Digital Image Processing (CSE 478) Lecture 19-20: Image Compression

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Motivation

- Consider a 2 hour, full HD video (resolution of 1920 × 1080)
- The storage space required per frame :1920 \times 1080 \times 24 bits = 6.22 MB
- Space required per second: $1920 \times 1080 \times 24 \times 30$ bits
- Space required for entire movie: $1920 \times 1080 \times 24 \times 30 \times 2 \times 60 \times 60$ bits = $1920 \times 1080 \times 3 \times 30 \times 2 \times 60 \times 60$ bytes = 1.34×10^{12} bytes= **1340 GB**

To put it on a 25 GB blu ray disc: required compression factor = 53.6

Redundancy

- Coding redundancy
- Spatial and Temporal redundancy
- Irrelevant Information (often perceptually irrelevant)

Compression is all about exploiting these redundancies!

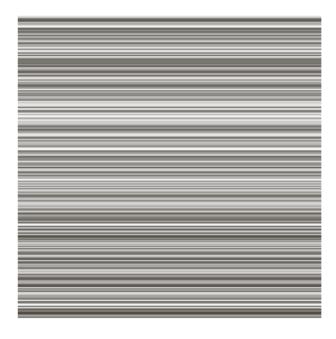
Coding redundancy



r_k	$p_r(r_k)$	Code 1	$l_1(r_k)$	Code 2	$l_2(r_k)$
$r_{87} = 87$	0.25	01010111	8	01	2
$r_{128} = 128$	0.47	10000000	8	1	1
$r_{186} = 186$	0.25	11000100	8	000	3
$r_{255} = 255$	0.03	11111111	8	001	3
r_k for $k \neq 87, 128, 186, 255$	0	_	8	_	0

Average encoding length?

Spatial and temporal redundancy



Spatial and temporal redundancy



frame t frame t+1

Spatial and temporal redundancy

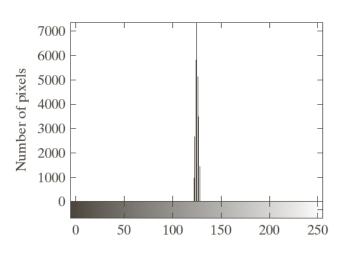


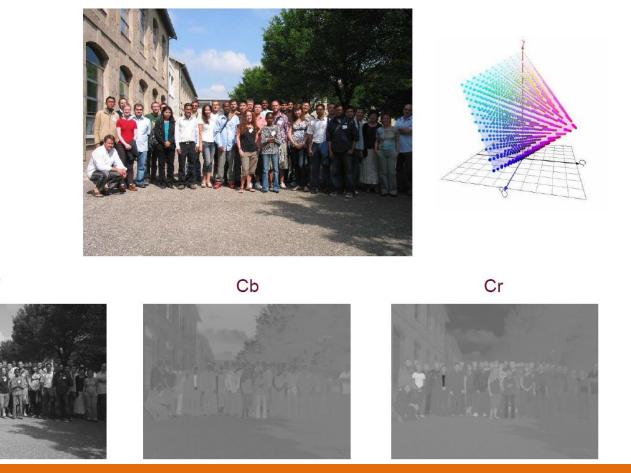




 Not all visual information is perceived by eye/brain, so throw away those that are not











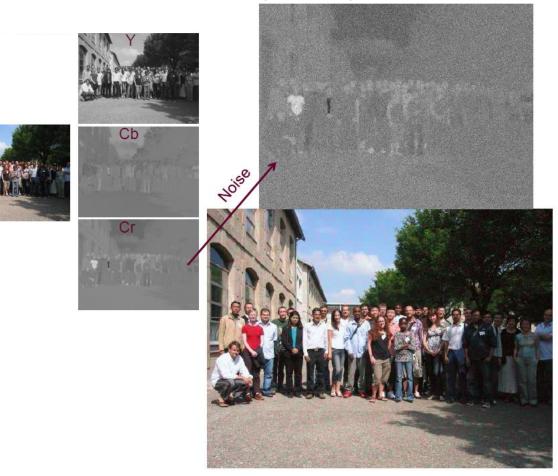












Compression types and evaluations

Two kinds:

- 1. Lossless
- 2. Lossy



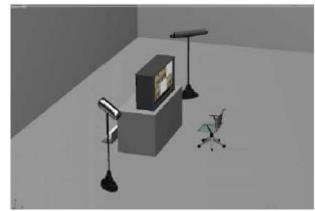
Quality measurement: judged by human viewers

- Five scale system on the degree of impairment
 - 1. Impairment is not noticeable
 - 2. Impairment is just noticeable
 - 3. Impairment is definitely noticeable, but not objectionable
 - 4. Impairment is objectionable
 - 5. Impairment is extremely objectionable

Advantages: relies on HVS

Drawbacks: time, viewing conditions, viewers?





