#### Håkan Sundell, Philippas Tsigas

# Fast and lock-free concurrent priority queue for multi-thread systems

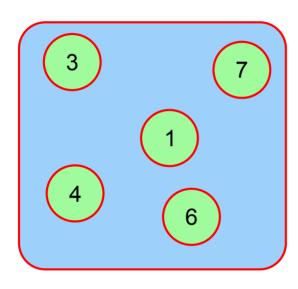
Presented by Andreas Haas
University of Salzburg
Concurrency and Memory Management Seminar
2010-12-16

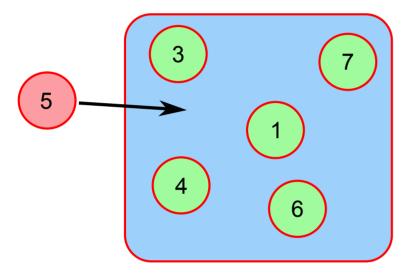
#### Overview

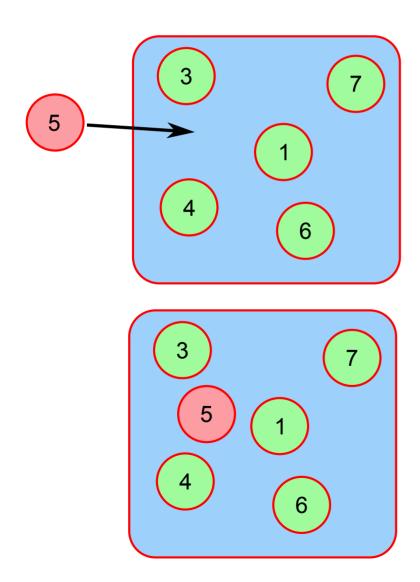
- Introduction
- Priority queue
- Skip list
- Non-blocking algorithm
- Correctness

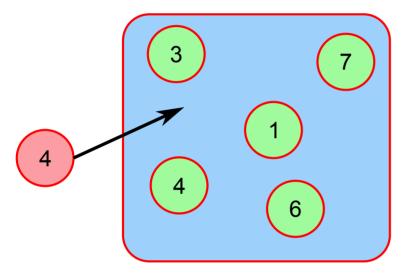
## Non-blocking

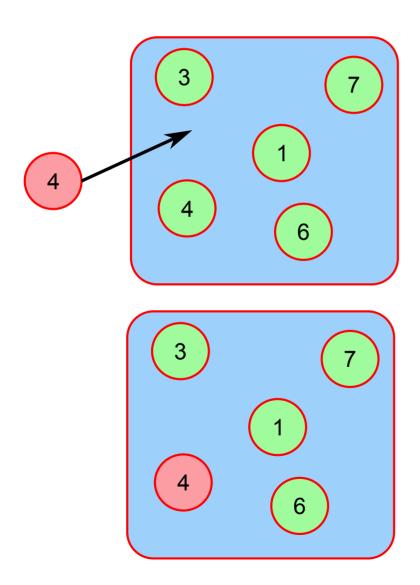
- No mutual exclusion
- Guaranteed progress of at least one operation
- Atomic operations:
  - Test-And-Set (TAS)
  - Fetch-And-Add (FAA)
  - Compare-And-Swap (CAS)



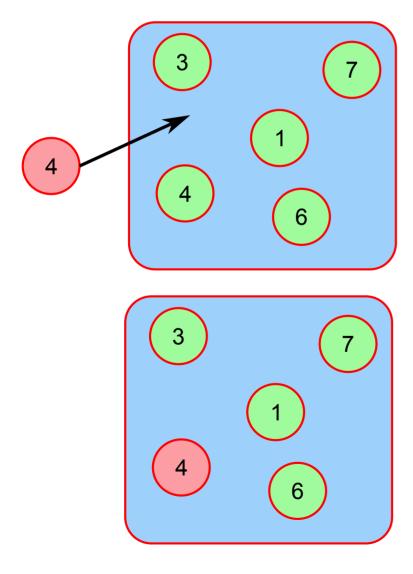




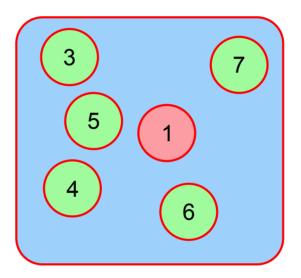




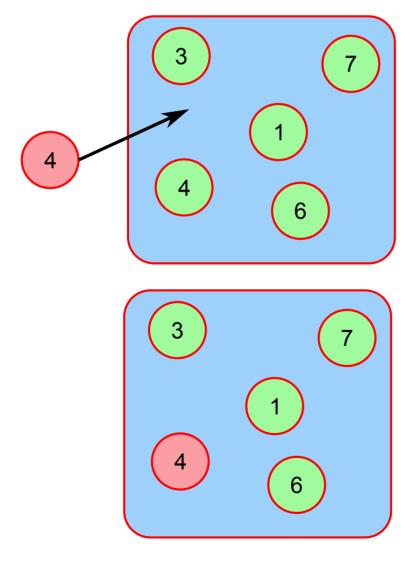
#### Insert



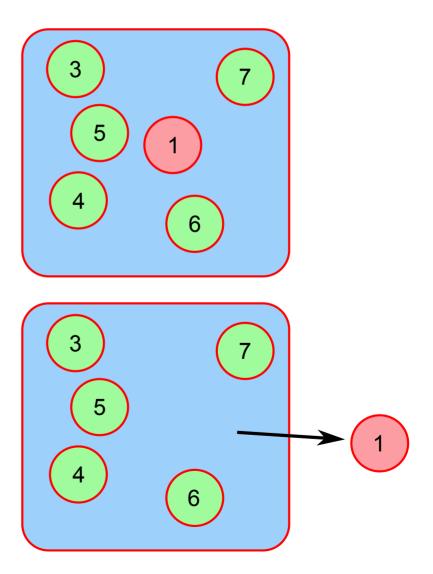
#### DeleteMin



#### Insert



#### **DeleteMin**

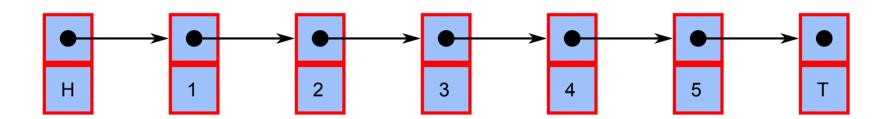


- 4
- **Applications**
- Scheduling
- O Discrete Simulation
- o Dijkstra's algorithm
- ...
- Common implementations
  - Sorted list
  - Unsorted list
  - Heap
  - Binary search tree

