately from that fact that for any function  $f(x)$ ,  $\lim_{x \rightarrow x_0} f(x) = L$  implies  $\limsup_{x \rightarrow x_0} f(x) = L$ .  $\square$

**Corollary 1.2** Let  $f, g : (a, b) \rightarrow \mathbb{R}$  and fix  $x_0 \in (a, b)$ . Then

$$f = o(g) \implies f = \mathcal{O}(g).$$

*Proof.* We know  $f = o(g)$ , Let  $k = 1$ . Then  $\exists \delta > 0$  such that  $|x - x_0| < \delta$  implies  $|f(x)| < k|g(x)|$ . Since we have found the pair  $(k, \delta)$ , then we also have  $f = \mathcal{O}(g)$ . This is also very clear from the alternative definition of the orders via limsups.  $\square$

There are some important properties of the orders that comes in handy when dealing with different situations. These are summarized in the observation box below.

**Observation  $\odot$  1.1.2** The followings are some important properties of the order symbol that is very useful for quick calculations.

- $f = \mathcal{O}(1)$  as  $x \rightarrow x_0$  implies  $f$  is **bounded** in an open neighborhood of  $x_0$ .  
This is because, from the definition, we know that  $f$  will become eventually (i.e. if  $x$  is close enough to  $x_0$ ) smaller than some scaled version of  $g(x) \equiv 1$  function, thus implies  $f$  is bounded near  $x_0$ .
- $f = o(1)$  as  $x \rightarrow x_0$  implies  $f \rightarrow 0$  as  $x \rightarrow x_0$ .  
This is because  $f = o(g)$  implies  $f$  will be smaller than any scaled version of  $g$  when  $x$  is close enough to  $x_0$ . Thus  $f = o(1)$  as  $x \rightarrow x_0$  literally translates to  $\forall \epsilon > 0$ , there exists  $\delta > 0$  such that  $|x - x_0| < \delta$  implies  $|f(x)| < \epsilon$ , hence  $f(x) \rightarrow 0$  as  $x \rightarrow x_0$ .

**Lemma 1.1** Let  $g : U \rightarrow \mathbb{R}$ . Then we have

$$o(g + o(g)) = o(g) \quad \text{as } x \rightarrow x_0.$$

*Proof.* Let  $f$  be a function such that  $f = o(g + o(g))$  as  $x \rightarrow x_0$ . Fix  $k$ . Then  $\exists \delta > 0$  such that  $0 < |x - x_0| < \delta$  implies  $|f(x)| < k/2|g(x) + o(g(x))|$ . Then from the triangle inequality, we can write

$$|f(x)| < k/2|g(x)| + k/2|o(g(x))|.$$



Also, we can find  $\delta_1 > 0$  such that  $0 < |x - x_0| < \delta_1$  implies  $|o(g(x))| < 1|g(x)|$ . Let  $\hat{\delta} = \min\{\delta, \delta_1\}$ . Then  $0 < |x - x_0| < \hat{\delta}$  implies

$$|f(x)| < k|g(x)|.$$

This completes the proof.  $\square$

The following definition box introduces two very important symbols that will be used extensively through out the text.

**Definition 1.4 — A tale of two symbols.** The symbol  $\ll$  used for two functions, i.e.  $f(x) \ll g(x)$  as  $x \rightarrow x_0$  means  $f = o(g)$  as  $x \rightarrow x_0$ . The symbol  $\approx$  dose not have a precise definition and it is used to designate an approximate numerical value. For instance  $\pi \approx 3.1415$ .

### 1.1.1 Asymptotic Approximation

Here in this section we will go through more formal definitions of the idea of approximating a function. The following definition is a key.

**Definition 1.5 — Asymptotic approximation.** Let  $f, \varphi : U \rightarrow \mathbb{R}^n$ , where  $U$  is an open subset of  $\mathbb{R}^n$ . Let  $x_0 \in U$ . We say  $\varphi$  is *asymptotic approximation* of the function  $f$ , if we have

$$f(x) = \varphi(x) + o(\varphi) \quad \text{as} \quad x \rightarrow x_0.$$

We denote this as  $f \sim \varphi$  as  $x \rightarrow x_0$ .

**Proposition 1.1 — Asymptotic approximation defines an equivalence relation.** Let  $\sim$  be a relation on the set  $C^\infty(U)$  where  $U \subset \mathbb{R}$ , given by

$$f \sim g \quad \text{then} \quad f = g + o(g) \quad \text{as} \quad x \rightarrow x_0.$$

i.e.  $g$  is an asymptotic approximation of  $f$  as  $x \rightarrow x_0$ . The relation  $\sim$  defines an equivalence relation.

