

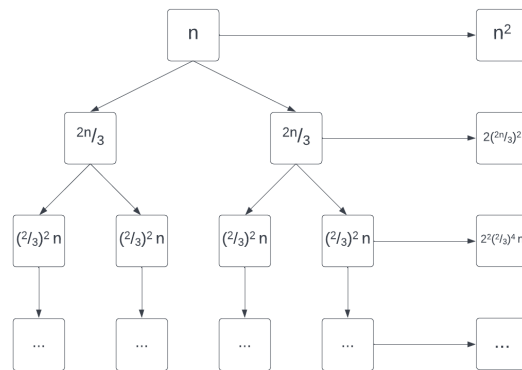
### Homework#3

#### Problem 1 (5 points):

Solve the recurrence  $T(n) = 2T\left(\frac{2}{3}n\right) + n^2$  first by using a recursion tree and then using the Master theorem. Show work.

Recursion Tree:

First, we calculate cost at each level:



$$\frac{2n}{3} + \frac{2n}{3} \rightarrow \left(\frac{2}{3}\right)^2 n^2 + \left(\frac{2}{3}\right)^2 n^2 \rightarrow 2\left(\frac{2}{3}n\right)^2$$

$$\left(\frac{2}{3}\right)^2 n^2 + \left(\frac{2}{3}\right)^2 n + \dots \rightarrow 2^2 \left(\frac{2}{3}\right)^4 n^2$$

Assuming k levels, at the kth level we see:

$$\left(\frac{2}{3}\right)^k n = 1, \text{ and so } k = \log_{2/3} n \text{ levels}$$

Finally, identify asymptotic bound:

$$T(n) = n^2 \left[ 2^0 * \left(\frac{2}{3}\right)^0 + 2^1 * \left(\frac{2}{3}\right)^1 + 2^2 * \left(\frac{2}{3}\right)^2 + \dots \text{to } k \right]$$

$$T(n) = n^2 \left[ \left(\frac{8}{9}\right)^0 + \left(\frac{8}{9}\right)^1 + \left(\frac{8}{9}\right)^2 + \dots \text{to } k \right]$$

$$T(n) = n^2 \left[ \frac{1 * \left(1 - \left(\frac{8}{9}\right)^{\log_{3/2} n + 1}\right)}{1 - \frac{8}{9}} \right] = n^2 \left[ 1 - \frac{8}{9} * n^{\log_{3/2} \left(\frac{8}{9}\right)} \right]$$

$$\therefore \text{As } n^{\log_{3/2} \left(\frac{8}{9}\right)} \rightarrow 0, \text{ We see that } T(n) = \theta(n^2)$$

### Homework#3

Master Theorem:

$$T(n) = aT\left(\frac{n}{b}\right) + f(n)$$

...where  $f(n)$  is of the form  $n^c$  in which  $c \geq 0$

$$T(n) \in \begin{cases} \text{if } f(n) = O(n^{\log_b a - \varepsilon}) & T(n) = O(n^{\log_b a}) & \text{Case 1} \\ \text{if } f(n) = \Omega(n^{\log_b a + \varepsilon}) & T(n) = \theta(f(n)) & \text{Case 2} \\ \text{if } f(n) = \theta(n^{\log_b a} * (\log n)^k) & T(n) = \theta(n^{\log_b a} * (\log n)^{k+1}) & \text{Case 3} \end{cases}$$

First, we can rearrange:

$$T(n) = 2T\left(\frac{2}{3}n\right) + n^2$$

We see that  $f(n) = n^2$ , and  $n^{\log_b a} = n^{\log_{3/2} 2}$

$$f(n) = \Omega(n^{\log_b a + \varepsilon})$$

$$\text{when } f(n) \geq n^{\frac{\log_{3/2} 2}{2}},$$

**$\therefore$  Therefore  $T(n) = \theta(n^2)$**

Homework#3

**Problem 2 (5 points):**

Give asymptotic upper and lower bounds for the recurrence  $T(n) = 3T\left(\frac{n}{2}\right) + \frac{n}{\log n}$ . Justify your answer.