# Skip list

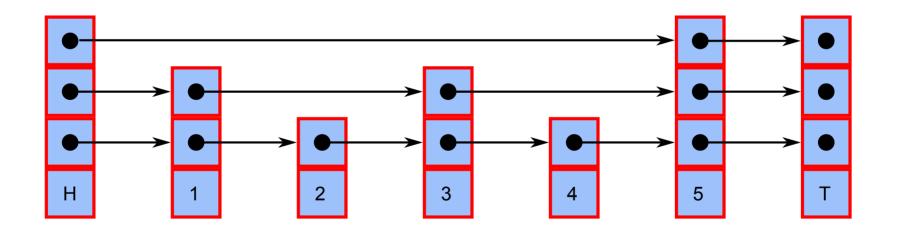


Singly-linked list

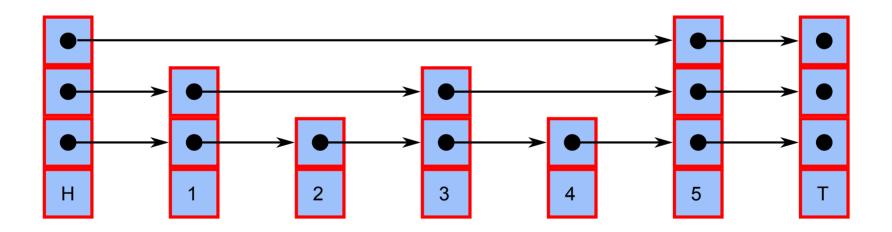


# Skip list

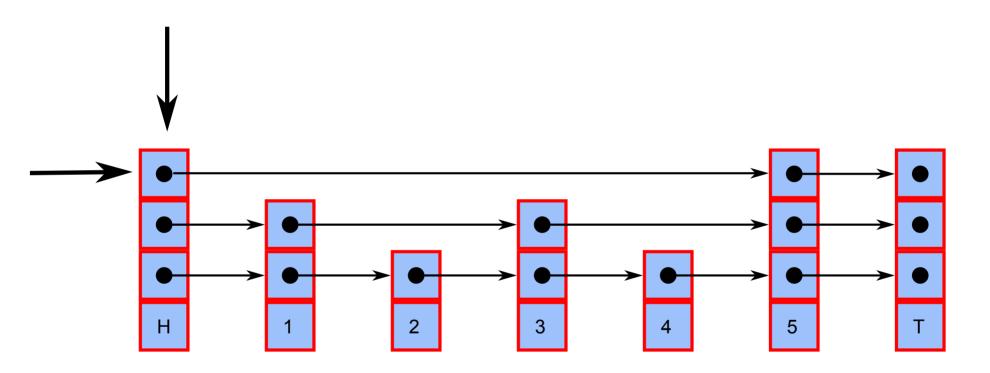
- Singly-linked list
- With shortcuts



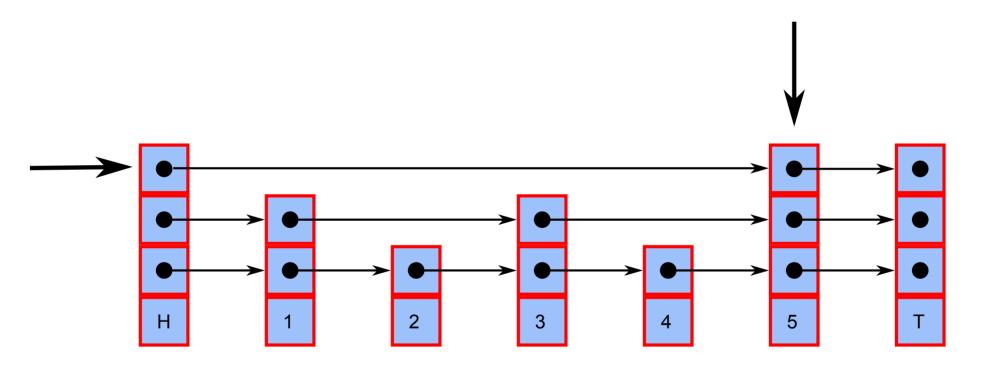




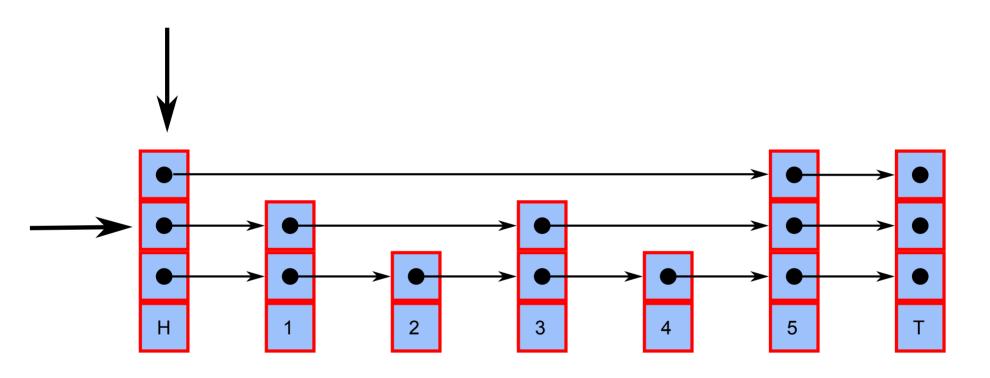




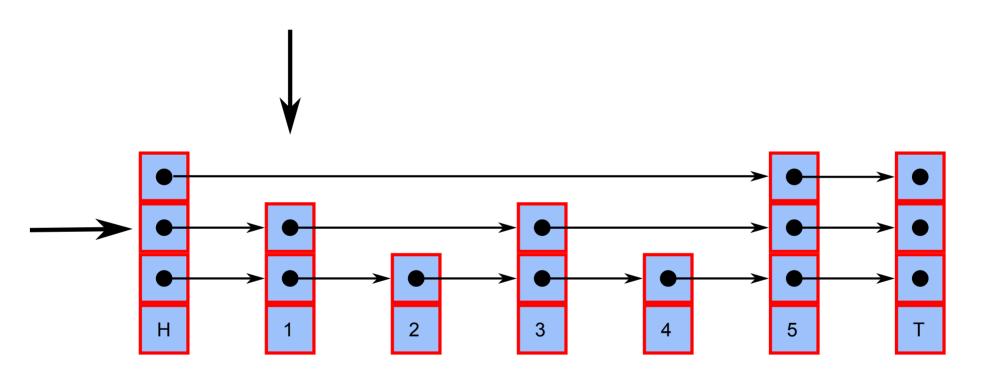




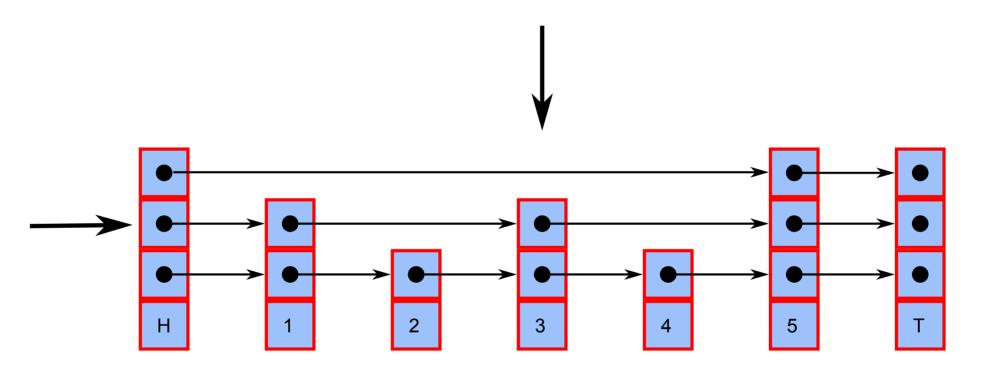




