GB21802 - Programming Challenges

Week 6 - Math Problems

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Results for the Previous Week

Here are the results for last week:

Week 5: Graphs II
Deadline: 6/2/2017, 11:59:59 PM (1 day, 14:16 hours from now)
Problems Solved -- 0P:31, 1P:6, 2P:8, 3P:7, 5P:2, 8P:1,

#	Name	Sol/Sub/Total	My Status
1	Wormholes	24/26/55	
2	Meeting Prof. Miguel...	16/16/55	
3	Full Tank?	6/7/55	
4	Degrees of Separation	9/10/55	
5	Avoiding Your Boss	2/2/55	
6	Software Allocation	2/2/55	
7	Sabotage	1/1/55	
8	Gopher II	1/1/55	

Math problems in programming Competitions

There is a large variety of Math problems

- Programming is usually easy;
- "Solving" the problem is usually the hard part;
- BigNum, Rounding, Reducing fractions

A sample of math topics in programming challenges:

- Ad-hoc (Simulation, Probability)
- Number Theory (Primality, Divisibility, Modulo)
- Combinatorics (Closed forms, recurrences)

Ad-hoc Problems

What is ad-hoc?

ad-hoc means “single purpose”.

Each problem has a unique algorithm, using basic maths

Common categories:

- Simulation: Follow a (simple) formula and report the result;
- Find the Pattern: Start from the data and discover the background rule;

Ad Hoc Example: Probability problems

Dice Throwing

If you have n dice, what is the chance of rolling a total above m?

- Example: For $n = 3$, $m = 16$, what is the probability?

Ad Hoc Example: Probability problems

Dice Throwing

If you have n dice, what is the chance of rolling a total above m?

- Example: For $n = 3$, $m = 16$, the chance is $10/216$
- All combinations of 3 dice: $6 * 6 * 6 = 216$
- Combinations above 16:

- 6,6,6
- 6,6,5
- 6,5,6
- 5,6,6
- 6,5,5
- 5,6,5
- 5,5,6
- 4,6,6
- 6,4,6
- 6,6,4

- What algorithm do you use?

Ad Hoc example: Probabilty Problems

The dice problem

If I have n dice, what is the chance of rolling a total above m?

Solving with DP

- For $n = 0$, we have only one result: $r = 0$
- For $n = 1$, we have 6 results: $r = \{1, 2, 3, 4, 5, 6\}$
- The result for $n = i$ and $r_{n-1} = k$ is $r_n = k + \{1, 2, 3, 4, 5, 6\}$

- With a state table (dice,result), we can count the number of dice combination above a certain value;

Ad Hoc example: Probability Problems

Example Code

```
int count(int dice_left, int score_left) {
   if (score_left < 1) return 1;
   if (dice_left == 0) return 0;
   if (result[dice_left][score_left] != -1)
      return result[dice_left][score_left];
   int sum = 0;
   for (int i = 0; i < 6; i++)
      sum += count(dice_left-1, score_left-(i+1))
   result[dice_left][score_left] = sum;
   return sum;
}

prob = count(n,m)/6**n;
```

Dealing with big numbers

Many math problems require very large numbers, how do we deal with them?

- C++ unsigned int = unsigned long = 2^{32} (9-10 digits)
- C++ unsigned long long = 2^{64} (19-20 digits)
- Factorial of 23: $>$ 20 digits!

What do you do?

- C++: No bigint in standard library
- Python: Big integers supported natively
- Java: java.math.BigInteger

I usually recommend C++, but Java is better for BigInt problems!

Problem Example: 10925 – Krakovia

Problem Description: Sum *n* numbers, and divide by *f* friends. Just that... but the numbers are huge!

```
import java.util.Scanner;
import java.math.BigInteger;
class Main {
  public static void main(String[] args) {
    Scanner sc = new Scanner(System.in);
    int caseNo = 1;
    while (true) {
      int N = sc.nextInt(), F = sc.nextInt();
      if (N == 0 && F == 0) break;
      BigInteger sum = BigInteger.ZERO;     // Constant!
      for (int i = 0; i < N; i++) {
        BigInteger V = sc.nextBigInteger(); // I/O!
        sum = sum.add(V); }
      System.out.println("Bill #" + (caseNo++)
        + " costs " + sum + ": each friend should pay "
        + sum.divide(BigInteger.valueOf(F)) + "\n" );}
  }
}
```

More functions from Java.math.BigInteger

Algebraic functions

BigInteger.add(), .subtract(), .multiply(), .divide(), .pow(), .mod(), .remainder()

Output – Prints number in any base!

```
BI = BigInteger(10); System.println(BI.toString(2))
// Result: 1010
```

Primality – Probabilistic primality test

```
isPrime = BI.isProbablePrime(int certainty)
// Chance of being correct is 1 - (1/2)^certainty
```

Other cool functions

BigInteger.gcd(BI) BigInteger.modPow(BI exponent, BI m)

Number Theory

Number Theory studies the integer numbers and sets.