Quality measurement: Signal to noise ratio

$$e(x,y) = f(x,y) - g(x,y).$$
 $E_{\text{ms}} = \frac{1}{MN} \sum_{x=0}^{M-1} \sum_{y=0}^{N-1} e(x,y)^2$

$$SNR_{ms} = 10 \log_{10} \left(\frac{\sum_{x=0}^{N} \sum_{y=0}^{N} g(x, y)^2}{MN \cdot E_{ms}} \right)$$

$$PSNR = 10\log_{10}\left(\frac{255^2}{E_{\rm ms}}\right)$$

Information theory: Self energy

- Information is defined as knowledge, fact, and news
- It can be measured quantitatively
- The carriers of information are symbols. Consider a symbol with an occurrence probability p. The amount of information contained in the symbol is defined as:

$$I = \log_2 \frac{1}{p}$$
 bits or $I = -\log_2 p$

Information theory: Entropy

- Consider a source that contains L possible symbols {s,i=0,1,2,...,L-1}
- With corresponding occurrence probabilities defined as $\{p_i, i=0,1,2,...,L-1\}$

Entropy

$$H = -\sum_{i=0}^{L-1} p_i \log_2 p_i$$

r_k	$p_r(r_k)$	Code 1	$l_1(r_k)$	Code 2	$l_2(r_k)$
$r_{87} = 87$	0.25	01010111	8	01	2
$r_{128} = 128$	0.47	10000000	8	1	1
$r_{186} = 186$	0.25	11000100	8	000	3
$r_{255} = 255$	0.03	11111111	8	001	3
r_k for $k \neq 87, 128, 186, 255$	0	_	8	_	0

$$log(0.47) = -1.09$$

 $log(0.03) = -5.06$

Information theory: Shannon's theorem

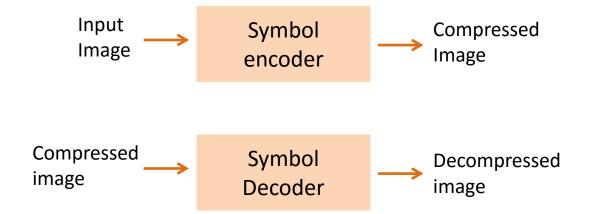
- Shannon's lossless source coding theorem states that for a discrete, memoryless, stationary information source, the minimum bit rate required to encode a symbol on average is equal to the entropy of the source.
- In other words: we can't do better than the entropy
- Lets understand with an example

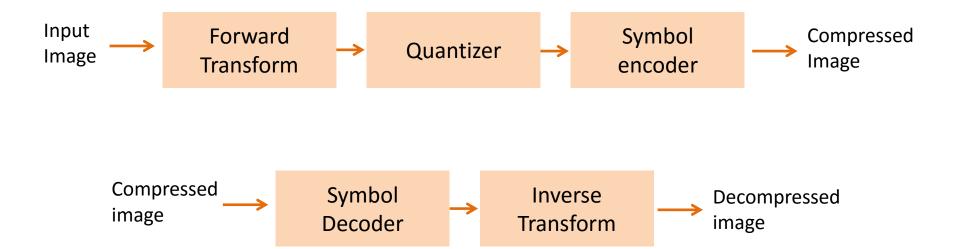
r_k	$p_r(r_k)$	Code 1	$l_I(r_k)$	Code 2	$l_2(r_k)$
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r_k for $k \neq 87, 128, 186, 255$	0	_	8	_	0

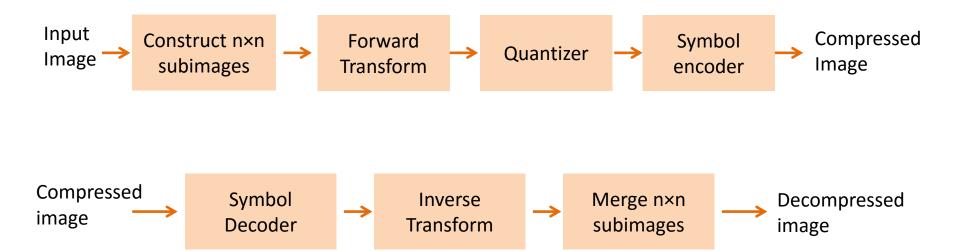
Validity of the code?