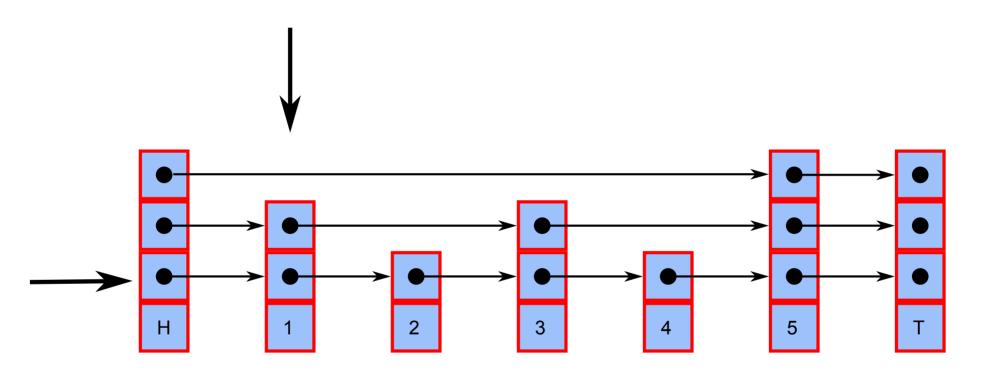




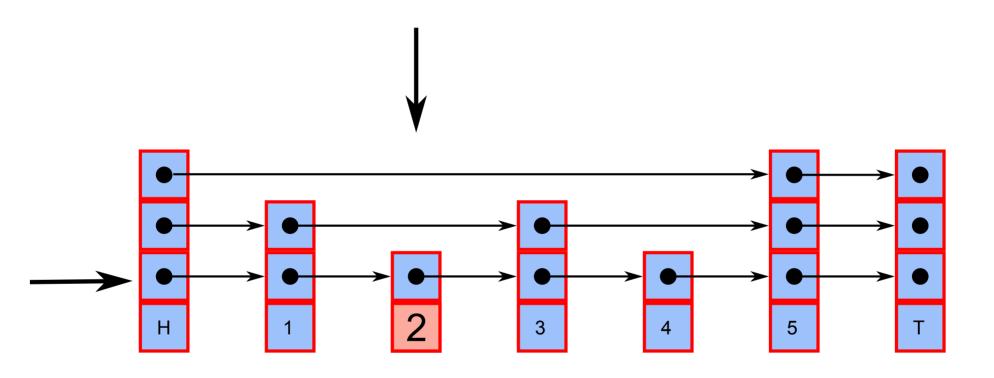




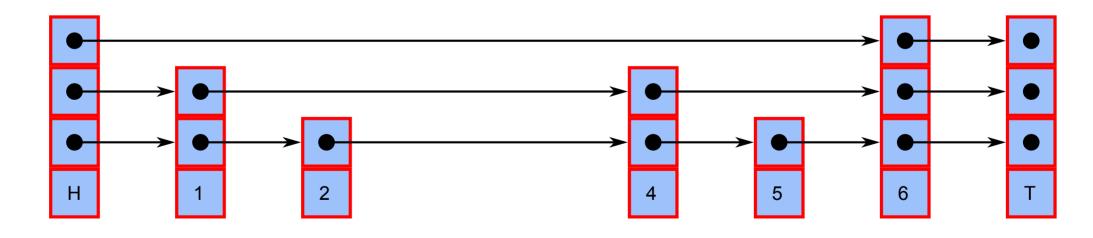


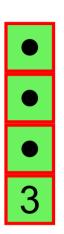




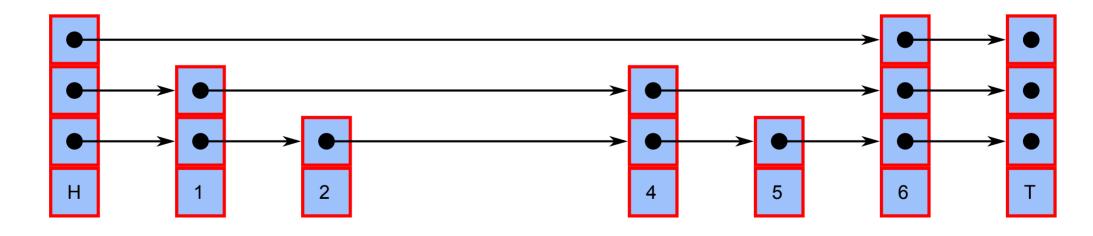


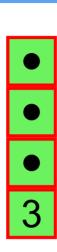
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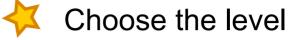