**Using the Master Theorem:**

$$T(n) = aT\left(\frac{n}{b}\right) + f(n), \text{ where } a \geq 1, b \geq 1, f \text{ is asymptotically positive}$$

$$a = 3, b = 2, d = 1$$

**Case #1:  $d > \log_b a$**

$$f(n) = \theta(n^d) \text{ when } d > \log_b a$$

$$T(n) \in \theta(f(n))$$

**Case #2:  $d = \log_b a$**

$$f(n) = \theta(n^d \log n) \text{ when } d = \log_b a$$

$$\text{Therefore, } T(n) \in \theta(n^d \log n)$$

**Case #3:  $d < \log_b a$**

$$f(n) = \theta(n^{\log_b a}) \text{ when } d < \log_b a$$

$$T(n) \in \theta(n^{\log_b a})$$

$$T(n) = 3T\left(\frac{n}{2}\right) + \frac{n}{\log n}$$

$$\text{When } a = 3, b = 2, \text{ and } f(n) = \frac{n}{\log n} < n$$

$\therefore$  We can say that  $f(n) = \theta(n)$ , and using case #1,  $T(n) \in \theta(n^{\log_2 3})$

Homework#3

**Problem 3 (10 = 5 + 5 points):**

**Give asymptotic upper and lower bounds for each of the following recurrences. Justify your answer.**

**(a)  $T(n) = \sqrt{n}T(\sqrt{n}) + n$**

$$T(n) = \sqrt{n} * T(\sqrt{n}) + n$$

$$T(n) = n^{\frac{1}{2}} * T\left(n^{\frac{1}{2}}\right) + n$$

$$T(n) = n^{\frac{1}{2}} * \left( n^{\frac{1}{4}} T\left( n^{\frac{1}{4}} + n^{\frac{1}{2}} \right) \right) + n$$

$$T(n) = n^{\frac{3}{4}} * T\left( n^{\frac{1}{2}} \right) + 2n$$

$$T(n) = n^{\frac{7}{8}} * T\left( n^{\frac{1}{8}} \right) + 3n$$

$$T(n) = n^{1-\frac{1}{2^k}} * T\left( n^{\frac{1}{2^k}} \right) + kn$$

*In the event that  $n^{\frac{1}{2^k}}$  is less than 2, then  $k > \log \log n$*

**$\therefore$  Therefore,  $T(n) = \theta(n \log \log n)$**

*Alternatively we can also say that*

$$S(n) = \frac{T(n)}{n}$$

*In this case, the recurrence now becomes:*

$$S(n) = S(\sqrt{n}) + 1$$

*Via recursion tree, we can see that  $S(n) = \theta(\log \log n)$*

**$\therefore$  Therefore,  $T(n) = \theta(n \log \log n)$**

### Homework#3

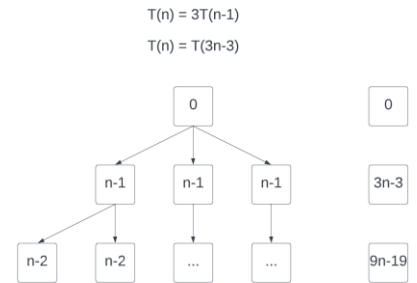
**(b)  $T(n) = 3T(n - 1)$**

*Via Recursion Tree, we can say that:*

$$T(n) = 3T(n - 1)$$

$$T(n) = 3T(n - 1) + 0$$

$$\begin{aligned} T(n) &= T(3n - 3) + 0 \\ &= 3^0n + 3^1n + 3^2n + 3^3n + \dots \\ \therefore \text{Therefore, } T(n) &= \theta(3^n) \end{aligned}$$



*In addition:*

$$T(n) = 3(3T(n - 2))$$

$$T(n) = 3^2T(n - 2)$$

$$T(n) = 3^2(3T(n - 3)) = 3^3(T(n - 3))$$

*Given a value  $k$ , we can say that*

$$T(n) = 3^k(T(n - k))$$

$$T(n) = 3^n(T(0))$$

$$\therefore \text{Therefore, } T(n) = \theta(3^n)$$

### Homework#3

#### Problem 4: Programming Assignment (10 points):

- **Maximum Product Subarray**
- Given an integer array `nums`, find a contiguous non-empty subarray within the array that has the largest product, and return *the product*.
- The test cases are generated so that the answer will fit in a 32-bit integer.
- A subarray is a contiguous subsequence of the array.

