```
class Solution {
public:
    int maxProfit(vector<int>& prices) {
        int n = prices.size();
        int hold[n];
        int not_hold[n];
        memset(hold, 0, sizeof(hold));
        memset(not_hold, 0, sizeof(not_hold));
        hold[0] = -prices[0];
        for (int i = 1; i < n; i++) {
            hold[i] = max(not_hold[i - 1] - prices[i], hold[i - 1]);
            not_hold[i] = max(not_hold[i - 1], hold[i - 1] + prices[i]);
        }
        return not_hold[n - 1];
    }
};
```

Propostion: hold[i] and not_hold[i] is always the maximum profit at ith day. hold[i] means the maximum profit you can earn at ith day if you hold a stock at ith day. not_hold[i] means the maximum profit you can earn at ith day if you do not hold a stock at ith day.

If the above proposition is proof correct, then we could just output $not_hold[n - 1]$ cause it is the maximum profit we can earn at last day.

Proof by Induction:

Basecase:

when i == 0, if you buy a stock, then you need to pay prices[0], hold[0] = -prices[0], it is obvious the max profit you could earn. And if you don't buy, then of course the profit is 0, $not_hold[0] = 0$.

Induction Step:

Suppose the proposition is correct for hold[k] and $not_hold[k]$, then at k + 1 day, there will be two cases:

- 1. case 1: you hold a stock at k + 1 day. It can be the result by 1. you buy a stock at k + 1 day. 2. you hold the stock at k day, and you don't sell it at k + 1 day. Because hold[k] and not_hold[k] is the maximum profit at k day, then the maximum profit is the maximum profit for case 1 is not_hold[k] prices[k + 1], case 2 is hold[k]. The maximum profit at k + 1 day is the larger one among the 2 cases. I.e. hold[k + 1] = max(not_hold[k] prices[k + 1], hold[k]). The algorithm is correct.
- 2. case 2: you do not hold a stock at k + 1 day. It can be the result by 1. you sell a stock at k + 1 day. 2. you do not hold the stock at k day, and you don't buy one at k + 1 day. Because hold[k] and not_hold[k] is the maximum profit at k day, then the maximum profit is the maximum profit for case 1 is hold[k] + prices[k + 1], case 2 is not_hold[k]. The maximum profit at k + 1 day is the larger one among the 2 cases. I.e. not_hold[k + 1] = max(not_hold[k] + prices[k + 1], not_hold[k]). The algorithm is correct.

Therefore hold[k + 1] and $not_hold[k + 1]$ is the maximum profit at k + 1 day. The proposition is correct.

Time complexity:

We traverse the prices vector for one time, so it is O(n) where n is the length of the prices.