- Primality;
- Division and Remainders;
- Sequences of numbers;

Number Theory: Primality Testing

Prime Numbers: Only divisible by 1 and itself:

2,3,5,7,11,13...

How do you test if a number N is prime?

- Full search: For each $f \in 2..N-1$, test if $N\%f == 0$
  $O(N)$
- A little Pruning: For each $f \in 2..\text{floor}(\sqrt{N})$, test if $N\%f == 0$
  $O(\sqrt{(N)})$
- Can you do it in $O(\sqrt{n}/\log(n))$?

Number Theory: Primality Testing

The Prime Number Theorem (simplified)

The probability of $i < N$ is prime is $1/\log(N)$

collorary[1] **1**: There are $N/\log(N)$ primes $< N$
collorary 2: We just need to test the **primes** between 1 and $\sqrt{N}$

But how do we find all primes between 1 and $\sqrt{N}$ fast?

[1] "Collorary" means "consequence"

Sieve of Eratosthenes

Idea

- Start with a set from 2 to $\sqrt{N}$.
- Test if each i in the set is prime.
- If i is prime, remove all multiples mi.

```
def sieve(k):                        ## Find all primes up to k
    primes = []
    sieve = [1]*(k+1)      ## all numbers start in the list
    sieve[0] = sieve[1] = 0                  ## except 0 and 1
    for i in range(k+1):                              ## O(N)
        if (sieve[i] == 1):
            primes.append(i)                ## new prime found
            j = i*i    ## why can i start from i*i, not i*2?
            while (j < k+1):                     ## O(loglogN)
                sieve[j] = 0
                j += i                        ## next multiple
    return primes
```

Sieve of Eratosthenes

Amortized Complexity

- The complexity of the Sieve is $O(N \log \log N)$
- If we do the Sieve every time we test for primes, we are not saving much.
- But we can do the Sieve one time, and test many primes later!

When we do an expensive operation once, we call it **amortized complexity**

Finding Prime Factors

Any natural number N can be expressed as a unique set of prime numbers:

$$N = 1 p_1^{e_1} p_2^{e_2} \dots p_n^{e_n}$$

These are the Prime Factors of N. From this set, we can also obtain the set of Factors of N (all numbers i where $i|N$).

Factorization is a key issue in cryptography

Very Naive approach – Test all numbers!

For every $i \in 1..N/2$, test $i|N$ and isPrime(i).

Very Expensive!

Naive approach – Test all primes

Calculate a list of primes i up to N/2, test if $i|N$.

Wrong Answer, why?

Prime factorization: Divide and conquer approach

Recursive Idea

The prime factorization of N is equal to the union of p_i and the prime factorization of N/p_i, where p_i is the smallest prime factor of N.

The set of all factors is composed of all combinations of the set of prime factors (including repetitions).

```
def primefactors(n):
   primes = sieve(int(np.sqrt(n))+1)
   c = 0, i = n, factors = []
   while i > 1:
      if (i%primes[c] == 0):
          i = i/primes[c]
          factors.append(primes[c])
      else:
          c = c+1
   return factors
```

Working with Prime Factors: 10139 – Factovisors

Problem description

Calculate whether m divides $n!$ $(1 \leq m, n \leq 2^{31} - 1)$

Factorial of 22 is already bigint! But we can break down these numbers into their factors, which are all $\leq 2^{30}$.

- F_m: primefactors(m)
- $F_{n!}$: $\cup$(primefactors(1), primefactors(2)...,primefactors(n))

Having the factor sets, m divides $n!$ if $F_m \subset F_{n!}$.