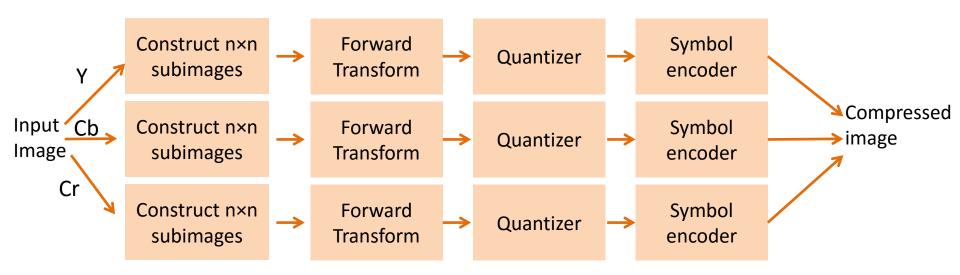
• Lets take an example

Symbol	Probability	Code1	Code2	Code3	Code4
s1	1/2	0	0	0	0
s2	1/4	0	1	10	01
s3	1/8	1	00	110	011
s4	1/8	10	11	111	0111



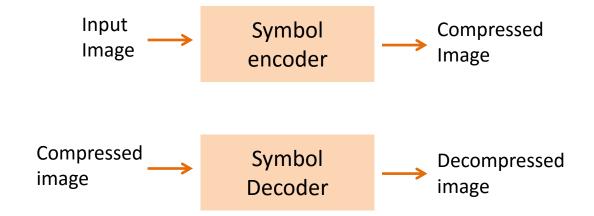






Lossless compression

Lets begin with simplest case: Lossless compression



Lossless compression: Huffman coding

- Already discussed in class
- Quick example : ABRAAKADABRAA

Lossless compression: Run Length coding

- Already discussed in class
- 15 0's, 11 1's, 8 0's, 6 1's
- 1111101110000110

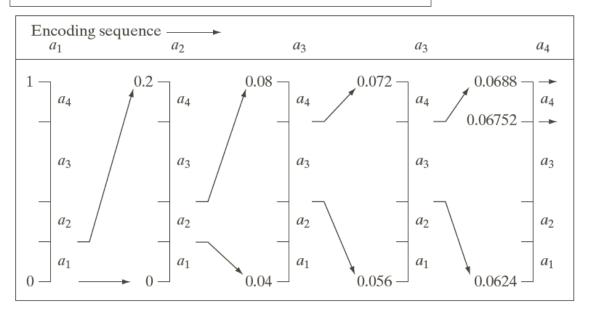
How many bits to store the count?

Give a scenario where run length coding will be extremely effective?

Lossless compression: Arithmetic coding

Source Symbol	Probability	Initial Subinterval
a_1	0.2	[0.0, 0.2)
a_2	0.2	[0.2, 0.4)
a_3	0.4	[0.4, 0.8)
a_4	0.2	[0.8, 1.0)

Input sequence: $a_1 a_2 a_3 a_3 a_4$



Final code: 0.068 (could be anything between the computed range)

3 decimal digits for 5 symbols = 3/5 digits per symbol

How many bits per symbol?

Lossless compression: Arithmetic coding

Source Symbol	Probability	Initial Subinterval
a_1	0.2	[0.0, 0.2)
a_2	0.2	[0.2, 0.4)
a_3	0.4	[0.4, 0.8)
a_4	0.2	[0.8, 1.0)

