

Homework #9

Problem 1. Longest common subsequence problem (10 pts)

Given two strings $x_1...x_n, y_1...y_m$ wish to find the length of their longest common subsequence, that is, the largest k for which there exist indices $i_1 < ... < i_k$ and $j_1 < ... < j_k$ such that $x_{i_1}...x_{i_k} = y_{j_1}...y_{j_k}$. Show how to do this in time $O(nm)$.

- The longest common subsequence is defined as the longest subsequence found across two or more sequences, regardless of position within the individual sequences
- One standard method of tackling such an item within the confines of dynamic programming is as follows:
 - We can keep a record of the longest common suffixes by storing them
 - As we store these, we maintain that the longest common suffix for $x[0, 1, \dots, i-1]$ and $y[0, 1, \dots, i-1]$
 - We can begin by iterating over the sequence in the form of a loop over the length of m , building them in a bottom-up approach
 - We then enter a second loop over the sequence over the length of n
 - When the condition of i and j is met, we get the longest common suffix
- Since we **iterate over m and n** respectively, the algorithm gives us a **time complexity of $O(nm)$**

```
def LongestCommonSubsequence(x, y):
    """
    Returns the Longest Common Subsequence
    @input x: Input String Sequence x
    @input y: Input String Sequence y
    """

    # Calculate the lengths of the sequences
    m, n = len(x), len(y)
    strd_A = [[None]*(n + 1) for i in range(m + 1)]

    # Enter first loop over range m+1
    for i in range(m + 1):
        # Enter second loop over range n+1
        for j in range(n + 1):
            # Condition when == 0
            if i == 0 or j == 0 :
                strd_A[i][j] = 0
            elif x[i-1] == y[j-1]:
                strd_A[i][j] = strd_A[i-1][j-1]+1
            else:
                strd_A[i][j] = max(strd_A[i-1][j], strd_A[i][j-1])

    # Return lengths
    return strd_A[m][n]
```

- Using this method, if we use the input values of "AGTCCGCGAA", and "GAA", we will see that GAA is the longest common value with a length of 3.

Problem 2. Longest subsequence problem with linear space complexity (10 pts)

Find a space complexity of your solution for Problem 1.

- In the algorithm above, we use the variable `strd_A` to store m and n , we can determine that the space complexity would be **$O(n*m)$** , similar to the time complexity.

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Problem 3. Longest increasing subsequence problem (10 pts)

Given a string $x_1 \dots x_n$ we wish to find the length of its longest increasing subsequence, that is, the largest k for which there exist indices $i_1 < \dots < i_k$ such that $x_{i_1} < x_{i_2} < \dots < x_{i_k}$. Show how to do this in time $O(n^2)$.

- We can define the longest increasing subsequence as the length of a sequence by which the values remain to be increasing within a given input list. For example, in the list [1, 2, 4, 3, 7, 9], we get an LIS of length 5 from [1, 2, 4, 7, 9].
- We can use a dynamic programming approach here by storing a list for LIS as part of the process
- Algorithm:
 - We begin by determining the length of the given input array
 - Next, we create a list to store the values
 - We then create a variable to remember the maximum value visited thus far
 - Next, we iterate over the array in the form of loop, and then iterate over the individual element
 - We check to see if the value of the first loop is greater than the second, thereby checking to determine if its incrementing, essentially seeing of $x[i] > x[j]$
 - If the condition is met, we store that value in our previously created array
 - At the end of both loops, we can check the maximum value that was stored in the process in order to return that length as part of the question.
 - Since **two nested loops** are involved in the process, we will see a time complexity of **$O(n^2)$**

```
def LongestIncreasingSubsequence(x):  
    """  
    Function that determines LIS for given array  
    @input x: Input array of values  
    """  
  
    # Determine length and create placeholder array  
    x_len = len(x)  
    strd_arr = [1]*x_len  
    max_lis = 1  
  
    # Enter double loop to determine LIS and store in array  
    for i in range(1, x_len):  
        for j in range(0, i):  
            # Compare values i and j and populate array  
            if x[i] > x[j] and strd_arr[i] < strd_arr[j] + 1:  
                strd_arr[i] = strd_arr[j]+1  
  
    # Determine the maximum from placeholder array  
    for i in range(x_len):  
        max_lis = max(max_lis, strd_arr[i])  
  
    # Return maximum  
    return max_lis
```

- If we add an input value of [1, 2, 3, 2, 5, 6, 1, 7], we will see the LIS have a length of 6 represented by the array of [1,2,3,5,6,7]
- Since we used a nested loop here, we will see the time complexity at $O(n^2)$

Resources:

[1] <https://northeastern.instructure.com/courses/117409/pages/module-9>

[2] https://en.wikipedia.org/wiki/Longest_increasing_subsequence

[3] Introduction to Algorithms, Cormen, Third Edition. (CLRS)

[4] https://en.wikipedia.org/wiki/Longest_common_subsequence_problem