

# CHISELVERIFY: A VERIFICATION LIBRARY FOR CHISEL

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# OUTLINE

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- Background: Verification & Chisel.
- Current solutions: Current ways to verify chisel designs.
- Motivation: Need for Verification tools for Chisel.
- Our solution: bringing verification to Chisel directly in Scala.
- Evaluation: Comparing verification with ChiselVerify to verification with UVM.

# BACKGROUND

# BACKGROUND: VERIFICATION

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- Verification of digital systems:
  - Testing a design before it has been taped out.
  - Verification tools allow to verify that the tests we are writing “test the right things”.
- Universal Verification Methodology (UVM):
  - Testing methodology implemented in SystemVerilog, a popular solution for verification.
  - Allows for the writing of standardised test-benches, which can be reused on multiple designs.

# BACKGROUND: VERIFICATION TOOLS

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- Functional Coverage:
  - Qualitative and fine-grained coverage metric.
  - Gives progress on “which features have been tested?”.
- Constrained Random Verification (CRV):
  - Enables the creation of random variables which generate values that satisfy a set of constraints.
  - Allows for the generation of constantly valid inputs, making a small test case capable of covering many of the component’s functionalities.

# BACKGROUND: CHISEL

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- Chisel:
  - **Hardware Construction Language** embedded in Scala.
  - Allows for the use of **Object-Oriented** and **Functional** programming constructs in the scope of hardware description.
  - Compiles through an intermediate representation called **FIRRTL** before generating output **Verilog**.
  - Can be either **synthesised** or **simulated** using:
    - Output **Verilog** (e.g. with **verilator**)
    - Intermediate **FIRRTL** with **Treadle** (a custom simulator and testing engine).

# CURRENT SOLUTIONS

# OVERVIEW OF THE CURRENT SOLUTIONS

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- For Chisel:

- **ChiselTest**: “traditional” test-benches with peek-poke-expect interfaces and forking, lacks verification features.
- **ScalaTest**: Software testing framework, not ideal for hardware, doesn’t simulate the hardware, only checks the Chisel code itself.

- For Verilog:

- **SystemVerilog**: Extension of Verilog that enables object oriented programming and verification features inside of the test-benches.
- **UVM**: Verification Methodology, enables a standardised testing method that can be reused for many different DUTs.



# CURRENT SOLUTIONS: WHAT'S MISSING?

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- Other types of solutions:
  - RFuzz: Fuzzer for RTL circuits. Only allows fuzzing, no verification tools.
  - Chisel Formal: Formal Verification in Chisel, static verification, doesn't solve the same problems.
- What's missing in these solutions:
  - **ChiselTest**: For test-benches, not really verification.
  - **ScalaTest**: Not made for Hardware.
  - **SystemVerilog & UVM**:
    - Too verbose, requires ~800 LOC for a test-bench.
    - Requires multiple languages to test a Chisel design.

# MOTIVATION

# MOTIVATION: CURRENT LACK OF VERIFICATION TOOLS

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- So far Chisel doesn't have any “native” tools that allow the easy use of verification functionalities.
- Current solutions all require external tools to function.
- We need tools which are built for Chisel.
- We also want tools which are **efficient** to use and **easy to learn**, not the case of current popular solution like SystemVerilog with UVM.

**OUR SOLUTION: CHISELVERIFY**

# CHISELVERIFY: OVERVIEW

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- Hardware Verification library for Chisel, inspired by UVM.
- Powered entirely by Scala and ChiselTest.
- ChiselVerify brings the following to the Chisel ecosystem:
  - Functional Coverage
  - Constrained Random Verification
  - Bus Functional Models
  - Timed Assertions

# CHISELVERIFY: FUNCTIONAL COVERAGE

# FUNCTIONAL COVERAGE: WHAT IS IT?

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## Verification Plan

Representation of the DUT's expected features.

### CoverGroup

Set of DUT ports that will be sampled together, represents a "feature".

### CoverPoint

Set of values that a port is expected to have to verify a feature, represents a feature of single port.

#### Bins

Definition of a set of values that a port should reach during testing, done in two ways:

**Range:** first to last  
(both included)

**Conditional:**  
Arbitrary  
(**portValues**) =>  
**Boolean** function

#### Cross

Defines a relation between two bins in a CoverPoint, i.e. how many value pairs, within the defined cross set, have these two points reached during testing.