Please find a zip file attached with the code provided in the form of a Jupyter notebook. Using **Python3**, you can install Jupyter Notebook via 'pip install jupyter notebook'. Once installed, you can run the commands you will see below.

The following figure shows the function **maximumSubArrayProductFinder** which takes one argument, an array **nums**. The function will check for an empty or too large of an array, before going through the given array to find the largest product of a sub array.

```
[47]: def maximumSubArrayProductFinder(nums):  
    """  
    Returns the largest product of a subarray found in  
    the input array nums  
  
    Input (array) : An array such as [1, 2, 3, -5]  
  
    Returns (int) : An integer representing the largest product of a subarray  
    """  
  
    # Check for empty array  
    if len(nums) < 1:  
        return "Empty Array"  
    elif len(nums) > 20000:  
        return "Too large"  
  
    # Reverse order of array  
    nums_reversed = nums[::-1]  
  
    print("Original: ", nums)  
    print("Reversed: ", nums_reversed)  
  
    # Iterate over the array nums  
    for num in range(1, len(nums)):  
  
        # Alter nums[num] with the product  
        nums[num] *= nums[num - 1] or 1  
  
        # Alter nums_reversed[num] with the product  
        nums_reversed[num] *= nums_reversed[num - 1] or 1  
  
    # Complete array:  
    full_array = nums + nums_reversed  
  
    # Find maximum:  
    maximum = max(full_array)  
  
    # Return the maximum of the array  
    return maximum
```

### Homework#3

The following figure show 11 test cases developed to demonstrate the both the examples included in the assignments specification, as well as a few others. Each test case has an **input\_array\_n**, the **function** being executed with that array, as well as the **output**. These test cases show the utility of the application to handle the ability of finding sub array products, as well as ensuring that edge cases are handled properly.

```
input_array_1 = [2,3,-2,4]
maximumSubArrayProductFinder(input_array_1)
```

```
Original: [2, 3, -2, 4]
Reversed: [4, -2, 3, 2]
6
```

```
input_array_2 = [-2, 0, -1]
maximumSubArrayProductFinder(input_array_2)
```

```
Original: [-2, 0, -1]
Reversed: [-1, 0, -2]
0
```

```
input_array_3 = [-10, 5, -5, 0, 10]
maximumSubArrayProductFinder(input_array_3)
```

```
Original: [-10, 5, -5, 0, 10]
Reversed: [10, 0, -5, 5, -10]
250
```

```
input_array_4 = [-1, 2, -3, 4, 5]
maximumSubArrayProductFinder(input_array_4)
```

```
Original: [-1, 2, -3, 4, 5]
Reversed: [5, 4, -3, 2, -1]
120
```

```
input_array_5 = [-10, 5, -5, 0, 10, 2, 5, 7, 10, 8, 5, 1, -10]
maximumSubArrayProductFinder(input_array_5)
```

```
Original: [-10, 5, -5, 0, 10, 2, 5, 7, 10, 8, 5, 1, -10]
Reversed: [-10, 1, 5, 8, 10, 7, 5, 2, 10, 0, -5, 5, -10]
280000
```

```
input_array_6 = [-2, 0, 0, 0, 0, -55]
maximumSubArrayProductFinder(input_array_6)
```

```
Original: [-2, 0, 0, 0, 0, -55]
Reversed: [-55, 0, 0, 0, 0, -2]
0
```

```
input_array_7 = [-2, -55, 0]
maximumSubArrayProductFinder(input_array_7)
```

```
Original: [-2, -55, 0]
Reversed: [0, -55, -2]
110
```

```
input_array_8 = [0, -2, -1]
maximumSubArrayProductFinder(input_array_8)
```

```
Original: [0, -2, -1]
Reversed: [-1, -2, 0]
2
```

```
input_array_9 = [0]
maximumSubArrayProductFinder(input_array_9)
```

```
Original: [0]
Reversed: [0]
0
```

```
input_array_10 = []
maximumSubArrayProductFinder(input_array_10)
```

'Empty Array'

```
input_array_11 = [5]*20001
maximumSubArrayProductFinder(input_array_11)
```

'Too large'