Another sequence: a_1a_1 a_3

Lossless compression: Dictionary coding

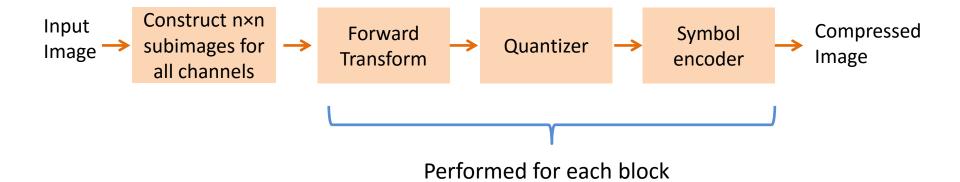
- Important in presence of recurring patterns
- Static and adaptive
- Static example: a b r a c a d a b r a
- Dynamic dictionary
 - Build during compression after observing the data
 - Rebuild at the decompression step
 - LZW is the commonly used algorithm

$A = \frac{1}{2}$	{a,	b,	С,	d,	r	ļ
<i>7</i> 1 —	ιu,	υ,	υ,	α,	•	١

Code	Entry
000	а
001	b
010	С
011	d
100	r
101	ab
110	ac
111	ad

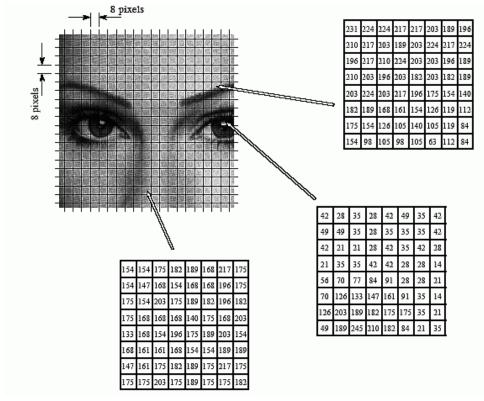
Lossy compression (with a case study of JPEG)

Lossy compression: JPEG



Block transform coding

- Partition the image into small non overlapping n×n blocks
 - 8×8 blocks in JPEG



Block Transform coding

- Process blocks using 2D transforms
- General forward transform of image g, size n×n:

$$T(u,v) = \sum_{x=0}^{n-1} \sum_{v=0}^{n-1} g(x,y) r(x,y,u,v)$$

Inverse transform

$$g(x,y) = \sum_{x=0}^{n-1} \sum_{v=0}^{n-1} T(u,v) s(x,y,u,v)$$

• r(x, y, u, v) and s(x, y, u, v) are basis functions or transformation kernels

Block Transform coding

$$T(u,v) = \sum_{x=0}^{n-1} \sum_{y=0}^{n-1} g(x,y) r(x,y,u,v) \qquad g(x,y) = \sum_{x=0}^{n-1} \sum_{v=0}^{n-1} T(u,v) s(x,y,u,v)$$

•
$$r(x, y, u, v) = e^{-j2\pi(ux+vy)/n}$$
 and $s(x, y, u, v) = \frac{1}{n^2}e^{j2\pi(ux+vy)/n}$

•
$$r(x, y, u, v) = s(x, y, u, v) = \frac{1}{n} (-1)^{\sum_{i=0}^{m-1} \lfloor b_i(x) p_i(u) + b_i(y) p_i(v) \rfloor}$$

•
$$r(x, y, u, v) = s(x, y, u, v) = \alpha(u)\alpha(v)\cos\left[\frac{(2x+1)u\pi}{2n}\right]\cos\left[\frac{(2y+1)v\pi}{2n}\right]$$

Block Transform coding: which transform to use?