*Example:*

*Cross(A, B, 1 to 1, 1 to 1) =>  
Have ports A and B reached the value 1 at the same time?*

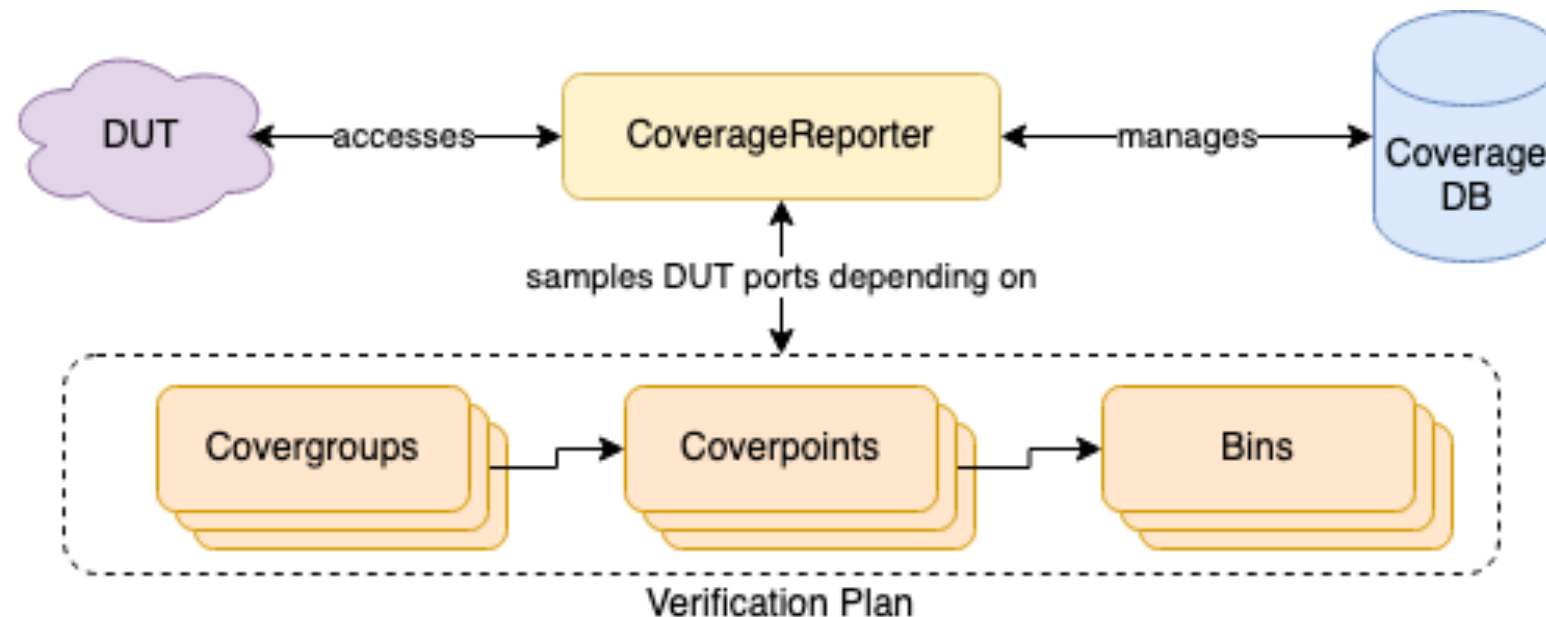
#### Timed

Same idea as the cross bin, but with an added time constraint, i.e. have the two points hit a cross set within a given number of cycles from each other?

There are 3 ways to define a time constraint:  
Always, Eventually and Never

# FUNCTIONAL COVERAGE: HOW DID WE DO IT?

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- **CoverageDB:** DataBase that maintains the values gathered for all of the bins across multiple tests in a test suite.
- **Coverage Reporter:** Handles the registration of CoverPoints and Bins to the DB, samples the bin values and creates the coverage report.
  - This is used to create the verification plan.



# FUNCTIONAL COVERAGE: API

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## Two main constructs:

- Cover: Represents a feature of the DUT. Can be represented by both a single and multiple DUT ports. Its behavior is defined by the **bins** it contains.

```
case class cover(pointName: String, ports: Data*)  
  def apply(b: Bin*)  
  def apply(delay: DelayType)(b: Bin*)
```

- Bins: Definition of a part of a feature. Defined using value ranges or condition, and can contain timing information.

```
def bin(name: String, range: Option[Range] = None, condition: Option[Seq[BigInt] => Boolean] = None, expectedHits: BigInt = 0)
```

# FUNCTIONAL COVERAGE: USING IT

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- Create the coverage reporter and verification plan.

```
val cr = new CoverageReporter(dut)
cr.register(
    cover("accu", dut.io.outA)( //CoverPoint 1
        bin("lo10", 0 until 10), bin("First100", 0 until 100)),
    cover("test", dut.io.outB)( //CoverPoint 2
        bin("testLo10", 0 until 10)),
    //Declare cross points
    cover("accuAndTest", dut.io.outA, dut.io.outB)(
        cross("both1", Seq(1 to 1, 1 to 1)))
)
```