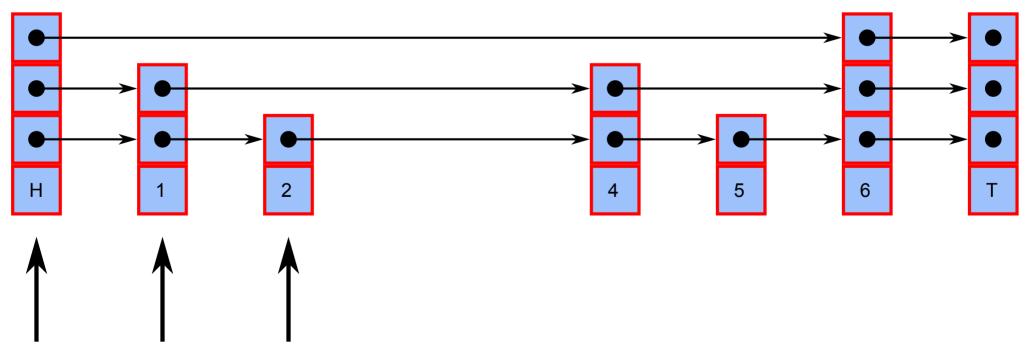


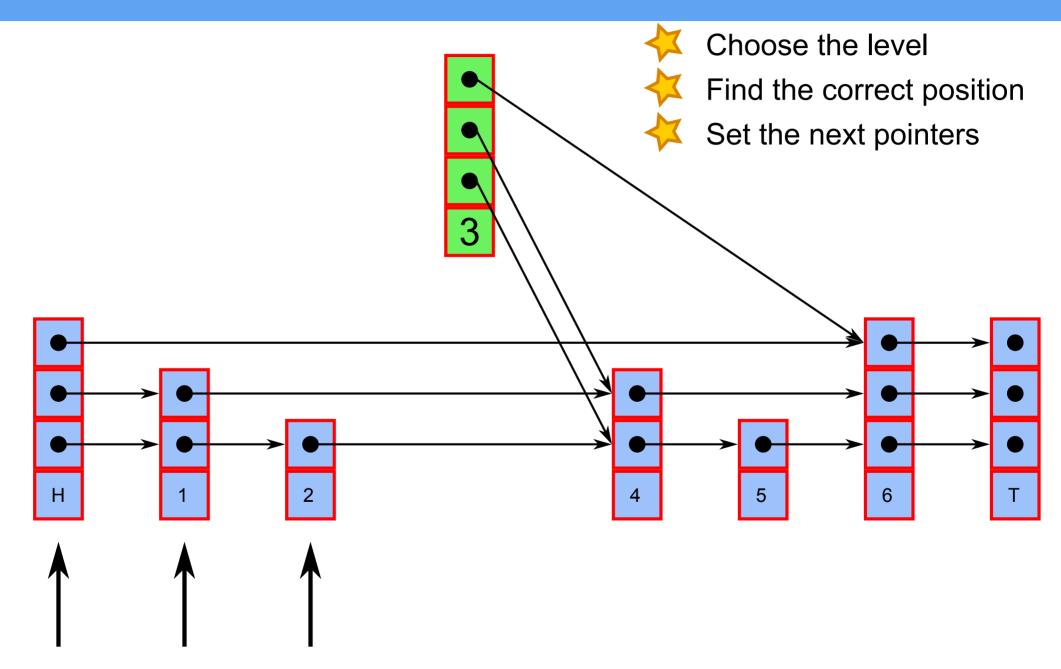


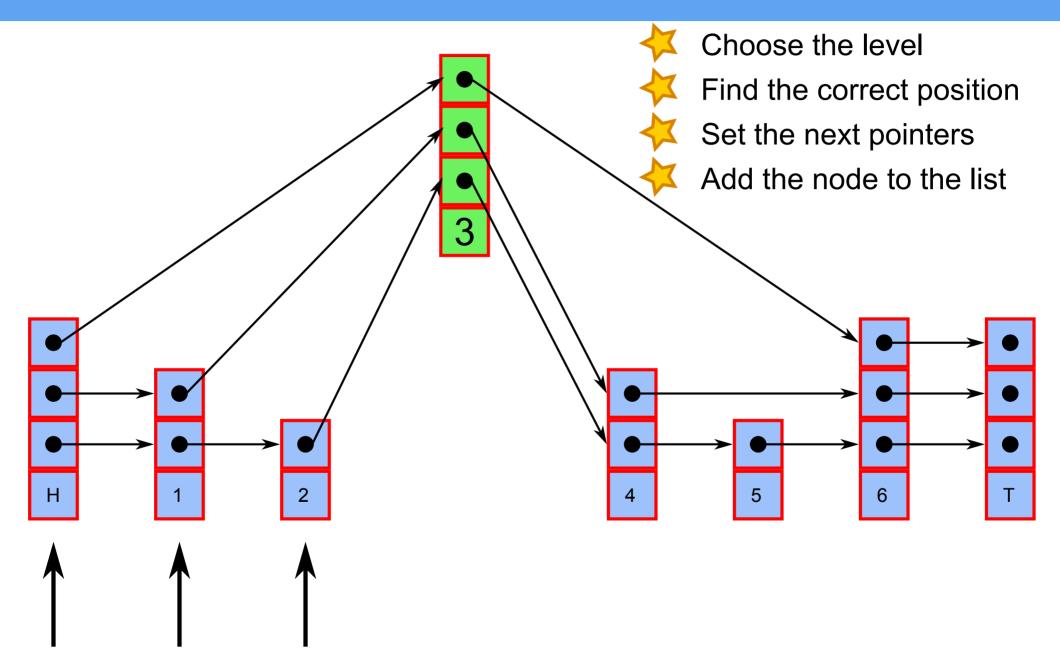


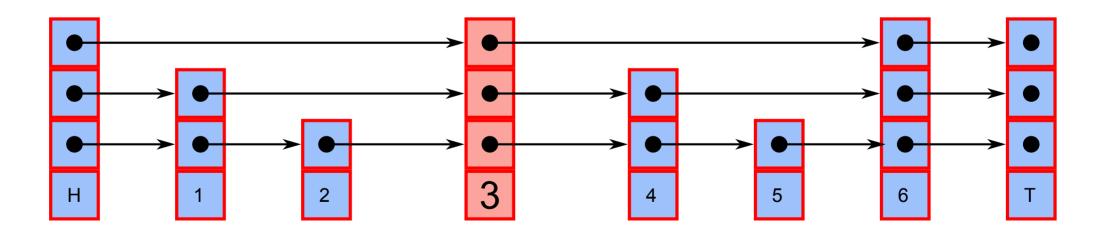






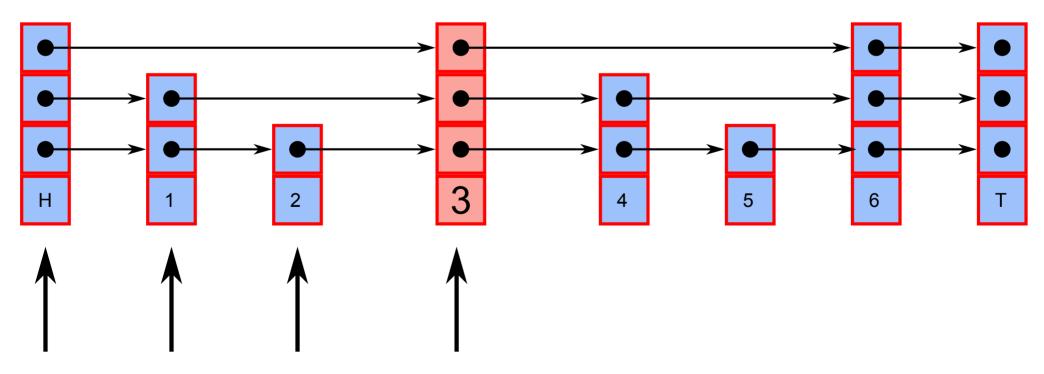


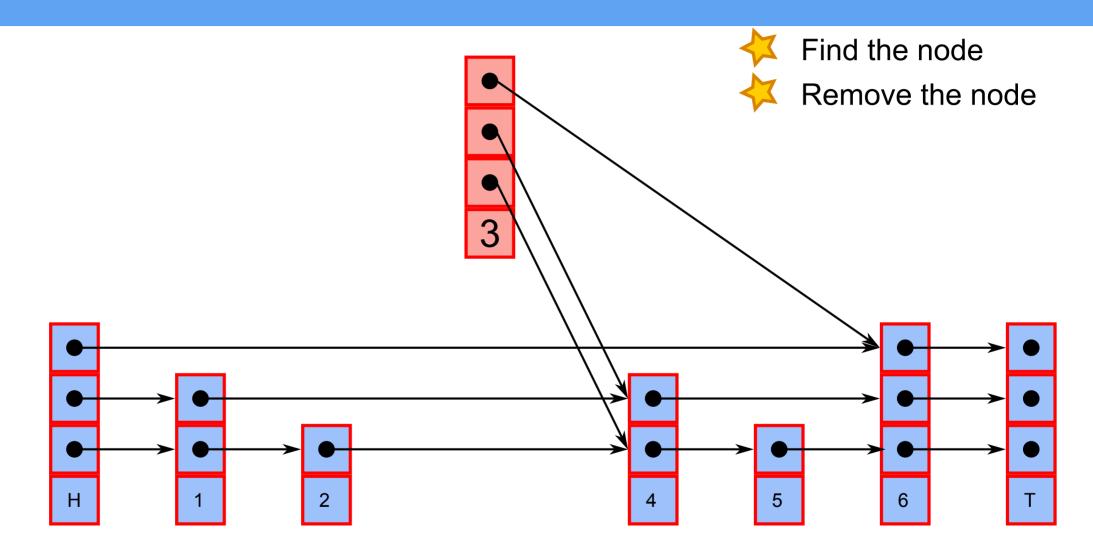






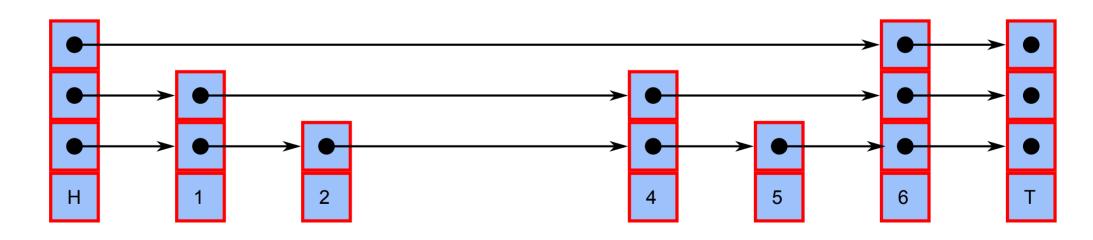