*Proof. Reflexivity.* Since the constant function  $O(x) \equiv 0$  is little-O of any function, then, for any  $C^\infty(U \ni x_0)$  function we can write

$$f(x) = f(x) + o(0) \quad \text{as} \quad x \rightarrow x_0.$$

Thus the reflexivity holds, i.e.  $f \sim f$  for all  $f \in C^\infty(U)$

**Symmetry.** Let  $f, g \in C^\infty(U)$  such that  $f \sim g$ . Then we can write  $f(x) = g(x) + o(g(x))$  as  $x \rightarrow x_0$ . This implies  $o(f(x)) = o(g(x) + o(g(x)))$ . For the right hand side, from [Lemma 1.1](#) we have

$$o(g + o(g)) = o(g).$$

Thus we have  $o(f) = o(g)$ . Finally we can get

$$\varphi = f + o(\varphi) = f + o(f).$$

This shows that the relation is symmetric, i.e.  $f \sim g$  implies  $g \sim f$ .

**Transitivity.** Let  $f \sim g$  and  $g \sim h$ . Then

$$f = g + o(g), \quad g = h + o(h).$$

By substituting the second expression in the first one we will get

$$f = h + o(h + o(h)) = h + o(h).$$

This  $f \sim h$ . This completes the proof for  $\sim$  being an equivalence relation.  $\square$

**Proposition 1.2** Let  $f, g : U \rightarrow \mathbb{R}$  where  $x_0 \in U \subset \mathbb{R}$ , and assume  $g(x) \neq 0$  for all  $x$  in an open neighborhood of  $x_0$ . Then

$$\lim_{x \rightarrow x_0} \frac{f(x)}{g(x)} = 1 \quad \implies \quad f \sim g \text{ as } x \rightarrow x_0.$$

*Proof.* We have  $\lim_{x \rightarrow x_0} \frac{f(x)}{g(x)} = 1$ . Then, from the definition of limit we know that  $\forall \epsilon > 0$  there exists  $\delta > 0$  such that  $0 < |x - x_0| < \delta$  implies  $|\frac{f(x)}{g(x)} - 1| < \epsilon$ . Then we can write

$$|f(x) - g(x)| < \epsilon |g(x)|.$$

This is literally the definition of little-O. Thus we have

$$f(x) - g(x) = o(g(x)) \quad \text{as } x \rightarrow x_0.$$

By re-arranging the terms we will have

$$f(x) = g(x) + o(g(x)) \quad \text{as } x \rightarrow x_0.$$

$\square$

## 1.2 Solved Problems

■ **Problem 1.1** What values of  $\alpha$ , if any, yield  $f = O(x^\alpha)$  as  $x \rightarrow 0$ .

- (a)  $f(x) = \sqrt{1+x^2}$ .
- (b)  $f(x) = x \sin(x)$ .
- (c)  $f(x) = (1 - e^x)^{-1}$ .
- (d)  $f(x) = \ln(1+x)$ .
- (e)  $f(x) = x \ln(x)$ .
- (f)  $f(x) = \sin(1/x)$ .
- (g)  $f(x) = \sqrt{x+a}$  where  $0 \leq a \leq 1$ .

**Solution** (a) In order to have  $f = O(x^\alpha)$  as  $x \rightarrow x_0$ , we need to have

$$\lim_{x \rightarrow 0} \frac{\sqrt{1+x^2}}{x^\alpha}.$$

Since the numerator goes to 1 as  $x \rightarrow 0$ , then we must have  $\alpha < 0$ . Then we will have

$$\lim_{x \rightarrow 0} \frac{\sqrt{1+x^2}}{x^\alpha} = 0 \quad \text{for } \alpha < 0.$$

Thus for the negative values of  $\alpha$  we have  $f = o(g)$  as  $x \rightarrow x_0$ , which also implies  $f = O(g)$  as  $x \rightarrow x_0$ .

(b) We need to have

$$\lim_{x \rightarrow 0} \frac{x \sin(x)}{x^\alpha} = L < \infty.$$

From the Taylor's theorem with remainder we have

$$\sin(x) = x - \cos(\xi)x^3/3! \quad \text{for some } \xi \in (0, x).$$

Thus

$$\lim_{x \rightarrow 0} \frac{x \sin(x)}{x^\alpha} = \lim_{x \rightarrow 0} \frac{x^2 - \cos(\xi)x^4/3!}{x^\alpha}.$$