- Sample the CoverPoints inside of the test.

```
cr.sample()
```

# RESULT: FUNCTIONAL COVERAGE REPORT

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- Create and print the coverage report

```
//Generate report  
val report = cr.report  
print(report.serialize)
```

- Example result:

```
===== COVERAGE REPORT =====  
===== GROUP ID: 1 =====  
COVER_POINT PORT NAME: accu  
BIN lo10 COVERING Range 0 until 10 HAS 10 HIT(S) = 100,00%  
BIN First100 COVERING Range 0 until 100 HAS 50 HIT(S) = 50,00%  
=====  
COVER_POINT PORT NAME: test  
BIN testLo10 COVERING Range 0 until 10 HAS 4 HIT(S) = 40,00%  
=====  
CROSS_POINT accuAndTest FOR POINTS accu AND test  
BIN both1 COVERING Range 1 to 1 CROSS Range 1 to 1 HAS 1 HIT(S) = 100,00%  
=====
```

# **CHISELVERIFY: CONSTRAINED RANDOM VERIFICATION**

# CONSTRAINED RANDOM VERIFICATION: WHAT IS IT?

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- Model tests using randomness:
  - Give random inputs to the DUT and expect values using a golden model.
- Guide the randomness using constraints:
  - We don't want the randomness to be uniformly distributed => use constraints to describe the random distribution.
- Idea: Add a Constraint Programming Language to Scala/Chisel
  - Create Random Objects, then define Random Variables and Constraints in it.
  - Random Objects define a **Constraint Satisfaction Problem (CSP)**. This problem is solved using a CSP Solver.
  - We use **JaCoP** as a CSP Solver.

# CONSTRAINED RANDOM VERIFICATION: HOW DID WE DO IT?

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- Same ideas as in SV:
  - Random objects created by extending the **RandObj** trait using a given **Model** with a seed.
  - Random fields are added as **Rand/Randc** values inside the class.
  - Constraints are defined using either:
    - Single constraints: using “#” (e.g. `val lenConstraint = len #> 2` )
    - **ConstraintGroups** which are equivalent to SV

*Example:*

```
class Frame extends RandObj(new Model) {  
  val pkType: Rand = new Rand(0, 3)  
  val len: Rand = new Rand(0, 10)  
  val noRepeat: Randc = new Randc(0, 1)  
  
  val legal: ConstraintGroup = new ConstraintGroup {  
    len #>= 2  
    len #<= 5  
  }  
}
```

# CONSTRAINED RANDOM VERIFICATION: WHAT DO WE ADD?

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- The list of operator used to construct constraint is the following: `#<`, `#<=`, `#>`, `#>=`, `#=`, `div`, `#*`, `mod`, `#+`, `-`, `#\=`, `#^`, `in`, `inside`.
- We also added conditional constraints using:
  - **IfCon**: If a condition is met, then use the constraint
  - **ElseC**: If a condition is met, then use the constraint, else use an other.

*Example:*

```
val constraint1: crv.Constraint = IfCon(len #= 1) {  
    payload.size #= 3  
} ElseC {  
    payload.size #= 10  
}
```

# CONSTRAINED RANDOM VERIFICATION: USING IT

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- `dist` can also be used to define custom distributions on a random variable by associating ranges to weights using the `:=` operator.
- The `randomize()` method returns `true` only if the CSP Solver found a set of values that satisfy the current constraints.

*Example:*

```
val myPacket = new Frame(new Model)
assert(myPacket.randomize)
```

