

$$T(n) = 2T\left(\frac{2}{3}n\right) + n^2$$

Approach 1.

$$\begin{aligned}
 T(n) &= 2T\left(\frac{2}{3}n\right) + n^2 \\
 &= 2 \cdot \left(2T\left(\frac{4}{9}n\right) + \frac{4}{9}n^2\right) + n^2 \\
 &= 4T\left(\frac{4}{9}n\right) + \left(1 + \frac{8}{9}\right)n^2 \\
 &= 8T\left(\frac{8}{27}n\right) + \left(1 + \frac{8}{9} + \frac{64}{81}\right)n^2 \\
 &= 16T\left(\frac{16}{81}n\right) + \left(1 + \frac{8}{9} + \frac{64}{81} + \frac{512}{729}\right)n^2 \\
 &= \dots \log_2^{\frac{3}{2}} n \cdot T(1) + 9 \cdot \left(1 - \left(\frac{8}{9}\right)^{\log_2^{\frac{3}{2}} n - 1}\right)n^2 \\
 &\leq 2^{\frac{\log_2 n}{\log_2 \frac{3}{2}}} T(1) + 9n^2 \\
 &= O(n^2)
 \end{aligned}$$

Approach 2.

$$T(n) = 2T\left(\frac{2}{3}n\right) + n^2$$

$$a=2 \quad b=\frac{3}{2} \quad f(n)=n^2$$

$$\log_b a = \log_{\frac{3}{2}} 2$$

For case 3 in Master Theorem.

$f(n) = \Omega(n^{\log_{\frac{3}{2}} 2 + \epsilon})$ when $0 < \epsilon < 2 - \log_{\frac{3}{2}} 2$
so $f(n)$ grows polynomially faster than $n^{\log_{\frac{3}{2}} 2}$

and $2f\left(\frac{2}{3}n\right) \leq C f(n)$

$$\Leftrightarrow \frac{8}{9}n^2 \leq Cn^2 \text{ when } C \geq \frac{8}{9}$$

Therefore, $T(n) = \Theta(f(n)) = \Theta(n^2) = O(n^2)$