Find the node





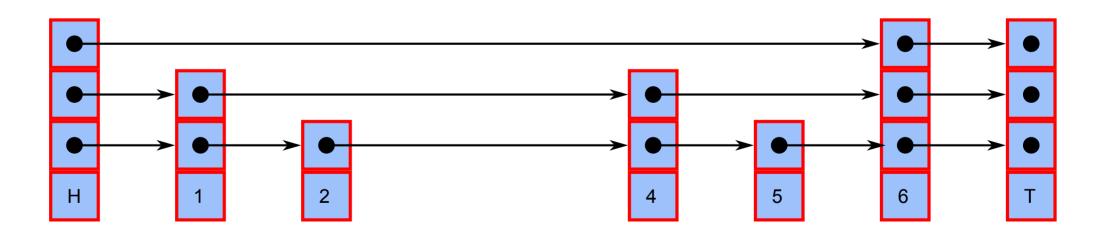
Find the node
Remove the node
Delete the node

3



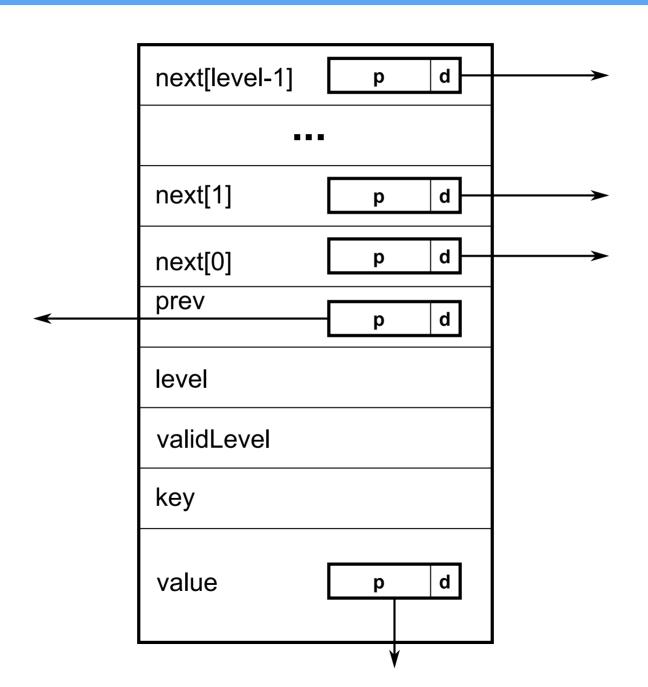
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Find the node
Remove the node
Delete the node
Return the value