Apply transform to each 8×8 block

Keep highest 50% of the coefficients in each block

Reconstruct using the inverse transform on each block

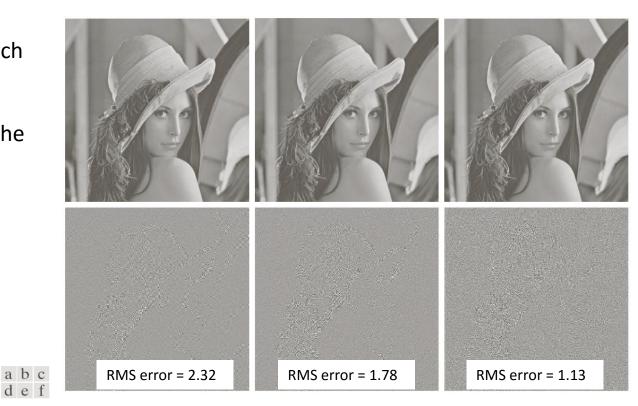
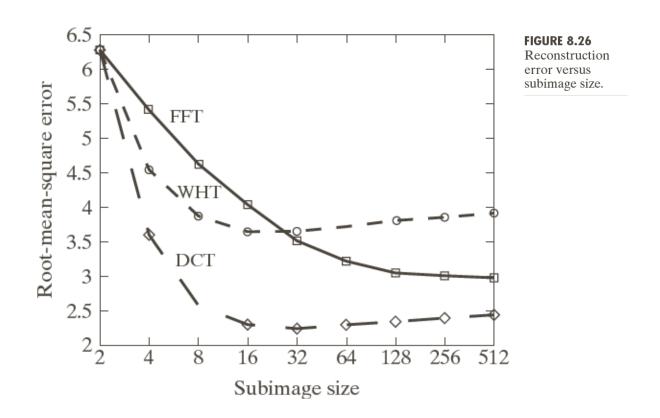
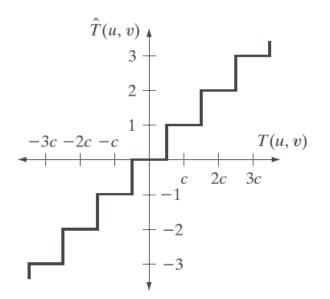


FIGURE 8.24 Approximations of Fig. 8.9(a) using the (a) Fourier, (b) Walsh-Hadamard, and (c) cosine transforms, together with the corresponding scaled error images in (d)–(f).

Block Transform coding: which transform to use?



Quantization



16	11	10	16	24	40	51	61
12	12	14	19	26	58	60	55
14	13	16	24	40	57	69	56
14	17	22	29	51	87	80	62
18	22	37	56	68	109	103	77
24	35	55	64	81	104	113	92
49	64	78	87	103	121	120	101
72	92	95	98	112	100	103	99

Different coefficients quantized with different step-size

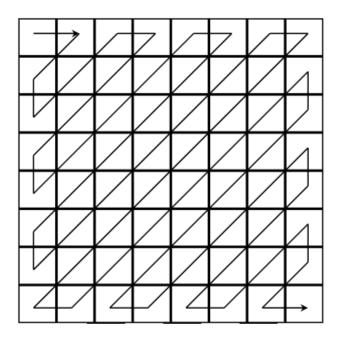
Finally encode the quantized output!

Quantization (example)

-415	-29	-62	25	55	-20	-1	3		-26	-3	-6	2	2	0	0	0
7	-21	-62	9	11	-7	-6	6		1	-2	-4	0	0	0	0	0
-46	8	77	-25	-30	10	7	-5	0	-3	1	5	-1	-1	0	0	0
-50	13	35	-15	-9	6	0	3	u	-4	1	2	-1	0	0	0	0
11	-8	-13	-2	-1	1	-4	1 -	\rightarrow	1	0	0	0	0	0	0	0
-10	1	3	-3	-1	0	2	-1		0	0	0	0	0	0	0	0
-4	-1	2	-1	2	-3	1	-2		0	0	0	0	0	0	0	0
-1	-1	-1	-2	-1	-1	0	-1		0	0	0	0	0	0	0	0
	•	•			•						•	•	•			

16	11	10	16	24	40	51	6
12	12	14	19	26	58	60	5:
14	13	16	24	40	57	69	50
14	17	22	29	51	87	80	62
18	22	37	56	68	109	103	7
24	35	55	64	81	104	113	92
49	64	78	87	103	121	120	10
72	92	95	98	112	100	103	99

Symbol encoding (Zigzag ordering)

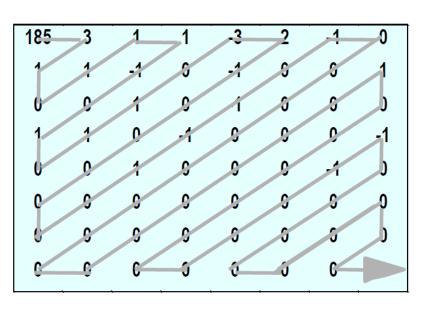


0	1	5	6	14	15	27	28
2	4	7	13	16	26	29	42
3	8	12	17	25	30	41	43
9	11	18	24	31	40	44	53
10	19	23	32	39	45	52	54
20	22	33	38	46	51	55	60
21	34	37	47	50	56	59	61
35	36	48	49	57	58	62	63

JPEG uses run length encoding!