Examples:

- $m = 48$ and $n = 6$
  $F_m = \{2, 2, 2, 2, 3\} F_{n!} = \{2, 3, 2, 2, 5, 2, 3\}$
- $m = 25$ and $n = 6$
  $F_m = \{5, 5\} F_{n!} = \{2, 3, 2, 2, 5, 2, 3\}$

Modulo Operation

We can use modulo arithmetic to operate on very large numbers without calculating the entire number.

Remember that:

1. $(a + b)\%s = ((a\%s) + (b\%s) + s)\%s$
2. $(a * b)\%s = ((a\%s) * (b\%s))\%s$
3. $(a^n)\%s = ((a^{n/2}\%s) * (a^{n/2}\%s) * (a^{n\%2}\%s))\%s$

Modulo Operation – UVA 10176, Ocean Deep!

Problem summary

Test if a binary number n (up to 100000 digits) is divisible by 131071

- The problem wants to know if $n\%13107 == 0$
- But n is too big!
- Use the recurrence in the previous slide to break down each digit to a reasonable value.

Euclid Algorithm and Extended Euclid Algorithm

- Euclid Algorithm gives us the greatest common divisor D of a, b;
- Extended Euclid Algorithm also gives us x, y so that $ax + by = D$;
- Both are extremely simple to code:

```
int gcd(int a, int b) {return (a == 0?b:gcd(b%a,a));}

int x, y;
int egcd(int a, int b) {
   if (a==0)
      {x = 0; y = 1; return b;}         // stop condition
   int d = egcd(b%a, a);
   int tx = x;                          // gcd recurrence
   x = y - (b/a)*tx; y = tx; return d; }    // update x,y
```

Using EGCD: The Diophantine Equation

Problem Example (variations of this problem are common)

You have 839 yen. Xhoco candy costs 25 yen, Yanilla candy costs 18 yen. How many candies can we buy?

The equation $xA + yB = C$ is called the Linear Diophantine Equation. It has infinite solutions if GCD(A,B)|C, but none if it does not.

The first solution (x_0, y_0) can be derived from the extended GCD, and other solutons can be found from: expressed as:

- $x = x_0 + (b/d)n$
- $y = y_0 - (a/d)n$

Where d is GCD(A,B) and n is an integer.

Using EGCD: The Diophantine Equation

Problem Example (variations of this problem are common)

You have 839 yen. Xhoco candy costs 25 yen, Yanilla candy costs 18 yen. How many candies can we buy?

- EGCD gives us: $x = -5, y = 7, d = 1$ or $25(-5) + 18(7) = 1$
- Multiply both sides by 839: $25(-4195) + 18(5873) = 839$
- So: $x_n = -4195 + 18n$ and $y_n = 5873 - 25n$
- We have to find n so that both x_n, y_n are > 0.
- $-4195 + 18n \geq 0$ and $5873 - 25n \geq 0$
- $n \geq 4195/18$ and $5873/25 \geq n$
- $4195/18 \leq n \leq 5873/25$
- $233.05 \leq n \leq 234.92$

Combinatorics problems

Definition

Combinatorics is the branch of mathematics concerning the study of countable discrete structures.

Combinatory problems involve understanding a sequence, and figuring one of:

- Recurrence: A formula that calculates the n^{th} member of a sequence, based on the value of previous members;
- Closed form: A formula that calculates the n^{th} member of a sequence independently from other members;

It is not uncommon to use Dynamic Programming or Bignum to solve combinatoric related problems.

Example: Triangular Numbers

Definition

The triangular numbers is the sequence where the n^{th} value is composed of the sum of all integers from 1 to n

- S(1) = 1
- S(2) = 1+2 = 3
- S(3) = 1+2+3 = 6
- ...
- S(7) = 1+2+3+4+5+6+7 = 28

What are the recurrence and the closed form for this sequence?

Example: Triangular Numbers

- S(1) = 1, S(2) = 3, S(3) = 6

Recurrence

The recursive form of a sequence:

$$S(n) = S(n-1) + n; S(1) = 1$$

Closed Form

The non-recursive form of a sequence:

$$S(n) = \frac{n(n+1)}{2}$$

Problem: Calculate the first triangle number with more than 500 factors!

A more famous sequence: Fibonacci Numbers

Definition – very famous sequence

Each number is the sum of the two numbers before it.

F() = 0,1,1,2,3,5,8,13,21,34...

The recurrence is well known

$$F(0) = 0, F(1) = 1, F(n) = F(n-1) + F(n-2)$$

When implementing the recurrence, don't forget the memoization table!