## Non-blocking algorithm

#### Node Structure



#### Memory management

- No node should be reclaimed and then later re-allocated while some other process is traversing that node
  - ABA Problem
  - Reference Counting is used

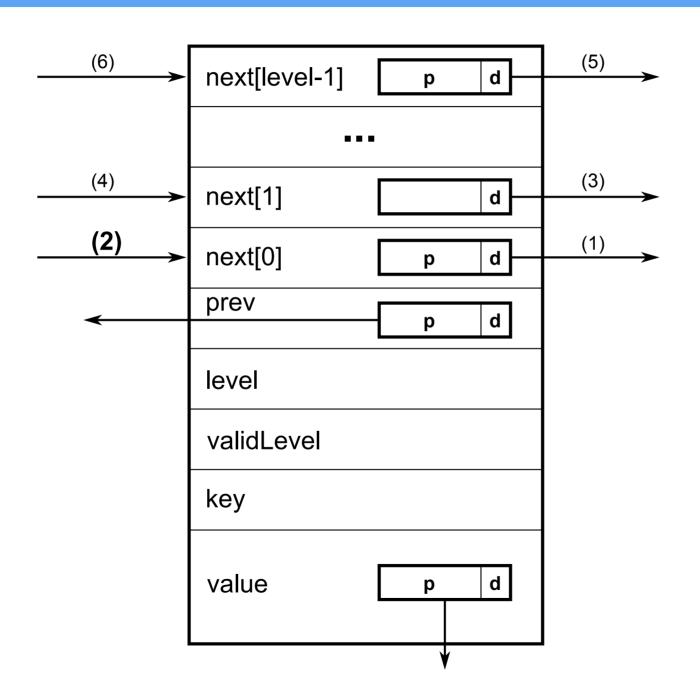
#### Key insight

- Next pointers have to be changed consistently, but not necessarily all at once (atomically)
  - A node is interpreted to be inserted when inserted at the lowest level
- Use a deletion mark to mark a node that is about to be deleted
  - A marked node is interpreted as deleted
  - © Every thread removes all marked nodes it encounters
    - Helping strategy

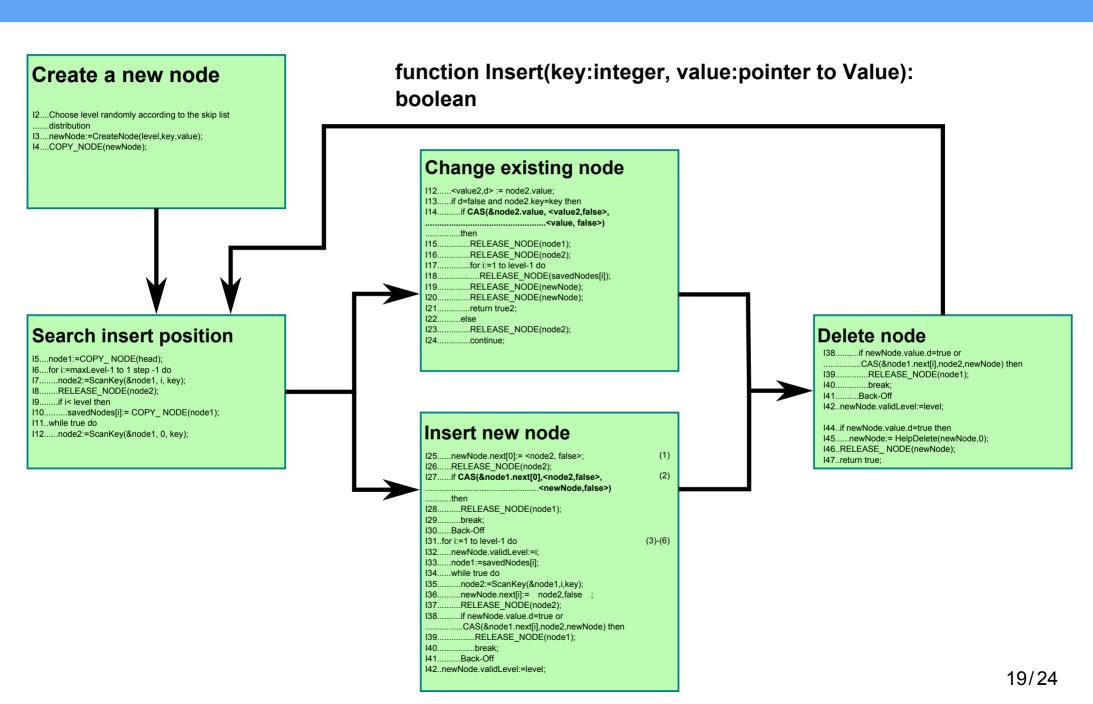
#### Non-blocking find node

- While traversing the skip list, all nodes are deleted which are flagged with the deletion mark.
  - Helping to achieve non-blocking

### Non-blocking insert



### Non-blocking Insert



# Create a new node

```
I2....Choose level randomly according to the skip list ......distribution
I3....newNode:=CreateNode(level,key,value);
I4....COPY_NODE(newNode);
```

# Search insert position

```
I5....node1:=COPY_ NODE(head);
I6....for i:=maxLevel-1 to 1 step -1 do
I7......node2:=ScanKey(&node1, i, key);
I8......RELEASE_NODE(node2);
I9......if i< level then
I10......savedNodes[i]:= COPY_ NODE(node1);
I11..while true do
I12.....node2:=ScanKey(&node1, 0, key);
```