Symbol coding example

Zigzag scan (additional example)



Run length coding



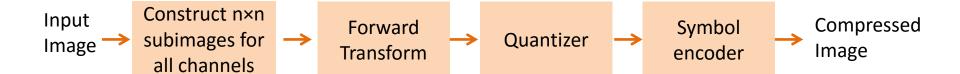
Mean of Block: 185

(0,3) (0,1) (1,1) (0,1) (0,1) (0,1) (0,-1) (1,1)

(1,1) (0,1) (1,-3) (0,2) (0,-1) (6,1) (0,-1) (0,-1)

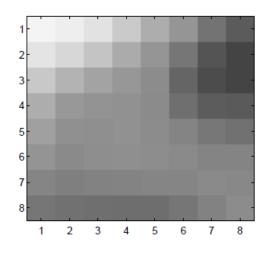
(1,-1) (14,1) (9,-1) (0,-1) EOB

Lossy compression: JPEG



Lest understand the entire procedure with an example

Consider a single 8×8 pixel block B:



- Intensity range → [0 255]
- Subtract 127 from each entry and computer 2D DCT

Forward transform and quantization

DCT of image block

$$\hat{\mathbf{B}} = \begin{pmatrix} 118.9 & 187.7 & -17.7 & 16.8 & 14.4 & 2.4 & 5.3 & 3.5 \\ 104.1 & 187.1 & -30.8 & 10.0 & -1.0 & -4.7 & 0.6 & 0.3 \\ 46.3 & 10.4 & 9.1 & -9.0 & -15.7 & 0 & -1.3 & -2.7 \\ 76.8 & -12.1 & -10.7 & -0.2 & -10.4 & 4.8 & 2.7 & -3.3 \\ 6.4 & -15.3 & 1.7 & -1.7 & -1.1 & 2.5 & 1.1 & -2.5 \\ 10.6 & -5.6 & -6.5 & -0.6 & 2.6 & 0.9 & -1.4 & 2.4 \\ 0.4 & -2.3 & 1.2 & -1.7 & 2.3 & -0.5 & 0.1 & -0.1 \\ 3.2 & -0.7 & -0.9 & 2.6 & -1.1 & 1.5 & -1.8 & 0.2 \end{pmatrix}$$

Quantization and rounding

16	11	10	16	24	40	51	61
12	12	14	19	26	58	60	55
14	13	16	24	40	57	69	56
14	17	22	29	51	87	80	62
18	22	37	56	68	109	103	77
24	35	55	64	81	104	113	92
49	64	78	87	103	121	120	101
72	92	95	98	112	100	103	99

More than 75% entries are zero (notice their placement)

Encoding

Zigzag scan

0	1	5	6	14	15	27	28
2	4	7	13	16	26	29	42
3	8	12	17	25	30	41	43
9	11	18	24	31	40	44	53
10	19	23	32	39	45	52	54
20	22	33	38	46	51	55	60
21	34	37	47	50	56	59	61
35	36	48	49	57	58	62	63

[**7** 17 9 3 16 -2 1 -2 1 5 0 -1 1 1 1 0 0 0 0 -1 EOB]

Lets try to reconstruct

Reconstruction: Decoding + Dequantization

[**7** 17 9 3 16 -2 1 -2 1 5 0 -1 1 1 1 0 0 0 0 -1 EOB]

Dequantization

16	11	10	16	24	40	51	61
12	12	14	19	26	58	60	55
14	13	16	24	40	57	69	56
14	17	22	29	51	87	80	62
18	22	37	56	68	109	103	77
24	35	55	64	81	104	113	92
49	64	78	87	103	121	120	101
72	92	95	98	112	100	103	99

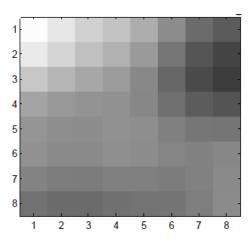
Decoding

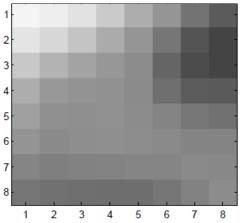
Compute IDCT and add 127

```
254
       233
             211
                    197
                           175
                                 142
                                        110
                                               93
                                               70
236
       216
             194
                    179
                           156
                                 120
                                        87
201
       184
             169
                           138
                                 105
                                               62
                    158
                                        76
166
       155
             149
                           137
                                 114
                                        94
                                               86
                    147
151
       144
             143
                    147
                           144
                                 130
                                        120
                                              119
                                              137
147
       140
             140
                    145
                           143
                                 135
                                        132
133
       127
             126
                    130
                           130
                                 126
                                        131
                                              140
                                               141/
116
       109
             109
                    114
                           117
                                 118
                                        128
```

