

# Programming Challenges

## Week 9 - Geometry

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# Introduction

## Geometry

Problems related to points, lines, angles and circles.  
Usually there will be more than one way to calculate.

## Geometrical Constructs

Lines, Segments, Planes, Circles, Convex Polygons, Concave Polygons, etc...

## Geometrical Computing

Our main concern, are **degeneracies** and **instability**

# Degeneracies

## Numerical Instability

$$\arcsin(\sin(\pi/4)) \neq \pi/4$$

Don't forget that operations with real numbers are not guaranteed to be precise;

## Degeneracies

Special cases for geometric calculations. Normally caused by divisions by zero. But sometimes have other sources.

$$\tan(\pi/2) = \frac{\sin(\pi/2)}{\cos(\pi/2)} = \frac{1}{0}$$

# Line (1)

## Characteristics

- Infinite;
- Divide a plan into two;
- Segment – a limited line

## Line (2)

### Representation – Two Points

- A line can be described by two points;
- $(x_0, y_0), (x_1, y_1)$

### Problems with this representation

- Not unique: We can have two identical lines represented by different points
- Calculating extra points requires interpolation;

# Line (3)

## Representation – Point and angle

$$y = mx + b$$

- $m$  (slope):  $\frac{y_1 - y_0}{x_1 - x_0}$
- $b$  (y-intercept): the point where  $x = 0$ ;

Problem: When the line is vertical, we have a degeneration (division by zero on the slope)

# Line (4)

## Representation – Point and angle 2

$$ay + bx + c = 0$$

or

$$x = c$$

When the line is vertical

```
p2l(double[] p1, double[] p2):  
    if (p1[0] == p2[0]): // vertical line  
        l.a = 1;  
        l.b = 0;  
        l.c = -p1[0];  
    else:  
        l.b = 1;  
        l.a = -(p1[1]-p2[1]) / (p1[0]-p2[0]);  
        l.c = -(l.a*p1[0]) - (l.b*p1[1]);
```

# Line Intersection (1)

## Line Intersection

We can calculate if two lines are parallel quickly, by checking if their inclination is the same. Note the Epsilon!

```
parallelQ(line l1, line l2):  
    return ((abs(l1.a-l2.a) <= EPSILON)&&  
            (abs(l1.b-l2.b) <= EPSILON))
```



# Line Intersection (2)

## Line Intersection Point

If the lines are not parallel, they have one intersection point.

$$x = \frac{b_2 - b_1}{m_1 - m_2}, y = m_1 \frac{b_2 - b_1}{m_1 - m_2} + b_1$$

```
intersection_point(line l1, line l2):
    if (!(parallelQ(l1,l2))):
        p[0] = (l2.b*l1.c - l1.b*l2.c)/
                (l2.a*l1.b - l1.a*l2.b);
        if (abs(l1.b) > EPSILON): // Vertical?
            p[1] = - (l1.a*(p[X])+l1.c)/l1.b;
        else:
            p[1] = - (l2.a*(p[X])+l2.c)/l2.b;
```

## Line Intersection (3)

### Angle between two lines

Two non parallel lines will always intersect at a given angle. If the lines are in the  $ax + by + c = 0$  format, we can calculate their angles as follows:

```
intersection_angle(line l1, line l2):  
    num = l1.a*l2.b - l2.a*l1.b  
    den = l1.a*l2.a * l1.b*l2.b  
    return(tan(num/den))
```

# Line Intersection (4)

## Closest Point

- The closest point  $p_l$  in a line  $l$  to point  $p$ , is the point where the line  $(p, p_l)$  intersects  $l$ .
- Closest point can be used to find the distance between a line and a point ( $d(p, p_l)$ );
- Degenerate/Easy cases:  $p$  is in  $l$ ,  $l$  is vertical,  $l$  is horizontal;
- The slope  $m$  of the line  $(p, p_l)$  is  $\frac{1}{l.a}$ ;
- Calculate the intersection between  $(p, p_l)$  and  $l$ ;

# Line Segments (1)

Line segments are lines delimited by start and end points;

```
typedef struct {  
    point p1,p2  
} segment;
```

## Line Segments (2)

### Degenerative Cases

- Are the Segments in the same line? (test for same points)
  - Are the Segments parallel? (no intersection)
- 
- Calculate the intersecting point between the lines.
  - Test if this point is within a rectangle defined by each line segment.