This implies  $\alpha \leq 2$  so that the limit above is finite. Since otherwise, by applying the l'Hopital's rule two times we can see that the limit diverges to  $\infty$ . In a nutshell

$$f = O(g) \text{ as } x \rightarrow 0, \quad \text{when } \alpha \leq 2$$

and also

$$f = o(g) \text{ as } x \rightarrow 0, \quad \text{when } \alpha < 2$$

(c) First, we want to calculate the series expansion of  $f$  at 0 (which will be its Laurent expansion). To do so, we can write

$$f(x) = \frac{1}{1 - e^{-x}} = \frac{1}{1 - (1 - x + x^2/2 + O(x^3))} = \frac{1}{x - x^2/2 + O(x^3)} = \frac{1}{x} \cdot \frac{1}{1 - x/2 + O(x^2)}$$

Now we can simplify the term above using the geometric series

$$f(x) = \frac{1}{x}(1 + x/2 + \mathcal{O}(x^2)) = \frac{1}{x} + \frac{1}{2} + \mathcal{O}(x) \quad \text{as } x \rightarrow 0.$$

In order to  $f = \mathcal{O}(x^\alpha)$ , then  $x^\alpha$  should go to infinity faster than  $f$ . Thus  $\alpha \leq 0$ . To write this mathematically, we want the following limit to be finite. Then

$$\lim_{x \rightarrow 0} \frac{1/x + 1/2 + \mathcal{O}(x)}{x^\alpha} = \lim_{x \rightarrow 0} x^{-1-\alpha} + 1/2x^{-\alpha} + \mathcal{O}(x^{1-\alpha}).$$

Thus we need to have

$$-1 - \alpha \geq 0 \quad \wedge \quad -\alpha \geq 0 \quad \wedge \quad 1 - \alpha \geq 0$$

The intersection of the solutions of the inequalities above is

$$\alpha \leq -1.$$

Thus in summary

$$f = \mathcal{O}(x^\alpha) \quad \text{as } x \rightarrow 0 \quad \text{when } \alpha \leq -1,$$

and

$$f = o(x^\alpha) \quad \text{as } x \rightarrow 0 \quad \text{when } \alpha < -1.$$

(d) From the Taylor's theorem with remainder, we have

$$\ln(1+x) = x - x^2/2 + \frac{x^3}{3(1+\xi)^3} \quad \text{for some } \xi \in (0, x).$$

Then for  $\alpha \leq 1$  we have  $f = \mathcal{O}(x^\alpha)$ , and for  $\alpha < 1$  we have  $f = o(x^\alpha)$ .

(e) Still thinking how to solve.

(f) The function given by  $f(x) = \sin(1/x)$  oscillates between -1 and 1 very rapidly as  $x \rightarrow 0$ . Thus for any positive value of  $\alpha$ , i.e.  $\alpha > 0$ , for any  $\delta > 0$  there are infinity many points, denoted by  $Z_\delta$  such that  $\sin(x) = 1$  for all  $x \in Z_\delta$ . Thus  $f \neq \mathcal{O}(x^\alpha)$  for any positive value of  $\alpha$ . However, for  $\alpha \leq 0$ , we have  $f = \mathcal{O}(x^\alpha)$ . That is because, in the case  $\alpha = 0$ , then since  $\sin(1/x)$  is bounded, then it let  $k = 2, \delta = 1$ . In the case where  $\alpha < 0$  let  $\delta = 1$  and  $k = 1$ . Thus  $\sin(1/x) = \mathcal{O}(x^\alpha)$  for all  $\alpha < 0$ . Note that for  $\alpha = 0$ ,  $f \neq o(x^\alpha)$ . However, for  $\alpha < 0$  we also have  $f = o(x^\alpha)$ , i.e. in the case where  $\alpha < 0$  the function given by  $\sin(1/x)$  will eventually be smaller than any scaled version of  $x^\alpha$  when  $x$  gets sufficiently close to 0.

■ **Problem 1.2** Define the function  $f : \mathbb{R} \rightarrow \mathbb{R}$  as follows

$$f(x) = \begin{cases} 0 & x \leq 0, \\ e^{-1/x} & x > 0. \end{cases}$$

Show that  $f = o(x^n)$  for all  $n \in \mathbb{N}$ .

**Solution** Since  $f(x) \rightarrow 0$  as  $x \rightarrow 0$ , and  $g(x) = x^n \rightarrow 0$  as well (for all  $n \in \mathbb{N}$ ) we can use the limit criteria to determine the order. We compute

$$\lim_{x \rightarrow 0} \frac{f(x)}{g(x)} = \lim_{x \rightarrow 0} \frac{e^{-1/x}}{x^n} = \lim_{y \rightarrow \infty} \frac{e^{-y}}{(1/y)^n} = \lim_{y \rightarrow \infty} \frac{y^n}{e^y}$$

Now we can apply the l'Hopital law for  $n$  times then we will get

$$\lim_{x \rightarrow 0} \frac{f(x)}{g(x)} = 0.$$

This proves that  $f = o(g)$ .