Closed Form

The Fibonacci numbers also have a less well known closed form:

$$F(n) = \frac{1}{\sqrt{5}}\left(\left(\frac{1+\sqrt{5}}{2}\right)^n - \left(\frac{1-\sqrt{5}}{2}\right)^n\right)$$

Square roots introduce floating point errors. What is the maximum n this can calculate with less than 0.1 error?

Fibonacci Facts

Zeckendorf's theorem

Every positive integer can be written in a unique way as a sum of one or more distinct fibonacci numbers, which are not consecutive.

```
def zeckenfy(n):
    fibs = []
    f = greatest fib =< n; fibs.append(f)
    fibs.append(zeckenfy(n-f))
    return fibs
```

Pisano's period

The last digits of the Fibonacci sequence repeat!

The last one/two/three/four digits repeat with a period of 60/300/1500/15000.
F(6) = 8
F(66) = 27777890035288
F(366) = 1380356705549181797202918793682511
3333650564850089197542855968899086435571688

Binomial Coefficients

Definition

Binomial Coefficients are the number series that correspond to the coefficients of the expansion of a binomial:

Binom(3) = $(a+b)^3$ = $1a^3 + 3ab^2 + 3ab^2 + b^3$ = $\{1, 3, 3, 1\}$

We are usually interested in the k^{th} coefficient of the n^{th} binomial:

$C(n,k) = C(3,2) = \{1, 3, 3, 1\} = 3$

Pascal's Triangle gives us a good representation of C(n,n):

0	1	0	0	0	0	0	0	0	0
0	1	1	0	0	0	0	0	0	0
0	1	2	1	0	0	0	0	0	0
0	1	3	3	1	0	0	0	0	0
0	1	4	6	4	1	0	0	0	0
0	1	5	10	10	5	1	0	0	0
0	1	6	15	20	15	6	1	0	0
0	1	7	21	35	35	21	7	1	0
0	1	8	28	56	70	56	28	8	1

Uses for the Binomial Coefficient

The value of $C(n, k)$ tells us how many ways we can choose n items, k at a time.

Some use cases:

- Probabilities: What is the probability of winning a loto when you choose 5 numbers out of 60? $1/C(60, 5)$
- Grids: How many ways are there to go from the bottom left end of a *mn* grid to the top right, if you can only go up and right? $C(m + n, n)$

Calculating the Binomial Coefficient

Closed form of C(n,k)

$$C(n,k) = \frac{n!}{(n-k)!k!}$$

Problem: Multiplying factorials tends to generate huge numbers even for small n and k.

Recurrence for C(n,k)

- C(n,0) = C(n,n) = 1;
- C(n,k) = C(n-1,k-1) + C(n-1,k)

Using a memoization table will cut the calculation time by half. In this case, top-down DP will usually be faster than bottom-up.

Another useful sequence: Catalan Numbers

The Catalan sequence

$$C(n) = 1, 1, 2, 5, 14, 42, 132, 429, 1430$$

The Recurrence

$$C(n) = \sum_{k=0}^{n-1} C(k)C(n-1-k)$$

Closed Form

$$C(n) = \frac{1}{n+1}\binom{2n}{n}$$

Catalan Numbers – Uses

- Number of ways that you can match n parenthesis.
  C(3):((())),()(()),(())(),()()(),(()())
- Number of ways that you can triangulate a poligon with $n+2$ sides
- Number of monotonic paths on an *nxn* grid that do not pass above the diagonal.
- Number of distinct binary trees with n vertices
- Etc...

Integer Partition

f(5,5) = (5),(4,1),(3,2),(3,1,1),(2,2,1),(2,1,1,1),(1,1,1,1,1)

Definition and calculation

$f(n,k)$ – number of ways that we can sum n, using integers equal or less than k.

Recurrence:

- $f(n,k) = f(n-k,k) + f(n,k+1)$
- $f(1,1) = 1$; $f(n,k) = 0$ if $k > n$

Class Summary

- Math Problems
- Java's Big Integer class
- Primality
- Modulo arithmetic
- GCD and Diophantine Equations
- Combinatorics

Next week: Geometry problems!

This Week's Problems

- How Many Trees?
- Dice Throwing
- Self-describing Sequence
- Triangle Counting
- Summation of Four Primes
- Divisibility of Factors
- Marbles
- Winning Streak

To Learn More

Euler Project: Mathematical questions using computers:
`http://projecteuler.net`