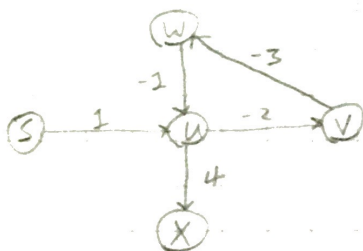


# Algorithms Illuminated Ch. 18

## Problem 18.1



Manually execute Bellman-Ford.

	S	u	v	w	x
	1	2	3	4	5
0	0	$\infty$	$\infty$	$\infty$	$\infty$
1	0	1	$\infty$	$\infty$	$\infty$
2	0	1	-1	$\infty$	5
3	0	1	-1	-4	5
4	0	-5	-1	-4	5
5	0	-5	-7	-4	-1

recurrence

$$AC[i][v] = \min($$

$$AC[i-1][v],$$

$$\min_{(u,v) \in E} \{AC[i-1][u] + l_{uv}\})$$

negative cycle, not stable

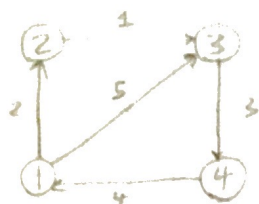
## Problem 18.2

This is not true. Problem 18.1, vertex u, provides a counter-example. The shortest path stays at 1 for  $1 \leq i \leq 3$ , but then decreases to -5.

## Problem 18.3

In this case, Bellman-Ford only requires  $K$  outer iterations. Each iteration must check every incoming edge for every node, which sums to  $m$ .  $C, O(Km)$

# Problem 18.4



Manually execute Floyd-Warshall

(base case)

Step 0:  $K=0$

Source

	1	2	3	4
1	0	2	5	$\infty$
2	$\infty$	0	1	$\infty$
3	$\infty$	$\infty$	0	3
4	4	$\infty$	$\infty$	0

SAR

Step 1:  $K=1$

	1	2	3	4
1	0	2	5	$\infty$
2	$\infty$	0	1	$\infty$
3	$\infty$	$\infty$	0	3
4	4	6	9	0

Step 2:  $K=2$

	1	2	3	4
1	0	2	3	$\infty$
2	$\infty$	0	1	$\infty$
3	$\infty$	$\infty$	0	3
4	4	6	7	0

Step 3:  $K=3$

	1	2	3	4
1	0	2	3	6
2	$\infty$	0	1	4
3	$\infty$	$\infty$	0	3
4	4	6	7	0

Step 4:  $K=4$

	1	2	3	4
1	0	2	3	6
2	8	0	1	4
3	7	9	0	3
4	4	6	7	0

## Problem 18.5

(base case)

Step 0:  $K=0$

	1	2	3	4
1	0	2	5	$\infty$
2	$\infty$	0	1	$\infty$
3	$\infty$	$\infty$	0	-3
4	-4	$\infty$	$\infty$	0

Step 1:  $K=1$

	1	2	3	4
1	0	2	5	$\infty$
2	$\infty$	0	1	$\infty$
3	$\infty$	$\infty$	0	-3
4	-4	-2	1	0

Step 2:  $K=2$

	1	2	3	4
1	0	2	3	$\infty$
2	$\infty$	0	1	$\infty$
3	$\infty$	$\infty$	0	-3
4	-4	-2	-1	0

Step 3:  $K=3$

	1	2	3	4
1	0	2	3	0
2	$\infty$	0	1	-2
3	$\infty$	$\infty$	0	-3
4	4	-2	-1	-4

Step 4:  $K=4$

	1	2	3	4
1				
2				
3				
4				

etc., I get the idem.