# Triangles (1)

- Polygon defined by three line segments;
- Characterized by the relationship between its angles and the line segment sizes;
- Commonly used to represent more complex polygons;

## Manipulating angles

- Angles can be represented by radians (0 to  $2\pi$ ) or degrees (0 to 360);
- Mixing the two of them is an easy way to insert bugs in your code;
- Make sure what is the usual input for your library's trigonometric functions;

# Triangles (2)

## Basic Triangle Facts

- Three angles, summing to a total of 180 degrees ( $\pi$  radians);
- Law of sines (A,B,C are angles; a,b,c are opposite edges):

$$\frac{a}{\sin A} = \frac{b}{\sin B} = \frac{c}{\sin C}$$

- Law of Cosines:

$$a^2 = b^2 + c^2 - 2bc \cos(A)$$

# Triangles (3)

## Right Triangles

A right triangle has one angle with 90 degrees ( $\pi/2$  radians). It has many neat properties;

For  $\alpha$  a non-right angle, with an *opposite* side and an *adjacent* side;

- $\cos(\alpha) = \frac{|adjacent|}{|hypotenuse|}$
- $\sin(\alpha) = \frac{|opposite|}{|hypotenuse|}$
- $\tan(\alpha) = \frac{|opposite|}{|adjacent|}$



# Triangles (4)

## Common problems with triangles

- Given two angles and a side, find the rest;
- Given two sides and an angle, find the rest;
- Given a side and a height, find the rest;
- Etc;

# Triangles (5)

## Area of a Triangle

$a$  is the altitude,  $h$  is the base;

$$A(T) = (1/2)ah$$

## Signed area

$a, b, c$  are the points of a triangle.

$$(a_x b_y - a_y b_x + a_y c_x - a_x c_y + b_x c_y - c_x b_y)/2$$

- Negative signed area:  $a, b, c$  are clockwise;
- Positive signed area:  $a, b, c$  are counterclockwise;
- Zero signed area:  $a, b, c$  are collinear;

# Circle 1

## Representation

- Center point and radius;
- Three boundary points;

## Measures

- Area:  $\pi r^2$
- Circumference:  $2\pi r$

# Circle 2

## Intersection between line and circle

Radius  $r$  and distance between center and line  $d$ ;

- $d > r$  – no intersection;
- $d == r$  – tangent, one intersection;
- $d < r$  – two intersection points;

## Intersection between two circles

- Two circles will intersect if the distance to their centers  $\leq r_1 + r_2$
- The points of intersection form triangles with determined sides. Angles and coordinates can be calculated as needed.

# Polygons

## Definition

Let's define a polygon as a closed chain of non-intersecting line segments. We can represent polygons by listing the  $n$  vertices in order around its boundary.

```
typedef struct {  
    int n;  
    point p[MAXPOLY]  
} polygon
```

- We can represent the “last” segment by  $(p[(n-1)\%n], p[n\%n])$

# Convex Polygons

## Definition

A polygon  $P$  is convex if any line segment defined by two points within  $P$  are contained in  $P$

- All internal angles in a convex polygon must be  $< \pi$  radians;
- The sum of all angles in a convex polygon is  $2\pi$ ;
- We can test a polygon by convexity by checking that all its angles turn to the same side. (ccw a,b,c)

# The Convex Hull

The convex hull is a basic algorithm often used to organize unstructured data.

# The Graham Scan

## Simple algorithm to create a convex hull

- select leftmost and lowest point as starting points;
- sort points by angle direction from the starting points;
- add the first point to the hull, and repeat.

How to avoid degeneracy? (Wrap around, collinear points, repeated points)



# Area of a polygon

## Convex Polygon

Add all signed triangular areas: The negative areas will compensate the positives.

## Concave Polygon

Picasso Algorithm: Remove “ears” (triangles) from the polygon, adding to the total area.

# Testing if a point is inside a polygon

# Problems

- Dog and Gopher
- Rope Crisis in Ropeland
- Herding Frosh
- Chainsaw Massacre