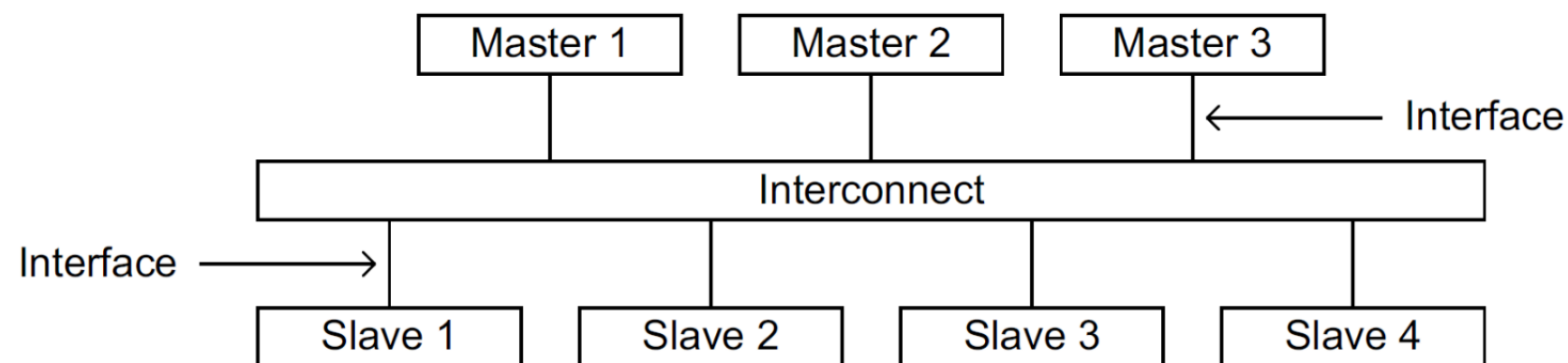


# **CHISELVERIFY: BUS FUNCTIONAL MODELS**

# BUS FUNCTIONAL MODEL: WHAT IS IT?

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- Abstraction of the inner workings of a standardised interface.
  - Allows for the use of data transfer via **Transactions**, rather than having to deal with the inner wiring manually.
  - Software abstraction, is useful for faster verification.
- 
- We chose to create a first BFM for the AXI4 Bus:



# BUS FUNCTIONAL MODEL: HOW TO USE IT

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- Create a **FunctionalMaster** for your AXI interfaced DUT.

```
val master = new FunctionalMaster(dut)
```

- Create transactions:

- Write:

```
master.createWriteTrx(0, Seq(42), size = 2)

var resp = master.checkResponse()
while (resp == None) {
    resp = master.checkResponse()
    dut.clock.step()
}
```

- Read:

```
master.createReadTrx(0, size = 2)

var data = master.checkReadData()
while (data == None) {
    data = master.checkReadData()
    dut.clock.step()
}
```

# CHISELVERIFY: TIMED ASSERTIONS

# TIMED ASSERTIONS: QUICK OVERVIEW

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- Allows to create predicated assertions that take into account certain timing delays.
- **Types of delays:** Given a delay of  $x$  cycles:
  - **Exactly:** Assertion is true exactly in  $x$  cycles.
  - **Eventually:** Assertion is true at least once in the next  $x$  cycles.
  - **Always:** Assertion is true every cycle for the next  $x$  cycles.
  - **Never:** Assertion isn't true at any cycle for the next  $x$  cycles.
- Two types of Timed Assertions:
  - **ExpectTimed:** Uses *ChiselTest*'s `expect` for the assertion.
  - **AssertTimed:** Uses a software `assert(_)` for the assertion.

# TIMED ASSERTIONS: USING IT

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```
def exact[T <: Module](d: Int, msg: String = s"EXACTLY ASSERTION FAILED")
    (cond: => Boolean)(implicit dut: T): Unit =

def never[T <: Module](d: Int = 100, msg: String = s"NEVER ASSERTION FAILED")
    (cond: => Boolean)(implicit dut: T): Unit =

def eventually[T <: Module](d: Int = 100, msg: String = s"EVENTUALLY ASSERTION FAILED")
    (cond: => Boolean)(implicit dut: T) : Unit =

def always[T <: Module](d: Int = 100, msg: String = s"ALWAYS ASSERTION FAILED")
    (cond: => Boolean)(implicit dut: T): Unit =
```

*Example:*

```
eventually(2, "aEqb expected timing is wrong") { dut.io.a.peek() === dut.io.b.peek() }
exact(2, "aEqb expected timing is wrong") { dut.io.a.peek() === dut.io.b.peek() }
always(2, "aEqb expected timing is wrong") { dut.io.a.peek() === dut.io.b.peek() }
never(2, "aEqb expected timing is wrong") { dut.io.a.peek() === dut.io.b.peek() }
```

# CONCLUSION

# CONCLUSION

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- ChiselVerify brings verification to the Chisel ecosystem.
- High-Level Functional backend (i.e. Scala) allows for much more efficiency during verification process (in comparison to SystemVerilog with UVM).
- Can be used to verify non-Chisel designs as well thanks to Chisel Blackboxes.



# REFERENCES

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- *Xilinx AXI reference guide: [https://www.xilinx.com/support/documentation/ip\\_documentation/axi\\_ref\\_guide/latest/ug1037-vivado-axi-reference-guide.pdf](https://www.xilinx.com/support/documentation/ip_documentation/axi_ref_guide/latest/ug1037-vivado-axi-reference-guide.pdf)*

# GETTING STARTED USING CHISELVERIFY

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- Current Project repository:

<https://github.com/chiselverify/chiselverify/>

- Project Wiki (Good way to get started):

<https://github.com/chiselverify/chiselverify/wiki/>

- ChiselVerify is published on Maven. To use it, add following line to your `build.sbt`:

```
libraryDependencies += "io.github.chiselverify" % "chiselverify" % "0.1"
```

# QUESTIONS?