## Change existing node

```
I12.....<value2,d> := node2.value;
I13.....if d=false and node2.key=key then
114.....if CAS(&node2.value, <value2,false>,
                .....<value, false>)
 .....then
I15.....RELEASE NODE(node1);
I16.....RELEASE NODE(node2);
117.....for i:=1 to level-1 do
I18.....RELEASE NODE(savedNodes[i]);
I19.....RELEASE NODE(newNode);
I20.....RELEASE NODE(newNode);
I21....return true2;
122.....else
I23.....RELEASE NODE(node2);
124.....continue;
```

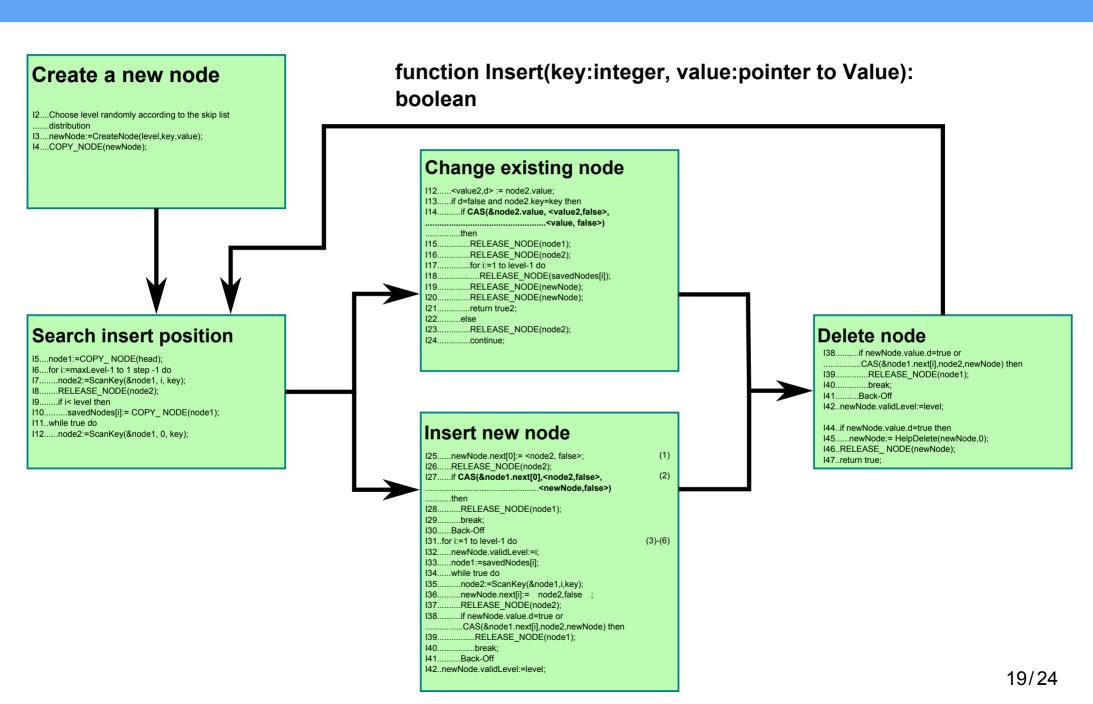
#### Insert new node

```
(1)
125.....newNode.next[0]:= <node2, false>;
I26.....RELEASE_NODE(node2);
                                                      (2)
127.....if CAS(&node1.next[0],<node2,false>,
  .....<newNode,false>)
then
128.....RELEASE_NODE(node1);
129.....break:
I30.....Back-Off
                                                   (3)-(6)
131..for i:=1 to level-1 do
I32.....newNode.validLevel:=i;
I33.....node1:=savedNodes[i];
134.....while true do
I35.....node2:=ScanKey(&node1,i,key);
I36.....newNode.next[i]:= node2,false ;
I37.....RELEASE NODE(node2);
I38.....if newNode.value.d=true or
.....CAS(&node1.next[i],node2,newNode) then
I39.....RELEASE NODE(node1);
I40.....break;
I41.....Back-Off
I42..newNode.validLevel:=level;
```

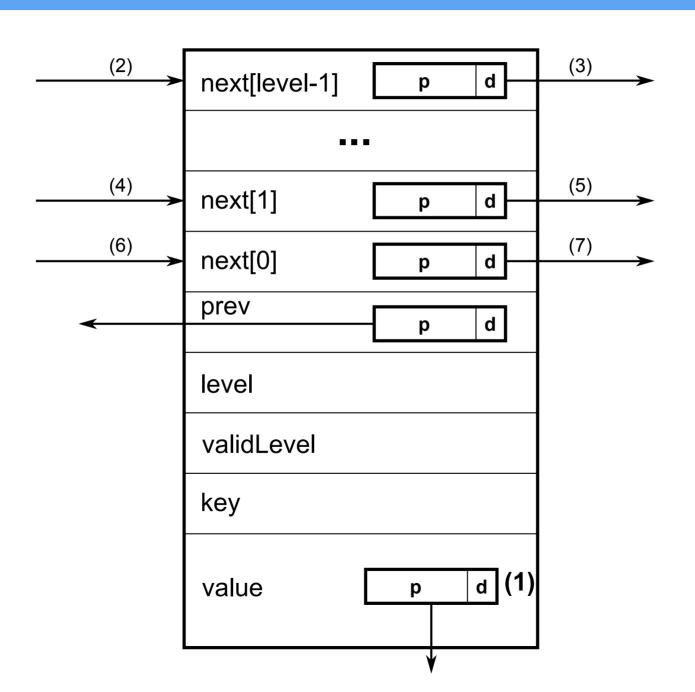
### Delete node

```
138.....if newNode.value.d=true or
......CAS(&node1.next[i],node2,newNode) then
I39.....RELEASE NODE(node1);
I40.....break;
I41.....Back-Off
I42..newNode.validLevel:=level;
144..if newNode.value.d=true then
I45.....newNode:= HelpDelete(newNode,0);
I46..RELEASE NODE(newNode);
147..return true;
```