Compare with original

/245	239	227	203	174	150	116	92 \
229	216	197	172	150	119	85	69
201	180	164	152	141	102	77	69
174	153	148	146	140	112	93	91
161	145	144	146	141	133	120	114
149	139	143	144	142	139	133	133
134	128	131	132	134	134	139	137
119	114	112	111	111	119	131	141/
	229 201 174 161 149 134	229 216 201 180 174 153 161 145 149 139 134 128	229 216 197 201 180 164 174 153 148 161 145 144 149 139 143 134 128 131	229 216 197 172 201 180 164 152 174 153 148 146 161 145 144 146 149 139 143 144 134 128 131 132	229 216 197 172 150 201 180 164 152 141 174 153 148 146 140 161 145 144 146 141 149 139 143 144 142 134 128 131 132 134	229 216 197 172 150 119 201 180 164 152 141 102 174 153 148 146 140 112 161 145 144 146 141 133 149 139 143 144 142 139 134 128 131 132 134 134	229 216 197 172 150 119 85 201 180 164 152 141 102 77 174 153 148 146 140 112 93 161 145 144 146 141 133 120 149 139 143 144 142 139 133 134 128 131 132 134 134 139

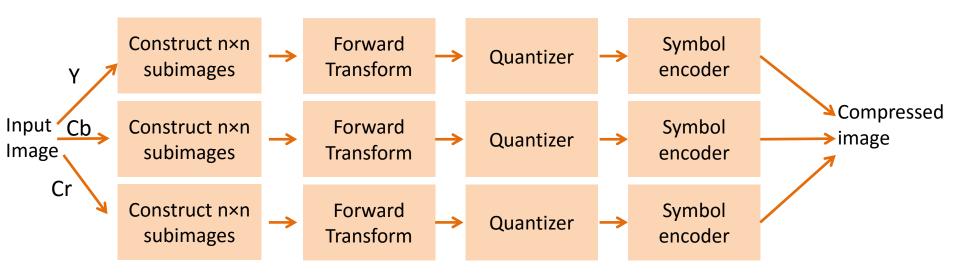




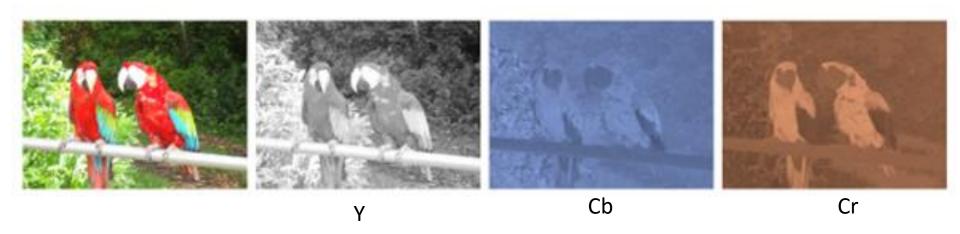
Summary JPEG

- Divide into 8×8 subimages
- Compute DCT on each
- Quantize the coefficients
- Order coefficients in zigzag pattern
- Encode 1D sequence using run-length coding and Huffman coding

Color Images



Color Images



- Different quantization matrices for chrominance and luminance
- Chroma subsampling (use reduced resolution of chroma channels)

Quantization matrices

16	11	10	16	24	40	51	61
12	12	14	19	26	58	60	55
14	13	16	24	40	57	69	56
14	17	22	29	51	87	80	62
18	22	37	56	68	109	103	77
24	35	55	64	81	104	113	92
49	64	78	87	103	121	120	101
72	92	95	98	112	100	103	99

17	18	24	47	99	99	99	99
18	21	26	66	99	99	99	99
24	26	56	99	99	99	99	99
47	66	99	99	99	99	99	99
99	99	99	99	99	99	99	99
99	99	99	99	99	99	99	99
99	99	99	99	99	99	99	99
99	99	99	99	99	99	99	99

Luminance

Chrominance

These matrices are scaled for higher compression!