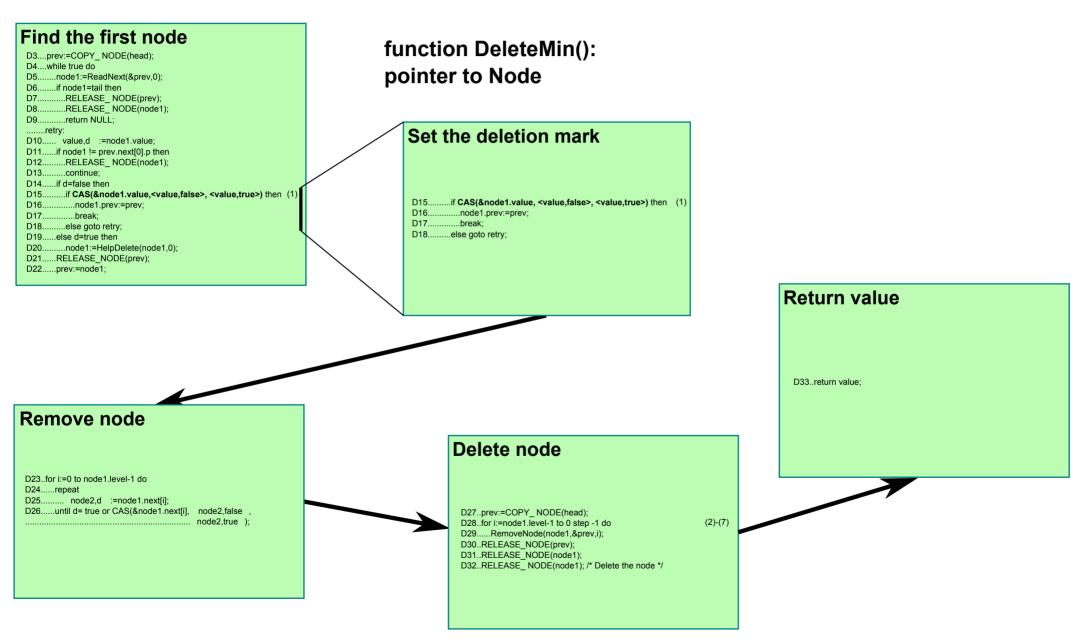
### Non-blocking Insert



#### Non-blocking DeleteMin



### Non-blocking DeleteMin



#### Find the first node

```
D3....prev:=COPY_ NODE(head);
D4...while true do
D5.....node1:=ReadNext(&prev,0);
D6.....if node1=tail then
D7.....RELEASE NODE(prev);
D8.....RELEASE NODE(node1);
D9.....return NULL;
....retry:
D10..... value,d :=node1.value;
D11.....if node1 != prev.next[0].p then
D12.....RELEASE_ NODE(node1);
D13.....continue;
D14.....if d=false then
D15.....if CAS(&node1.value,<value,false>, <value,true>) then (1)
D16.....node1.prev:=prev;
D17.....break;
D18.....else goto retry;
D19.....else d=true then
D20.....node1:=HelpDelete(node1,0);
D21.....RELEASE_NODE(prev);
D22.....prev:=node1;
```

#### Set the deletion mark

```
D15.....if CAS(&node1.value, <value, false>, <value, true>) then (1) D16......node1.prev:=prev; D17......break; D18.....else goto retry;
```

#### Remove node

```
D23..for i:=0 to node1.level-1 do
D24.....repeat
D25...... node2,d :=node1.next[i];
D26.....until d= true or CAS(&node1.next[i], node2,false , node2,true );
```

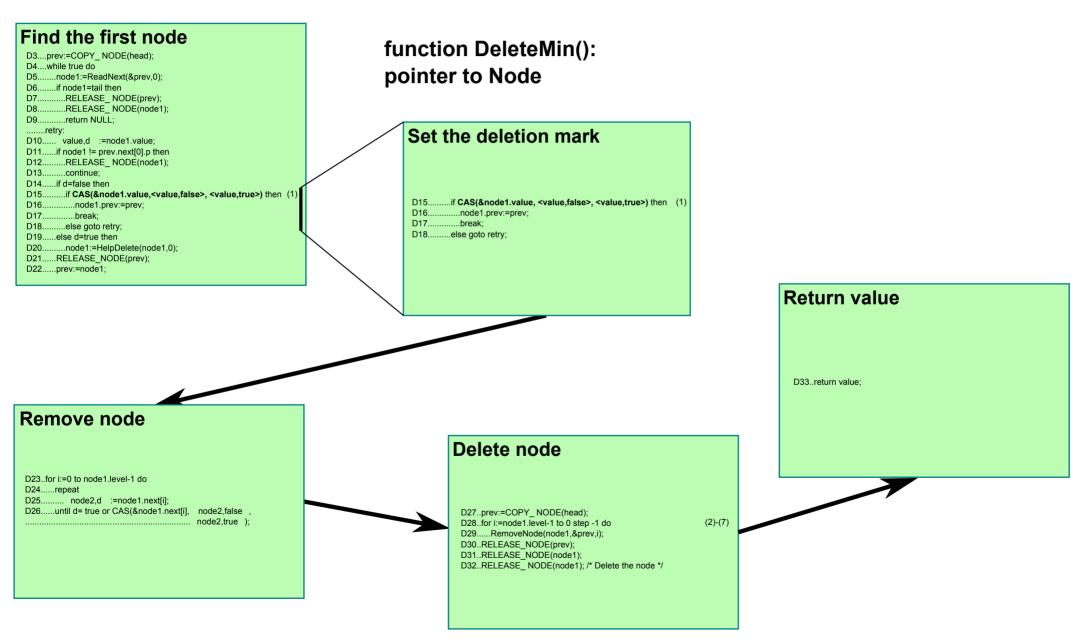
#### Delete node

```
D27..prev:=COPY_ NODE(head);
D28..for i:=node1.level-1 to 0 step -1 do (2)-(7)
D29......RemoveNode(node1,&prev,i);
D30..RELEASE_NODE(prev);
D31..RELEASE_NODE(node1);
D32..RELEASE_NODE(node1); /* Delete the node */
```

#### Return value

D33..return value;

### Non-blocking DeleteMin



#### Correctness

- Changes to the skip list are done in a single CAS instruction
  - Insert: The next pointer pointing to the new node
  - Oblighted Delete Min: The deletion mark of the node with the minimal key is set
- The data structure is always in a consistent state
- Linearizable since changes occur atomically

#### Conclusion

- Priority Queue based on a sorted linked list
  - Use shortcuts (skip list) as optimization
    - The shortcuts are not part of the consistent state of the priority queue
- Delayed deletion using deletion marks
- Helping for non-blocking
- Correctness is achieved by changing the consistent state of all threads by single CAS operations

#### Thank you for your attention



Håkan Sundell, Philippas Tsigas: Fast and lock-free concurrent priority queue for multi-thread systems, JPDC 65, 2005